

# AI4010: Online Learning

## Second Midterm Exam

### Oct 2025

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**Instructions:**

- The total number of marks is 20.
  - The total duration of the exam is 90 minutes. No electronic aids are allowed. You can keep a maximum of one sheet of paper with formulas/notes.
  - All questions are mandatory. A yes/no answer without proper proof or justification will be given zero marks even if it is correct.
  - Any plagiarism case, if detected, will attract F grade in the course irrespective of overall performance.
  - Use  $0 \log(0) = 0$  wherever required.
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**Problem 0.1** (2 marks). Consider that the loss functions are bounded in range  $[0, 1]$ . Show that the regret of FTL algorithm is upper bounded by the number of times the leader<sup>1</sup> is changed during the sequence of plays.

**Problem 0.2** (3 marks). (Batch vs Online convex optimization) Consider a standard batch convex optimization problem: you want to find a minimum of the convex function  $f : \mathbb{R} \rightarrow \mathbb{R}_+ \setminus \{\infty\}$  over a bounded convex set  $\mathcal{K} \subset \mathbb{R}$  to within an accuracy  $\varepsilon > 0$ . In other words, you must output  $x \in \mathcal{K}$  satisfying  $f(x) \leq \min_{y \in \mathcal{K}} f(y) + \varepsilon$ .

You are given an online algorithm ALG with the following property. For any number of rounds  $t \geq 1$  and any sequence of non-negative loss functions  $\{f_s\}_{s=1}^t$  from a family of convex functions  $\mathcal{F}$ , the algorithm's regret over<sup>2</sup> a convex set  $\mathcal{K}$  is at  $t^\alpha$  for some  $\alpha \in [0, 1)$ ; i.e. ALG satisfies sublinear regret guarantee. How will you accomplish batch optimization objective using ALG? Show your work.

*Hint: Use Jensen's inequality.*

**Problem 0.3** (2+2+2 = 6 marks). Compute the Bregman Divergence of

$$1. R(x) = \frac{1}{2} \|x\|_2^2 \text{ for } x \in \mathbb{R}^d$$

$$2. R(x) = -2 \sum_{i=1}^d \sqrt{x_i} \text{ for } x \in (0, +\infty)^d$$

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<sup>1</sup>The leader at time  $t$  is a choice  $x_{t+1}$  i.e. optimal choice of the next round.

<sup>2</sup>The regret is computed with respect to any single point in  $\mathcal{K}$ .

$$3. R(x) = \sum_{i=1}^d x_i(\log(x_i) - 1) \text{ for } x \in \Delta_d.$$

**Problem 0.4 (Dual Function).** In class we saw that the OMD algorithm<sup>3</sup> performs the update in the dual space defined by the gradients. This can be interpreted (slightly more elaborately) as follows. OMD first maps the point  $x_t$  from primal space (i.e.  $x$ -space) to the dual space defined by  $\nabla R(\cdot)$  i.e. computes  $\theta_t = \nabla R(x_t)$ , performs the (gradient) update in gradient space i.e. computes  $\theta_{t+1} = \theta_t - \eta_t \nabla_t$ , maps the updated point using  $(\nabla R)^{-1}$  back to the original space i.e. computes  $y_{t+1}$  such that  $\nabla R(y_{t+1}) = \theta_{t+1}$  i.e.  $y_{t+1} = (\nabla R)^{-1}(\theta_{t+1})$  and then take the Bregman projection onto the convex set i.e. compute  $x_{t+1} = \arg \min_{x \in \mathcal{K}} B_R(x || y_{t+1})$ . Define a function  $R^*(\theta) = \sup_{x \in \mathbb{R}^d} (\langle x, \theta \rangle - R(x))$ . The function  $R^*(\cdot)$  is called as the Fenchel dual/conjugate of  $R$ . In this question we will show that  $\nabla R^* = (\nabla R)^{-1}$ .

1. Write the Fenchel conjugates of below functions. [2+2 = 4 marks]

(a)  $R(x) = \log(\sum_{i=1}^d e^{x_i})$  (use  $0 \log(0) = 0$ )

(b)  $R(x) = \frac{1}{2} x^T Q x$  where  $Q \in \mathbb{R}^{n \times n}$  is a positive definite matrix.

2. Show that the following two conditions are equivalent. [2 marks]

(a)  $R(u) + R^*(v) = \langle u, v \rangle$

(b)  $v = \nabla R(u)$

3. Finally, using the above result and also using the fact that  $R^{**} = R$ , show that  $u = \nabla R^*(\nabla R(u))$  and  $u = \nabla R(\nabla R^*(u))$ . [3 marks]

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<sup>3</sup>Here, we will consider all assumptions we made for OMD.