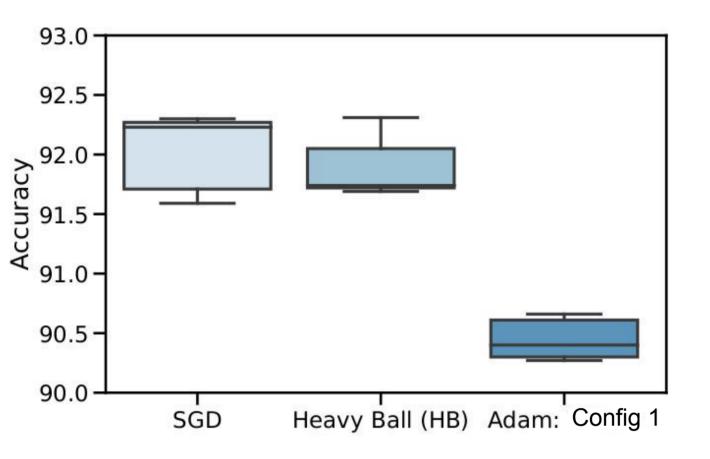
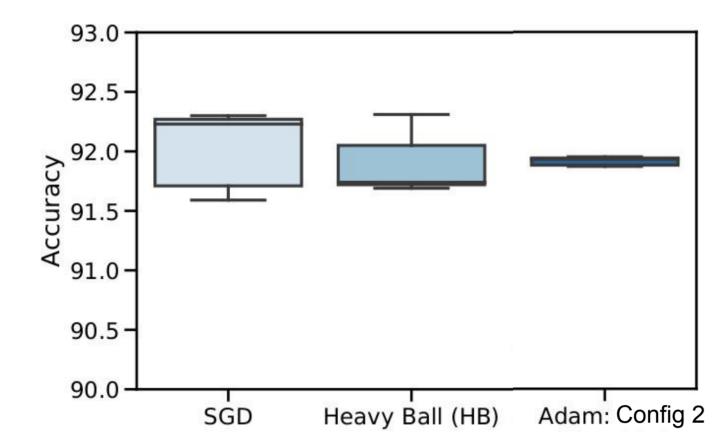
Hyperparameter Optimization Is Deceiving Us, and How to Stop It

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We want to prevent our conclusions about algorithm performance from depending on the underlying configuration of the hyperparameter optimization (HPO) that we perform

We do not know the ground truth





non-adaptive optimizers outperform adaptive optimizers

AND

non-adaptive optimizers outperform adaptive optimizers

→ **FALSE**

It is **fine** to **accept either** to form conclusions

It is **fine** to **accept neither** to form no conclusions

It is not fine to accept both to form inconsistent conclusions

We do not want it to be **possible** to form inconsistent conclusions because we want to derive reliable knowledge about algorithm performance

This is challenging because

- the process of picking HPO configurations to test is vague
- whether we believe our conclusions or not is also vague

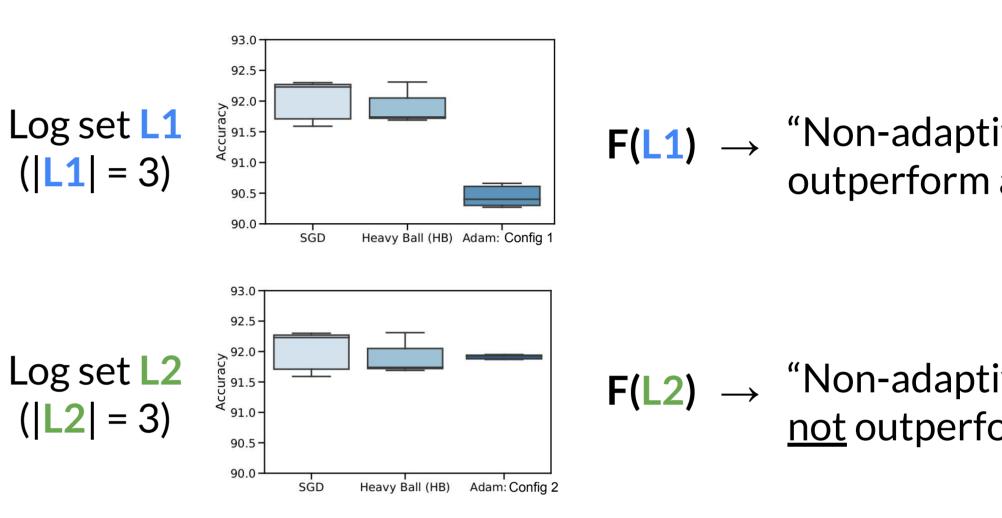
We remove the vagueness from the problem by pinning down

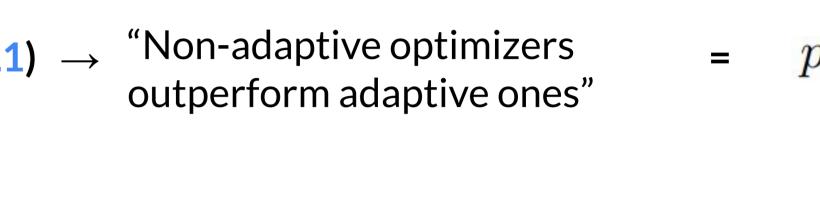
- a concrete formalization for the
- possible outcomes of running HPO
- A concrete formalization for our **belief in** conclusions

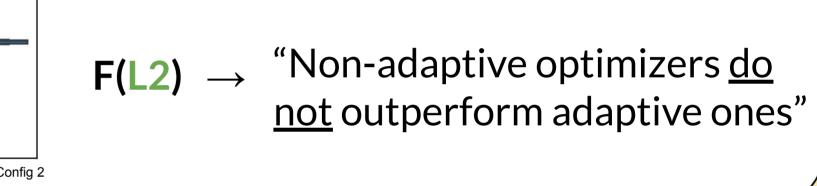
To do so, we define HPO to return a **log**, which records all the choices and measurements made during an HPO run (enabling reproducibility)

We then formalize the process of drawing conclusions from empirical studies using HPO:

Epistemic Hyperparameter Optimization (EHPO) takes a set of HPO procedures and a **function F**, which maps a set of HPO logs to conclusions about algorithm performance







We imagine a demon trying to deceive us about algorithm performance via EHPO

The demon maintains a set of HPO logs that it can modify, and presents us with a final log set, from which we can draw conclusions

The demon *could* produce **L1** or *could* produce

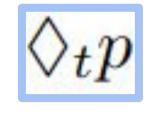


L2, and then could discard L2

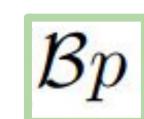
Modal logic is the standard way to formalize **could** by extending propositional logic to allow us to reason about possibility by introducing an additional operator

 $\Diamond \phi$ reads, "It is **possible** that ϕ "

We combine two modal operators so that we can capture the idea that it is not possible to adopt inconsistent beliefs about experimental outcomes



means a demon could adopt strategy guaranteed to cause **desired outcome** *p*, taking time at most **t** in expectation; a set of EHPO output logs models that it is possible **p** in time **t**



means we believe/conclude **p**; a set of EHPO output logs models our belief in **p**

With both of these operators, we can formalize the problem of hyperparameter deception

t-non-deceptive axiom



If it is **possible** for the demon can get us to **believe p** in time **t**, then it is **not possible** for the demon to get us to **believe** not p in time t

We use this formalization to prove non-trivial theorems about whether a HPO procedure is defended against deception

We can do this by proving that an EHPO satisfies our t-non-deceptive axiom

Intuitively, our defense works as follows: Given some **naive** reasoner, we construct a defended reasoner that is always more skeptical than the naive reasoner

If the naive reasoner is t-non-deceptive, then any more skeptical reasoner is also *t*-non-deceptive

The important takeaway is that it is **always possible** to construct a t-defended EHPO

We describe a **defended variation of** random search, which is defended against deception for **t**, but it is able to draw conclusions using up compute budgets that are O(√t)

	\boldsymbol{p}	$\neg p$	$1-\delta$	Conclude
SGD vs. Adam	0.213	0.788	0.75	$\neg p$
			0.8	Nothing
			0.9	Nothing
HB vs. Adam	0.168	0.832	0.75	$\neg p$
			0.8	$\neg p$
			0.9	Nothing

It makes conclusions using much fewer resources than the total compute budget for which it is defended against deception