

Bitwise Operators && Functions

Special class

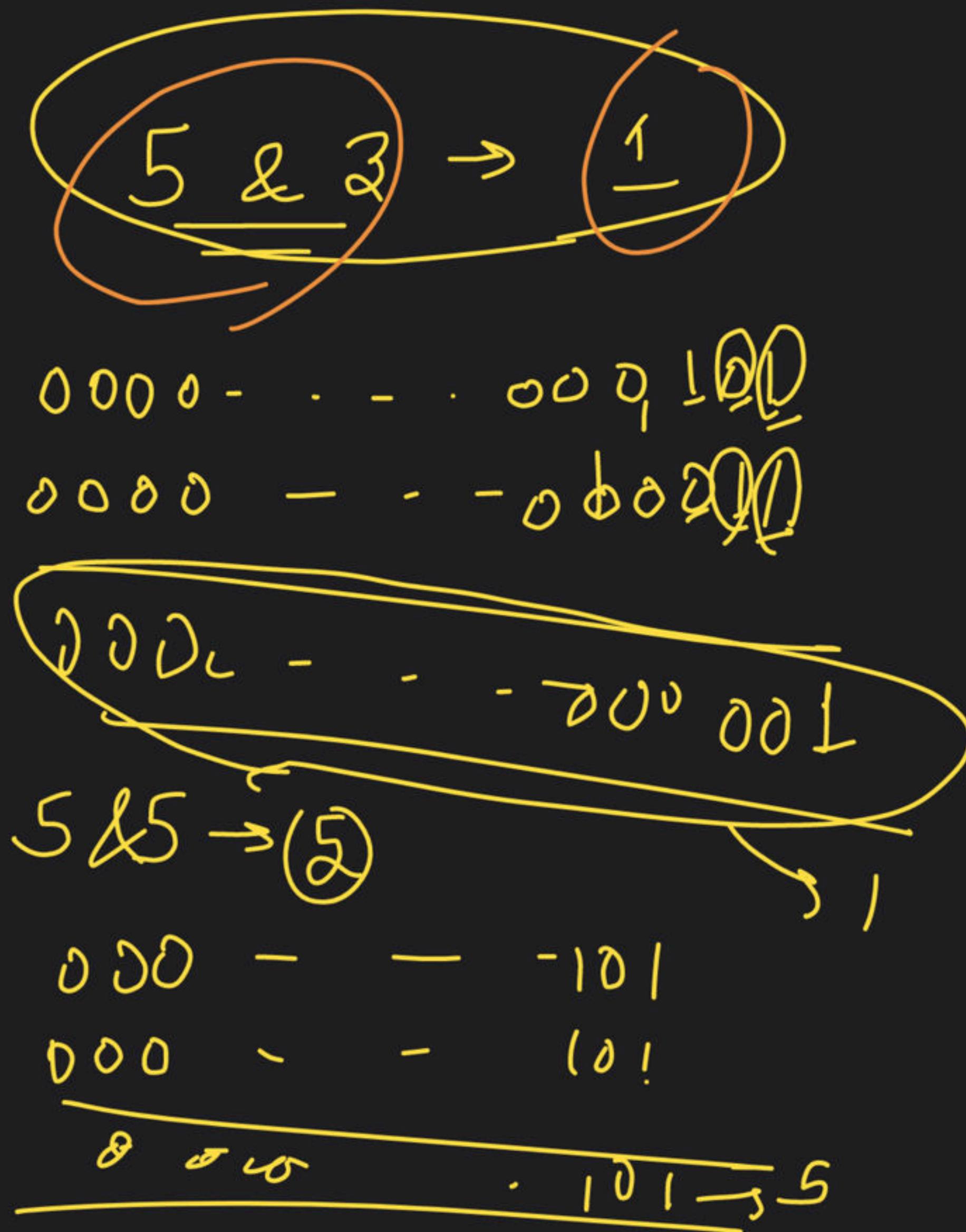
→ Bitwise Operators:

- & (and)
- | (or)
- ^ (xor)
- ~ (not)

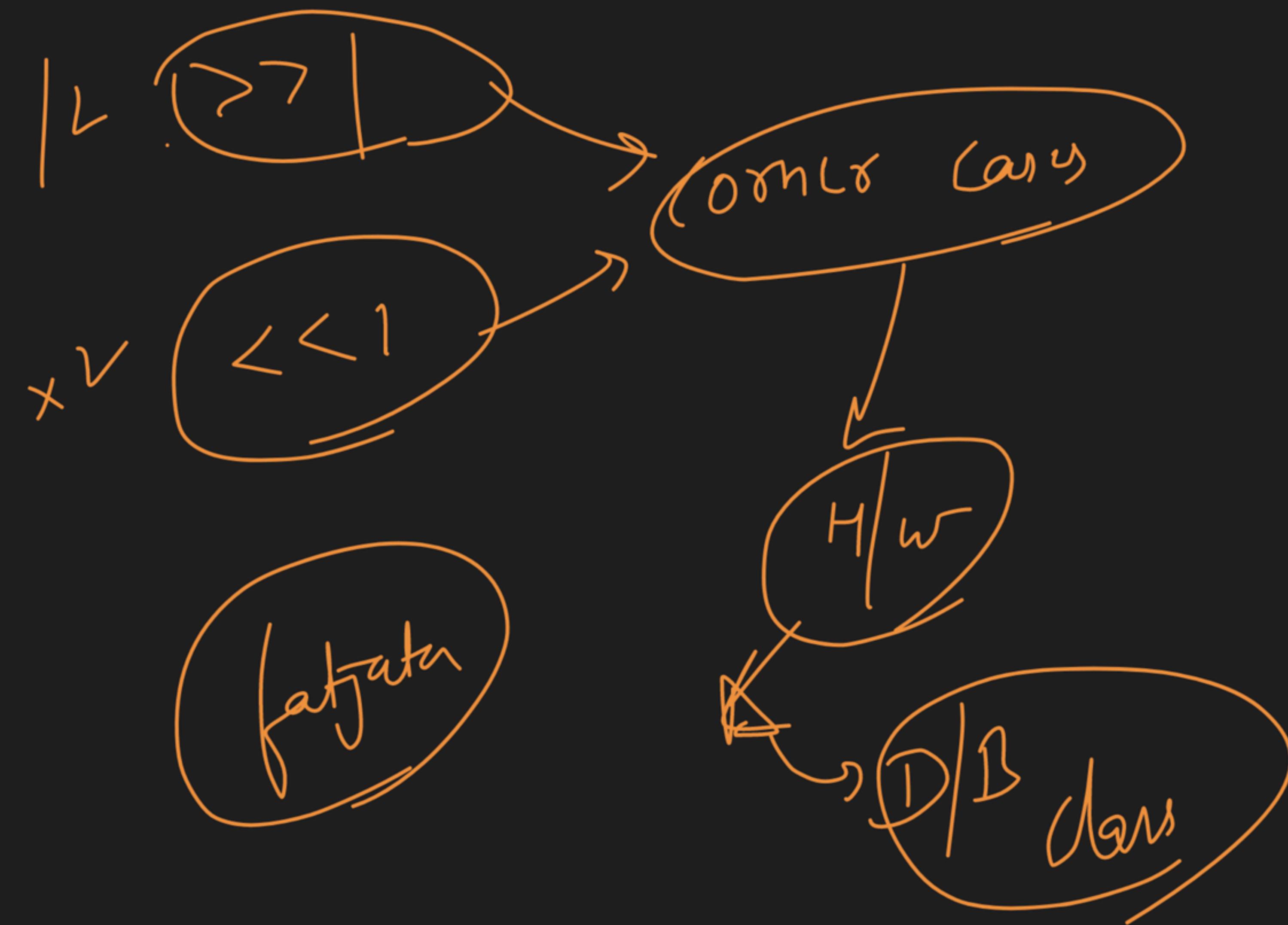
0 1
~ ~
 $<<$ (left shift)
 $>>$ (right shift)

0 &

a	b	o/p
0	0	0
0	1	0
1	0	0
1	1	1



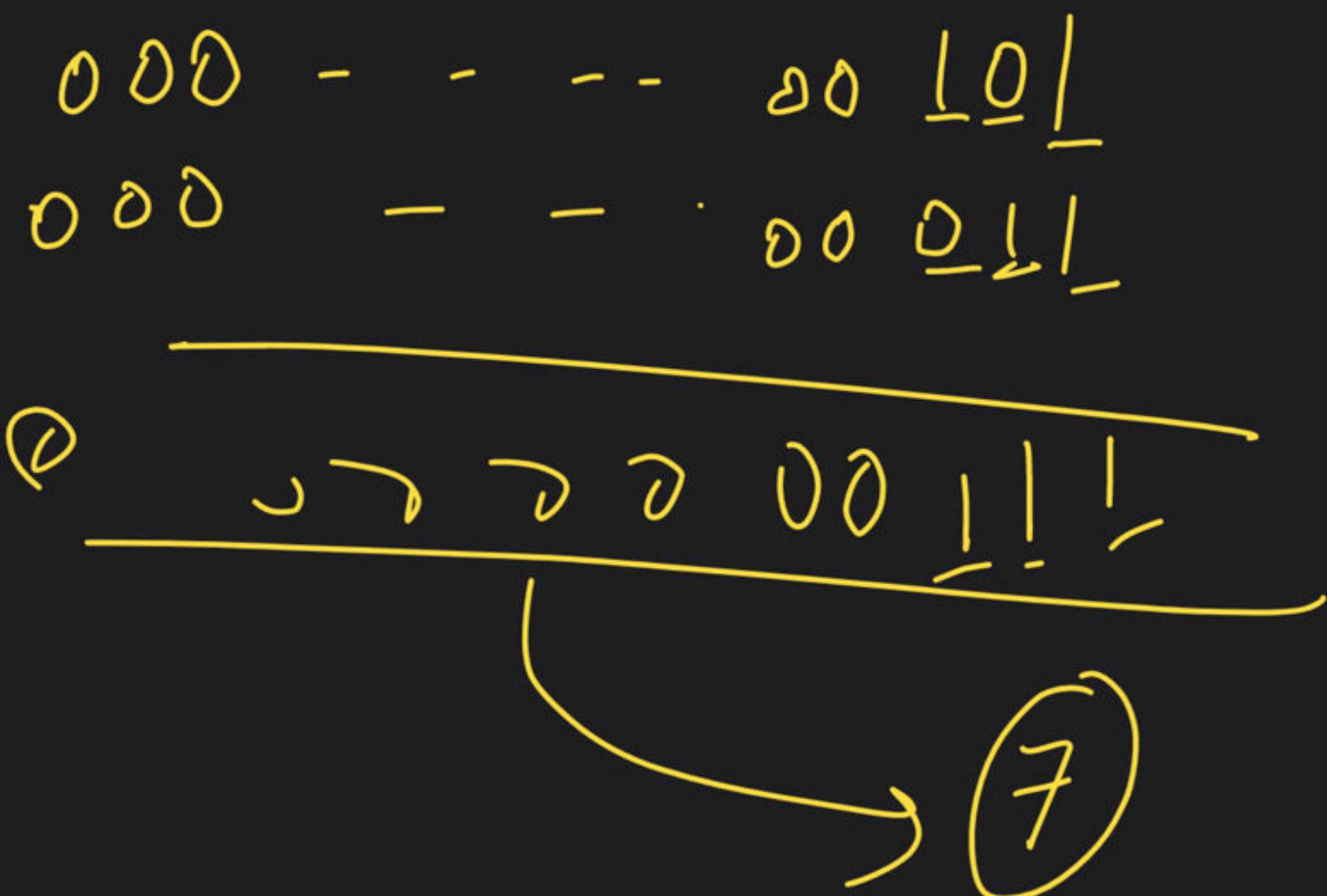




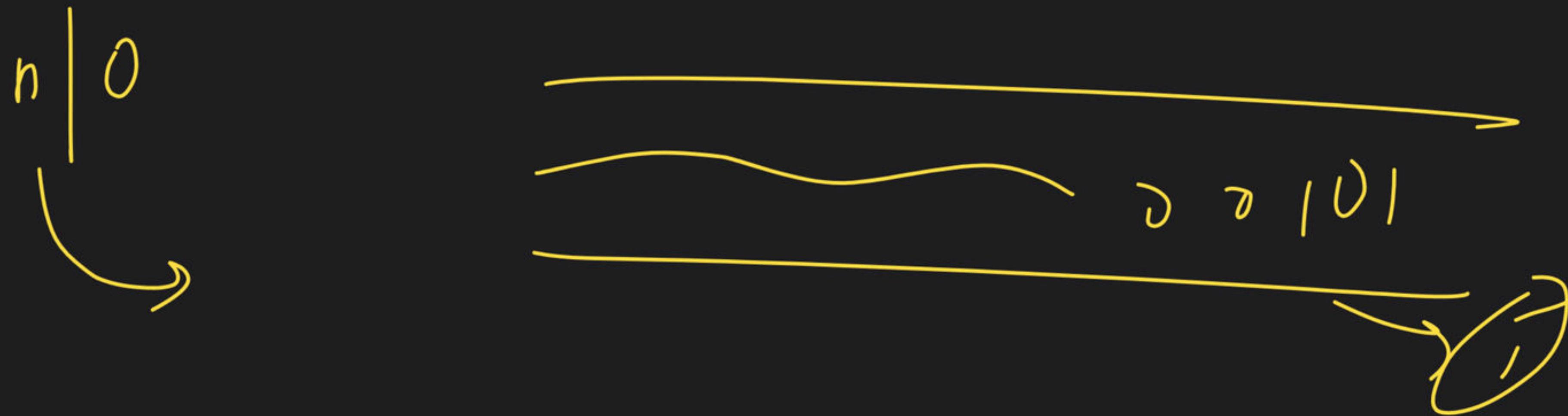
(2)

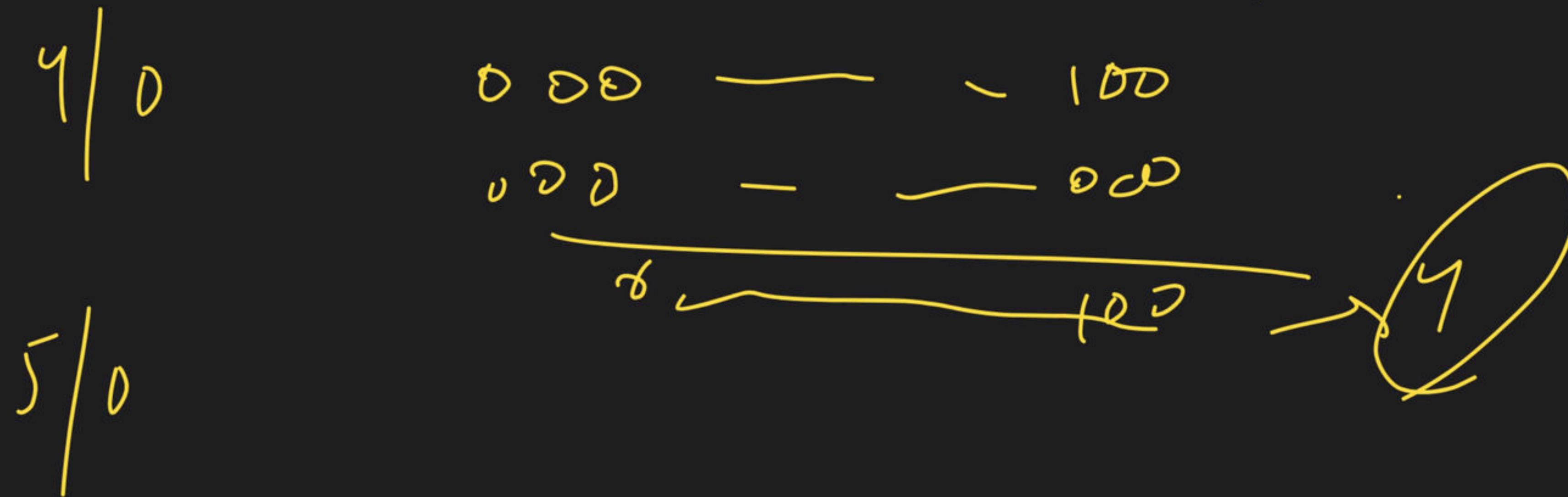
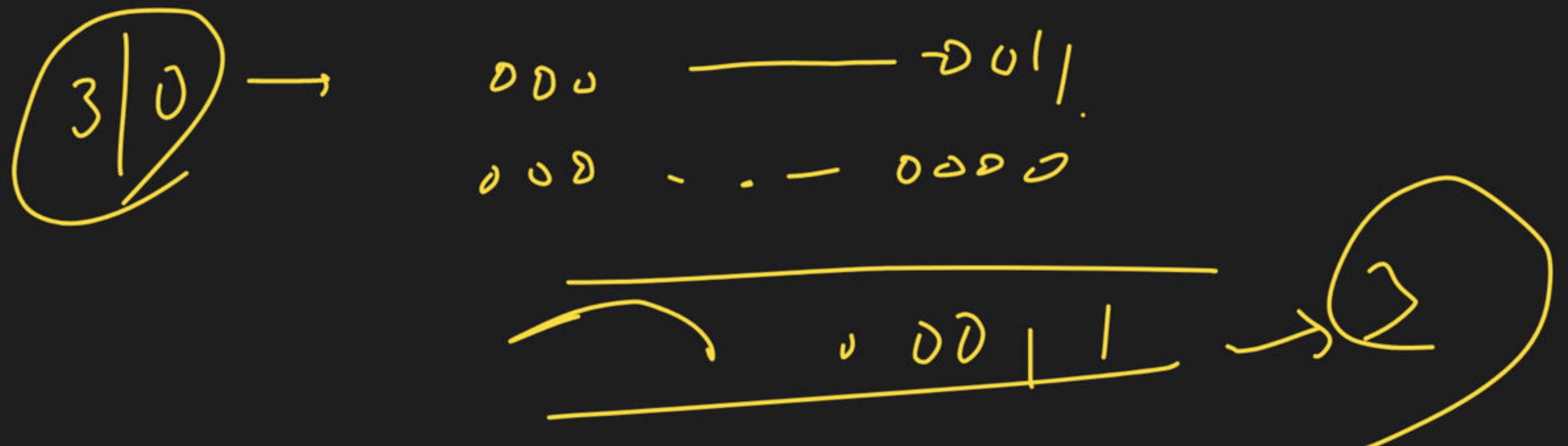
a	b	d/p
0	0	0
Q	1	1
1	0	1
1	1	1

$$\begin{array}{r} 5 \\ \underline{\quad} \\ 3 \end{array}$$



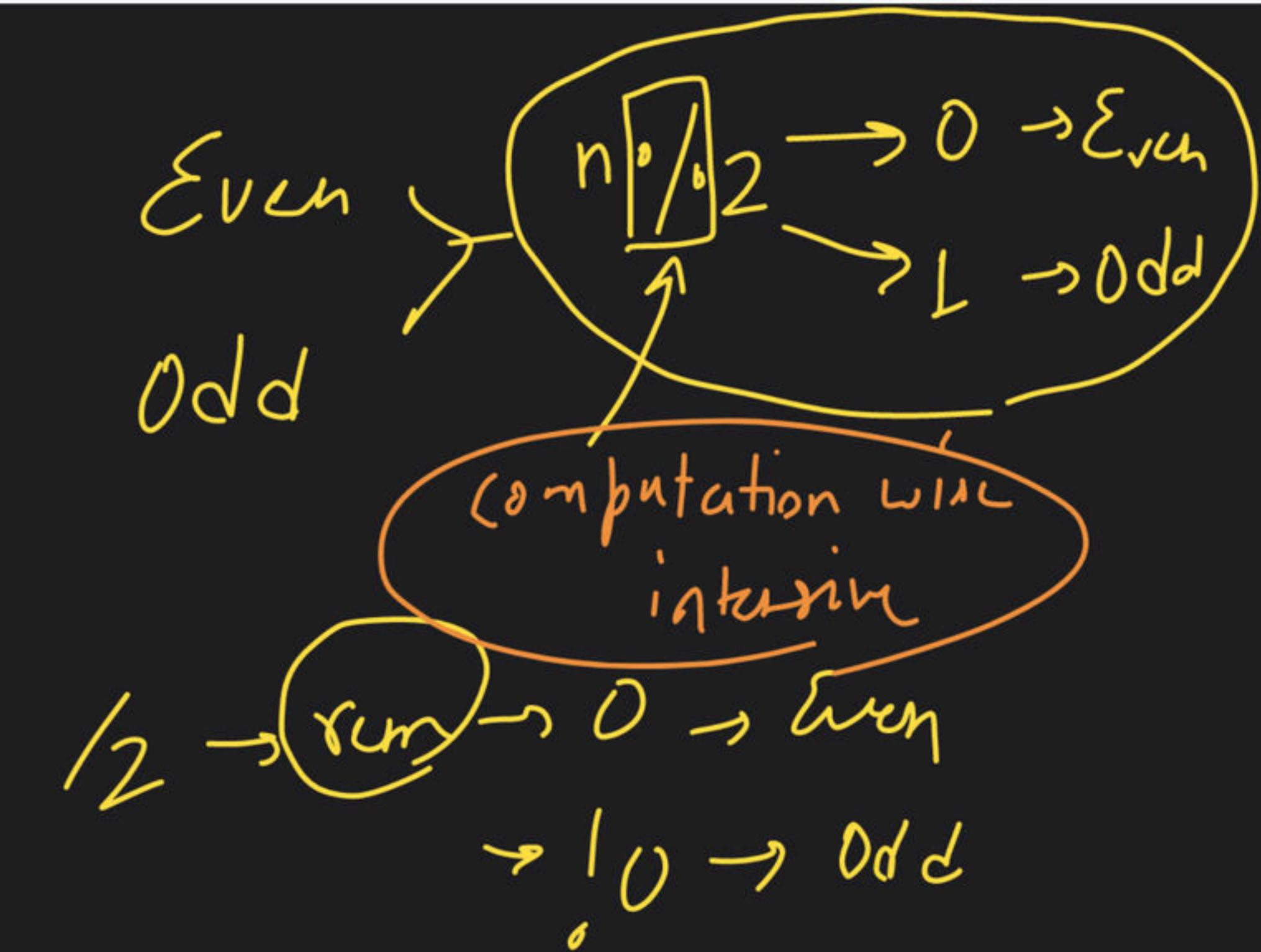
$$\begin{array}{r} 5 \quad | \quad 5 \\ 000 \quad . \quad - \quad - \quad - \quad 00 \quad | \quad 0 \quad | \\ 000 \quad - \quad - \quad - \quad - \quad 00 \quad | \quad 0 \quad | \\ \hline & & & & & 00 \quad | \quad 0 \quad | \end{array}$$





$\&$ (Bitwise AND)

$n \& 1 \rightarrow$



$\gamma \& 1$ \rightarrow

00000 - - - 00100

00000 - - . 00001

$\alpha\beta$ 2 0 0000

$\beta\alpha\beta$ \rightarrow

000 - - - 0010

000 - - - 0001

$\beta\alpha\beta$

00000

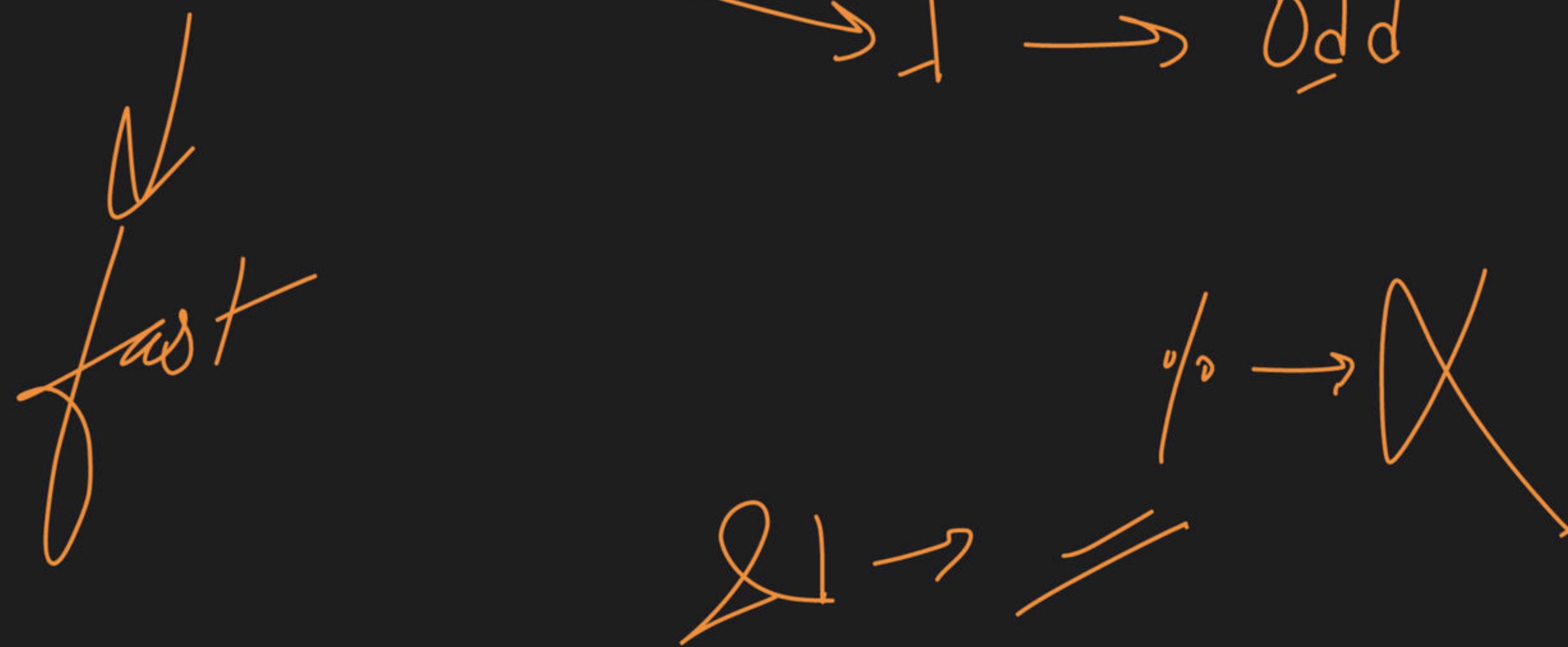
00000

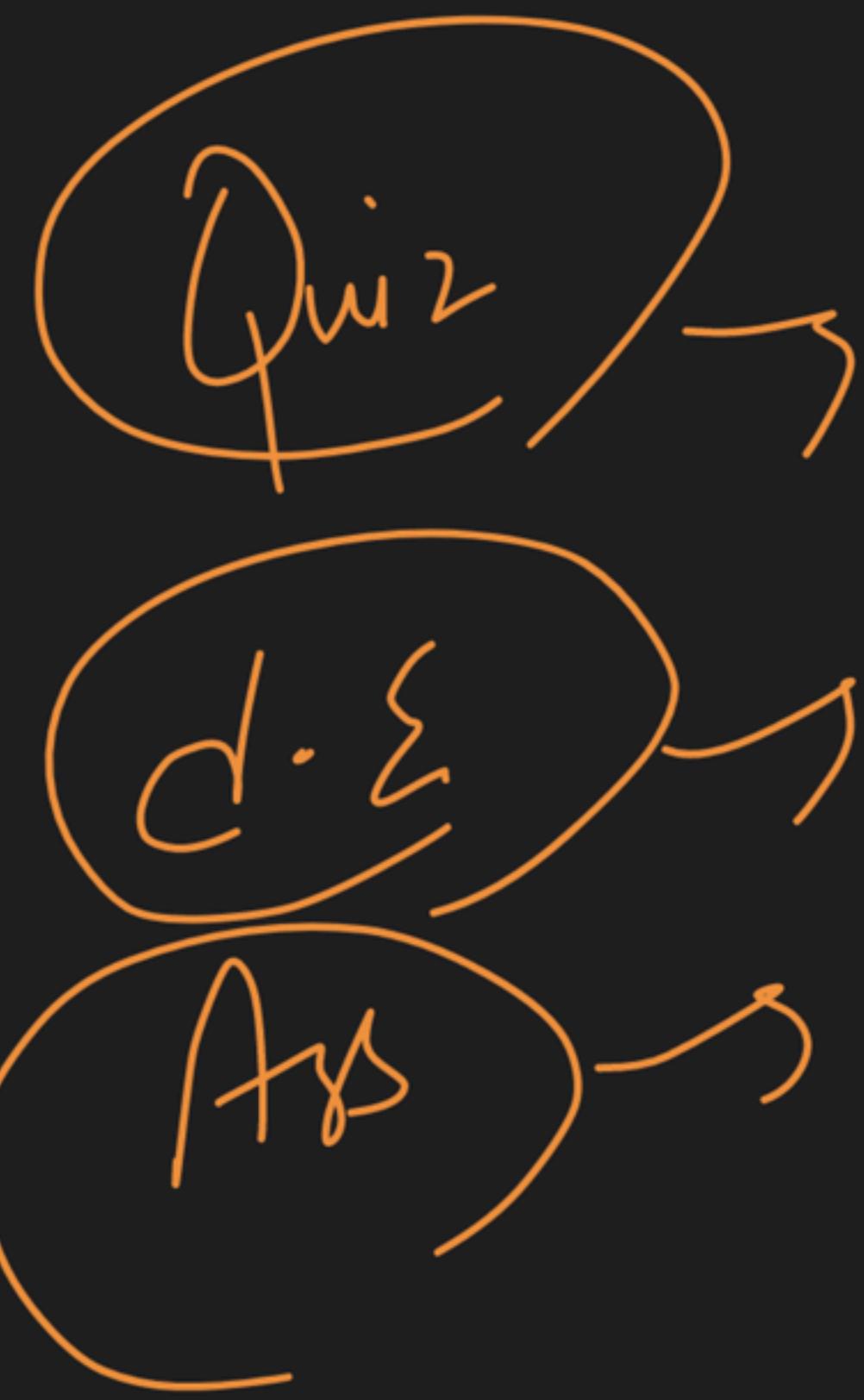
$\gamma \& 1$

$\beta\alpha\beta$

$\beta\alpha\beta$

$\beta\alpha\beta$



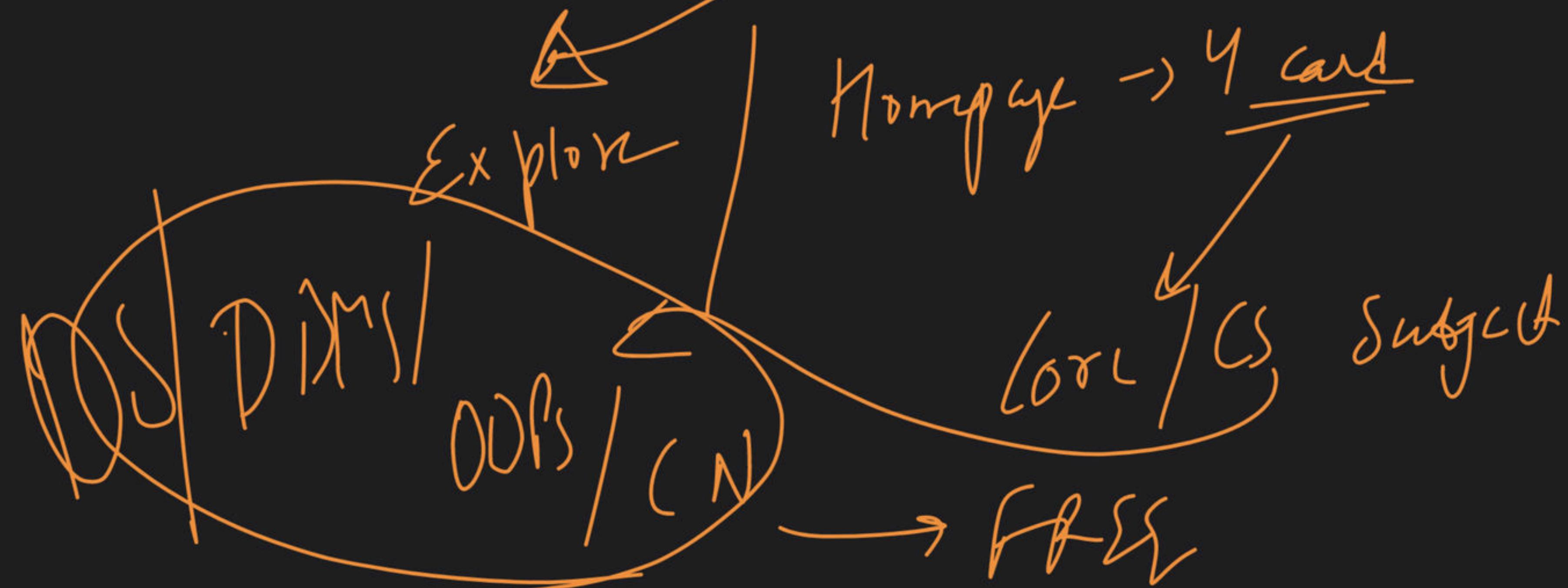


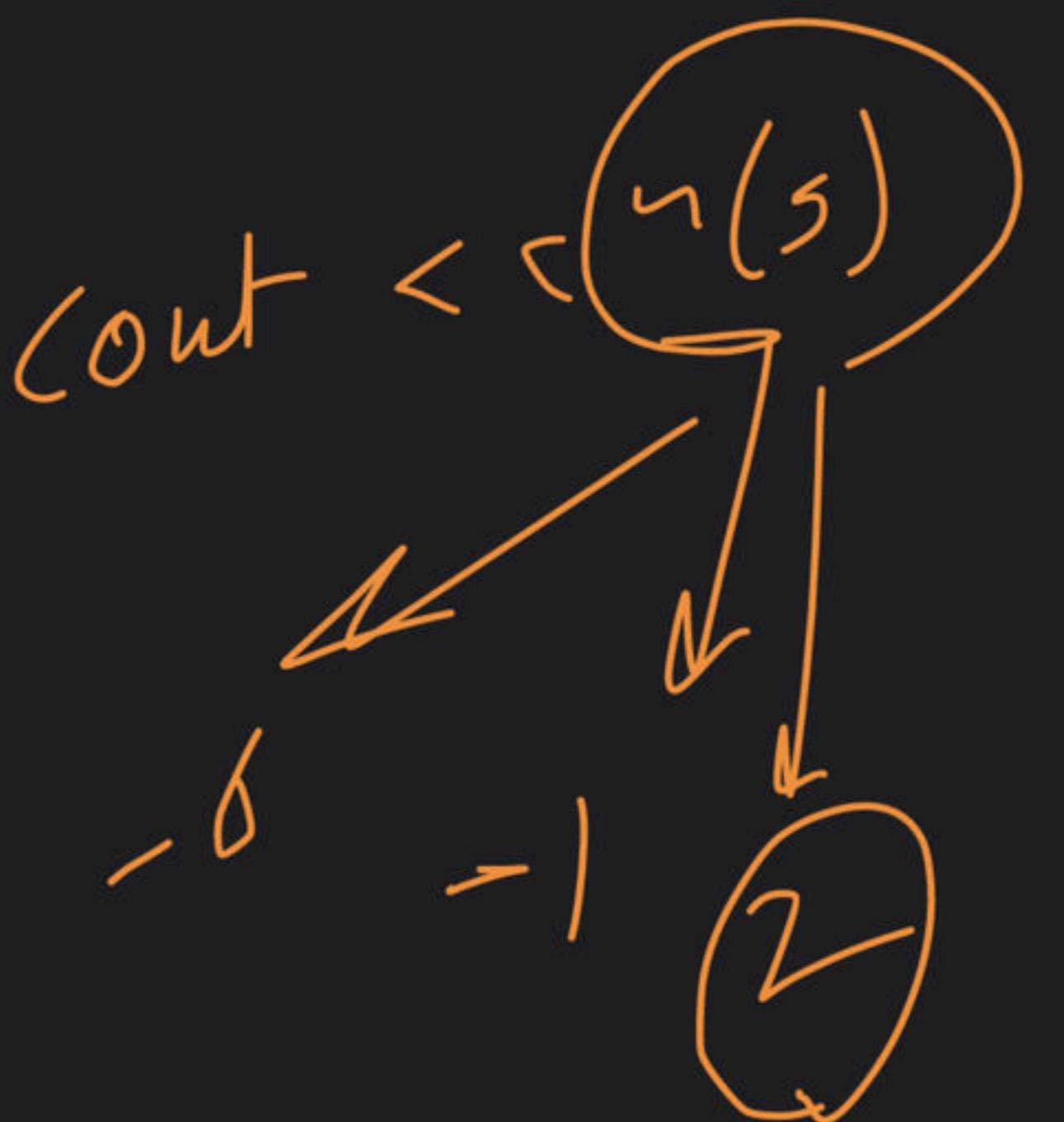
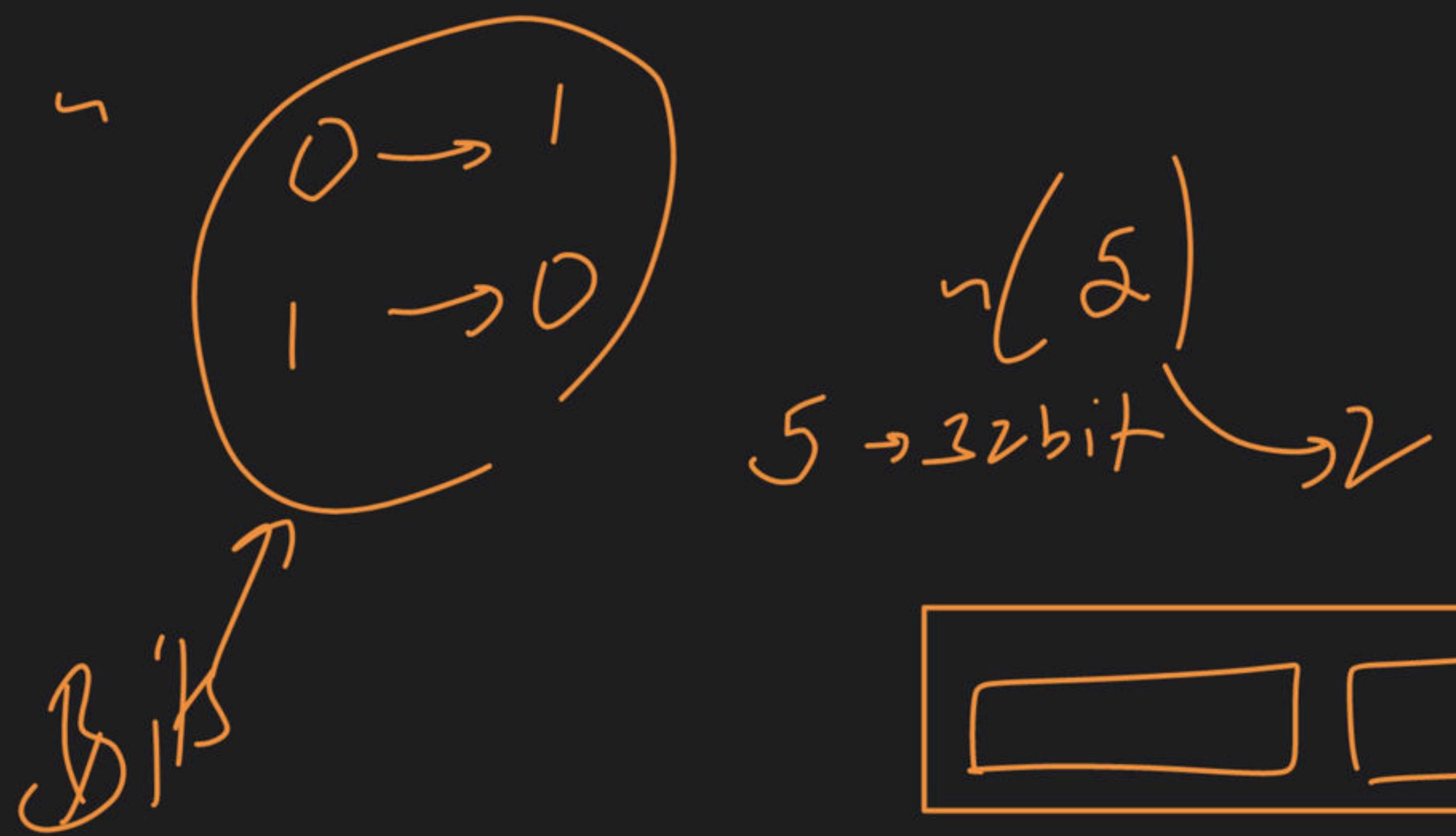
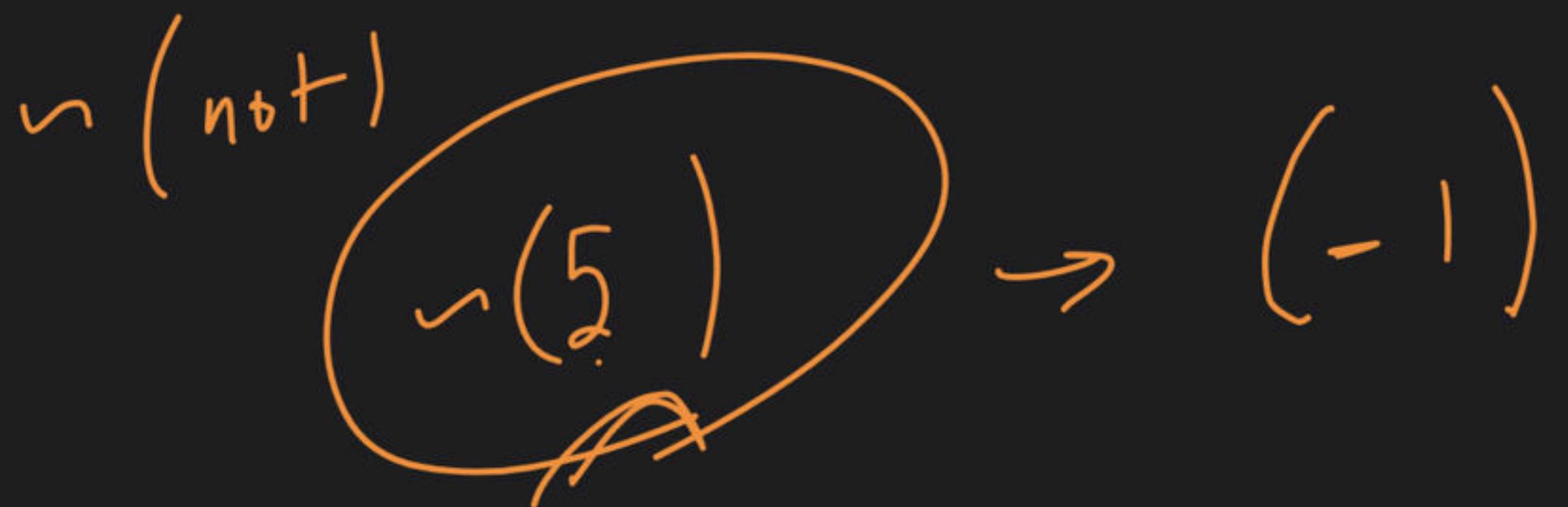
0 Even or Odd

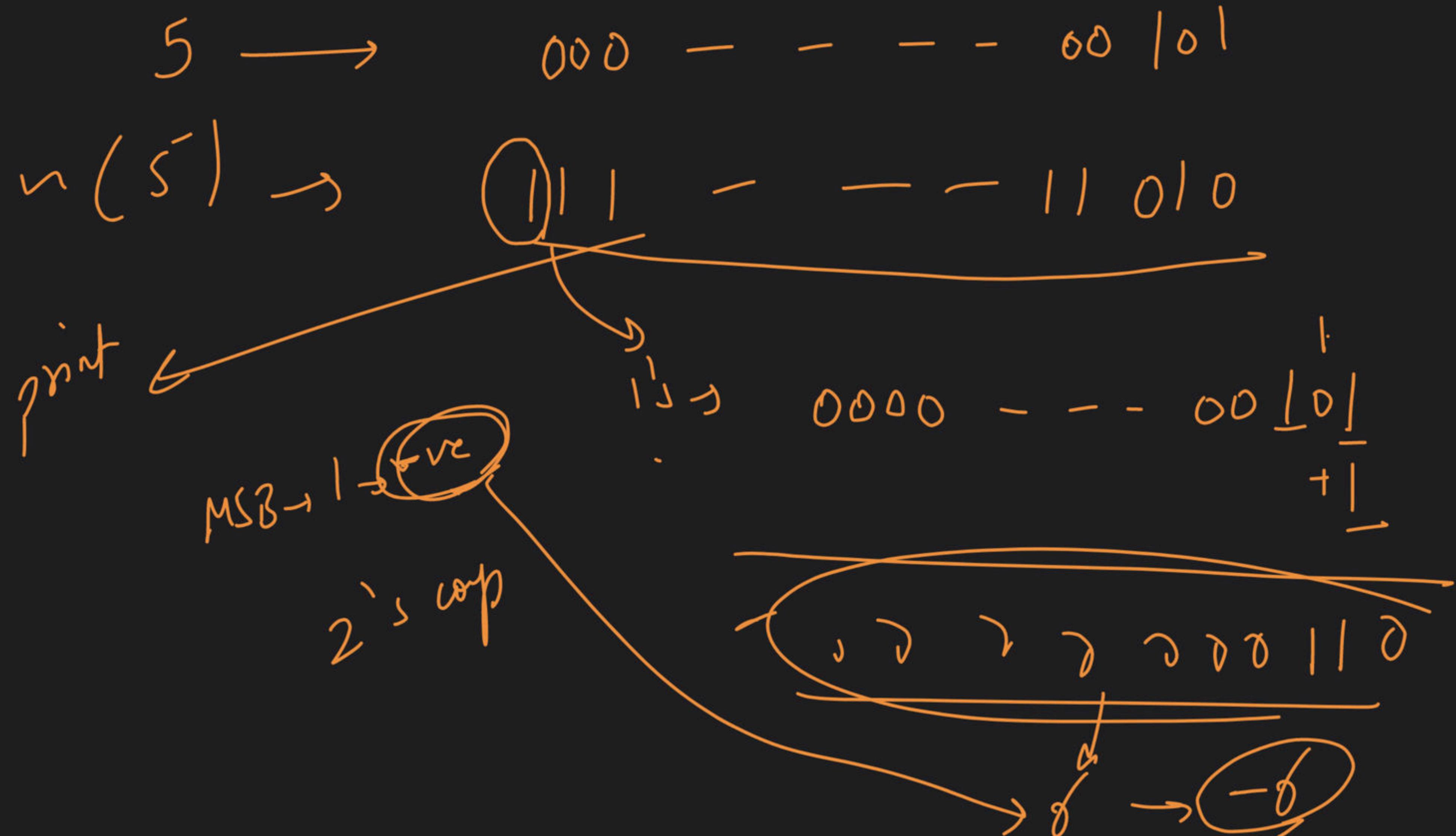
$\lambda I \rightarrow 0 \rightarrow \text{Even}$

$\lambda I \rightarrow [0 \rightarrow \text{odd}]$

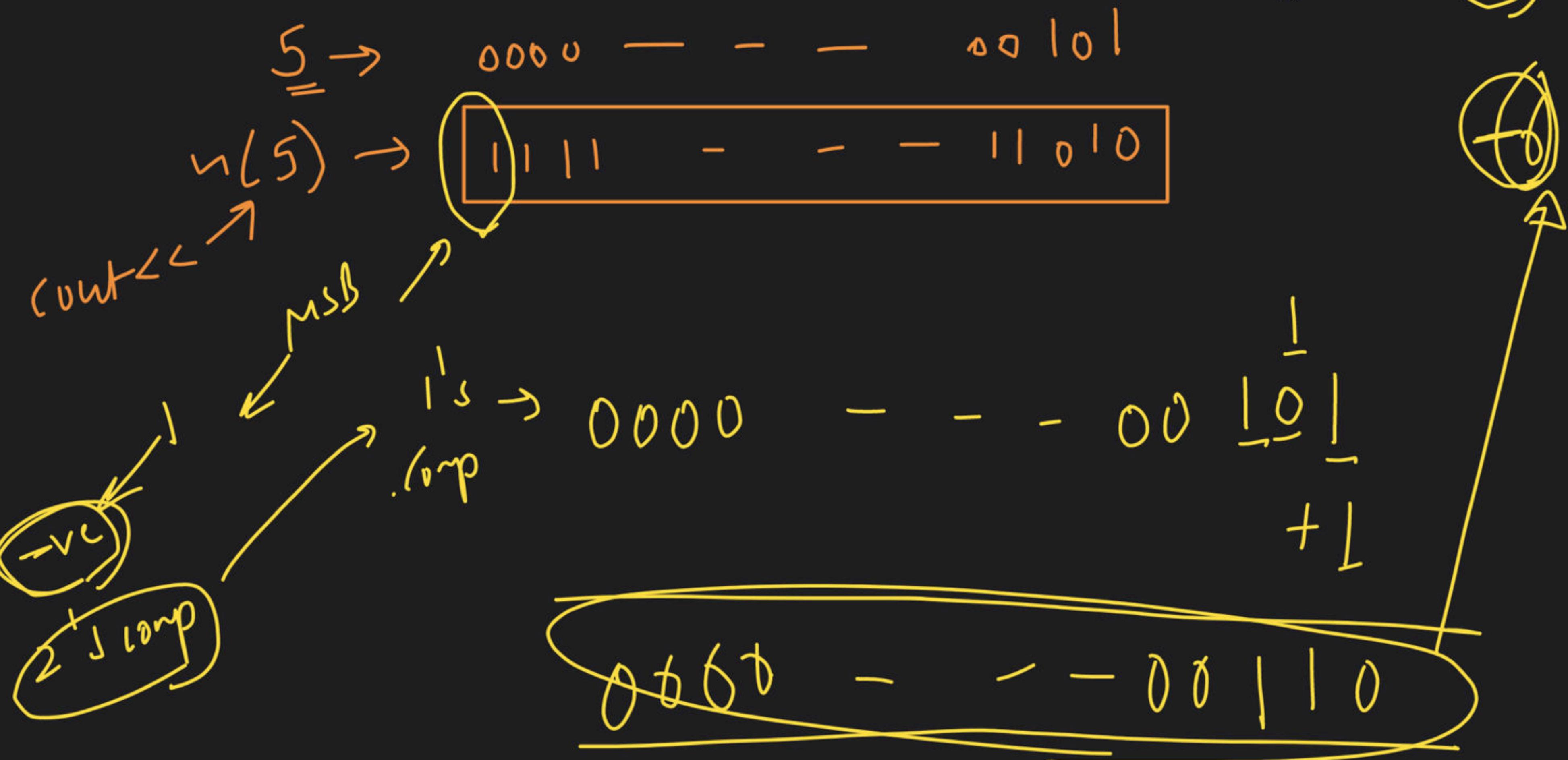
Note - Making Tool → 1 week







$\sim(5) \rightarrow$



XOR

8 ans \rightarrow 0

diff \rightarrow 1

find unique number

a	b	o/b
0	0	\rightarrow 0
0	1	\rightarrow 1
1	0	\rightarrow 1
1	1	\rightarrow 0

$$\frac{3}{7}, \frac{8}{1}, \frac{7}{1}, \frac{4}{1}, \frac{12}{7}, \frac{3}{1}, \frac{12}{1}, \frac{8}{1}.$$

7 ans

The diagram illustrates the assembly of a 32-bit word from four 8-bit bytes. The process involves shifting and concatenating the bytes.

Input bytes (labeled t):

- Byte 1: 00000000
- Byte 2: 00000000
- Byte 3: 00000000
- Byte 4: 00000100

Operations:

- Left Shift:** The first byte (00000000) is shifted left by 8 bits. This is labeled "left shift" with a circled arrow.
- Concatenation:** The result of the left shift is combined with the second byte (00000000). This is labeled "5 << 1" with a circled arrow.
- Left Shift:** The result of the previous step is shifted left by 8 bits again. This is labeled "left shift" with a circled arrow.
- Concatenation:** The result of the second left shift is combined with the third byte (00000000). This is labeled "5 << 1" with a circled arrow.
- Left Shift:** The result of the previous step is shifted left by 8 bits again. This is labeled "left shift" with a circled arrow.
- Concatenation:** The result of the third left shift is combined with the fourth byte (00000100). This is labeled "5 << 1" with a circled arrow.
- Final Result:** The final concatenated value is 00000000 00000000 00000000 00000100. The byte 00000100 is underlined, indicating it is the least significant byte.

b_1

$b_1 \ll 2$

0000 0000

5

0000 0000

0000 0000

0000 0010

$S \ll 2$

$Q \ll 2$

$S \times 2 \times L$

$= 20$

0000 0000

0000 0000

0000 0000

0001 0100

20

5 << n

n m

<< n

n ≠ 2ⁿ

5 * 2ⁿ

number = A

shift by n bits



$A * 2^n$

K 5

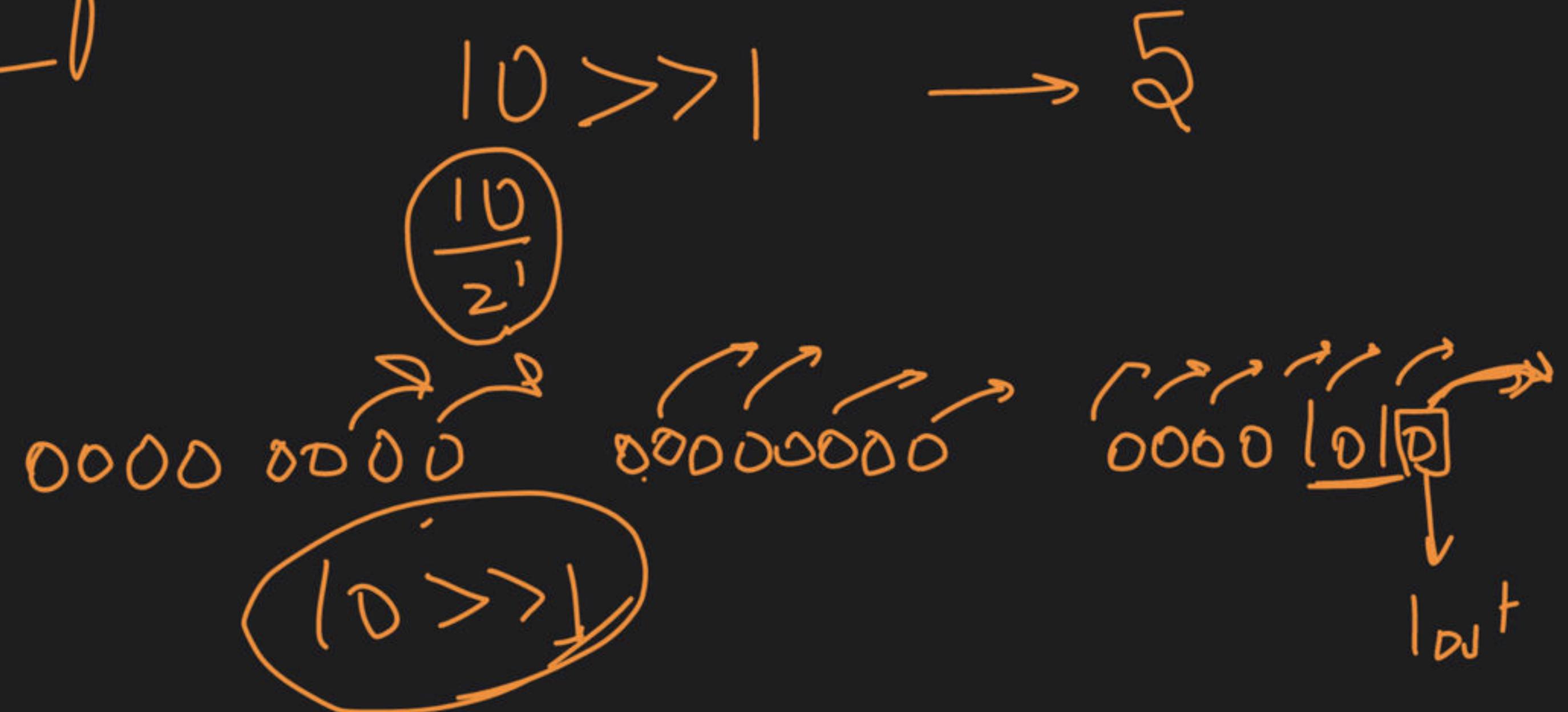
number = 5

Shift 3 bits

5×2^3

>> Right shift

10
10
with sign 0



00000000 00000000 00000000 00000101

5

$10 \rightarrow$

00000000

00000000

00000000

$10 >> 2 \rightarrow 2$

$$\frac{10}{2^2} = \frac{10}{4} = 2$$

000001010

$>> 2$

00000010

00000000

00000000

00000000

$\swarrow 2$

number $\rightarrow A$

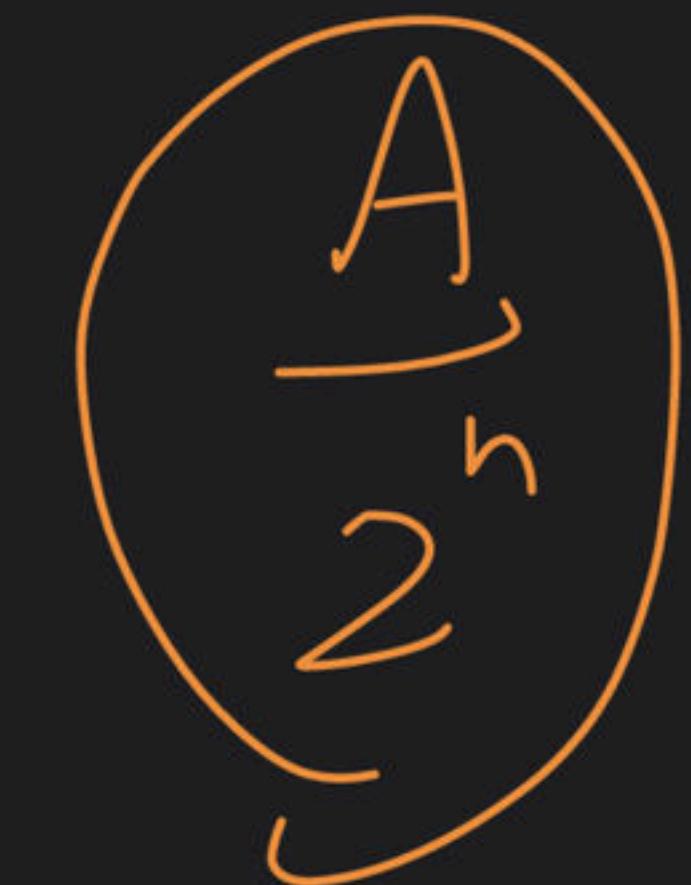
number = 20

right shift by n bits

R.s. by $\rightarrow 3$ bits



$$\begin{aligned} 2 &\gg 2 \\ \Sigma \quad : \quad \Sigma - 1 &= \frac{\Sigma - 1}{2} \\ 2 & \end{aligned}$$



$$\frac{20}{2^3} = \frac{20}{8}$$

$$= 2$$

7

5

000000

0001



Count of
set bit \rightarrow



5

0000

- - - -

0000



&1

—
—
—
→

set
bit

(0+1)



000

-

—

— 000



&1

→ x



000

—



0000

&1

0+1+1



3 ↳

0000

- - - -

>> |

0000 0

- - - -

loop

>> |

0000

- - - -

0000 0

= 0 → Bulk gate

00 00 11

& |

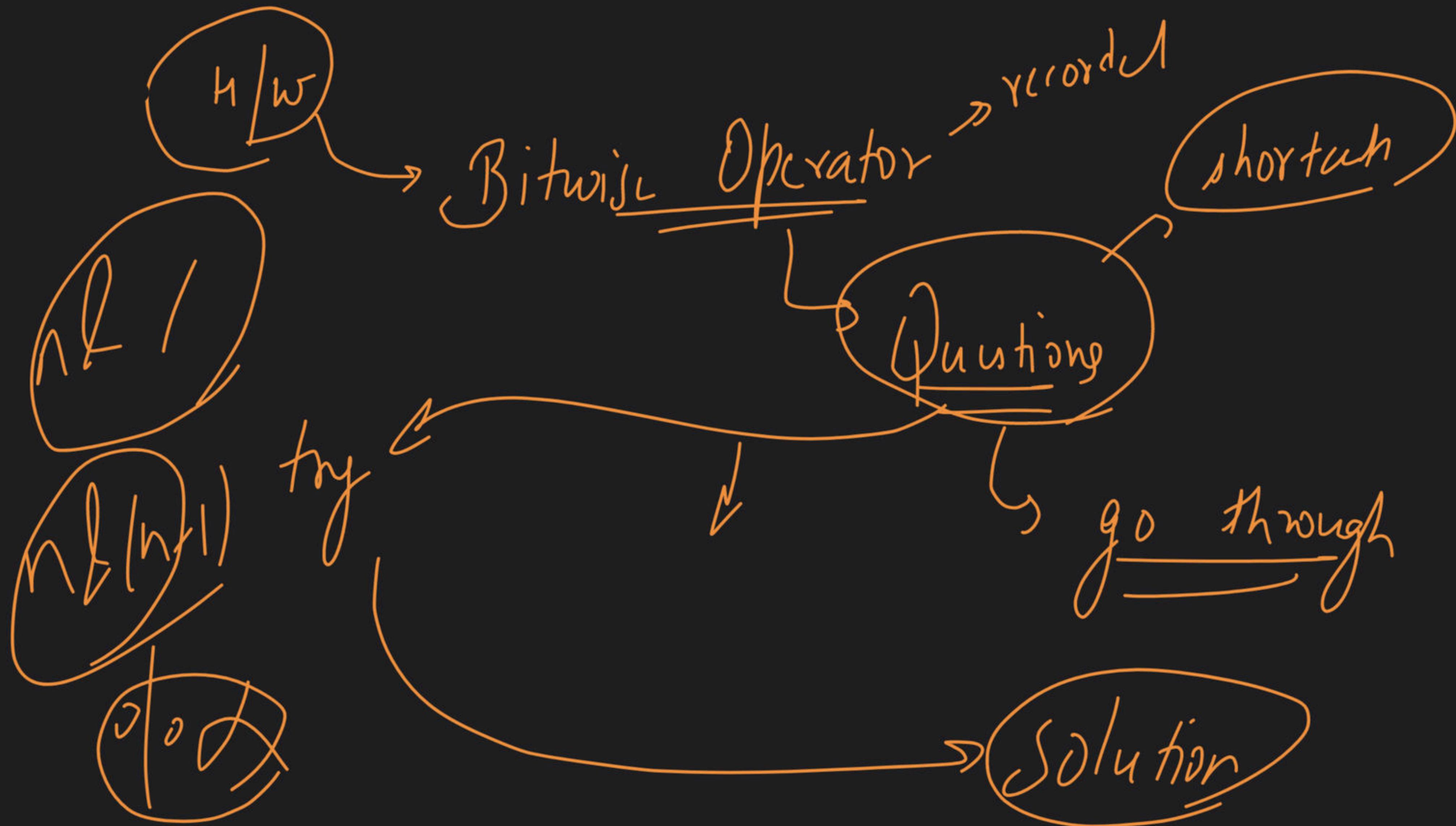
00 00 01

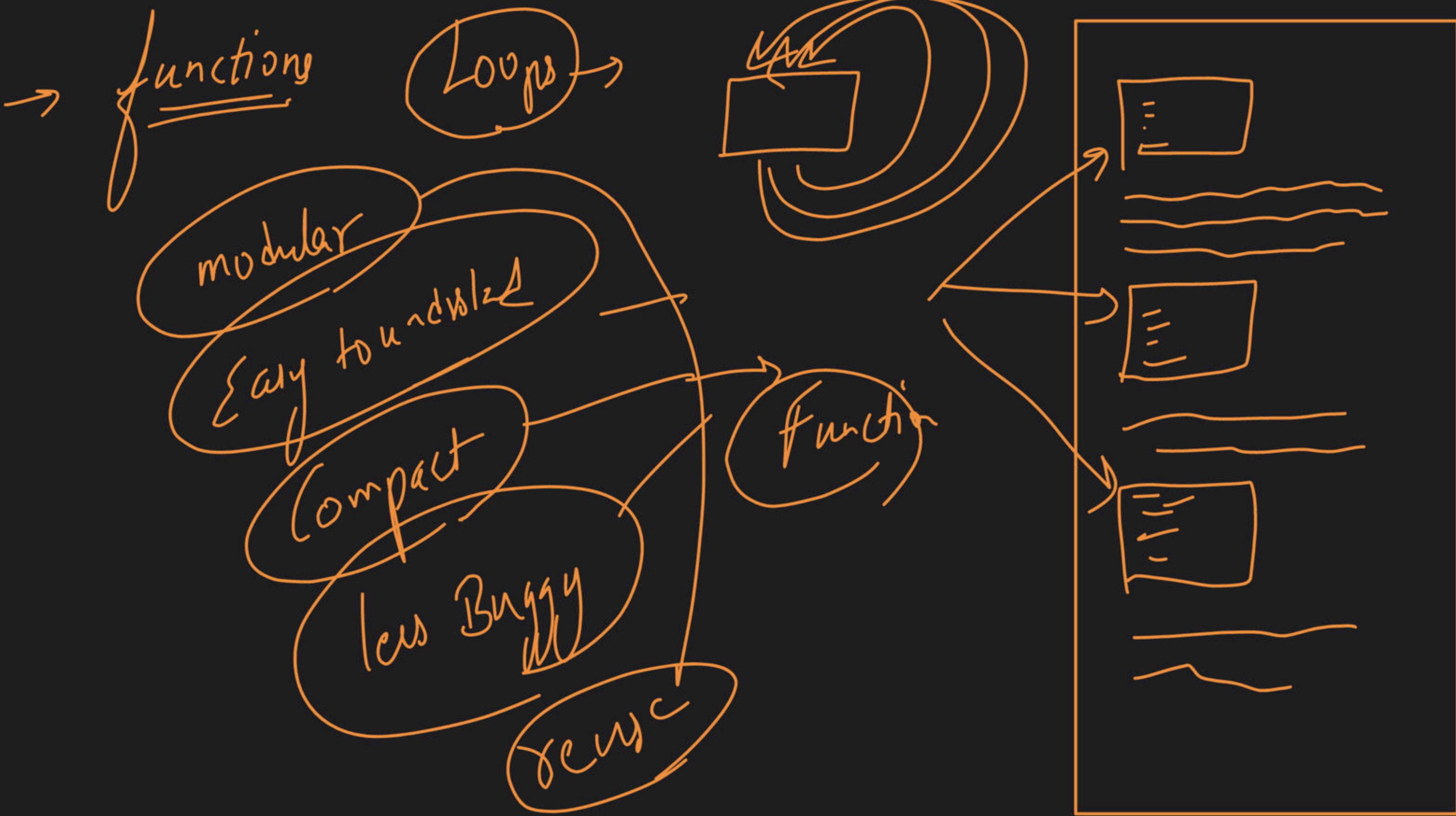
↓

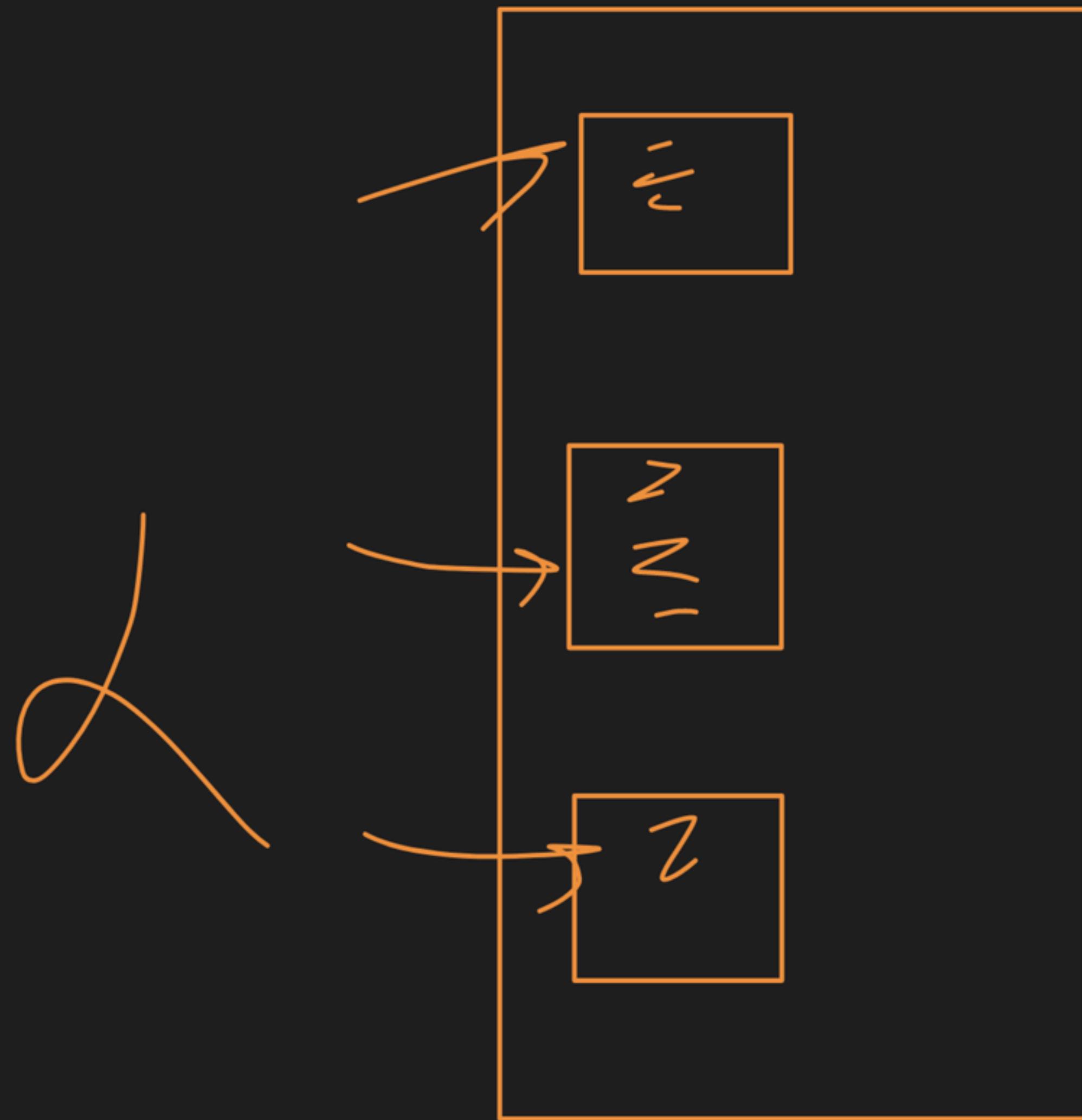
Ist set bit

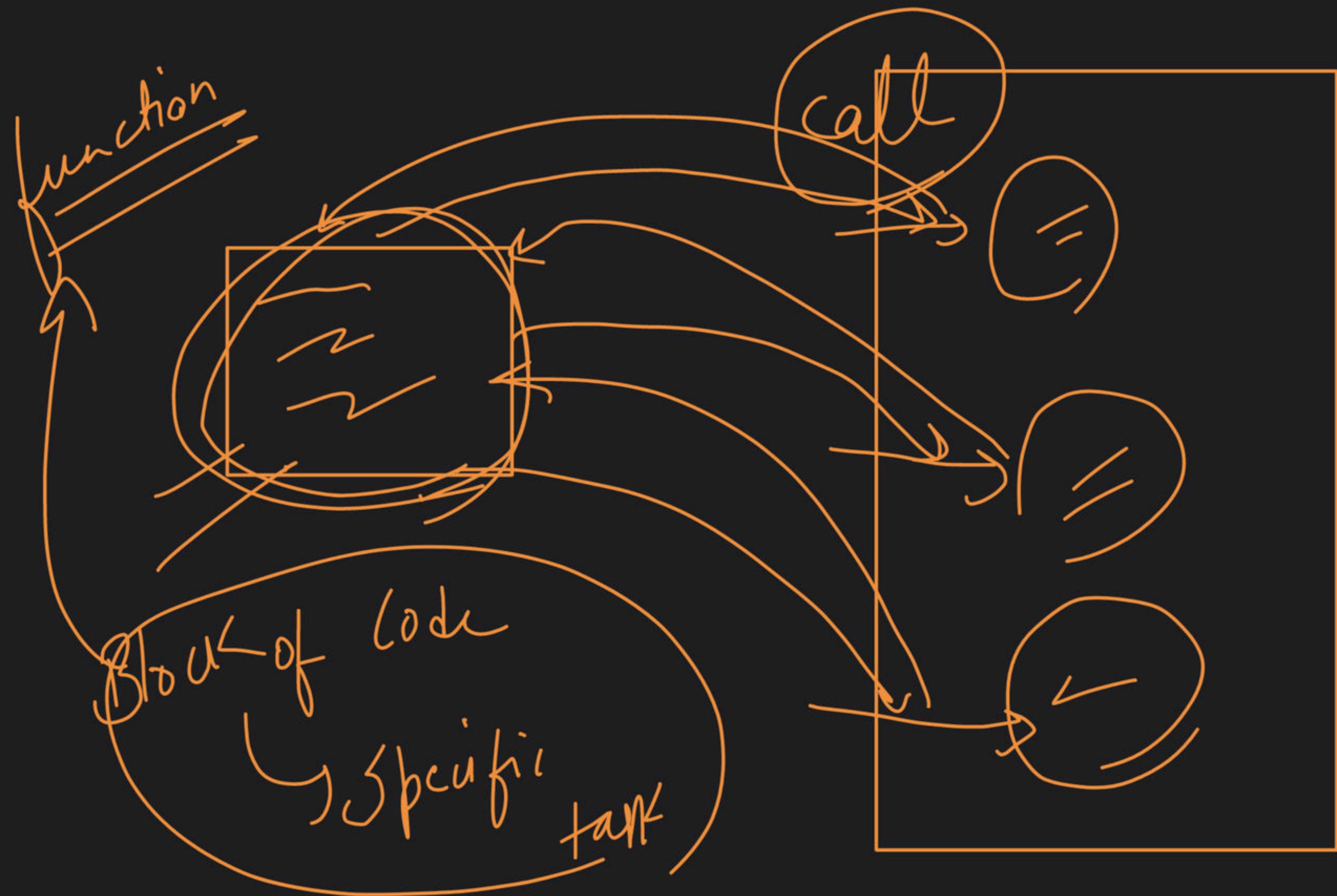
IInd set bit

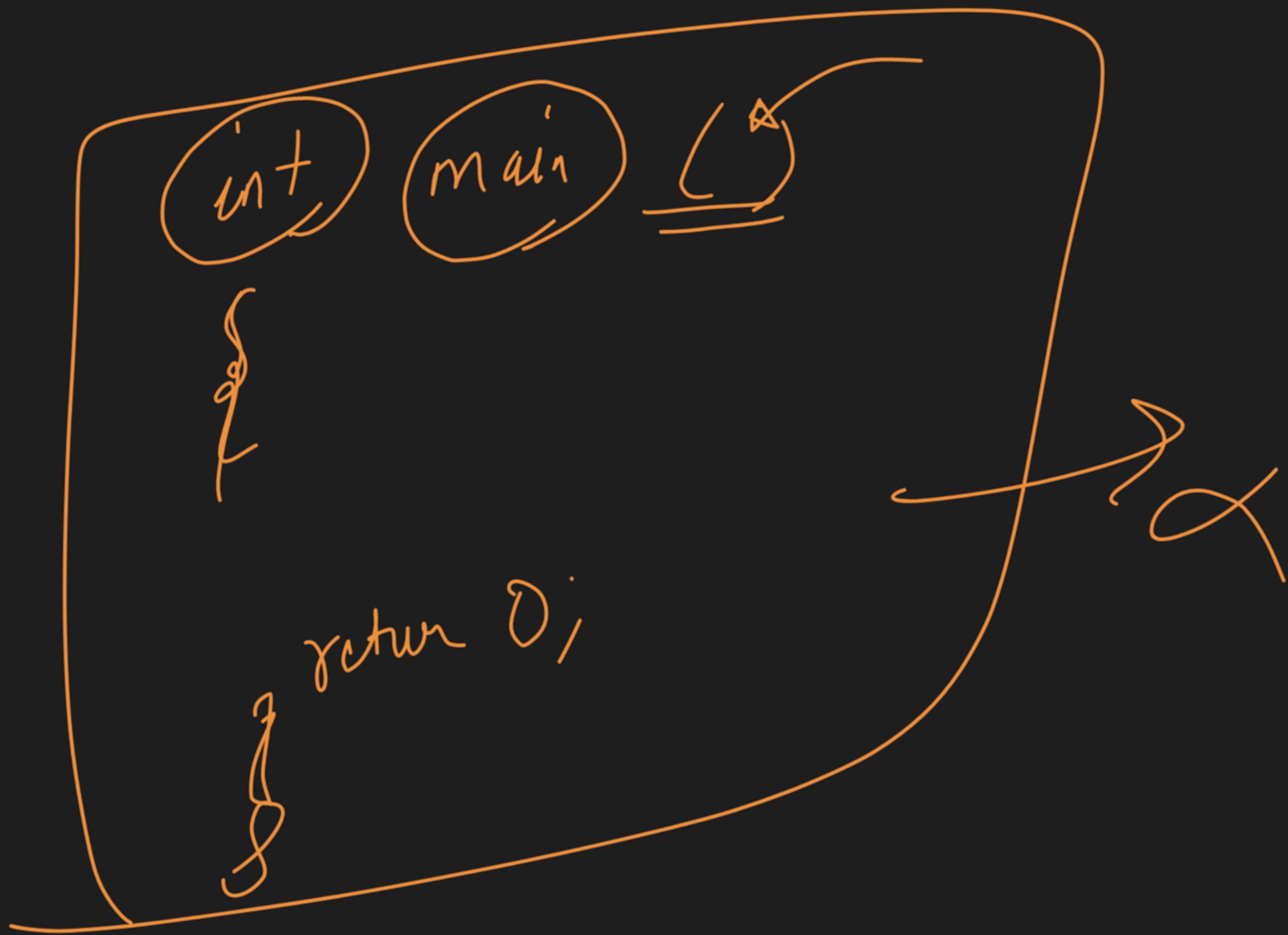
28th bit



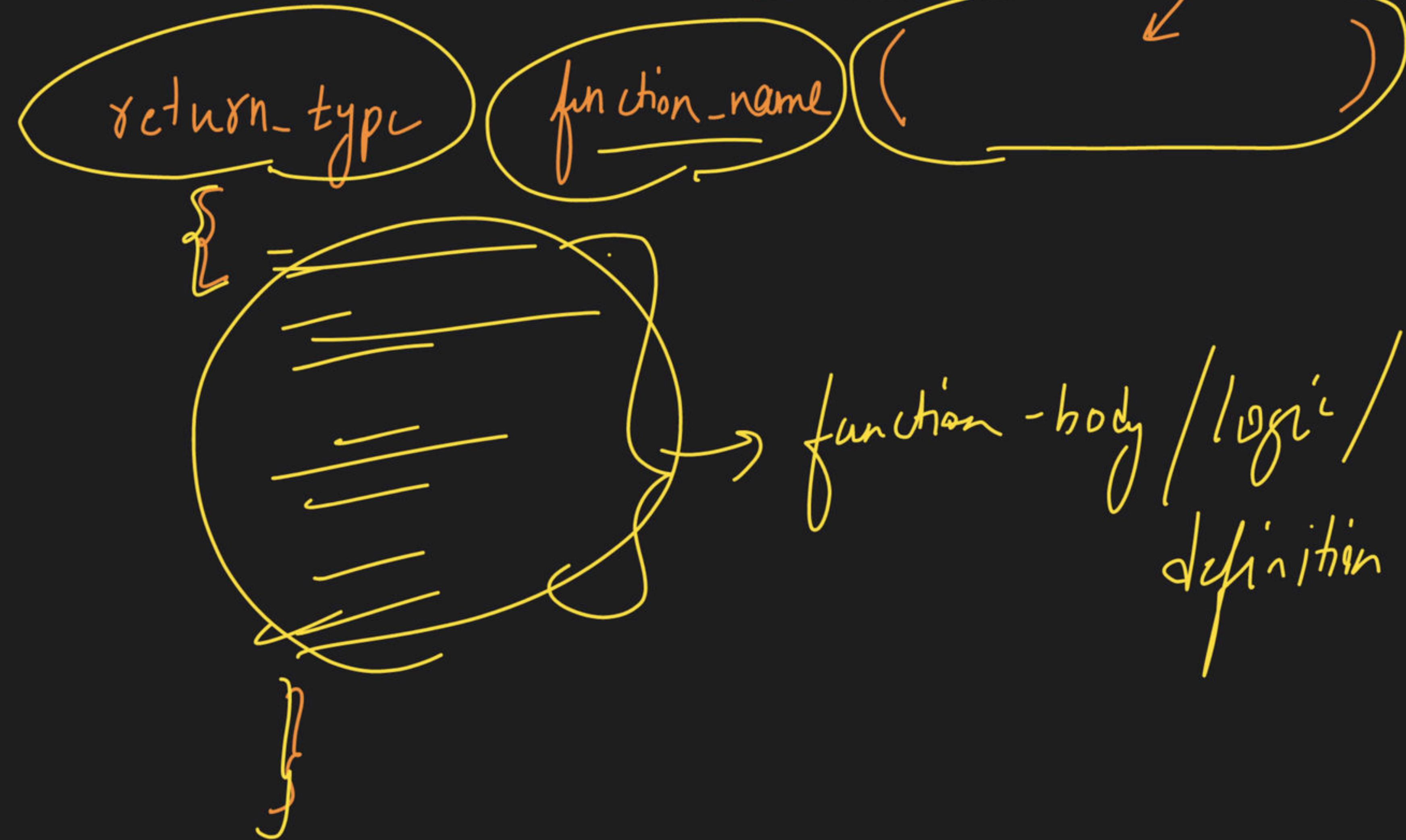




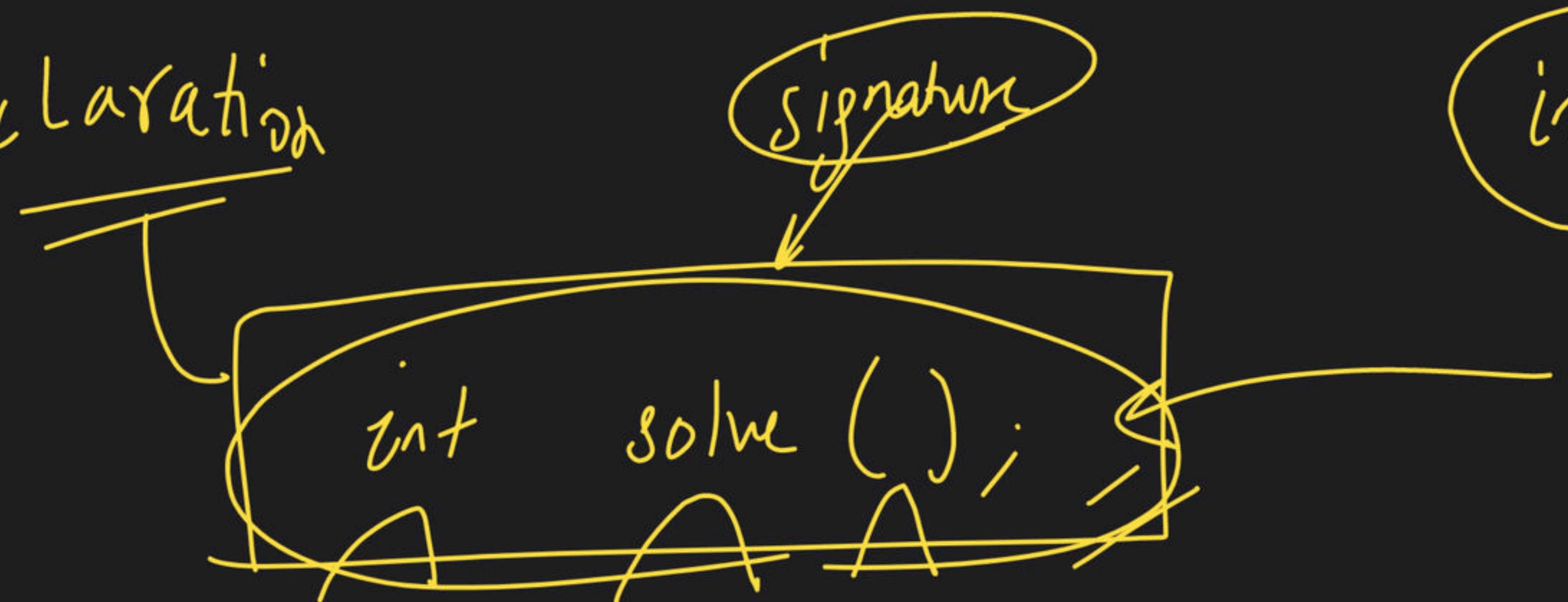




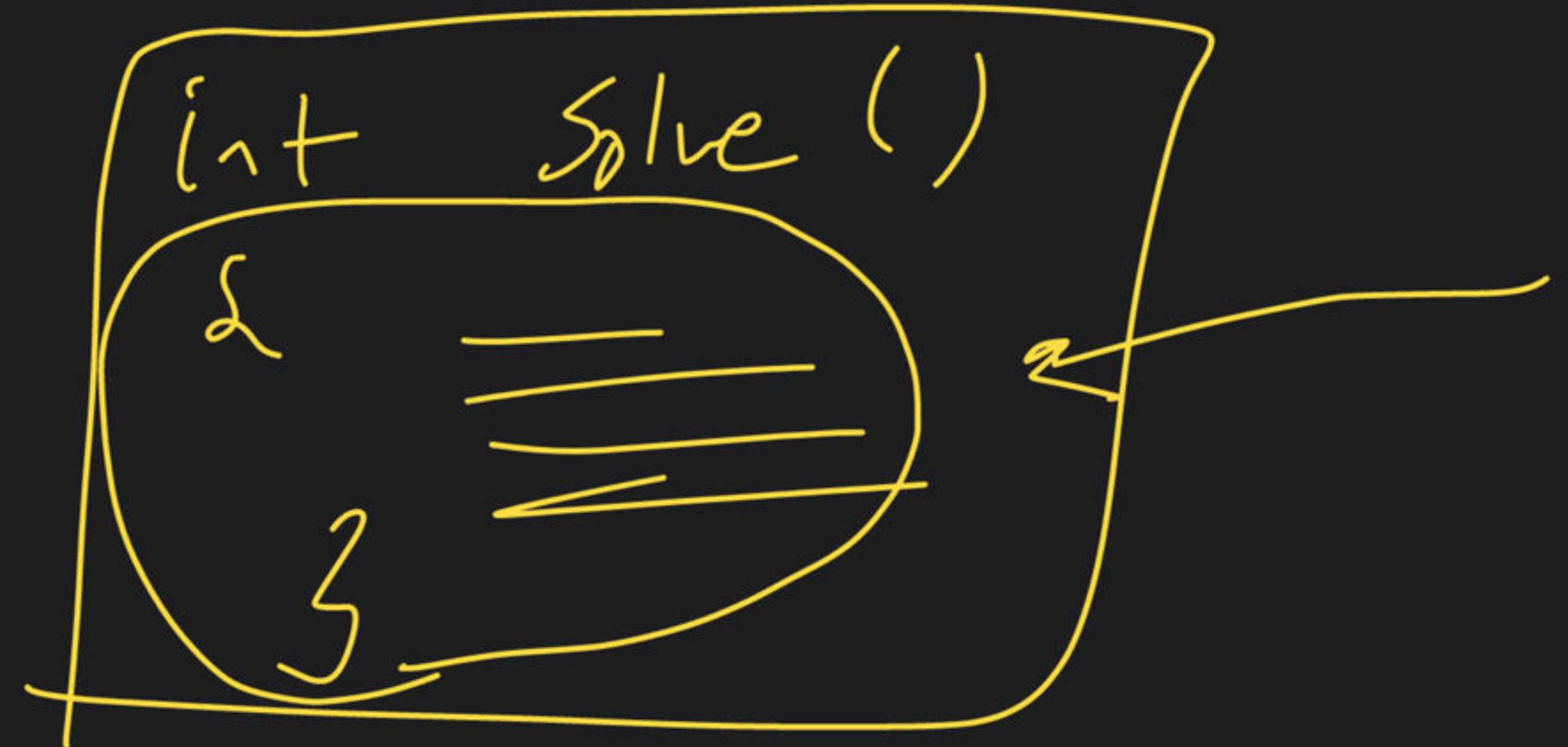
Syntax:-

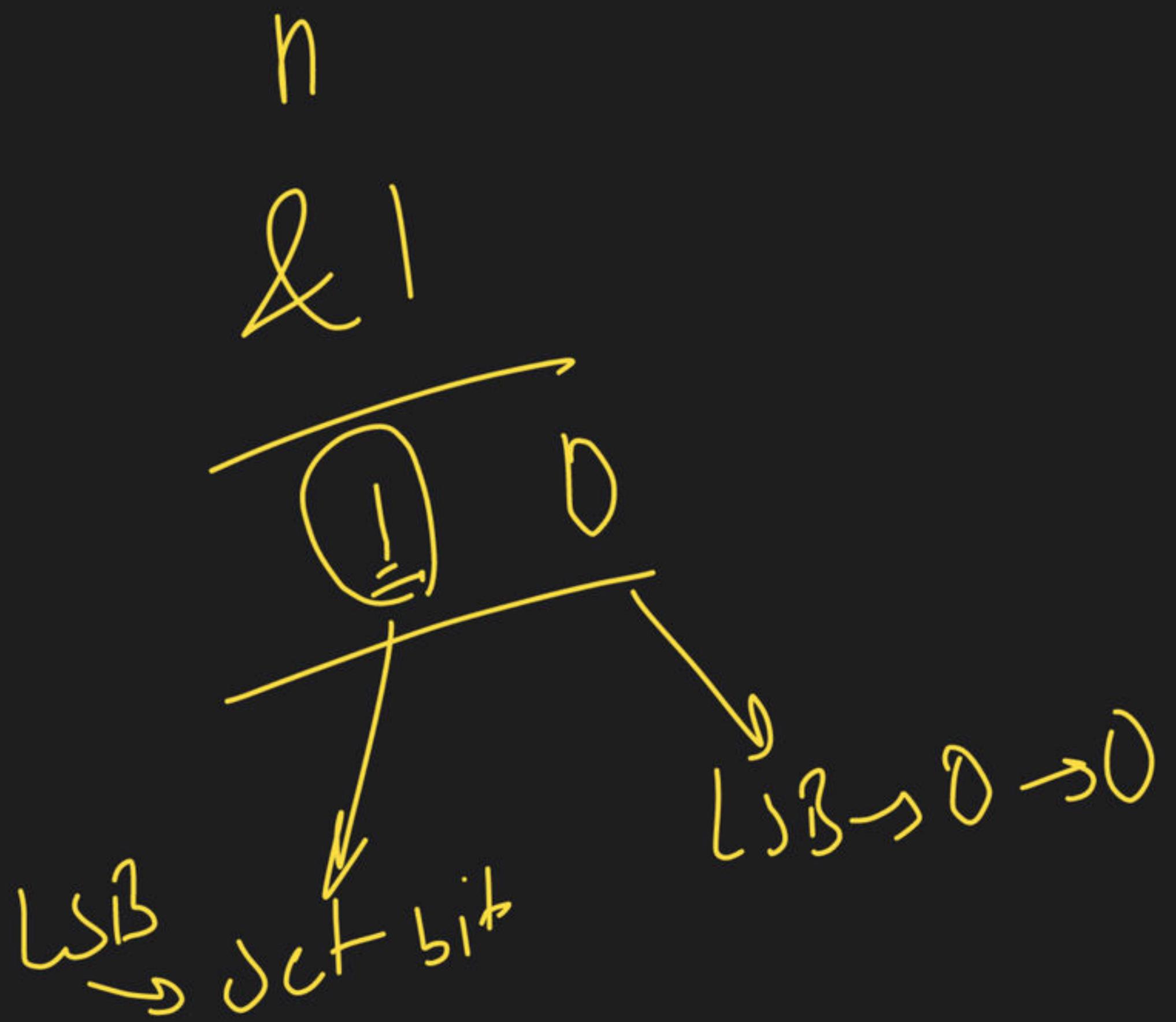


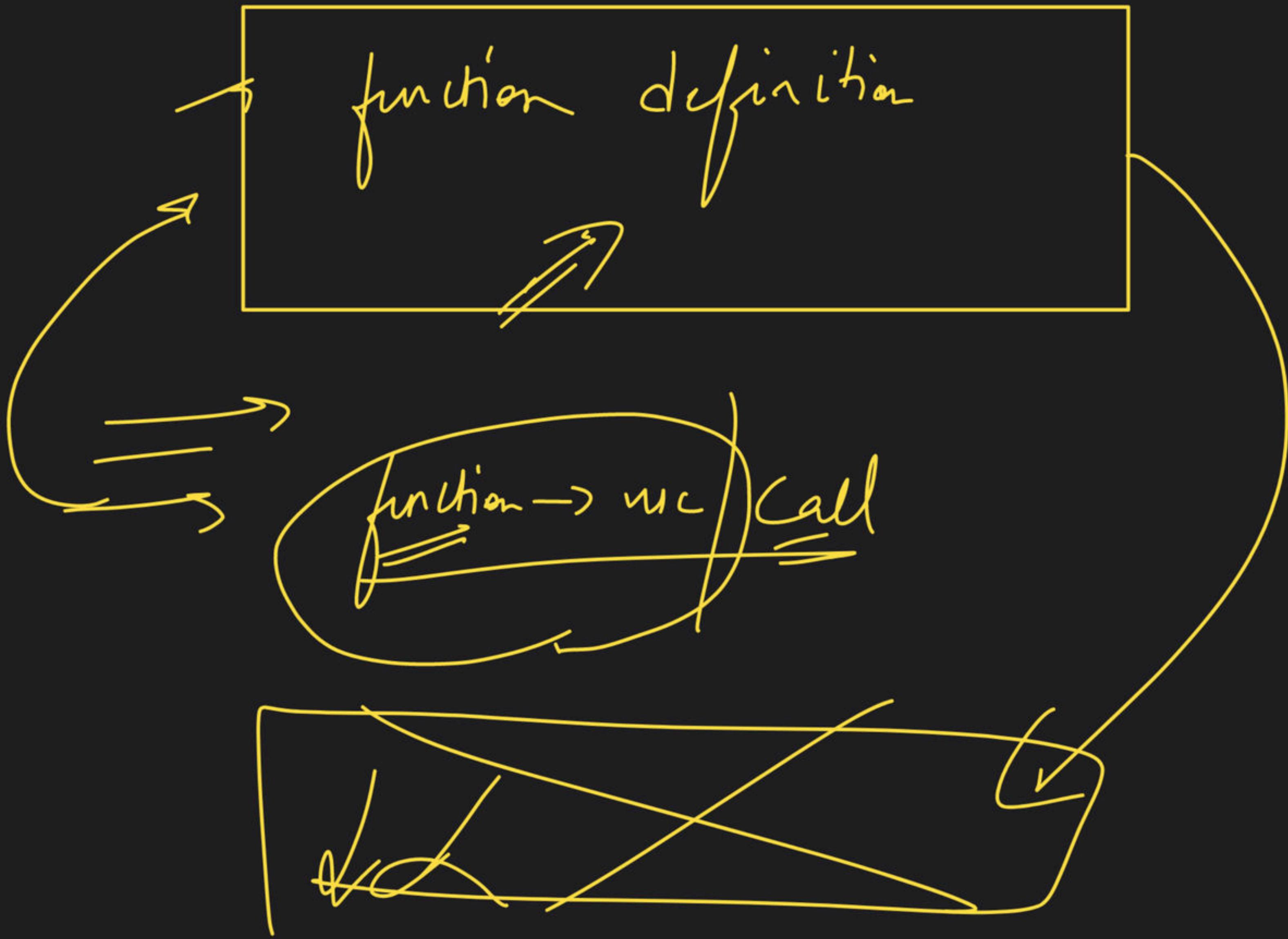
declaration



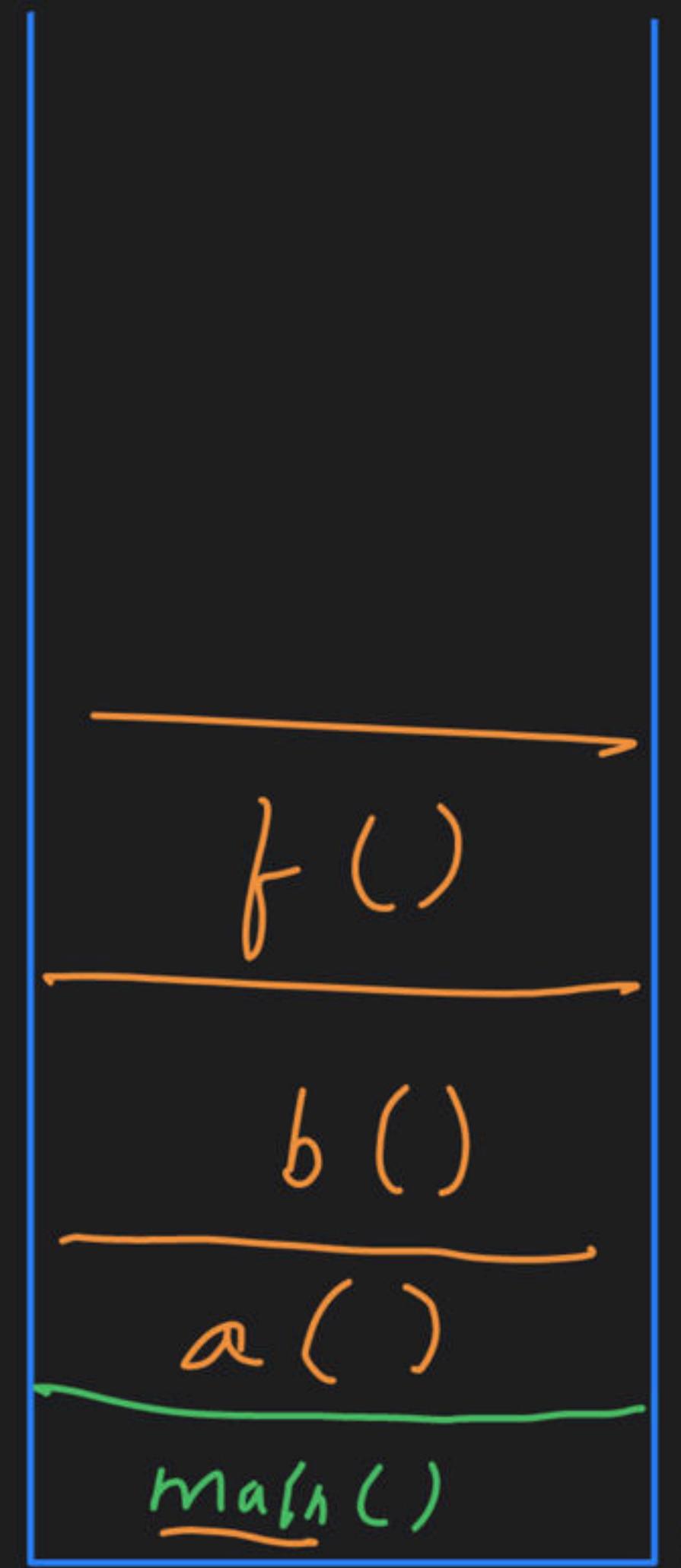
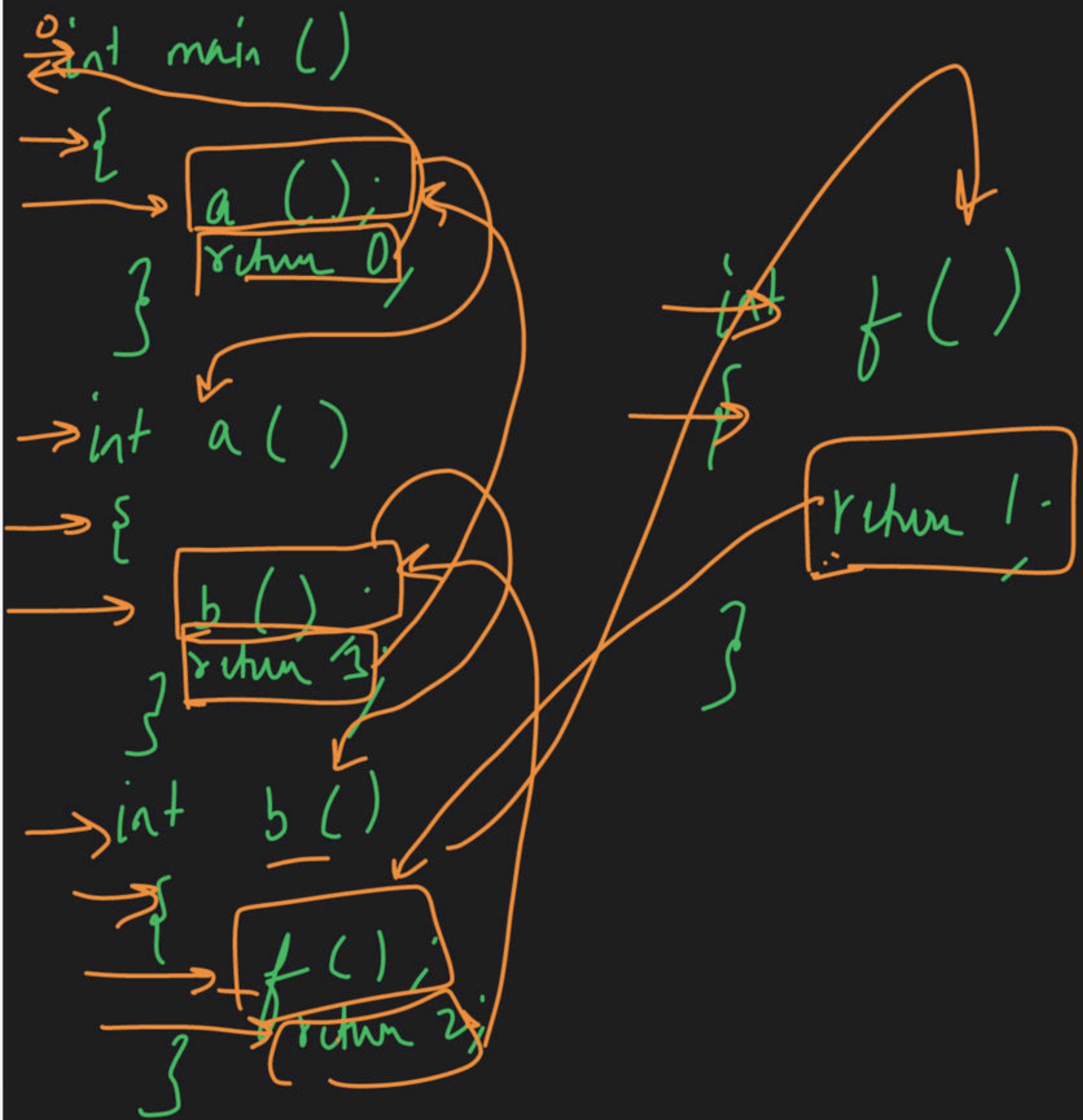
definition



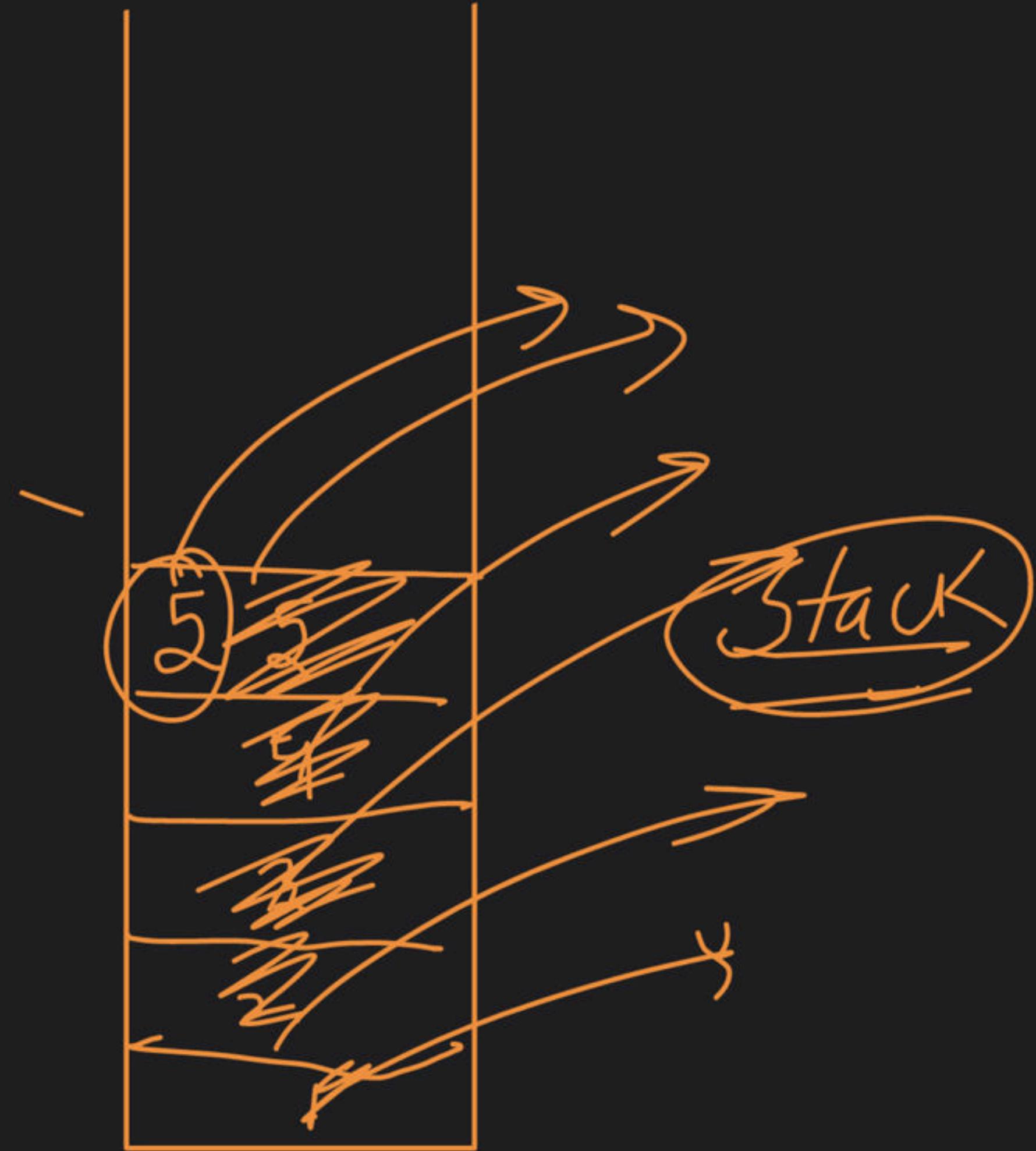
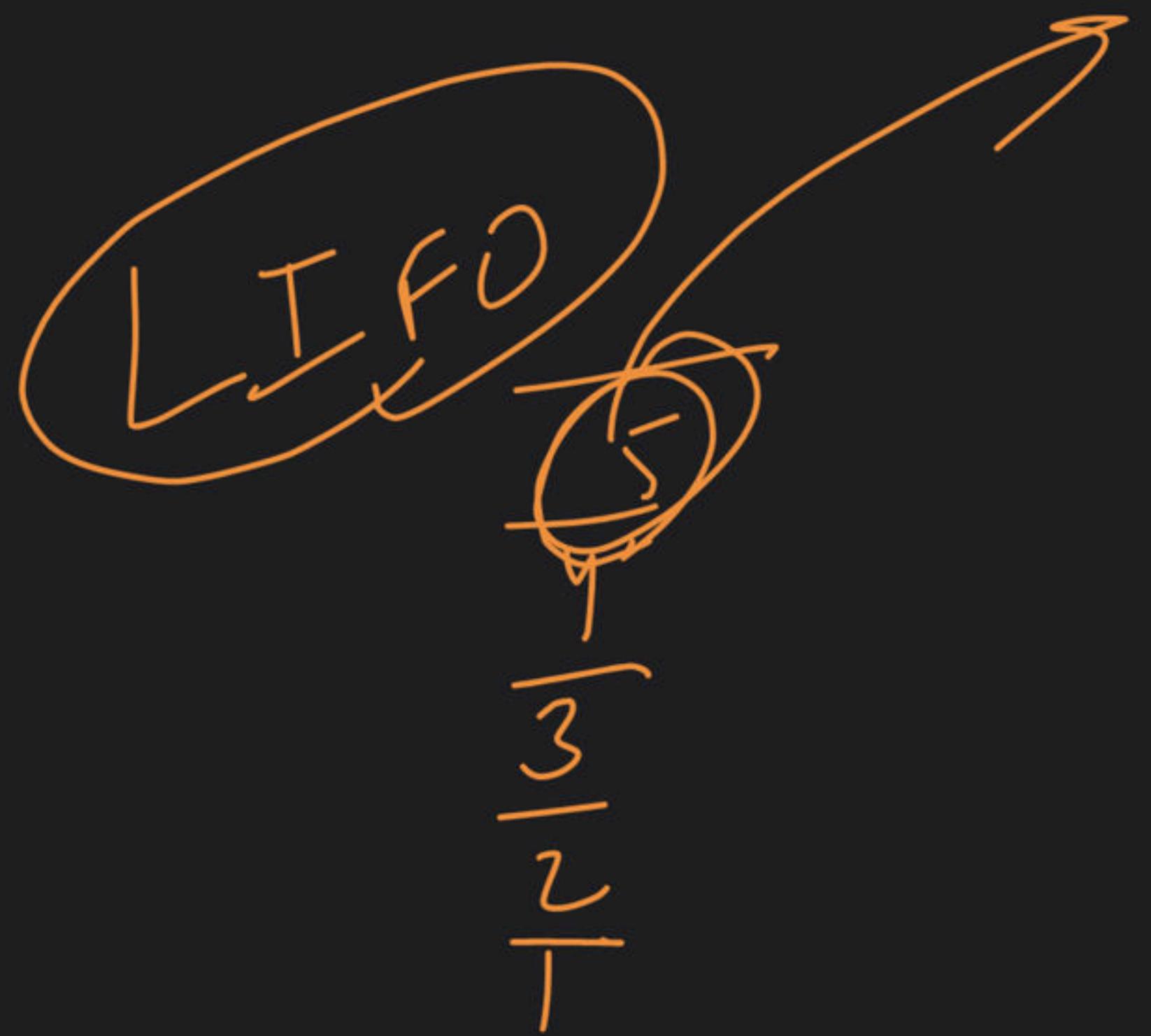




```
① void SayMyName ()  
② {  
③     cout << "Babbar";  
④ } function ends here  
⑤ int main ()  
⑥ {  
⑦     sayMyName (); function call  
⑧     return 0;  
⑨ }
```



Call
function Stack



function call

int main()

int a = 5;

int b = 10;

int ans =

~~solve(a, b);~~

cout << ans;

return 0;

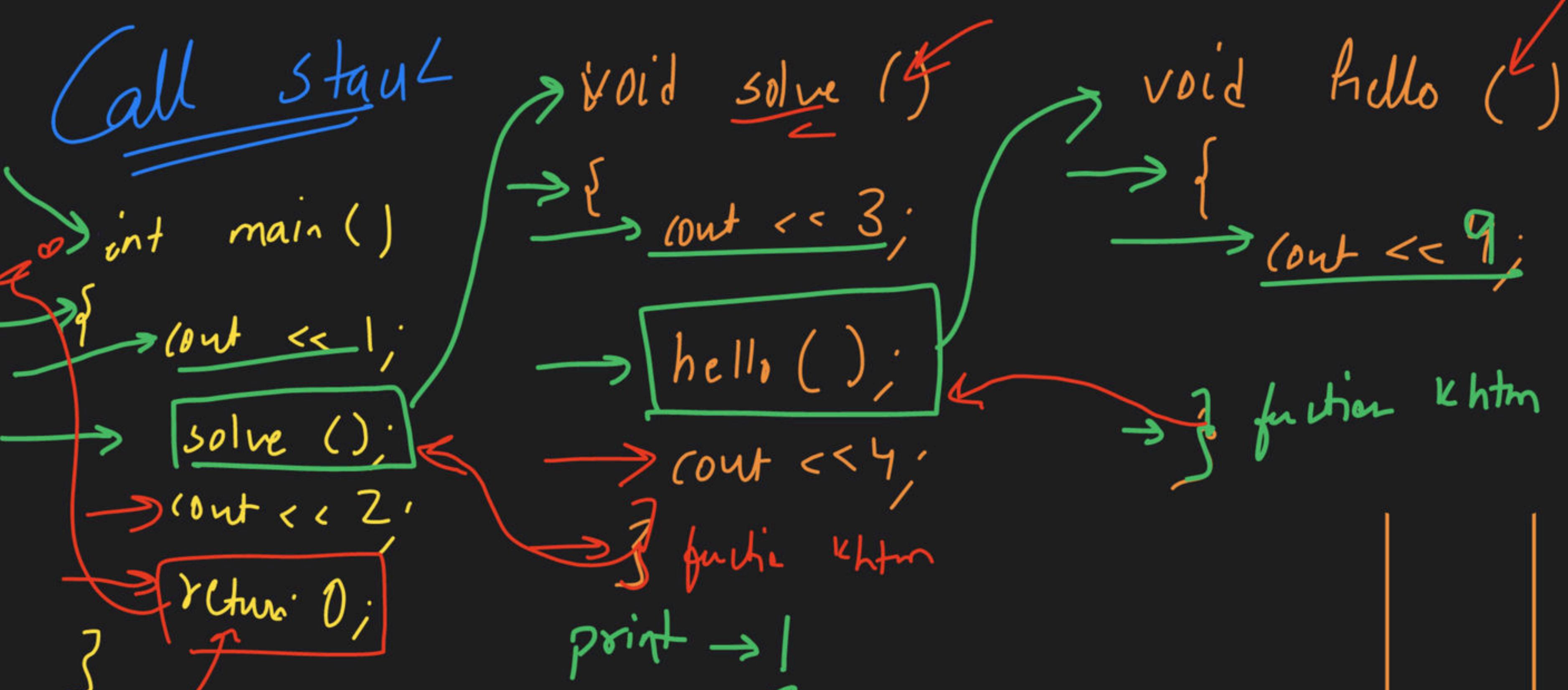
int solve

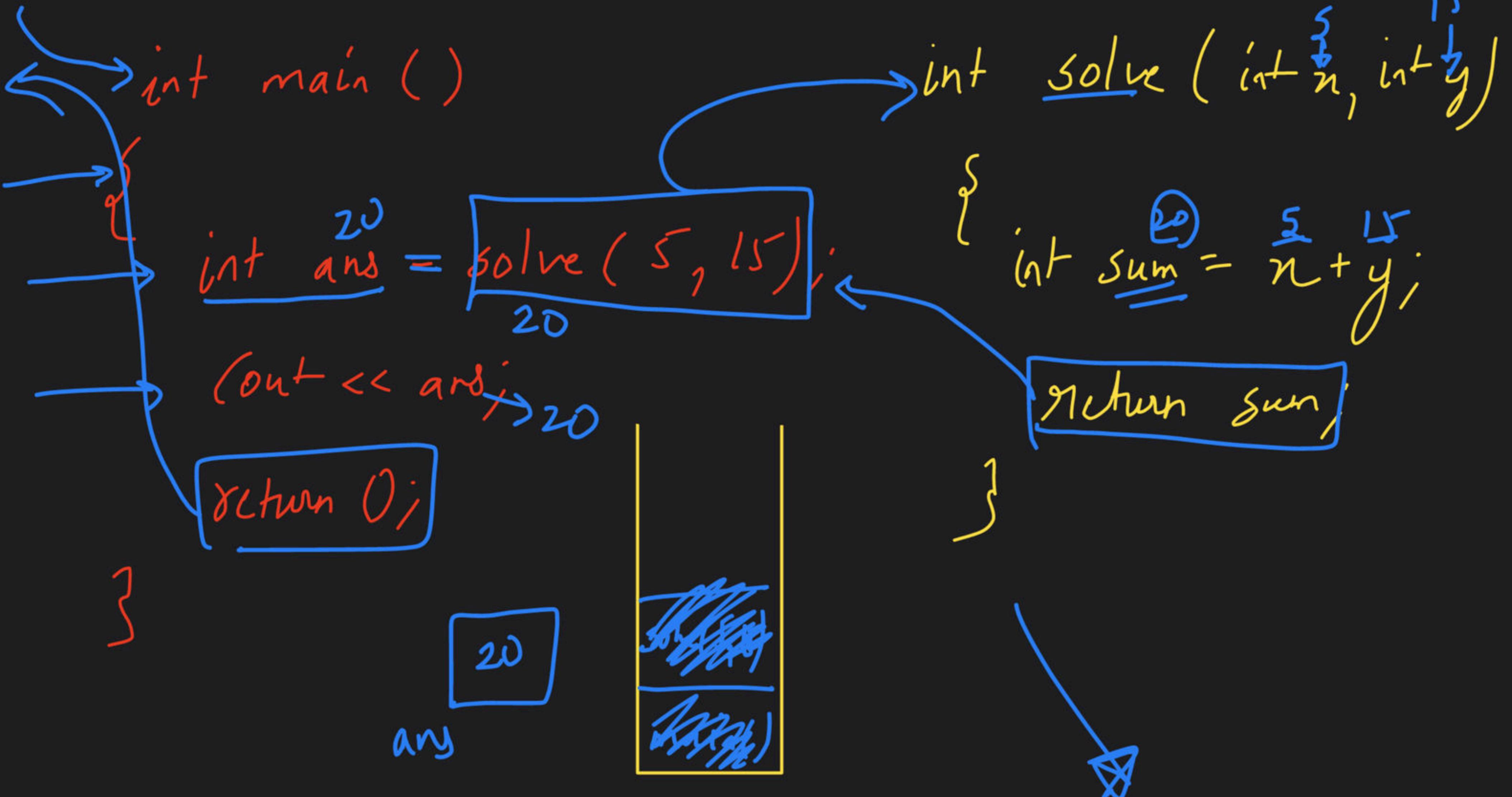
f.

int sum = a + b;

return sum;

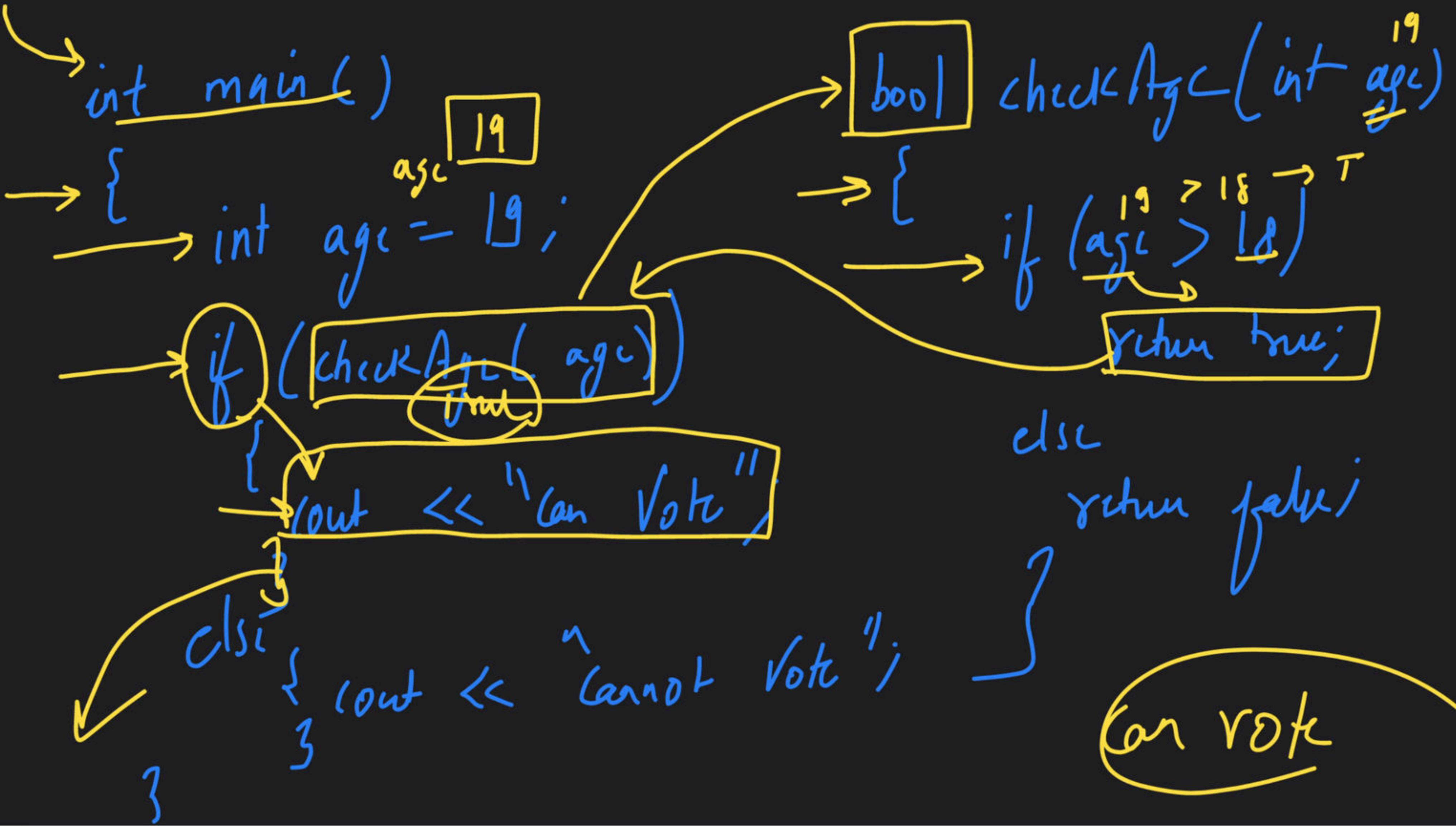
ans





dynamic memory allocation





$$a^b$$

int main

Sum of AP

2, 4, 6, 8, 10

$$n=5$$

$$a=2$$

$$l=10$$

$$\text{Sum} = \frac{n}{2} (a+l)$$
$$= \frac{5 \times 12}{2} = \frac{60}{2} = \frac{5}{2} (2+10)$$
$$= 30$$

$n \rightarrow$ no. of terms

$a \rightarrow$ first term

$l \rightarrow$ last term

Prime no check



$i/p \rightarrow n \rightarrow 12$

$n = 12 \rightarrow 15$

(2 → 1)

1, 15

True → prime

false → not prime

sum → 0 → not prime
!v → prime

~~*~~
-2
-3
-4
~~*~~

$$5 \cdot 1 \cdot 2 \rightarrow 1 \\ 5 \cdot 1 \cdot 3 \rightarrow 2 \\ 5 \cdot 1 \cdot 4 \rightarrow 1$$

$n = 2$ prime

~~X~~
 $n=6$

2

$$6 \cdot 1 \cdot 2 \rightarrow 0$$

not
prime

3

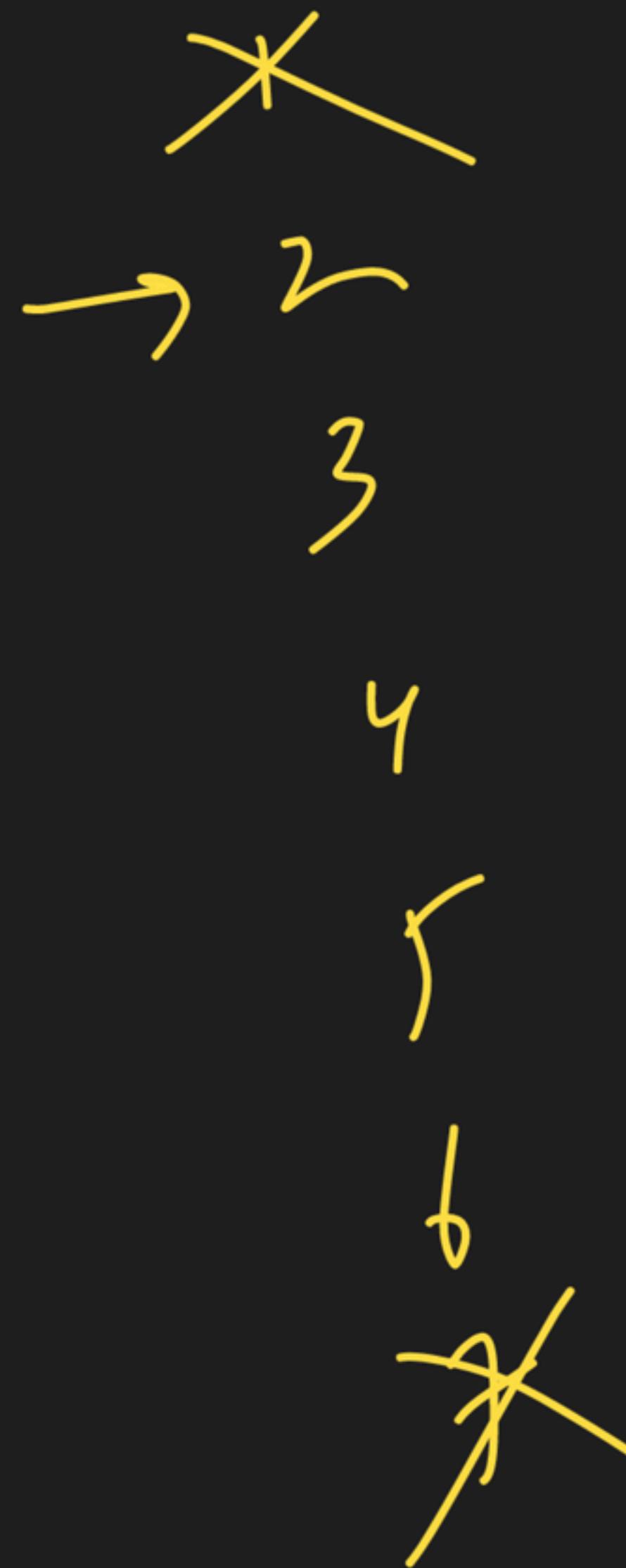
4

5

~~X~~

$n=7$

prime
no



$$7 \cdot 1 \cdot 2 \rightarrow 1$$

$$7 \cdot 1 \cdot 3 \rightarrow 1$$

$$7 \cdot 1 \cdot 4 \rightarrow 3$$

$$7 \cdot 1 \cdot 5 \rightarrow 2$$

$$7 \cdot 1 \cdot 6 \rightarrow 1$$

~~x~~

2

3

4

$\delta \cdot 1 \cdot 2 \rightarrow 0$

$h = \delta$

not prime

~~y~~

~~*~~

-2 $\rightarrow 9 \cdot 1 \cdot 2 \rightarrow 1$

3 $9 \cdot 1 \cdot 3 \rightarrow 0$

1

5

6

7

~~8~~

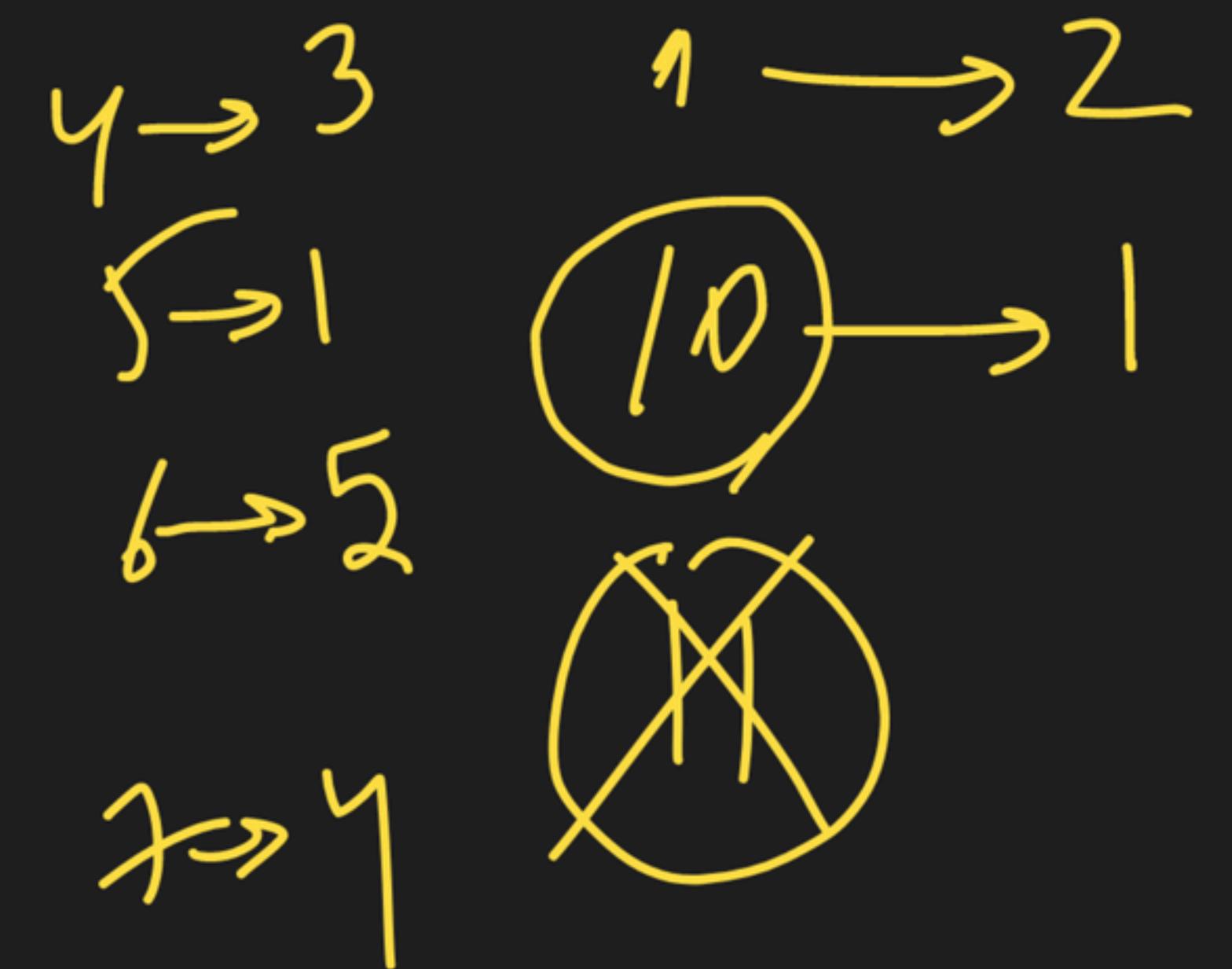
$n = 9$

not prime



$h=11$

$3 \rightarrow 2$ $8 \rightarrow 3$



$\alpha^0 \rightarrow \rho^{\text{dim}}$

$\alpha \rightarrow \gamma_{\text{main}}$

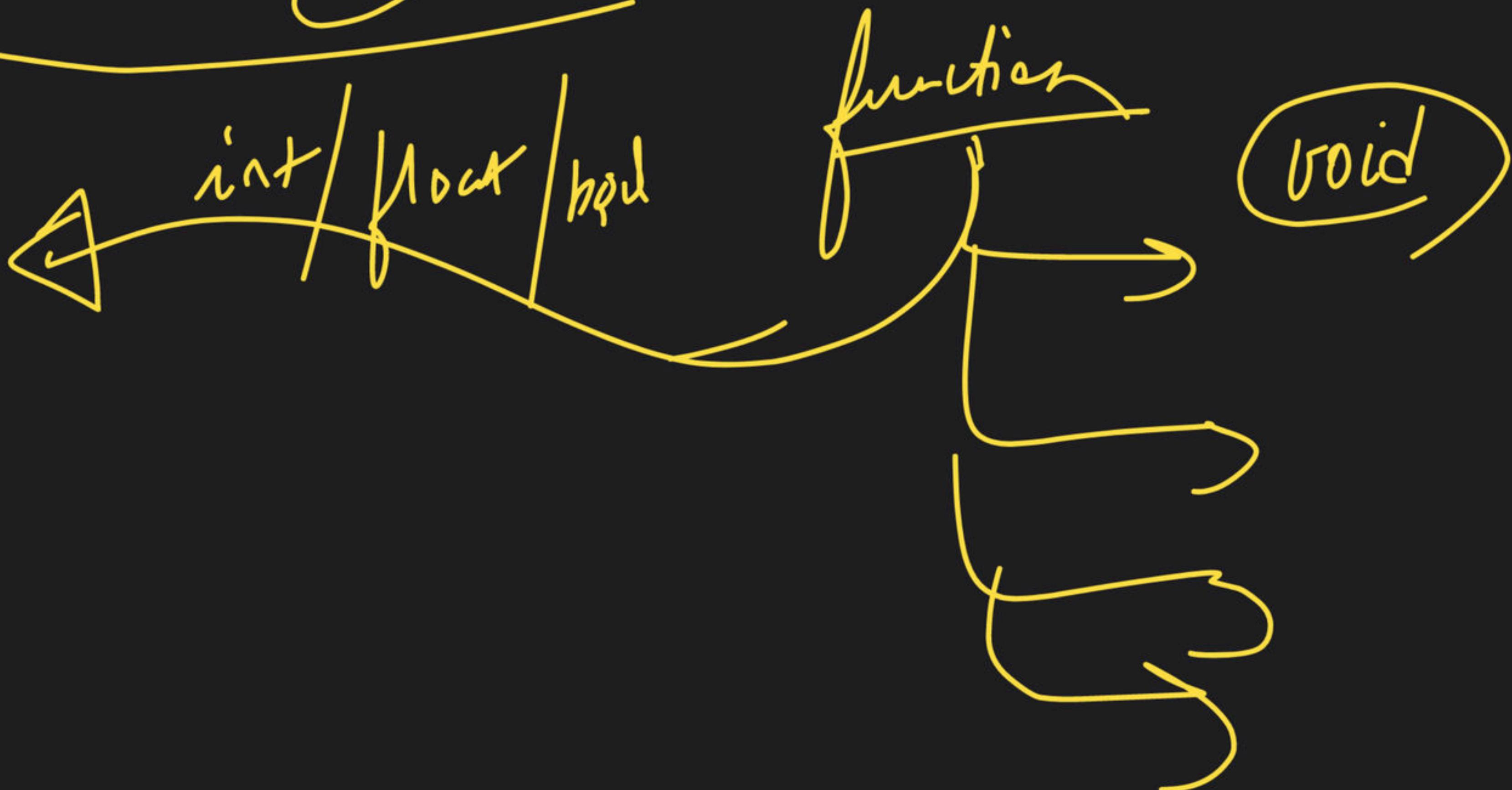
$n = 13$

for ($i = 2 \rightarrow \leq (n-1)$)



Odd no. →

2 min Break



-ve

0 1 1 1

-100 → -50

0010

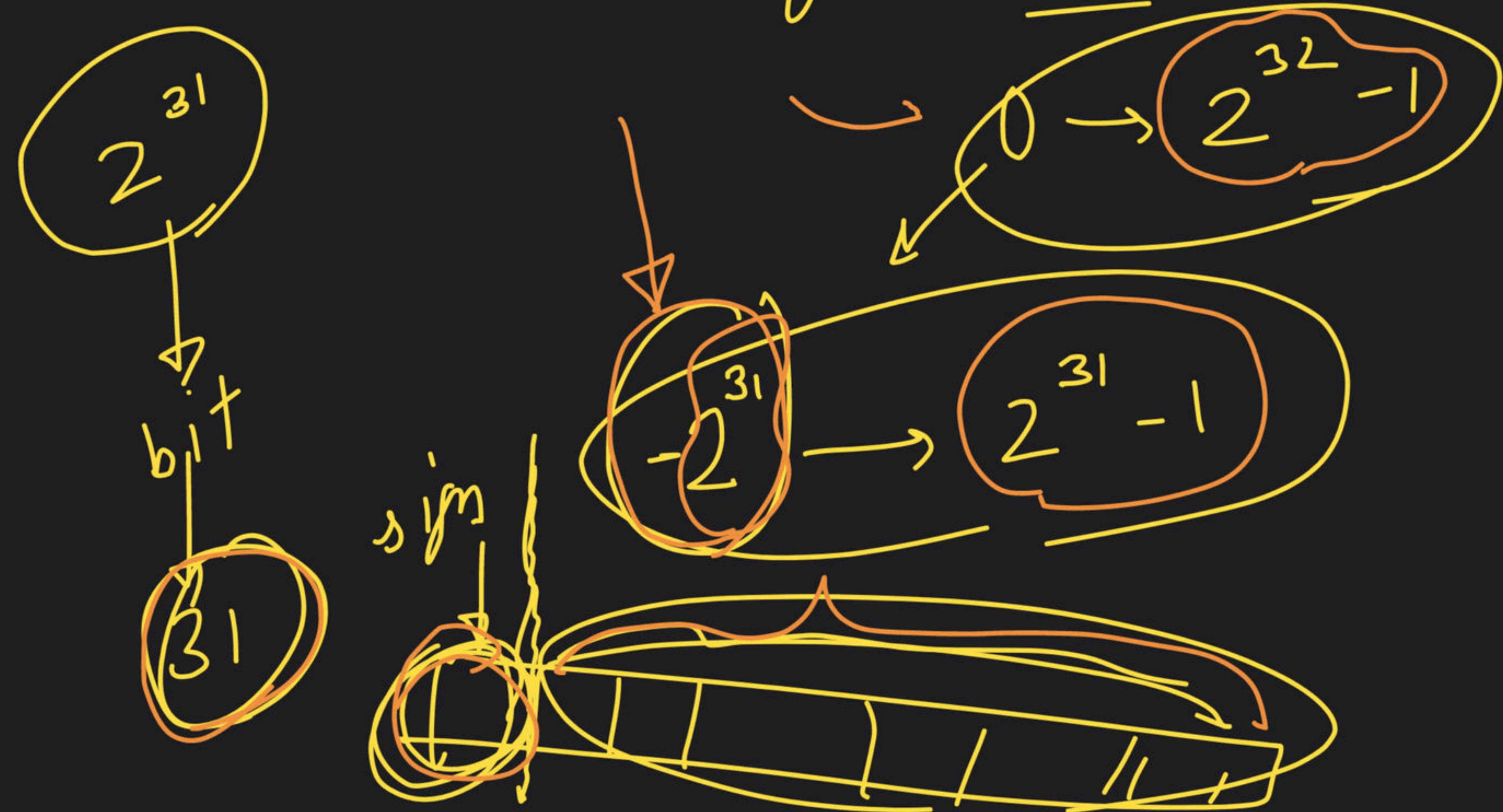
>> 1

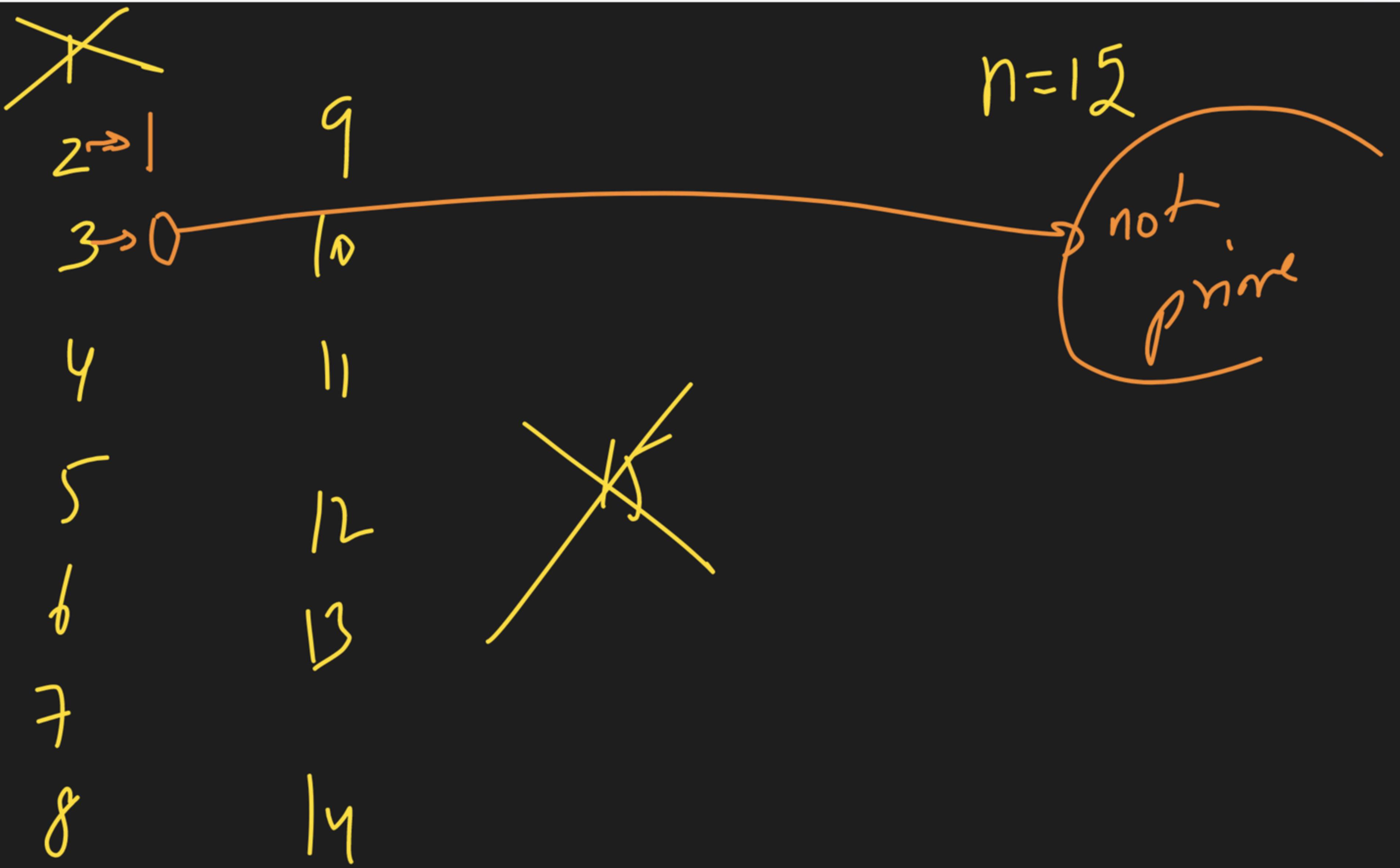
0 1 1 1

0 001

ste

int \rightarrow 4 byte \rightarrow 32 bit





row → n → Y

obj

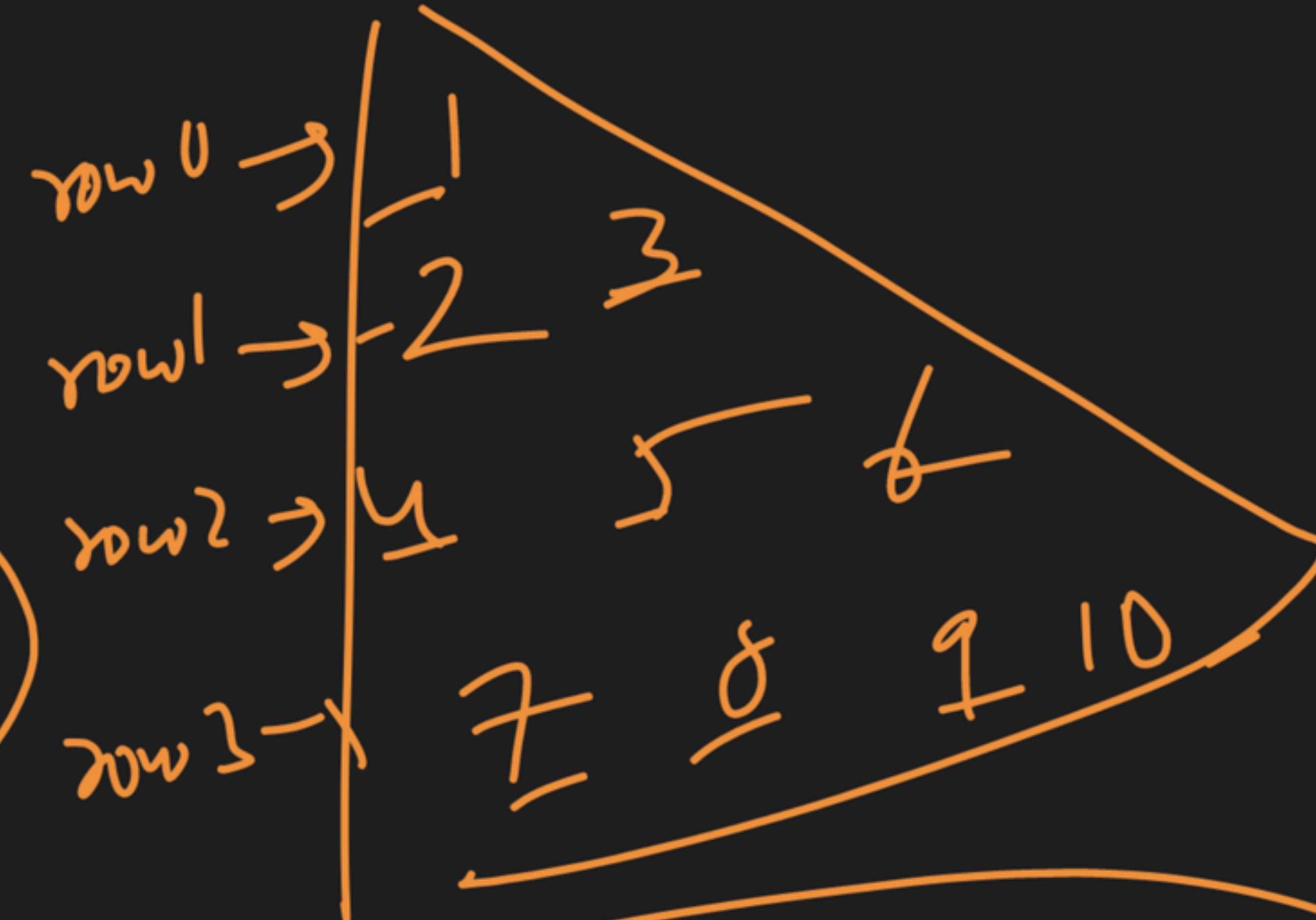
row 0 → 1 no

row 1 → 2 no

row 2 → 3 no

row 3 → 4 no

int
number = 1



for ($i=0 \rightarrow <4$)
{
 for ($c=0 \rightarrow <no[i]$)
 // logic cout < number
 num++

```
int num = 1  
for (r → 0 → <n)  
{  
    for (c=0 → <num + 1)  
    {  
        cout << num;  
        num++;  
    }  
    cout << endl;  
}
```

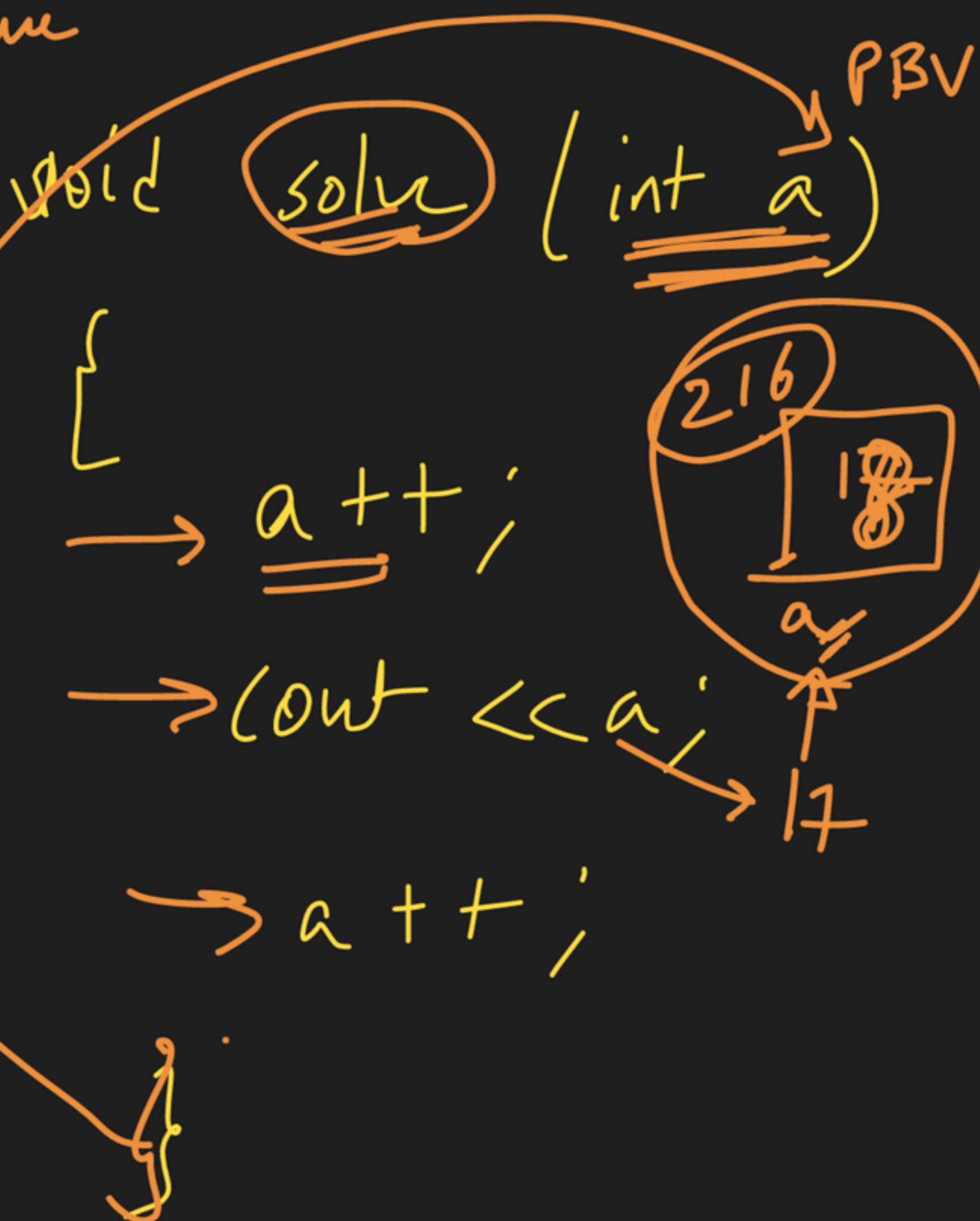
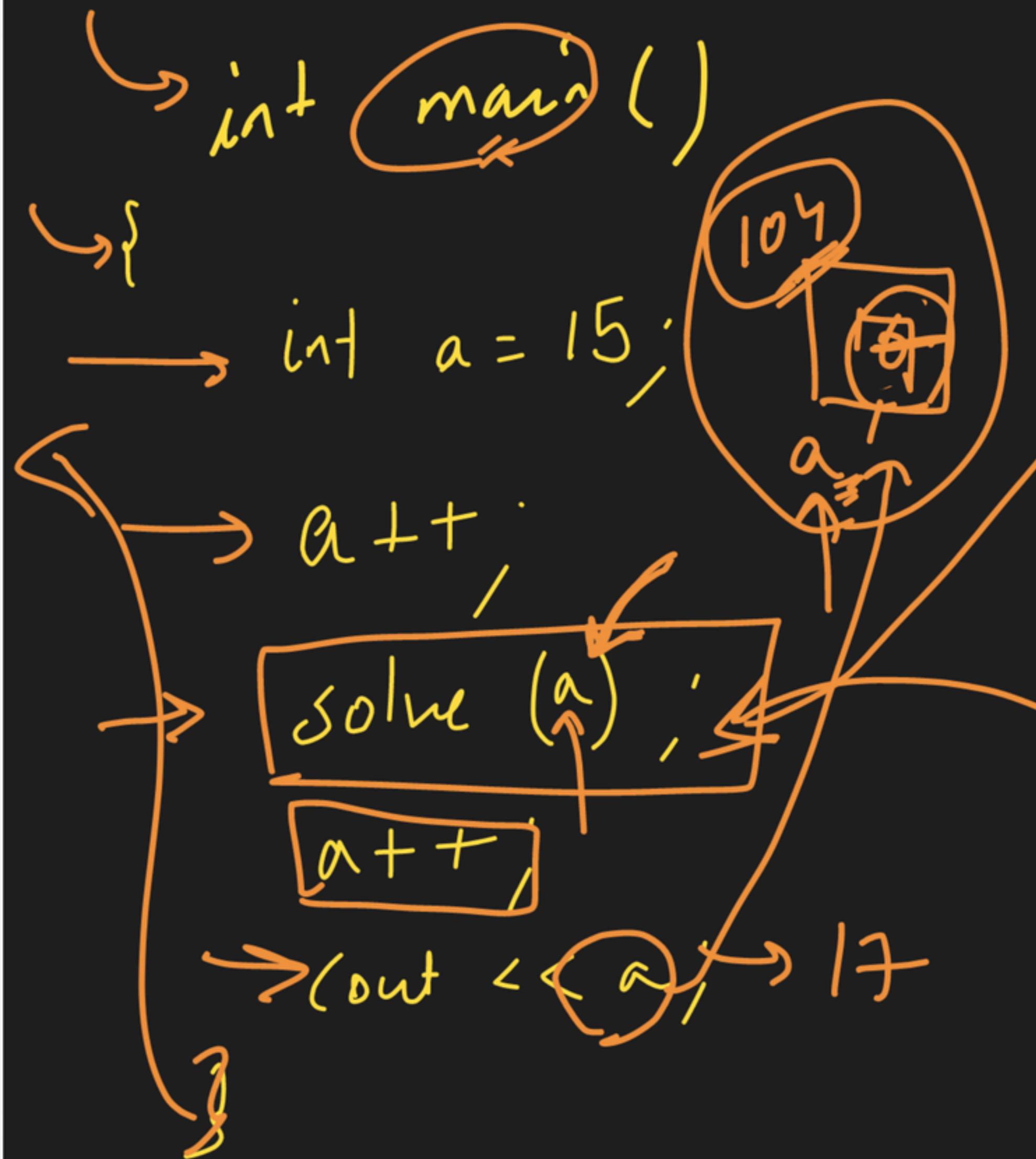
→ Pass by value

copy banti
hai

Pass by reference

actual value
me change
hota hai

Pass by value



PBV

→ int

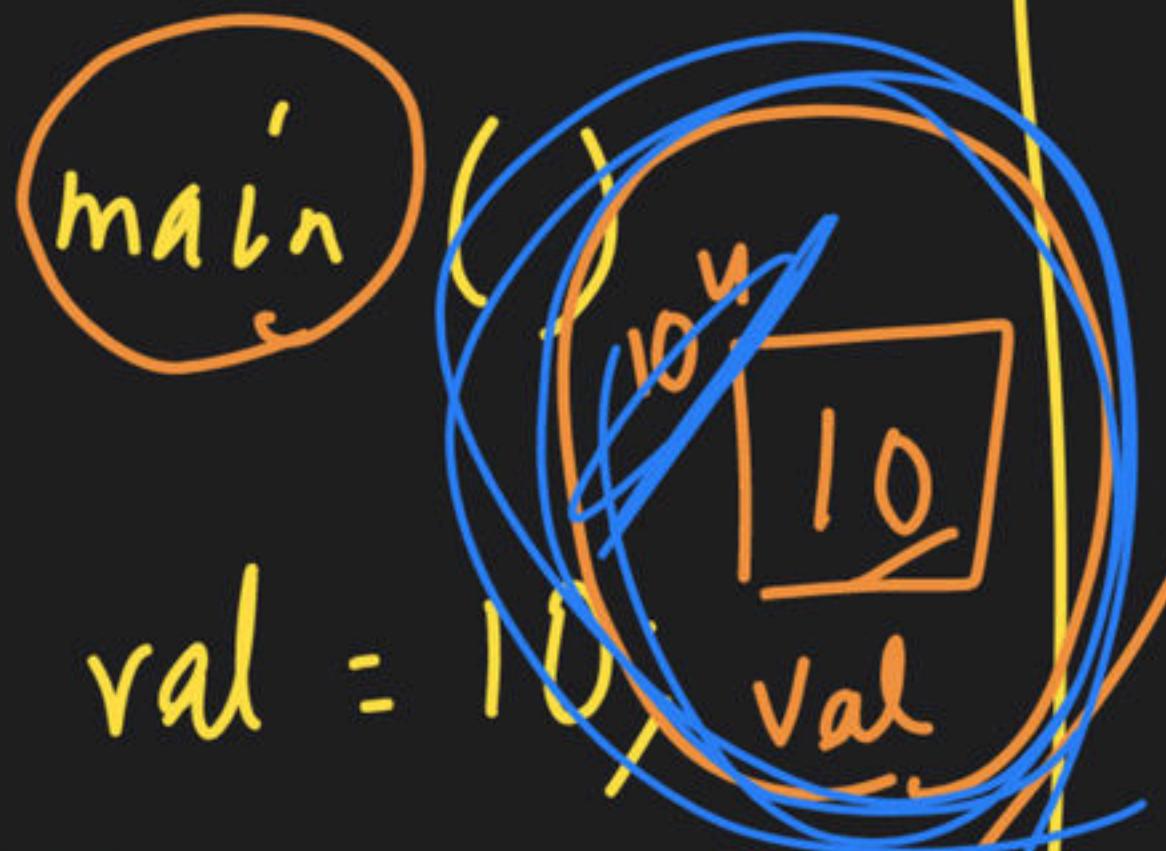
→ {

→ int val = 10;

→ solve (val);

→ cout << val;

}



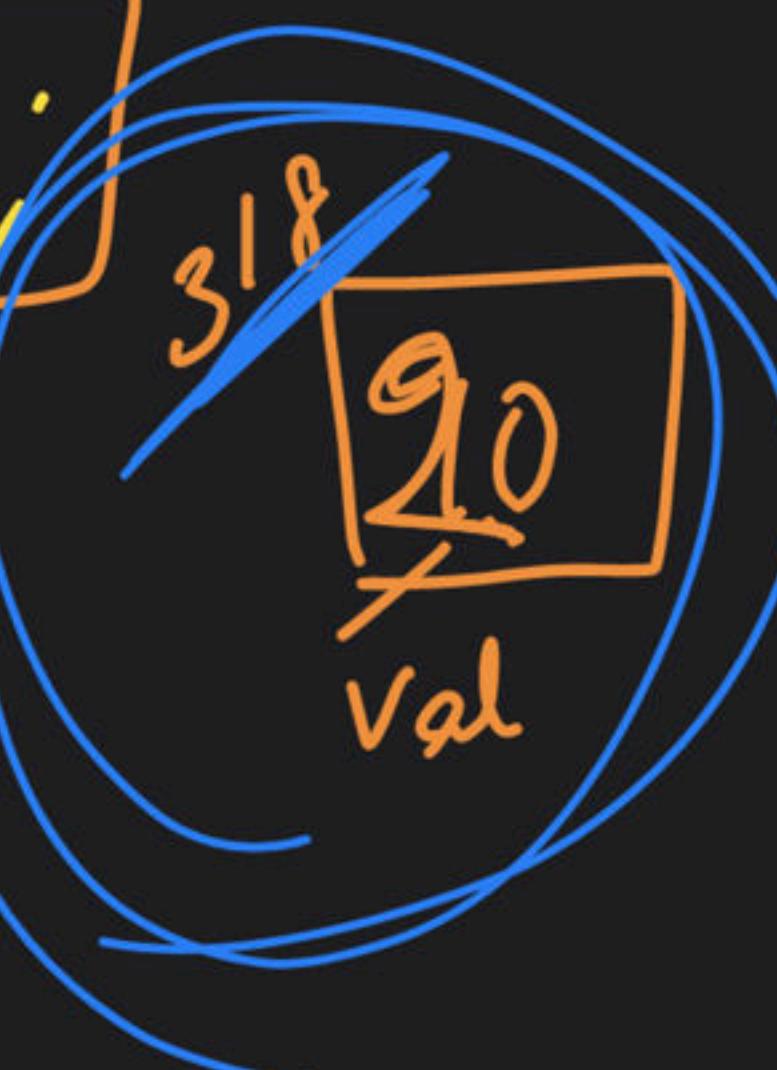
void

→ {

→ Val = val + 10;

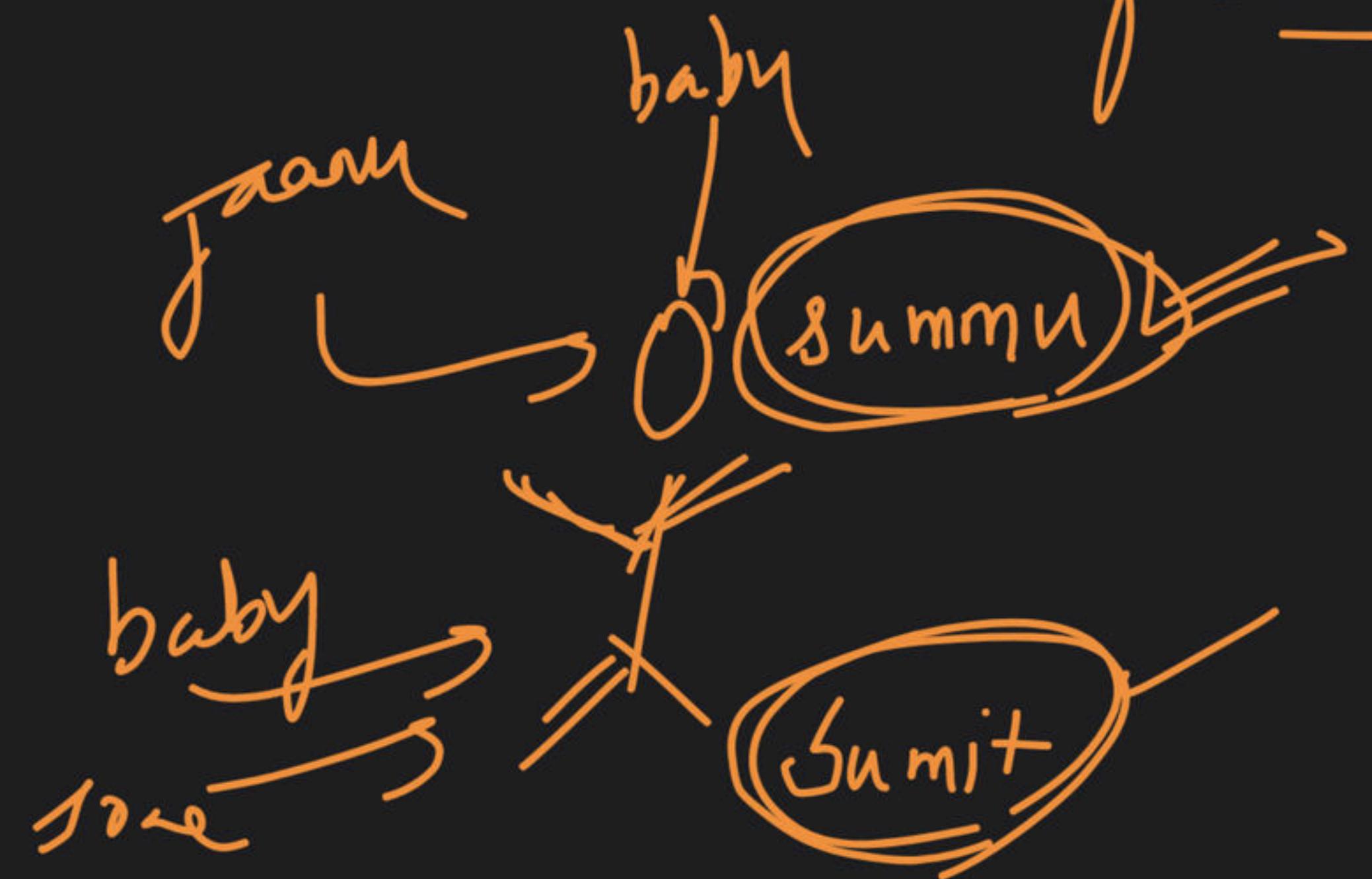
return;

}



Pass by reference

Reference Value



Pass by reference

(12, 22, 21)

int

main()

{
 int a = 10;

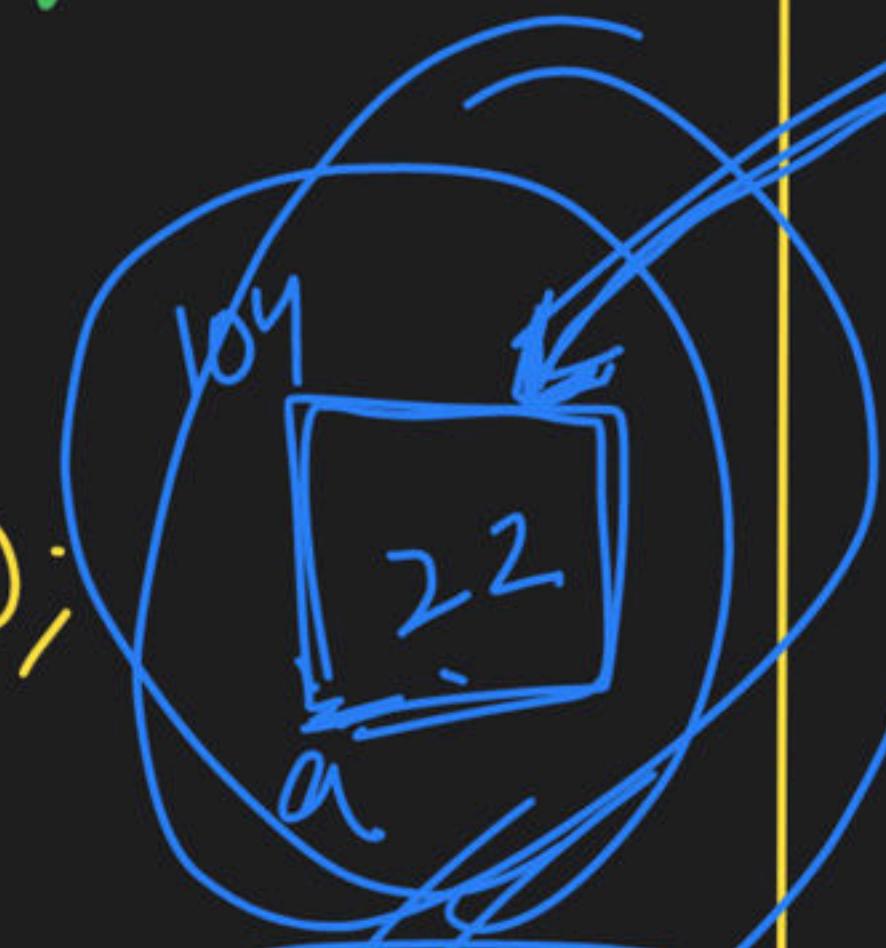
 a++;

 solve(a);

 a++;

 cout << a;

}



void

solve

(int a)

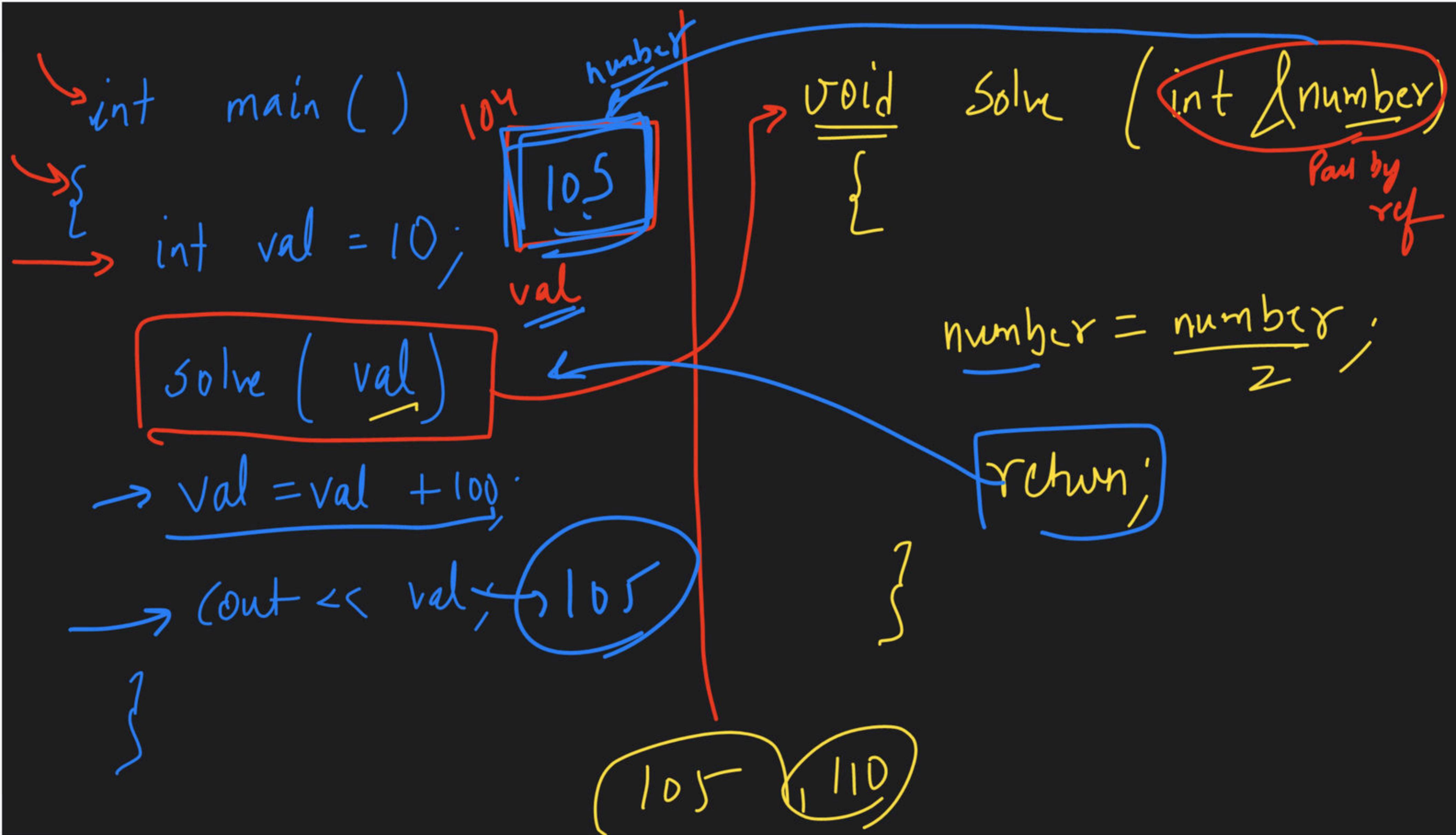
{

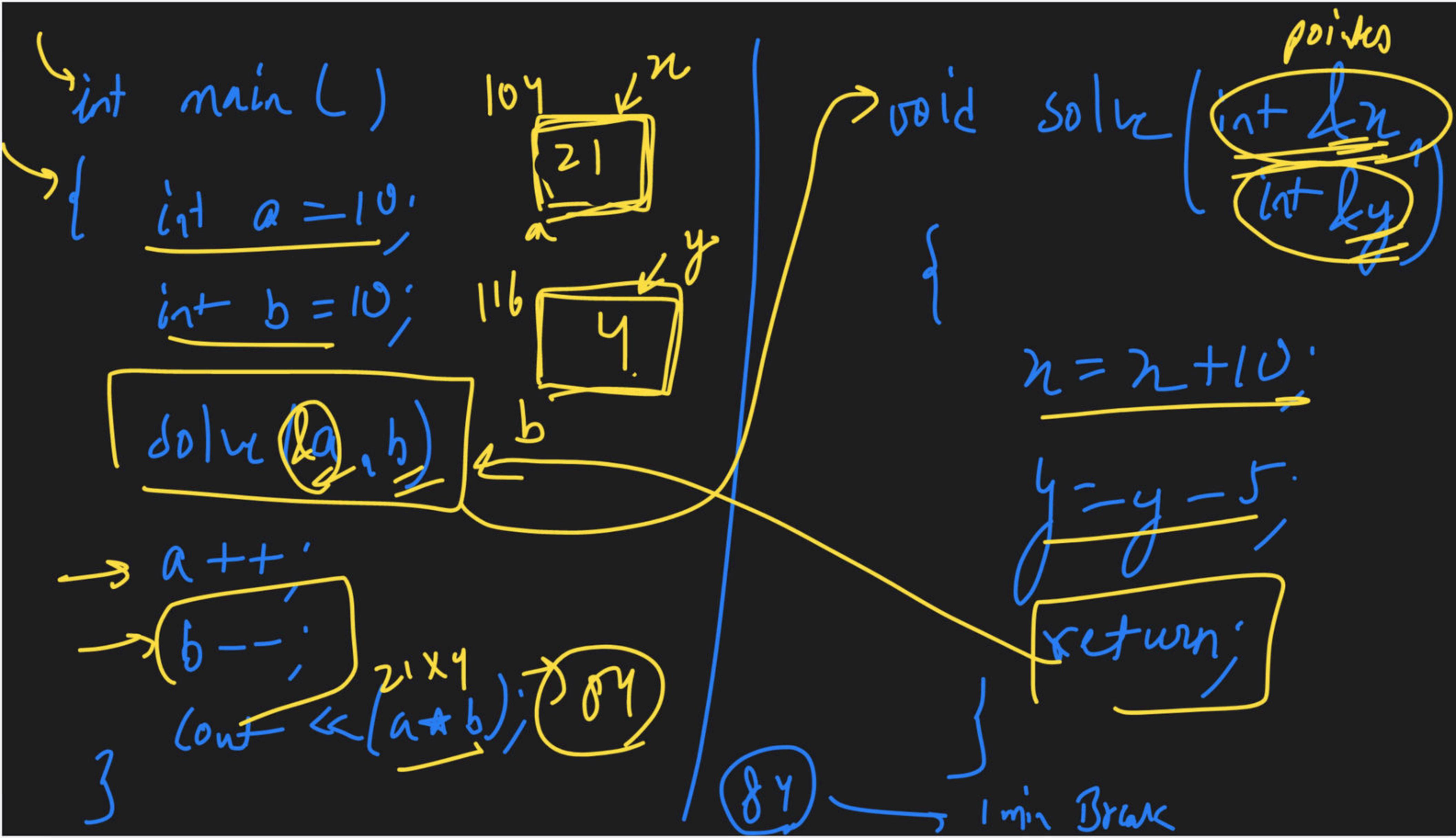
a = a + 10;

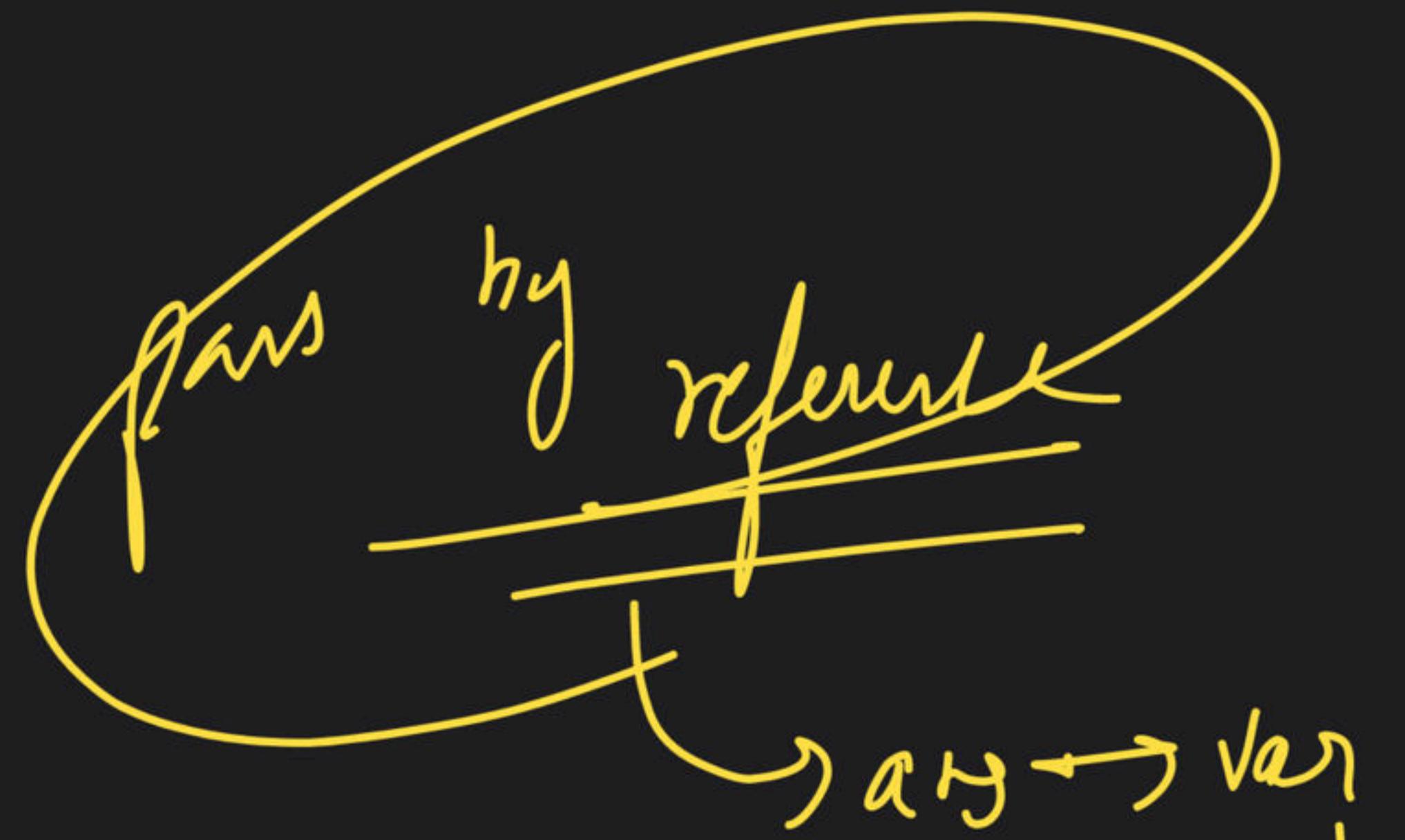
return;

}

Pass by reference

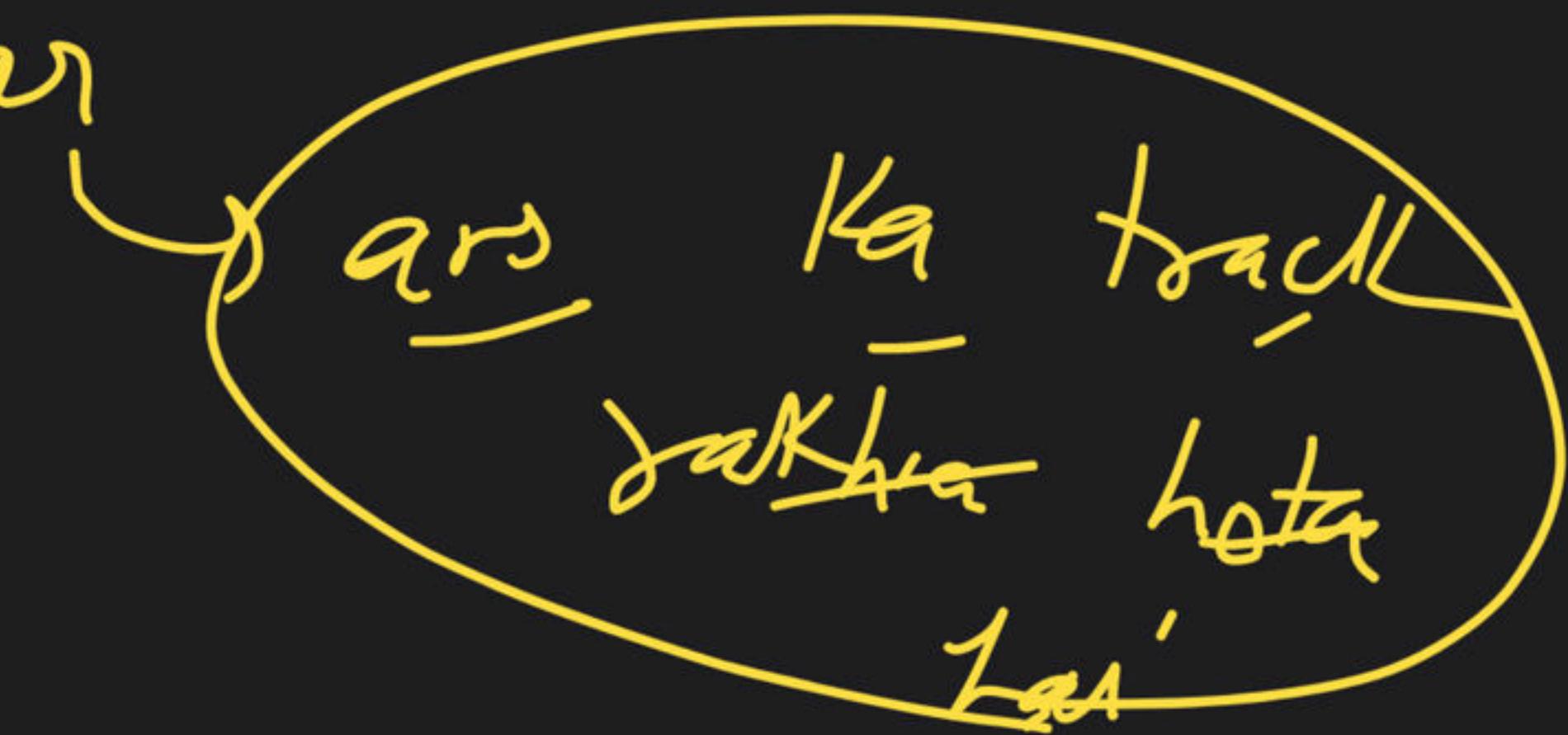




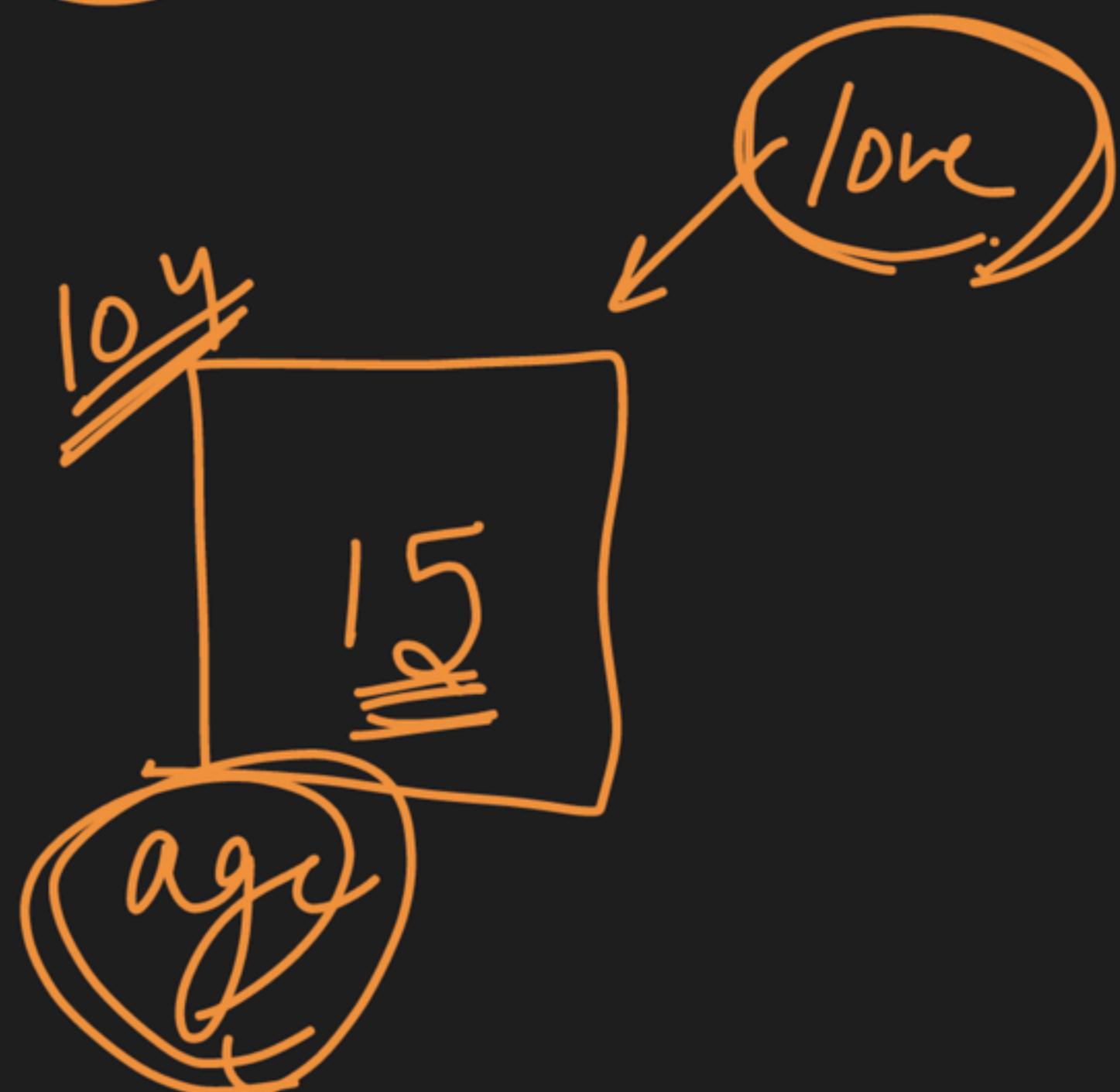
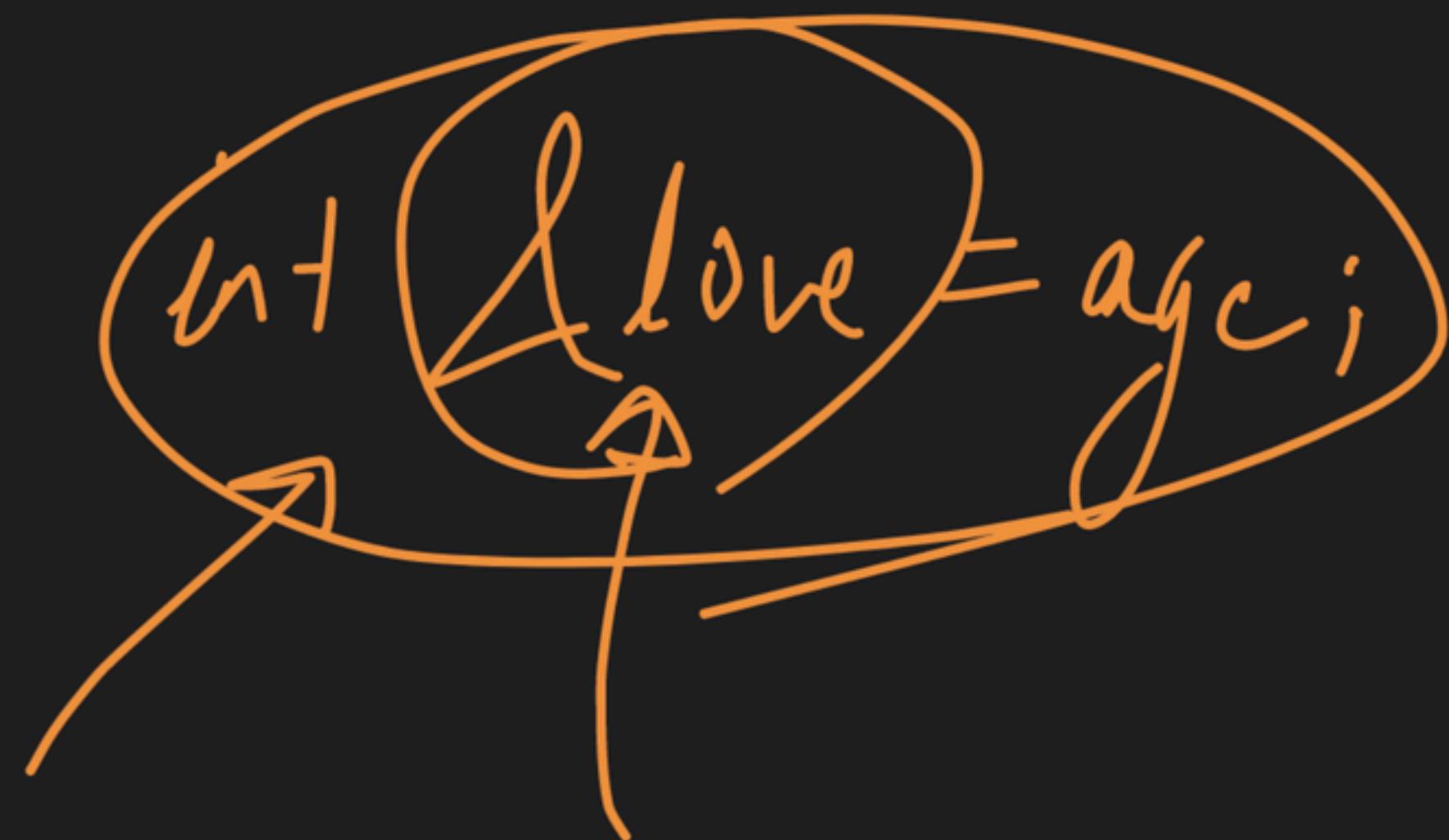


Pass by value

normally use



int age = 15



&

int love

↓
reference
variable

int val = 15 →

int demo = val

val = val + 10

(out << val) → 21

(out << demo) → 21

val --;

cout << demo; → 14

demo --

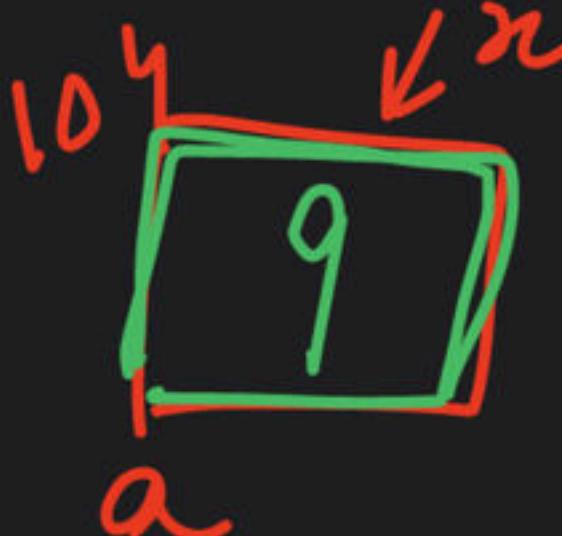
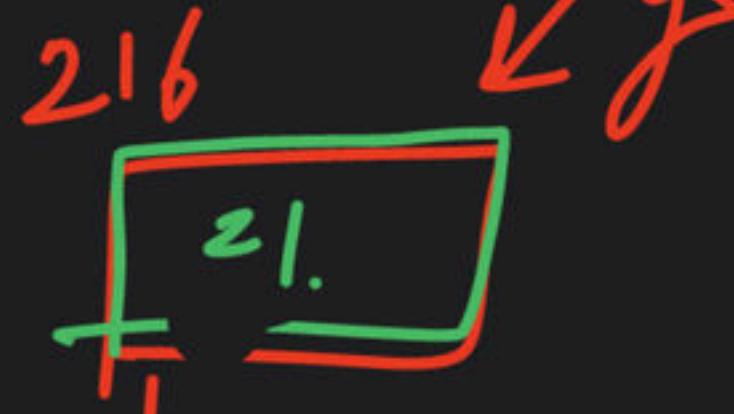
cout << val - -

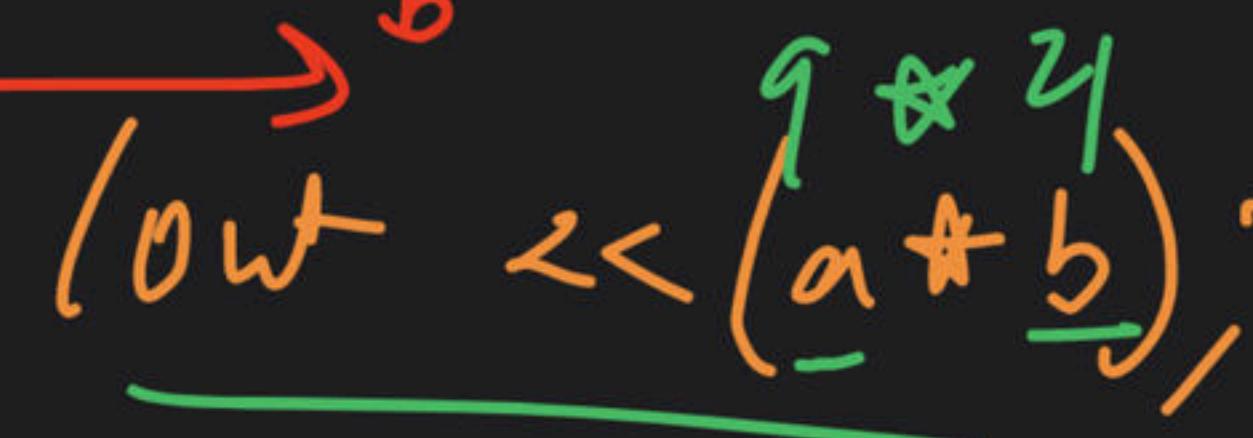
demo --;

dem

104
21
val

21

int a = 10; →  

int b = 20 → 

int n = a;

int &y = b;

n = ;

b++ ;

189









