

# Dynamic Programming - 2

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## Contents

16. Best Time to Buy and Sell Stock
17. Best Time to Buy and Sell Stock II
18. Best Time to Buy and Sell Stock III
19. Best Time to Buy and Sell Stock IV
20. Best Time to Buy and Sell Stock with Cooldown
21. Best Time to Buy and Sell Stock with Transaction Fee
22. Jump Game
23. Jump Game II
24. Reach a given Score
25. Applications of Catalan Numbers
26. Nth Catalan Number
27. Number of Valid Parenthesis Expression
28. Unique Binary Search Trees

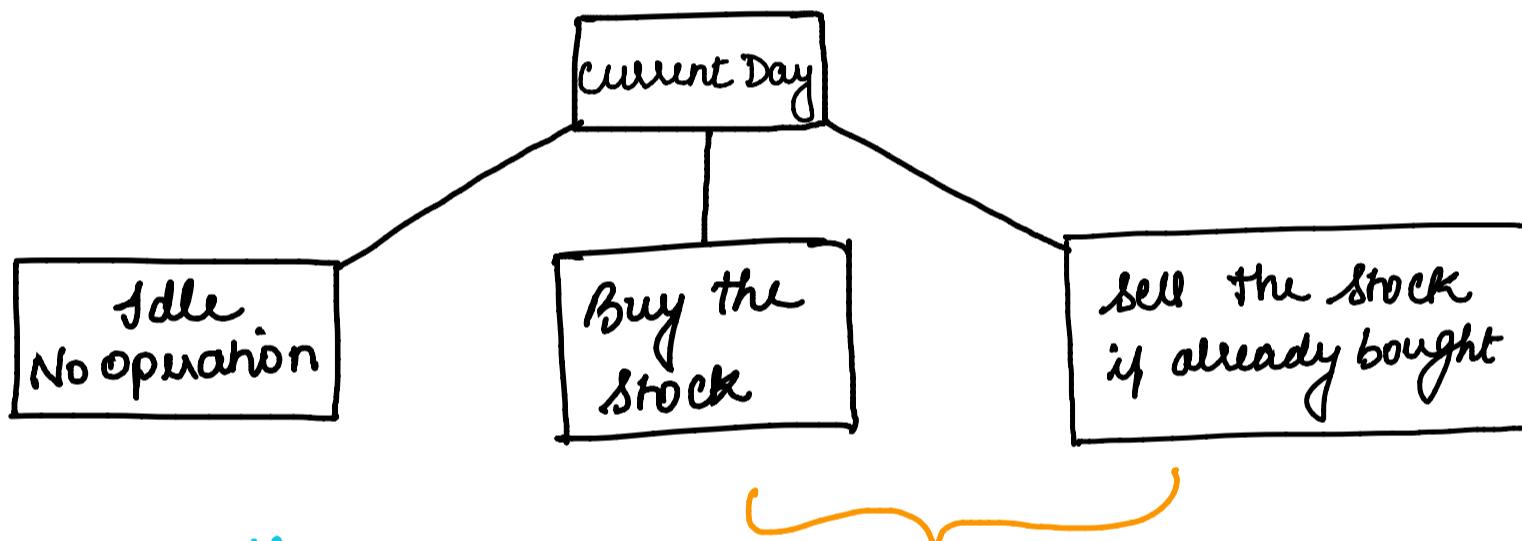
## 16 Best time to Buy & Sell Stock

Given an array of prices, find the max profit if we are allowed to do one transaction

Eg

prices = [7, 1, 5, 3, 6, 4] → we get maxProfit when we buy at day 0 & sell on day 4  
0 1 2 3 4 5  
⇒ profit =  $6 - 1 = \underline{\underline{5}}$ .

Let's look at choices we have,

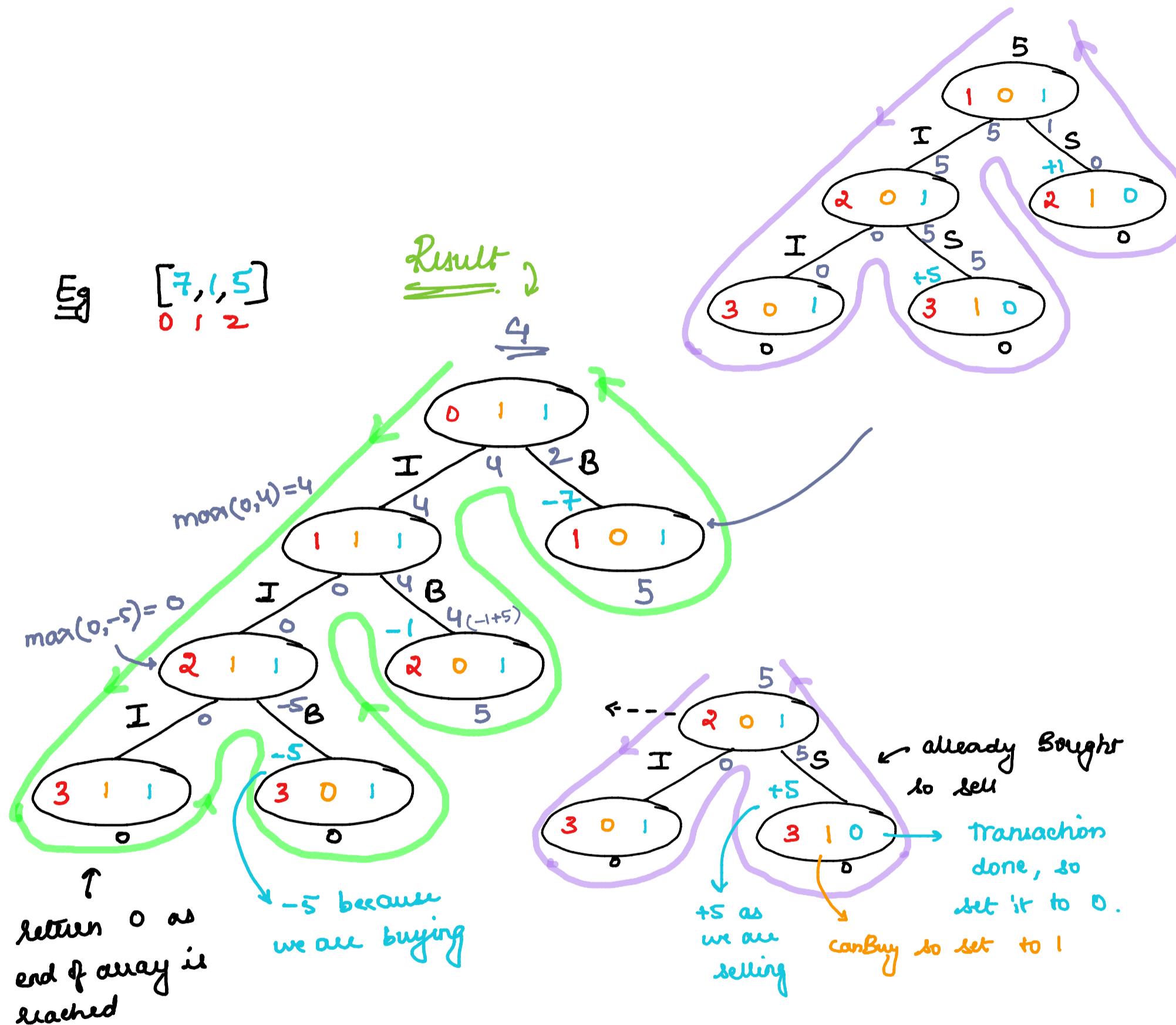


→ to handle the case that transaction could occur once, we use a variable called  $\text{transaction} = 1$ .

→ to handle these cases, we use a variable called  $\text{canBuy}$ .  
→ once bought  $\text{canBuy} = \text{false}$   
→ once sold  $\text{canBuy} = \text{true}$

∴ Our recursive structure would be as follows →

current Day, canBuy, transaction



code →



```
1 class Solution {
2 public:
3     int find(vector<int> &prices, int currDay, int k, bool canBuy, vector<vector<int>> &memo){
4
5         if(currDay >= prices.size() || k<=0 ) return 0;
6
7         if(memo[currDay][canBuy] != -1) return memo[currDay][canBuy];
8
9         if(canBuy)
10     {
11         int idle = find(prices, currDay+1, k, canBuy, memo);
12         int buy = -prices[currDay] + find(prices, currDay+1, k, !canBuy, memo);
13         return memo[currDay][canBuy] = max(buy, idle);
14     }
15     else
16     {
17         int idle = find(prices, currDay+1, k, canBuy, memo);
18         int sell = prices[currDay] + find(prices, currDay+1, k-1, !canBuy, memo);
19         return memo[currDay][canBuy] = max(sell, idle);
20     }
21 }
22 int maxProfit(vector<int>& prices) {
23     int n = prices.size();
24     vector<vector<int>> memo(n, vector<int> (2, -1));
25     // canBuy = true and transaction as k = 1
26     return find(prices,0,1,true,memo);
27 }
28 };
```

## 17 Best time to Buy & Sell Stock - II →

→ In this we can have many transactions that can be done.

Ex → prices = [<sup>0 1 2 3 4 5</sup><sub>7, 1, 5, 3, 6, 4</sub>]

↳ Buy on 1 & sell on 2 profit =  $5 - 1 = 4$

Buy on 3 & sell on 4 profit =  $6 - 3 = 3$

---

Total Profit = 7 Ans

Code →

Remove the parameter K i.e transaction limit.

```
● ● ●
1 class Solution {
2 public:
3     int find(vector<int> &prices, int currDay, bool canBuy, vector<vector<int>> &memo){
4
5         if(currDay >= prices.size()) return 0;
6
7         if(memo[currDay][canBuy] != -1) return memo[currDay][canBuy];
8
9         if(canBuy)
10        {
11             int idle = find(prices, currDay+1, canBuy, memo);
12             int buy = -prices[currDay] + find(prices, currDay+1, !canBuy, memo);
13             return memo[currDay][canBuy] = max(buy, idle);
14         }
15         else
16        {
17             int idle = find(prices, currDay+1, canBuy, memo);
18             int sell = prices[currDay] + find(prices, currDay+1, !canBuy, memo);
19             return memo[currDay][canBuy] = max(sell, idle);
20         }
21     }
22     int maxProfit(vector<int>& prices) {
23         int n = prices.size();
24         vector<vector<int>> memo(n, vector<int> (2, -1));
25         // canBuy = true and transaction are infinite so ignore k
26         return find(prices, 0, true, memo);
27     }
28 }
```

## 18 Best time to Buy & Sell Stock - III →

In this maximum profit has to be achieved by making atmost 2 transactions.

Eg prices = [3, 3, 5, 0, 0, 3, 1, 4]

↳ Buy on 4 & sell on 5      profit = 3 - 0 = 3

Buy on 6 & sell on 7      profit = 4 - 1 = 3

Total Profit = 6 Ans

code →

In the base condition is no. of transactions  $\geq 2$   
then return 0.

(Line 6)

↳ ie possible transactions  
are when it is = 0, 1

```
● ● ●
1 class Solution {
2 public:
3     int find(vector<int> &prices, int currDay, int transaction, bool canBuy,
4             vector<vector<vector<int>>> &memo){
5
6         if(currDay >= prices.size() || transaction >= 2) return 0;
7
8         if(memo[currDay][canBuy][transaction] != -1) return memo[currDay][canBuy][transaction];
9
10        if(canBuy)
11        {
12            int idle = find(prices, currDay+1, transaction, canBuy, memo);
13            int buy = -prices[currDay] + find(prices, currDay+1, transaction, !canBuy, memo);
14            return memo[currDay][canBuy][transaction] = max(buy, idle);
15        }
16        else
17        {
18            int idle = find(prices, currDay+1, transaction, canBuy, memo);
19            int sell = prices[currDay] + find(prices, currDay+1, transaction+1, !canBuy, memo);
20            return memo[currDay][canBuy][transaction] = max(sell, idle);
21        }
22    }
23    int maxProfit(vector<int>& prices) {
24        int n = prices.size();
25        vector<vector<vector<int>>> memo(n, vector<vector<int>>(2, vector<int>(2, -1)));
26        // canBuy = true and transactions are allowed 2 times
27        return find(prices, 0, 0, true, memo);
28    }
29};
```

## ⑯ Best time to Buy & Sell Stock - IV →

This is a generalised version of previous problem, instead of limiting it to 2 transactions, we need to allow atmost k transactions.

code →

Pass k as an argument & use it to limit transaction in base condition. (Line 6)

```
● ● ●
1 class Solution {
2 public:
3     int find(vector<int> &prices, int currDay, int transaction, int k, bool canBuy,
4             vector<vector<vector<int>>> &memo){
5
6         if(currDay >= prices.size() || transaction>=k) return 0;
7
8         if(memo[currDay][canBuy][transaction] != -1) return memo[currDay][canBuy][transaction];
9
10        if(canBuy)
11        {
12            int idle = find(prices, currDay+1, transaction, k, canBuy, memo);
13            int buy = -prices[currDay] + find(prices, currDay+1, transaction, k, !canBuy, memo);
14            return memo[currDay][canBuy][transaction] = max(buy, idle);
15        }
16        else
17        {
18            int idle = find(prices, currDay+1, transaction, k, canBuy, memo);
19            int sell = prices[currDay] + find(prices, currDay+1, transaction+1, k, !canBuy, memo);
20            return memo[currDay][canBuy][transaction] = max(sell, idle);
21        }
22    }
23    int maxProfit(int k, vector<int>& prices) {
24        int n = prices.size();
25        vector<vector<vector<int>>> memo(n ,vector<vector<int>>(2, vector<int>(k+1, -1)));
26        // canBuy = true and transactions are allowed atmost k times
27        return find(prices, 0, 0, k, true, memo);
28    }
29};
```

## ⑩ Best time to Buy & Sell Stock with CoolDown →

In this, cooldown means that we cannot buy a stock on the immediate day after it is sold.

⇒ The day after sold should be skipped.

code →

To skip day after sell, increment the currDay by 2. (Line 18)

```
● ● ●
1 class Solution {
2 public:
3     int find(vector<int> &prices, int currDay, bool canBuy, vector<vector<int>> &memo) {
4
5         if(currDay >= prices.size()) return 0;
6
7         if(memo[currDay][canBuy] != -1) return memo[currDay][canBuy];
8
9         if(canBuy)
10        {
11             int idle = find(prices, currDay+1, canBuy, memo);
12             int buy = -prices[currDay] + find(prices, currDay+1, !canBuy, memo);
13             return memo[currDay][canBuy] = max(buy, idle);
14        }
15        else
16        {
17            int idle = find(prices, currDay+1, canBuy, memo);
18            int sell = prices[currDay] + find(prices, currDay+2, !canBuy, memo);
19            return memo[currDay][canBuy] = max(sell, idle);
20        }
21    }
22    int maxProfit(vector<int>& prices) {
23        int n = prices.size();
24        vector<vector<int>> memo(n, vector<int> (2, -1));
25        // canBuy = true & transaction = infinite so ignore k & while sell, currDay +=2
26        return find(prices, 0, true, memo);
27    }
28};
```

## 21) Best time to Buy & Sell Stock with Transaction Fee →

In this variation, we don't have limit on transaction but while making a transaction i.e. selling it, some fee has to be paid i.e. transaction fee.

Code →

Deduct the fee from the selling day's amount.

(Line 18)

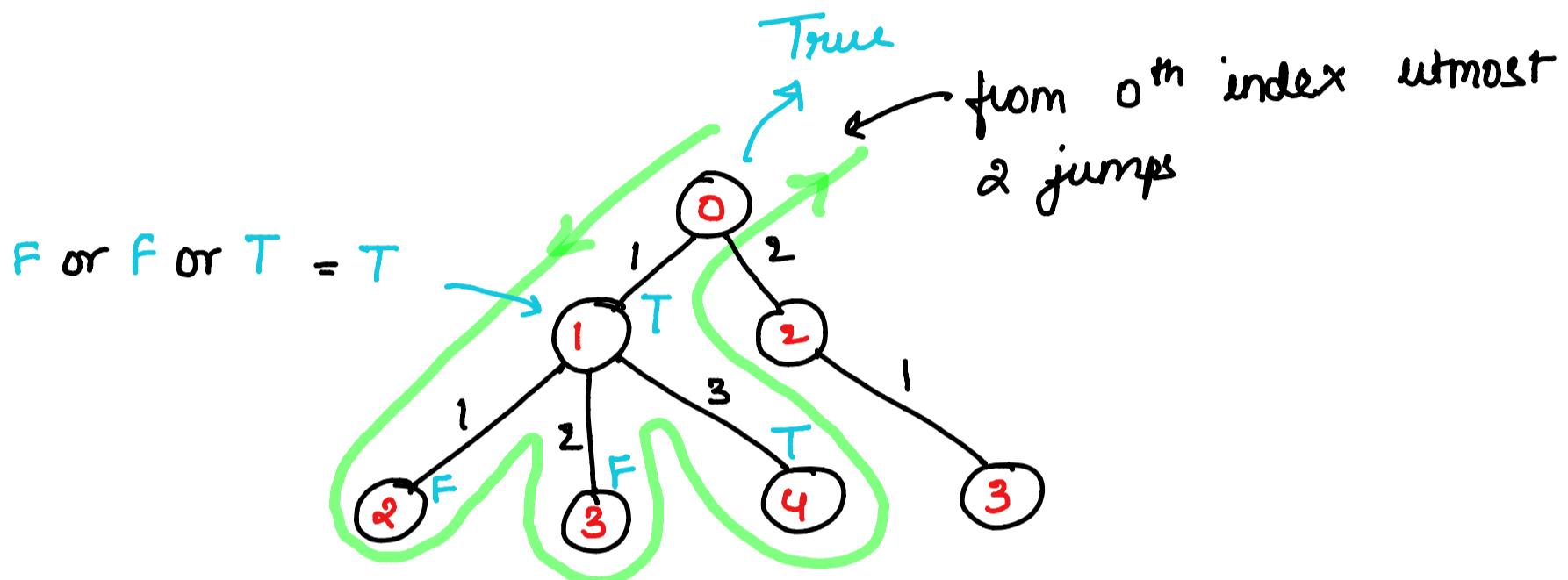


```
1 class Solution {
2 public:
3     int find(vector<int> &prices, int currDay, int fee, bool canBuy, vector<vector<int>> &memo){
4
5         if(currDay >= prices.size()) return 0;
6
7         if(memo[currDay][canBuy] != -1) return memo[currDay][canBuy];
8
9         if(canBuy)
10     {
11         int idle = find(prices, currDay+1, fee, canBuy, memo);
12         int buy = -prices[currDay] + find(prices, currDay+1, fee, !canBuy, memo);
13         return memo[currDay][canBuy] = max(buy, idle);
14     }
15     else
16     {
17         int idle = find(prices, currDay+1, fee, canBuy, memo);
18         int sell = (prices[currDay]-fee) + find(prices, currDay+1, fee, !canBuy, memo);
19         return memo[currDay][canBuy] = max(sell, idle);
20     }
21 }
22
23     int maxProfit(vector<int>& prices, int fee) {
24         int n = prices.size();
25         vector<vector<int>> memo(n, vector<int> (2, -1));
26         // canBuy = true & transaction = infinite so ignore k & while selling deduct fee
27         return find(prices, 0, fee, true, memo);
28     }
29 };
```

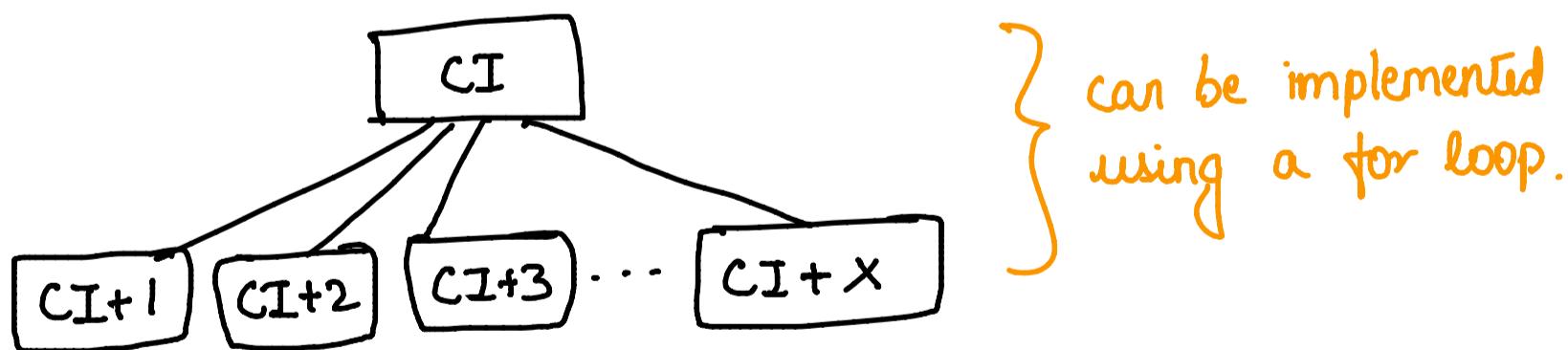
## 22 Jump Game →

Given array of nums which indicate max number of jump from any index. Return true if you can reach last index.

Eg.  $\text{nums} = [2, 3, 1, 1, 4]$



Therefore,  $[ - - \frac{x}{CI} - - - ]$



Note: submitting DP solution gives TLE. This is just for understanding. Optimal solution involves Greedy approach.

$$TC \rightarrow O(\max(\text{nums}[i]) \times n)$$

*max time for for loop.*

Code →

```

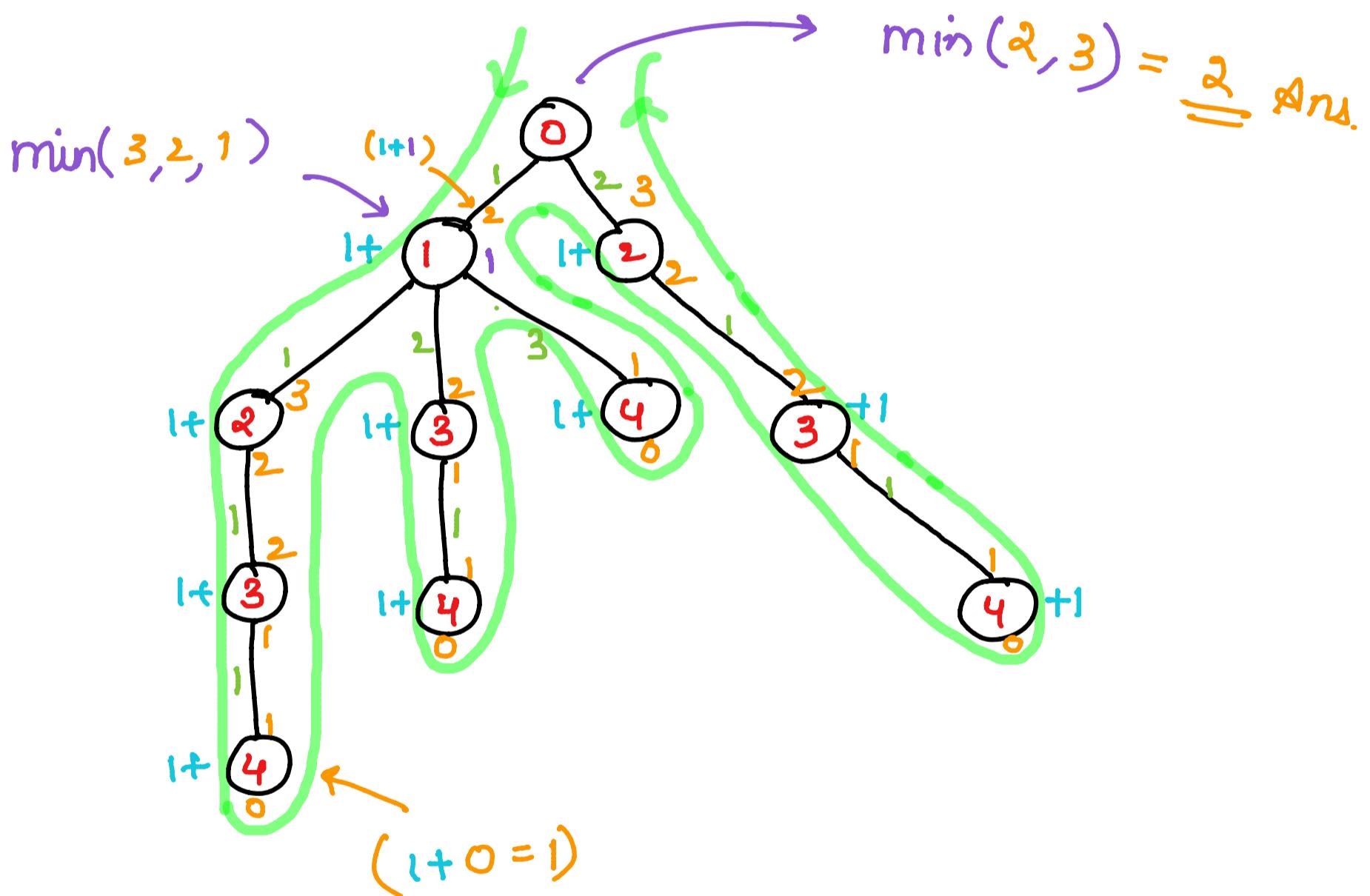
1 class Solution {
2 public:
3     bool isPossible(vector<int>& nums, int curr, unordered_map<int,bool>& memo)
4     {
5         if(curr >= nums.size()-1) return true;
6
7         int currKey = curr;
8
9         if(memo.find(currKey)!=memo.end()) return memo[currKey];
10
11        int currJump = nums[curr];
12
13        if(currJump >= nums.size() - curr) return true;
14
15        bool ans = false;
16
17        for(int i=1; i<=currJump; i++){
18            bool tempAns = isPossible(nums,curr+i,memo);
19            ans = ans || tempAns;
20        }
21        return memo[currKey] = ans;
22    }
23
24    bool canJump(vector<int>& nums){
25        unordered_map<int,bool> memo;
26        return isPossible(nums, 0, memo);
27    }
28 };

```

## ②③ Jump Game II →

Given array of nums which indicate max number of jump from any index. Reach last index in minimum number of moves.

Eg  $\text{nums} = [2, 3, 1, 1, 4]$   
 0 1 2 3 4



→ if  $\text{currentIndex} \geq \text{lastIndex}$   
 then return 0.

while returning add 1 for counting ways!

## Code →

```
1 class Solution {
2 public:
3
4     int minJumps(vector<int>& nums,int curr,vector<int>&memo)
5     {
6         if( curr >= nums.size()-1) return 0;
7
8         int currKey = curr;
9         if(memo[currKey]!=-1) return memo[currKey];
10
11         int currJump = nums[curr];
12
13         // some large value
14         int ans = 10001;
15
16         for(int i=1;i<=currJump;i++){
17             int tempans = 1 + minJumps(nums,curr+i,memo);
18             ans = min(ans, tempans);
19         }
20         return memo[currKey] = ans;
21     }
22
23     int jump(vector<int>& nums) {
24         vector<int> memo(nums.size()+1,-1);
25         return minJumps(nums, 0, memo);
26     }
27 };
```

## ②4) Reach a given score →

given 3 scores  $[3, 5, 10]$  & 'n'.

Return total number of ways to create n using the scores.

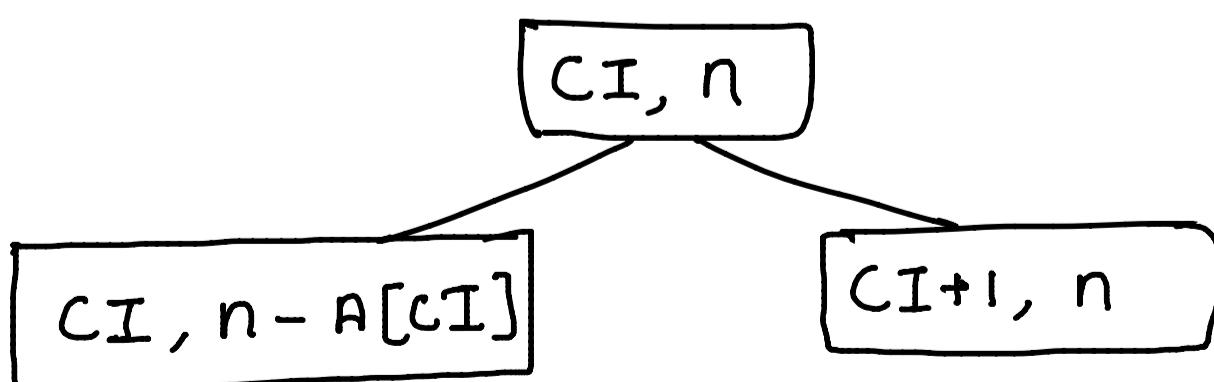
Eg  $n=8$  then no. of ways to create 8 from  $[3, 5, 10]$   
is 1.  $(3+5)$

$n=13$  then no. of ways to create 13 from  $[3, 5, 10]$   
is 2  $(3+5+5)$  &  $(3+10)$

$n=20$  then no. of ways to create 20 from  $[3, 5, 10]$   
is 4  $(3+3+3+3+3+5)$  &  $(5+5+5+5)$   
&  $(5+5+10)$  &  $(10+10)$

∴ let say  $A = [3, 5, 10]$  then

↑  
CI



Code →

```
1  typedef long long LL;
2
3  LL ways(int curr, LL n, vector<int>&score, vector<vector<int>>&vec)
4  {
5      if(n==0) return 1;
6
7      if(curr>=score.size()) return 0;
8
9      if(vec[curr][n]!=-1) return vec[curr][n];
10
11     LL consider = 0;
12
13     if(score[curr]<=n)
14         consider = ways(curr,n-score[curr],score,vec);
15
16     LL notconsider = ways(curr+1,n,score,vec);
17
18     return vec[curr][n] = consider + notconsider;
19 }
20
21 LL count(LL n)
22 {
23     vector<int> score{3,5,10};
24     vector<vector<int>> vec(score.size(),vector<int>(1001,-1));
25     return ways(0,n,score,vec);
26 }
```

25) Applications of Catalan Number →

Catalan Numbers are defined using the formula

$$C_n = \frac{(2n)!}{(n+1)! n!} = \prod_{k=2}^n \frac{n+k}{k} \text{ for } n \geq 0$$

This can be used recursively as follows,

$$C_{n+1} = \sum_{i=0}^n C_i C_{i-1} \quad \left. \right\} \quad n \geq 0 \text{ & } C_0 = 1$$

$$\rightarrow C_0 = \underline{\underline{1}}.$$

$$\rightarrow C_1 = \underline{\underline{1}}.$$

$$\rightarrow C_2 = C_0 \cdot C_1 + C_1 \cdot C_0 = 1 \cdot 1 + 1 \cdot 1 = \underline{\underline{2}}.$$

$$\rightarrow C_3 = C_0 C_2 + C_1 C_1 + C_2 C_0 = 1 \cdot 2 + 1 \cdot 1 + 2 \cdot 1 = \underline{\underline{5}}.$$

$$\begin{aligned} \rightarrow C_4 &= C_0 C_3 + C_1 C_2 + C_2 C_1 + C_3 C_0 \\ &= 1 \cdot 5 + 1 \cdot 2 + 2 \cdot 1 + 5 \cdot 1 = \underline{\underline{14}}. \end{aligned}$$

dpp^n's →

1. No. of possible BST with  $n$  keys.
2. No. of valid combinations for  $N$  pair of parenthesis.

## 26 $N^{th}$ Catalan Number

To find  $N^{th}$  Catalan Number we can use formula

$$C_{n+1} = \sum_{i=0}^n C_i C_{i-1} \quad \left. \right\} \quad n \geq 0 \quad \& \quad C_0 = 1$$

↳ this can be implemented by

- i) having base condition for  $n == 0 \& n == 1$
- ii) using a loop to sum values from  $i = 0$  to  $n$ .

Code →

```

● ● ●

1 class Solution
2 {
3     public:
4     cpp_int ncatalan(int n, vector<cpp_int>& memo) {
5         if(n == 0 || n == 1) return 1;
6
7         int curr = n;
8         if(memo[curr] != -1) return memo[curr];
9
10        cpp_int catalan = 0;
11
12        for(int i=0;i<n;i++) {
13            catalan += ncatalan(i, memo)*ncatalan(n-i-1, memo);
14        }
15
16        memo[curr] = catalan;
17        return memo[curr];
18    }
19
20    cpp_int findCatalan(int n)
21    {
22        vector<cpp_int> memo(1001,-1);
23        return ncatalan(n, memo);
24    }
25}

```

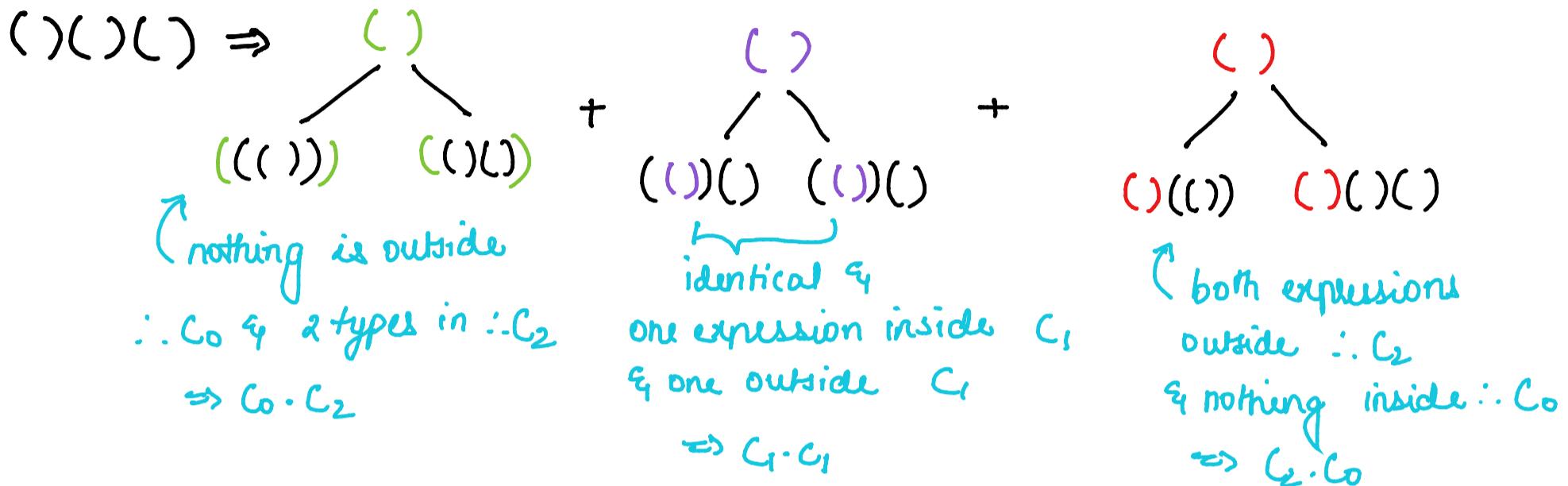
27

## Number of valid Parenthesis Expression →

Given  $N$ , find total number of ways in which we can arrange  $N$  pair of parenthesis in a balanced way.

Eg

$$N=2 \Rightarrow ( ) ( ) ( ), ( ) ( ( ) ), (( )) ( ), (( ( ) )) \therefore \text{res} = 4$$



$$\Rightarrow C_0 \cdot C_2 + C_1 \cdot C_1 + C_2 \cdot C_0 = C_3 \Rightarrow \text{for } n \text{ we need to find ncatalan}(n/2)$$

Code →

```

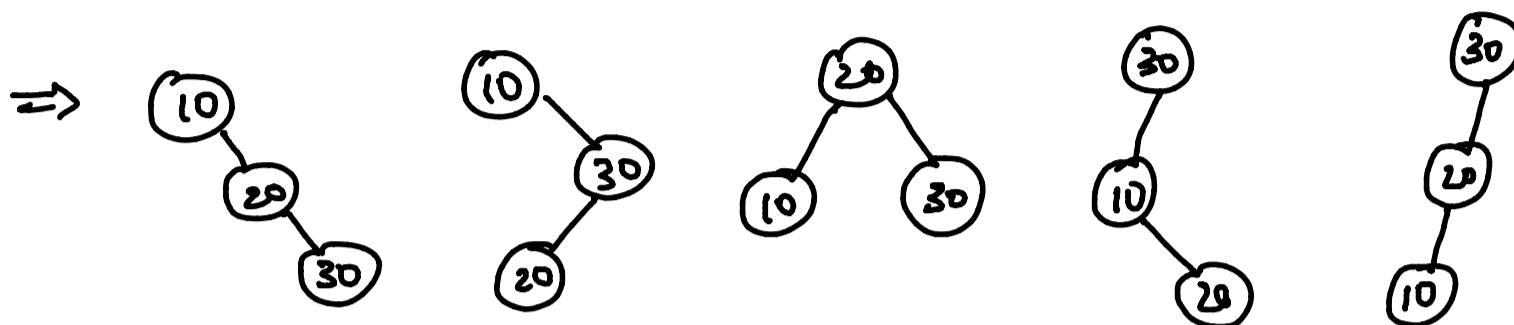
● ● ●

1 #include<bits/stdc++.h>
2 using namespace std;
3
4 int ncatalan(int n, unordered_map<int,int>& memo) {
5     if(n == 0 || n == 1) return 1;
6
7     int curr = n;
8     if(memo[curr]!=-1) return memo[curr];
9
10    int catalan = 0;
11
12    for(int i=0;i<n;i++) {
13        catalan += ncatalan(i, memo)*ncatalan(n-i-1, memo);
14    }
15
16    memo[curr] = catalan;
17    return memo[curr];
18 }
19
20 int countValidParenthesis(int n)
21 {
22     unordered_map<int,int> memo;
23     return ncatalan(n/2, memo);
24 }
25
26 int main(){
27     int n;
28     cin>>n;
29     cout<<countValidParenthesis(n);
30 }
```

## 28 Unique Binary Search Trees →

given integer  $n$ , returns no. of unique BST that can be formed.

Eg  $n=3$  & let's say elements are  $[10, 20, 30]$



∴ For  $n=3$ , the result is 5.

∴ The catalan number gives us the result.

code →

```
● ○ ●  
1 class Solution {  
2 public:  
3  
4     int uniqueBST(int n, vector<int>& memo)  
5     {  
6         if(n==0 || n==1) return 1;  
7  
8         if( memo[n]!=-1) return memo[n];  
9  
10        int ans = 0;  
11        for(int i=0;i<n;i++)  
12            ans += uniqueBST(i,memo)*uniqueBST(n-i-1,memo);  
13  
14        return memo[n] = ans;  
15    }  
16  
17    int numTrees(int n) {  
18        vector<int> memo(n+1,-1);  
19        return uniqueBST(n, memo);  
20    }  
21};
```

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