AI ASSISTED CODING

LAB-12.4

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BATCH:03

TASK-01:

Implementing Bubble Sort with AI Comments.

PROMPT:

Generate a python code to implement bubble sort.

CODE:

```
def bubble_sort(arm);

def bubble_sort(arm);

Sorts a list in ascending order using the Bubble Sort algorithm.

sparm arr: A list of comparable elements.

"""

n - lm(arm)

# Outer loop to iterate through the list. This represents the number of passes.

# After each pass "I", the 1-th largest element will be in its correct final position.

for i in range(n):

# A flag to optimize the sort. If no swaps occur in a full pass,

# a list get to applicable sorted, and we can terminate early.

**Swapped False

# Inner loop for comparing adjacent elements.

# The range is "n-1-I'b because the last "i" elements are already sorted

# mad in their correct final positions.

for j in range(n) - 1 - 1):

# Compare the current element with the next one.

if #rm(j) arm(j=1) - arm(j=1) | arm(j=1) |
```

```
PROBLEMS OUTPUT DEBUG CONSOLE TERMINAL PORTS

PS C:\Users\ramch\OneDrive\Desktop\ai> & C:\Users\ramch\AppData/Local/Programs\Python\Python312\python.exe c:\Users\ramch\OneDrive\Desktop\ai\lab12.4\12.4.1.py
#Driginal list: [64, 34, 25, 12, 22, 11, 90]
Sorted list: [11, 12, 22, 25, 34, 64, 90]

Original list (already sorted): [1, 2, 3, 4, 5]
Sorted list: [1, 2, 3, 4, 5]
SPS C:\Users\ramch\OneDrive\Desktop\ai>
```

Bubble Sort repeatedly compares and swaps adjacent elements if they are in the wrong order, moving the largest element to the end in each pass — like bubbles rising to the top.

TASK-02:

Optimizing Bubble Sort → Insertion Sort

PROMPT:

Convert the following bubble sort code to the insertion sort.

CODE:

```
PROBLEMS OUTPUT DEBUG COMSOLE TERMINAL PORTS

PS C:\USers\ramch\OneDrive\Desktop\aix & C:\Users\ramch\AppData/Local/Programs/Python/Python312/python.exe c:\Users\ramch\OneDrive\Desktop\aix/lab12.4/12.4.1.py

**Original list: [64, 34, 25, 12, 22, 11, 90]

Original list: [11, 12, 22, 25, 34, 64, 90]

Original list: [1, 12, 3, 4, 5]

Sorted list: [1, 2, 3, 4, 5]

SPS C:\USers\ramch\UneDrive\Desktop\aix & C:\Users\ramch\AppData/Local/Programs/Python/Python312/python.exe c:\Users\ramch\UneDrive\Desktop\aix\lab12.4/12.4.2.py

Bubble Sort took: 0.000750 seconds

Insertion Sort took: 0.000750 seconds

OPS C:\USers\ramch\UneDrive\Desktop\aix\lab12.4/12.4.2.py

Bubble Sort took: 0.000750 seconds

OPS C:\USers\ramch\UneDrive\Desktop\aix\lab12.4/12.4.2.py
```

Optimization from Bubble Sort to Insertion Sort:

Instead of repeatedly swapping adjacent elements like in Bubble Sort, **Insertion Sort** shifts elements to insert each item directly into its correct position. This reduces unnecessary swaps and makes it faster, especially for nearly sorted data.

TASK-03:

Binary Search vs Linear Search

PROMPT:

Implement the linear search and binary search with comments.

CODE:

```
### A CALON 2.

| Import random
| Import random
| Import random
| Performs a linear search to find the target element in a list.
| Import random
| Import rand
```

```
> • 12.4.3.py > ...
def binary_search(arr: list, target: any) -> int:
                    - Pre-requisite: The input list `arr` MUST be sorted.

- Time Complexity:

- Best Case: O(1) (target is the middle element)

- Average Case: O(log n)

- Morst Case: O(log n)

- Borst Case: O(log n)

- Space Complexity: O(1) (iterative version) or O(log n) (recursive version due to call stack)

- Highly efficient for large, sorted datasets.
         while low <= high:
    mid = (low + high) // 2
    if arr[mid] == target:
        return mid
    elif arr[mid] < target:
        low = mid + 1
    else:
        high = mid - 1
return -1
         # Tangets for search
tanget present_start = 0
tanget_present_middle = LIST_SIZE // 2
tanget_present_end = LIST_SIZE - 1
tanget_not_present = LIST_SIZE + 100
          # Student Observation Table Header
print(f*('$cenario':<0) | {'Linear Search Time (s)':<25} | {'Binary Search Time (s)':<25}")
print("-" * 45)</pre>
         # Test 1: Linear Search on unsorted data (target present)
start_time = time.perf_counter()
linear_search(unsorted_data, target_present_middle)
end_time = time.perf_counter()
linear_time_unsorted_present = end_time - start_time
print(f"('Unsorted (Target Present)':<38) | ('Inear_time_unsorted_present:<25.8f) | {'N/A (Requires Sorted)':<25}")</pre>
      linear_time_unsorted_present = end_time - start_time
print(f"{'Unsorted_(Target Present)':<30} | {linear_time_unsorted_present:<25.8f} | {'N/A (Requires Sorted)':<25}")</pre>
     # Test 2: Linear Search on unsorted data (target not present)

Start_time = time.perf_counter()

linear_search(unsorted_data, target_not_present)

end_time = time.perf_counter()

linear_time = time.perf_counter()

linear_time_unsorted_not_present = end_time = start_time

print(f"{'Unsorted (Target Not Present)':<38} | {linear_time_unsorted_not_present:<25.8f} | {'N/A (Requires Sorted)':<25}")
     # Test Of Linear search on sorted data (target pres

Start_time = time.perf_counter()

linear_search(sorted_data, target_present_middle)

end_time = time.perf_counter()

linear_time_sorted_present = end_time - start_time
     # Test 4: Binary Search on sorted data (target present)
start_time = time.perf_counter()
binary_search(sorted data, target_present_middle)
end_time = time.perf_counter()
binary_time_sorted_present = end_time - start_time
print(f^{*(sorted_traget_present)*:<38) | (linear_time_sorted_present:<25.8f) | {binary_time_sorted_present:<25.8f}*)</pre>
     # Test 5: Linear Search on sorted data (target not present)
start_time = time.perf_counter()
linear_search(sorted_data, target_not_present)
end_time = time.perf_counter()
linear_time_sorted_not_present = end_time - start_time
     # Test 6: Binary Search on sorted data (target not present)
start_time = time.perf_counter()
binary_search(sorted_data, target_not_present)
end_time = time.perf_counter()
binary_time_sorted_not_present = end_time - start_time
print(f*('Sorted_(Target_Not_Present)':<30) | (linear_time_sorted_not_present:<25.8f) | (binary_time_sorted_not_present:<25.8f)")</pre>
     print("\nllote: Binary Search times for unsorted data are marked 'N/A' as it requires a sorted list.")
print("If the data is initially unsorted, the time to sort it must be added to Binary Search's total time.")
```

```
PS C:\Users\ramch\OneDrive\Desktop\ai> & C:\Users\ramch\AppBata/Local/Programs/Python/Python312/python.exe c:\Users\ramch\OneDrive\Desktop\ai/1ab12.4/12.4.3.py

--- Performance Comparison (List Size: 1800080) ---

Scenario | Linear Search Time (s) | Binary Search Time (s)

--- Insorted (Target Present) | 0.00364970 | N/A (Requires Sorted)
--- Unsorted (Target Not Present) | 0.0027820 | N/A (Requires Sorted)
--- Unsorted (Target Not Present) | 0.0027820 | N/A (Requires Sorted)
--- Sorted (Target Not Present) | 0.00211910 | 0.00000770
--- Sorted (Target Not Present) | 0.00224120 | 0.00000340

Note: Binary Search times for unsorted data are marked 'N/A' as it requires a sorted list.
--- If the data is initially unsorted, the time to sort it must be added to Binary Search's total time.
--- PS C:\Users\ramch\OneDrive\Desktop\ai>
```

Linear Search: Checks each element one by one until the target is found or the list ends. Works on **unsorted** data but is **slow (O(n))**.

Binary Search: Repeatedly divides a **sorted** list in half to find the target. Much **faster (O(log n))**, but requires the data to be sorted.

TASK-04:

Quick Sort and Merge Sort Comparison

PROMPT:

Implement the quick sort and merge sort using recursion.

CODE:

```
isolate > 12.4 py > Comege_got

1 import time
2 import random
3 import by:
5 # Increase recursion limit for large datasets, especially for Quick Sort's worst case.
5 yx.setrecursionlimit(2000)
6 syx.setrecursionlimit(2000)
7 def merge_sort(arr: list) -> list:
6 syx.setrecursionlimit(2000)
7 def merge_sort(arr: list) -> list:
7 def merge_sort(arr: list) -> list:
8 feege Sort is a divide and conquer algorithm. It works by recursively
8 dividing the input list into the halves, calling itself for the two halves,
9 dividing the input list into the halves, calling itself for the two halves,
15 dividing the input list into the halves, calling itself for the two halves,
16 Args:
17 arr. (list). The list of elements to be sorted.
18 Returns:
19 | Inst. A new list containing the sorted elements.
19 | Performance Roses:
10 | Inst. Completing the sorted elements.
10 | Returns:
11 | Inst. Completing the sorted elements.
10 | Requires additional space to hold the merged sub-arrays.
10 | Requires additional space to hold the merged sub-arrays.
10 | Performance Rose | Completing | Completing |
10 | Requires additional space to hold the merged sub-arrays.
11 | Inst. Limit = merge_sort(arr[inid]) | right[half = merge_sort(arr[inid
```

```
# Append remaining elements
sorted_list.extend(left[i:])
sorted_list.extend(right[j:])
return sorted_list
       Performance Notes:

- Time Complexity:

- Best Case: O(n log n) (pivot is always the median)

- Average Case: O(n log n)

- Norst Case: O(n*2) (pivot is always the smallest or largest element,
| which occurs with already sorted or reverse-sorted data).

- Space Complexity: O(log n) on average (due to recursion stack),
| O(n) in the worst case.
def _quick_sort_recursive(arr, low, high):
    """Helper function for recursive calls."""
       ""Helper function for recursive calls.""

if low < high:

partition_index = _partition(arr, low, high)
_quick_sort_recursive(arr, low, partition_index - 1)
_quick_sort_recursive(arr, partition_index + 1, high)
                                                                                                                                                                                                                                                                                                               Q Ln 17, Col 55 Spaces: 4 UTF-8 C
 def _partition(arr, low, high):
         if arr[j] <= pivot:
    i += 1
    arr[i], arr[j] = arr[j], arr[i]
arr[i + 1], arr[high] = arr[high], arr[i + 1]
return i + 1</pre>
# --- Performance Comparison
if __name__ == "__main__":
    LIST_SIZE = 1000
        # Generate data
random_randint(0, LIST_SIZE) for _ in range(LIST_SIZE)]
sorted_data = list(range(LIST_SIZE))
reverse_sorted_data = list(range(LIST_SIZE, 0, -1))
        datasets = {
    "Random": random_data,
    "Sorted": sorted_data,
                "Reverse-Sorted": reverse sorted data
        # We pass a copy because quick_sort sorts in-place qs_data = data.copy()
```

OUTPUT:

start_time = time.perf_counter()
quick_sort(qs_data)
end_time = time.perf_counter()
qs_time = end_time - start_time

We pass a copy to be consistent
ms_data = data.copy()
start_time = time.perf_counter()
merge_sort(ms_data)
end_time = time.perf_counter()
ms_time = end_time - start_time

 $print("\nNote: Quick Sort's O(n'2) worst-case on sorted data is clearly visible.") print("Merge Sort's O(n log n) performance is consistent across all data types.")$

o124 > ♥ 1244.py > ♡ merge_sort 42 def _merge(left: list, right: list) -> list:



Quick Sort: Uses a **pivot** to partition the array into smaller and larger elements, then sorts each part recursively. It's **faster on average** ($O(n \log n)$) but may degrade to $O(n^2)$ in the worst case.

Merge Sort: Divides the array into halves, sorts them, and then **merges** them. It always runs in **O(n log n)** time but uses **extra memory** for merging.

TASK-05:

AI-Suggested Algorithm Optimization

PROMPT:

Generate the python code which implements the duplicate search.

CODE:

```
| Section | Sect
```

```
PSC:\Users\ramch\OneDrivo\Desktop\ai> & C:\Users\ramch\AppOuta\Local\Programs\Python\Python312\python.exe c:\Users\ramch\OneDrivo\Desktop\ai/lab12.4/12.4.5.py

**Propriet = Finder Performance Comparison (List Size: 5000) ---

Brute-Force Algorithm:
Time taken: 0.413785 seconds
Found 1461 unique duplicates.

---

Optimized Algorithm (using sets):
Time taken: 0.005961 seconds
Found 1461 unique duplicates.

Observation: The optimized version is significantly faster for large lists.

Observation: The optimized version is significantly faster for large lists.

OFS C:\Users\ramch\OneOrive\Desktop\ai>
```

The task involves first writing a naive duplicate-finding algorithm using nested loops, which has $O(n^2)$ complexity. Then, AI can optimize it by using a set or dictionary to track seen elements, reducing the complexity to O(n). Students compare execution times on large inputs and explain that the optimization improves efficiency by avoiding repeated comparisons.