CS 33

Signals Part 2

Job Control

```
$ who
    foreground job
$ multiprocessProgram
    foreground job
^Z
stopped
$ bg
[1] multiprocessProgram &

    multiprocessProgram becomes background job 1

$ longRunningProgram &
[2]
$ fg %1
multiprocessProgram

    multiprocessProgram is now the foreground job

^C
```

Process Groups

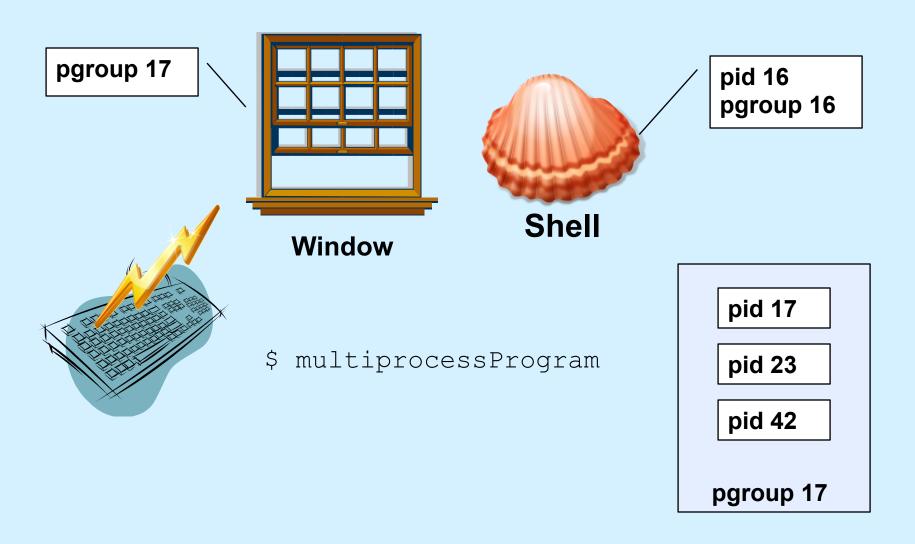
- Set of processes sharing the window/ keyboard
 - sometimes called a job
- Foreground process group/job
 - currently associated with window/keyboard
 - receives keyboard-generated signals
- Background process group/job
 - not currently associated with window/keyboard
 - doesn't currently receive keyboard-generated signals

Keyboard-Generated Signals

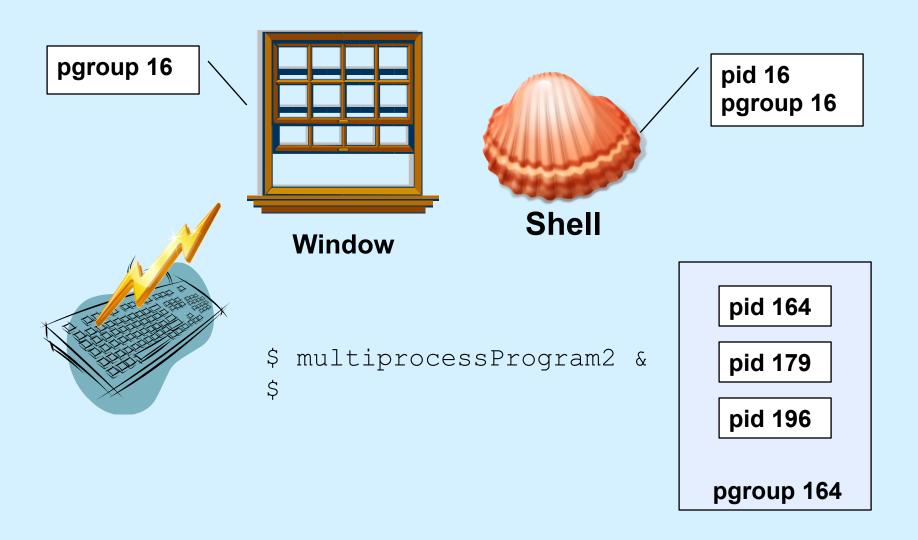
- You type ctrl-C
- How does the system know which process(es) to send the signal to?



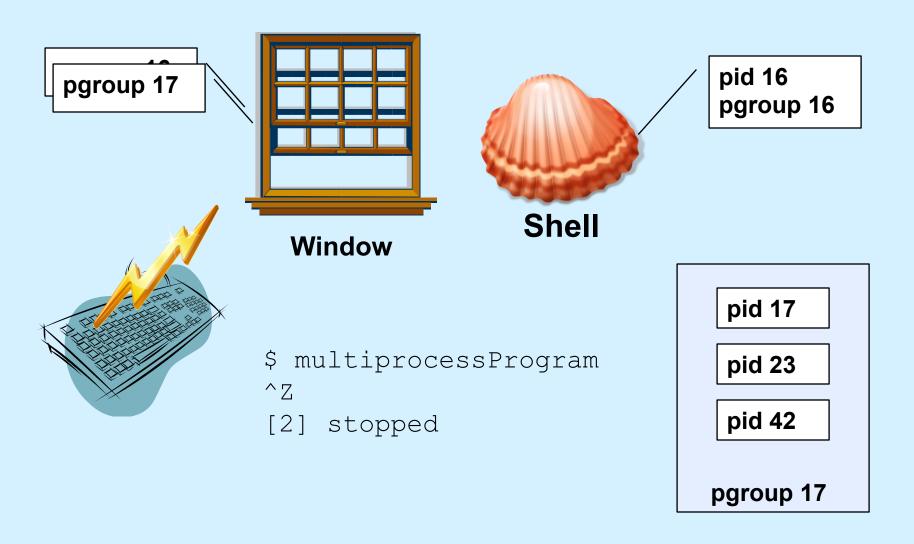
Foreground Job



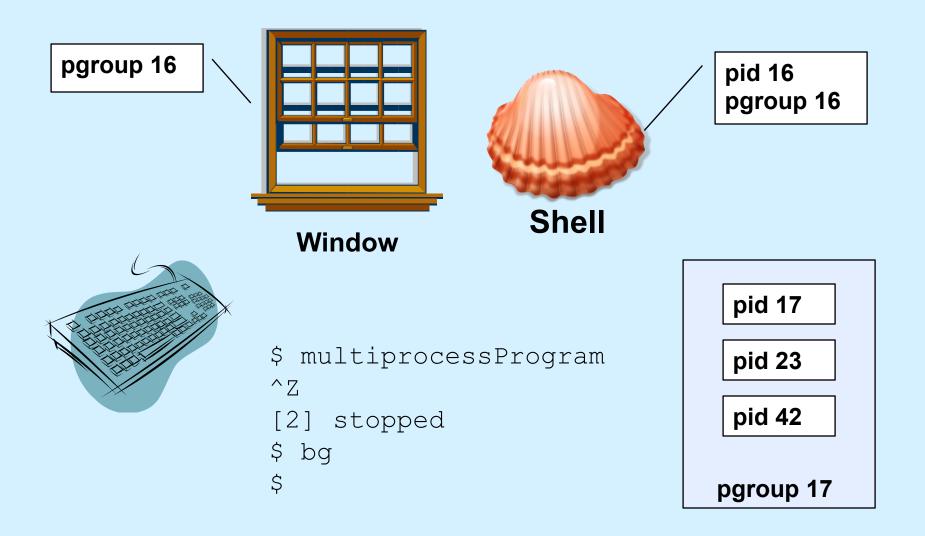
Background Job



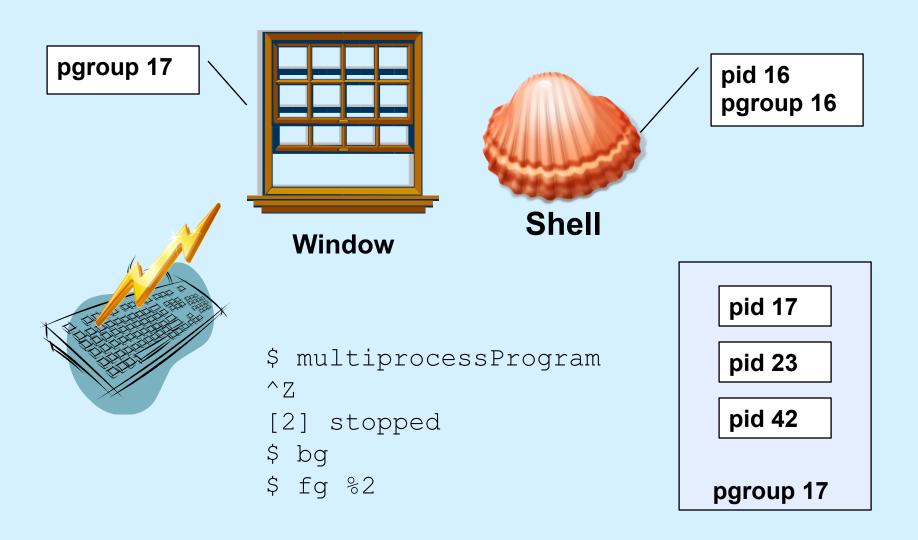
Stopping a Foreground Job



Backgrounding a Stopped Job



Foregrounding a Job



Quiz 1

```
$ long_running_prog1 &
$ long_running_prog2
^Z
[2] stopped
$ ^C
Which property
```

Which process group receives the SIGINT signal?

- a) the one containing the shell
- b) the one containing long_running_prog1
- c) the one containing long_running_prog2

Creating a Process Group

```
if (fork() == 0) {
  // child
  setpgid(0, 0);
     /* puts current process into a
        new process group whose ID is
        the process's pid.
        Children of this process will be in
        this process's process group.
     * /
  execv(...);
// parent
```

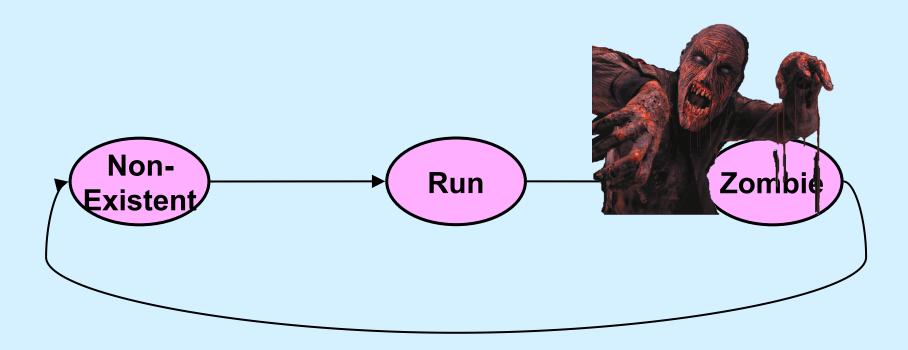
Setting the Foreground Process Group

```
tcsetpgrp(fd, pgid);
  // sets the process group of the
  // terminal (window) referenced by
  // file descriptor fd to be pgid
```

Kill: Details

- int kill(pid_t pid, int sig)
 - if pid > 0, signal sig sent to process pid
 - if pid == 0, signal sig sent to all processes in the caller's process group
 - if pid == -1, signal sig sent to all processes in the system for which sender has permission to do so
 - if pid < −1, signal sig is sent to all processes in process group −pid

Process Life Cycle



Reaping: Zombie Elimination

- · Shell must call waitpid on each child
 - easy for foreground processes
 - what about background?

- wait (&status) is equivalent to waitpid (-1, &status, 0)

(continued)

```
pid_t waitpid(pid_t pid, int *status, int options);
```

- options are some combination of the following
 - » WNOHANG
 - return immediately if no child has exited (returns 0)
 - » WUNTRACED
 - also return if a child has stopped (been suspended)
 - **» WCONTINUED**
 - also return if a child has been continued (resumed)

When to Call waitpid

- Shell reports status only when it is about to display its prompt
 - thus sufficient to check on background jobs just before displaying prompt

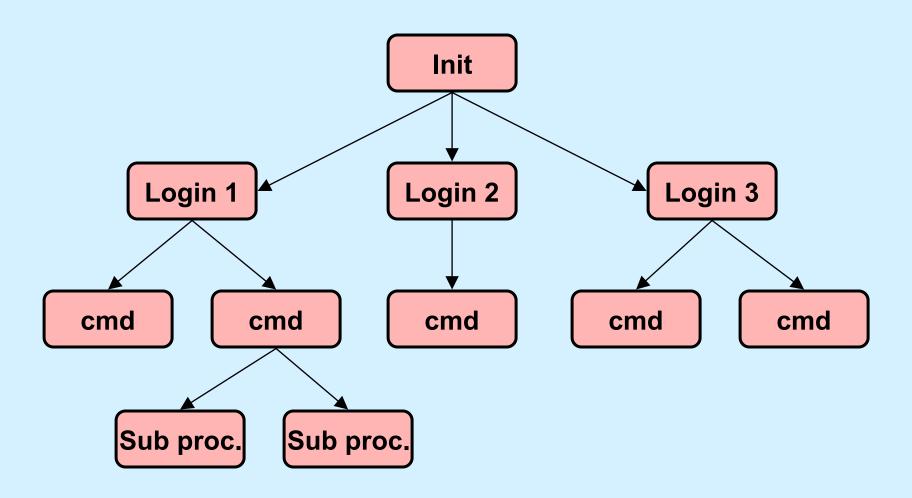
waitpid status

- WIFEXITED(*status): 1 if the process terminated normally and 0 otherwise
- WEXITSTATUS(*status): argument to exit
- WIFSIGNALED(*status): 1 if the process was terminated by a signal and 0 otherwise
- WTERMSIG(*status): the signal which terminated the process if it terminated by a signal
- WIFSTOPPED(*status): 1 if the process was stopped by a signal
- WSTOPSIG(*status): the signal which stopped the process if it was stopped by a signal
- WIFCONTINUED(*status): 1 if the process was resumed by SIGCONT and 0 otherwise

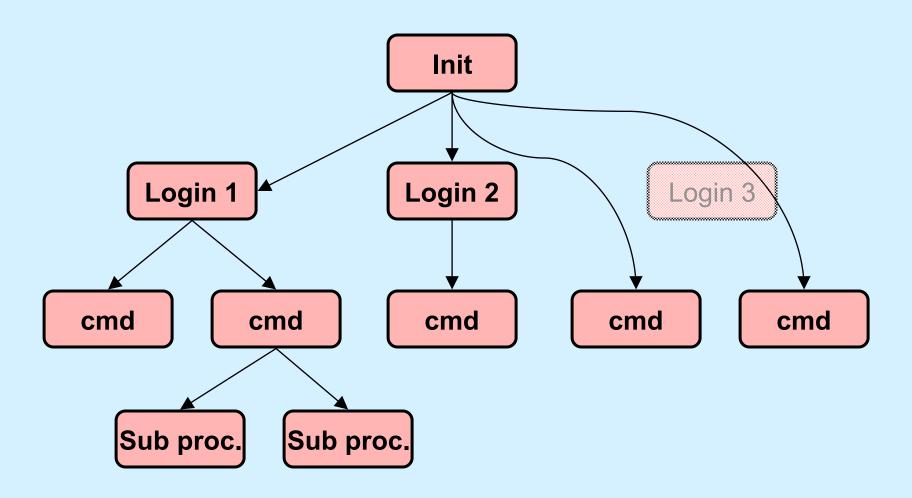
Example (in Shell)

```
int wret, status;
while ((wret = waitpid(-1, &wstatus, WNOHANG|WUNTRACED)) > 0) {
  // examine all children who've terminated or stopped
  if (WIFEXITED(wstatus)) {
    // terminated normally
  if (WIFSIGNALED(wstatus)) {
    // terminated by a signal
  if (WIFSTOPPED(wstatus)) {
    // stopped
```

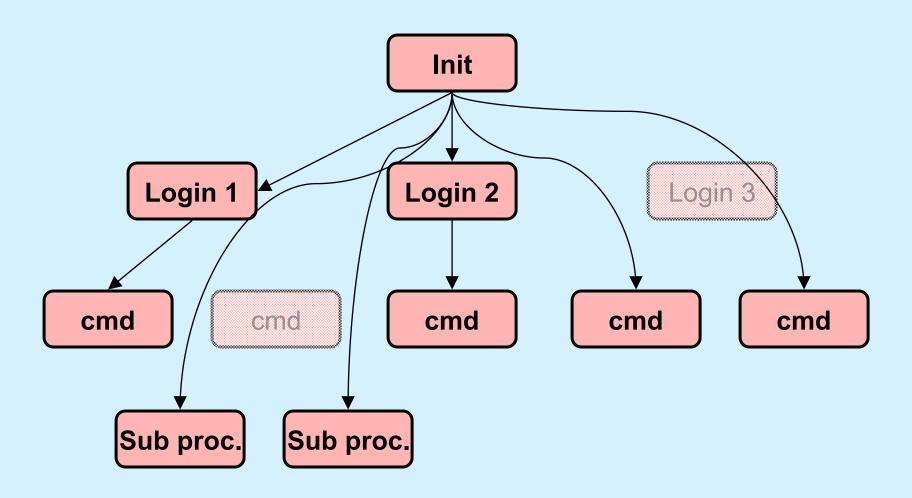
Process Relationships (1)



Process Relationships (2)



Process Relationships (3)



Signals, Fork, and Exec

```
// set up signal handlers ...
if (fork() == 0) {
   // what happens if child gets signal?
   signal (SIGINT, SIG IGN);
   signal(SIGFPE, handler);
   signal(SIGQUIT, SIG DFL);
   execv("new prog", argv, NULL);
   // what happens if SIGINT, SIGFPE,
   // or SIGQUIT occur?
```

Dealing with Failure

- fork, execv, wait, kill directly invoke the operating system
- Sometimes the OS says no
 - usually because you did something wrong
 - sometimes because the system has run out of resources
 - system calls return −1 to signify a problem

Reporting Failure

Integer error code placed in global variable errno

```
int errno;
```

- "man 3 errno" lists all possible error codes and meanings
- to print out meaning of most recent error

```
perror("message");
```

Fork

```
int main() {
 pid_t pid;
 while(1) {
    if ((pid = fork()) == -1) {
      perror("fork");
      exit(1);
```

Exec

```
int main() {
  if (fork() == 0) {
      char *argv[] = {"garbage", 0};
      execv("/garbage", argv);
      /* if we get here, there was an
         error! */
      perror("execv");
      exit(1);
```

Signals and Blocking System Calls

- What if a signal is generated while a process is blocked in a system call?
 - 1) deal with it when the system call completes
 - 2) interrupt the system call, deal with signal, resume system call

or

3) interrupt system call, deal with signal, return from system call with indication that something happened

Interrupted System Calls

```
while (read(fd, buffer, buf_size) == -1) {
   if (errno == EINTR) {
      /* interrupted system call - try again */
      continue;
   }
   /* the error is more serious */
   perror("big trouble");
   exit(1);
}
```

Timed Out, Revisited

```
void timeout(int sig) {}
                                       if (read(2, password,
                                              128)) == -1) {
int main() {
                                           if (errno == EINTR) {
  struct sigaction act;
                                             fprintf(stderr,
  sigemptyset(&act.sa mask);
                                               "Timed out\n");
  act.sa flags = 0;
                                             exit(1);
  act.sa handler = timeout;
  sigaction (SIGALRM, &act,
                                            perror("read");
    NULL);
                                            exit(1);
  alarm(10);
                                       alarm(0);
  char password[128];
                                       UsePassword(password);
                                       return 0;
```

Quiz 2

```
int ret;
char buf[128] = fillbuf();
ret = write(1, buf, 128);
```

- The value of ret is:
 - a) either -1 or 128
 - b) either -1, 0, or 128
 - c) any integer in the range [-1, 128]

Interrupted While Underway

```
if (num xfrd < remaining) {</pre>
remaining = total count;
bptr = buf;
                                   /* interrupted after the
for (;;) {
                                      first step */
                                   remaining -= num xfrd;
  num xfrd = write(fd, bptr,
                                   bptr += num xfrd;
      remaining);
                                   continue;
  if (num xfrd == -1) {
    if (errno == EINTR) {
                                /* success! */
     /* interrupted early */
                                 break:
     continue;
    perror("big trouble");
    exit(1);
```

Automatic Restart

```
void reap_child(int sig) {
   printf("bye bye\n");
}

struct sigaction act;
act.sa_handler = reap_child;
sigemptyset(&act.sa_mask);
act.sa_flags = SA_RESTART;
sigaction(SIGCHLD, &act, 0);
```

```
remaining = total count;
bptr = buf;
while ((num xfrd =
    write(fd, bptr, remaining))
      != remaining) {
  if (num xfrd == -1) {
    /* no EINTR from SIGCHLD */
    break;
  /* still must deal with
     partial completions */
  remaining -= num xfrd;
  bptr += num xfrd;
```

Asynchronous Signals (1)

```
main() {
  void handler(int);
  signal(SIGINT, handler);
   ... /* long-running buggy code */
void handler(int sig) {
   ... /* die gracefully */
  exit(1);
```

Asynchronous Signals (2)

Asynchronous Signals (3)

Asynchronous Signals (4)

```
char buf[BSIZE];
int pos;
void myput(char *str) {
  int i;
  int len = strlen(str);
  for (i=0; i<len; i++, pos++) {
   buf[pos] = str[i];
    if ((buf[pos] == '\n') || (pos == BSIZE-1)) {
      write(1, buf, pos+1);
     pos = -1;
```

Async-Signal Safety

 Which library routines are safe to use within signal handlers?

_	abort	_	dup2	_	getppid	_	readlink	_	sigemptyset	_	tcgetpgrp
_	accept	_	execle	_	getsockname	_	recv	_	sigfillset	_	tcsendbreak
_	access	_	execve	_	getsockopt	_	recvfrom	_	sigismember	_	tcsetattr
_	aio_error	_	_exit	_	getuid	_	recvmsg	_	signal	_	tcsetpgrp
_	aio_return	_	fchmod	_	kill	-	rename	_	sigpause	_	time
_	aio_suspend	_	fchown	_	link	_	rmdir	_	sigpending	_	timer_getoverrun
_	alarm	_	fcntl	_	listen	_	select	_	sigprocmask	_	timer_gettime
_	bind	_	fdatasync	_	lseek	-	sem_post	_	sigqueue	_	timer_settime
_	cfgetispeed	_	fork	_	lstat	-	send	_	sigsuspend	_	times
_	cfgetospeed	_	fpathconf	_	mkdir	-	sendmsg	_	sleep	_	umask
-	cfsetispeed	-	fstat	_	mkfifo	_	sendto	-	sockatmark	_	uname
_	cfsetospeed	_	fsync	_	open	_	setgid	_	socket	_	unlink
_	chdir	_	ftruncate	_	pathconf	-	setpgid	_	socketpair	_	utime
_	chmod	_	getegid	_	pause	-	setsid	_	stat	_	wait
_	chown	_	geteuid	_	pipe	-	setsockopt	_	symlink	_	waitpid
-	clock_gettime	-	getgid	_	poll	_	setuid	-	sysconf	_	write
_	close	_	getgroups	_	posix_trace_even	t–	shutdown	_	tcdrain		
_	connect	_	getpeername	_	pselect	_	sigaction	_	tcflow		
_	creat	_	getpgrp	_	raise	-	sigaddset	_	tcflush		
_	dup	_	getpid	_	read	-	sigdelset	_	tcgetattr		

Quiz 3

Printf is not required to be async-signal safe. Can it be implemented so that it is?

- a) no, it's inherently not async-signal safe
- b) yes, but it would be so complicated, it's not done
- c) yes, it can be easily made async-signal safe