CS 33

Architecture and Optimization (1)

Many of the slides in this lecture are either from or adapted from slides provided by the authors of the textbook "Computer Systems: A Programmer's Perspective," 2nd Edition and are provided from the website of Carnegie-Mellon University, course 15-213, taught by Randy Bryant and David O'Hallaron in Fall 2010. These slides are indicated "Supplied by CMU" in the notes section of the slides.

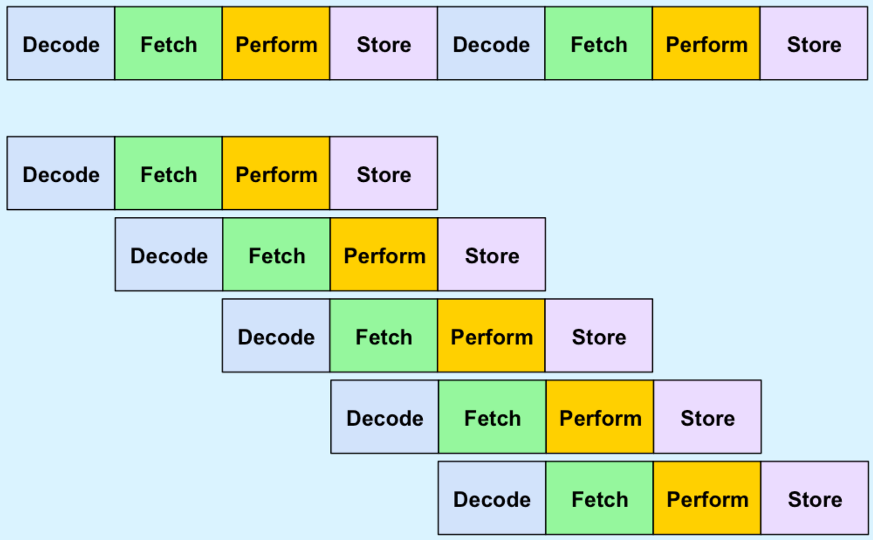
Simplistic View of Processor

```
while (true) {
  instruction = mem[eip];
  execute(instruction);
}
```

Some Details ...

```
void execute(instruction_t instruction) {
  decode(instruction, &opcode, &operands);
  fetch(operands, &in_operands);
  perform(opcode, in_operands, &out_operands);
  store(out_operands);
}
```

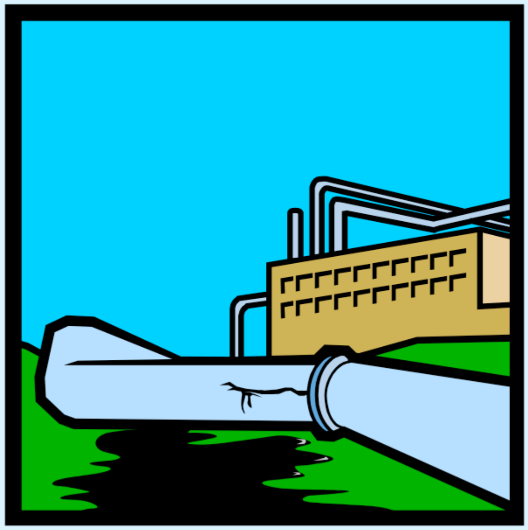
Pipelines


Decode
Fetch
Perform
Store
Decode
Fetch
Perform
Store
Decode
Fetch
Perform
Store
Decode
Fetch
Perform
Store
Decode
Fetch
Perform
Store
Decode
Fetch
Perform
Store
Decode
Fetch
Perform
Store

Analysis

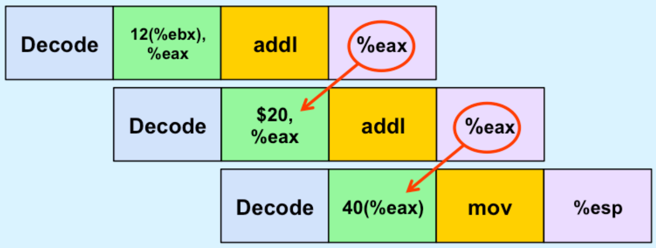
- **Not pipelined**
  - each instruction takes, say, 320 nanoseconds
    - » 320 ns latency
  - 3.125 billion instructions/second (GIPS)
- **Pipelined**
  - each instruction still takes 320 ns
    - » latency still 320 ns
  - an instruction completes every 80 ns
    - » 12.5 GIPS throughput

Hazards ...

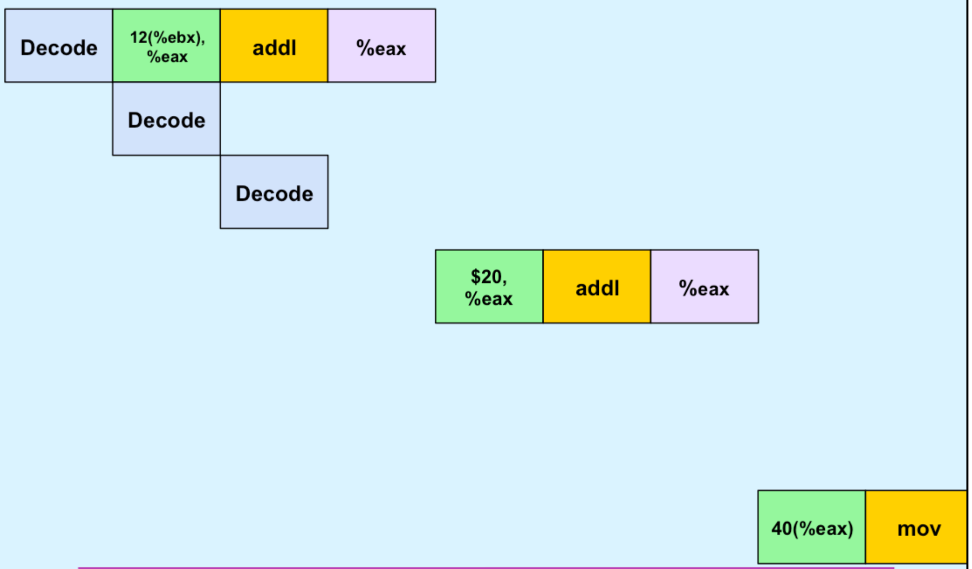


Data Hazards

```
addl 12(%ebx), %eax
addl $20, %eax
movl 40(%eax), %esp
```



Coping



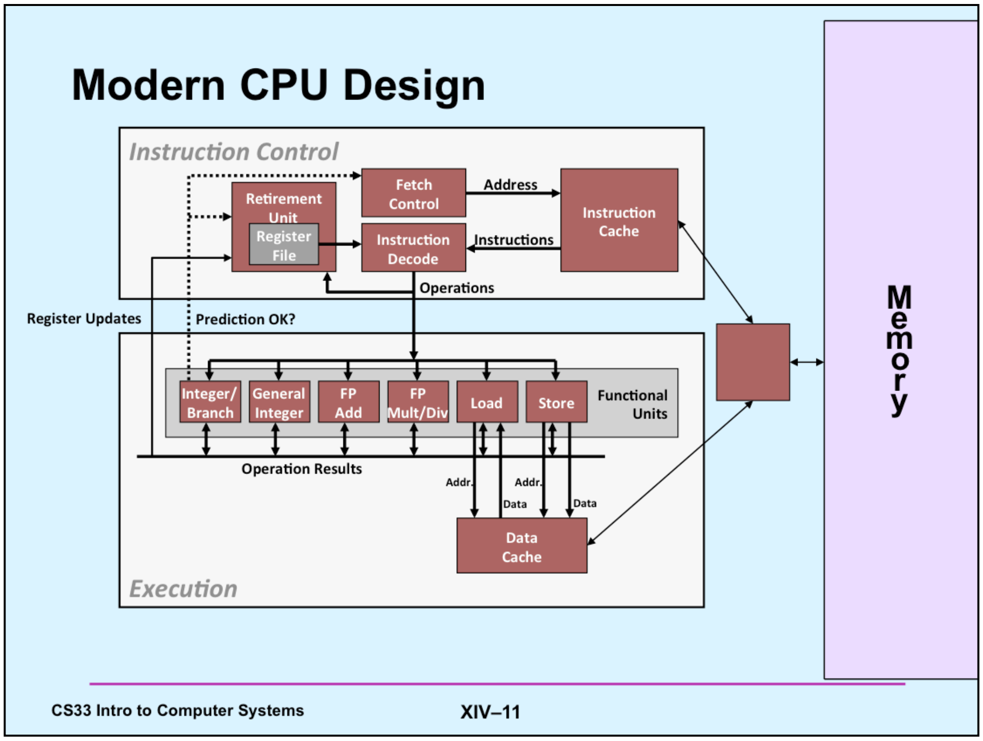
Decode
12(%ebx),
%eax
addl
%eax
Decode
Decode
$20,
%eax
addl
%eax
40(%eax)
mov

Control Hazards

```
  movl $0, %ecx
.L2:
  movl %edx, %eax
  andl $1, %eax
  addl %eax, %ecx
  shrl $1, %edx
  jne  .L2 # what goes in the pipeline?
  movl %ecx, %eax
  ...
```


Coping: Guess ...

- Branch prediction
  - assume, for example, that conditional branches are always taken
  - but don't do anything to registers or memory until you know for sure



Adapted from slide supplied by CMU.

Performance Realities

There's more to performance than asymptotic complexity

- **Constant factors matter too!**
  - easily see 10:1 performance range depending on how code is written
  - must optimize at multiple levels:
    - » algorithm, data representations, procedures, and loops
- **Must understand system to optimize performance**
  - how programs are compiled and executed
  - how to measure program performance and identify bottlenecks
  - how to improve performance without destroying code modularity and generality

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Optimizing Compilers

- Provide efficient mapping of program to machine
  - register allocation
  - code selection and ordering (scheduling)
  - dead code elimination
  - eliminating minor inefficiencies
- Don’t (usually) improve asymptotic efficiency
  - up to programmer to select best overall algorithm
  - big-O savings are (often) more important than constant factors
    - » but constant factors also matter
- Have difficulty overcoming “optimization blockers”
  - potential memory aliasing
  - potential procedure side-effects

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Limitations of Optimizing Compilers

- Operate under fundamental constraint
  - must not cause any change in program behavior
  - often prevents it from making optimizations that would only affect behavior under pathological conditions
- Behavior that may be obvious to the programmer can be obfuscated by languages and coding styles
  - e.g., data ranges may be more limited than variable types suggest
- Most analysis is performed only within procedures
  - whole-program analysis is too expensive in most cases
- Most analysis is based only on *static* information
  - compiler has difficulty anticipating run-time inputs

- When in doubt, the compiler must be conservative

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Generally Useful Optimizations

- **Optimizations that you or the compiler should do regardless of processor / compiler**

- **Code Motion**
  - reduce frequency with which computation performed
    - » if it will always produce same result
    - » especially moving code out of loop

```
void set_row(long *a, long *b,
   long i, long n){
    long j;
    for (j = 0; j < n; j++)
        a[n*i+j] = b[j];
}
```

→

```
    long j;
    int ni = n*i;
    for (j = 0; j < n; j++)
        a[ni+j] = b[j];
```


Supplied by CMU.

Compiler-Generated Code Motion

```
void set_row(long *a, long *b,
   long i, long n){
    long j;
    for (j = 0; j < n; j++)
         a[n*i+j] = b[j];
}
```

```
    long j;
    long ni = n*i;
    long *rowp = a+ni;
    for (j = 0; j < n; j++)
         rowp[j] = b[j];
```

```
set_row:
        testq   %rcx, %rcx               # Test n
        jle     .L1                      # If 0, goto done
        imulq   %rcx, %rdx               # i *= n
        leaq    (%rdi,%rdx,8), %rdi      # rowp = A + n*i*8
        movl    $0, %eax                 # j = 0
.L3:                                    # loop:
        movq    (%rsi,%rax,8), %rdx      # t = b[j]
        movq    %rdx, (%rdi,%rax,8)      # rowp[j] = t
        addq    $1, %rax                 # j++
        cmpq    %rcx, %rax               # Compare n:j
        jg      .L3                      # If >, goto loop
.L1:                                    # done:
        rep ; ret
```


Supplied by CMU, updated for current gcc.

Reduction in Strength

- Replace costly operation with simpler one
- Shift, add instead of multiply or divide

  `16*x      -->   x << 4`

  – utility is machine-dependent

  – depends on cost of multiply or divide instruction

    » on Intel Nehalem, integer multiply requires 3 CPU cycles
- Recognize sequence of products

```
for (i = 0; i < n; i++)
  for (j = 0; j < n; j++)
    a[n*i + j] = b[j];
```

→

```
int ni = 0;
for (i = 0; i < n; i++) {
  for (j = 0; j < n; j++)
    a[ni + j] = b[j];
  ni += n;
}
```


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Share Common Subexpressions

- Reuse portions of expressions
- Compilers often not very sophisticated in exploiting arithmetic properties

```
/* Sum neighbors of i,j */
up =    val[(i-1)*n + j  ];
down =  val[(i+1)*n + j  ];
left =  val[i*n      + j-1];
right = val[i*n      + j+1];
sum = up + down + left + right;
```

```
long inj = i*n + j;
up =    val[inj - n];
down =  val[inj + n];
left =  val[inj - 1];
right = val[inj + 1];
sum = up + down + left + right;
```

3 multiplications: i*n, (i–1)*n, (i+1)*n

1 multiplication: i*n

```
leaq   1(%rsi), %rax   # i+1
leaq   -1(%rsi), %r8   # i-1
imulq  %rcx, %rsi      # i*n
imulq  %rcx, %rax      # (i+1)*n
imulq  %rcx, %r8       # (i-1)*n
addq   %rdx, %rsi      # i*n+j
addq   %rdx, %rax      # (i+1)*n+j
addq   %rdx, %r8       # (i-1)*n+j
```

```
imulq    %rcx, %rsi   # i*n
addq     %rdx, %rsi   # i*n+j
movq     %rsi, %rax   # i*n+j
subq     %rcx, %rax   # i*n+j-n
leaq     (%rsi,%rcx), %rcx # i*n+j+n
```


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Quiz 1

The fastest means (on the Intel Nehalem) for evaluating

```
n*n + 2*n + 1
```

requires exactly:

a) 2 multiplies and 2 additions
b) one multiply and two additions
c) one multiply and one addition
d) three additions

Optimization Blocker #1: Procedure Calls

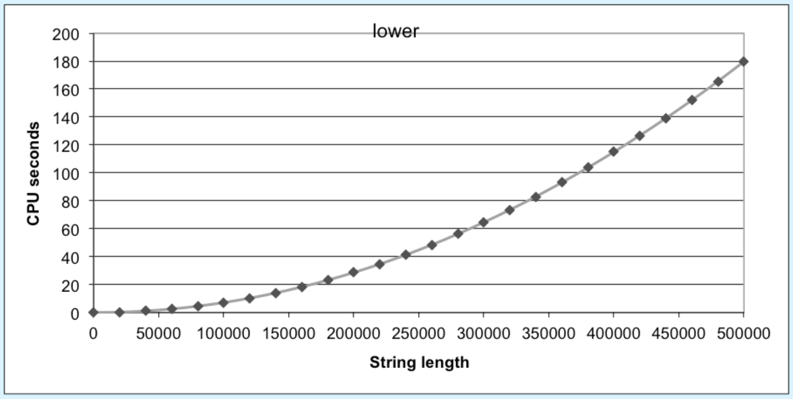
- **Procedure to convert string to lower case**

```
void lower(char *s){
  int i;
  for (i = 0; i < strlen(s); i++)
    if (s[i] >= 'A' && s[i] <= 'Z')
      s[i] -= ('A' - 'a');
}
```


Supplied by CMU.

Lower Case Conversion Performance

- Time quadruples when string length doubles
- Quadratic performance



Supplied by CMU.

Convert Loop To Goto Form

```
void lower(char *s){
   int i = 0;
   if (i >= strlen(s))
     goto done;
 loop:
   if (s[i] >= 'A' && s[i] <= 'Z')
       s[i] -= ('A' - 'a');
   i++;
   if (i < strlen(s))
     goto loop;
 done:
}
```

- **`strlen` executed every iteration**

Supplied by CMU.

Calling Strlen

```
size_t strlen(const char *s){
    size_t length = 0;
    while (*s != '\0') {
       s++;
       length++;
    }
    return length;
}
```

- strlen performance
  - only way to determine length of string is to scan its entire length, looking for null character
- Overall performance, string of length N
  - N calls to strlen
  - overall $O(N^2)$ performance

Supplied by CMU.

Improving Performance

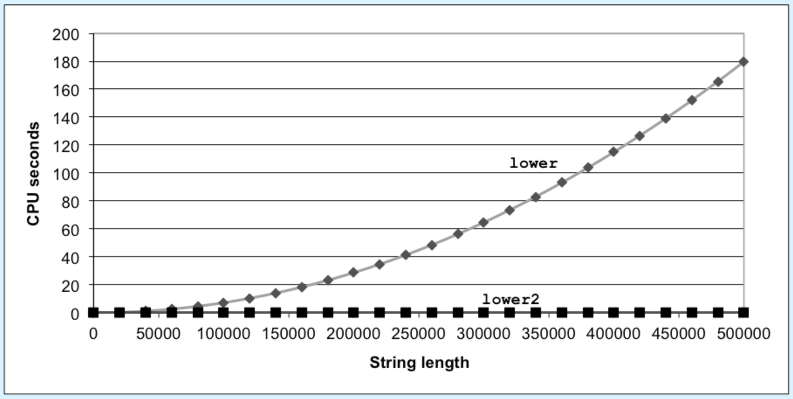
```
void lower2(char *s){
  int i;
  int len = strlen(s);
  for (i = 0; i < len; i++)
    if (s[i] >= 'A' && s[i] <= 'Z')
      s[i] -= ('A' - 'a');
}
```

- **Move call to `strlen` outside of loop**
  - **since result does not change from one iteration to another**
  - **form of code motion**

Supplied by CMU.

Lower-Case Conversion Performance

- **Time doubles when string-length doubles**
  - **linear performance of lower2**



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Optimization Blocker: Procedure Calls

- ***Why couldn't compiler move `strlen` out of inner loop?***
  - procedure may have side effects
    - » alters global state each time called
  - function may not return same value for given arguments
    - » depends on other parts of global state
    - » procedure `lower` could interact with `strlen`
- **Warning:**
  - compiler treats procedure call as a black box
  - weak optimizations near them
- **Remedies:**
  - use of `inline` functions
    - » gcc does this with –O2
  - do your own code motion

```
int lencnt = 0;
size_t strlen(const char *s){
    size_t length = 0;
    while (*s != '\0') {
       s++; length++;
    }
    lencnt += length;
    return length;
}
```


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Memory Matters

```
/* Sum rows of n X n matrix a
   and store result in vector b  */
void sum_rows1(long *a, long *b, long n) {
    long i, j;
    for (i = 0; i < n; i++) {
        b[i] = 0;
        for (j = 0; j < n; j++)
            b[i] += a[i*n + j];
    }
}
```

```
# sum_rows1 inner loop
.L3:
   movq    (%rdi), %rcx         # rcx = *aptr
   addq    %rcx, (%rsi,%rax,8)  # b[i] += rcx
   addq    $8, %rdi             # aptr++
   cmpq    %r8, %rdi
   jne     .L3
```

- **Code updates `b[i]` on every iteration**
- **Why couldn't compiler optimize this away?**

Supplied by CMU, updated for current gcc.

Note that *a* is passed as a 1-D array, but interpreted as a 2-D array. This isn't terribly good programming style (gcc, fortunately, refrains from commenting on one's style), but it is definitely the sort of program that gcc must be prepared to deal with.

Memory Aliasing

```
/* Sum rows of n X n matrix a
   and store result in vector b  */
void sum_rows1(int *a, int *b, long n) {
    long i, j;
    for (i = 0; i < n; i++) {
        b[i] = 0;
        for (j = 0; j < n; j++)
            b[i] += a[i*n + j];
    }
}
```

```
int A[9] =
  { 0,   1,   2,
    4,   8,  16,
   32,  64, 128};

int *B = &A[3];

sum_rows1(A, B, 3);
```

Value of B:

```
init:  [4, 8, 16]
```

```
i = 0: [3, 8, 16]
```

```
i = 1: [3, 22, 16]
```

```
i = 2: [3, 22, 224]
```

- **Code updates `b[i]` on every iteration**
- **Must consider possibility that these updates will affect program behavior**

Supplied by CMU, updated for current gcc.

Removing Aliasing

```
/* Sum rows of n X n matrix a
   and store result in vector b  */
void sum_rows2(int *a, int *b, int n) {
    long i, j;
    for (i = 0; i < n; i++) {
         int val = 0;
         for (j = 0; j < n; j++)
             val += a[i*n + j];
         b[i] = val;
    }
}
```

```
# sum_rows2 inner loop
.L4:
   addq    (%rdi), %rax
   addq    $8, %rdi
   cmpq    %rcx, %rdi
   jne     .L4
```

- **No need to store intermediate results**

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Optimization Blocker: Memory Aliasing

- Aliasing
  - two different memory references specify single location
  - easy to have happen in C
    - » since allowed to do address arithmetic
    - » direct access to storage structures
  - get in habit of introducing local variables
    - » accumulating within loops
    - » your way of telling compiler not to check for aliasing

Supplied by CMU.

C99 to the Rescue

- **New attribute**
  - **restrict**
    - » applied to a pointer, tells the compiler that the object pointed to will be accessed only via this pointer
    - » compiler thus doesn't have to worry about aliasing
    - » but the programmer does ...
    - » syntax

```
int *restrict pointer;
```

Memory Matters, Fixed

```
/* Sum rows of n X n matrix a
   and store result in vector b  */
void sum_rows3(long *restrict a, long *restrict b, long n) {
    long i, j;
    for (i = 0; i < n; i++) {
        b[i] = 0;
        for (j = 0; j < n; j++)
            b[i] += a[i*n + j];
    }
}
```

```
# sum_rows1 inner loop
.L3:
   addq    (%rdi), %rax
   addq    $8, %rdi
   cmpq    %rcx, %rdi
   jne     .L3
```

- **Code doesn't update `b[i]` on every iteration**

Note: we must give gcc the flag "-std=gnu99" for this to be compiled.

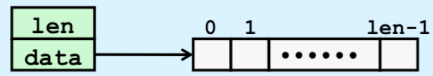
Exploiting Instruction-Level Parallelism

- Need general understanding of modern processor design
  - hardware can execute multiple instructions in parallel
- Performance limited by data dependencies
- Simple transformations can have dramatic performance improvement
  - compilers often cannot make these transformations
  - lack of associativity and distributivity in floating-point arithmetic

Supplied by CMU.

Benchmark Example: Datatype for Vectors

```
/* data structure for vectors */
typedef struct{
    int len;
    data_t *data;
} vec_t, *vec_ptr_t;
```




```
/* retrieve vector element and store at val */
int get_vec_element(vec_ptr_t v, int idx, data_t *val){
    if (idx < 0 || idx >= v->len)
        return 0;
    *val = v->data[idx];
    return 1;
}

/* return length of vector */
int vec_length(vec_ptr_t v) {
     return v->len;
}
```


Supplied by CMU.

Benchmark Computation

```
void combine1(vec_ptr_t v, data_t *dest){
    long int i;
    *dest = IDENT;
    for (i = 0; i < vec_length(v); i++) {
      data_t val;
      get_vec_element(v, i, &val);
      *dest = *dest OP val;
    }
}
```

Compute sum or product of vector elements

- Data Types
  - use different declarations for `data_t`
    - » `int`
    - » `float`
    - » `double`
- Operations
  - use different definitions of `OP` and `IDENT`
    - » `+, 0`
    - » `*, 1`

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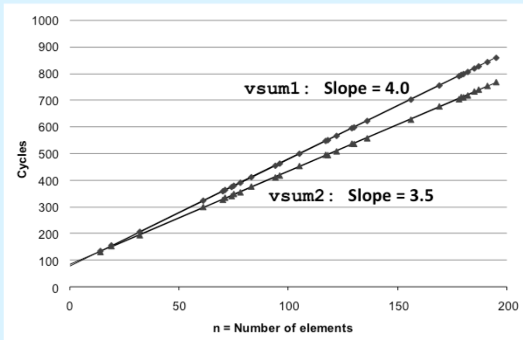
Cycles Per Element (CPE)

- Convenient way to express performance of program that operates on vectors or lists
- Length = n
- T = CPE*n + Overhead
  - CPE is slope of line



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Benchmark Performance

```
void combine1(vec_ptr_t v, data_t *dest){
    long int i;
    *dest = IDENT;
    for (i = 0; i < vec_length(v); i++) {
      data_t val;
      get_vec_element(v, i, &val);
      *dest = *dest OP val;
    }
}
```

Compute sum or product of vector elements

Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Combine1 unoptimized	29.0	29.2	27.4	27.9
Combine1 –O1	12.0	12.0	12.0	13.0

Supplied by CMU.

Move vec_length

```
void combine2(vec_ptr_t v, data_t *dest){
    long int i;
    long int length = vec_length(v);
    *dest = IDENT;
    for (i = 0; i < length; i++) {
       data_t val;
       get_vec_element(v, i, &val);
       *dest = *dest OP val;
    }
}
```

Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Combine1 unoptimized	29.0	29.2	27.4	27.9
Combine1 –O1	12.0	12.0	12.0	13.0
Combine2	8.03	8.09	10.09	12.08

Supplied by CMU.

Eliminate Procedure Calls

```
void combine3(vec_ptr_t v, data_t *dest){
    long int i;
    long int length = vec_length(v);
    data_t *data = get_vec_start(v);
    *dest = IDENT;
    for (i = 0; i < length; i++) {
       *dest = *dest OP data[i];
    }
}
```

```
data_t *get_vec_start(
     vec_ptr v) {
  return v->data;
}
```

Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Combine2	8.03	8.09	10.09	12.08
Combine3	6.01	8.01	10.01	12.02

Supplied by CMU.

Eliminate Unneeded Memory References

```
void combine4(vec_ptr_t v, data_t *dest){
  int i;
  int length = vec_length(v);
  data_t *d = get_vec_start(v);
  data_t t = IDENT;
  for (i = 0; i < length; i++)
    t = t OP d[i];
  *dest = t;
}
```

Method	Integer		Double FP	
Operation	Add	Mult	Add	Mult
Combine1 –O1	12.0	12.0	12.0	13.0
Combine4	2.0	3.0	3.0	5.0

Supplied by CMU.

Quiz 2

Combine4 is pretty fast; we've done all the "obvious" optimizations. How much faster will we be able to make it? (Hint: it involves taking advantage of pipelining and multiple functional units on the chip.)

a) 1× (it's already as fast as possible)
b) 2× – 4×
c) 16× – 64×