
STACK-NUMBER IS NOT BOUNDED BY QUEUE-NUMBER

Vida Dujmović[‡], David Eppstein[§], Robert Hickingbotham[¶], Pat Morin[‡], David R. Wood[‡]

November 2, 2020

ABSTRACT. We describe a family of graphs with queue-number at most 4 but unbounded stack-number. This resolves open problems of Heath, Leighton and Rosenberg (1992) and Blankenship and Oporowski (1999).

1 Introduction

Stacks and queues are fundamental data structures in computer science, but which is more powerful? In 1992, Heath, Leighton and Rosenberg [23, 24] introduced an approach for answering this question by defining the graph parameters *stack-number* and *queue-number* (defined below), which respectively measure the power of stacks and queues for representing graphs. The following fundamental problems, implicit in [23, 24], were made explicit by Dujmović and Wood [17]¹:

- Is stack-number bounded by queue-number?
- Is queue-number bounded by stack-number?

If stack-number is bounded by queue-number but queue-number is not bounded by stack-number, then stacks would be considered to be more powerful than queues. Similarly, if the converse holds, then queues would be considered to be more powerful than stacks. Despite extensive research on stack- and queue-numbers, these fundamental questions have remained unsolved.

We now formally define stack- and queue-number. Let G be a graph and let $<$ be a total order on $V(G)$. Two disjoint edges $vw, xy \in E(G)$ with $v < w$ and $x < y$ *cross* with respect to $<$ if $v < x < w < y$ or $x < v < y < w$, and *nest* with respect to $<$ if $v < x < y < w$ or $x < v < w < y$. Let $\varphi : E(G) \rightarrow \{1, \dots, k\}$ for some integer $k \geq 1$.

[‡]School of Computer Science and Electrical Engineering, University of Ottawa, Ottawa, Canada (vida.dujmovic@uottawa.ca). Research supported by NSERC.

[§]Department of Computer Science, University of California, Irvine, California, USA (eppstein@uci.edu).

[¶]School of Mathematics, Monash University, Melbourne, Australia (robert.hickingbotham@monash.edu).

[‡]School of Computer Science, Carleton University, Ottawa, Canada (morin@scs.carleton.ca). Research supported by NSERC.

[‡]School of Mathematics, Monash University, Melbourne, Australia (david.wood@monash.edu). Research supported by the Australian Research Council.

¹A *graph parameter* is a function α such that $\alpha(G) \in \mathbb{R}$ for every graph G and such that $\alpha(G_1) = \alpha(G_2)$ for all isomorphic graphs G_1 and G_2 . A graph parameter α is *bounded* by a graph parameter β if there exists a function f such that $\alpha(G) \leq f(\beta(G))$ for every graph G .

Then (\prec, φ) is a k -stack layout of G if vw and xy do not cross for all edges $vw, xy \in E(G)$ with $\varphi(vw) = \varphi(xy)$. Similarly, (\prec, φ) is a k -queue layout of G if vw and xy do not nest for all edges $vw, xy \in E(G)$ with $\varphi(vw) = \varphi(xy)$. The smallest integer s for which G has an s -stack layout is called the *stack-number* of G , denoted $\text{sn}(G)$. The smallest integer q for which G has a q -queue layout is called the *queue-number* of G , denoted $\text{qn}(G)$.

Given a k -stack layout (\prec, φ) of a graph G , for each $i \in \{1, \dots, k\}$, the set $\varphi^{-1}(i)$ behaves like a stack, in the sense that each edge $vw \in \varphi^{-1}(i)$ with $v \prec w$ corresponds to an element in a sequence of stack operations, such that if we traverse the vertices in the order of \prec , then vw is pushed onto the stack at v and popped off the stack at w . Similarly, each set $\varphi^{-1}(i)$ in a queue layout behaves like a queue. In this way, the stack-number and queue-number respectively measure the power of stacks and queues to represent graphs.

Note that stack layouts are equivalent to book embeddings (first defined by Ollmann [26] in 1973), and stack-number is also known as *page-number*, *book-thickness* or *fixed outer-thickness*. Stack and queue layouts have other applications including computational complexity [7, 8, 15, 21], RNA folding [22], graph drawing in two [1, 2, 31] and three dimensions [12–14, 32], and fault-tolerant multiprocessing [9, 28–30]. See [3–5, 10, 11, 16, 18, 25, 34, 35] for bounds on the stack- and queue-number for various graph classes.

Is Stack-Number Bounded by Queue-Number?

This paper considers the first of the above questions. In a positive direction, Heath et al. [23] showed that every 1-queue graph has a 2-stack layout. On the other hand, they described graphs that need exponentially more stacks than queues. In particular, n -vertex ternary hypercubes have queue-number $O(\log n)$ and stack-number $\Omega(n^{1/9-\epsilon})$ for any $\epsilon > 0$.

Our key contribution is the following theorem, which shows that stack-number is not bounded by queue-number.

Theorem 1. *For every $s \in \mathbb{N}$ there exists a graph G with $\text{qn}(G) \leq 4$ and $\text{sn}(G) > s$.*

This demonstrates that stacks are not more powerful than queues for representing graphs.

Cartesian Products

As illustrated in Figure 1, the graph G in Theorem 1 is the cartesian product² $S_b \square H_n$, where S_b is the star graph with root r and b leaves, and H_n is the dual of the hexagonal grid, defined by

$$\begin{aligned} V(H_n) := \{1, \dots, n\}^2 \quad \text{and} \quad E(H_n) := & \{(x, y)(x+1, y) : x \in \{1, \dots, n-1\}, y \in \{1, \dots, n\}\} \\ & \cup \{(x, y)(x, y+1) : x \in \{1, \dots, n\}, y \in \{1, \dots, n-1\}\} \\ & \cup \{(x, y)(x+1, y+1) : x, y \in \{1, \dots, n-1\}\} . \end{aligned}$$

²For graphs G_1 and G_2 , the *cartesian product* $G_1 \square G_2$ is the graph with vertex set $\{(v_1, v_2) : v_1 \in V(G_1), v_2 \in V(G_2)\}$, where $(v_1, v_2)(w_1, w_2) \in E(G_1 \square G_2)$ if $v_1 = w_1$ and $v_2 w_2 \in E(G_2)$, or $v_1 w_1 \in E(G_1)$ and $v_2 = w_2$. The *strong product* $G_1 \boxtimes G_2$ is the graph obtained from $G_1 \square G_2$ by adding the edge $(v_1, v_2)(w_1, w_2)$ whenever $v_1 w_1 \in E(G_1)$ and $v_2 w_2 \in E(G_2)$. Note that Pupyrev [27] independently suggested using graph products to show that stack-number is not bounded by queue-number.

Do we want a figure showing an example of a stack layout and a queue layout here? H_4 would be a good example. Then we could refer to it when we show $\text{sqn}(H_n) \leq 3$.

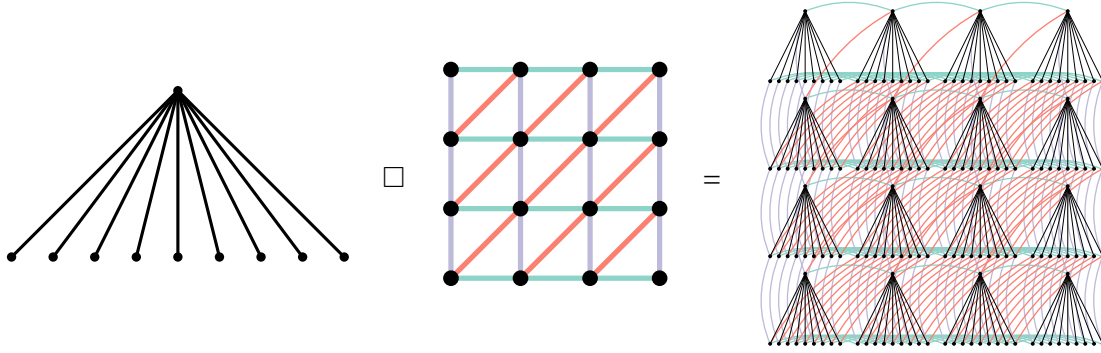


Figure 1: $S_9 \square H_4$.

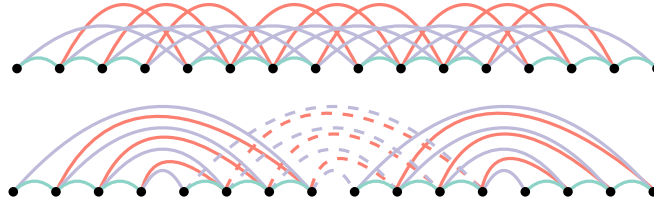


Figure 2: A strict 3-queue layout and a 2-stack layout of H_4

We prove the following:

Theorem 2. *For every $s \in \mathbb{N}$, if b and n are sufficiently large compared to s , then*

$$\text{sn}(S_b \square H_n) > s.$$

We now show that $\text{qn}(S_b \square H_n) \leq 4$, which with Theorem 2 implies Theorem 1. We need the following definition due to Wood [33]. A queue layout $(\varphi, <)$ is *strict* if for every vertex $u \in V(G)$ and for all neighbours $v, w \in N_G(u)$, if $u < v, w$ or $v, w < u$, then $\varphi(uv) \neq \varphi(uw)$. Let $\text{sqn}(G)$ be the minimum integer k such that G has a strict k -queue layout. To see that $\text{sqn}(H_n) \leq 3$, order the vertices row-by-row and then left-to-right within a row, with vertical edges in one queue, horizontal edges in one queue, and diagonal edges in another queue (see Figure 2). Wood [33] proved that for all graphs G_1 and G_2 ,

$$\text{qn}(G_1 \square G_2) \leq \text{qn}(G_1) + \text{sqn}(G_2). \quad (1)$$

Of course, S_b has a 1-queue layout (since no two edges are nested for any vertex-ordering). Thus $\text{qn}(S_b \square H_n) \leq 4$.

Bernhart and Kainen [4] implicitly proved a result similar to (1) for stack layouts. Let $\text{dsn}(G)$ be the minimum integer k such that G has a k -stack layout $(<, \varphi)$ where φ is a proper edge-colouring of G ; that is, $\varphi(vx) \neq \varphi(vy)$ for any two edges vx, vy with a common endpoint. Then for every graph G_1 and every bipartite graph G_2 ,

$$\text{sn}(G_1 \square G_2) \leq \text{sn}(G_1) + \text{dsn}(G_2). \quad (2)$$

The key difference between (1) and (2) is that G_2 is assumed to be bipartite in (2). Theorem 2 says that this assumption is essential, since it is easily seen that $\text{dsn}(H_n)$ is bounded, but $\text{sn}(S_b \square H_n)$ is unbounded by Theorem 2.

Subdivisions

A noteworthy consequence of Theorem 1 is that it resolves a conjecture of Blankenship and Oporowski [6]. A graph G' is a *subdivision* of a graph G if G' can be obtained from G by replacing the edges vw of G by internally disjoint paths P_{vw} with endpoints v and w . If each P_{vw} has exactly k internal vertices, then G' is the k -subdivision of G . If each P_{vw} has at most k internal vertices, then G' is a $(\leq k)$ -subdivision of G . Blankenship and Oporowski [6] conjectured that the stack-number of $(\leq k)$ -subdivisions (k fixed) is not much less than the stack-number of the original graph. More precisely:

Conjecture 1 ([6]). *There exists a function f such that for every graph G and integer k , if G' is any $(\leq k)$ -subdivision of G , then $\text{sn}(G) \leq f(\text{sn}(G'), k)$.*

Dujmović and Wood [17] established a connection between this conjecture and the question of whether stack-number is bounded by queue-number. In particular, they showed that if Conjecture 1 were true, then stack-number would be bounded by queue-number. Since Theorem 1 shows that stack-number is not bounded by queue-number, Conjecture 1 is false. The proof of Dujmović and Wood [17] is based on the following key lemma: every graph G has a 3-stack subdivision with $1 + 2\lceil \log_2 \text{qn}(G) \rceil$ division vertices per edge. Applying this result to the graph $G = S_b \square H_n$ in Theorem 1, the 5-subdivision of $S_b \square H_n$ has a 3-stack layout. If Conjecture 1 were true, then $\text{sn}(S_b \square H_n) \leq f(3, 5)$, contradicting Theorem 1.

Is Queue-Number Bounded by Stack-Number?

It remains open whether queues are more powerful than stacks; that is, whether queue-number is bounded by stack-number. Several results are known about this problem. Heath et al. [23] showed that every 1-stack graph has a 2-queue layout. Dujmović et al. [11] showed that planar graphs have bounded queue-number. (Note that graph products also feature heavily in this proof.) Since 2-stack graphs are planar, this implies that 2-stack graphs have bounded queue-number. It is open whether 3-stack graphs have bounded queue-number. In fact, the case of three stacks is as hard as the general question. Dujmović and Wood [17] proved that queue-number is bounded by stack-number if and only if 3-stack graphs have bounded queue-number. Moreover, if this is true then queue-number is bounded by a polynomial function of stack-number.

2 Proof of Theorem 2

We now turn to the proof of our main result, the lower bound on $\text{sn}(G)$, where $G := S_b \square H_n$. Consider a hypothetical s -stack layout $(\varphi, <)$ of G where n and b are chosen sufficiently large compared to s as detailed below. We begin with three lemmata that, for sufficiently large b , provide a large subgraph S_d of S_b for which the induced stack layout of $S_d \square H_n$ is highly structured.

For each node v of S_b , define π_v as the permutation of $\{1, \dots, n\}^2$ in which (x_1, y_1) appears before (x_2, y_2) if and only if $(v, (x_1, y_1)) < (v, (x_2, y_2))$. The following lemma is an immediate consequence of the Pigeonhole Principle:

Lemma 1. *There exists a permutation π of $\{1, \dots, n\}^2$ and a set L_1 of leaves of S_b of size $a \geq b/(n^2)!$ such that $\pi_v = \pi$ for each $v \in L_1$.*

For each leaf v in L_1 , let φ_v be the edge colouring of H_n defined by $\varphi_v(xy) := \varphi((v, x)(v, y))$ for each $xy \in E(H_n)$. Since H_n has maximum degree 6 and is not 6-regular, it has fewer than $3n^2$ edges. Therefore there are fewer than s^{3n^2} edge colourings of H_n using s colours. Another application of the Pigeonhole Principle proves the following:

Lemma 2. *There exists a subset $L_2 \subseteq L_1$ of size $c \geq a/s^{3n^2}$ and an edge colouring $\phi : E(H_n) \rightarrow \{1, \dots, s\}$ such that $\varphi_v = \phi$ for each $v \in L_2$.*

Let S_c be the subgraph of S_b induced by $L_2 \cup \{r\}$. The preceding two lemmata ensure that, for distinct leaves v and w of S_c , the stack layouts of the isomorphic graphs $G[\{(v, p) : p \in V(H_n)\}]$ and $G[\{(w, p) : p \in V(H_n)\}]$ are identical. The next lemma is a statement about the relationships between the stack layouts of $G[\{(v, p) : v \in V(S_c)\}]$ and $G[\{(v, q) : v \in V(S_c)\}]$ for distinct $p, q \in V(H_n)$. It does not assert that these two layouts are identical but it does state that they fall into one of two categories.

Lemma 3. *There exists a sequence $u_1, \dots, u_d \in L_2$ of length $d \geq c^{1/2^{n^2-1}}$ such that, for each $p \in V(H_n)$, either $(u_1, p) < (u_2, p) < \dots < (u_d, p)$ or $(u_1, p) > (u_2, p) > \dots > (u_d, p)$.*

Proof. Let p_1, \dots, p_{n^2} denote the vertices of H_n in any order. Begin with the sequence $S_1 := v_{1,1}, \dots, v_{1,c}$ that contains all c elements of L_2 ordered so that $(v_{1,1}, p_1) < \dots < (v_{1,c}, p_1)$. For each $i \in \{2, \dots, n^2\}$, the Erdős-Szekeres Theorem [19] implies that S_{i-1} contains a subsequence $S_i := v_{i,1}, \dots, v_{i,|S_i|}$ of length $|S_i| \geq \sqrt{|S_{i-1}|}$ such that $(v_{i,1}, p_i) < \dots < (v_{i,|S_i|}, p_i)$ or $(v_{i,1}, p_i) > \dots > (v_{i,|S_i|}, p_i)$. It is straightforward to verify by induction on i that $|S_i| \geq c^{1/2^{i-1}}$ resulting in a final sequence $S_{n^2} =: L_3$ of length at least $c^{1/2^{n^2-1}}$. \square

For the rest of the proof we work with the star S_d whose leaves are u_1, \dots, u_d described in Lemma 3. Consider the (improper) colouring of H_n obtained by colouring each vertex $p \in V(H_n)$ *red* if $(u_1, p) < \dots < (u_d, p)$ and colouring p *blue* if $(u_1, p) > \dots > (u_d, p)$. We need the following famous Hex Lemma [20].

Lemma 4 ([20]). *Every vertex 2-colouring of H_n contains a monochromatic path on n vertices.*

Apply Lemma 4 with the above-defined colouring of H_n . We obtain a path $R := p_1, \dots, p_n$ in H_n that, without loss of generality, consists entirely of red vertices; thus $(u_1, p_j) < \dots < (u_d, p_j)$ for each $j \in \{1, \dots, n\}$. Let X be the subgraph $S_d \square R$ of G .

Lemma 5. *X contains a set of at least $\min\{\lfloor d/2^n \rfloor, \lceil n/2 \rceil\}$ pairwise crossing edges with respect to $<$.*

Proof. Extend the total order $<$ to a partial order over subsets of $V(G)$, where for all $V, W \subseteq V(G)$, we have $V < W$ if and only if $v < w$ for each $v \in V$ and each $w \in W$. We abuse notation slightly by using $<$ to compare elements of $V(G)$ and subsets of $V(G)$ so that, for $v \in V(G)$ and $V \subseteq V(G)$, $v < V$ denotes $\{v\} < V$. We will define sets $A_1 \supseteq \dots \supseteq A_n$ of leaves of S_d so that each A_i satisfies the following conditions:

- (C1) A_i contains $d_i \geq d/2^{i-1}$ leaves of S_d .
- (C2) Each leaf $v \in A_i$ defines an i -element vertex set $Z_{i,v} := \{(v, p_j) : j \in \{1, \dots, i\}\}$. For any distinct $v, w \in A_i$, the sets $Z_{i,v}$ and $Z_{i,w}$ are *separated* with respect to $<$; that is, $Z_{i,v} < Z_{i,w}$ or $Z_{i,v} > Z_{i,w}$.

Before defining A_1, \dots, A_n we first show how the existence of the set A_n implies the lemma. To avoid triple-subscripts, let $d' := d_n \geq d/2^{n-1}$. The set A_n defines vertex sets $Z_{n,v_1} < \dots < Z_{n,v_{d'}}$ (see Figure 3). Recall that r is the root of S_b so it is adjacent to each of $v_1, \dots, v_{d'}$ in S_d . Therefore, for each $j \in \{1, \dots, n\}$ and each $i \in \{1, \dots, d'\}$, the edge $(r, p_j)(v_i, p_j)$ is in X . Therefore, (r, p_j) is adjacent to an element of each of $Z_{n,v_1}, \dots, Z_{n,v_{d'}}$.

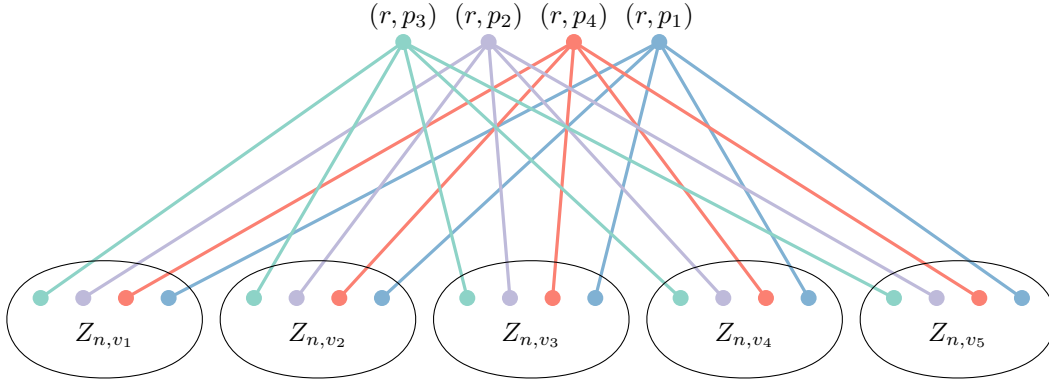


Figure 3: The sets $Z_{n,v_1}, \dots, Z_{n,v_{d'}}$ where $n = 4$ and $d' = 5$.

Since $Z_{n,v_1}, \dots, Z_{n,v_{d'}}$ are separated with respect to $<$, when viewed from afar, this situation looks like a complete bipartite graph $K_{n,d'}$ with the root vertices $L := \{(r, p_j) : j \in \{1, \dots, n\}\}$ in one part and the groups $R := Z_{n,v_1} \cup \dots \cup Z_{n,v_{d'}}$ in the other part. Any linear ordering of $K_{n,d'}$ has a large set of pairwise crossing edges so, intuitively, the induced subgraph $X[L \cup R]$ should also have a large set of pairwise crossing edges. We can formalize this as follows: Label the vertices in L as r_1, \dots, r_n so that $r_1 < \dots < r_n$. Then at least one of the following two cases applies (see Figure 4):

1. $Z_{n,\lfloor d'/2 \rfloor} < r_{\lceil n/2 \rceil}$ in which case the graph between $r_{\lceil n/2 \rceil}, \dots, r_n$ and $Z_{n,1}, \dots, Z_{n,\lfloor d'/2 \rfloor}$ has a set of at least $\min\{\lfloor d'/2 \rfloor, \lceil n/2 \rceil\}$ pairwise-crossing edges.
2. $r_{\lceil n/2 \rceil} < Z_{\lceil d'/2 \rceil+1}$ in which case the graph between $r_1, \dots, r_{\lceil n/2 \rceil}$ and $Z_{\lceil d'/2 \rceil+1}, \dots, Z_{d'}$ has a set of $\min\{\lfloor d'/2 \rfloor, \lceil n/2 \rceil\}$ pairwise-crossing edges.

Since, by (C1), $d' \geq d/2^{n-1}$, either case results in a set of pairwise-crossing edges of size at least $\min\{\lfloor d/2^n \rfloor, \lceil n/2 \rceil\}$, as claimed.

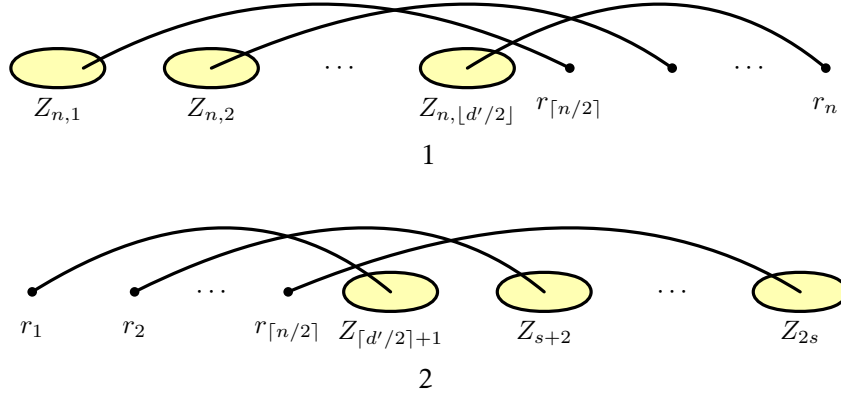


Figure 4: The two cases in the proof of Lemma 5.

All that remains is to define the sets $A_1 \supseteq \dots \supseteq A_n$ that satisfy (C1) and (C2). Let A_1 be the set of all the leaves of S_d . For each $i \in \{2, \dots, n\}$, the set A_i is defined as follows: Let $Z_1, \dots, Z_{|A_{i-1}|}$ denote the sets $Z_{i-1,v}$ for each $v \in A_{i-1}$ ordered so that $Z_1 < \dots < Z_{|A_{i-1}|}$. By Property (C2), this is always possible. Label the vertices of A_{i-1} as $v_1, \dots, v_{|A_{i-1}|}$ so that $(v_1, p_{i-1}) < \dots < (v_r, p_{i-1})$. (This is equivalent to naming them so that $(v_j, p_{i-1}) \in Z_j$ for each $j \in \{1, \dots, |A_{i-1}|\}$.) Define the set $A_i := \{v_{2k+1} : k \in \{0, \dots, \lfloor (|A_{i-1}| - 1)/2 \rfloor\}\} = \{v_j \in A_{i-1} : j \text{ is odd}\}$. This completes the definition of A_1, \dots, A_n .

All that remains is to verify that A_i satisfies (C1) and (C2) for each $i \in \{1, \dots, n\}$. We do this by induction on i . The base case $i = 1$ is trivial so we assume from this point on that $i \in \{2, \dots, n\}$. To see that A_i satisfies (C1) just observe that $|A_i| = \lceil |A_{i-1}|/2 \rceil \geq |A_{i-1}|/2 \geq d/2^{i-1}$, where the final inequality follows by applying the inductive hypothesis $|A_{i-1}| \geq d/2^{i-2}$. Now all that remains is to show that A_i satisfies (C2).

Recall that, for each $v \in A_{i-1}$, the edge $e_v := (v, p_{i-1})(v, p_i)$ is in X . We have the following properties:

- (P1) By Lemma 2, $\phi(e_v) = \phi(p_{i-1}p_i)$ for each $v \in A_{i-1}$.
- (P2) Since p_{i-1} and p_i are both red, $(v, p_{i-1}) < (w, p_{i-1})$ if and only if $(v, p_i) < (w, p_i)$ for each $v, w \in A_{i-1}$.
- (P3) By Lemma 1, $(v, p_{i-1}) < (v, p_i)$ for every $v \in A_{i-1}$ or $(v, p_{i-1}) > (v, p_i)$ for every $v \in A_{i-1}$.

We claim that these three conditions imply that the vertex sets $\{(v, p_{i-1}) : v \in A_{i-1}\}$ and $\{(v, p_i) : v \in A_{i-1}\}$ interleave perfectly with respect to $<$. More precisely:

Claim 1. $(v_1, p_{i-1+t}) < (v_1, p_{i-t}) < (v_2, p_{i-1+t}) < (v_2, p_{i-t}) \dots < (v_r, p_{i-1+t}) < (v_r, p_{i-t})$ for some $t \in \{0, 1\}$.

Proof of Claim 1. By (P3) we may assume, without loss of generality, that $(v, p_{i-1}) < (v, p_i)$ for each $v \in A_{i-1}$, in which case we are trying to prove the claim for $t = 0$. Therefore, it is sufficient to show that $(v_j, p_i) < (v_{j+1}, p_{i-1})$ for each $j \in \{1, \dots, r-1\}$. For the sake of contradiction, suppose $(v_j, p_i) > (v_{j+1}, p_{i-1})$ for some $j \in \{1, \dots, r-1\}$. By the labelling of A_{i-1} , $(v_j, p_{i-1}) < (v_{j+1}, p_{i-1})$ so, by (P2), $(v_j, p_i) < (v_{j+1}, p_i)$. Therefore

$$(v_j, p_{i-1}) < (v_{j+1}, p_{i-1}) < (v_j, p_i) < (v_{j+1}, p_i) .$$

Therefore the edges $e_{v_j} = (v_j, p_{i-1})(v_j, p_i)$ and $e_{v_{j+1}} = (v_{j+1}, p_{i-1})(v_{j+1}, p_i)$ cross with respect to $<$. But this is a contradiction since, by (P1), $\varphi(e_{v_j}) = \varphi(e_{v_{j+1}}) = \phi(p_{i-1}p_i)$. This contradiction completes the proof of Claim 1. \square

We now complete the proof that A_i satisfies (C2). Apply Claim 1 and assume without loss of generality that $t = 0$, so that

$$(v_1, p_{i-1}) < (v_1, p_i) < (v_2, p_{i-1}) < (v_2, p_i) \cdots < (v_r, p_{i-1}) < (v_r, p_i) .$$

For each $j \in \{1, \dots, r-2\}$, we have $(v_{j+1}, p_{i-1}) \in Z_{j+1} < Z_{j+2}$, so $(v_j, p_i) < (v_{j+1}, p_{i-1}) < Z_{j+2}$. Therefore $Z_j \cup \{(v_j, p_i)\} < Z_{j+2}$. By a symmetric argument, $Z_j \cup \{(v_j, p_i)\} > Z_{j-2}$ for each $j \in \{3, \dots, r\}$. Finally, since $(v_j, p_i) < (v_{j+2}, p_i)$ for each odd $i \in \{1, \dots, r\}$, we have $Z_j \cup \{(v_j, p_i)\} < Z_{j+2} \cup \{(v_{j+2}, p_i)\}$ for each odd $j \in \{1, \dots, r-2\}$. Thus A_i satisfies (C2) since the sets $Z_1 \cup \{(v_1, p_i)\}, Z_3 \cup \{(v_3, p_i)\}, \dots, Z_{2\lfloor (r-1)/2 \rfloor + 1} \cup \{(v_{2\lfloor (r-1)/2 \rfloor + 1}, p_i)\}$ are precisely the sets $Z_{i,1}, \dots, Z_{i,d_i}$ determined by our choice of A_i . \square

Proof of Theorem 2. Let $G := S_b \square H_n$, where $n := 2s+1$ and $b := (n^2)!s^{3n^2}((s+1)2^n)^{2^{n^2-1}}$. Suppose that G has an s -stack layout $(\varphi, <)$. In particular, there are no $s+1$ pairwise crossing edges in G with respect to $<$. By Lemmas 1 to 3, we have $a \geq b/(n^2)! = s^{3n^2}((s+1)2^n)^{2^{n^2-1}}$ and $c \geq a/s^{3n^2} \geq ((s+1)2^n)^{2^{n^2-1}}$ and $d \geq c^{1/2^{n^2-1}} \geq (s+1)2^n$. By Lemma 5, the graph X , which is a subgraph of G , contains $\min\{\lfloor d/2^n \rfloor, \lceil n/2 \rceil\} = s+1$ pairwise crossing edges with respect to $<$. This contradiction shows that $\text{sn}(G) > s$. \square

3 Open Problems

Recall that every 1-queue graph has a 2-stack layout [23] and we proved that there are 4-queue graphs with unbounded stack-number. The following questions remain open: Do 2-queue graphs have bounded stack-number? Do 3-queue graphs have bounded stack-number?

Do we want to make a conjecture here? I would say “no” for both questions.

Given the role of cartesian products in our proof, it is natural to ask when is $\text{sn}(G_1 \square G_2)$ bounded? Note that H_n is a subgraph of a planar Hamiltonian graph (namely, H_{2n}), so $\text{sn}(H_n) \leq 2$. So $\text{sn}(G_1 \square G_2)$ can be unbounded even when G_1 is a star and $\text{sn}(G_2) \leq 2$. Since $\text{sn}(G_2) \leq 1$ if and only if G_2 is outerplanar, the following questions naturally arise: Is $\text{sn}(S \square H)$ bounded for every star S and outerplanar graph H with bounded degree? Is $\text{sn}(T \square H)$ bounded for every tree T and outerplanar graph H with bounded degree? The assumption that H has bounded degree is needed since $S_n \square S_n$ contains the 1-subdivision of $K_{n,n}$, which has unbounded stack-number [5].

Since $H_n \subseteq P \boxtimes P$ where P is the n -vertex path, Theorem 1 implies that $\text{sn}(S \boxtimes P \boxtimes P)$ is unbounded for stars S and paths P . It is easily seen that $\text{sn}(S \boxtimes P)$ is bounded [27]. The following question naturally arises (independently asked by Pupyrev [27]): Is $\text{sn}(T \boxtimes P)$ bounded for every tree T and path P ? We conjecture the answer is “no”.

References

- [1] PATRIZIO ANGELINI, GIUSEPPE DI BATTISTA, FABRIZIO FRATI, MAURIZIO PATRIGNANI, AND IGNAZ RUTTER. [Testing the simultaneous embeddability of two graphs whose intersection is a biconnected or a connected graph](#). *J. Discrete Algorithms*, 14:150–172, 2012.
- [2] MICHAEL BAUR AND ULRIK BRANDES. [Crossing reduction in circular layouts](#). In *Proc. 30th International Workshop on Graph-Theoretic Concepts in Computer Science (WG '04)*, vol. 3353 of *Lecture Notes in Computer Science*, pp. 332–343. Springer, 2004.
- [3] MICHAEL A. BEKOS, HENRY FÖRSTER, MARTIN GRONEMANN, TAMARA MCHEDLIDZE, FABRIZIO MONTECCHIANI, CHRYSANTHI N. RAFTOPOULOU, AND TORSTEN UECKERDT. [Planar graphs of bounded degree have bounded queue number](#). *SIAM J. Comput.*, 48(5):1487–1502, 2019.
- [4] FRANK R. BERNHART AND PAUL C. KAINEN. [The book thickness of a graph](#). *J. Combin. Theory Ser. B*, 27(3):320–331, 1979.
- [5] ROBIN BLANKENSHIP. [Book embeddings of graphs](#). Ph.D. thesis, Department of Mathematics, Louisiana State University, U.S.A., 2003.
- [6] ROBIN BLANKENSHIP AND BOGDAN OPOROWSKI. [Drawing subdivisions of complete and complete bipartite graphs on books](#). Tech. Rep. 1999-4, Department of Mathematics, Louisiana State University, U.S.A., 1999.
- [7] JEAN BOURGAIN. [Expanders and dimensional expansion](#). *C. R. Math. Acad. Sci. Paris*, 347(7-8):357–362, 2009.
- [8] JEAN BOURGAIN AND AMIR YEHUDAYOFF. [Expansion in \$SL_2\(\mathbb{R}\)\$ and monotone expansion](#). *Geometric and Functional Analysis*, 23(1):1–41, 2013.
- [9] FAN R. K. CHUNG, F. THOMSON LEIGHTON, AND ARNOLD L. ROSENBERG. [Embedding graphs in books: a layout problem with applications to VLSI design](#). *SIAM J. Algebraic Discrete Methods*, 8(1):33–58, 1987.
- [10] GIUSEPPE DI BATTISTA, FABRIZIO FRATI, AND JÁNOS PACH. [On the queue number of planar graphs](#). *SIAM J. Comput.*, 42(6):2243–2285, 2013.
- [11] VIDA DUJMOVIĆ, GWENAËL JORET, PIOTR MICEK, PAT MORIN, TORSTEN UECKERDT, AND DAVID R. WOOD. [Planar graphs have bounded queue-number](#). *J. ACM*, 67(4):22, 2020.
- [12] VIDA DUJMOVIĆ, PAT MORIN, AND DAVID R. WOOD. [Layout of graphs with bounded tree-width](#). *SIAM J. Comput.*, 34(3):553–579, 2005.
- [13] VIDA DUJMOVIĆ, PAT MORIN, AND DAVID R. WOOD. [Layered separators in minor-closed graph classes with applications](#). *J. Combin. Theory Ser. B*, 127:111–147, 2017.
- [14] VIDA DUJMOVIĆ, ATTILA PÓR, AND DAVID R. WOOD. [Track layouts of graphs](#). *Discrete Math. Theor. Comput. Sci.*, 6(2):497–522, 2004.
- [15] VIDA DUJMOVIĆ, ANASTASIOS SIDIROPOULOS, AND DAVID R. WOOD. [Layouts of expander graphs](#). *Chicago J. Theoret. Comput. Sci.*, 2016(1), 2016.
- [16] VIDA DUJMOVIĆ AND DAVID R. WOOD. [On linear layouts of graphs](#). *Discrete Math. Theor. Comput. Sci.*, 6(2):339–358, 2004.
- [17] VIDA DUJMOVIĆ AND DAVID R. WOOD. [Stacks, queues and tracks: Layouts of graph subdivisions](#). *Discrete Math. Theor. Comput. Sci.*, 7:155–202, 2005.
- [18] VIDA DUJMOVIĆ AND DAVID R. WOOD. [Graph treewidth and geometric thickness parameters](#). *Discrete Comput. Geom.*, 37(4):641–670, 2007.

-
- [19] PAUL ERDŐS AND GEORGE SZEKERES. [A combinatorial problem in geometry](#). *Compositio Math.*, 2:463–470, 1935.
- [20] DAVID GALE. [The game of Hex and the Brouwer fixed-point theorem](#). *Amer. Math. Monthly*, 86(10):818–827, 1979.
- [21] ZVI GALIL, RAVI KANNAN, AND ENDRE SZEMERÉDI. [On 3-pushdown graphs with large separators](#). *Combinatorica*, 9(1):9–19, 1989.
- [22] CHRISTIAN HASLINGER AND PETER F. STADLER. [RNA structures with pseudo-knots: Graph-theoretical, combinatorial, and statistical properties](#). *Bull. Math. Biology*, 61(3):437–467, 1999.
- [23] LENWOOD S. HEATH, F. THOMSON LEIGHTON, AND ARNOLD L. ROSENBERG. [Comparing queues and stacks as mechanisms for laying out graphs](#). *SIAM J. Discrete Math.*, 5(3):398–412, 1992.
- [24] LENWOOD S. HEATH AND ARNOLD L. ROSENBERG. [Laying out graphs using queues](#). *SIAM J. Comput.*, 21(5):927–958, 1992.
- [25] MICHAEL KAUFMANN, MICHAEL A. BEKOS, FABIAN KLUTE, SERGEY PUPYREV, CHRYSANTHI N. RAFTOPOULOU, AND TORSTEN UECKERDT. [Four pages are indeed necessary for planar graphs](#). *J. Comput. Geom.*, 11(1):332–353, 2020.
- [26] L. TAYLOR OLLMANN. On the book thicknesses of various graphs. In FREDERICK HOFFMAN, ROY B. LEVOW, AND ROBERT S. D. THOMAS, eds., *Proc. 4th Southeastern Conference on Combinatorics, Graph Theory and Computing*, vol. VIII of *Congr. Numer.*, p. 459. Utilitas Math., 1973.
- [27] SERGEY PUPYREV. [Book embeddings of graph products](#). 2020, arXiv:2007.15102.
- [28] ARNOLD L. ROSENBERG. The DIOGENES approach to testable fault-tolerant arrays of processors. *IEEE Trans. Comput.*, C-32:902–910, 1983.
- [29] ARNOLD L. ROSENBERG. Book embeddings and wafer-scale integration. In *Proc. 17th Southeastern International Conf. on Combinatorics, Graph Theory, and Computing*, vol. 54 of *Congr. Numer.*, pp. 217–224. 1986.
- [30] ARNOLD L. ROSENBERG. [DIOGENES, circa 1986](#). In *Proc. VLSI Algorithms and Architectures*, vol. 227 of *Lecture Notes in Comput. Sci.*, pp. 96–107. Springer, 1986.
- [31] FARHAD SHAHROKHI, ONDREJ SÝKORA, LÁSZLÓ A. SZÉKELY, AND IMRICH VŘTO. [Book embeddings and crossing numbers](#). In ERNST W. MAYR, GUNTHER SCHMIDT, AND GOTTFRIED TINHOFFER, eds., *Proc. 20th International Workshop on Graph-Theoretic Concepts in Computer Science (WG '94)*, *Lecture Notes in Comput. Sci.*, pp. 256–268. Springer, 1994.
- [32] DAVID R. WOOD. [Bounded degree book embeddings and three-dimensional orthogonal graph drawing](#). In PETRA MUTZEL, MICHAEL JÜNGER, AND SEBASTIAN LEIPERT, eds., *Proc. 9th International Symposium on Graph Drawing (GD '01)*, vol. 2265 of *Lecture Notes in Computer Science*, pp. 312–327. Springer, 2001.
- [33] DAVID R. WOOD. [Queue layouts of graph products and powers](#). *Discrete Math. Theor. Comput. Sci.*, 7(1):255–268, 2005.
- [34] MIHALIS YANNAKAKIS. [Embedding planar graphs in four pages](#). *J. Comput. System Sci.*, 38(1):36–67, 1989.
- [35] MIHALIS YANNAKAKIS. [Planar graphs that need four pages](#). *J. Combin. Theory Ser. B*, 145:241–263, 2020.
-