# Lecture 02—Amdahl's Law, Modern Hardware ECE 459: Programming for Performance

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#### **Limitations**

Our main focus is parallelization.

- Most programs have a sequential part and a parallel part; and,
- Amdahl's Law answers, "what are the limits to parallelization?"

## Formulation (1)

S: fraction of serial runtime in a serial execution.

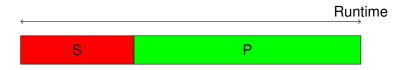
P: fraction of parallel runtime in a serial execution.

Therefore, S + P = 1.

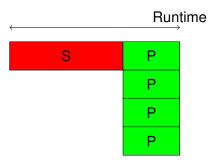
With 4 processors, best case, what can happen to the following runtime?



## Formulation (1)



We want to split up the parallel part over 4 processors



## Formulation (2)

 $T_s$ : time for the program to run in serial N: number of processors/parallel executions  $T_p$ : time for the program to run in parallel

Under perfect conditions, get N speedup for P

$$T_p = T_s \cdot (S + \frac{P}{N})$$

## Formulation (3)

How much faster can we make the program?

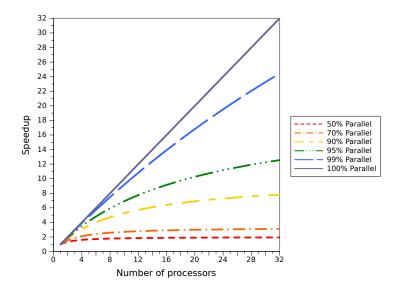
$$speedup = \frac{T_s}{T_p}$$

$$= \frac{T_s}{T_S \cdot (S + \frac{P}{N})}$$

$$= \frac{1}{S + \frac{P}{N}}$$

(assuming no overhead for parallelizing; or costs near zero)

## **Scaling with Fraction of Parallel Code**



## Amdahl's Law

Replace *S* with (1 - P):

$$\textit{speedup} = \tfrac{1}{(1-P) + \frac{P}{N}}$$

maximum speedup = 
$$\frac{1}{(1-P)}$$
, since  $\frac{P}{N} \to 0$ 

As you might imagine, the asymptotes in the previous graph are bounded by the maximum speedup.

#### Amdahl's Law Generalization

The program may have many parts, each of which we can tune to a different degree.

Let's generalize Amdahl's Law.

$$f_1, f_2, \dots, f_n$$
: fraction of time in part  $n$   $S_{f_1}, S_{f_n}, \dots, S_{f_n}$ : speedup for part  $n$ 

$$\textit{speedup} = \frac{1}{\frac{f_1}{\mathcal{S}_{f_1}} + \frac{f_2}{\mathcal{S}_{f_2}} + \ldots + \frac{f_n}{\mathcal{S}_{f_n}}}$$

# **Application (1)**

Consider a program with 4 parts in the following scenario:

		Speedup	
Part	Fraction of Runtime	Option 1	Option 2
1	0.55	1	2
2	0.25	5	1
3	0.15	3	1
4	0.05	10	1

We can implement either Option 1 or Option 2. Which option is better?

# **Application (2)**

"Plug and chug" the numbers:

## Option 1

$$speedup = \frac{1}{0.55 + \frac{0.25}{5} + \frac{0.15}{3} + \frac{0.05}{5}} = 1.53$$

## Option 2

$$speedup = \frac{1}{\frac{0.55}{2} + 0.45} = 1.38$$

## **Empirically estimating parallel speedup P**

Useful to know, don't have to commit to memory:

$$P_{\text{estimated}} = \frac{\frac{1}{speedup} - 1}{\frac{1}{N} - 1}$$

- Quick way to guess the fraction of parallel code
- Use P<sub>estimated</sub> to predict speedup for a different number of processors

## **Another Example**

We run a program in serial and find it spends 12.5% of its execution on serial code and 87.5% on parallel code. How many processors do we need to get within 10% of the perfect parallel runtime?

# Summary of Amdahl's Law

Important to focus on the part of the program with most impact.

#### Amdahl's Law:

- estimates perfect performance gains from parallelization; but,
- only applies to solving a fixed problem size in the shortest possible period of time

## **Gustafson's Law: Formulation**

*n*: problem size

S(n): fraction of serial runtime for a parallel execution

P(n): fraction of parallel runtime for a parallel execution

$$T_p = S(n) + P(n) = 1$$
  
 $T_s = S(n) + N \cdot P(n)$ 

$$speedup = \frac{T_s}{T_p}$$

#### **Gustafson's Law**

$$speedup = S(n) + N \cdot P(n)$$

Assuming the fraction of runtime in serial part decreases as n increases, the speedup approaches N.

 Yes! Large problems can be efficiently parallelized. (Ask Google.)

## **Driving Metaphor**

#### Amdahl's Law

Suppose you're travelling between 2 cities 90 km apart. If you travel for an hour at a constant speed less than 90 km/h, your average will never equal 90 km/h, even if you energize at your destination.

#### Gustafson's Law

Suppose you've been travelling at a constant speed less than 90 km/h. Given enough distance, you can bring your average up to 90 km/h.