Lecture 04—Pthreads and Simple Locks (v2)

ECE 459: Programming for Performance

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Background

Recall the difference between processes and threads:

- Threads are basically light-weight processes which piggy-back on processes' address space.

Traditionally (pre-Linux 2.6) you had to use `fork` (for processes) and `clone` (for threads)

History

`clone` is not POSIX compliant.
Developers mostly used `fork` in the past, which creates a new process.

- Drawbacks?
- Benefits?

Benefit: `fork` is Safer and More Secure Than Threads

1. Each process has its own virtual address space:
   - Memory pages are not copied, they are copy-on-write—
   - Therefore, uses less memory than you would expect.
2. Buffer overruns or other security holes do not expose other processes.
3. If a process crashes, the others can continue.

Example: In the Chrome browser, each tab is a separate process.

Drawback of Processes: Threads are Easier and Faster

- Interprocess communication (IPC) is more complicated and slower than interthread communication.
  - Need to use operating system utilities (pipes, semaphores, shared memory, etc) instead of thread library (or just memory reads and writes).
- Processes have much higher startup, shutdown and synchronization cost.
- And, Pthreads fix the issues of `clone` and provide a uniform interface for most systems **(focus of Assignment 1)**.

Appropriate Time to Use Processes

If your application is like this:

- Mostly independent tasks, with little or no communication.
- Task startup and shutdown costs are negligible compared to overall runtime.
- Want to be safer against bugs and security holes.

Then processes are the way to go.

For performance reasons, along with ease and consistency, we'll use Pthreads.

fork Usage Example (OS refresher)

```
pid = fork();
if (pid < 0) {
        fork_error_function();
} else if (pid == 0) {
        child_function();
} else {
        parent_function();
}
```

fork produces a second copy of the calling process, which starts execution after the call.

The only difference between the copies is the return value: the parent gets the pid of the child, while the child gets 0.

Threads Offer a Speedup of 6.5 over `fork`

Here's a benchmark between `fork` and Pthreads on a laptop, creating and destroying 50,000 threads:

```
jon@riker examples master % time ./create_fork
0.18s user 4.14s system 34% cpu 12.484 total
jon@riker examples master % time ./create_pthread
0.73s user 1.29s system 107% cpu 1.887 total
```

Clearly Pthreads incur much lower overhead than `fork`.

Assumptions

First, we'll see how to use threads on "embarrassingly parallel problems".

- mostly-independent sub-problems (little synchronization); and
- strong locality (little communication).

Later, we'll see:

- which problems can be parallelized (dependencies)
- alternative parallelization patterns
  (right now, just use one thread per sub-problem)

POSIX Threads

- Available on most systems
- Windows has Pthreads Win32, but I wouldn't use it; use Linux for this course
- API available by `#include <pthread.h>`
- Compile with pthread flag (`gcc -pthread prog.c -o prog`)

Creating Threads

```
int pthread_create(pthread_t* thread,
                   const pthread_attr_t* attr,
                   void* (*start_routine)(void*),
                   void* arg);
```

thread: creates a handle to a thread at pointer location
attr: thread attributes (NULL for defaults, more details later)
start_routine: function to start execution
arg: value to pass to start_routine

returns 0 on success, error number otherwise
(contents of *thread are undefined)

Creating Threads—Example

```
#include <pthread.h>
#include <stdio.h>

void* run(void*) {
  printf("In run\n");
}

int main() {
  pthread_t thread;
  pthread_create(&thread, NULL, &run, NULL);
  printf("In main\n");
}
```

Simply creates a thread and terminates
(usage isn't really right, as we'll see.)

Waiting for Threads

```
int pthread_join(pthread_t thread,
                 void** retval)
```

thread: wait for this thread to terminate (thread must be joinable).
retval: stores exit status of thread (set by `pthread_exit`) to the location pointed by *retval. If cancelled, returns `PTHREAD_CANCELED`. `NULL` is ignored.

returns 0 on success, error number otherwise.

Only call this one time per thread! Multiple calls on the same thread leads to undefined behaviour.

Waiting for Threads—Example

```
#include <pthread.h>
#include <stdio.h>

void* run(void*) {
  printf("In run\n");
}

int main() {
  pthread_t thread;
  pthread_create(&thread, NULL, &run, NULL);
  printf("In main\n");
  pthread_join(thread, NULL);
}
```

This now waits for the newly created thread to terminate.

Passing Data to Threads. . . Wrongly

Consider this snippet:

```
int i;
for (i = 0; i < 10; ++i)
  pthread_create(&thread[i], NULL, &run, (void*)&i);
```

This is a terrible idea. Why?

Passing Data to Threads...Wrongly

Consider this snippet:

```
int i;
for (i = 0; i < 10; ++i)
  pthread_create(&thread[i], NULL, &run, (void*)&i);
```

This is a terrible idea. Why?

1. The value of `i` will probably change before the thread executes
2. The memory for `i` may be out of scope, and therefore invalid by the time the thread executes

Passing Data to Threads

What about:

```
int i;
for (i = 0; i < 10; ++i)
  pthread_create(&thread[i], NULL, &run, (void*)i);

...

void* run(void* arg) {
  int id = (int)arg;
```

This is suggested in the book, but should carry a warning:

Passing Data to Threads

What about:

```
int i;
for (i = 0; i < 10; ++i)
  pthread_create(&thread[i], NULL, &run, (void*)i);

...

void* run(void* arg) {
  int id = (int)arg;
```

This is suggested in the book, but should carry a warning:

- Beware size mismatches between arguments: no guarantee that a pointer is the same size as an int, so your data may overflow.
- Sizes of data types change between systems. For maximum portability, just use pointers you got from `malloc`.

Detached Threads

Joinable threads (the default) wait for someone to call `pthread_join` before they release their resources.

Detached threads release their resources when they terminate, without being joined.

```
int pthread_detach(pthread_t thread);
```

thread: marks the thread as detached

returns 0 on success, error number otherwise.

Calling `pthread_detach` on an already detached thread results in undefined behaviour.

Thread Termination

```
void pthread_exit(void *retval);
```

retval: return value passed to function that calls `pthread_join`

start_routine returning is equivalent to calling `pthread_exit` with that return value;

`pthread_exit` is called implicitly when the `start_routine` of a thread returns.

Attributes

By default, threads are *joinable* on Linux, but a more portable way to know what you're getting is to set thread attributes. You can change:

- Detached or joinable state
- Scheduling inheritance
- Scheduling policy
- Scheduling parameters
- Scheduling contention scope
- Stack size
- Stack address
- Stack guard (overflow) size

Attributes—Example

```
size_t stacksize;
pthread_attr_t attributes;
pthread_attr_init(&attributes);
pthread_attr_getstacksize(&attributes, &stacksize);
printf("Stack size = %i\n", stacksize);
pthread_attr_destroy(&attributes);
```

Running this on a laptop produces:

```
jon@riker examples master % ./stack_size
Stack size = 8388608
```

Setting a thread state to joinable:

```
pthread_attr_setdetachstate(&attributes,
                            PTHREAD_CREATE_JOINABLE);
```

Detached Threads: Warning!

```
#include <pthread.h>
#include <stdio.h>

void* run(void*) {
  printf("In run\n");
}

int main() {
  pthread_t thread;
  pthread_create(&thread, NULL, &run, NULL);
  pthread_detach(thread);
  printf("In main\n");
}
```

When I run it, it just prints “In main”, why?

Detached Threads: Solution to Problem

```
#include <pthread.h>
#include <stdio.h>

void* run(void*) {
  printf("In run\n");
}

int main() {
  pthread_t thread;
  pthread_create(&thread, NULL, &run, NULL);
  pthread_detach(thread);
  printf("In main\n");
  pthread_exit(NULL); // This waits for all detached
                      // threads to terminate
}
```

Make the final call `pthread_exit` if you have any detached threads.

Three Useful Routines

```
pthread_t pthread_self(void);

int pthread_equal(pthread_t t1, pthread_t t2);

int pthread_once(pthread_once_t* once_control,
                 void (*init_routine)(void));
pthread_once_t once_control = PTHREAD_ONCE_INIT;
```

`pthread_self` returns the handle of the currently running thread.

Use `pthread_equal` if you're comparing 2 threads.

If you want to run a section of code once, you need `pthread_once` (it's well-named). It will run only once per `once_control`.

Threading Challenges

- Be aware of scheduling (you can also set affinity with pthreads on Linux).
- Make sure the libraries you use are **thread-safe**:
  - Means that the library protects its shared data.
- glibc reentrant functions are also safe: a program can have more than one thread calling these functions concurrently.
- **Example:** In Assignment 1, we'll use `rand_r`, not `rand`.

Mutual Exclusion

Mutexes are the most basic type of synchronization.

- Only one thread can access code protected by a mutex at a time.
- All other threads must wait until the mutex is free before they can execute the protected code.

Creating Mutexes—Example

```
pthread_mutex_t m1 = PTHREAD_MUTEX_INITIALIZER;
pthread_mutex_t m2;

pthread_mutex_init(&m2, NULL);
...
pthread_mutex_destroy(&m1);
pthread_mutex_destroy(&m2);
```

- Two ways to initialize mutexes: statically and dynamically
- If you want to include attributes, you need to use the dynamic version

Mutex Attributes

- **Protocol**: specifies the protocol used to prevent priority inversions for a mutex
- **Prioceiling**: specifies the priority ceiling of a mutex
- **Process-shared**: specifies the process sharing of a mutex

You can specify a mutex as *process shared* so that you can access it between processes. In that case, you need to use shared memory and `mmap`, which we won't get into.

Using Mutexes: Example

```
// code
pthread_mutex_lock(&m1);
// protected code
pthread_mutex_unlock(&m1);
// more code
```

- Everything within the `lock` and `unlock` is protected.
- Be careful to avoid deadlocks if you are using multiple mutexes.
- Also you can use `pthread_mutex_trylock`, if needed.

Data Race Example

Recall that dataraces occur when two concurrent actions access the same variable and at least one of them is a **write**

```
...
static int counter = 0;

void* run(void* arg) {
    for (int i = 0; i < 100; ++i) {
        ++counter;
    }
}

int main(int argc, char *argv[])
{
    // Create 8 threads
    // Join 8 threads
    printf("counter = %i\n", counter);
}
```

Is there a datarace in this example? If so, how would we fix it?

Example Problem Solution

```
...
static pthread_mutex_t mutex = PTHREAD_MUTEX_INITIALIZER;
static int counter = 0;

void* run(void* arg) {
    for (int i = 0; i < 100; ++i) {
        pthread_mutex_lock(&mutex);
        ++counter;
        pthread_mutex_unlock(&mutex);
    }
}

int main(int argc, char *argv[])
{
    // Create 8 threads
    // Join 8 threads
    pthread_mutex_destroy(&mutex);
    printf("counter = %i\n", counter);
}
```