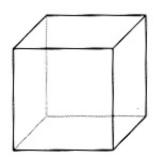
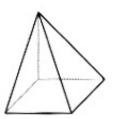
Triangle meshes

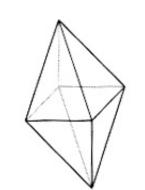
CS 4620 Lecture II

Notation

- $n_T = \# \text{tris}; n_V = \# \text{verts}; n_F = \# \text{edges}$
- Euler: $n_V n_F + n_T = 2$ for a simple closed surface
 - and in general sums to small integer
 - argument for implication that $n_T:n_F:n_V$ is about 2:3:1





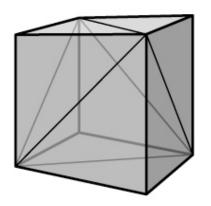


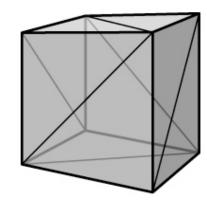
Validity of triangle meshes

- in many cases we care about the mesh being able to bound a region of space nicely
- in other cases we want triangle meshes to fulfill assumptions of algorithms that will operate on them (and may fail on malformed input)
- two completely separate issues:
 - topology: how the triangles are connected (ignoring the positions entirely)
 - geometry: where the triangles are in 3D space

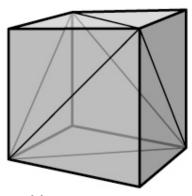
Topology/geometry examples

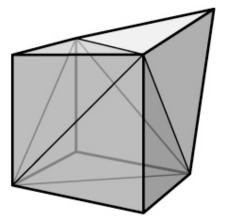
same geometry, different mesh topology:





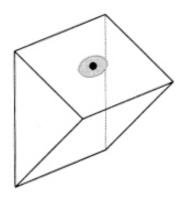
same mesh topology, different geometry:

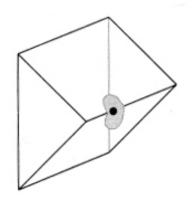


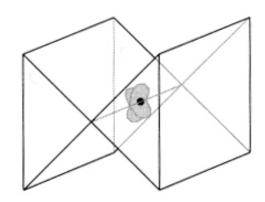


Topological validity

- strongest property, and most simple: be a manifold
 - this means that no points should be "special"
 - interior points are fine
 - edge points: each edge should have exactly 2 triangles
 - vertex points: each vertex should have one loop of triangles
 - not too hard to weaken this to allow boundaries



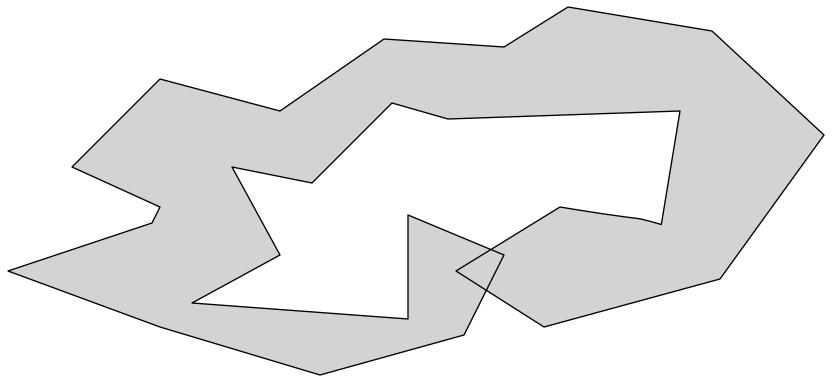




[Foley et al.]

Geometric validity

- generally want non-self-intersecting surface
- hard to guarantee in general
 - because far-apart parts of mesh might intersect



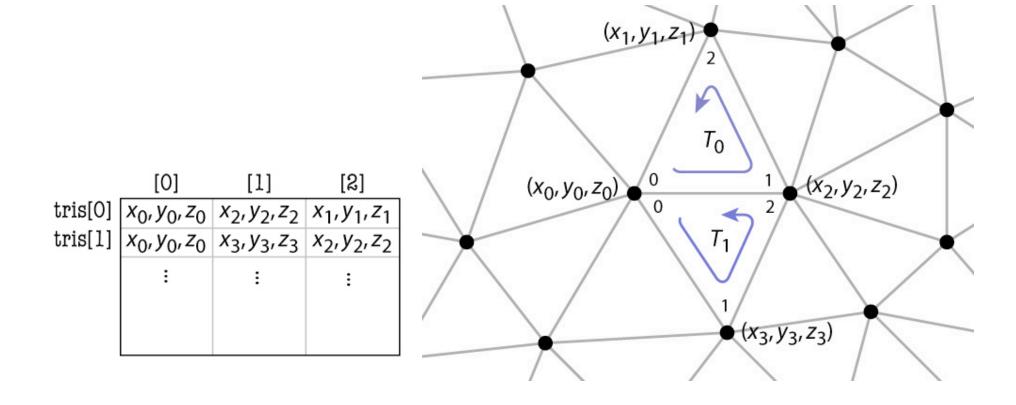
Representation of triangle meshes

- Compactness
- Efficiency for rendering
 - enumerate all triangles as triples of 3D points
- Efficiency of queries
 - all vertices of a triangle
 - all triangles around a vertex
 - neighboring triangles of a triangle
 - (need depends on application)
 - finding triangle strips
 - computing subdivision surfaces
 - mesh editing

Representations for triangle meshes

- Separate triangles
- Indexed triangle set
 - shared vertices
- Triangle strips and triangle fans
 - compression schemes for transmission to hardware
- Triangle-neighbor data structure
 - supports adjacency queries
- Winged-edge data structure
 - supports general polygon meshes

Separate triangles



Separate triangles

- array of triples of points
 - float $[n_T][3][3]$: about 72 bytes per vertex
 - 2 triangles per vertex (on average)
 - 3 vertices per triangle
 - 3 coordinates per vertex
 - 4 bytes per coordinate (float)
- various problems
 - wastes space (each vertex stored 6 times)
 - cracks due to roundoff
 - difficulty of finding neighbors at all

- Store each vertex once
- Each triangle points to its three vertices

```
Triangle {
    Vertex vertex[3];
}

Vertex {
    float position[3]; // or other data
    }

// ... or ...

Mesh {
    float verts[nv][3]; // vertex positions (or other data)
    int tInd[nt][3]; // vertex indices
    }
```

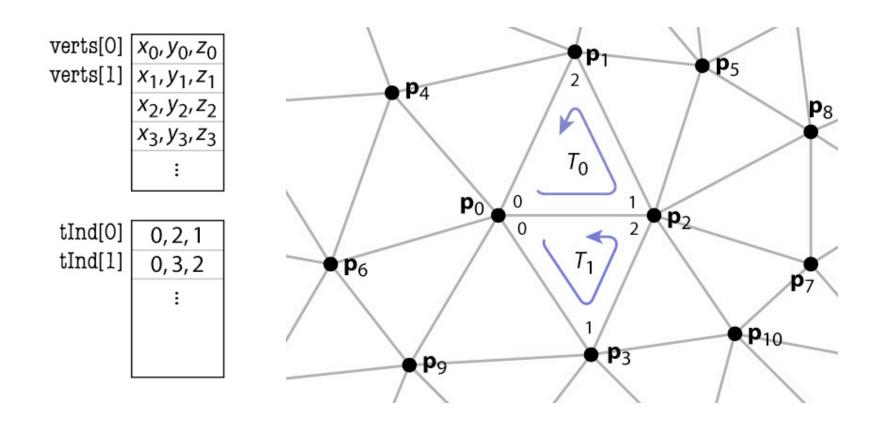
- Store each vertex once
- Each triangle points to its three vertices

```
Triangle {
    Vertex vertex[3];
}

Vertex {
    float position[3]; // or other data
    }

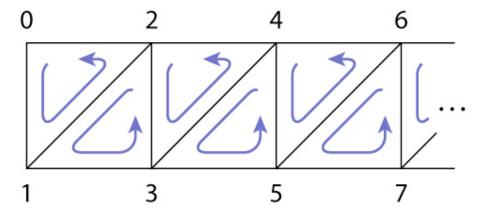
// ... or ...

Mesh {
    float verts[nv][3]; // vertex positions (or other data)
    int tInd[nt][3]; // vertex indices
    }
```

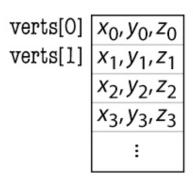


- array of vertex positions
 - float[n_V][3]: I2 bytes per vertex
 - (3 coordinates x 4 bytes) per vertex
- array of triples of indices (per triangle)
 - $int[n_T][3]$: about 24 bytes per vertex
 - 2 triangles per vertex (on average)
 - (3 indices x 4 bytes) per triangle
- total storage: 36 bytes per vertex (factor of 2 savings)
- represents topology and geometry separately
- finding neighbors is at least well defined

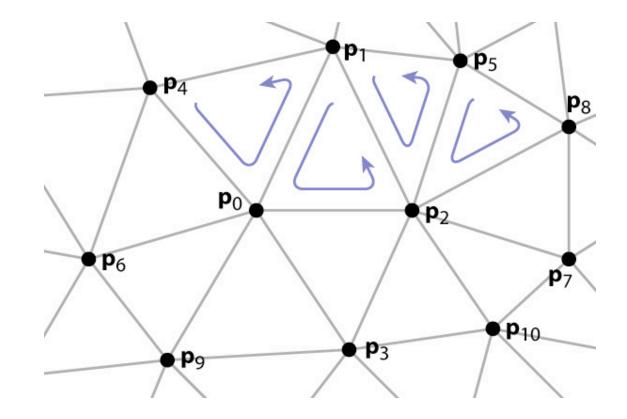
- Take advantage of the mesh property
 - each triangle is usually adjacent to the previous

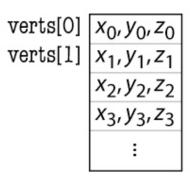


- let every vertex create a triangle by reusing the second and third vertices of the previous triangle
- every sequence of three vertices produces a triangle (but not in the same order)
- e. g., 0, 1, 2, 3, 4, 5, 6, 7, ... leads to
 (0 1 2), (2 1 3), (2 3 4), (4 3 5), (4 5 6), (6 5 7), ...
- for long strips, this requires about one index per triangle

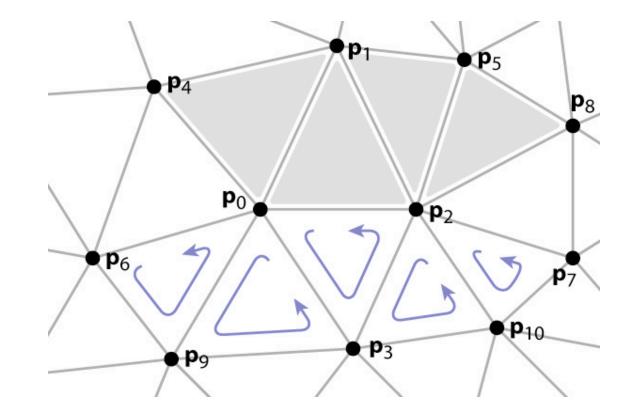


tStrip[0] 4, 0, 1, 2, 5, 8 tStrip[1] 6, 9, 0, 3, 2, 10, 7





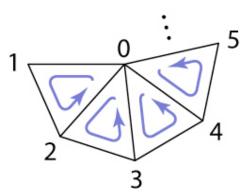
tStrip[0] 4, 0, 1, 2, 5, 8 tStrip[1] 6, 9, 0, 3, 2, 10, 7



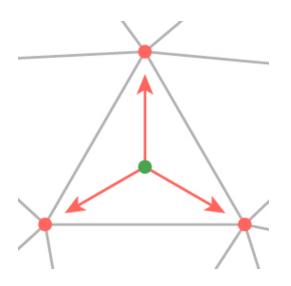
- array of vertex positions
 - float[n_V][3]: I2 bytes per vertex
 - (3 coordinates x 4 bytes) per vertex
- array of index lists
 - $int[n_S][variable]: 2 + n indices per strip$
 - on average, $(I + \varepsilon)$ indices per triangle (assuming long strips)
 - 2 triangles per vertex (on average)
 - about 4 bytes per triangle (on average)
- total is 20 bytes per vertex (limiting best case)
 - factor of 3.6 over separate triangles; I.8 over indexed mesh

Triangle fans

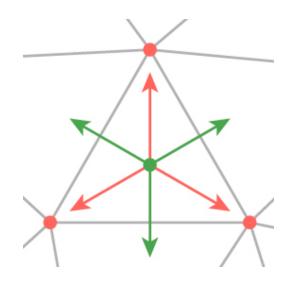
- Same idea as triangle strips, but keep oldest rather than newest
 - every sequence of three vertices produces a triangle
 - e. g., 0, 1, 2, 3, 4, 5, ... leads to(0 1 2), (0 2 3), (0 3 4), (0 3 5),
 - for long fans, this requires
 about one index per triangle
- Memory considerations exactly the same as triangle strip



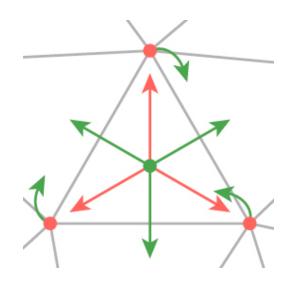
- Extension to indexed triangle set
- Triangle points to its three neighboring triangles
- Vertex points to a single neighboring triangle
- Can now enumerate triangles around a vertex



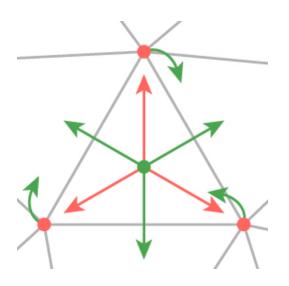
- Extension to indexed triangle set
- Triangle points to its three neighboring triangles
- Vertex points to a single neighboring triangle
- Can now enumerate triangles around a vertex

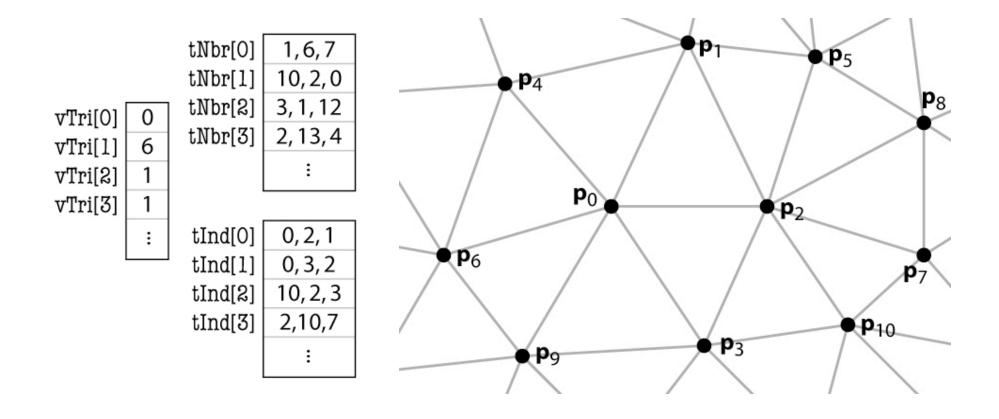


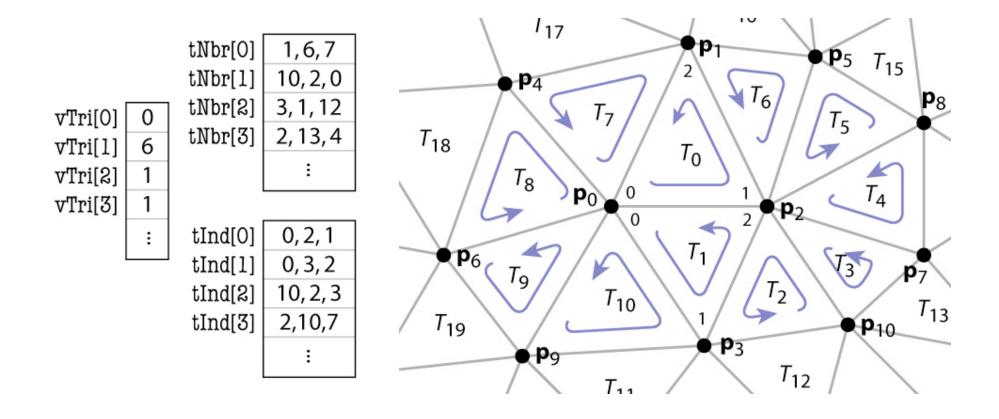
- Extension to indexed triangle set
- Triangle points to its three neighboring triangles
- Vertex points to a single neighboring triangle
- Can now enumerate triangles around a vertex

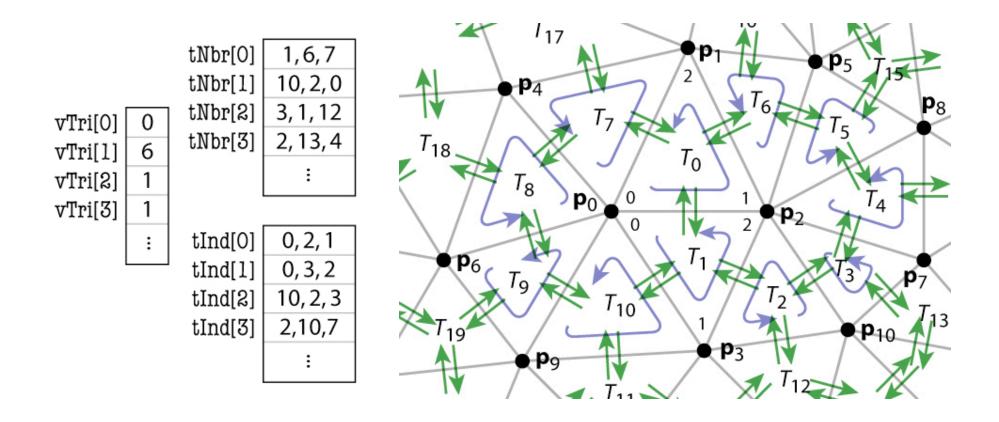


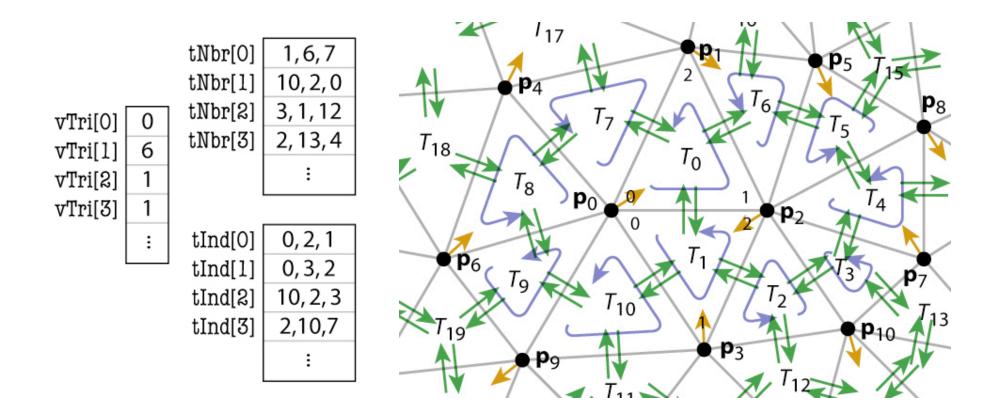
```
Triangle {
   Triangle nbr[3];
   Vertex vertex[3];
// t.neighbor[i] is adjacent
// across the edge from i to i+1
Vertex {
  // ... per-vertex data ...
   Triangle t; // any adjacent tri
// ... or ...
Mesh {
  // ... per-vertex data ...
  int tInd[nt][3]; // vertex indices
   int tNbr[nt][3]; // indices of neighbor triangles
   int vTri[nv]; // index of any adjacent triangle
```

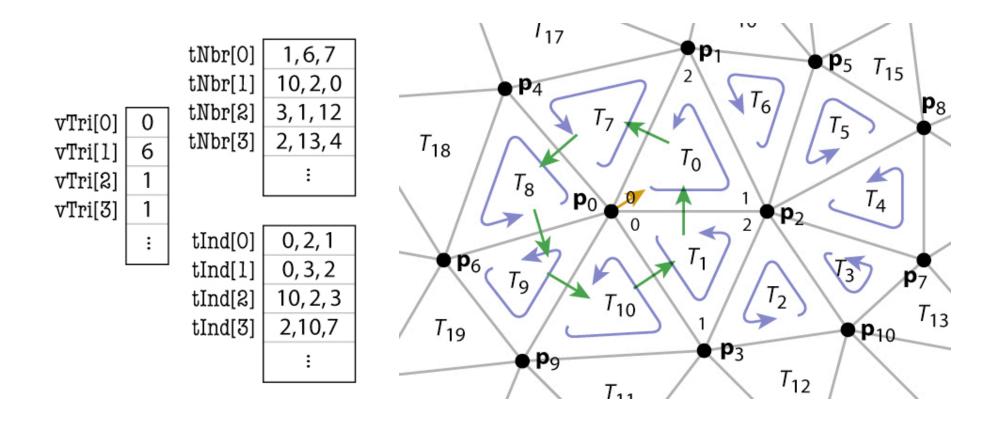






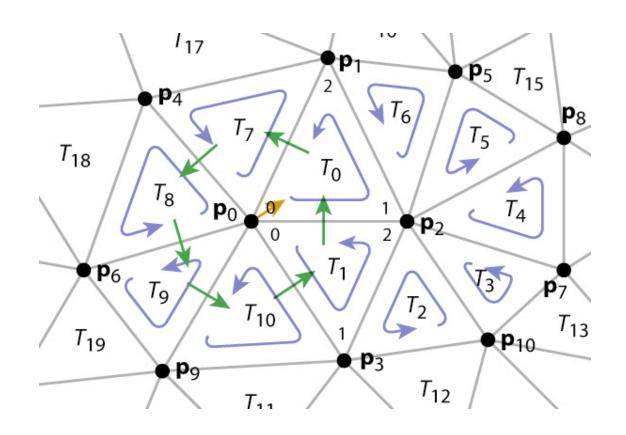






```
TrianglesOfVertex(v) {
    t = v.t;
    do {
        find t.vertex[i] == v;
        t = t.nbr[pred(i)];
        } while (t != v.t);
    }

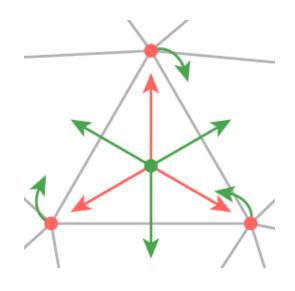
pred(i) = (i+2) % 3;
succ(i) = (i+1) % 3;
```



- indexed mesh was 36 bytes per vertex
- add an array of triples of indices (per triangle)
 - $int[n_T][3]$: about 24 bytes per vertex
 - 2 triangles per vertex (on average)
 - (3 indices x 4 bytes) per triangle
- add an array of representative triangle per vertex
 - int $[n_V]$: 4 bytes per vertex
- total storage: 64 bytes per vertex
 - still not as much as separate triangles

Triangle neighbor structure—refined

```
Triangle {
   Edge nbr[3];
   Vertex vertex[3];
// if t.nbr[i].i == j
// then t.nbr[i].t.nbr[j] == t
Edge {
  // the i-th edge of triangle t
   Triangle t;
   int i; // in \{0,1,2\}
  // in practice t and i share 32 bits
Vertex {
  // ... per-vertex data ...
   Edge e; // any edge leaving vertex
```

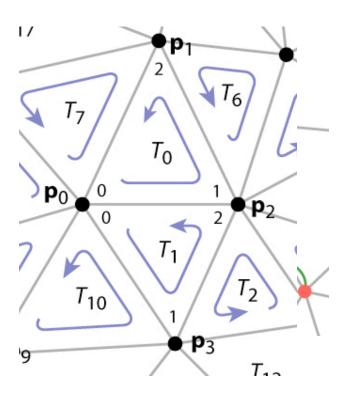


$$T_0.nbr[0] = \{ T_1, 2 \}$$

 $T_1.nbr[2] = \{ T_0, 0 \}$
 $V_0.e = \{ T_1, 0 \}$

Triangle neighbor structure—refined

```
Triangle {
   Edge nbr[3];
   Vertex vertex[3];
// if t.nbr[i].i == j
// then t.nbr[i].t.nbr[j] == t
Edge {
  // the i-th edge of triangle t
   Triangle t;
   int i; // in \{0,1,2\}
  // in practice t and i share 32 bits
Vertex {
  // ... per-vertex data ...
   Edge e; // any edge leaving vertex
```

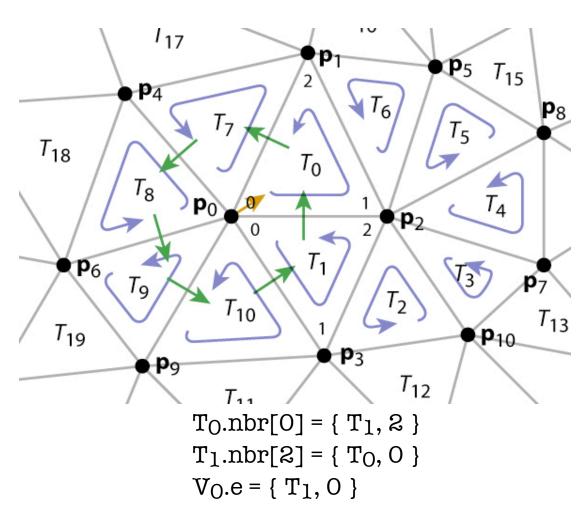


$$T_0.nbr[0] = \{ T_1, 2 \}$$

 $T_1.nbr[2] = \{ T_0, 0 \}$
 $V_0.e = \{ T_1, 0 \}$

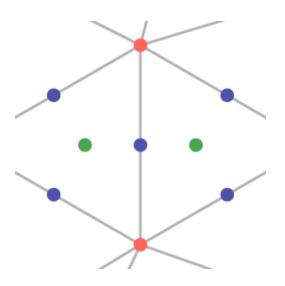
```
TrianglesOfVertex(v) {
    {t, i} = v.e;
    do {
        {t, i} = t.nbr[pred(i)];
      } while (t!= v.t);
    }

pred(i) = (i+2) % 3;
succ(i) = (i+1) % 3;
```



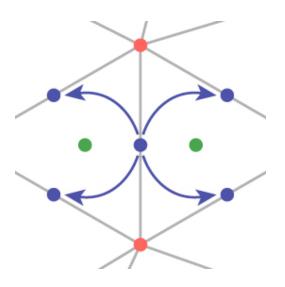
Winged-edge mesh

- Edge-centric rather than face-centric
 - therefore also works for polygon meshes
- Each (oriented) edge points to:
 - left and right forward edges
 - left and right backward edges
 - front and back vertices
 - left and right faces
- Each face or vertex points to one edge



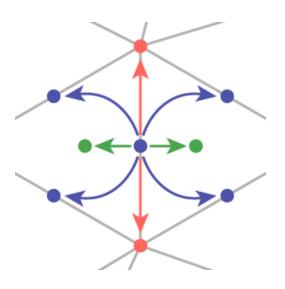
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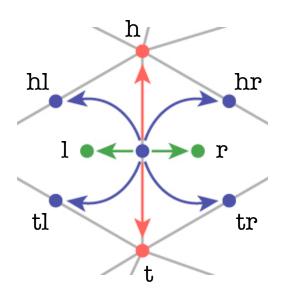
Winged-edge mesh

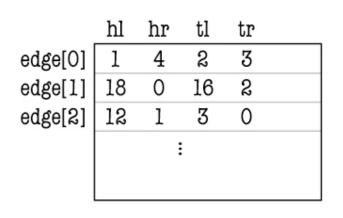
- Edge-centric rather than face-centric
 - therefore also works for polygon meshes
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 - left and right forward edges
 - left and right backward edges
 - front and back vertices
 - left and right faces
- Each face or vertex points to one edge

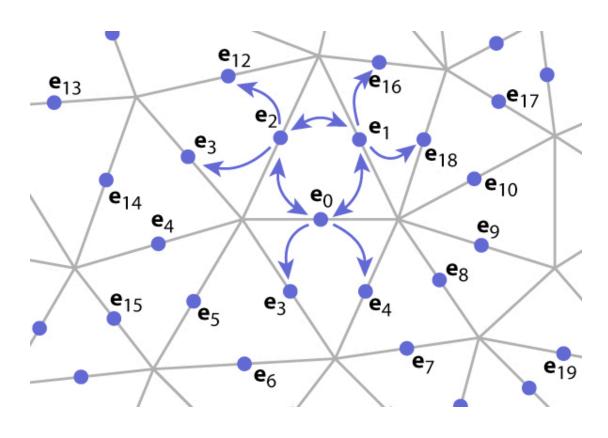


Winged-edge mesh

```
Edge {
  Edge hl, hr, tl, tr;
  Vertex h, t;
  Face l, r;
Face {
  // per-face data
  Edge e; // any adjacent edge
Vertex {
  // per-vertex data
  Edge e; // any incident edge
```

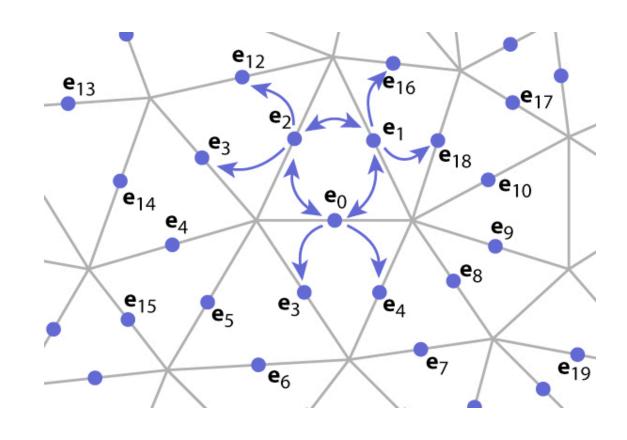






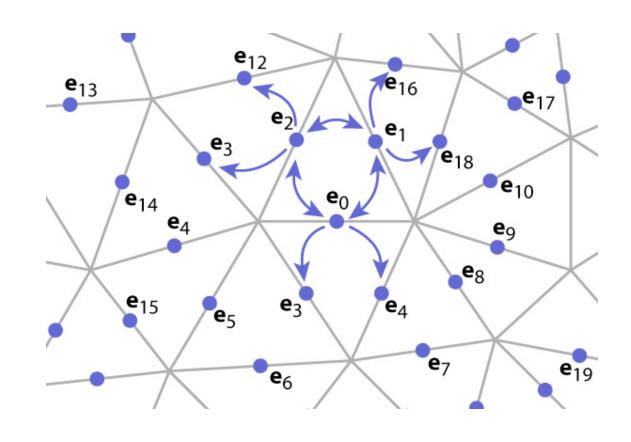
```
EdgesOfFace(f) {
    e = f.e;
    do {
        if (e.l == f)
            e = e.hl;
        else
            e = e.tr;
        } while (e != f.e);
    }
```

	hl	hr	tl	tr	
edge[O]	1	4	2	3	
edge[1]	18	0	16	2	
edge[0] edge[1] edge[2]	12	1	3	0	
		:			



```
EdgesOfVertex(v) {
    e = v.e;
    do {
        if (e.t == v)
            e = e.tl;
        else
            e = e.hr;
        } while (e != v.e);
    }
```

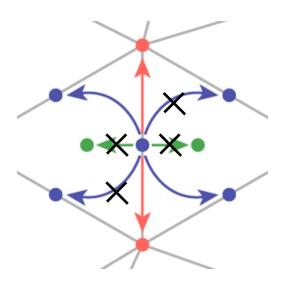
	hl	hr	tl	tr	
edge[O]	1	4	2	3	
edge[1]	18	0	16	2	
edge[2]	12	1	3	0	
		:			



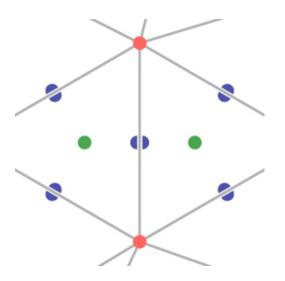
- array of vertex positions: I2 bytes/vert
- array of 8-tuples of indices (per edge)
 - head/tail left/right edges + head/tail verts + left/right tris
 - $int[n_F][8]$: about 96 bytes per vertex
 - 3 edges per vertex (on average)
 - (8 indices x 4 bytes) per edge
- add a representative edge per vertex
 - $int[n_V]$: 4 bytes per vertex
- total storage: I I 2 bytes per vertex
 - but it is cleaner and generalizes to polygon meshes

Winged-edge optimizations

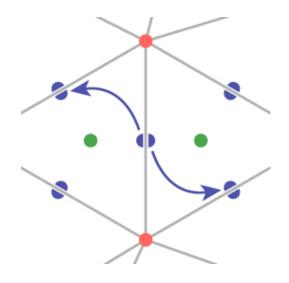
- Omit faces if not needed
- Omit one edge pointer on each side
 - results in one-way traversal



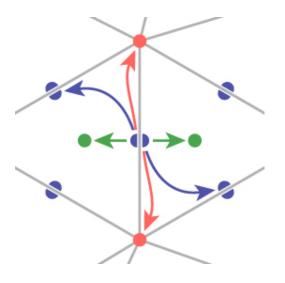
- Simplifies, cleans up winged edge
 - still works for polygon meshes
- Each half-edge points to:
 - next edge (left forward)
 - next vertex (front)
 - the face (left)
 - the opposite half-edge
- Each face or vertex points to one half-edge



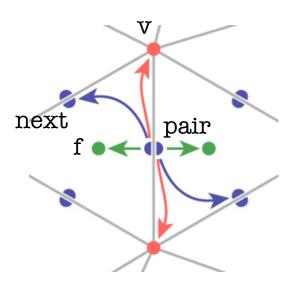
- Simplifies, cleans up winged edge
 - still works for polygon meshes
- Each half-edge points to:
 - next edge (left forward)
 - next vertex (front)
 - the face (left)
 - the opposite half-edge
- Each face or vertex points to one half-edge

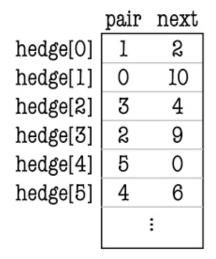


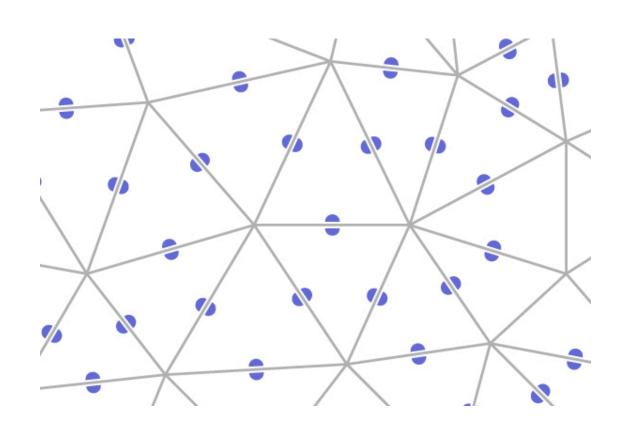
- Simplifies, cleans up winged edge
 - still works for polygon meshes
- Each half-edge points to:
 - next edge (left forward)
 - next vertex (front)
 - the face (left)
 - the opposite half-edge
- Each face or vertex points to one half-edge

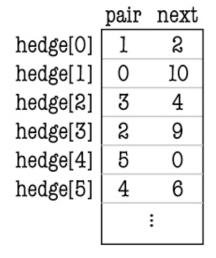


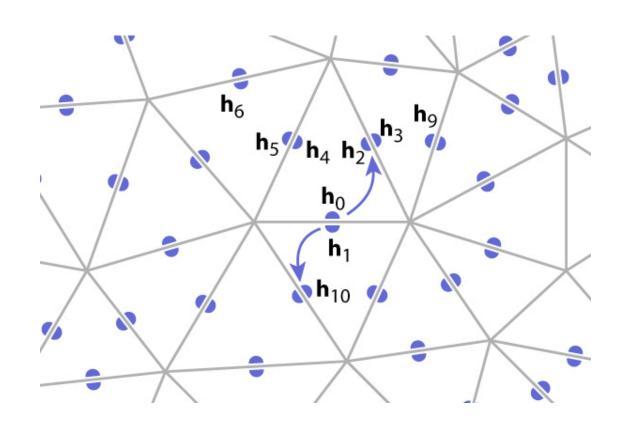
```
HEdge {
  HEdge pair, next;
  Vertex v;
  Face f;
Face {
  // per-face data
  HEdge h; // any adjacent h-edge
Vertex {
  // per-vertex data
  HEdge h; // any incident h-edge
```

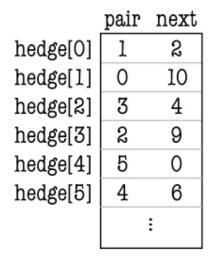


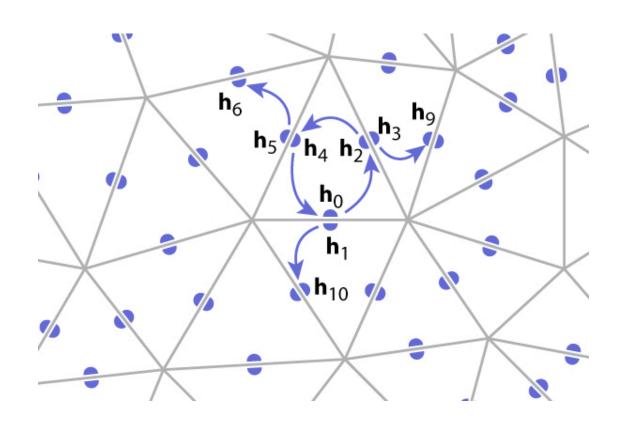






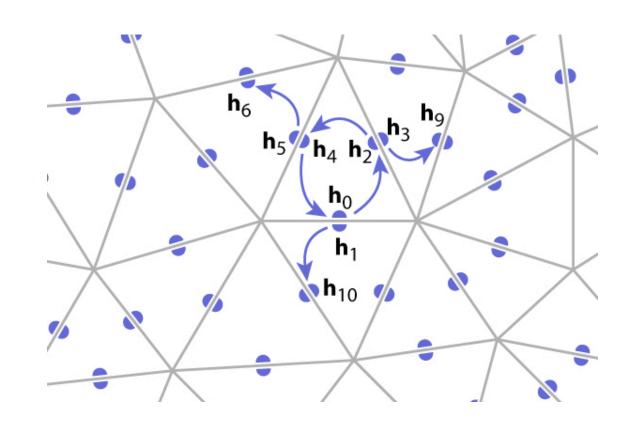






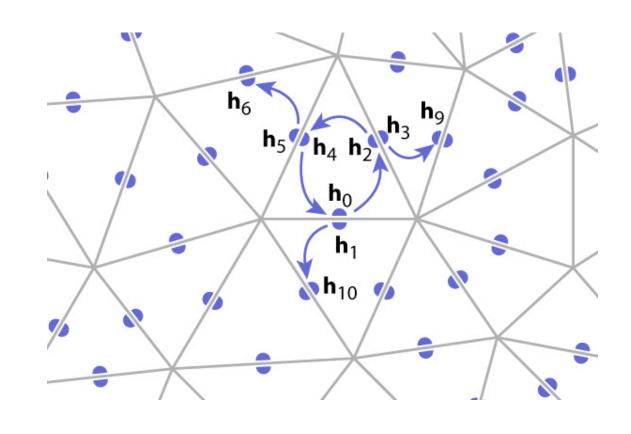
```
EdgesOfFace(f) {
    h = f.h;
    do {
        h = h.next;
    } while (h != f.h);
}
```

	pair	next
hedge[0]	1	2
hedge[1]	0	10
hedge[2]	3	4
hedge[3]	2	9
hedge[4]	5	0
hedge[5]	4	6
	:	



```
EdgesOfVertex(v) {
    h = v.h;
    do {
        h = h.pair.next;
    } while (h != v.h);
}
```

	pair	next
hedge[0]	1	2
hedge[1]	0	10
hedge[2]	3	4
hedge[3]	2	9
hedge[4]	5	0
hedge[5]	4	6
	:	



- array of vertex positions: I2 bytes/vert
- array of 4-tuples of indices (per h-edge)
 - next, pair h-edges + head vert + left tri
 - $-\inf[2n_F][4]$: about 96 bytes per vertex
 - 6 h-edges per vertex (on average)
 - (4 indices x 4 bytes) per h-edge
- add a representative h-edge per vertex
 - $-\inf[n_v]$: 4 bytes per vertex
- total storage: I I 2 bytes per vertex

Half-edge optimizations

- Omit faces if not needed
- Use implicit pair pointers
 - they are allocated in pairs
 - they are even and odd in an array

