

Greedy Heuristics

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Code Repository

<https://github.com/patrickmolina1/Evolutionary-Computation>

Problem Description

The task is to select a subset of nodes and form an optimal Hamiltonian cycle minimizing the total cost and travel distance. Each node is defined by three attributes: x-coordinate, y-coordinate, and cost. Exactly 50% of the nodes must be selected (rounded up if the total number is odd). The goal is to minimize the sum of the cycle length and the total cost of the selected nodes. Distances between nodes are computed using the Euclidean distance, rounded to the nearest integer. A distance matrix is precomputed after reading each instance and used throughout the optimization process, allowing instances to be represented solely by distance values

Methods

Four constructive heuristics are implemented and adapted to this problem:

1. **Random Solution** – Generates a feasible cycle by randomly selecting nodes.
2. **Nearest Neighbor (End Insertion)** – Adds the next node giving the best improvement when inserted only at the **end** of the path.
3. **Nearest Neighbor (Flexible Insertion)** – Adds the next node at any position (beginning, end, or inside) that best improves the objective function.
4. **Greedy Cycle** – Builds the cycle by repeatedly inserting nodes that minimize the total increase in distance and cost.

Here, “nearest” is understood as the **best improvement in the objective value**, not just geometric closeness. For each heuristic, **200 solutions** are generated from each starting node, and **200 random solutions** are also produced for comparison.

Computational experiment

CPU Specs

AMD Ryzen 7 5800H with Radeon Graphics

| Cores | Threads | Clockspeed | Turbo Speed |
|-------|---------|------------|-------------|
| 8 | 16 | 3.2 GHz | 4.4 GHz |

Pseudocode

Algorithm: Random Solution

Input: instance

Output: solution

1. Get $n \leftarrow$ number of nodes in instance
2. Calculate $\text{numToSelect} \leftarrow \lceil n / 2 \rceil$
3. Create shuffled copy of all nodes
4. Select first numToSelect nodes from shuffled list
5. Create order list with IDs of selected nodes
6. Shuffle order list randomly
7. Calculate totalDistance :
 - For each consecutive pair in order (including wrap-around):
 - Add distance from current to next node
8. Calculate $\text{totalNodeCost} \leftarrow$ sum of costs of selected nodes
9. Calculate $\text{totalCost} \leftarrow \text{totalDistance} + \text{totalNodeCost}$
10. Return solution with selected nodes, order, costs

Algorithm: Nearest Neighbor End Only

Input: instance

Output: solution

1. Select random startNode from all nodes
2. Initialize selected \leftarrow [startNode]
3. Initialize order \leftarrow [startNode.id]
4. Initialize remaining \leftarrow all nodes except startNode
5. Calculate numToSelect \leftarrow [n / 2]
6. While selected.size < numToSelect AND remaining is not empty:
 - a. Get lastNode \leftarrow last node in selected
 - b. Get firstNode \leftarrow first node in selected
 - c. Initialize bestCandidate \leftarrow null, minIncrease \leftarrow ∞
 - d. For each candidate in remaining:
 - i. Calculate distToCandidate \leftarrow distance[lastNode][candidate]
 - ii. Calculate distCandidateToFirst \leftarrow distance[candidate][firstNode]
 - iii. Calculate distLastToFirst \leftarrow distance[lastNode][firstNode]
 - iv. Calculate distanceIncrease \leftarrow distToCandidate + distCandidateToFirst - distLastToFirst
 - v. Calculate objectiveIncrease \leftarrow distanceIncrease + candidate.cost
 - vi. If objectiveIncrease < minIncrease:
 - Update minIncrease \leftarrow objectiveIncrease
 - Update bestCandidate \leftarrow candidate
 - e. If bestCandidate \neq null:
 - i. Add bestCandidate to end of selected
 - ii. Add bestCandidate.id to end of order
 - iii. Remove bestCandidate from remaining
7. Calculate totalDistance (sum distances between consecutive nodes in cycle)
8. Calculate totalNodeCost \leftarrow sum of costs of selected nodes
9. Calculate totalCost \leftarrow totalDistance + totalNodeCost
10. Return solution with selected nodes, order, costs

Algorithm: Nearest Neighbor All Positions

Input: instance

Output: solution

1. Select random startNode from all nodes
2. Initialize selected \leftarrow [startNode]
3. Initialize order \leftarrow [startNode.id]
4. Initialize remaining \leftarrow all nodes except startNode
5. Calculate numToSelect \leftarrow $\lceil n / 2 \rceil$
6. While selected.size < numToSelect AND remaining is not empty:
 - a. Initialize bestCandidateToAdd \leftarrow null, minSelectionMetric \leftarrow ∞
 - b. For each candidate in remaining:
 - i. Find minDistanceToTour \leftarrow minimum distance from candidate to any node in selected
 - ii. Calculate selectionMetric \leftarrow minDistanceToTour + candidate.cost
 - iii. If selectionMetric < minSelectionMetric:
 - Update minSelectionMetric \leftarrow selectionMetric
 - Update bestCandidateToAdd \leftarrow candidate
 - c. If bestCandidateToAdd \neq null:
 - i. Initialize bestPosition \leftarrow -1, minIncrease \leftarrow ∞
 - ii. For each position i in order:
 - Get prevNodeId \leftarrow order[i]
 - Get nextNodeId \leftarrow order[(i + 1) % order.size]
 - Calculate distPrevToNext \leftarrow distance[prevNodeId][nextNodeId]
 - Calculate distPrevToCandidate \leftarrow distance[prevNodeId][bestCandidateToAdd]
 - Calculate distCandidateToNext \leftarrow distance[bestCandidateToAdd][nextNodeId]
 - Calculate distanceIncrease \leftarrow distPrevToCandidate + distCandidateToNext - distPrevToNext
 - Calculate objectiveIncrease \leftarrow distanceIncrease + bestCandidateToAdd.cost
 - If objectiveIncrease < minIncrease:
 - Update minIncrease \leftarrow objectiveIncrease
 - Update bestPosition \leftarrow i + 1
 - iii. Insert bestCandidateToAdd at bestPosition in selected
 - iv. Insert bestCandidateToAdd.id at bestPosition in order
 - v. Remove bestCandidateToAdd from remaining
 - d. Else: break
7. Calculate totalDistance (sum distances between consecutive nodes in cycle)
8. Calculate totalNodeCost \leftarrow sum of costs of selected nodes
9. Calculate totalCost \leftarrow totalDistance + totalNodeCost
10. Return solution with selected nodes, order, costs

Algorithm: Greedy Cycle**Input:** instance, startNode**Output:** solution

1. Initialize selected \leftarrow [startNode]
2. Initialize order \leftarrow [startNode.id]
3. Initialize remaining \leftarrow all nodes except startNode
4. Calculate numToSelect \leftarrow $\lfloor n / 2 \rfloor$
5. While selected.size < numToSelect AND remaining is not empty:
 - a. Initialize bestCandidate \leftarrow null, bestPosition \leftarrow -1, minIncrease \leftarrow ∞
 - b. For each candidate in remaining:
 - i. For each position pos in order:
 - Get prevNodeId \leftarrow order[pos]
 - Get nextNodeId \leftarrow order[(pos + 1) % order.size]
 - Calculate distPrevToNext \leftarrow distance[prevNodeId][nextNodeId]
 - Calculate distPrevToCandidate \leftarrow distance[prevNodeId][candidate]
 - Calculate distCandidateToNext \leftarrow distance[candidate][nextNodeId]
 - Calculate distanceIncrease \leftarrow distPrevToCandidate + distCandidateToNext - distPrevToNext
 - Calculate objectiveIncrease \leftarrow distanceIncrease + candidate.cost
 - If objectiveIncrease < minIncrease:
 - Update minIncrease \leftarrow objectiveIncrease
 - Update bestCandidate \leftarrow candidate
 - Update bestPosition \leftarrow pos + 1
 - c. If bestCandidate \neq null:
 - i. Insert bestCandidate at bestPosition in selected
 - ii. Insert bestCandidate.id at bestPosition in order
 - iii. Remove bestCandidate from remaining
6. Calculate totalDistance (sum distances between consecutive nodes in cycle)
7. Calculate totalNodeCost \leftarrow sum of costs of selected nodes
8. Calculate totalCost \leftarrow totalDistance + totalNodeCost
9. Return solution with selected nodes, order, costs

Results

| Method | Instance A | Instance B |
|--------------------------------|-----------------------------|-----------------------------|
| Random Solution | 264419.52 (230885 - 300429) | 213914.41 (191803 - 235581) |
| Nearest Neighbor End Only | 103674.65 (90323 - 117672) | 69828.78 (62606 - 77415) |
| Nearest Neighbor All Positions | 72334.99 (71515 - 73823) | 48994.88 (47295 - 51030) |
| Greedy Cycle | 72598.91 (71488 - 74410) | 51532.22 (49001 - 57324) |

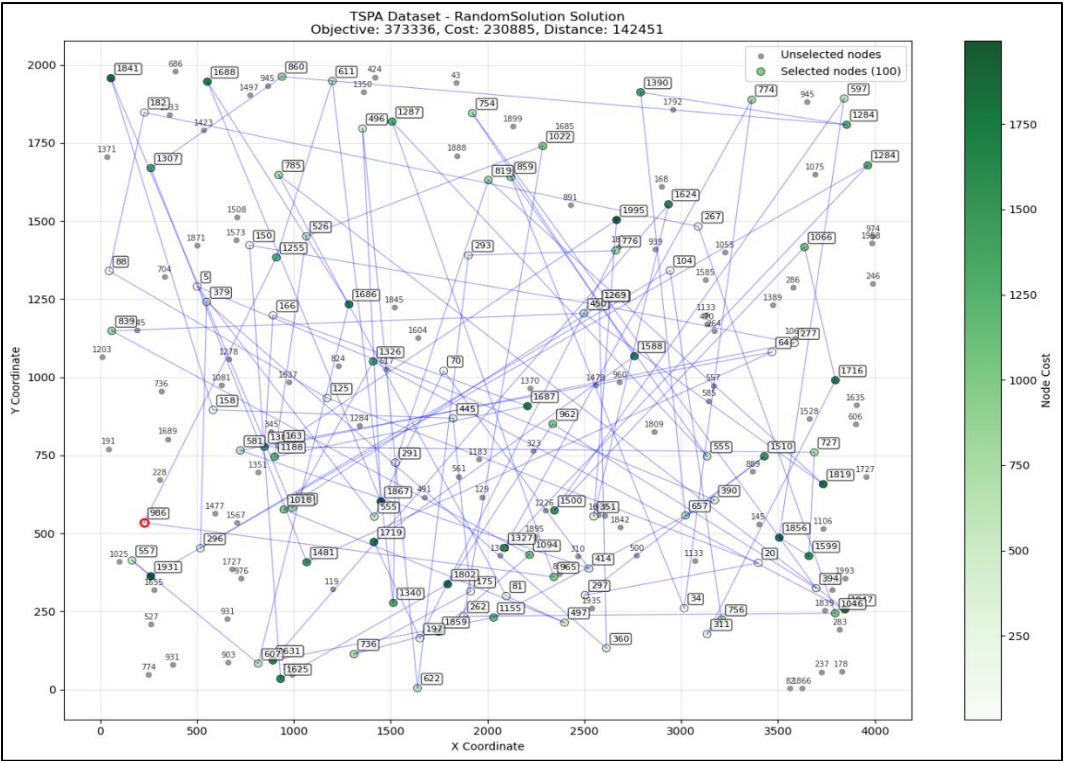
Running Times(ms)

| Method | Instance A | Instance B |
|--------------------------------|---------------|---------------|
| Random Solution | 0.06 (0 - 3) | 0.07 (0 - 4) |
| Nearest Neighbor End Only | 0.25 (0 - 3) | 0.30 (0 - 3) |
| Nearest Neighbor All Positions | 1.78 (0 - 13) | 2.08 (0 - 13) |
| Greedy Cycle | 2.32 (1 - 22) | 2.37 (1 - 23) |

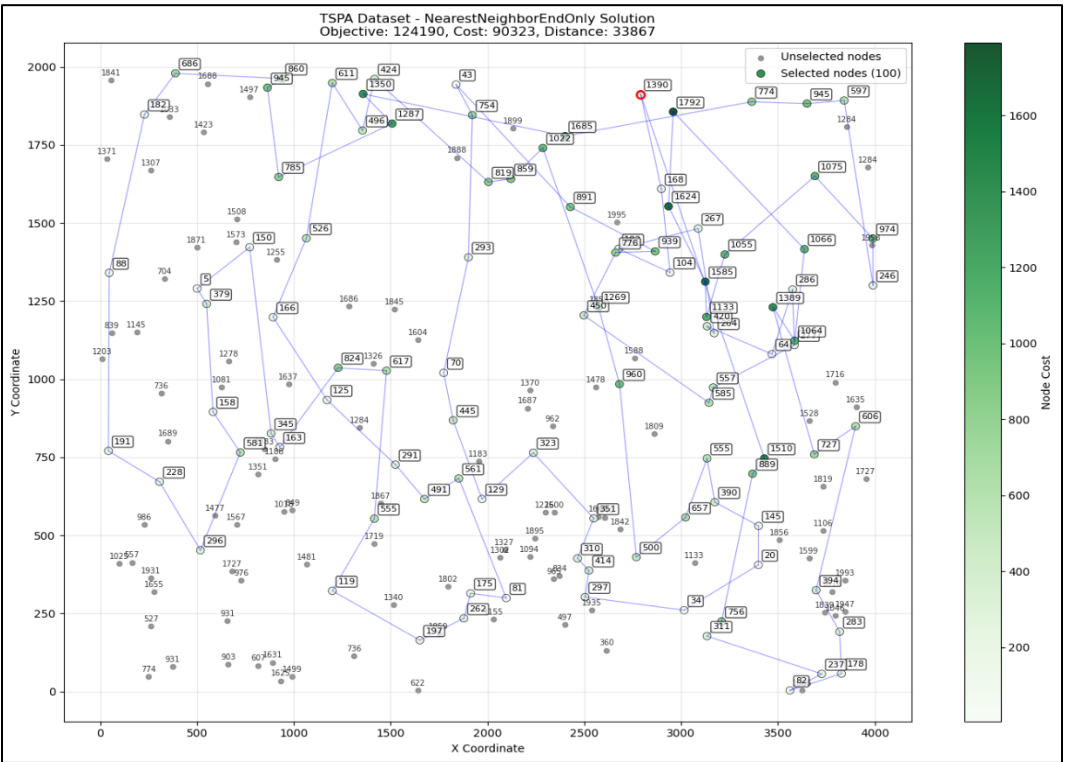
Best Solution

Instance A

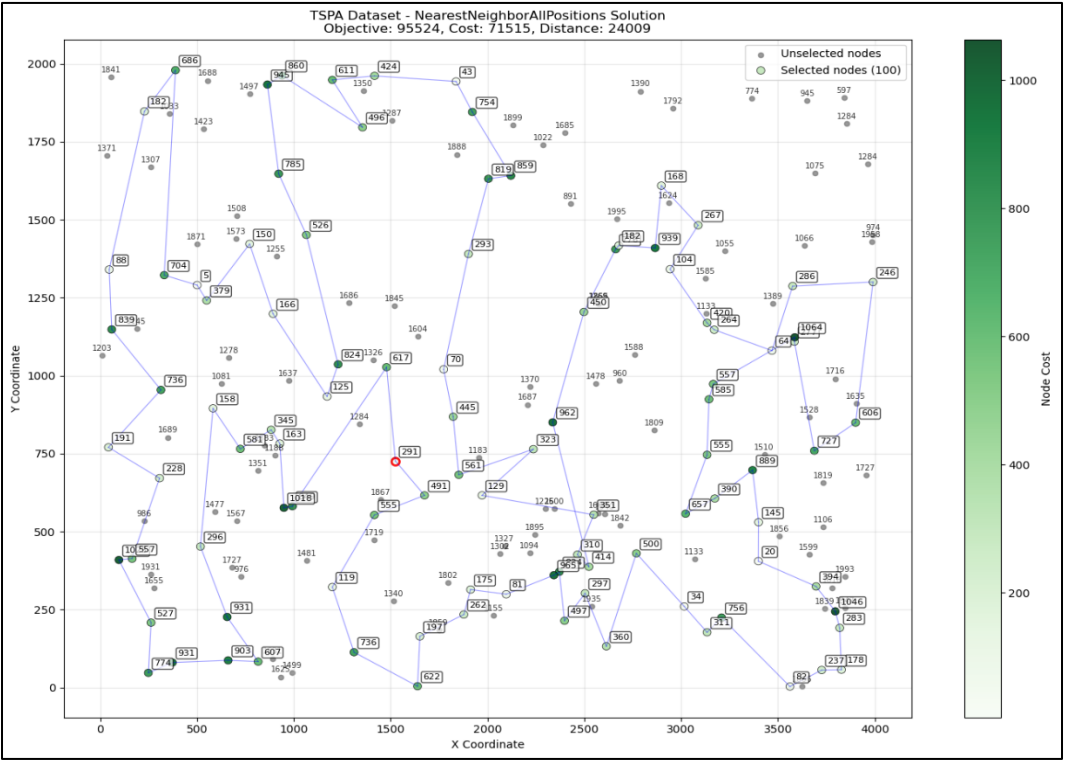
Random Solution



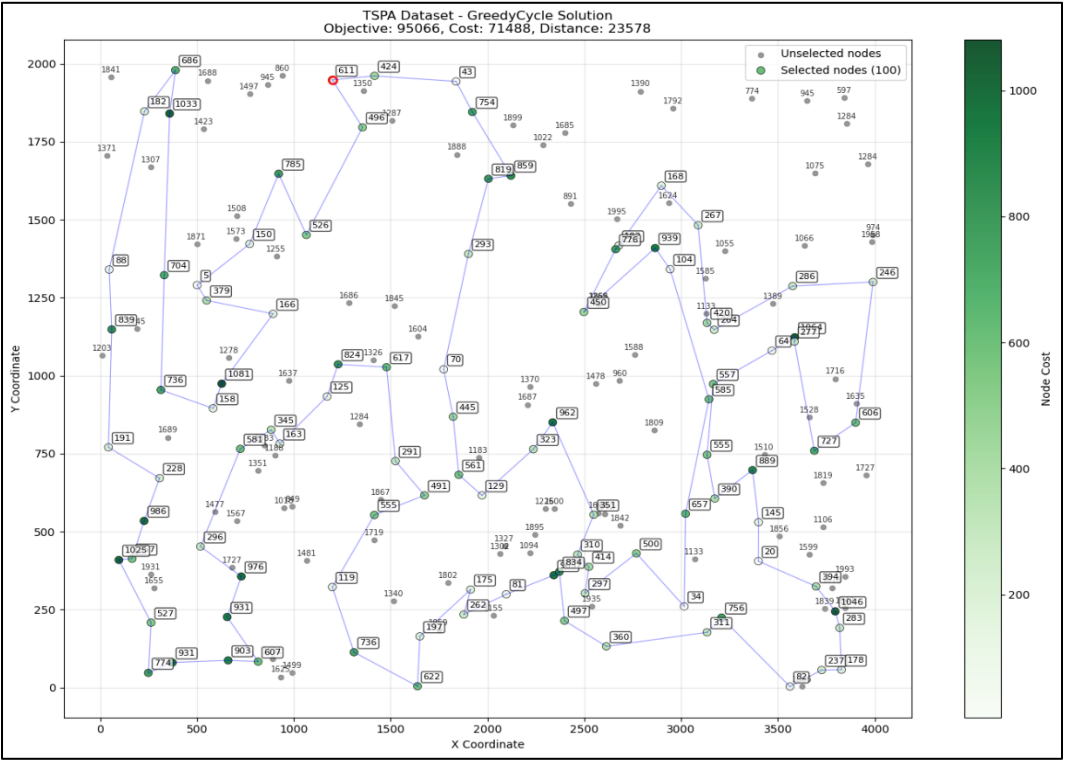
Nearest Neighbor End Only



Nearest Neighbor All Positions

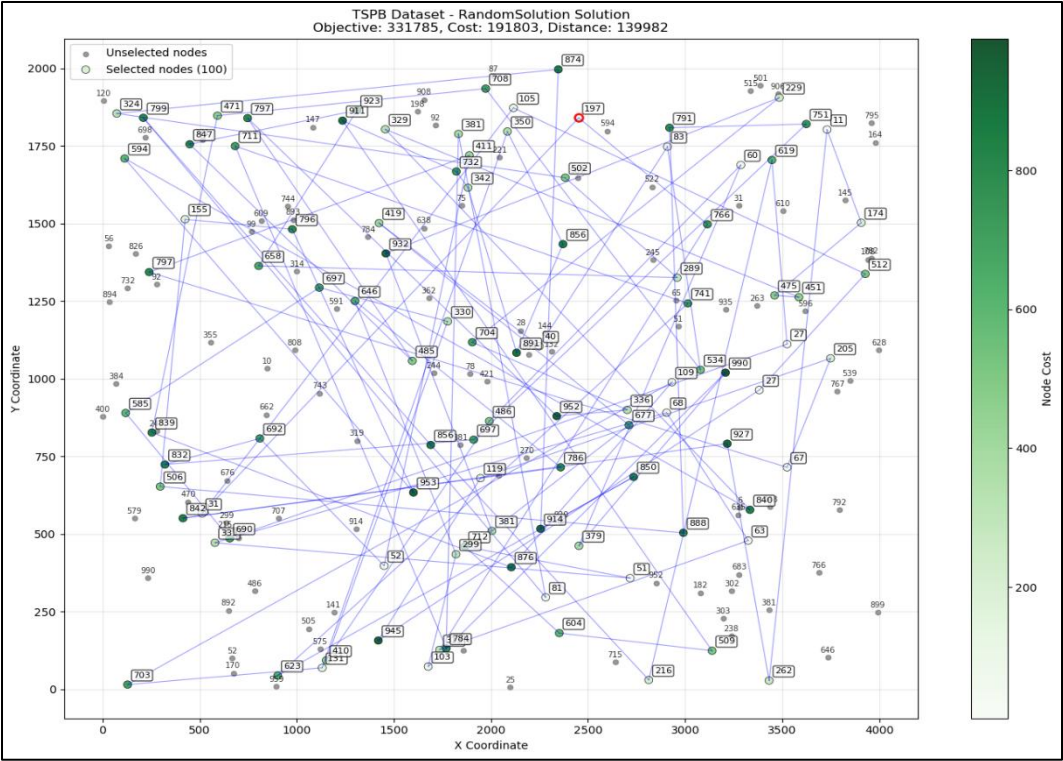


Greedy Cycle

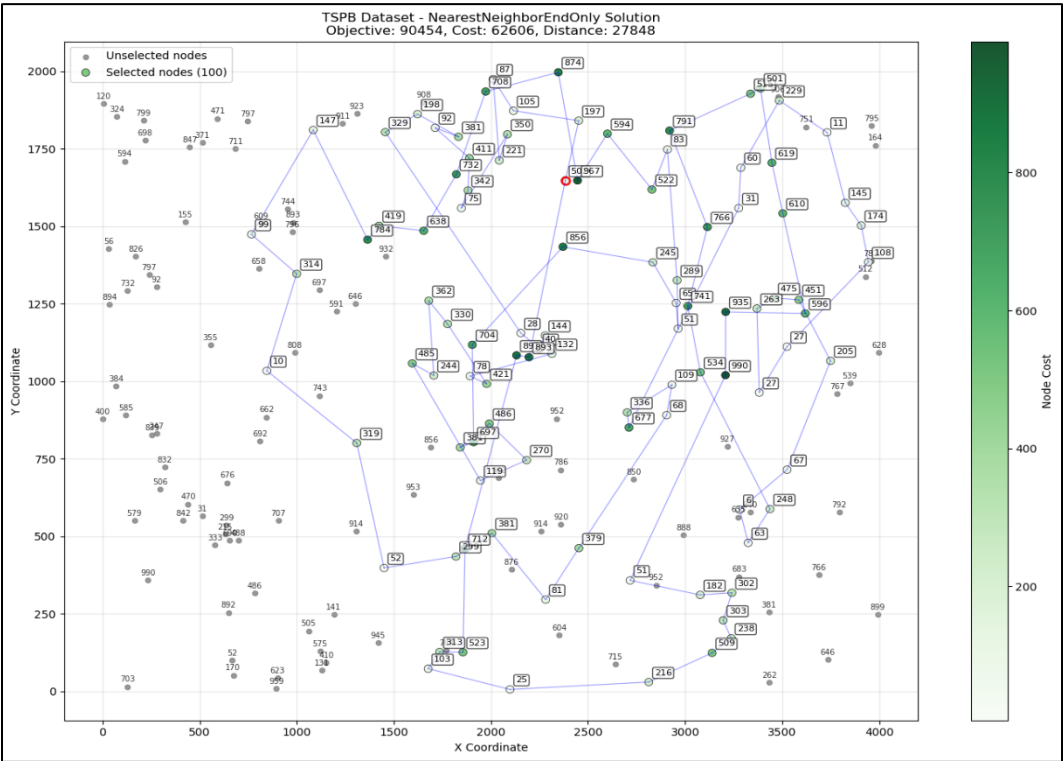


Instance B

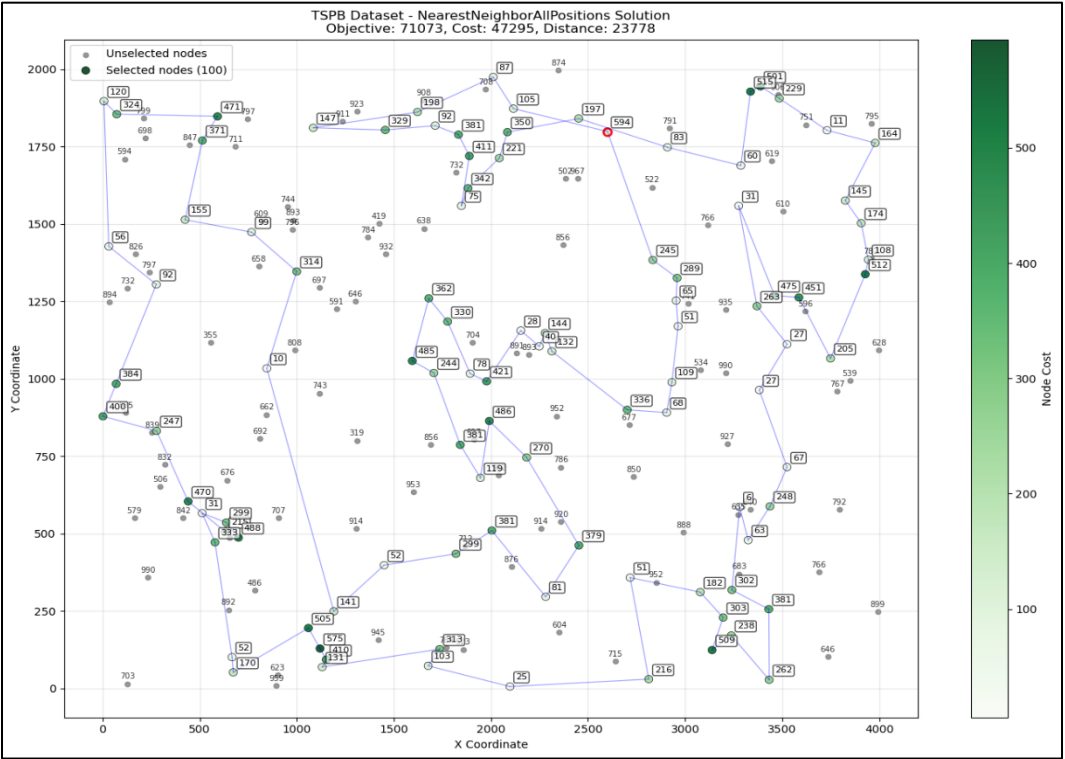
Random Solution



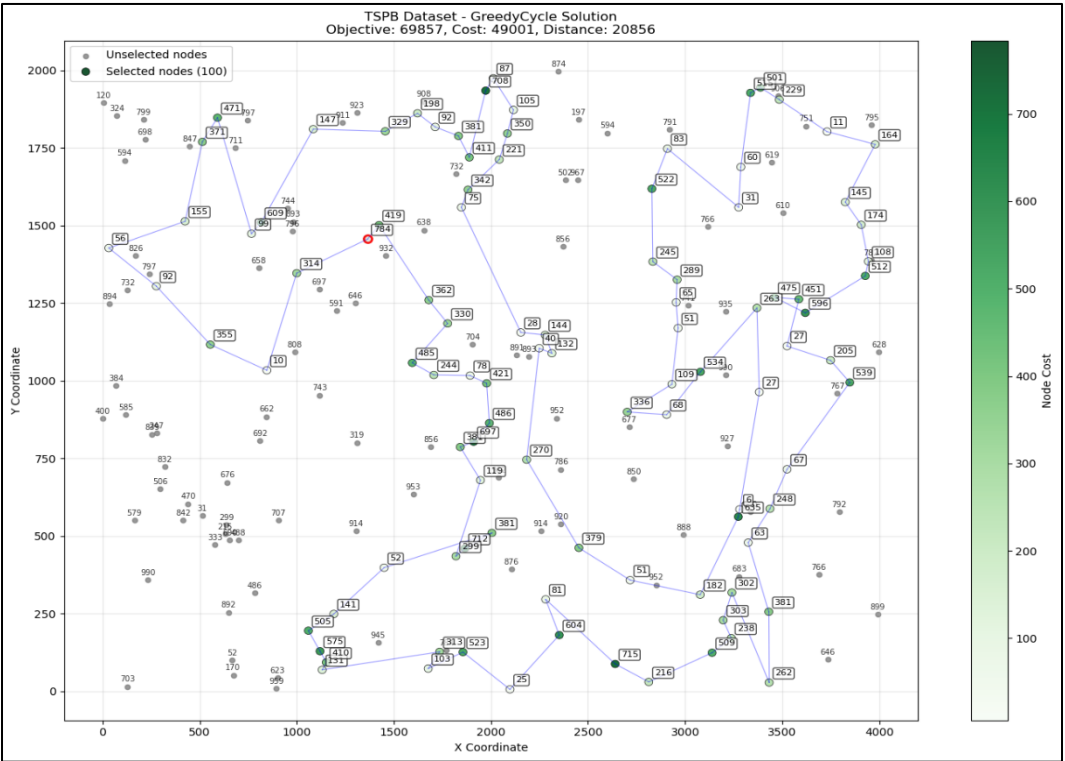
Nearest Neighbor End Only



Nearest Neighbor All Positions



Greedy Cycle



Cycles

Instance A

Random Solution

48-198-117-194-151-23-17-148-146-121-6-127-180-80-42-134-149-109-50-68-53-86-131-33-75-41-184-49-25-44-174-39-19-164-168-38-196-105-115-130-47-46-114-154-56-14-18-22-59-78-101-91-129-73-58-93-199-24-185-43-1-176-162-193-92-158-126-155-152-61-65-40-37-139-156-173-124-31-89-128-161-135-9-137-70-0-45-67-60-112-54-147-71-182-21-120-186-57-153-100

Nearest Neighbor End Only

73-144-49-62-14-106-178-185-165-40-52-55-148-9-102-15-183-89-137-176-80-63-94-152-97-1-101-120-78-145-92-57-129-2-167-37-114-186-23-143-0-117-46-115-59-151-133-79-53-180-154-135-123-162-51-118-65-116-139-193-41-42-43-184-160-34-22-18-108-93-140-68-153-170-64-21-7-164-90-27-95-138-3-32-155-132-39-119-8-196-81-31-113-175-16-171-44-25-179-91

Nearest Neighbor All Positions

151-51-149-131-65-116-43-42-184-84-112-4-190-10-177-30-54-160-34-181-146-22-18-108-159-193-41-139-115-59-118-46-68-140-93-0-117-143-183-89-186-23-137-176-80-79-94-63-152-97-1-124-148-9-62-102-144-14-49-178-106-185-165-90-81-196-40-119-52-55-57-129-92-179-145-78-31-56-113-175-171-16-25-44-120-2-75-101-86-26-100-53-180-154-135-70-127-123-162-133

Greedy Cycle

117-0-46-68-139-193-41-115-5-42-181-159-69-108-18-22-146-34-160-48-54-30-177-10-190-4-112-84-35-184-43-116-65-59-118-51-151-133-162-123-127-70-135-180-154-53-100-26-86-75-44-25-16-171-175-113-56-31-78-145-179-92-57-52-185-119-40-196-81-90-165-106-178-14-144-62-9-148-102-49-55-129-120-2-101-1-97-152-124-94-63-79-80-176-137-23-186-89-183-143

Instance B

Random Solution

155-174-61-113-106-193-153-87-97-165-192-126-195-39-69-15-175-108-124-58-197-10-161-128-152-0-134-159-82-2-28-49-143-37-115-182-44-150-125-55-189-107-84-50-119-162-83-178-79-132-77-17-157-187-199-185-164-16-98-11-78-144-140-81-110-36-21-74-41-30-64-137-20-47-130-99-6-88-100-171-120-122-156-129-53-59-56-24-19-177-13-32-166-179-8-7-86-66-3-117

Nearest Neighbor End Only

189-155-3-70-145-15-168-195-13-169-132-188-6-29-0-109-35-33-160-11-139-138-182-8-111-144-104-56-49-69-34-18-62-55-152-170-184-167-84-161-126-43-134-85-147-90-51-121-25-177-21-82-77-81-106-124-143-159-183-140-28-20-148-47-94-185-86-95-130-99-179-166-176-113-194-128-83-174-53-4-149-199-9-22-181-110-153-163-103-89-127-165-187-141-36-61-91-87-39-12

Nearest Neighbor All Positions

184-34-55-18-62-124-106-143-35-109-0-29-160-33-11-139-182-138-104-8-144-111-81-77-82-21-177-5-121-51-90-122-133-10-107-40-63-135-38-27-1-198-31-73-54-117-193-190-80-45-142-78-175-61-36-141-187-153-163-89-165-127-137-114-103-176-113-194-166-86-185-95-183-130-99-179-66-94-47-148-60-20-28-149-4-140-152-155-15-145-195-168-13-132-169-6-147-188-70-3

Greedy Cycle

85-51-121-131-135-63-122-133-10-90-191-147-6-188-169-132-13-161-70-3-15-145-195-168-29-109-35-0-111-81-153-163-180-176-86-95-128-106-143-124-62-18-55-34-170-152-183-140-4-149-28-20-60-148-47-94-66-22-130-99-185-179-172-166-194-113-114-137-103-89-127-165-187-146-77-97-141-91-36-61-175-78-142-45-5-177-82-87-21-8-104-56-144-160-33-138-182-11-139-134

Conclusion

- After the experiment it reveals a clear performance of **Greedy Cycle** and **Nearest Neighbor with All Positions** having similar results but **NN-all positions** has achieved the best solution quality for both TSPA and TSPB instances.
- **Nearest Neighbor with All Positions** insertion demonstrated significantly better performance than its **End Only variant**, highlighting the importance of considering multiple insertion points during solution construction rather than restricting nodes to path endpoints.
- The **Random Solution** method, while computationally trivial, produced substantially inferior results, serving primarily as a baseline to demonstrate the value of intelligent heuristic design.