

Quantification

Handout 2

Patrick D. Elliott

November 8, 2022

Contents

1	Reading	1
2	Quantifiers: the basics	1
3	The type of quantificational NPs	2
3.1	Subset-to-superset inferences	3
3.2	Law of non-contradiction	3
4	Predicates of predicates	3
5	Determiners	4
6	Generalized quantifiers	5

1 Reading

- Chapter 6 of *Invitation to Formal Semantics*, especially the sections on quantifiers.

2 Quantifiers: the basics

We'd translate the following sentence into predicate logic as $\forall x, \text{smiled}(x)$

- (1) Everybody smiled.

What, exactly, is the contribution of *everybody*? Informally, it looks like some kind of template:

$$(2) \quad \forall x. _ (x)$$

The predicate **smiled** fills in the slot.

We'll implement this idea in the lambda calculus by using a functional abstraction:

$$(3) \quad \mathbf{everybody} := \lambda P. \forall x. P(x)$$

That means that the type of *everybody* is $(E \rightarrow T) \rightarrow T$, i.e., a function from predicates to truth-values.

This is the type of a **quantifier**; as we'll see later on, this can also be the type of NPs more generally!

We can define the other quantifiers from predicate logic in a similar way:

$$(4) \quad \mathbf{something} := \lambda P. \exists x. P(x)$$

$$(5) \quad \mathbf{nothing} := \lambda P. \neg \exists x. P(x)$$

3 The type of quantificational NPs

Quantificational NPs in general seem to pattern with type E expressions, in terms of their distribution.

(6) A singer loves Frida.

(7) Frida loves a singer.

See also: *someone, everybody, nobody, some linguist, at least one linguist, at most one linguist, no linguist, few linguists*, etc.

Can quantificational NPs be of type E ? In fact, we can in fact show that they can't be.

3.1 Subset-to-superset inferences

- (8) Susan came yesterday morning.
 \Rightarrow Susan came yesterday.

Why does this hold?

Is the inference valid in the following case?

- (9) At most one letter came yesterday morning. $\Rightarrow?$ At most one letter came yesterday.

Exercise: figure out which (if any) quantificational NPs validate subset to superset inferences.

3.2 Law of non-contradiction

- (10) Mont Blanc is higher than 4000m, and Mont Blanc is not higher than 4000m.
- (11) More than two mountains are higher than 4000m, and more than two mountains are not higher than 4000m.

Exercise: again, figure out which (if any) quantificational NPs validate the law of non-contradiction.

4 Predicates of predicates

The solution, as we've alluded to, is to treat quantifiers as expressions of type $(E \rightarrow T) \rightarrow T$.

What does this mean for their meaning?

A helpful way to think about what quantifiers do exploits a correspondence between functions $f : \mathbf{Dom}_\sigma \mapsto \{\mathbf{true}, \mathbf{false}\}$, and *sets* of things in \mathbf{Dom}_σ .

We can define the following translation; $Set(f)$ is called the *set characterized by f* , and we can freely switch between sets and functions without losing information:

$$\text{Set}(f) = \{ x \mid f(x) = \mathbf{true} \}$$

This allows us to think of predicates, for example, as denoting *sets of individuals*.

Quantifiers, on the other hand, denote *sets of predicates*.

As an example, consider *everything* and *nothing*:

$$(12) \quad \llbracket \mathbf{everything} \rrbracket = \{ f \mid \forall x \in \mathbf{Dom}_E, f(x) = \mathbf{true} \}$$

Alternatively, if we think of the functions in the denotation of /everything as themselves sets:

$$(13) \quad \llbracket \mathbf{everything} \rrbracket = \{ X \mid \mathbf{Dom}_E \subseteq X \}$$

Exercise: what about *something* and *nothing*?

5 Determiners

Based on what we know about the type of NPs, and the type of quantifiers, we can conclude what the type of a determiner such as *some*, *no*, and *every* should be:

$$(E \rightarrow T) \rightarrow (E \rightarrow T) \rightarrow T$$

Determiners denote functions from predicates to quantifiers.

In quasi-predicate-logic notation, we can write the following:

$$(14) \quad \mathbf{every} := \lambda R. \lambda P. \forall x [R(x) \rightarrow P(x)]$$

6 Generalized quantifiers

What is the meaning of *every cat*?

In predicate logic, it would be something like the following:

$$(15) \quad \lambda P. \forall x[\mathbf{cat}(x) \rightarrow P(x)]$$

What set of sets does this quantifier characterize?

$$(16) \quad \{ P \subseteq \mathbf{Dom}_E \mid \{ x \mid x \text{ is a cat} \} \subseteq P \}$$

some dog:

$$(17) \quad \lambda P. \exists x[\mathbf{dog}(x) \wedge P(x)]$$

$$(18) \quad \{ P \subseteq \mathbf{Dom}_E \mid \{ x \mid x \text{ is a dog} \} \cap P \neq \emptyset \}$$

We can derive *everything* and *something*, by replacing *cat* and *dog* with \mathbf{Dom}_E .

The strategy of characterizing quantifiers in terms of sets of sets generalizes to any quantificational expression.

$$(19) \quad \llbracket \mathbf{nothing} \rrbracket = \{ P \subseteq \mathbf{Dom}_E \mid P = \emptyset \}$$

$$(20) \quad \llbracket \mathbf{exactlyTwoThings} \rrbracket = \{ P \subseteq \mathbf{Dom}_E \mid \mathbf{Card}(P) = 2 \}$$

$$(21) \quad \llbracket \mathbf{atleastTwoThings} \rrbracket = \{ P \subseteq \mathbf{Dom}_E \mid \mathbf{Card}(P) \geq 2 \}$$

Sets of sets of entities are called **generalized quantifiers**.

Determiners characterize curried *relations* between sets, we call such relations between sets **determiner relations**.

$$(22) \quad \llbracket \mathbf{every} \rrbracket = \{ (R, P) \mid R \subseteq P \}$$

$$(23) \quad \llbracket \mathbf{some} \rrbracket = \{ (R, P) \mid R \cap P \neq \emptyset \}$$