

Propositional logic and recursive datatypes

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- Haskell's type system can be used to represent grammars as datatypes.
- Grammars of non-trivial formal languages/natural languages typically involve *recursive* rules for well-formed formulas/sentences
- An example from a CFG for natural language:

S -> NP VP

VP -> V S

- "She is upset."
- "Zach said she is upset."
- "Jenna believes Zach said she is upset."

- In order to implement recursive grammars in Haskell, we use *recursive datatypes*.
 - Operations that apply to inhabitants of recursive datatypes are (typically) recursive functions.
 - Today we'll learn about recursive datatypes/functions in Haskell, using the simple example of propositional logic.

- Next week we'll talk about semantics, again using propositional logic as an example.
- For homework, read chapter 5 of (van Eijck, Jan and Unger, Christina, 2010) "Formal semantics for fragments".

Propositional logic

Sentences of propositional logic:

- $p \wedge q$
- $p \wedge (q \rightarrow r)$
- $p \vee \neg(q \wedge p)$
- $\perp \rightarrow (\top \vee (q \wedge r))$

To state a grammar for propositional logic, we first need a set of variables

Var := $\{ p, q, r, \dots \}$, and a set of constants $\{ \top, \perp \}$.

- If $p \in \mathbf{Var} \cup \{ \top, \perp \}$, then p is a sentence of propositional logic.
- If ϕ, ψ are sentences of propositional logic, then $(\phi \wedge \psi)$ is a sentence of propositional logic.
- If ϕ, ψ are sentences of propositional logic, then $(\phi \vee \psi)$ is a sentence of propositional logic.
- If ϕ, ψ are sentences of propositional logic, then $(\phi \rightarrow \psi)$ is a sentence of propositional logic.
- If ϕ is a sentence of propositional logic, then so is $\neg\phi$

The grammar as a datatype

We'll use a *sum type* to implement the grammar of propositional logic. We can start with just the following:

```
data PropL = PVar String deriving (Eq, Show)
```

```
ghci> :t (PVar "p1")
```

```
PropL
```

```
ghci> :t (PVar "p3")
```

```
PropL
```


Adding negation: recursive datatypes!

```
data PropL = PVar String | PNot PropL deriving (Eq, Show)
```

Note that we *reuse* **PropL** in the constructor for negative sentences; this means that if an expression **p** is of type **PropL**, then **PNot p** is of type **PropL**.

```
ghci> :t (PNot (PVar "p1"))
PropL
ghci> :t (PNot (PNot (PNot (PVar "p1"))))
PropL
```

Note that this already gives us infinite inhabitants of type **PropL**.

Adding connectives

```
data PropL = PVar String | PNot PropL | PAnd PropL PropL | POr
  ↳ PropL PropL deriving (Eq, Show)
```

```
ghci> :t (PAnd (PVar "p1") (POr (PVar "p2") (PVar "p3")))
PropL
ghci> :t (PAnd (PVar "p1"))
error
```

Bonus: infix constructors!

Functions that take two arguments can be used as *infix operators* by enclosing them in backticks. This also goes for constructors:

```
ghci> :t PAnd
      PropL -> PropL -> PropL
ghci> :t ((PVar "p1") `PAnd` (PVar "p2"))
      PropL
```

You can even use infix constructors in the data declaration:

```
data PropL = PVar String | PNot PropL | PropL `PAnd` PropL | PropL
  <-> `POr` PropL deriving (Eq, Show)
```

Implementing a custom **Show** instance

Implementing a custom **Show** instance for **PropL** simply amounts to defining a function **show** of type **PropL -> String**.

```
data PropL = PVar String | PNot PropL | PropL `PAnd` PropL | PropL
  ⇨ `POr` PropL deriving Eq
```

```
instance Show PropL where
```

```
  show (PVar s) = s
```

```
  show (PNot p) = "~" ++ show p
```

```
  show (p `PAnd` q) = "(" ++ show p ++ " & " ++ show q ++ ")"
```

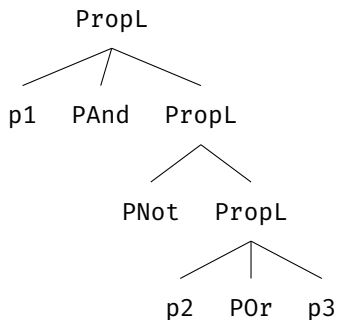
```
  show (p `POr` q) = "(" ++ show p ++ " | " ++ show q ++ ")"
```

The **Show** instance we just declared will automatically be used by `ghci`.

```
ghci> ((PVar "p1") `PAnd` (PNot ((PVar "p1") `POr` (PVar "p3"))))  
(p1 & ~(p2 | p3))
```

We can also use it explicitly by calling **show** on something of type **PropL**.

Recursive datatypes are used to create an Abstract Syntax Tree (AST) for sentences of propositional logic.



Let's say that we want to compute the number of operators in a formula. In order to do so we'll need a recursive function **opsNr**.

First, we define the base of the recursion (where the recursion halts):

```
opsNr :: PropL -> Int
opsNr (PVar _) = 0
```

For all other cases we need recursion:

```
opsNr :: PropL -> Int
opsNr (PVar _) = 0
opsNr (PNot p) = 1 + opsNr p
opsNr (PAnd p q) = 1 + opsNr p + opsNr q
opsNr (POr p q) = 1 + opsNr p + opsNr q
```


Another recursive function: depth

```
depth :: PropL -> Int
depth (PVar _) = 0
depth (PNot p) = 1 + depth p
opsNr (PAnd p q) = undefined
opsNr (POr p q) = undefined
```

```
depth :: PropL -> Int
depth (PVar _) = 0
depth (PNot p) = 1 + depth p
depth (PAnd p q) = 1 + max (depth p) (depth q)
depth (POr p q) = 1 + max (depth p) (depth q)
```

Exercise: gather names

- **Exercise:** write a recursive function that returns a list of all of the variables that occur in a formula.
- As a bonus, remove duplicates and sort the output alphabetically.

Fin

van Eijck, Jan and Unger, Christina (2010). *Computational Semantics with Functional Programming*, Cambridge University Press.