

Ellipsis, binding, & Logical Form

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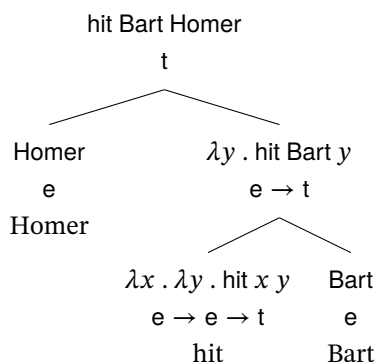
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1 Overview

- Pronouns: the semantics of assignment-sensitivity.
 - The presentation of assignment-sensitivity will be largely based on Charlow (2018); see also Buring (2005) for another excellent reference on this topic.
- Strict vs. sloppy readings: the classic account (Sag 1976, Williams 1977).
- Focus: the semantics of alternatives.
- Rooth's account of strict vs. sloppy readings and ellipsis identity.

2 The semantics of pronominals

- Background:
 - Semantic types:
 - * $e := \text{Homer} \mid \text{Bart} \mid \text{Burns} \mid \dots$
 - * $t := 1 \mid 0$
 - * If a and b are types then $a \rightarrow b$ is the type of a function from inhabitants of a to inhabitants of b .
 - * There is one basic rule of semantic composition – *function application*.
 - A basic computation:¹



- Sentences containing free pronouns aren't true simpliciter.

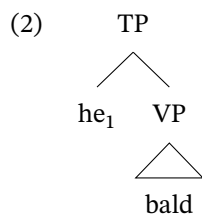
¹ I assume a one-to-one mapping between *merge* in the syntactic computation, and *function application* in the semantic computation.

The direction of application is flexible: $\llbracket [\alpha \beta] \rrbracket = \llbracket [\alpha] \rrbracket (\llbracket [\beta] \rrbracket)$ or $\llbracket [\beta] \rrbracket (\llbracket [\alpha] \rrbracket)$, whichever is defined.

Generally, complements and specifiers are the arguments of their sisters, although this parallelism will break down when we consider quantificational DPs.

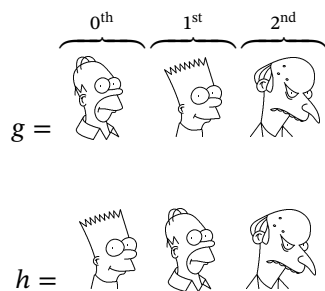
(1) He is bald.

- Intuitively, (1) is *true* iff the individual who the speaker intended this tokening *he* to pick out is bald.
- First, we need a formal device for distinguishing the intended reference of pronoun tokenings at the level of the semantics.
- In order to do this, let's suppose that pronouns are assigned an *index* $n \in \mathbb{N}$.



- In formal semantics, the standard way of modelling the context of utterance is via *assignments*. Assignments are ways of mapping indices/variables to entities in the domain.
- Formally, we can treat an *assignment* as an n -long *sequence* of individuals, given a set of indices $\{i \mid 0 \leq i \leq n\}$.²
- Here are two example assignments, given a set of indices $\{0, 1, 2\}$

² We'll adopt the convention whereby initial member of the sequence is the 0th.



- We can think of the meaning of a pronoun as picking out the n^{th} entity in the assignment g .
- Typically, the context of utterance is represented as a parameter on the interpretation function $\llbracket \cdot \rrbracket$ (see, e.g., Heim & Kratzer 1998).

$$\llbracket \text{he}_n \rrbracket^g = g_n$$

- A sentence, such as he_1 is bald ends up as true or false *relative* to the assignment.

$$(3) \quad a. \quad \llbracket \text{he}_1 \text{ is bald} \rrbracket^g = \text{bald } g_1 = \text{bald } \begin{array}{c} \text{Bart} \\ \text{Simpson} \end{array} = 0$$

$$b. \quad \llbracket \text{he}_1 \text{ is bald} \rrbracket^h = \text{bald } h_1 = \text{bald } \begin{array}{c} \text{Homer} \\ \text{Simpson} \end{array} = 1$$

- Equivalently, we can think of sentence meanings *characteristic functions* of assignments.

$$(4) \quad \llbracket \text{he}_1 \text{ is bald} \rrbracket = \lambda g . \text{bald } g_1$$

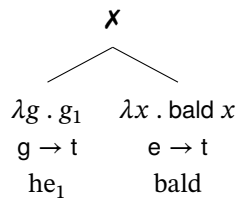
- The sentence *he₁ is bald* maps an assignment *g* to *true*, if and only if the 1st member of *g* is bald.
- When *g* is fed in as the argument of the function, the result is 0, since the 1st member of *g* is Bart, and Bart is not bald. When *h* is fed in as the argument, the result is 1, since the 1st member of *h* is Homer, and Homer is bald.
- Recall from intro semantics that we can retrieve a *set* from a function to truth values. We can equivalently think of the meaning of an assignment-sensitive sentence as a set of assignments.

$$(5) \llbracket \text{he}_1 \text{ is bald} \rrbracket = \left(\begin{array}{ccc} \text{Homer} & \text{Homer} & \text{Homer} , \\ \text{Homer} & \text{Homer} & \text{Bart} , \\ \text{Homer} & \text{Homer} & \text{Bart} , \\ \text{Bart} & \text{Homer} & \text{Homer} , \\ \text{Bart} & \text{Bart} & \text{Homer} , \\ \text{Bart} & \text{Bart} & \text{Bart} , \\ \text{Bart} & \text{Bart} & \text{Bart} , \\ \text{Bart} & \text{Bart} & \text{Bart} , \\ \text{Bart} & \text{Bart} & \text{Bart} , \\ \dots & \dots & \dots \end{array} \right)$$

- Notice that the 1st column *only* features Homer and Bart. Since only Homer and Bart are bald.
- How do we derive such meanings compositionally?
- For us, pronouns will denote functions from assignments to individuals, i.e., *assignment sensitive individuals*.

$$\llbracket \text{pro}_n \rrbracket = \lambda g . g_n$$

- How do we get assignment-sensitive individuals to play nicely with ordinary compositional apparatus?



- We need two type-shifters – η and \oplus – in order to integrate assignment sensitive meanings (Charlow 2018).³
 - η shifts an ordinary meaning to a trivially *assignment sensitive* meaning.
 - \oplus is an enriched form of function application that we use to compose assignment-sensitive meanings.

³ Together with the type-constructor for assignment-sensitive meanings G defined below, η and \oplus form an *applicative functor* (see Charlow 2018 for a demonstration).

$$(6) \quad G a := g \rightarrow a$$

Applicative functors provide a useful technique for modelling *enriched* semantic composition.

$$(7) \quad \begin{array}{ll} \text{a.} & a^\eta = \lambda g . a \\ \text{b.} & f \oplus x = \lambda g . f g (x g) \end{array}$$

- Let's see how these shifters work in practice:

$$\begin{array}{c} \lambda g . \text{bald } g_1 \\ \quad \quad \quad \downarrow \beta\text{-reduce} \\ \lambda g . [\lambda x . \text{bald } x] g_1 \\ \quad \quad \quad \downarrow \beta\text{-reduce} \\ \lambda g . [[\lambda g . \lambda x . \text{bald } x] g] ([\lambda g . g_1] g) \\ \quad \quad \quad \downarrow \\ \quad \quad \quad \oplus \\ \quad \quad \quad \swarrow \quad \searrow \\ \lambda g . g_1 \quad \lambda g . \lambda x . \text{bald } x \\ \quad \quad \quad \quad \quad \downarrow \eta \\ \quad \quad \quad \quad \quad \lambda x . \text{bald } x \end{array}$$

- Pronominals can be interpreted as *bound variables*. The truth of (8) does not depend on what the 1st member of the assignment is, i.e. (8) is *true* if Lisa said that Bart teased Lisa, and $g_1 = \text{Homer}$.

$$(8) \quad \begin{array}{l} \text{Someone}^1 \text{ said that Bart teased them}_1. \\ \exists x [x \text{ said } (\text{teased } x \text{ Bart})] \end{array}$$

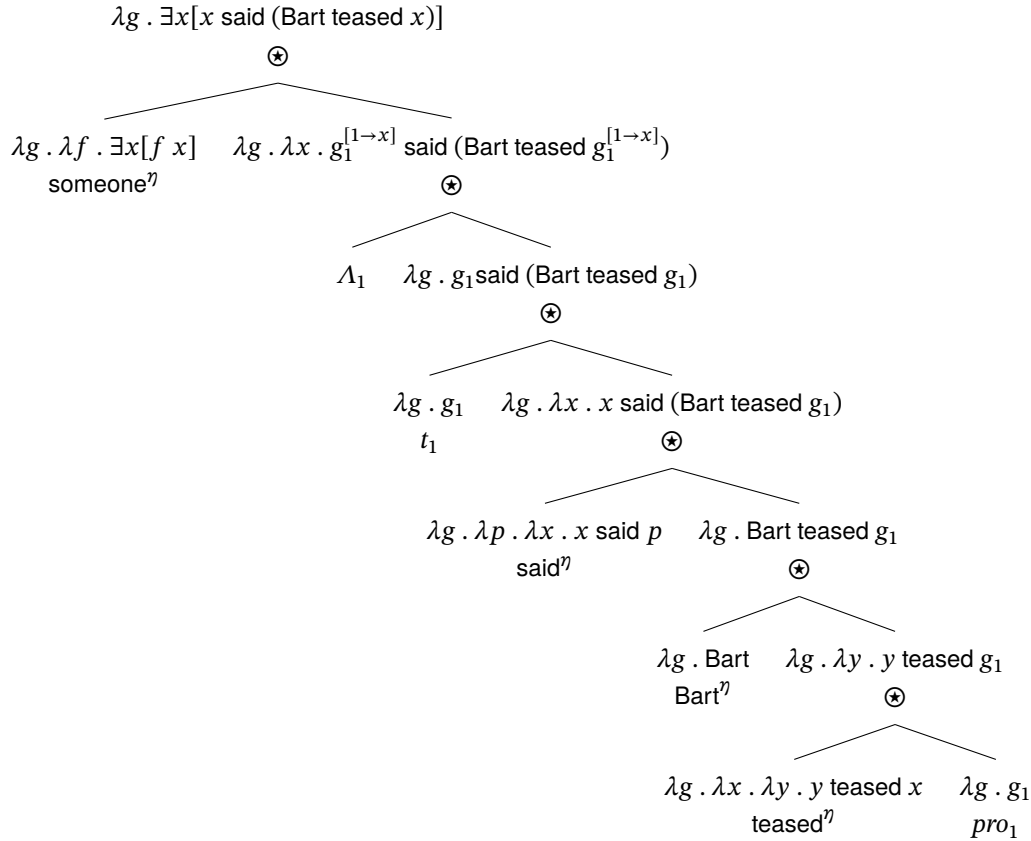
- In order to account for bound variable readings of pronominals, we're going to introduce a new operator Λ_n . Λ_n composes with an assignment-sensitive expression and *abstracts* over pronominals with a matching index pro_n .

$$(9) \quad \Lambda_n \mathfrak{t} = \lambda g . \lambda x . \mathfrak{t} g^{[n \rightarrow x]}$$

- We can take quantificational DPs such as *someone* to simply denote generalised quantifiers.

(10) $\text{someone} = \lambda f . \exists x[f\ x]$

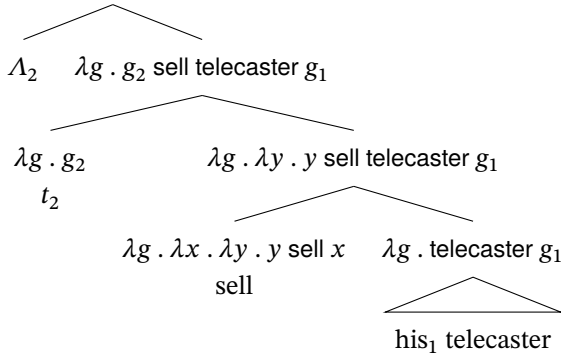
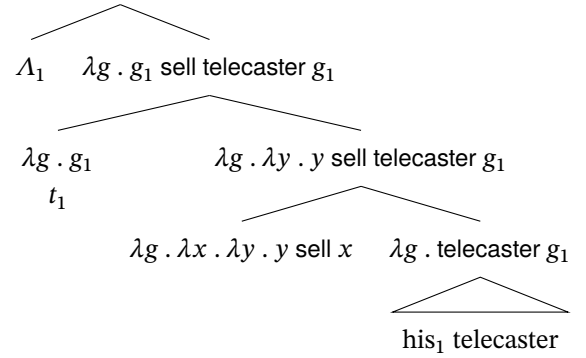
- Now we have all the apparatus we need in order to derive *bound readings* of pronominals.



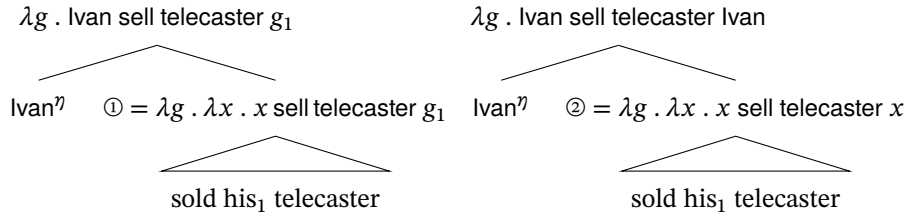
- Note that we can treat traces as having the same semantics as pronouns – namely, they’re assignment-sensitive individuals.

3 *Strict vs. sloppy readings*

- There are two conceivable LF’s for the following sentence, under the reading *Ivan sold Ivan’s telecaster*:
- We’ll refer to the LF in (11a) as the *accidental coreference* LF. Here, the pronominal just happens to pick out the same individual as the subject, given an assignment *g*.

①: *accidental coreference* $\lambda g . \lambda x . x \text{ sell telecaster } g_1$ ②: *binding* $\lambda g . \lambda x . x \text{ sell telecaster } x$ 

- Note that ① and ② denote distinct functions:
 - ① denotes a function from an assignment g , to a function from individuals x to *true* iff x sold g_1 's telecaster. In the utterance context $g_1 = \text{Ivan}$. ① is *assignment sensitive* because it contains a free pronoun, i.e., a pronoun not bound by a matching Λ .
 - ② denotes a function from an assignment g , to a function from individuals x to *true* iff x sold x 's telecaster. ② is *assignment insensitive*, because every occurrence of a pronoun/trace is bound by a matching Λ .
- When the VP's compose with the type e expression *Ivan*, the resulting propositions are identical – just so long as $g_1 = \text{Ivan}$.



Question: how is the *identity condition* on ellipsis met under the sloppy reading?

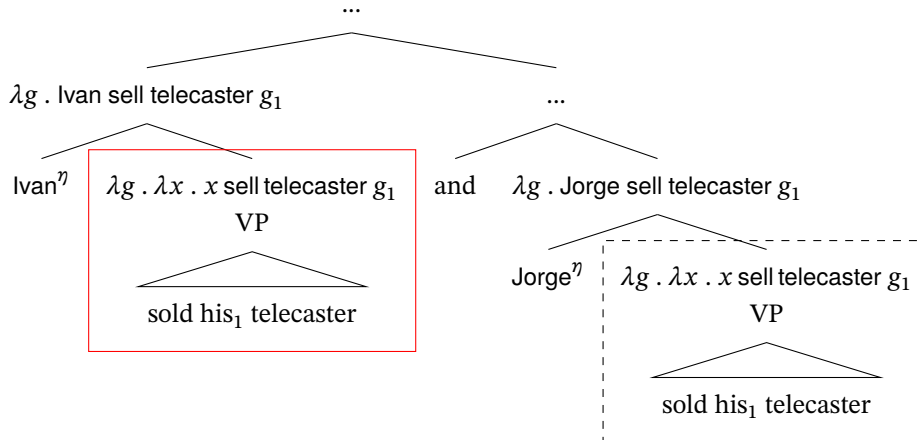
(15) *The parallelism condition*

Ellipsis of EC requires semantic identity with with an antecedent constituent AC.

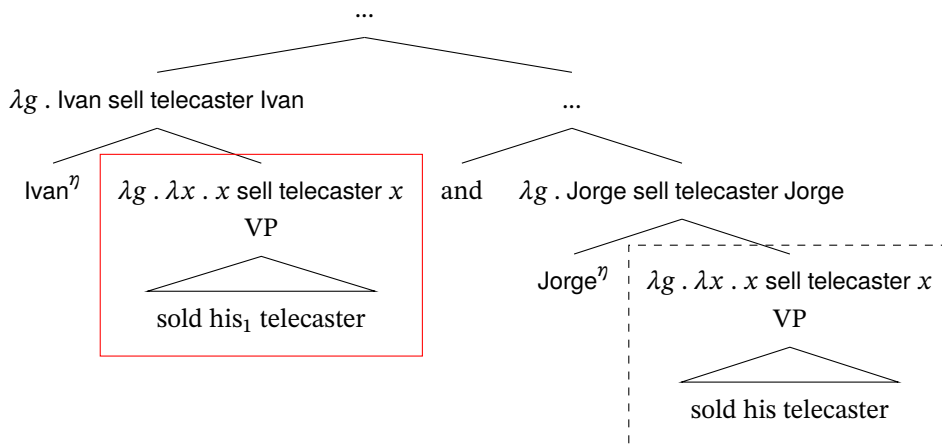
Following Sag (1976) and Williams (1977), positing a VP-internal Λ -operator allows us to capture the strict/sloppy ambiguity as an ambiguity in the *antecedent*.

- With an *accidental coreference* antecedent, we get the *strict* reading.
- With a *binding* antecedent, we get the *sloppy* reading.

Strict reading:



Sloppy reading:



- In the next section, we'll introduce a new dimension of meaning – as well as dealing with *assignment-sensitive* meanings, we're going to be dealing with *alternative-sensitive* meanings.

References

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