General entanglement-assisted transformation for bipartite pure quantum states

Wei Song,[1] Yan Huang,[2] Nai-Le Liu,[1] and Zeng-Bing Chen[1, 3]

[1]Hefei National Laboratory for Physical Sciences at Microscale and Department of Modern Physics,
University of Science and Technology of China, Hefei, Anhui 230026, China
[2]Department of Computer Science and Technology,
University of Science and Technology of China, Hefei, Anhui 230026, China
[3]Physikalisches Institut, Universität Heidelberg, Philosophenweg 12, D-69120 Heidelberg, Germany

We introduce the *general catalysts* for pure entanglement transformations under local operations and classical communications in such a way that we disregard the profit and loss of entanglement of the catalysts *per se*. As such, the possibilities of pure entanglement transformations are greatly expanded. We also design an efficient algorithm to detect whether a $k \times k$ general catalyst exists for a given entanglement transformation. This algorithm can as well be exploited to witness the existence of standard catalysts.

Entanglement plays a central role in quantum information processing (QIP) tasks, such as quantum communication [1], quantum superdense coding [2] and quantum computation [3]. With the development of quantum information science, people have realized that quantum entanglement is a kind of physical resource in nature, like energy. To implement certain QIP tasks, measuring, manipulating and purifying entanglement [4] by local operations and classical communications (LOCC) are unavoidable. An important problem concerns the entanglement transformation between bipartite states under LOCC. This problem arises as a consequence of the fundamental question of how we can convert one type of physical resource into another. There have been considerable efforts devoted to this problem [5, 6, 7, 8, 9, 10, 11, 12, 13, 14]. Bennett *et al.* made a first step on this problem [5] and proposed an entanglement concentration protocol which solves the entanglement transformation problem in the asymptotic case. Another significant advance for finite cases was made by Nielsen [6], who connected the entanglement transformation with the theory of majorization [15, 16] in mathematics. Let $|\psi\rangle = \sum_{i=1}^{n} \sqrt{\alpha_i} |i\rangle |i\rangle$ and $|\varphi\rangle = \sum_{i=1}^{n} \sqrt{\beta_i} |i\rangle |i\rangle$ be pure bipartite states with ordered Schmidt coefficient (OSC) vectors $\psi = (\alpha_1, ..., \alpha_n)$ and $\varphi = (\beta_1, ..., \beta_n)$, where $\alpha_1 \geq \cdots \geq \alpha_n \geq 0$ and $\beta_1 \geq \cdots \geq \beta_n \geq 0$. Then there exists a transformation that converts $|\psi\rangle$ to $|\varphi\rangle$ with 100% probability under LOCC iff $\psi \prec \varphi$, where $\prec$ denotes a majorization relation, namely, for $1 \leq l \leq n$

$$\sum_{i=1}^{l} \alpha_i \leq \sum_{i=1}^{l} \beta_i. \tag{1}$$

Nielsen's theorem provides us with a convenient tool for investigating entanglement transformation. Shortly after Nielsen's work, a surprising phenomenon discovered by Jonathan and Plenio [8] is that sometimes an extra entangled state can allow otherwise impossible entanglement transformation to become realizable. The extra state acts just like a catalyst in a chemical reaction, remaining what it was before the transformation.

A simple example introduced by Jonathan and Plenio is that $|\psi\rangle \nrightarrow |\varphi\rangle$ but $|\psi\rangle \otimes |\chi\rangle \rightarrow |\varphi\rangle \otimes |\chi\rangle$, where $|\psi\rangle = (0.4, 0.4, 0.1, 0.1)$, $|\varphi\rangle = (0.5, 0.25, 0.25, 0)$, and $|\chi\rangle = (0.6, 0.4)$. Here, we have used an OSC vector to represent a bipartite state. But the state $|\chi\rangle$ cannot always act as an assistant in the above way, e.g., it is not capable of catalyzing the transformation $|\psi\rangle \rightarrow |\varphi'\rangle$ where $|\varphi'\rangle = (0.48, 0.27, 0.25, 0)$.

However, if we allow some entanglement of the catalyst state $|\chi\rangle$ to be consumed, then we will greatly improve the possibilities of entanglement transformations. For example, if we choose $|\chi'\rangle = (2/3, 1/3)$, then, by using Nielsen's theorem, one can easily verify that the transformation $|\psi\rangle \otimes |\chi\rangle \rightarrow |\varphi'\rangle \otimes |\chi'\rangle$ can be realized under LOCC. During this process, some entanglement of $|\chi\rangle$ is consumed, i.e., $E(|\chi\rangle) > E(|\chi'\rangle)$, where E is the von Neumann entanglement entropy.

Formally, suppose that $|\psi\rangle$ and $|\varphi\rangle$ are bipartite pure states with $|\psi\rangle \nrightarrow |\varphi\rangle$ under LOCC, and that another auxiliary bipartite pure entangled state $|\chi\rangle$ is standby. If there exists $|\chi'\rangle$ such that $|\psi\rangle \otimes |\chi\rangle \rightarrow |\varphi\rangle \otimes |\chi'\rangle$, we call $|\chi\rangle$ a *general catalyst* for the entanglement transformation from $|\psi\rangle$ to $|\varphi\rangle$, in the sense that we do not care whether the entanglement of catalyst $|\chi\rangle$ is reduced or increased during the process. When the entanglement of $|\chi'\rangle$ keeps the same as that of $|\chi\rangle$, i.e., $E(|\chi\rangle) = E(|\chi'\rangle)$, the catalyst $|\chi\rangle$ reduces to the standard one defined by Jonathan and Plenio [8]; if $E(|\chi\rangle) < E(|\chi'\rangle)$, the catalyst $|\chi\rangle$ becomes the so called *supercatalyst* introduced by Bandyopadhyay *et al.* [11]; if $E(|\chi\rangle) > E(|\chi'\rangle)$, some entanglement of the catalyst is consumed, which is illustrated by the above example; in this case we term $|\chi\rangle$ a *subcatalyst*.

An important question arises naturally: Given bipartite pure states $|\psi\rangle$ and $|\varphi\rangle$ with $|\psi\rangle \nrightarrow |\varphi\rangle$ under LOCC, what states can be general catalysts for the above entanglement transformation? Another question concerns how could we decide whether or not a $k \times k$ general catalyst exists for certain entanglement transformation. On the other hand, as entanglement is a very scarce resource (because it cannot be generated by local means and will

unavoidably be degraded by decoherence when transmitted in a noisy environment), we hope that the entanglement of the catalyst consumed during the entanglement-assisted transformation is as little as possible. This evokes us to investigate the properties of general catalysts, which will be the first part of this paper. In the following we shall start with the simplest cases, i.e., entanglement transformations between 2×2 bipartite pure states.

Consider 2×2 bipartite pure states $|\psi\rangle^\downarrow = (\alpha_1, \alpha_2)$ and $|\varphi\rangle^\downarrow = (\beta_1, \beta_2)$, where we have used $|\psi\rangle^\downarrow$ to denote a state with Schmidt coefficient vectors being sorted in a nonincresing order, and $|\psi\rangle \nrightarrow |\varphi\rangle$ under LOCC. Assume that we are provided with another 2×2 entangled bipartite pure state $|\chi\rangle^\downarrow = (x, 1-x)$, where $0.5 \le x \le 1$. The following theorem provides a *sufficient and necessary* condition for $|\chi\rangle$ to be a general catalyst for realizing the transformation of $|\psi\rangle$ to $|\varphi\rangle$:

Theorem 1. The above pure state $|\chi\rangle$ is a general catalyst for the transformation of $|\psi\rangle$ to $|\varphi\rangle$ iff $x \le \beta_1/\alpha_1$.

Proof: Assume that there exists $|\chi'\rangle^\downarrow = (x', 1-x')$ such that $|\psi\rangle \otimes |\chi\rangle \to |\varphi\rangle \otimes |\chi'\rangle$. Using Nielsen's theorem we have $x\alpha_1 \le x'\beta_1$, and so $x \le x'\beta_1/\alpha_1 \le \beta_1/\alpha_1$. Conversely, we assume that $x \le \beta_1/\alpha_1$. Then we shall show that $|\chi\rangle$ is a general catalyst. Notice first that $\alpha_1 > \beta_1$, because $|\psi\rangle \nrightarrow |\varphi\rangle$. For convenience we divide the problem into two cases.

Case 1: $x \ge \alpha_1$. In this case we can sort the Schmidt coefficients of $|\psi\rangle \otimes |\chi\rangle$ in a nonincreasing order

$$(|\psi\rangle \otimes |\chi\rangle)^\downarrow = (\alpha_1 x, \alpha_2 x, \alpha_1 (1-x), \alpha_2 (1-x)). \quad (2)$$

If $x' > \beta_1$, then

$$(|\varphi\rangle \otimes |\chi'\rangle)^\downarrow = (\beta_1 x', \beta_2 x', \beta_1 (1-x'), \beta_2 (1-x')), \quad (3)$$

and Nielsen's theorem imposes the following inequalities

$$\begin{cases} \alpha_1 x \le \beta_1 x', \\ x \le x', \\ \alpha_1 + \alpha_2 x \le \beta_1 + \beta_2 x'. \end{cases} \quad (4)$$

These inequalities will be satisfied as long as $x' \ge \max\left\{\frac{\alpha_1}{\beta_1}x, 1-\frac{\alpha_2}{\beta_2}(1-x)\right\}$. If otherwise, i.e., $x' \le \beta_1$, then the second inequality in (4) should be replaced by $x \le \beta_1$ which, together with the assumption $x \ge \alpha_1$, contradicts with the requirement $\alpha_1 > \beta_1$.

Case 2: $x < \alpha_1$. If $x' > \beta_1$, then the inequalities imposed by Nielsen's theorem are

$$\begin{cases} \alpha_1 x \le \beta_1 x', \\ \alpha_1 \le x', \\ \alpha_1 + \alpha_2 x \le \beta_1 + \beta_2 x', \end{cases} \quad (5)$$

which hold as long as $x' \ge \max\left\{\frac{\alpha_1}{\beta_1}x, \alpha_1, 1-\frac{\alpha_2}{\beta_2}(1-x)\right\}$. If $x' \le \beta_1$, then the second inequality in (5) will be replaced by $\alpha_1 \le \beta_1$ which is a contradiction to the premise $\alpha_1 > \beta_1$. ■

Remark. $|\chi\rangle$ is always a subcatalyst since $x' > x$. If we choose the lower bound of x', then we could get an optimal $|\chi'\rangle$, in that there will be a minimum loss of entanglement of $|\chi'\rangle$. An extreme case on the other side is where the entanglement of the auxiliary state is completely consumed. Indeed, any $k \times k$ bipartite pure state $|\chi\rangle^\downarrow = (x_1, ..., x_k)$ is a general catalyst for transforming $|\psi\rangle$ to $|\varphi\rangle$ iff $x_1 \le \beta_1/\alpha_1$. This can be shown by simply using Nielsen's theorem and putting $|\chi'\rangle^\downarrow = (1)$, i.e., a separable state. If x_1 is strictly smaller than β_1/α_1, we can always find sufficiently small ε such that $|\psi\rangle \otimes |\chi\rangle \to |\varphi\rangle \otimes (1-\varepsilon, \varepsilon)$, thereby the auxiliary state will not be consumed completely. ■

Next, we consider 3×3 cases.

Theorem 2. Let 3×3 bipartite pure states $|\psi\rangle^\downarrow = (\alpha_1, \alpha_2, \alpha_3)$ and $|\varphi\rangle^\downarrow = (\beta_1, \beta_2, \beta_3)$ be incomparable, i.e., $|\psi\rangle \nleftrightarrow |\varphi\rangle$ under LOCC. If $x_1 \le \min\{\beta_1/\alpha_1, (\beta_1+\beta_2)/(\alpha_1+\alpha_2)\}$, then an arbitrary $k \times k$ bipartite pure state $|\chi\rangle^\downarrow = (x_1, ..., x_k)$ is a general catalyst for the transformation of $|\psi\rangle$ to $|\varphi\rangle$.

Proof: Suppose the entanglement of $|\chi\rangle$ is completely lost after the transformation. Then it suffices to consider the first two Schmidt coefficients of $|\psi\rangle \otimes |\chi\rangle$. Two separate cases should be considered in turn.

Case 1: $x_1 \le x_2 \alpha_1/\alpha_2$. In this case, the two largest Schmidt coefficients of $|\psi\rangle \otimes |\chi\rangle$ are $\alpha_1 x_1$ and $\alpha_1 x_2$. If there exists $|\chi\rangle$ such that $|\psi\rangle \otimes |\chi\rangle \to |\varphi\rangle \otimes |\chi'\rangle$ with $|\chi'\rangle$ being a separable state, then Nielsen's theorem imposes the conditions that $x_1 \le \beta_1/\alpha_1$ and $x_1 + x_2 \le (\beta_1 + \beta_2)/\alpha_1$. On the other hand, since $|\psi\rangle \nleftrightarrow |\varphi\rangle$, it follows from Nielsen's theorem that one of the following two possibilities must hold: either

$$\begin{cases} \alpha_1 > \beta_1 \\ \alpha_1 + \alpha_2 < \beta_1 + \beta_2 \end{cases} \quad (6a)$$

or

$$\begin{cases} \alpha_1 < \beta_1 \\ \alpha_1 + \alpha_2 > \beta_1 + \beta_2. \end{cases} \quad (6b)$$

In both cases, we have $\beta_1 + \beta_2 > \alpha_1$. Hence, the condition $x_1 + x_2 \le (\beta_1 + \beta_2)/\alpha_1$ always holds and so can be neglected.

Case 2: $x_1 > x_2 \alpha_1/\alpha_2$. By a similar procedure we can verify that, in order for $|\chi\rangle$ to be a general catalyst, x_1 must satisfy $x_1 \le \min\left\{\frac{\beta_1}{\alpha_1}, \frac{\beta_1+\beta_2}{\alpha_1+\alpha_2}\right\}$. If the above inequality is strict for all x_1, then there exist cases where the entanglement of $|\chi'\rangle$ is larger than zero. ■

The following Theorem 3, Theorem 4, and Example 1 and Example 2 show that sometimes the only possible choice is to use subcatalysts. This captures what we have emphasized that general catalysts greatly expand the possibilities of entanglement transformations.

Theorem 3. Let $|\psi\rangle$ and $|\varphi\rangle$ be incomparable states, where $|\psi\rangle^\downarrow = (\alpha_1, ..., \alpha_n)$ and $|\varphi\rangle^\downarrow = (\beta_1, ..., \beta_n)$. Suppose a 2×2 or 3×3 state $|\chi\rangle$ is a catalyst for the transformation of $|\psi\rangle$ to $|\varphi\rangle$. If $\alpha_1 > \beta_1$ and $\alpha_n < \beta_n$, then $|\chi\rangle$ must be a subcatalyst.

Proof: First, suppose $|\chi\rangle^{\downarrow} = (x, 1-x)$ and $|\chi'\rangle^{\downarrow} = (x', 1-x')$. By Nielsen's theorem we have $\alpha_1 x \leq \beta_1 x'$ which, together with the condition $\alpha_1 > \beta_1$, implies $x < x'$. Second, suppose $|\chi\rangle^{\downarrow} = (x_1, x_2, x_3)$ and $|\chi'\rangle^{\downarrow} = (x'_1, x'_2, x'_3)$. Using Nielsen's theorem we have $\alpha_1 x_1 \leq \beta_1 x'_1$ and $\alpha_n x_3 \geq \beta_n x'_3$ which, together with $\alpha_1 > \beta_1$ and $\alpha_n < \beta_n$, imply $x_1 < x'_1$ and $x_1 + x_2 < x'_1 + x'_2$. Consequently, we obtain $\chi \prec \chi'$. ∎

Theorem 4. Let $|\psi\rangle$ and $|\varphi\rangle$ be $2 \times n$-level states with $|\psi\rangle \nrightarrow |\varphi\rangle$, then there does not exist any standard catalyst or supercatalyst for the transformation of $|\psi\rangle$ to $|\varphi\rangle$.

Proof: Suppose there exists a catalyst $|\chi\rangle$ such that $|\psi\rangle \otimes |\chi\rangle \to |\varphi\rangle \otimes |\chi'\rangle$, with $E(|\chi\rangle) \leq E(|\chi'\rangle)$. Then we have $E(|\psi\rangle) + E(|\chi\rangle) \geq E(|\varphi\rangle) + E(|\chi'\rangle)$, and so $E(|\psi\rangle) \geq E(|\varphi\rangle)$. Recalling that there is an equivalence between $|\psi\rangle \to |\varphi\rangle$ and $E(|\psi\rangle) \geq E(|\varphi\rangle)$ for $2 \times n$-level states [6], we obtain $|\psi\rangle \to |\varphi\rangle$, which is a contradiction. ∎

Note that this theorem is compatible with Theorem 1.

Example 1. When $|\psi\rangle$ has fewer Schmidt coefficients than $|\varphi\rangle$, by using Nielsen's theorem it is evident that no standard catalyst exists for transforming $|\psi\rangle$ to $|\varphi\rangle$. However, in some situations the transformation may be realized by using a subcatalyst. Suppose $|\psi\rangle^{\downarrow} = (\alpha_1, \alpha_2)$, $|\varphi\rangle^{\downarrow} = (\beta_1, \beta_2, \beta_3)$, and $|\chi\rangle = (x, 1-x)$. To implement the transformation of $|\psi\rangle$ to $|\varphi\rangle$, it is obvious that the entanglement of $|\chi\rangle$ should be consumed completely. If $x \geq \alpha_1$, then the condition arising from Nielsen's theorem reads $x \leq \min\{\beta_1/\alpha_1, \beta_1 + \beta_2\}$; if $x < \alpha_1$, then the condition reads $x \leq \beta_1/\alpha$, $\alpha_1 \leq \beta_1 + \beta_2$. We conclude that under the condition $\alpha_1 \leq \beta_1 + \beta_2$, the pure state $|\chi\rangle = (x, 1-x)$ with $x \leq \{\beta_1/\alpha_1, \beta_1 + \beta_2\}$ is a subcatalyst for transforming $|\psi\rangle$ to $|\varphi\rangle$.

Example 2: Let $|\psi\rangle = \left(\frac{1}{3}, \frac{1}{3}, \frac{1}{6}, \frac{1}{6}\right)$, and $|\varphi\rangle = \left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{12}, \frac{1}{12}, \frac{1}{12}, \frac{1}{12}\right)$, we are provided with another auxiliary 4×4 entangled bipartite state $|\chi\rangle = \left(\frac{1}{4}, \frac{1}{4}, \frac{1}{4}, \frac{1}{4}\right)$ as catalyst. We could find an optimal state $|\chi'\rangle = \left(\frac{1}{2}, \frac{1}{2}\right)$ (i.e., the entanglement of the subcatalyst state $|\chi\rangle$ consumed during the transformation reach a minimum value) such that the transformation $|\psi\rangle \otimes |\chi\rangle \to |\varphi\rangle \otimes |\chi'\rangle$ is possible. Furthermore, it is easy to show that $|\psi\rangle \otimes |\chi\rangle \leftrightarrow |\varphi\rangle \otimes |\chi'\rangle$, since the Schmidt coefficients of $|\psi\rangle \otimes |\chi\rangle$ and $|\varphi\rangle \otimes |\chi'\rangle$ are the same. It means that we could also transform the state $|\varphi\rangle \otimes |\chi'\rangle$ to $|\psi\rangle \otimes |\chi\rangle$. Here, we call state $|\chi\rangle$ as a time-reverse subcatalyst in the above entanglement transformation process.

Next, we consider an interesting question. Let $\{|\psi\rangle, |\varphi\rangle\}$ and $\{|\chi\rangle, |\chi'\rangle\}$ be two incomparable state pairs. Can they assist each other mutually so as to realize the transformation $|\psi\rangle \otimes |\chi\rangle \to |\varphi\rangle \otimes |\chi'\rangle$ by LOCC? We shall demonstrate that this can be the case in some situations.

Example 3. Consider two incomparable state pairs $\{|\psi\rangle, |\varphi\rangle\}$ and $\{|\chi\rangle, |\chi'\rangle\}$, where $|\psi\rangle^{\downarrow} = (\alpha_1, \alpha_2, \alpha_3)$, $|\varphi\rangle^{\downarrow} = (\beta_1, \beta_2, \beta_3)$, $|\chi\rangle^{\downarrow} = (x_1, x_2, x_3)$, and $|\chi'\rangle^{\downarrow} = (x'_1, x'_2, x'_3)$. Suppose that

$$\begin{cases} \alpha_2 x_1 \geq \alpha_1 x_2 \geq \alpha_3 x_1 \geq \alpha_2 x_2, \alpha_3 x_2 \geq \alpha_1 x_3 \\ \beta_2 x'_1 \geq \beta_1 x'_2, \beta_2 x'_2 \geq \beta_1 x'_3, \beta_2 x'_3 \geq \beta_3 x'_1 \end{cases} \quad (7)$$

Then we can sort the Schmidt coeffcients of $|\psi\rangle \otimes |\chi\rangle$ and $|\varphi\rangle \otimes |\chi'\rangle$ in a nonincreasing order

$$(|\psi\rangle \otimes |\chi\rangle)^{\downarrow} = (\alpha_1 x_1, \alpha_2 x_1, \alpha_1 x_2, \alpha_3 x_1, \alpha_2 x_2, \alpha_3 x_2, \alpha_1 x_3, \\ \alpha_2 x_3, \alpha_3 x_3), \quad (8)$$

$$(|\varphi\rangle \otimes |\chi'\rangle)^{\downarrow} = (\beta_1 x'_1, \beta_2 x'_1, \beta_1 x'_2, \beta_2 x'_2, \beta_1 x'_3, \beta_2 x'_3, \beta_3 x'_1, \\ \beta_3 x'_2, \beta_3 x'_3). \quad (9)$$

Since $|\psi\rangle \nleftrightarrow |\varphi\rangle$, either the set of inequalities in Eq. (6a) or that in Eq. (6b) is satisfied. To be specific, we assume the former. Accordingly, in view of the fact that $|\chi\rangle \nleftrightarrow |\chi'\rangle$, in order for the desired entanglement transformation to be realizable, the following inequalities must be satisfied:

$$\begin{cases} x_1 + x_2 > x'_1 + x'_2, \\ \alpha_1 x_1 \leq \beta_1 x'_1, \\ (\alpha_1 + \alpha_2) x_1 \leq (\beta_1 + \beta_2) x'_1, \\ (\alpha_1 + \alpha_2) x_1 + \alpha_1 x_2 \leq (\beta_1 + \beta_2) x'_1 + \beta_1 x'_2, \\ x_1 + \alpha_1 x_2 \leq (\beta_1 + \beta_2)(x'_1 + x'_2), \\ x_1 + (\alpha_1 + \alpha_2) x_2 \leq \beta_1 + \beta_2 (x'_1 + x'_2), \\ x_1 + x_2 \leq \beta_1 + \beta_2, \\ x_1 + x_2 + \alpha_1 x_3 \leq \beta_1 + \beta_2 + \beta_3 x'_1, \\ x_1 + x_2 + (\alpha_1 + \alpha_2) x_3 \leq \beta_1 + \beta_2 + \beta_3 (x'_1 + x'_2). \end{cases} \quad (10)$$

To show these inequalities can be satisfied simultaneously, we choose

$$\begin{aligned} \alpha_1 &= 0.5, \quad \alpha_2 = 0.26, \quad \alpha_3 = 0.24, \\ \beta_1 &= 0.49, \quad \beta_2 = 0.48, \quad \beta_3 = 0.03, \\ x_1 &= 0.62, \quad x_2 = 0.3, \quad x_3 = 0.08. \end{aligned} \quad (11)$$

Then, the set of inequalities in Eq. (7) and Eq. (10) are equivalent to the following set:

$$\begin{cases} x'_1 \geq 31/49, \\ 0.97 x'_1 + 0.49 x'_2 \geq 0.6212, \\ 0.97 (x'_1 + x'_2) \geq 0.77, \\ 0.48 x'_1 \geq 0.49 x'_2, \\ 0.49 x'_1 + 0.97 x'_2 \geq 0.49, \\ 17 x'_1 + 16 x'_2 \leq 16, \\ x'_1 + x'_2 < 0.92. \end{cases} \quad (12)$$

We can picture the region of the independent parameters x'_1 and x'_2 which satisfy the above inequalities simultaneously in a diagram, Fig. 1. We find the region nonempty, the desired result. For example, we may choose $x'_1 = 0.81, x'_2 = 0.1, x'_3 = 0.09$. It should be noted that what $|\psi\rangle$ and $|\chi\rangle$ act as in this process are subcatalysts.

On the other hand, if $|\chi\rangle \nrightarrow |\chi'\rangle$ and $|\varphi\rangle$ is a maximally entangled state, then for any state $|\psi\rangle$ we have $|\psi\rangle \otimes |\chi\rangle \nrightarrow |\varphi\rangle \otimes |\chi'\rangle$, since otherwise we have $|\varphi\rangle \otimes |\chi\rangle \to |\psi\rangle \otimes |\chi\rangle \to$
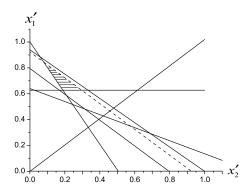
FIG. 1: For two incomparable state pairs $\{|\psi\rangle, |\varphi\rangle\}$ and $\{|\chi\rangle, |\chi'\rangle\}$ with $|\psi\rangle = (0.5, 0.26, 0.24)$, $|\varphi\rangle = (0.49, 0.48, 0.03)$, $|\chi\rangle = (0.62, 0.3, 0.08)$, and $|\chi'\rangle = (x_1', x_2', 1 - x_1' - x_2')$, if the transformation $|\psi\rangle \otimes |\chi\rangle \rightarrow |\varphi\rangle \otimes |\chi'\rangle$ is feasible, then the parameters x_1' and x_2' should lie in the shadow region.
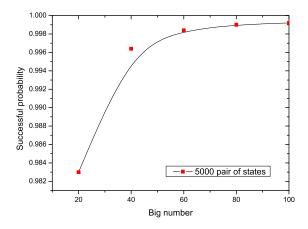

FIG. 2: The numerical results for the successful probability as a function of the big number we choose. It increases with the big number and reaches 99.92% when the big number is 100.

$|\varphi\rangle \otimes |\chi'\rangle$. But a maximally entangled state cannot act as a standard catalyst.

In the remainder of this paper, we will consider the following problem. Assume $|\psi\rangle^{\downarrow} = (\alpha_1, ..., \alpha_n)$ and $|\varphi\rangle^{\downarrow} = (\beta_1, ..., \beta_n)$ with $|\psi\rangle \nrightarrow |\varphi\rangle$ under LOCC. How could we decide whether or not a $k \times k$ general catalyst exists for converting $|\psi\rangle$ into $|\varphi\rangle$? Notice that it suffices to consider whether the process $|\psi\rangle \otimes |X\rangle \rightarrow |\varphi\rangle$ is possible, where $|X\rangle$ is a $k \times k$ maximally entangled state. For $k \geq n$, the above process is always possible, because a $k \times k$ maximally entangled state can always be transformed into any $k \times k$ entangled state under LOCC; for $k < n$, we only need to check whether the majorization relation $(|\psi\rangle \otimes |X\rangle)^{\downarrow} \prec |\varphi\rangle^{\downarrow}$ holds. It can be implemented by checking whether the $n - 1$ inequalities are satisfied. However, this method cannot be applied to decide the existence of $k \times k$ standard catalysts for certain entanglement transformation. We will propose a Monte Carlo algorithm to solve this problem. Firstly, we generate a group of $x_1 \geq \cdots \geq x_k \geq 0$ randomly which satisfy $\sum_{i=1}^{k} x_i = 1$. Then, we merge sort the Schmidt coefficients of $|\psi\rangle \otimes |\chi\rangle$ and $|\varphi\rangle \otimes |\chi\rangle$ in a nonincreasing order where $|\chi\rangle = (x_1, ..., x_k)$. Now the aim is to check whether the majorization relation $(|\psi\rangle \otimes |\chi\rangle)^{\downarrow} \prec (|\varphi\rangle \otimes |\chi\rangle)^{\downarrow}$ holds. After running this procedure M times, if we still cannot find the state $|\chi\rangle$ such that the above majorization relation holds, we say there does not exist a standard catalyst for this transformation. Of course, there is a failure probability when the algorithm gives a false output. But when the big number M is large enough, the successful probability of our algorithm will approach 1. The detailed description is as follows:

(i) For $i = 0$ to k do, $x_i \leftarrow rand[0, 1]$, where $x_1 \geq \cdots \geq x_k \geq 0$, and $\sum_{i=1}^{k} x_i = 1$

(ii) set count= 0

(iii) while count<BIGNUMBER

(iv) begin merge sort the Schmidt coefficients of $|\psi\rangle \otimes |\chi\rangle$ and $|\varphi\rangle \otimes |\chi\rangle$ in a nonincreasing order, respectively

(v) if there exists $|\chi\rangle = (x_1, ..., x_k)$ satisfying $(|\psi\rangle \otimes |\chi\rangle)^{\downarrow} \prec (|\varphi\rangle \otimes |\chi\rangle)^{\downarrow}$

(vi) then a $k \times k$ standard catalyst exists for this transformation

(vii) return success

(viii) else count=count+1 end begin

(xi) return failure

In Ref. [14] X. M. Sun et al also proposed a deterministic algorithm which runs in $O\left(n^{2k+3.5}\right)$ time. However, if k is a variable, employing their algorithm to determine the existence of standard catalyst will become a NP-hard problem. Suppose we choose the big number to be M, then it is easy to see that our algorithm runs in $O(Mnk)$ time, which is greatly improved than the deterministic one.

To show the effectiveness of this algorithm, we will give some examples in the following. We devise a program to generate 5000 pair of 8×8 states $\{|\psi\rangle, |\varphi\rangle\}$ which always have 4×4 standard catalysts, where $|\psi\rangle \nrightarrow |\varphi\rangle$. We run the above algorithm and find that, when the big number is chosen to be 100, the successful probability is 99.92%. We plot the result in Fig. 2.

In summary, by introducing the concept of the general catalyst, we can greatly expand the possibilities of entanglement-assisted transformations between pure entangled states. We consider the problem of how to decide the existence of a $k \times k$ general catalyst for certain entanglement transformation. We also propose a Monte Carlo algorithm for determining the existence of the standard catalyst. When the dimensions of the state and of the

potential catalyst are both very big numbers, our algorithm is far more efficient than the previous deterministic algorithm. We believe our results may have potential applications in future manipulations of quantum entanglement.

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