

Neither contextuality nor non-locality admits catalysts

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Abstract

We show that the resource theory of contextuality does not admit catalysts. As a corollary, we observe that the same holds for non-locality. This adds a further example to the list of “anomalies of entanglement”, showing that non-locality and entanglement behave differently as resources. We also show that catalysis remains impossible even if instead of classical randomness we allow some more powerful behaviors to be used freely in the free transformations of the resource theory.

Introduction. Contextuality [1] and non-locality [2, 3] play a prominent role in a wide variety of applications of quantum mechanics, with non-locality being used for example in quantum key-distribution [4], certified randomness [5] and randomness expansion [6]. Similarly, contextuality powers quantum computation in some computational models [7–12] and even increases expressive power in quantum machine learning [13]. Consequently, it is vital to understand how non-locality and contextuality behave as resources.

In this Letter, we show that neither contextuality nor non-locality admit catalysts: that is, there are no correlations that can be used to enable an otherwise impossible conversion between correlations and still recovered afterwards. To state this slightly more precisely, let us write $d, e, f \dots$ for various correlations (whether classical or not), $d \otimes e$ for having independent instances of d and e , and $d \rightsquigarrow e$ (read as “ d simulates e ”) for the existence of a conversion $d \rightarrow e$. Then our results state that, in suitably formalized resource theories of contextuality and non-locality, whenever $d \otimes e \rightsquigarrow d \otimes f$, then $e \rightsquigarrow f$ already. This gives a strong indication that contextuality (and non-locality) are resources that **get spent** when you use them: there is no way of using a correlation d to achieve a task you could not do otherwise while keeping d intact. As entanglement theory famously allows for catalysts [14], this can be seen as yet another “anomaly of non-locality” [15] and thus further testament to the fact that non-locality and entanglement are different resources.[16]

We do this by working in precisely defined resource theories of contextuality and non-locality. These are not strictly speaking quantum resource theories [17], but resource theories in a more general sense [18, 19], as we allow resources such as PR-boxes [20] that are not quantum realizable. The kinds of conversions between correlations we have in mind capture the intuitive idea of using one system to **simulate** another one, and have been studied in earlier literature [21–29]. These roughly correspond to the local operations and shared randomness-paradigm ((LOSR)) or to wirings and prior-to-input classical communication

(WPICC), depending on the precise definitions of these terms. However, existing formalizations of these in the literature are often limited to the bipartite or tripartite settings, and at times overlook some technical issues resulting in non-convex sets of transformations [30, Appendix].

More importantly, existing formalizations of the resource theory of non-locality tend to focus on the case where each party has a discrete set of measurements, of which they can perform at most one. However, this is false even in relatively simple situations: for instance, if Alice shares one PR-box with Bob and one with Charlie, then she has four measurements available, but is not restricted to only one measurement as she can choose a measurement for each of her boxes. In particular, she might first measure one of the boxes, and use the outcome (and possible auxiliary randomness) to choose what to measure next.

To overcome such issues, we work in the general approach to contextuality initiated in [31], and later extended in [32] to capture building some correlations from others using such **probabilistic** and **adaptive** means, resulting in a resource theory for contextuality. We then obtain the resource theory of non-locality from this via a general mathematical construction [33] that builds a resource theory of n -partite resources from a given resource theory. Working with such generality clarifies the relationship between resource theories of contextuality and non-locality, and captures the kinds of interconversions studied in earlier literature in the exact, single-shot regime [34] in a precise yet tractable manner.

We believe that our result could be phrased and proved in terms other approaches to contextuality [35–39], as long as one formalizes such adaptive measurement protocols and transformations between correlations within them. However, if one works with axiomatically rather than operationally defined transformations as in [40] the proof no longer carries over automatically, as it could happen that operational transformations form a proper subset of axiomatically defined ones, as happens for instance in resource theories of entanglement [41] and magic [42].

No catalysis for contextuality We begin by briefly reviewing the resource theory of contextuality as defined in [32]. To start, we formalize the idea of a **measurement scenario** S : we imagine a situation where there is a finite set X_S of measurements available, each measurement $x \in X_S$ giving rise to outcomes in some finite set $O_{S,x}$. However, only some measurements might be feasible to perform together — other combinations may be ruled out by practical limitations or excluded by physical theory. We collect all jointly compatible

measurements into a single set Σ_S , which we expect to satisfy two natural properties:

1. Any measurement $x \in X_S$ induces a compatible set $\{x\} \in \Sigma_S$
2. Any subset of a compatible set of measurements is compatible

Collecting all this data together results in a **measurement scenario** $S = \langle X_S, \Sigma_S, O_S \rangle$.

Given two scenarios S and T , we let $S \otimes T$ denote the scenario that represents having access to S and T in parallel, so that a joint measurement is possible precisely if its components in S and T are.

An **empirical model** over a scenario S is given by specifying for each compatible $\sigma \in \Sigma_S$ a joint probability distribution e_σ for measurements in σ . We only consider empirical models for which the behavior of a joint measurement does not depend on what is measured with it, if anything. Thus, whenever $\tau \subset \sigma \in \Sigma$, the distribution e_τ over τ can be obtained by marginalizing e_σ to τ . We express this generalization of the usual no-signalling conditions as $e_\sigma|_\tau = e_\tau$, so that more generally for a joint distribution d over outcomes of $Y \subset X_S$ and $Z \subset Y$, the expression $d|_Z$ denotes the marginal distribution on outcomes of Z .

If the impossibility of measuring everything together is only a practical limitation, one can contemplate the distribution d that would arise when measuring X_S . If d explained the model e we have at hand, we would expect it to satisfy $d|_\sigma = e_\sigma$ for every $\sigma \in \Sigma$. If such a distribution exists, we call e **non-contextual**. If no such distribution exists, we call e **contextual**, as we have reason to believe that is infeasible in principle to measure everything together, unless one accepts that the observed joint distribution for a subset $Y \subset X$ depends on what it is measured with i.e. its context.

A simple example of scenario and a contextual empirical model on it, discussed in [43], is given by three measurements, any two of which are compatible but not all three together. Each measurement takes outcomes in $\{0, 1\}$, and whenever two measurements are performed one observes $(0, 1)$ and $(1, 0)$ with equal probability, with the probability distributions for singletons fixed by this and resulting in observing the two outcomes with equal probability. Too see that this is contextual, note that there is no joint outcome for all three measurements that is consistent with the observed marginals.

More examples can be obtained from scenarios studied in non-locality, where one typically specifies a scenario by giving the number of parties, the number of measurements available to each of them and the size of the outcome sets, where it is then understood that maximal

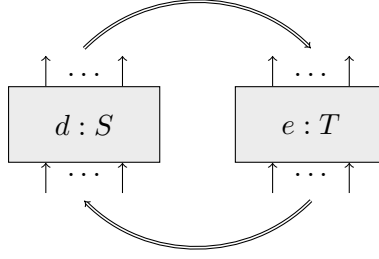


FIG. 1. Depiction of a deterministic simulation $d \rightarrow e$

compatible measurements are given by a choice of a single measurement by each party. This includes for instance the famous CHSH model [44] and the PR-box [20] which goes beyond what is allowed in quantum mechanics, both models arising in the scenario with two parties having access to two dichotomic measurements.

We now move on to transformations between scenarios and empirical models. We will build up to “wirings”, which in full generality capture the idea of **simulating** an empirical model from another adaptively and with the help of non-contextual randomness. These will be the free transformations of our resource theory, but we begin by considering the problem of building one scenario from another. A particularly simple way of building T from S , is by declaring that for each measurement $x \in X_T$ of T some measurement $\pi(x) \in X_S$ of S is to be performed instead. Moreover, each outcome o of $\pi(x)$ is to be interpreted as the outcome $\alpha_x(o)$ instead. If for each compatible $\sigma \in \Sigma_S$ the corresponding measurement $\pi(\sigma)$ is jointly compatible in S , the pair $\langle \pi, \alpha = (\alpha_x)_{x \in X} \rangle$ describes a way of building T from S , and we will denote this by writing $\langle \pi, \alpha \rangle : S \rightarrow T$.

Given such $\langle \pi, \alpha \rangle : S \rightarrow T$, any empirical model $e : S$ vs $e : S$ induces a model $\langle \pi, \alpha \rangle_*(e) : T$, which describes the statistics one would see if one was to observe the statistics given by e and then transform them according to $\langle \pi, \alpha \rangle$. The pair $\langle \pi, \alpha \rangle$ is defined to be a **deterministic simulation** of $e : T$ from $d : S$, denoted by $\langle \pi, \alpha \rangle : d \rightarrow e$, precisely when $\langle \pi, \alpha \rangle$ transforms d to e , i.e. if $\langle \pi, \alpha \rangle_*(d) = e$.

There are two ways in which these deterministic simulations are weaker than one would want in a general resource theory of contextuality: first of all, one might want to allow the usage of auxiliary (non-contextual) randomness, so that the dependence of measurements and their outcomes of T on those of S is stochastic. The usage of auxiliary randomness is captured in [45] and discussed further in [46], and in our terms could be defined by allowing

probabilistic mixtures of transformations. However, this viewpoint leaves out another important generalization—namely the possibility that a single measurement in T can depend on a joint measurement of S as in [47] or more generally on a measurement protocol on S that chooses which joint measurement to perform **adaptively**. The most general formulation of this idea would allow a measurement x in T to be simulated by a probabilistic and **adaptive** procedure, that first measures (depending on some classical randomness) something in S , then based on the outcome (and possibly further classical randomness) chooses what to measure next (if anything) and so on. This idea is formalized carefully in [32] in two stages, the first one adding adaptivity and the second adding randomness.

To model adaptivity, one builds from a scenario S a new scenario $\mathbf{MP}(S)$ of (deterministic) **measurement protocols** over S . A measurement protocol is a procedure that, at any stage, either stops and reports all of the the measurement results obtained so far, or, based on previously seen outcomes, performs a measurement in S that is compatible with all of the previous measurements. The measurements of $\mathbf{MP}(S)$ are given by such protocols over S , and a set of measurement protocols is compatible if they can be performed jointly without having to query measurements outside of Σ_S . Then one can define adaptive (but still deterministic) transformations $S \rightarrow T$ between scenarios as deterministic transformations $\mathbf{MP}(S) \rightarrow T$ as in Figure 2. In [32] we show that the assignment $S \mapsto \mathbf{MP}(S)$ defines a comonad on the category of scenarios: this abstract language is not needed here, but can be thought of as guaranteeing that one has a well-behaved way of composing $\mathbf{MP}(S) \rightarrow T$ with $\mathbf{MP}(T) \rightarrow U$ to obtain a map $\mathbf{MP}(S) \rightarrow U$. Intuitively, the composite is obtained by first interpreting each measurement in U as a measurement protocol over T , and each measurement in that measurement protocol as an measurement protocol over S , and then “flattening” the resulting measurement protocol of measurement protocols (i.e. a measurement of $\mathbf{MP}(\mathbf{MP}((S)))$) into a measurement protocol in S .

To model (non-contextual) randomness, one then defines a **simulation** $d \rightarrow e$ to be a deterministic simulation $\mathbf{MP}(d \otimes c) \rightarrow e$ for some non-contextual model c as in Figure 3. Again, such simulations can be shown to compose in a well-behaved manner [48].

We will denote the existence of such a simulation $d \rightarrow e$ by $d \rightsquigarrow e$, intended to be read as “ d simulates e ”. Simulations as thus defined interact well with contextuality. For instance,

- [32, Theorem 21] states that the non-contextual fraction, studied in [12] is a monotone: that is, if $d \rightsquigarrow e$ then $\text{NCF}(d) \leq \text{NCF}(e)$.

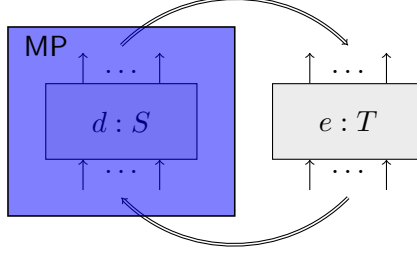


FIG. 2. Depiction of an adaptive simulation $d \rightarrow e$

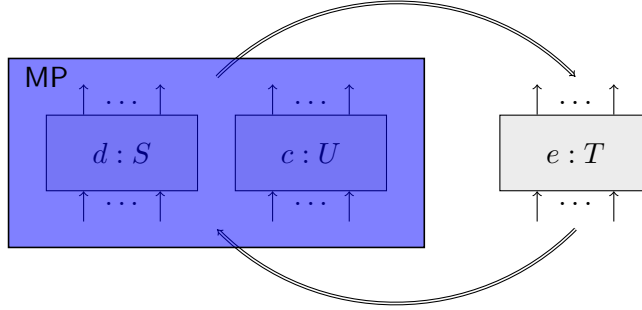


FIG. 3. General simulation $d \rightarrow e$, where we require c to be non-contextual.

- [47, Theorem 4.1] implies that an empirical model is non-contextual iff it can be simulated from a trivial model. In fact, the notions of logical and strong contextuality can be captured along similar lines [49], by relaxing the equality of probability distributions in the definition of simulation by equality (or inclusion) of supports of these distributions.

We can now state our main result for this resource theory of contextuality.

Theorem 1. *If $d \otimes e \rightsquigarrow d \otimes f$ in the resource theory of contextuality, then $e \rightsquigarrow f$.*

The key ideas of the proof are simple, even if the full details get technical: if d can catalyze a transformation $e \rightarrow f$ once, it can do so arbitrarily many times. Choosing a big enough number of copies of e to transform to copies of f , using the pigeonhole principle then ends up guaranteeing that one only needs a compatible subset of S_d (or rather, a compatible set of measurement protocols). Making this precise requires formalizing our framework more carefully, and we do this in the Appendix. Our result subsumes [32, Theorem 22], as setting $d = f$ and letting e be a trivial model implies that if $d \rightsquigarrow d \otimes d$ the model d must be non-contextual.

No-catalysis for non-locality We now explain how to interpret non-locality within this framework. At the level of scenarios and models on them, non-locality can be seen as a special case of contextuality: for non-locality, the measurement scenario typically arises by considering n parties, with the i -th party choosing one measurement from a set X_i of measurements available to them, with a measurement $x \in X_i$ giving outcomes in some outcome set $O_{i,x}$. Often one restricts the situation even further and assumes that each party has the same number of measurements available and each measurement takes outcomes in a set of the same size, so that the scenario is specified (up to isomorphism) by three numbers: the number of parties, the number of measurements available to each of them and the size of the outcome sets. Whether or not one imposes this further restriction, such scenarios are of the form $S = \bigotimes_{i=1}^n S_i$ where each S_i is just a discrete set of measurements (so only singleton measurements are possible in each S_i). In particular, a maximal measurement corresponds to a choice (x_1, \dots, x_n) of a measurement at each site, and a set of correlations can be given as a family $p(o_1, \dots, o_n | x_1, \dots, x_n)$ of conditional probabilities for each such measurement. If the family p is (fully) no-signalling, it corresponds to a unique empirical model $e : S$ (where the probabilities over non-maximal measurements are obtained by marginalization), and p is local iff e is non-contextual.

However, if we allow parties to share different non-local resources, we move away from the situation where each party chooses one measurement from a set of mutually exclusive measurements. For instance if Bob shares one box with Alice and one with Charlie, Bob is not limited to a single measurement: he can choose one for each box. As quantum theory allows for arbitrary joint measurability graphs in the case of PVMs [50] and arbitrary simplicial complexes in the case of POVMs [51], we impose **no restrictions** on the measurement scenarios available at each site. However, if we are thinking about a n -partite scenario, it is reasonable to expect that measurement choices done at one site do not affect the measurements available at another one. Thus we model an n -partite scenario as a tuple $(S_i)_{i=1}^n$ of scenarios, thought of as representing the n parties sharing the scenario $\bigotimes_{i=1}^n S_i$.

At the level of transformations between models and scenarios, non-locality is no longer a special case of contextuality, as observed for example in [30, Appendix A.1]. This is because our wirings are slightly too general as they allow the i -th party to wire some of their measurements to measurements belonging to other parties, whereas operationally speaking it is reasonable to require each party to have access only to measurement available to them

(and shared randomness). Nevertheless, the resource theories are very closely related, as one can obtain the resource theory for n -partite non-locality from that of contextuality via a general construction, used in the context of cryptography in [33], that builds a resource theory of n -partite resources from a given resource theory.

In our setting, this amounts to defining the resource theory of n -partite non-locality as follows: an n -partite scenario is an n -tuple $(S_i)_{i=1}^n$ of scenarios (some of which may be empty), and an n -partite empirical model $e : (S_i)_{i=1}^n$ is an empirical model on $\bigotimes_{i=1}^n S_i$. The parallel composite \boxtimes of scenarios is defined pointwise, i.e. by setting $(S_i)_{i=1}^n \boxtimes (T_i)_{i=1}^n = (S_i \otimes T_i)_{i=1}^n$. To define \boxtimes for two n -partite models $e : (S_i)_{i=1}^n$ and $d : (T_i)_{i=1}^n$, note that the scenarios $(\bigotimes_{i=1}^n S_i) \otimes (\bigotimes_{i=1}^n T_i)$ and $\bigotimes_{i=1}^n (S_i \otimes T_i)$ are canonically isomorphic so that the model $e \otimes d : (\bigotimes_{i=1}^n S_i) \otimes (\bigotimes_{i=1}^n T_i)$ induces a model $e \boxtimes d$ on $\bigotimes_{i=1}^n (S_i \otimes T_i)$ via this isomorphism. Finally, an n -partite simulation $d \rightarrow e$ is defined as an n -tuple of simulations $(\text{MP}(T_i \otimes P_i \rightarrow S_i))$ that, when taken together transform $d \boxtimes c$ to e where c is some non-contextual shared correlation. In this manner the resource theory of n -partite non-locality is derived from that of contextuality by keeping track of the n -partite nature of scenarios and transformations between them.

Our proof of Theorem 1 readily implies that non-locality admits no catalysts.

Theorem 2. *If $d \boxtimes e \rightsquigarrow d \boxtimes f$ in the resource theory of n -partite non-locality, then $e \rightsquigarrow f$.*

Quantum-assisted transformations and beyond We briefly discuss a further generalization of our main result. One could consider resource theories with even more expressive simulations than the one's defined above. In the above, we define simulations as deterministic simulations assisted by non-contextual randomness. As suggested in [32], one could allow more general correlations to be used in transformations, and study e.g. quantum-assisted simulations. More specifically, one could define **quantum-assisted simulations** $e \rightarrow f$ as deterministic simulations $\text{MP}(e \otimes q) \rightarrow f$ where q is a quantum-realizable empirical model. More generally, for any class \mathcal{X} of empirical models that is closed under \otimes one gets a well-defined notion of \mathcal{X} -assisted simulations [33], and we will write $d \rightsquigarrow_{\mathcal{X}} e$ to denote the existence of an \mathcal{X} -assisted simulation from d to e . If the class \mathcal{X} contains all non-local correlations, the resulting resource theory has no catalysts. A similar result holds for \mathcal{X} -assisted n -partite transformations between n -partite correlations, but we contend ourselves to state the theorem for contextuality.

Theorem 3. *Let \mathcal{X} be a class of empirical models that is closed under \otimes and contains all classical correlations. Then $d \otimes e \rightsquigarrow_{\mathcal{X}} d \otimes f$ implies $e \rightsquigarrow_{\mathcal{X}} f$.*

For instance, if we take \mathcal{X} to be the class of quantum-realizable empirical models, this result implies that, if by some miracle we got access to a single PR-box, we couldn't use it to catalyze a quantum-assisted transformation that was hitherto impossible. Put another way, no matter what class of correlations we can use freely, the only way to get mileage out of a box that goes beyond our powers is to spend it—so we must choose wisely.

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APPENDIX

In this Appendix we provide proofs of our main theorems and some technical details concerning the background that is needed in order to formalize these. The background is formalized in section A, and mostly explains the framework of [32]. The proofs take place in Appendix B, where we first prove Theorem 1 in B 1 and then extend the result to prove Theorem 2 in B 2 and Theorem 3 in B 3.

Appendix A: Background

We mostly follow the development of [32] with some notational changes. For a scenario S and subset $Y \subset X_S$ of measurements, we let $\mathcal{E}_S(Y)$ (or $\mathcal{E}(Y)$ when S is clear from context) denote the set of possible joint outcomes for the (not necessarily comeasurable) set Y . Formally, $\mathcal{E}_S(Y)$ is defined as the cartesian product $\prod_{x \in Y} O_{S,x}$.

We now define empirical models more carefully. Given a probability distribution d on $\mathcal{E}_S(Z)$, which we think of as a joint probability distribution for measurements in Z , and a

subset $Y \subset Z$, we denote by $d|_Y$ the marginalization of d to measurements in Y . Now, an empirical model $e : S$ consists of a family $(e_\sigma)_{\sigma \in \Sigma}$, where each e_σ is a probability distribution over the set $\mathcal{E}(\sigma)$, and whenever $\tau \subset \sigma \in \Sigma$ we have

$$e_\sigma|_\tau = e_\tau$$

The **parallel composite** $S \otimes T$ of S and T is defined by $X_{S \otimes T} = X_S \sqcup X_T$, $\Sigma_{S \otimes T} = \{\sigma \sqcup \tau | \sigma \in \Sigma_S, \tau \in \Sigma_T\}$ and setting the outcome set at $x \in X_S \sqcup X_T$ to be $O_{S,x}$ if $x \in X_S$ and $O_{T,x}$ if $x \in X_T$. Here $X \sqcup Y$ denotes the disjoint union of the sets X and Y , and can be defined e.g. as $X \sqcup Y := X \times \{0\} \cup Y \times \{1\}$.

A deterministic procedure $S \rightarrow T$ consists of a

- A simplicial function $\pi : X_T \rightarrow X_S$, i.e. a function π satisfying $\pi(\sigma) \in \Sigma_S$ for every $\sigma \in \Sigma_T$
- A family $\alpha = (\alpha_x : O_{S,\pi(x)} \rightarrow O_{T,x})_{x \in X}$ of functions between outcome sets.

A procedure $\langle \pi, \alpha \rangle : S \rightarrow T$ induces a mapping between empirical models so that each $e : S$ is pushed forward to $\langle \pi, \alpha \rangle_* e : T$. Mathematically, the empirical model $\langle \pi, \alpha \rangle_* e$ is defined at measurement $\sigma \in \Sigma_T$ by

$$\langle \pi, \alpha \rangle_* e)_\sigma = \alpha_*(e_{\pi\sigma})$$

where the right hand side denotes the pushforward of the probability distribution $e_{\pi\sigma}$ along the function $\mathcal{E}_S(\pi(\sigma)) \rightarrow \mathcal{E}_T(\sigma)$ whose x -th coordinate projection is given by $\mathcal{E}_S(\pi(x)) \xrightarrow{\alpha_x} \mathcal{E}_T(x)$.

If one imagines S as corresponding to some particular experimental setup, and a model $e : S$ as corresponding to empirically observed propensities of outcomes for a fixed state preparation, then one can interpret a procedure $\langle \pi, \alpha \rangle : S \rightarrow T$ as a way of (deterministically) building the experimental setup T out of S : the map π tells what measurement in S each measurement in T corresponds to, and α tells how to interpret outcomes in S as outcomes in T . Then the model $\langle \pi, \alpha \rangle_* e$ describes the statistics one would see if one was to observe the statistics given by e but transform them according to $\langle \pi, \alpha \rangle$. In other words, the probability $\langle \pi, \alpha \rangle_* e$ gives to some fixed outcome over the joint measurement $\sigma \in \Sigma_T$ is the sum the probabilities e gives to all outcomes of $\pi(\sigma)$ that are mapped to s by α .

A deterministic simulation $d \rightarrow e$ where $d : S$ and $e : T$ is a deterministic procedure $\langle \pi, \alpha \rangle : S \rightarrow T$ that transforms d to e , i.e. that satisfies $\langle \pi, \alpha \rangle_* e = d$.

We now formalize measurement protocols carefully. These were used in [38, Appendix D] to relate the sheaf-theoretic approach [31] to contextuality to the hypergraph approach [38], and then later used in [32] to extend the former. Intuitively, a (deterministic) measurement protocol is a set of rules that tells at each stage, what to measure next given the previous measurements and their outcomes. At any stage of the measurement protocol, we thus have a sequence $(x_i, o_i)_{i=1}^n$ of measurement-outcome pairs. A technical insight of [32] is that measurement protocols can be defined in terms of sets of such sequences: if you know all such sequences that could happen during a deterministic protocol, you also know the protocol itself.

Definition 4. A **run** on a measurement scenario S is a sequence $\bar{x} := (x_i, o_i)_{i=1}^n$ such that $x_i \in X_S$ are distinct, $\{x_1, \dots, x_n\} \in \Sigma_S$, and each $o_i \in O_{S, x_i}$. We denote the empty run by Λ .

A run \bar{x} determines a context $\sigma_{\bar{x}} := \{x_1, \dots, x_n\} \in \Sigma_S$ and a joint assignment on that context $s_{\bar{x}} : x_i \mapsto o_i \in \mathcal{E}(\sigma_{\bar{x}})$. Two runs \bar{x} and \bar{y} are said to be **consistent** if they agree on common measurements, i.e. for every $z \in \sigma_{\bar{x}} \cap \sigma_{\bar{y}}$ we have $s_{\bar{x}}(z) = s_{\bar{y}}(z)$.

Given runs \bar{x} and \bar{y} , we denote their concatenation by $\bar{x} \cdot \bar{y}$. Note that $\bar{x} \cdot \bar{y}$ might not be a run.

Definition 5. A **measurement protocol** on S is a non-empty set Q of runs satisfying the following conditions:

- (i) if $\bar{x} \cdot \bar{y} \in Q$ then $\bar{x} \in Q$;
- (ii) if $\bar{x} \cdot (x, o) \in Q$, then $\bar{x} \cdot (x, o') \in Q$ for every $o' \in O_{S, x}$;
- (iii) if $\bar{x} \cdot (x, o) \in Q$ and $\bar{x} \cdot (x', o') \in Q$, then $x = x'$.

Identifying a measurement protocol with its set of possible runs, the first condition guarantees that a prefix of a possible run is a possible run. The second condition ensures that, if x might be measured at some stage, then any outcome of x is an (in-principle) possible outcome at that stage, and the third condition guarantees that at any given stage, the next measurement (if any), is uniquely determined.

Definition 6. Given a scenario S , we build a scenario $\text{MP}(S)$ as follows:

- its set of measurements is the set $\mathbf{MP}(S)$ of measurement protocols on S ;
- the outcome set O_Q of a measurement protocol $Q \in \mathbf{MP}(S)$ is its set of maximal runs, i.e. those $\bar{x} \in Q$ that are not a proper prefix of any other $\bar{y} \in Q$;
- a set $\{Q_1, \dots, Q_n\}$ of measurement protocols is compatible whenever for any choice of pairwise consistent runs $\bar{x}_i \in Q_i$ with $i \in \{1, \dots, n\}$, we have $\bigcup_i \sigma_{\bar{x}_i} \in \Sigma$.

Given an empirical model $e : S$, we define the empirical model $\mathbf{MP}(e) : \mathbf{MP}(S)$ as follows. For a compatible set $\sigma := \{Q_1, \dots, Q_n\}$ of measurement protocols and an assignment $s : Q_i \mapsto \bar{x}_i \in \mathcal{E}(\sigma)$, we set

$$\mathbf{MP}(e)_\sigma(s) := \begin{cases} e_{\bigcup_i \sigma_{\bar{x}_i}}(\bigcup_i s_{\bar{x}_i}) & \text{if } \{\bar{x}_i\} \text{ pairwise consistent} \\ 0 & \text{otherwise.} \end{cases}$$

The intuition behind a set of measurement protocols being compatible is that they can be performed together, without ever having to perform a measurement not allowed by Σ_S . We next make this intuition precise.

Definition 7. We say that a measurement protocol P **contains implicitly** a protocol Q , if any outcome of P determines the outcome of Q , and we denote this by $P \succeq Q$. Formally, $P \succeq Q$ if for any maximal run $\bar{x} \in P$ there is a (necessarily unique) maximal run $\bar{y} \in Q$ such that the assignment $s_{\bar{y}}$ determined by \bar{y} is a restriction of $s_{\bar{x}}$. We say that a measurement protocol P contains implicitly a set of protocols $\{Q_1, \dots, Q_n\}$ if $P \succeq Q_i$ for each i .

Lemma 8. *If $\{Q_1, \dots, Q_n\}$ is a compatible set of measurement protocols, there is a measurement protocol P that implicitly contains them all.*

Proof. We define a protocol P that first performs Q_1 , then (whatever is left of) Q_2 and so on. Formally, we first define inductively a merge operation $*$ for compatible runs:

$$\begin{aligned} \bar{x} * \Lambda &:= x \\ \bar{x} * ((y, o) \cdot \bar{y}) &:= \begin{cases} \bar{x} * \bar{y} & \text{if } y \in \sigma_{\bar{x}} \\ (\bar{x} \cdot (y, o)) * \bar{y} & \text{otherwise.} \end{cases} \end{aligned}$$

We extend $*$ to all pairs of runs by setting $\bar{x} * \bar{y} = \Lambda$ whenever \bar{x} and \bar{y} are not compatible.

We then define P by taking the closure of the set

$$\{\bar{x}_1 * \dots * \bar{x}_n \mid \bar{x}_i \in O_{\text{MP}(S), Q_i}\}$$

under prefixes. Now property (i) is true by construction, and properties (ii) and (iii) follow from each Q_i being a measurement protocol. Finally, each outcome of P is of the form $\bar{x}_1 * \dots * \bar{x}_n$, so that $P \succeq Q_i$ for each Q_i . \square

Remark. If P is a measurement protocol over S and we fix the outcomes of some subset $Y \subset X_S$ of measurements to equal $t \in \mathcal{E}(Y)$, this determines a measurement protocol $P(t)$ that does not measure anything in Y as follows: it proceeds exactly as P except that whenever P is supposed to measure some $x \in Y$ it behaves as if $t(x)$ was observed and proceeds accordingly. In particular, if P is a measurement protocol over $S \otimes T$ and t is an element of $\mathcal{E}_T(X_T)$, then $P(t)$ is a measurement protocol over S .

We can describe $P(t)$ more formally as follows. Let us say that a run $\bar{x} \in P$ is compatible with $t \in \mathcal{E}(Y)$ if $s_{\bar{x}}$ agrees with t on common measurements, i.e. if $s_{\bar{x}}(z) = t(z)$ for every $z \in \sigma_{\bar{x}} \cap Y$. Given a run \bar{x} and $t \in \mathcal{E}(Y)$, we define $\bar{x} \setminus t$ to be the run that omits measurements in Y . Formally, we can define this inductively on \bar{x} by setting

$$\begin{aligned} \Lambda \setminus t &= \Lambda \\ (\bar{x} \cdot (y, o)) \setminus t &:= \begin{cases} \bar{x} \setminus t & \text{if } y \in Y \\ (\bar{x} \setminus t) \cdot (y, o) & \text{otherwise.} \end{cases} \end{aligned}$$

Then we can define $P(t)$ by

$$P(t) := \{\bar{x} \setminus t \mid \bar{x} \in P \text{ is compatible with } t\}$$

Appendix B: Proofs

1. Proof of Theorem 1

Consider models $d : S$, $e : T$ and $f : U$ and an adaptive simulation $d \otimes e \rightarrow d \otimes f$. We wish to produce a simulation $e \rightarrow f$. Heuristically speaking, our overall strategy is to show that the composite $d \otimes e \rightarrow d \otimes f \rightarrow f$ needs to use d only in a non-contextual manner. More precisely, we show that we only need a compatible set of measurement protocols over

S to carry out this simulation. There are several compatible sets one could use for this: we choose one that works both for the resource theory of contextuality and for non-locality.

We first reduce to the deterministic case: by definition, there is a deterministic map $d \otimes e \otimes c \rightarrow d \otimes f$ with c non-contextual. Let us set $g := e \otimes c$ and denote the scenario of $e \otimes c$ by R . We then have a deterministic map $d \otimes g \rightarrow d \otimes f$. Now, if d can catalyze the transformation $g \rightarrow f$ once without getting spent in the process, it can do so arbitrarily many times by first catalyzing the first copy of g to f , then the second one and so on. This results in a deterministic simulation $d \otimes g^{\otimes n} \rightarrow d \otimes f^{\otimes n}$ for any n with the property that simulating the i -th copy of f uses only d and the i -th copy of g .

Thus the composite $d \otimes g^{\otimes n} \rightarrow d \otimes f^{\otimes n} \rightarrow f^{\otimes n}$ can be equivalently described by n deterministic simulations $d \otimes g \rightarrow f$ with the property that they can be run in parallel **with a single copy of d** to obtain a valid simulation $d \otimes g^{\otimes n} \rightarrow f^{\otimes n}$. As there are only finitely many deterministic simulations $d \otimes g \rightarrow f$, we can force as many as we wish to coincide by choosing a large enough n . This implies that there is a fixed deterministic simulation $d \otimes g \rightarrow f$ that yields a simulation $d \otimes g^{\otimes n} \rightarrow f^{\otimes n}$ for any n .

Consider now a fixed measurement x and the measurement protocol $\pi(x)$ over $S \otimes R$ used to simulate it. Enumerate $\mathcal{E}_R(X_R)$ as t_1, \dots, t_k . Since all the copies of x in $U^{\otimes n}$ are compatible with each other, so are their images under the simulation $d \otimes g^{\otimes n} \rightarrow f^{\otimes n}$. When simulating different copies of x one queries a single copy of d but different copies of g , this means that whatever results one gets in g the measurements done in d are always compatible. This implies that $\{\pi(x)(t_i) | i = 1, \dots, k\}$ is a compatible set of measurement protocols, where $\pi(x)(t_i)$ is defined as in Remark A. Thus they can be combined to a single measurement protocol P_x by Lemma 8. In particular, when simulating the i -th copy of x we can first measure P_x obtaining an outcome s , and then perform $\pi(x)(s)$ (which is a measurement protocol over R) in the i -th copy of R . We can then replace d by the restriction \hat{d} of $\text{MP}(d)$ to $\text{MP}(S)|_{\{P_x | x \in X_U\}}$. Moreover, considering distinct measurements x belong to different copies of $f : U$, we see that this set has to be a compatible set of measurement protocols. As the set $\{P_x | x \in X_U\}$ is comeasurable, the model \hat{d} is non-contextual, so that the deterministic simulation $\hat{d} \otimes g = \hat{d} \otimes e \otimes c \rightarrow f$ gives a probabilistic simulation $e \rightarrow f$ as desired.

2. Proof of Theorem 2

We now show how the above proof implies that there are no catalysts in the resource theory of non-locality. Indeed, assume a simulation $(\pi_i, \alpha_i)_{i=1}^n: d \boxtimes e \rightarrow d \boxtimes f$ in the resource theory of non-locality. Then $\bigotimes (\pi_i, \alpha_i)$ defines a simulation $d \boxtimes e \rightarrow d \boxtimes f$ in the resource theory of contextuality. Pre- and postcomposing by the canonical isomorphisms

$$\left(\bigotimes_{i=1}^n S_i\right) \otimes \left(\bigotimes_{i=1}^n T_i\right) \cong \bigotimes_{i=1}^n (S_i \otimes T_i)$$

then implies that we have a simulation $d \otimes e \rightarrow d \otimes f$, so that the proof gives us a simulation $(\pi', \alpha'): e \rightarrow f$. However, we must check that this map is of the form $\bigotimes (\pi'_i, \alpha'_i)$ in order for it to give a simulation $e \rightarrow f$ in the resource theory of non-locality. That this is true follows by inspecting the preceding proof. In fact, when simulating a measurement x in the T one performs a single measurement protocol in S that captures everything one might need when measuring x , and then proceeds to R . If x is a measurement at the i -th measurement site, so is this “everything” and the protocol following it in R , so that the constructed simulation $e \rightarrow f$ is indeed of the required form.

3. Proof of Theorem 3

If \mathcal{X} is closed under \otimes and contains all non-contextual models, then \mathcal{X} -assisted simulation $d \otimes e \rightarrow d \otimes f$ is given by a deterministic simulation $d \otimes e \otimes c \rightarrow d \otimes f$ where $c \in \mathcal{X}$. We can then set $g := e \otimes c$ to obtain a deterministic simulation $d \otimes g \rightarrow d \otimes f$ and retrace the proof of Theorem 1 to obtain a deterministic simulation $\hat{d} \otimes g = \hat{d} \otimes e \otimes c \rightarrow f$. As \hat{d} is non-contextual, our assumptions on \mathcal{X} imply that $\hat{d} \otimes c \in \mathcal{X}$, so that there is an \mathcal{X} -assisted simulation $e \rightarrow f$.

The proof of Theorem 2 similarly gives a proof of no-catalysis for the resource theory of \mathcal{X} -assisted n -partite non-locality.

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