Hw2 Report

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Environment gcc 5.11,Cs 工作站,Ubuntu
Language c++
Code explanation
Part 1 Build k-d tree
the feature of each layer of the tree follow the sequence of 0,1,2,3,4,5,6,7,8,0,1,2,...
feature num of each node is presented by attr
calculate the median of the feature
dat node::calculatemedium(int n) //feature n
{
    vector<dat> tmp;
    tmp.assign(s.begin(),s.end());
    sortclass = n;
    sort(tmp.begin(),tmp.end(),cmp);
    return tmp[tmp.size()/2];
                                        //指定 feature 的中位數
}
Save the median data in the tree->med then delete it from following set
use this median to partition the dataset to 2 subset(data that its value of the feature
< median =>left else =>right)
then use 2 new subset to build left and right subtree(repeat build tree)
→ dctree* dctree::buildtree(dctree *tree,node set)
if we get the subset that is null(no data belongs to this subtree)
if(set.s.size()==0)
    {
         tree->attr = -1;  // -1->不算葉節點
                                //建構子指定 attr = -1(所以她的左右子樹也不
    tree->left = new dctree();
算節點)
    tree->right = new dctree();
         return tree;
    }
Subset 只剩一個 data -> 葉節點
    if(set.s.size()==1)
                                                        {
                                    //attr =9 ->葉節點
         tree->attr=9;
         tree->left = new dctree();
         tree->right = new dctree();
         tree->med = set.s[0];
         return tree;
```

Part2 implement knn by kd tree

Algorithm I perceived:

- 1. 找到輸入的 data 經過 kdtree 走到的末節點(searctree,此點為暫時最近點)
- 2. 將此暫時最近點的 data 加入 knn 這個 vector (getktmp) 有加入 knn 的點->attr 設為-1(代表不再是節點)
- 3. 追溯樹 直到他回到樹的 root

若是 tree 的 root 到輸入點距離>輸入點到 parent 分割之後形成的 hyperplane 距離=>兩者維度為 parent->attr 的值相減為點到平面距離

代表有可能在他的兄弟子樹(besttemp->bro)中,有更近的點(比輸入點到 parent 代表的點還要近)

需進入兄弟節點搜索完才能回到 parent, o.w 回到 parent 且因此兄弟節點沒有更近點->不用搜索->所以 attr 為-1(不再是樹的節點)

```
void getktmp(dctree *best,int k) //讓 knn 維持在前 k 個最近的
    {
        if(knn.size()<k) //未滿 k 個 加進去 sort
        {
        knn.push_back(best->med);
        sort(knn.begin(),knn.end(),cmp1);
        else if(best->med.dist<knn[knn.size()-1].dist) //已經滿 k 個 但出現更近
    的
        {
            knn.erase(knn.end());
                                          //刪掉最遠的
            knn.push back(best->med);
                                          //把新的加進去 sort
            sort(knn.begin(),knn.end(),cmp1);
        }
dctree* knnfunc(dctree *tree,int k,dat q)
{
    dctree *besttemp = tree->searchtree(q,tree); //找到末節點
    while(besttemp!=tree->parent)
                                //當未回溯到 root 完畢
    {
    dat besttmp = besttemp->med;
    besttemp->med.dist = distance1(q,besttmp);
    if(besttemp->attr!=-1||besttemp->parent == 0) //若 attr不為-1(還是節點) 代
表第一次當最佳點且有資料在裡面
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```
{
       getktmp(besttemp,k); //將暫時最佳點加進去 knn
       besttemp->attr = -1; //已加過 attr 設為-1(避免重複判斷)
   }
    else
            //attr == -1 已不是節點 所以直接回到上一層
    {besttemp = besttemp->parent; continue;
    }
    if(besttemp->parent!=tree->parent) //不是 root
        dat pp = besttemp->parent->med;
    int att = besttemp->parent->attr;
       if(besttemp->bro->attr>=0&&distance1(q,tree->med)
>fabs(pp.dimension[att]-q.dimension[att])) //另一邊可能有更近句 且兄弟節點存
在
       {
           besttemp = besttemp->bro;
           besttemp = knnfunc(besttemp,k,q); //進入兄弟節點查詢 查完之後
會回到原本的 besttemp 的 parent
       }
       else
       {
                                       //兄弟節點未存在更近點
           besttemp->bro->attr = -1;
           besttemp = besttemp->parent; //回到 parent
       }
    }
                                   //此最佳點為 root 此時回到原本的
    else
tree 的位置,root knn 也判斷完畢
    {
           besttemp = besttemp->parent;
       }
    return besttemp;
}
```

```
dctree* dctree::searchtree(dat q,dctree* tree) //搜尋末節點
{
    if(tree->left->attr == -1&&tree->right->attr == -1) //有可能為葉節點或者不是
節點
    {
                          //葉節點
        if(tree->attr>=0)
        return tree;
        else
        return searchtree(q,tree->bro); //不是節點 兄弟節點存在
    }
    if(tree->right->attr<0)
                                  //右節點不存在
        return searchtree(q,tree->left);
    else if(tree->left->attr<0) return searchtree(q,tree->right); //左節點不存在
    else if(q.dimension[tree->attr]<tree->med.dimension[tree->attr]) //比中位數小
    return searchtree(q,tree->left); //往左走
    else if(q.dimension[tree->attr]>=tree->med.dimension[tree->attr]) //比中位數
大
    {
        return searchtree(q,tree->right); // 往右走
    }
}
```

Part 3 predict

找到 knn 裡面出現頻率最高的那個 class -> 預測他 看他有沒有跟測資真正 class 相同

算測資裡面有預測正確的比率