

# **COMP5212: Machine Learning**

**Lecture 9**

**Minhao Cheng**

# From last time

## Shattered

- Given a set  $S = \{x^{(i)}, \dots, x^{(d)}\}$  (no relation to the training set) of points  $x^{(i)} \in \mathcal{X}$ , we say that  $\mathcal{H}$  shatters  $S$  if  $\mathcal{H}$  can realize any labeling on  $S$ . I.e, if for any set of labels  $\{y^{(i)}, \dots, y^{(d)}\}$ , there exist some  $h \in \mathcal{H}$  so that  $h(x^{(i)}) = y^{(i)}$  for all  $i = 1, \dots, d$

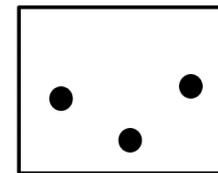
## Break point of $\mathcal{H}$

- If no data set of size  $k$  can be shattered by  $\mathcal{H}$ , then  $k$  is a break point for  $\mathcal{H}$ 
  - $m_{\mathcal{H}}(k) < 2^k$
  - VC dimension of  $\mathcal{H} : k - 1$  (the most points  $\mathcal{H}$  can shatter)

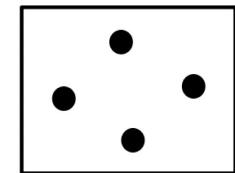
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- VC dimension of  $\mathcal{H}$  :  $k - 1$  (the most points  $\mathcal{H}$  can shatter)
- For 2-D perceptron:  $k = 4$ , VC dimension = 3

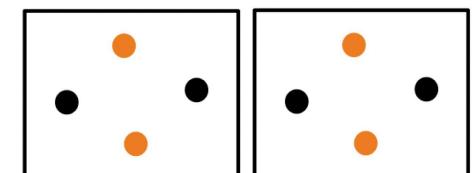
Shattered



Not Shattered



Can't generate



## Break point - examples

- Positive rays:  $m_{\mathcal{H}}(N) = N + 1$ 
  - Break point  $k = 2, d_{VC} = 1$

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  - Break point  $k = 3, d_{VC} = 2$

## Break point - examples

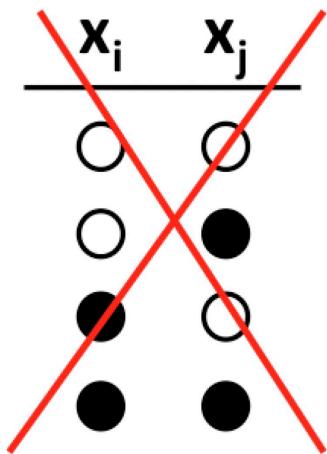
- Positive rays:  $m_{\mathcal{H}}(N) = N + 1$ 
  - Break point  $k = 2, d_{VC} = 1$
- Positive intervals:  $m_{\mathcal{H}}(N) = \frac{1}{2}N^2 + \frac{1}{2}N + 1$ 
  - Break point  $k = 3, d_{VC} = 2$
- Convex set:  $m_{\mathcal{H}}(N) = 2^N$ 
  - Break point  $k = \infty, d_{VC} = \infty$
- Connection to # of parameters

## We will show

- No break point  $\Rightarrow m_{\mathcal{H}}(N) = 2^N$
- Any break point  $\Rightarrow m_{\mathcal{H}}(N)$  is **polynomial** in  $N$

## Puzzle

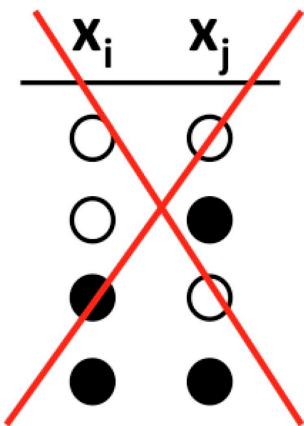
- Break point is  $k = 2$



$x_1$	$x_2$	$x_3$
○	○	○

## Puzzle

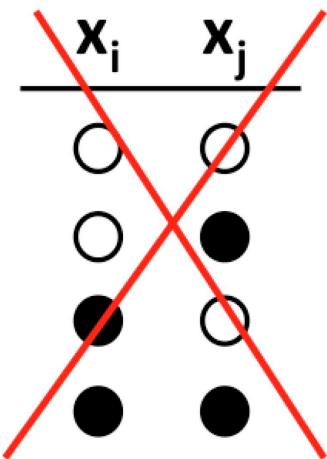
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$x_1$	$x_2$	$x_3$
○	○	○
○	○	●

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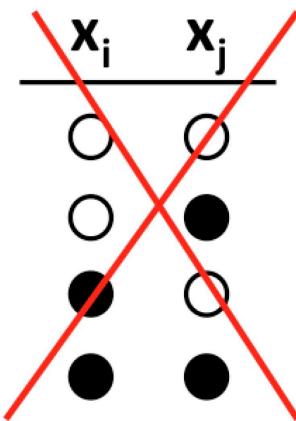
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$x_1$	$x_2$	$x_3$
○	○	○
○	○	●
○	●	○

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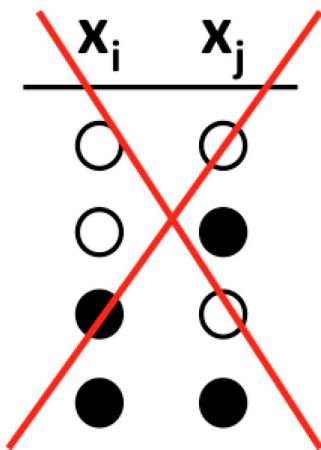
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$x_1$	$x_2$	$x_3$
○	○	○
○	○	●
○	●	○
○	●	●

## Puzzle

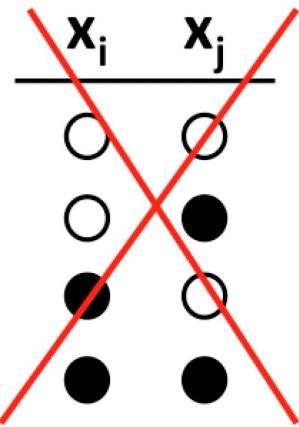
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$x_1$	$x_2$	$x_3$
○	○	○
○	○	●
○	●	○
○	●	●

## Puzzle

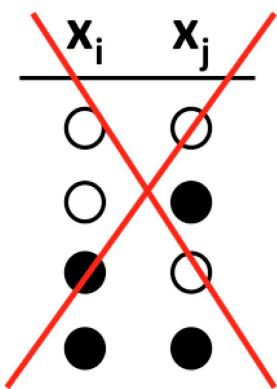
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○	○	○
○	○	●
○	●	○
●	○	○

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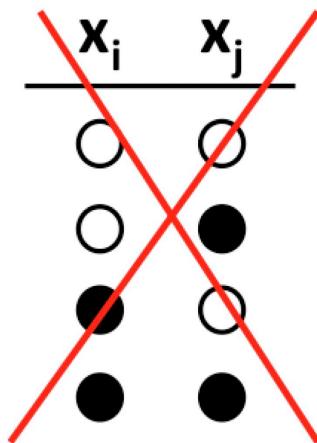
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$x_1$	$x_2$	$x_3$
○	○	○
○	○	●
○	●	○
●	○	○
●	○	●

## Puzzle

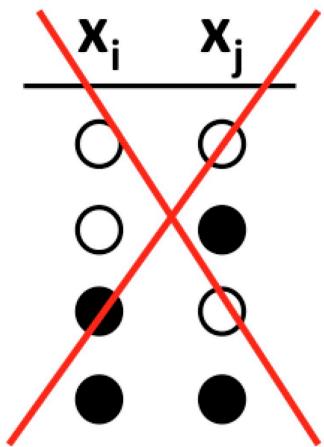
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	$x_1$	$x_2$	$x_3$
1	○	○	○
2	○	○	●
3	○	●	○
4	●	○	○
5	●	○	●

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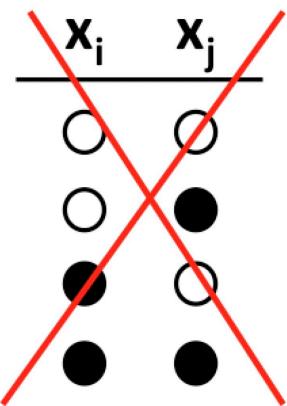
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$x_1$	$x_2$	$x_3$
○	○	○
○	○	●
○	●	○
●	○	○
●	●	○

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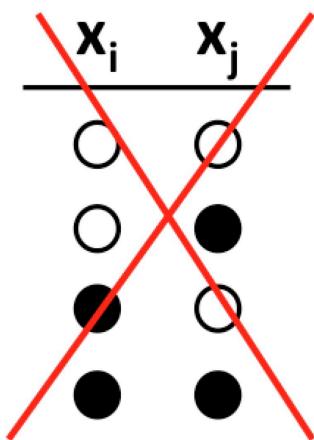
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$x_1$	$x_2$	$x_3$
○	○	○
○	○	●
○	●	○
●	○	○
●	●	○

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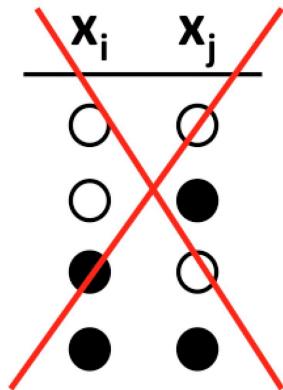
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	$x_1$	$x_2$	$x_3$
	○	○	○
	○	○	●
	○	●	○
	●	○	○
	●	●	●

## Puzzle

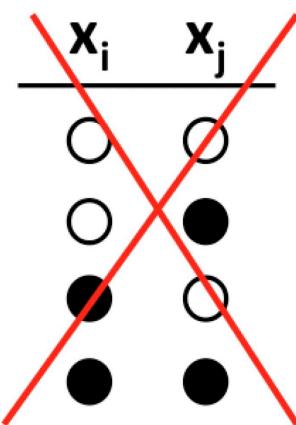
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$x_1$	$x_2$	$x_3$
○	○	○
○	○	●
○	●	○
●	○	○
●	●	●

## Puzzle

- Break point is  $k = 2$



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○	○	○
○	○	●
○	●	○
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## Bounding $m_{\mathcal{H}}(N)$

- Key quantity:
  - $B(N, k)$ : Maximum number of dichotomies on  $N$  points, with break  $k$

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- Key quantity:
  - $B(N, k)$ : Maximum number of dichotomies on  $N$  points, with break  $k$
  - If the hypothesis space has break point  $k$ , then
    - $m_{\mathcal{H}}(N) \leq B(N, k)$

## Recursive bound on $B(N, k)$

- For any “valid” set of dichotomies, reorganize rows by
  - $S_1$  : pattern of  $x_1, \dots, x_{N-1}$  only appears once
  - $S_2^+, S_2^-$  : pattern of  $x_1, \dots, x_{N-1}$  only appears twice

	# of rows	$\mathbf{x}_1$	$\mathbf{x}_2$	...	$\mathbf{x}_{N-1}$	$\mathbf{x}_N$
$S_1$	$\alpha$	+1	+1	...	+1	+1
		-1	+1	...	+1	-1
		:	:	...	:	:
		+1	-1	...	-1	-1
$S_2^+$	$\beta$	-1	+1	...	-1	+1
		+1	-1	...	+1	+1
		-1	-1	...	+1	+1
		:	:	...	:	:
$S_2^-$	$\beta$	+1	-1	...	+1	+1
		-1	-1	...	-1	+1
		:	:	...	:	:
		+1	-1	...	+1	-1
		-1	-1	...	-1	-1

## Recursive bound on $B(N, k)$

- Focus on  $x_1, \dots, x_{N-1}$  columns:
  - $\alpha + \beta \leq B(N - 1, k)$

		# of rows	$\mathbf{x}_1$	$\mathbf{x}_2$	$\dots$	$\mathbf{x}_{N-1}$	$\mathbf{x}_N$
$S_1$	$\alpha$		+1	+1	$\dots$	+1	+1
			-1	+1	$\dots$	+1	-1
			:	:	$\vdots$	$\vdots$	:
			+1	-1	$\dots$	-1	-1
			-1	+1	$\dots$	-1	+1
$S_2^+$	$\beta$		+1	-1	$\dots$	+1	+1
			-1	-1	$\dots$	+1	+1
			:	:	$\vdots$	$\vdots$	:
			+1	-1	$\dots$	+1	+1
			-1	-1	$\dots$	-1	+1
$S_2$	$\beta$		+1	-1	$\dots$	+1	-1
			-1	-1	$\dots$	+1	-1
			:	:	$\vdots$	$\vdots$	:
			+1	-1	$\dots$	+1	-1
			-1	-1	$\dots$	-1	-1
$S_2^-$	$\beta$		+1	-1	$\dots$	+1	-1
			-1	-1	$\dots$	+1	-1
			:	:	$\vdots$	$\vdots$	:
			+1	-1	$\dots$	+1	-1
			-1	-1	$\dots$	-1	-1

## Recursive bound on $B(N, k)$

- Now focus on the  $S_2 = S_2^+ \cup S_2^-$ :
- $\beta \leq B(N - 1, k - 1)$

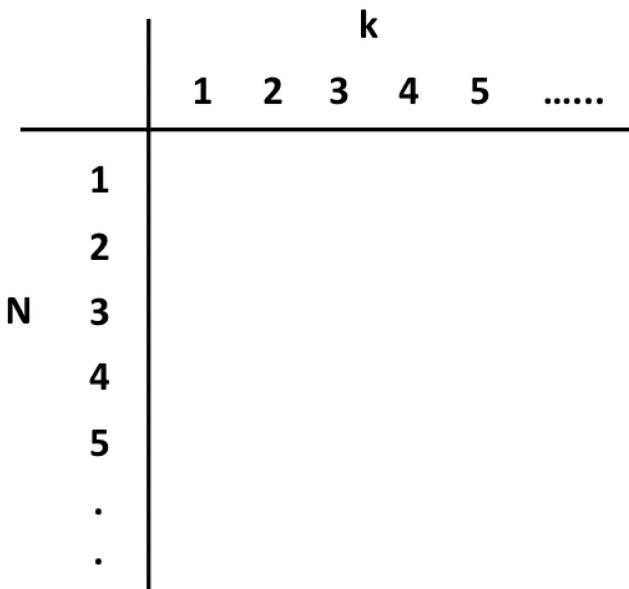
	# of rows	$\mathbf{x}_1$	$\mathbf{x}_2$	$\dots$	$\mathbf{x}_{N-1}$	$\mathbf{x}_N$
$S_1$	$\alpha$	+1	+1	$\dots$	+1	+1
		-1	+1	$\dots$	+1	-1
		$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$
		+1	-1	$\dots$	-1	-1
		-1	+1	$\dots$	-1	+1
$S_2^+$	$\beta$	+1	-1	$\dots$	+1	+1
		-1	-1	$\dots$	+1	+1
		$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$
		+1	-1	$\dots$	+1	+1
		-1	-1	$\dots$	-1	+1
$S_2^-$	$\beta$	+1	-1	$\dots$	+1	-1
		-1	-1	$\dots$	+1	-1
		$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$
		+1	-1	$\dots$	+1	-1
		-1	-1	$\dots$	-1	-1

## Recursive bound on $B(N, k)$

- $$B(N, k) = \alpha + \beta + \beta$$
- $\leq B(N - 1, k) + B(N - 1, k - 1)$
  - What's the upper bound for  $B(N, k)$ ?

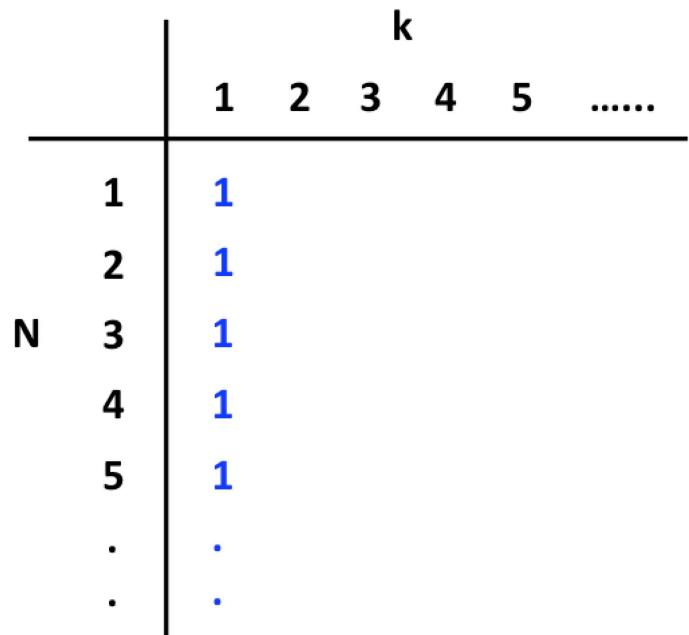
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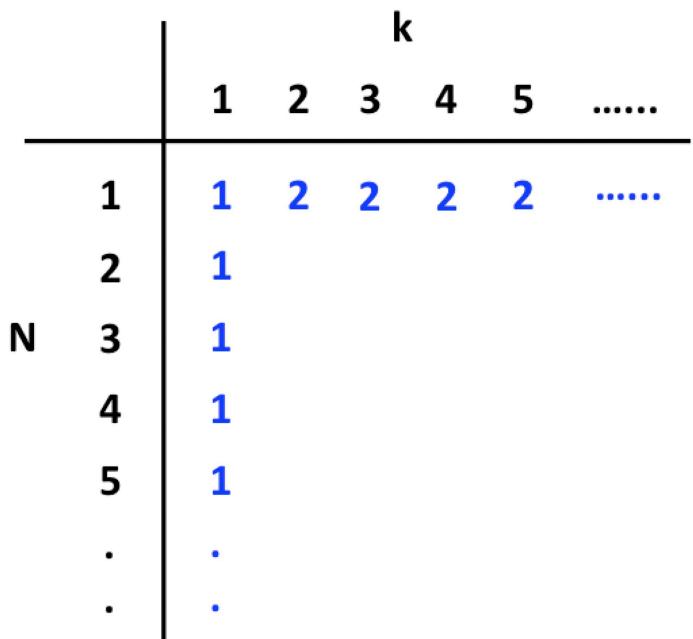
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$$\begin{aligned} B(N, k) &= \alpha + \beta + \beta \\ \bullet \quad &\leq B(N - 1, k) + B(N - 1, k - 1) \end{aligned}$$

	$k$	1	2	3	4	5	.....
$N$	1	1	2	2	2	2	.....
	2	1	3				
	3	1					
	4	1					
	5	1					
	.	.					
	.	.					

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		$k$	1	2	3	4	5	.....
		1	1	2	2	2	2	.....
		2	1	3	4	4	4	.....
N	3	1						
	4	1						
	5	1						
	.	.						
	.	.						

## Recursive bound on $B(N, k)$

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		k					
		1	2	3	4	5	.....
N		1	2	2	2	2	.....
1	1	1	2	2	2	2	.....
2	2	1	3	4	4	4	.....
3	3	1	4	7	8	8	.....
4	4	1	5	11	.....		
5	5	1	6	.	.		
.	.	.	.	.	.		
.	.	.	.	.	.		

## Analytic solution for $B(N, k)$ bound

- $B(N, k)$  is upper bounded by  $C(N, k)$ 
  - $C(N, 1) = 1, N = 1, 2, \dots$
  - $C(1, k) = 2, k = 2, 3, \dots$
  - $C(N, k) = C(N - 1, k) + C(N - 1, k - 1)$
  - Theorem:  $C(N, k) = \sum_{i=0}^{k=1} \binom{N}{i}$

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- Boundary conditions: (easy to check)

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  - $C(N, k) = C(N - 1, k) + C(N - 1, k - 1)$
- Sauer's Theorem:  $C(N, k) = \sum_{i=0}^{k-1} \binom{N}{i}$
- Boundary conditions: (easy to check)
- Induction:
  - $$\sum_{i=0}^{k-1} \binom{N}{i} = \underbrace{\sum_{i=0}^{k-1} \binom{N-1}{i}}_{\text{select } < k \text{ from } N \text{ items}} + \underbrace{\sum_{i=0}^{k-2} \binom{N-1}{i}}_{\text{N-th item not chosen}}$$

# It is polynomial!

- For a given  $\mathcal{H}$ , the break point  $k$  is fixed:

$$m_{\mathcal{H}}(N) \leq \underbrace{\sum_{i=0}^{k-1} \binom{N}{i}}_{\text{Polynomial with degree } k-1}$$

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- $\mathcal{H}$  is positive rays: (break point  $k = 2$ )
  - $m_{\mathcal{H}}(N) = N + 1$

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- $\mathcal{H}$  is 2D perceptrons: (break point  $k = 4$ )
  - $m_{\mathcal{H}}(N) = ?$

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- $\mathcal{H}$  is 2D perceptrons: (break point  $k = 4$ )

$$\bullet m_{\mathcal{H}}(N) \leq \frac{1}{6}N^3 + \frac{5}{6}N + 1$$

## Replace M by $m_{\mathcal{H}}(N)$

- Original bound:

- $\mathbb{P}[\exists h \in \mathcal{H} \text{ s.t. } |E_{\text{tr}}(h) - E(h)| > \epsilon] \leq 2\textcolor{red}{M}e^{-2\epsilon^2 N}$

- Replace M by  $m_{\mathcal{H}}(N)$

- $$\mathbb{P}[\exists h \in \mathcal{H} \text{ s.t. } |E_{\text{tr}}(h) - E(h)| > \epsilon] \leq 2 \cdot 2\textcolor{blue}{m}_{\mathcal{H}}(2N) \cdot e^{-\frac{1}{8}\epsilon^2 N}$$
  
$$\underbrace{\quad\quad\quad}_{BAD}$$

- Vapnik-Chervonenkis (VC) bound

# VC dimension

## Definition

- The VC dimension of a hypothesis set  $\mathcal{H}$ , denoted by  $d_{\text{VC}}(\mathcal{H})$ , is the largest value of  $N$  for which  $m_{\mathcal{H}}(N) = 2^N$ 
  - “The most points  $\mathcal{H}$  can shatter”

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  - “The most points  $\mathcal{H}$  can shatter”
  - $N \leq d_{\text{VC}}(\mathcal{H}) \Rightarrow \mathcal{H}$  can shatter  $N$  points
  - $k > d_{\text{VC}}(\mathcal{H}) \Rightarrow \mathcal{H}$  cannot be shattered
  - The smallest **break point** is 1 above VC-dimension

# VC dimension

## The growth function

- In terms of a break point  $k$ :

$$\bullet m_{\mathcal{H}}(N) \leq \sum_{i=0}^{k-1} \binom{N}{i}$$

- In terms of the VC dimension  $d_{\text{VC}}$ :

$$\bullet m_{\mathcal{H}}(N) \leq \sum_{i=0}^{d_{\text{VC}}} \binom{N}{i}$$

# VC dimension

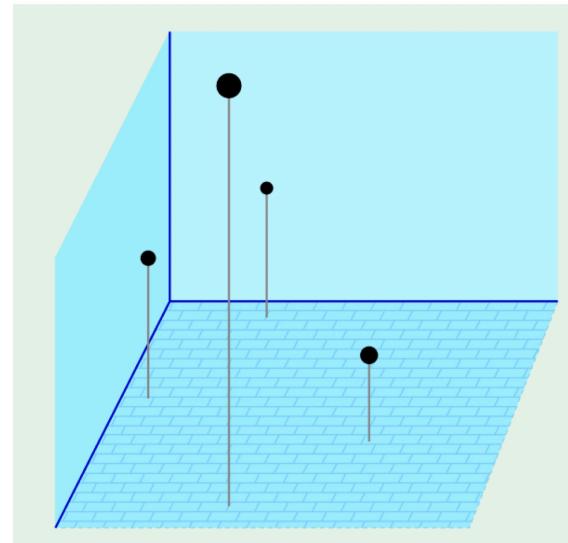
## VC dimension of linear classifier

- For  $d = 2$ ,  $d_{VC} = 3$

# VC dimension

## VC dimension of linear classifier

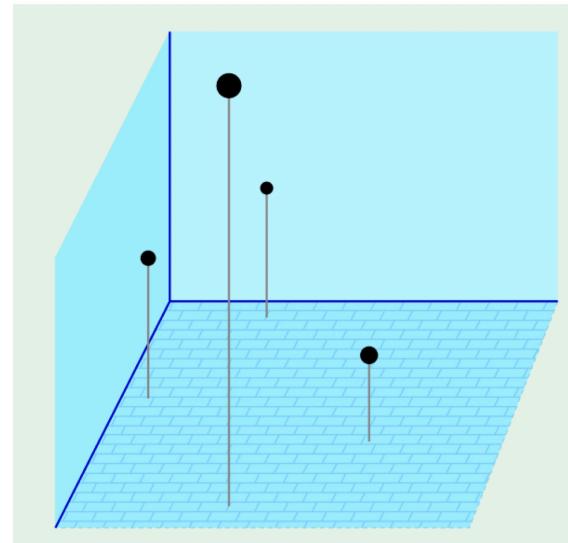
- For  $d = 2$ ,  $d_{VC} = 3$
- What if  $d > 2$ ?



# VC dimension

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- What if  $d > 2$ ?
- In general,
  - $d_{\text{VC}} = d + 1$



# VC dimension

## VC dimension of linear classifier

- For  $d = 2$ ,  $d_{VC} = 3$
- What if  $d > 2$ ?
- In general,
  - $d_{VC} = d + 1$
- We will prove  $d_{VC} \geq d + 1$  and  $d_{VC} \leq d + 1$

