

Ch2 Convex sets

- Affine line segments**: $y = \theta x_1 + (1-\theta)x_2$, $x_1, x_2 \in E^k$, $\theta \in [0, 1]$. $\text{relint } C = \{x \mid \theta x_1 + (1-\theta)x_2\}$.
- affine sets**: $C \rightarrow \text{affine}$, $x_1, \dots, x_k \in C$, $\theta_1 + \dots + \theta_k = 1 \Rightarrow \theta_1 x_1 + \dots + \theta_k x_k \in C$, $\theta_i \in E^k$.
- affine hull**: $\text{affhull } C = \{x \mid \theta_1 x_1 + \dots + \theta_k x_k \in C\} = \{x \mid \sum \theta_i x_i = \sum \theta_i x_i\}$, $\sum \theta_i = 1$.
- CE(R^n)**: $C = \{x \mid Ax = b\}$, $A \in E^{m,n}$, $b \in E^m$, $\text{affhull } C = \{x \mid Ax = b\}$.
- convex sets**: $C \rightarrow \text{convex}$, $x_1, \dots, x_k \in C$, $\theta_1 + \dots + \theta_k = 1 \Rightarrow \theta_1 x_1 + \dots + \theta_k x_k \in C$, $\theta_i \geq 0$.
- hyperplane & halfspaces**: $H = \{x \mid a^T x = b\}$, $a \in E^k$, $b \in E^1$, $\text{affhull } H = H$.
- Euclidean balls & ellipsoids**: $B(x_c, r) = \{x \mid \|x - x_c\|_2 \leq r\} = \{x \mid (x - x_c)^T (x - x_c) \leq r^2\} = \{x \mid x_c + ru \mid \|u\|_2 \leq 1\}$, $\varepsilon = \{x \mid (x - x_c)^T P^{-1} (x - x_c) \leq 1\} = \{x \mid x_c + Au \mid \|u\|_2 \leq 1\}$, length of $\varepsilon \sqrt{\lambda_1} \cdot A = P^T$, $P \in PD$, $P = P^T$.
- norm ball, norm cones**: $\text{norm ball} : G \{x \mid \|x - x_c\|_1 \leq r\}$, $\text{norm cone} : C = \{(x, t) \mid \|x\|_1 \leq t\} \subseteq E^{n+1}$, $\text{second-order cone} : C = \{(x, t) \in E^{n+1} \mid \|x\|_2 \leq t\} = \left\{ \begin{bmatrix} x \\ t \end{bmatrix} \mid \begin{bmatrix} x \\ t \end{bmatrix}^T \begin{bmatrix} I & 0 \\ 0 & t^2 \end{bmatrix} \begin{bmatrix} x \\ t \end{bmatrix} \leq 0, t \geq 0 \right\}, R^3, \{x(x, t) \mid (x^2 + y^2)/t^2 \leq 1\}$.

Operation preserve convexity

- Intersection**: $S_1, S_2 \text{ convex} \Rightarrow S_1 \cap S_2 \text{ convex}$.
- (Affine functions)**: $f(x) = Ax + b$, $f(S) = \{f(x) \mid x \in S\}$.
 - is convex**: $f'(S) = \{x \mid f(x) \in S\}$.
 - is linear**: $f'(S) = \{x \mid f(x) \in S\} = \{x \mid f(x) \geq 0\}$.
 - is scaling**: $f(S) = \{Ax \mid x \in S\}$.
 - translation**: $f(A) = \{x_0 \mid x \in S\}$.
 - projection**: $T = f(x) \in E^m \mid (x_1, x_2) \in S$.
 - Sum of two sets**: $S_1 + S_2 = \{x_1 + x_2 \mid x_1 \in S_1, x_2 \in S_2\}$.
 - linear function**: $f(x_1, x_2) = x_1 + x_2 \rightarrow S_1 + S_2$.
 - parallel sum**: $S = (x_1 + V_1) / (x_2 + V_2) = \{(x_1 + V_1) + (x_2 + V_2) \mid x_1 \in S_1, x_2 \in S_2\}$.

proper cones & generalized inequalities

- Acute K $\subseteq E^k$, proper cone**: K is convex, closed, solid, nonempty interior, K is pointed (notline $x \in K$), $x \in K$ iff $x \geq_K 0$.
- proper cone induce generalized inequalities**: $x \geq_K y \Leftrightarrow 0 - x \in K$, $x \leq_K y \Leftrightarrow y - x \in \text{int } K$.
- properties**:
 - $x \geq_K y, u \in E^k \rightarrow x + u \geq_K y + u$
 - $x \geq_K y, \forall j \in \{1, \dots, k\} \rightarrow x_j \geq_K y_j$
 - $x \geq_K y, x \geq_0, \alpha x \geq_K \alpha y$
 - $x \geq_K x$
 - $x \geq_K y, y \geq_K x \rightarrow x = y$
 - $x_i \geq_K y_i, i=1, \dots, n \rightarrow x \geq_K y$
 - $i \rightarrow \infty$

Show that maximum of a convex function f over the polyhedron P $\text{con}\{V_1, \dots, V_n\}$ is achieved at one of its vertices.

$$\sup_{x \in P} f(x) = \max_{x \in P} f(x)$$

minimum & minimal elements

- minimum**: every y , no point is $x \leq y$ "more" than it.
- minimal**: $x \in S$, no point is $y \in S$ "less" than it, only if $y = x$.
- minimum x** : $x \in S$ - x minimum iff $S \subseteq x + K$.
- minimal x** : $x \in S$ - x minimal iff $(x - K) \cap S = \{x\}$.

(Separating hyperplane theorem)

- CE(R^n)**: $x_0 \in E^k$, $x_0 \notin \text{affhull } C$, $x_0 \in \text{affhull } D$, $\exists a \in E^k$ s.t. $a^T x \leq b$ for all $x \in C$ and $a^T x \geq c$ for all $x \in D$.
- LMIs (Linear matrix inequalities)**: $A(x) = x_1 A_1 + \dots + x_n A_n + B$, $B, A_i \in E^{m,m}$, $\{x \mid A(x) \leq B\}$ is convex, $\{x \mid f(x) = B - A(x)\} = \{x \mid f(x) = B^T - A^T(x)\}$ is convex, $\{x \mid f(x) = P^{-1}(x)\}$ is convex, $\{x \mid (x - x_e)^T P^{-1}(x - x_e) \leq 1\}$ is convex, $\{x \mid f(x) = (P^T x, e^T x)\}$ is convex, $E = \{x \mid (x - x_e)^T P^{-1}(x - x_e) \leq 1\}$, $P \in S^n$, $P \succ 0$, $\{x \mid f(u) = P^T u + e^T x\} = \{x \mid \|u\|_2 \leq 1\}$.
- Dual generalized inequalities**: $K^* = \text{dual generalized inequality of } K$, $-x \geq_K y \iff -x^T \leq y^T \forall A \in E^{k,n}$, $-x \geq_K y \iff -x^T \leq y^T \forall A \in E^{k,n}$, $-x \geq_K y \iff -x^T \leq y^T \forall A \in E^{k,n}$, $-x \geq_K y \iff -x^T \leq y^T \forall A \in E^{k,n}$.
- Dual cones**: K is a cone, K^* is a dual cone, $K^* = \{y \mid x^T y \geq 0 \text{ for all } x \in K\}$.
- Linear-fractional & perspective functions**: $P: R^{n+1} \rightarrow R^n$, $\text{dom } P = R^{n+1} \setminus R_{+}$, $P(\beta, x) = \frac{x}{\beta}$, $C \subseteq \text{dom } P$, convex, $P(C) = \{P(x) \mid x \in C\}$, convex, $P^*(C) = \{P(x) \mid x \in C, x \neq 0\}$, convex.
- Linear-fractional function**: $f: R^n \rightarrow R^{n+1}$, $f(x) = \frac{Ax + b}{c^T x + d}$, $\text{dom } f = \{x \mid c^T x + d > 0\}$, $f(x) = \frac{(Ax + b)^T}{c^T x + d}$, $\text{dom } f = \{x \mid c^T x + d > 0\}$, $\text{image } f = \{y \mid y = f(x) \text{ for some } x \in \text{dom } f\}$, $\text{image } f$ is all convex.
- Orthogonal complement**: $V^* = \{y \mid y^T v = 0, \forall v \in V\}$.
- nonnegative without 0**: $y^T x \geq 0 \forall y \in V \setminus \{0\} \Rightarrow y \in V^*$.
- Self-dual**: $\text{dom } f = \text{dom } f^*$.
- Positive semidefinite cone**: $K = \{x \in E^{n+1} \mid x^T x \geq 0, \forall x \in E^n\}$, $K^* = \{Y \mid Y + (Y^T) \geq 0, \forall Y \in E^{n+1}\}$, $\text{dom } f = \text{dom } f^*$.
- Proof: self-dual**: $\text{dom } f = \text{dom } f^* \Rightarrow \text{dom } f = \text{dom } f^* \Rightarrow \text{dom } f = \text{dom } f^* \Rightarrow \text{dom } f = \text{dom } f^*$.
- Properties of dual cones**: $K^* \subseteq \text{closed \& convex}$, $K_1 \subseteq K_2 \rightarrow K_2^* \subseteq K_1^*$, $\text{cl } K$ is pointed, $\text{int } K^* \neq \emptyset$, $\text{int } K^* \neq \emptyset \rightarrow K^* \text{ is pointed}$, $K^* \text{ closure of conv } K$, K proper cone, K^* also proper, $K^{**} = K$.

Ch3 convex functions

(Definition)

$$f: \mathbb{R}^n \rightarrow \mathbb{R}$$

if dom f is convex set
 $x, y \in \text{dom } f \in [0, 1]$
 $\Rightarrow f(\theta x + (1-\theta)y) \leq \theta f(x) + (1-\theta)f(y)$
 $\Rightarrow (-f) \text{ convex}, f \text{ concave}$
 $\Rightarrow f \text{ is convex, i.f.f.}$

$g(t) = f(x+tV)$ is convex
 $\Leftrightarrow t \in [0, 1] \text{ dom } f \text{ (line)}$

(1st-order condition)

f differentiable

f convex i.f.f.

dom f convex

$$\Rightarrow f'(x) \geq 0 \text{ (convex)}$$

(concave)

f differentiable twice

f convex i.f.f.

dom f convex

$\Rightarrow f''(x) \geq 0$ (positive)

(concave)

Ch4 Convex problem

Basic terminology

minimize $f_0(x)$
subject to $f_i(x) \leq 0, i=1, \dots, m$
 $h_i(x) = 0, i=1, \dots, p$

(optimal value)

$p^* = \inf\{f_0(x) | f_i(x) \leq 0, i=1, \dots, m$
 $\quad \quad \quad h_i(x) = 0, i=1, \dots, p\}$

$p^* = \infty$ if problem infeasible
 $p^* = -\infty$ if problem unbounded below

D.T.W. $D = \bigcap_{i=1}^m \text{dom } f_i \cap \bigcap_{i=1}^p \text{dom } h_i$

(optimal & locally optimal points)

$\rightarrow X_{\text{opt}} = \{x | f_0(x) \leq 0, h_i(x) = 0, f_i(x) = p^*, i=1, \dots, m, i=h, \dots, p\}$

$\rightarrow X @ f_0(x) \leq p^* + \epsilon, \epsilon \text{-suboptimal}$

$\rightarrow f_0(x) = \inf\{f_0(y) | f_i(y) \leq 0, h_i(y) = 0$
 $\quad \quad \quad \text{if } \exists R > 0 \quad \|y-x\|_2 \leq R\}$ (locally problem)

\rightarrow optimal value attained
achieved
infeasible \Rightarrow solvable
unbounded \Rightarrow ∞

find x subject to $f_i(x) \leq 0, h_i(x) = 0$
 \rightarrow feasibility problem

implicit constraints $x \in D = \bigcap_{i=1}^m \text{dom } f_i \cap \bigcap_{i=1}^p \text{dom } h_i$
explicit constraints $h_i(x) = 0, f_i(x) \leq 0$

Convex Optimization

minimize $f_0(x)$
subject to $f_i(x) \leq 0, i=1, \dots, m$ convex
 $a_i^T x = b_i, i=1, \dots, p$ affine

feasible set $D = \bigcap_{i=0}^m \text{dom } f_i$

\rightarrow minimize a convex objective function
over
a convex set

Local & global optima

locally optima = global optima

$\rightarrow x$ is locally optimal $\Leftrightarrow x$ feasible

$f_0(x) = \inf\{f_0(y) | y \text{ feasible}, \|y-x\|_2 \leq R\}$
 $R > 0$

\rightarrow proof:

if x not globally optimal,

$\exists y \quad f_0(y) < f_0(x), \|y-x\|_2 > R$

also

$$g = (1-\theta)x + \theta y \quad \theta = \frac{R}{\|y-x\|_2} \rightarrow \|g-x\|_2 = \frac{R}{2} < R$$

$$\therefore f_0(g) \leq (1-\theta)f_0(x) + \theta f_0(y) < \theta f_0(x)$$

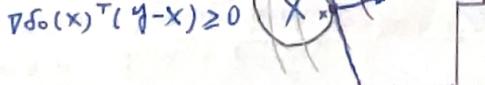
↑ contradicts

$$f_0(x) = \inf\{f_0(y) | y \text{ feasible}, \|y-x\|_2 \leq R\}$$

(optimality criterion)

recall $f_0(y) \geq f_0(x) + \nabla f_0(x)^T(y-x)$

x is optimal $\Leftrightarrow \nabla f_0(x) = 0$



more on vector optimization

Scalarization & \mathbb{R}^n
for any $\lambda \neq 0$, if \tilde{x} is an optimal point
for the scalar optimization problem below

minimize $\lambda^T f_0(x)$
subject to $f_i(x) \leq 0, i=1, \dots, m$
 $h_i(x) = 0, i=1, \dots, p$

\rightarrow then \tilde{x} is pareto optimal for the
vector optimization problem

\rightarrow for every pareto optimal point x^P ,

$\exists \lambda \geq 0, \lambda \neq 0$, such that

\tilde{x} is an optimal point of
scalarized problem

(equivalent convex problems)

(Eliminating equality constraints)

minimize $f_0(x)$
subject to $f_i(x) \leq 0, i=1, \dots, m$
 $h_i(x) = 0, i=1, \dots, p$

\downarrow
minimize $f_0(Fg+x_0)$
subject to $f_i(Fg+x_0) \leq 0, i=1, \dots, m$

\downarrow
introducing equality constraints
minimize $f_0(Ax+b_0)$
subject to $f_i(Ax+b_i) \leq 0, i=1, \dots, m$

\downarrow
minimize $f_0(y)$
subject to $f_i(y) \leq 0, i=1, \dots, m$
 $y = Ax+b_i, i=0, 1, \dots, m$

(introducing slack variables for linear inequalities)

minimize $f_0(x)$
subject to $a_i^T x \leq b_i, i=1, \dots, m$

\downarrow
minimize $f_0(x)$
subject to $A_i^T x + s_i = b_i, i=1, \dots, m$
 $s_i \geq 0, i=1, \dots, m$

(epigraph problem form)

minimize t
 $x+t$
subject to $f_0(x)-t \leq 0$
 $f_i(x) \leq 0, i=1, \dots, m$
 $a_i^T x + b_i \leq t, i=1, \dots, p$

(minimizing over some variables)

minimize $f_0(x_1, x_2)$
subject to $f_i(x_1) \leq 0, i=1, \dots, n$

\downarrow
minimize $\tilde{f}_0(x_1)$
subject to $f_i(x_1) \leq 0, i=1, \dots, m$
where $\tilde{f}_0(x_1) = \inf_{x_2} f_0(x_1, x_2)$

(Quasiconvex function)

minimize $f_0(x)$
subject to $f_i(x) \leq 0, i=1, \dots, m$
 $Ax=b$

w/ $f_0: \mathbb{R}^n \rightarrow \mathbb{R}$ quasiconvex
 f_1, \dots, f_m convex

$f_0(x) \leq t \Leftrightarrow \partial f_0(x) \leq 0$
 t -sublevel
 $\subseteq \mathbb{R}^n$

for w.l.o.g. find x
subject to $\partial f_0(x) \leq 0$
 $\rightarrow f_i(x) \leq 0, i=1, \dots, m$

\downarrow
suppose f_0 is differentiable
Let \bar{X} be the feasible set.
if $x \in \bar{X} \wedge \nabla f_0(x)^T(y-x) > 0$

$\forall \text{ all } y \in \bar{X} \setminus \{x\} :$
 x is optimal

(Linear Optimization Problem)

minimize $C^T x + d$ affine
subject to $Gx \leq h$ affine
 $\rightarrow Ax = b$ affine

(feasible set polyhedron)

\downarrow

Linear-fractional program

minimize $f_0(x)$
subject to $Gx \leq h$
 $Ax = b$

\downarrow
 $f_0(x) = \frac{C^T x + d}{B^T x + b}$ dom $f_0 = \{x | B^T x + b > 0\}$

is equivalent

\downarrow
minimize $C^T x + d$
subject to $Gy \leq h$
 $Ay = b$

\downarrow
 $e^T y + d = 1$

\downarrow
 $e^T y + h = 1, i=1, \dots, p$

\downarrow
 $f_0(y) = \log \left(\sum_{k=1}^{K_0} e^{a_k^T y + b_k} \right)$

\downarrow
subject to $\sum_{k=1}^{K_0} e^{a_k^T y + b_k} \leq 1, i=1, \dots, m$

\downarrow
 $e^{a_k^T y + b_k} = 1, i=1, \dots, P$

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subject to $\sum_{k=1}^{K_0} e^{a_k^T y + b_k} = 1, i=1, \dots, m$

\downarrow
 $e^{a_k^T y + b_k} = 1, i=1, \dots, P$

\downarrow
 $f_0(y) = \log \left(\sum_{k=1}^{K_0} e^{a_k^T y + b_k} \right)$

\downarrow
subject to $\sum_{k=1}^{K_0} e^{a_k^T y + b_k} = 1, i=1, \dots, m</math$

Algorithms (cont'd) Ch 9, 10, 11 Ch 9-10, 11
equality constrained minimization

minimize $f(x)$ s.t. $R \rightarrow R$
subject to Ax = b $A \in \mathbb{R}^{m \times n}$, $\text{rank } A = P$
- suppose Slater's conditions hold
i.e., $\exists x \in \text{relint } D$ such that
 $AX = b$

\tilde{x} is optimal i.f.f. $\exists \tilde{v}$ such that
(\tilde{x}, \tilde{v}) satisfies KKT!

$\begin{cases} A\tilde{x} = b \\ \nabla f(\tilde{x}) + A^T \tilde{v} = 0 \end{cases}$

solve ① = solve ②

Newton's method for
equality constraints

- Derivation of Newton Step

① @ x ,

• addition of v : $x+v$

• $x+v$ should be feasible

$A(x+v) = b$

$\nabla v = 0$

② $f(x+v) = f(x) + \nabla f(x)^T v + \frac{1}{2} v^T \nabla^2 f(x) v$

from ① & ② we get

minimize $f(x) + \nabla f(x)^T v + \frac{1}{2} v^T \nabla^2 f(x) v$

subject to $A v = 0$ → ③

③ \tilde{v} is optimal iff. $\exists w$ such that
(\tilde{v}, w) satisfies KKT!

$\begin{cases} A\tilde{v} = 0 \\ \nabla f(x) + \nabla^2 f(x)\tilde{v} + A^T w = 0 \end{cases}$

$\Leftrightarrow \begin{bmatrix} \nabla^2 f(x) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} \tilde{v} \\ w \end{bmatrix} = \begin{bmatrix} -\nabla f(x) \\ 0 \end{bmatrix}$

$\Delta x_{nt} = \tilde{v}$

$\Leftrightarrow \begin{bmatrix} \nabla^2 f(x) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} \Delta x_{nt} \\ w \end{bmatrix} = \begin{bmatrix} -\nabla f(x) \\ 0 \end{bmatrix}$

- Derivation of New Step (Again!)

④ @ x ,

• addition of v : $x+v$

• ($x+v, w$) satisfies ③ (KKT)!

$\begin{cases} A(x+v) = b \\ \nabla f(x+v) + A^T w = 0 \end{cases}$

$\Leftrightarrow \begin{cases} A(x+v) = b \\ \nabla f(x) + \nabla^2 f(x)v + A^T w = 0 \end{cases}$

$\Leftrightarrow \begin{bmatrix} \nabla^2 f(x) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} v \\ w \end{bmatrix} = \begin{bmatrix} -\nabla f(x) \\ 0 \end{bmatrix}$

$\Delta x_{nt} = \tilde{v}$

$\Leftrightarrow \begin{bmatrix} \nabla^2 f(x) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} \Delta x_{nt} \\ w \end{bmatrix} = \begin{bmatrix} -\nabla f(x) \\ 0 \end{bmatrix}$

The Newton decrement

- pretty similar to unconstrained case

$\lambda(x) = (\Delta x_n^T \nabla^2 f(x) \Delta x_n)^{0.5}$

- descent direction

- from ④

$\nabla f(x) \Delta x_{nt} + A^T w = -\nabla f(x)$

- for descent direction

$\nabla f(x)^T \Delta x_{nt} < 0$

$\Rightarrow -(\nabla^2 f(x) \Delta x_{nt} + A^T w)^T \Delta x_{nt} < 0$

$= -\Delta x_{nt}^T \nabla^2 f(x) \Delta x_{nt} + A^T w^T \Delta x_{nt} < 0$

$= -\Delta x_{nt}^T \nabla^2 f(x) \Delta x_{nt} + 0 < 0$

$= -\lambda^2(x) < 0$

- we got central point
- we got dual feasible at $x^*(t)$
with $\lambda^*(t), \nu^*(t)$

justification of stopping criterion

$$\begin{aligned} f(x) - p^* &= f(x) - \inf \{f(y) | Ay = b\} \\ &\Rightarrow f(x) - \inf_y \{f(y) | Ay = b\} \\ &= f(x) - \inf_v \{f(x+tv) | A(x+v) = b\} \\ &= f(x) - \inf_v \{f(x+tv) | Av = 0\} \\ &= f(x) - \inf_v \{f(x) + \nabla f(x)^T tv + \frac{1}{2} t^2 \nabla^2 f(x) v^2 | Av = 0\} \\ &= f(x) - \inf_v \{f(x) + \nabla f(x)^T tv + \frac{1}{2} t^2 \nabla^2 f(x) v^2 | Av = 0\} \end{aligned}$$

$$\begin{aligned} \text{from ④ } v &= \Delta x_{nt} \\ &= f(x) - f(x) - \nabla f(x)^T \Delta x_{nt} - \frac{1}{2} \Delta x_{nt}^T \nabla^2 f(x) \Delta x_{nt} \\ &= -(-\lambda^2(x)) - \frac{1}{2} \lambda(x)^2 \\ &= \frac{1}{2} \lambda(x)^2 \end{aligned}$$

Newton's method
(given) $x^{(k)}$ Edomif

$$\begin{aligned} Ax &= b && (1: n, :) \\ \text{(repeat)} \quad a. \Delta x^{(k+1)} &= \begin{bmatrix} \nabla^2 f(x) & A^T \\ A & 0 \end{bmatrix}^{-1} \begin{bmatrix} -\nabla f(x) \\ 0 \end{bmatrix} \\ x^{(k+1)} &= \nabla f(x^{(k)})^T \nabla^2 f(x^{(k)})^{-1} \nabla f(x^{(k)}) \\ b. \text{(Until)} \quad (x^{(k)})^2 / 2 &\leq \epsilon \\ c. \text{backtracking line search} \quad \text{get } t^{(k)} \\ d. x^{(k+1)} &= x^{(k)} + t^{(k)} \Delta x^{(k)} \end{aligned}$$

convergence analysis
similar to unconstrained

Interior Point Method

minimize $f_0(x)$ → ①

subject to $f_i(x) \leq 0$ $i=1 \dots m$
 $AX = b$ $A \in \mathbb{R}^{m \times n}$, $\text{rank } A = P$

- assume strictly feasible
 $\exists x \in D$ such that $f_i(x) < 0$ $i=1 \dots m$
 $AX = b$

- transfer inequality to equality

Logarithmic Barrier

- $I-(u) = \begin{cases} 0 & \text{if } u \leq 0 \\ -\infty & \text{if } u > 0 \end{cases}$

- minimize $f_0(x) + \sum_{i=1}^m I-(f_i(x))$
such that $AX = b$ → ②

- define $\phi(x) = -\sum_{i=1}^m \log(-f_i(x))$

- where:

$$\nabla \phi(x) = \sum_{i=1}^m \frac{1}{-f_i(x)} \nabla f_i(x)$$

$$\begin{aligned} \nabla^2 \phi(x) &= \sum_{i=1}^m \frac{1}{-f_i(x)^2} \nabla f_i(x) \nabla f_i(x)^T \\ &\quad + \sum_{i=1}^m \frac{1}{-f_i(x)} \nabla^2 f_i(x) \end{aligned}$$

- minimize $t f_0(x) + \phi(x)$ → ③

subject to $AX = b$
 $\phi(x) = -\sum_{i=1}^m \log(-f_i(x))$

Central Path

- $\frac{1}{t} \phi(x) + t \uparrow$, closer to $I-(u)$

- optimal solution of ③

$x^*(t)$ central point

- central path: $\{x^*(t) | t > 0\}$

- Question:

How good is $x^*(t)$?

A. minimize $t f_0(x) + \phi(x)$ → ④

subject to $AX = b$

$x^*(t) \rightarrow$ strictly feasible:

$Ax^*(t) = b$ $f_i(x^*(t)) < 0$

• $\exists \lambda \in \mathbb{R}^P$ such that

$$+ \nabla f_0(x^*(t)) + \nabla \phi(x^*(t)) + A^T \lambda = 0$$

$$\Rightarrow + \nabla f_0(x^*(t)) + \sum_{i=1}^m \frac{1}{-f_i(x^*(t))} \nabla f_i(x^*(t)) + A^T \lambda = 0$$

$$\Rightarrow \nabla f_0(x^*(t)) + \sum_{i=1}^m \lambda_i \frac{1}{-f_i(x^*(t))} \nabla f_i(x^*(t)) + A^T \lambda = 0$$

$$\Rightarrow x^*(t) = \arg \min L(x, \lambda, t) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \frac{1}{2} t^2 \lambda^T \lambda$$

$$\Rightarrow \phi(x^*(t), \lambda^*(t)) = f_0(x^*(t)) + \sum_{i=1}^m \lambda_i^*(t) f_i(x^*(t)) + \frac{1}{2} t^2 \lambda^*(t)^T (A x^*(t) - b)$$

$$= f_0(x^*(t)) + \sum_{i=1}^m \frac{1}{-f_i(x^*(t))} f_i(x^*(t)) + 0 = f_0(x^*(t)) - \frac{m}{t}$$

hyperbolic constraints

- $x^T x \leq y^2$ $y \geq 0$ $y \geq 0$

- $x^T x \leq 4y^2$ $y \geq 0$ $y \geq 0$

- $4x^T x + y^2 + z^2 \leq y^2 + z^2 + 4y^2$ $y \geq 0$ $z \geq 0$

- $4x^T x + y^2 + z^2 - 2y^2 \leq y^2 + z^2 + 2y^2$ $y \geq 0$ $z \geq 0$

- $4x^T x + (y-z)^2 \leq (y+z)^2$ $y \geq 0$ $z \geq 0$

- $\|x\|_2 \leq \|y\|_2$ $y \geq 0$ $z \geq 0$

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Linear Algebra Basics :

- 4 ways to write a linear system
- $\begin{cases} 2x_1 + 3x_2 - 2x_3 = 7 \\ x_1 - x_2 - 3x_3 = 5 \end{cases}$ system of equations
- $\left(\begin{array}{ccc|c} 2 & 3 & -2 & 7 \\ 1 & -1 & -3 & 5 \end{array} \right)$ augmented matrix
- $x_1 + x_2 + x_3 = b$
- $x_1 + x_2 + x_3 = b$ vector equation
- $\begin{pmatrix} 2 & 3 & -2 \\ 1 & -1 & -3 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 7 \\ 5 \end{pmatrix}$
- $Ax = b$ \exists soln i.f.f. b is in the span of column of A

Schurz complement

$$\begin{bmatrix} A & B \\ B^T & C \end{bmatrix} \succeq 0$$

$$\downarrow$$

$$C - B^T A^{-1} B \succeq 0$$

Matrix manipulation

$$\begin{aligned} (AB)^{-1} &= B^{-1}A^{-1} \\ (ABC)^{-1} &= C^{-1}B^{-1}A^{-1} \\ (A^T)^{-1} &= (A^{-1})^T \\ (A+B)^T &= A^T + B^T \\ (AB)^T &= B^T A^T \\ (ABC)^T &= C^T B^T A^T \end{aligned}$$

$$\text{tr}(A) = \sum_i \text{A}_{ii}$$

$$\text{tr}(A) = \sum_i \lambda_i; \lambda_i = \text{eig}(A)$$

$$\text{tr}(A) = \text{tr}(A^T)$$

$$\text{tr}(AB) = \text{tr}(BA)$$

$$\text{tr}(A+B) = \text{tr}(A) + \text{tr}(B)$$

$$\text{tr}(ABC) = \text{tr}(BCA) = \text{tr}(CAB)$$

$$\text{tr}(AA^T) = \text{tr}(AAT)$$

$$\det(A) = \prod_i \lambda_i; \lambda_i = \text{eig}(A)$$

$$\det(cA) = c \det(A)$$

$$\det(A^T) = \det(A)$$

$$\det(AB) = \det(A)\det(B)$$

$$\det(A^{-1}) = 1/\det(A)$$

$$\det(A^n) = \det(A)^n$$

$$\det(I+uv^T) = 1+u^T v$$

$$\nabla f(x) = \begin{bmatrix} \frac{\partial f(x)}{\partial x_1} \\ \frac{\partial f(x)}{\partial x_2} \\ \vdots \\ \frac{\partial f(x)}{\partial x_n} \end{bmatrix}$$

$$\nabla^2 f(x) = \begin{bmatrix} \frac{\partial^2 f(x)}{\partial x_1^2} & \frac{\partial^2 f(x)}{\partial x_1 \partial x_2} & \dots & \frac{\partial^2 f(x)}{\partial x_1 \partial x_n} \\ \frac{\partial^2 f(x)}{\partial x_2 \partial x_1} & \frac{\partial^2 f(x)}{\partial x_2^2} & \dots & \frac{\partial^2 f(x)}{\partial x_2 \partial x_n} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{\partial^2 f(x)}{\partial x_n \partial x_1} & \frac{\partial^2 f(x)}{\partial x_n \partial x_2} & \dots & \frac{\partial^2 f(x)}{\partial x_n^2} \end{bmatrix}$$

PD matrix M

$\text{eig}(M); > 0$, symmetric

convert QP to SOCP

$$\begin{aligned} \text{minimize} \quad & x^T A x + a^T x \\ \text{subject to} \quad & Bx \leq b \\ & s.t. \end{aligned}$$

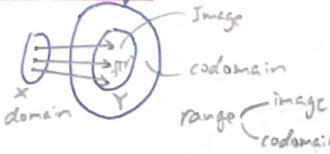
$$\begin{aligned} \Rightarrow \text{minimize} \quad & y + a^T x \\ \text{s.t.} \quad & Bx \leq b \\ & x^T A x \leq y \end{aligned} \Rightarrow \begin{aligned} & y \geq x^T A x \\ & \Rightarrow 0 \geq x^T A x - y \\ & \Rightarrow 0 \geq 4x^T A x - 4y \\ & \Rightarrow 0 \geq 4x^T A x + (1-y)^2 - (1+y)^2 \\ & \Rightarrow (1+y)^2 \geq 4x^T A x + (1-y)^2 \end{aligned}$$

Grammatic

$$\begin{bmatrix} a & b & c \\ a^2 & ab & ac \\ b^2 & bc & bc \\ c^2 & ac & bc \end{bmatrix}$$

for R^3

domain, range



Integration by change of variables.

$$\int_0^x f(t) dt \quad t=sx \quad \int_{\theta(b)}^{f(b)} f(u) du \quad u=\theta(x)$$

$$\int_a^b f(sx) ds \quad \int_a^b f(g(x)) g'(x) dx$$

$$\Delta B(x, \epsilon) = \{y \in R^n | \|y-x\|_2 \leq \epsilon\}$$

interior point: $\exists \epsilon > 0$

$$x \in \text{int } B(x, \epsilon) \subset C$$

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$$\textcircled{O} \setminus \textcircled{x} \rightarrow \text{int } C$$

limit point of C

if $\forall \epsilon > 0$, excluding

$$(B(x, \epsilon) \setminus \{x\}) \cap C \neq \emptyset$$

or

x is limit point of set S

if $\forall \epsilon > 0, \exists y \in S \setminus \{x\}$

$$w \delta(x, y) < \epsilon$$

closure

$$cl(C) = C \cup L(C) \quad \text{includes } bd(C)$$

cl(C), closed

cl(C), smallest closed set

contains C

$$C \subseteq S,$$

$$cl(C) \subseteq S$$

set C is closed i.f.f.

$$C = cl(C)$$

Boundary

$$bd(C) = cl(C) \setminus \text{int}(C)$$

$$\text{int}(C) \subseteq C \subseteq cl(C)$$

$$C \text{ open i.f.f. } C \cap bd(C) = \emptyset$$

$$C \text{ closed i.f.f. } bd(C) \subseteq C$$

$$\Delta f_0(x) = \frac{1}{2} x^T P x + q^T x + r_0$$

$$\nabla f_0(x) = \frac{1}{2} (P + P^T) x + q$$

$$\Delta \begin{bmatrix} 1 & 1 & 1 \\ 1 & 2 & 3 \\ 1 & 3 & 6 \end{bmatrix} \quad \text{eg. of SPD}$$

$\Delta \det[\lambda I - A]$ polynomial \rightarrow find eigenvalues

$\det[\lambda B - A] \quad \dots \rightarrow$ find generalized eigenvalues

linear-fractional programming \rightarrow quasiconvex function

$$\begin{aligned} \text{minimize} \quad & f(x) = \frac{c^T x + d}{e^T x + f} \quad \text{dom } f = \{x | e^T x + f > 0\} \\ \text{subject to} \quad & Gx \leq h \\ & Ax = b \\ & x \geq 0 \end{aligned}$$

$$\text{if } \{x | Gx \leq h, Ax = b, e^T x + f > 0\} \neq \emptyset$$

$$\begin{aligned} \text{minimize} \quad & C^T y + d \bar{y} \\ \text{subject to} \quad & Gy - h \bar{y} \leq 0 \\ & Ay - b \bar{y} = 0 \\ & \bar{y}^T y + f \bar{y} = 1 \\ & \bar{y} \geq 0 \end{aligned} \quad \begin{aligned} y &= \frac{x}{e^T x + f} \\ \bar{y} &= \frac{1}{e^T x + f} \end{aligned}$$

minimize $y + a^T x$

$$\text{s.t.} \quad \left\| \begin{bmatrix} 2A^T x \\ 1-y \end{bmatrix} \right\|_2 \leq 1+y$$

$$Bx \leq b$$

Descent Method

$$\Delta x^{(k+1)} = x^{(k)} + t^{(k)} \Delta x^{(k)}$$

step size
 direction
 search length

step
 (search direction)
 search length search direction

$\Delta f(x^{(k+1)}) < f(x^{(k)})$

Δ General Descent method

1. (given) $x^{(0)} \in \text{dom } f$

2. (repeat)

a. get descent direction Δx

$\Delta x^{(k)}$ Search direction

b. line search $t^{(k)}$ ——————

$t^{(k)}$ search length

1. Exact Line Search

$$t = \underset{s \geq 0}{\operatorname{argmin}} f(x + s\Delta x)$$

c. update

$$x := x + t\Delta x$$

3. (until)

$$\|\nabla f(x)\|_2 \leq \eta \quad (\text{oftenly})$$

2. Backtracking Line Search

1. given - $\Delta x^{(k)}$ @ $f(x^{(k)})$

$$x^{(k)} \in \text{dom } f$$

$$\alpha \in (0, 0.5)$$

$$\beta \in (0, 1)$$

2. $t := 1$

3. (while)

$$f(x + t\Delta x) > f(x) + \alpha t \nabla f(x)^T \Delta x$$

(do)

$$t := \beta t$$

Gradient Descent Method

1. (given) $x^{(0)} \in \text{dom } f$

2. (repeat)

a. $\Delta x^{(k)} = -\nabla f(x^{(k)})$

b. backtracking line search get $t^{(k)}$

c. $x^{(k+1)} = x^{(k)} + \Delta x^{(k)}$

3. (until)

$$\|\nabla f(x)\|_2 \leq \eta$$

Newton's Method

1. (given) $x^{(0)} \in \text{dom } f$

2. (repeat)

a. $\Delta x^{(k)} := -\nabla^2 f(x^{(k)})^{-1} \nabla f(x^{(k)})$

$$x^{(k+1)} := x^{(k)} + \nabla f(x^{(k)})^T \nabla^2 f(x^{(k)})^{-1} \nabla f(x^{(k)})$$

b. (until) $\frac{\|x^{(k+1)} - x^{(k)}\|_2}{2} \leq \epsilon$

c. backtracking line search get $t^{(k)}$

d. $x^{(k+1)} = x^{(k)} + t^{(k)} \Delta x^{(k)}$

Newton's method

equality constrained

1. (given) $x^{(0)} \in \text{dom } f$
 $Ax = b$

2. (repeat)

a. $\Delta x_{nt}^{(k)} := \begin{bmatrix} \nabla^2 f(x) & A^T \\ A & 0 \end{bmatrix}^{-1} \begin{bmatrix} -\nabla f(x) \\ w \end{bmatrix} \quad (1:n, :)$
 $\lambda^{(k)^2} := \nabla f(x^{(k)})^T \nabla^2 f(x^{(k)})^{-1} \nabla f(x^{(k)})$

b. (until) $\lambda^{(k)^2}/2 \leq \epsilon$

c. backtracking line search get
 $t^{(k)}$

d. $x^{(k+1)} = x^{(k)} + t^{(k)} \Delta x^{(k)}$

Newton's method (infeasible start)

$$\begin{bmatrix} \nabla^2 f(x) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} \Delta x_{pd} \\ \Delta v_{pd} \end{bmatrix} = - \begin{bmatrix} r_{\text{dual}} \\ r_{\text{pri}} \end{bmatrix} = - \begin{bmatrix} \nabla f(x) + A^T v \\ Ax - b \end{bmatrix} \quad \textcircled{\#}$$

1. (given) $x^{(0)} \in \text{dom } f$

v

$G > 0$

$\alpha \in (0, 0.5)$

$\beta \in (0, 1)$

2. (repeat)

a. $\Delta x_{nt}^{(k)} := \text{from } \textcircled{\#}$

$\Delta v_{nt}^{(k)} := \text{from } \textcircled{\#}$

b. backtracking line search on $\|r\|_2$

1. $t := 1$

2. while $\|r(x + t \Delta x_{nt}, v + t \Delta v_{nt})\|_2 > (1-\alpha t) \|r(x, v)\|_2$

$t := \beta t$

c. $x := x + t \Delta x_{nt}$

$v := v + t \Delta v_{nt}$

3. (until)

$Ax = b$

$\|r(x, v)\|_2 \leq \epsilon$

Barrier method with logarithm

inequality constrained

original problem:

$$\begin{aligned} & \text{minimize } f_0(x) \\ & \text{subject to } f_i(x) \leq 0 \quad i=1, \dots, m \\ & Ax = b \end{aligned}$$



new problem:

$$\text{minimize } T f_0(x) + \phi(x)$$

$$\text{subject to } Ax = b$$

$$\phi(x) = -\sum_{i=1}^m \log(-f_i(x))$$

1. (given) $x^{(0)}$ (feasible)

$$T^{(0)} := T^{(0)} > 0$$

$$\mu > 1$$

$$\epsilon > 0$$

2. (repeat)

1. solve $x^*(\tau)$ of $T f_0 + \phi$ subject to $Ax = b$
with " τ "

$$2. x := x^*(\tau)$$

3. (until)

$$\frac{\|x\|}{\tau} < \epsilon$$

$$4. \tau := \mu \tau$$

$$\nabla \phi(x) = \sum_{i=1}^m -\frac{1}{-f_i(x)} \nabla f_i(x)$$

$$\nabla^2 \phi(x) = \sum_{i=1}^m \frac{1}{f_i(x)^2} \nabla f_i(x) \nabla f_i(x)^T + \sum_{i=1}^m \frac{1}{-f_i(x)} \nabla^2 f_i(x)$$

Phase I via infeasible start Newton method

inequality constrained
w/ infeasible
start.

original problem:

$$\text{minimize } f_0(x)$$

$$\text{subject to } f_i(x) \leq 0 \quad i=1, \dots, m$$

$$Ax = b$$



$$\text{minimize } f_0(x)$$

$$\text{subject to } f_i(x) \leq S, \quad i=1, \dots, m$$

$$Ax = b$$

$$S = 0$$



$$\text{minimize } t^{10} f_0(x) - \sum_{i=1}^m \log(S - f_i(x))$$

$$\text{subject to } Ax = b$$

$$S = 0$$