

# Routing Protocols for Ad-Hoc Networks

Ad-hoc On-Demand Distance  
Vector Routing

&

DSR: The Dynamic Source Routing  
Protocol for Multi-Hop Wireless  
Ad Hoc Networks

# Outline

- ❑ **Ad-Hoc networks**
- ❑ **Ad-hoc routing algorithms**
- ❑ **Ad-Hoc on-demand Distance Vector Routing (AODV)**
- ❑ **Dynamic Source Routing (DSR)**
- ❑ **Comparison of AODV and DSR**

# Ad Hoc Networks

Wireless networks can be divided in two fundamental categories:

- ❑ **Infrastructure-based**

Wireless clients connecting to a base-station (APs, Cell Towers) that provides all the traditional network services (routing, address assignment)

- ❑ **Infrastructure-less**

The clients themselves must provide all the traditional services to each other

# Ad Hoc Networks

Ad-hoc networks main features:

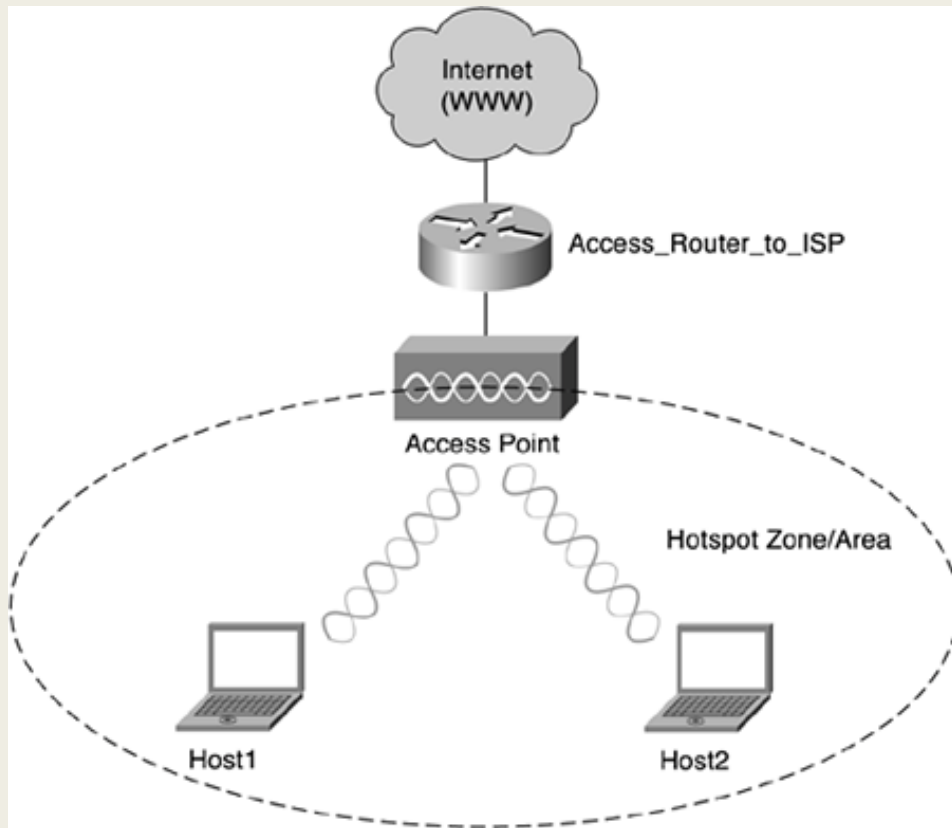
- ❑ Decentralized
- ❑ Do not rely on preexisting infrastructure
- ❑ Each node participates in routing by forwarding data to neighbor nodes
- ❑ Fast network topology changes due to nodes' movement

# Ad Hoc Networks

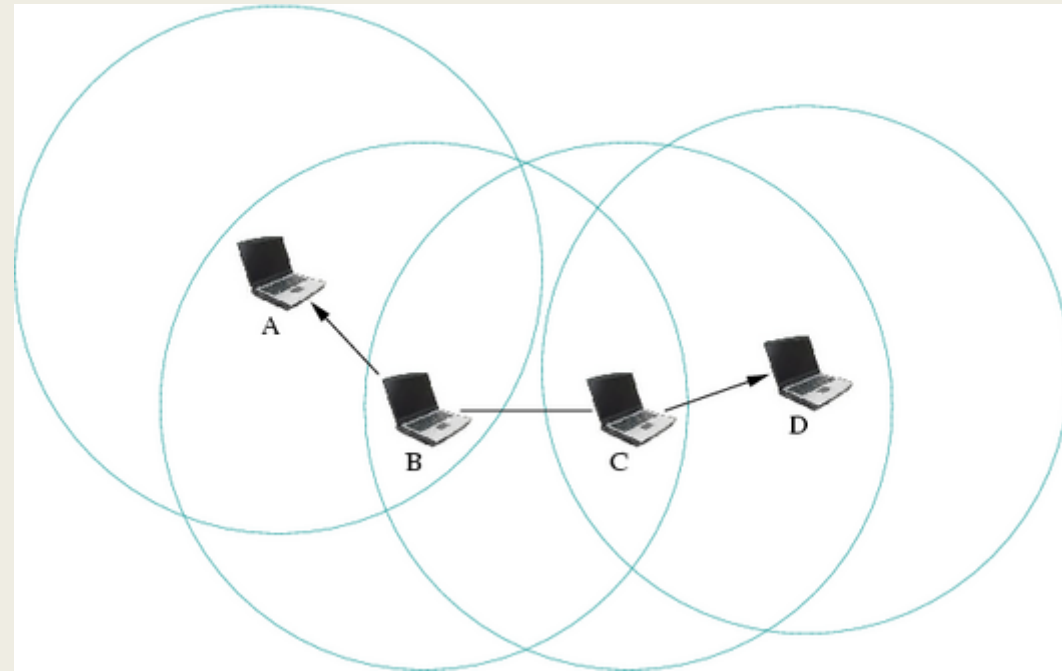
Why do we need ad-hoc networks?

- ❑ More laptop users
  - ❑ More smartphones users (e.g.. Android phones, iPhones)
  - ❑ More devices with Wi-Fi-support (e.g.. televisions, hi-fi, home-theaters, media servers etc.)
  - ❑ Moving users, vehicles, etc.
  - ❑ Outdoors places
- ✓ In all these occasions there is no centralized infrastructure (such APs)
- ✓ So ad-hoc network is a necessity

# Ad Hoc Networks



An infrastructure wireless network



An Ad-hoc network

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# Ad-hoc routing algorithms

Hottest routing algorithm categories:

## ❑ Pro-active (table-driven) routing

Maintains fresh lists of destinations & their routes by periodically distributing routing tables

Disadvantages:

1. Respective amount of data for maintenance
2. Slow reaction on restructuring and failures  
(e.g. OSLR, DSDV)

## ❑ Reactive (on-demand) routing

On demand route discovery by flooding the network with Route Request packets

Disadvantages:

1. High latency time in route finding
2. Flooding can lead to network clogging  
(e.g. AODV, DSR)



# Ad-hoc routing algorithms

Discuss and comparison

1. Ad-Hoc on-demand Distance Vector Routing (AODV)
2. Dynamic Source Routing (DSR)

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- Ad-Hoc networks
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- **Ad-Hoc on-demand Distance Vector Routing (AODV)**
  - **General info**
  - **Path Discovery**
  - **Path Maintenance**
  - **Local Connectivity Maintenance**
  - **Conclusion**
- **Dynamic Source Routing (DSR)**
- **Comparison of AODV and DSR**

# (AODV) General info

- ❑ Reactive algorithms like AODV create routes on-demand. They must however, reduce as much as possible the acquisition time
- ❑ We could largely eliminate the need of periodically system-wide broadcasts
- ❑ AODV uses symmetric links between neighboring nodes. It does not attempt to follow paths between nodes when one of the nodes can not hear the other one

# (AODV) General info

- Nodes that have not participate yet in any packet exchange (inactive nodes), they do not maintain routing information
- They do not participate in any periodic routing table exchanges

# (AODV) General info

- Each node can become aware of other nodes in its neighborhood by using **local** broadcasts known as hello messages
  
- neighbor routing tables organized to :
  1. optimize response time to local movements
  2. provide quick response time for new routes requests

# (AODV) General info

AODV main features:

- Broadcast route discovery mechanism
- Bandwidth efficiently (small header information)
- Responsive to changes in network topology
- Loop free routing

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# (AODV) Path Discovery

- Initiated when a source node needs to communicate with another node for which it has no routing info
- Every node maintains two counters:
  - node\_sequence\_number
  - broadcast\_id
- The source node broadcast to the neighbors a **route request packet** (called RREQ)



# (AODV) Path Discovery

- RREQ structure  
    <src\_addr, src\_sequence\_#, broadcast\_id, dest\_addr,  
        dest\_sequence\_#, hop\_cnt>
- **src\_addr** and **broadcast\_id** uniquely identifies a RREQ
- broadcast\_id is incremented whenever source node issues a RREQ
- Each neighbor either satisfy the RREQ, by sending back a routing reply (RREP), or rebroadcast the RREQ to its own neighbors after increasing the hop\_count by one.

# (AODV) Path Discovery

- If a node receives a RREQ that has the same `<src_addr, broadcast_id>` with a previous RREQ it drops it immediately
- If a node cannot satisfy the RREQ, stores:
  - Destination IP
  - Source IP
  - `broadcast_id`
  - Expiration time (used for reverse path process)
  - `src_sequence_#`

# (AODV) Path Discovery

## 1. Reverse Path Setup

- In each RREQ there are:
  - `src_sequence_#`
  - the last `dest_sequence_#`
- `src_sequence_#` used to maintain freshness information about the reverse route to the source
- `dest_sequence_#` indicates how fresh a route must be, before it can be accepted by the source

# (AODV) Path Discovery

## 1. Reverse Path Setup (continue)

- As RREQ travels from source to many destinations, it automatically sets up the reverse path, from all nodes back to the source.

But how does it work?

- Each node records the address of the neighbor from which it received the first copy of the RREQ
- These entries are maintained for at least enough time, for the RREQ to traverse the network and produce a reply

# (AODV) Path Discovery

## 1. Reverse Path Setup (continue)

RREQ reached destination  
Reversed path is fully set up  
From which RREP can travel  
back to S

S Source node  
D Destination node  
W — Z Neighbor nodes

Z, V, U can not satisfy RREQ  
i. Set up reverse path  
ii. Rebroadcast RREQ to  
neighbors

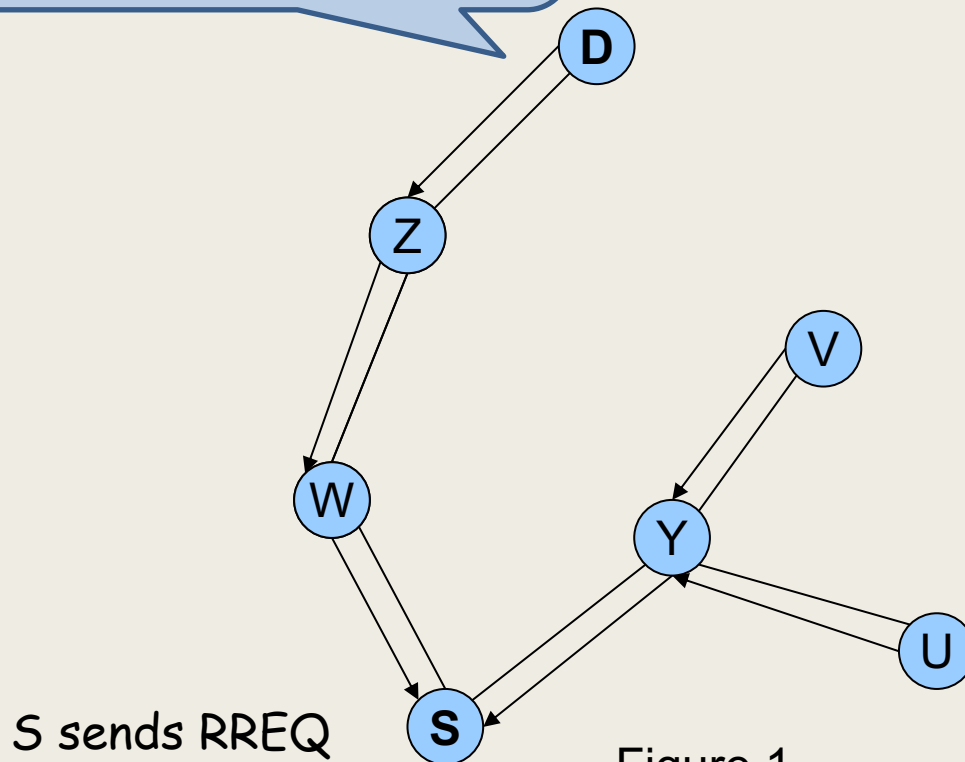


Figure 1

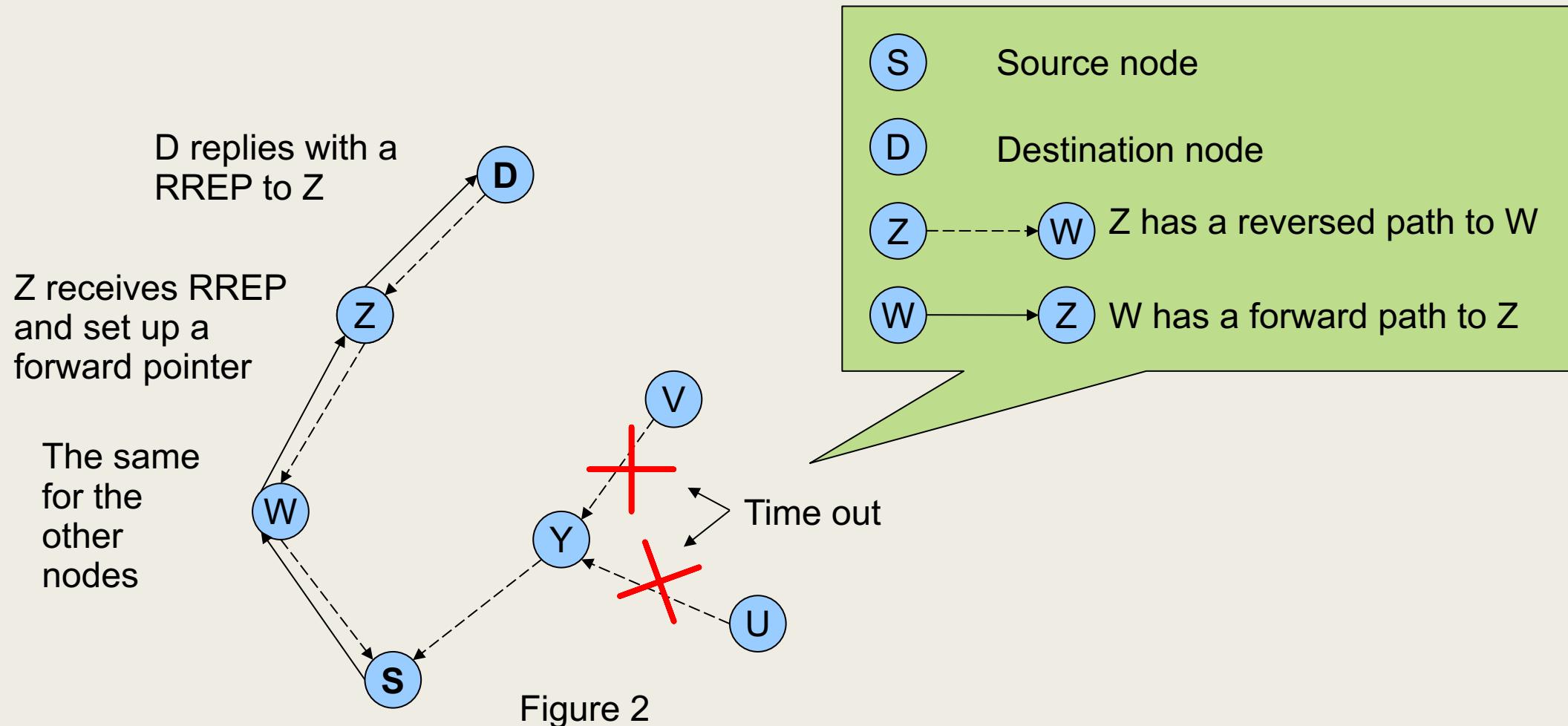
# (AODV) Path Discovery

## 2. Forward Path Setup

- A node receiving a RREP propagates the first RREP for a given source towards the source (using the reverse path that has already established)
- Nodes that are not in the path determined by the RREP will time out after 3000 ms and will delete the reverse pointers

# (AODV) Path Discovery

## 2. Forward Path Setup (continue)



# (AODV) Path Discovery

## 2. Forward Path Setup (Conclusion)

- Minimum number of RREPs towards source
- The source can begin data transmission as soon as the first RREP received and update later its routing information if it learns of a better route



# Outline

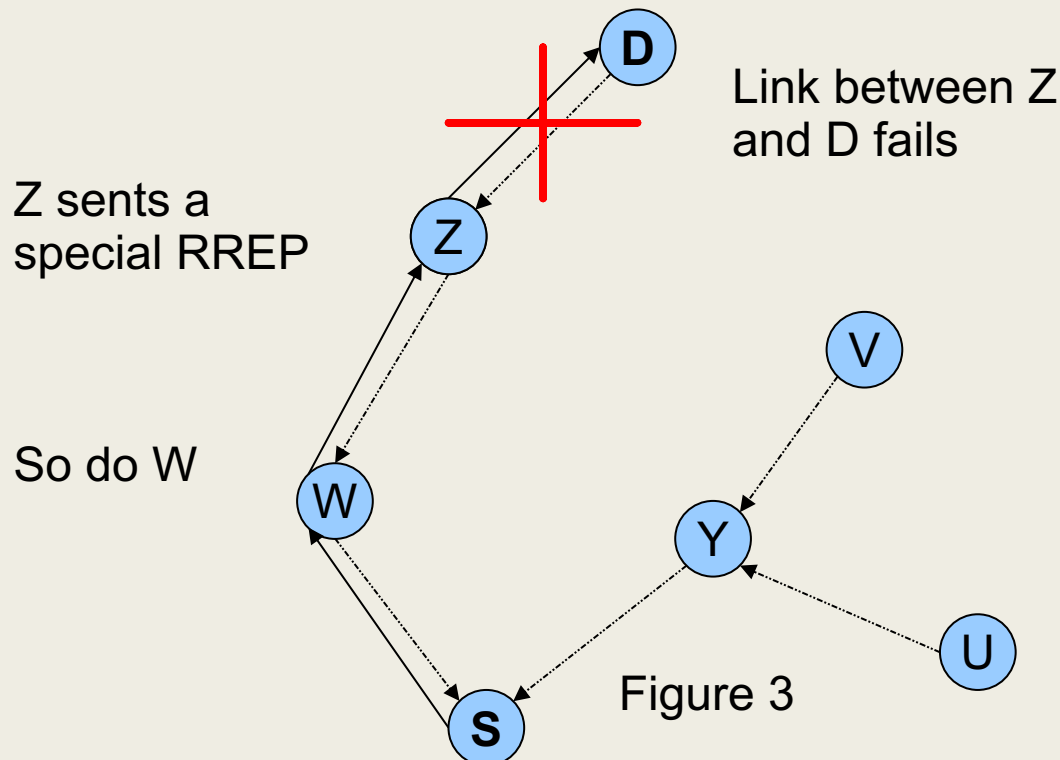
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# (AODV) Path Maintenance

- Movement of nodes not lying along an active path does NOT affect the route to that path's destination
- If the source node moves, it can simply re-initiate the route discovery procedure
- If the destination or some intermediate node moves, a special RREP is sent to the affected nodes
- To find out nodes movements periodic hello messages can be used, or (LLACKS) link-layer acknowledgments (far less latency)

# (AODV) Path Maintenance

- When a node is unreachable the special RREP that is sent back towards the source, contains a new sequence number and hop count of  $\infty$



So now source must find a new path. To do that, it sends a RREQ with a new greater sequence number

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# (AODV) Local Connectivity Maintenance

- Nodes learn of their neighbors in one or two ways:
  1. Whenever a node receives a broadcast from a neighbor it update its local connectivity information about this neighbor
  2. If a neighbor has not sent any packets within hello\_interval it broadcasts a hello message, containing its identity and its sequence number

# (AODV) Local Connectivity Maintenance

How hello messages work:

- Hello messages do not broadcasted outside the neighborhood because the contain a TTL (time to leave) value of 1.
- Neighbors that receive the hello message update their local connectivity information to the node that have broadcasted the hello message

# (AODV) Local Connectivity Maintenance

How hello messages work: (continue)

- Receiving a hello from a new neighbor, or failing to receive `allowed_hello_loss` (typically 2) consecutive hello messages from a node previously in the neighborhood, indicates that the local connectivity has changed

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# (AODV) Conclusion

## AODV main features:

- ❑ Nodes store only the routes they need
- ❑ Need for broadcast is minimized
- ❑ Reduces memory requirements and needless duplications
- ❑ Quick response to link breakage in active routes
- ❑ Loop-free routes maintained by use of destination sequence numbers
- ❑ Scalable to large populations of nodes

# AODV



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# (DSR) General

Two main mechanisms that work together to allow the discovery and maintainance of source routes:

- Route discovery
- Route maintainance

# (DSR) General

## Route discovery:

- Is the mechanism by which a source node  $S$ , obtains a route to a destination  $D$
- Used only when  $S$  attempt to send a packet to  $D$  and does not already knows a route to  $D$

# (DSR) General

## Route maintenance:

- Is the mechanism by which source node  $S$  is able to detect if the network topology has changed and can no longer use its route to  $D$
- If  $S$  knows another route to  $D$ , use it
- Else invoke route discovery process again to find a new route
- Used only when  $S$  wants to send a packet to  $D$

# (DSR) General

- Each mechanism operate entirely on demand
- DSR requires no periodic packets of any kind at any level
- Uni-directional and asymmetric routes support  
(e.g. send a packet to a node D through a route and receive a packet D from another route)

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# (DSR) Basic Route Discovery

When S wants to sent a packet to D:

- it places in the header of the packet a source route giving the sequence of hops that the packet should follow on its way to D
- S obtains a suitable source route by searching its route table
- If no route found for D, S initiate the Route Discovery protocol to dynamically find a new route to D

# (DSR) Basic Route Discovery

## Sender

- ❑ Broadcasts a Route Request Packet (RREQ)
- ❑ RREQ contains a unique **Request ID** and the address of the sender

## Receiver

- ❑ If this node is the destination node, or has route to the destination send a Route Reply packet (RREP)
- ❑ Else if is the source, drop the packet
- ❑ Else if is already in the RREQ's route table, drop the packet
- ❑ Else append the node address in the RREQ's route table and broadcast the updated RREQ

# (DSR) Basic Route Discovery

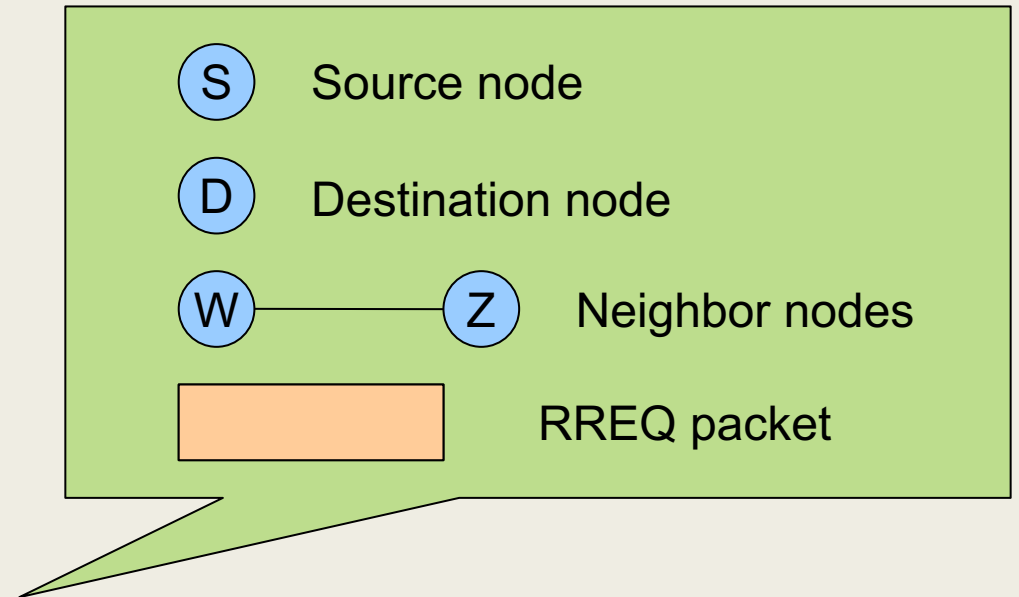
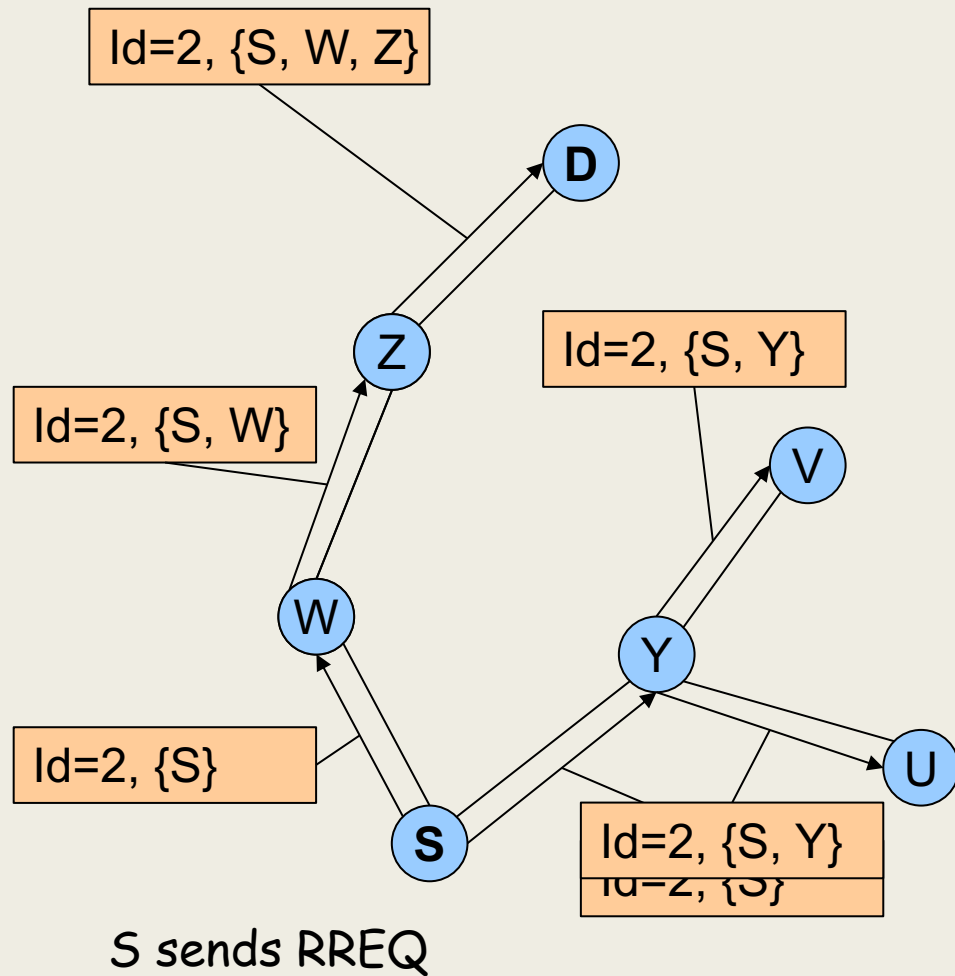


Figure 4

# (DSR) Basic Route Discovery

When a RREQ reaches the destination node, a RREP must be sent back to source

The destination node:

- Examine its own Route Cache for a route back to source
  - If found, it use this route to send back the RREP
  - Else, the destination node starts a new Route Discovery process to find a route towards source node
- ✓ In protocols that require bi-directional links like 802.11, the reversed route list of the RREQ packet can be used, in order to avoid the second Route Discovery

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# (DSR) Basic Route Maintenance

## Each node transmitting a packet:

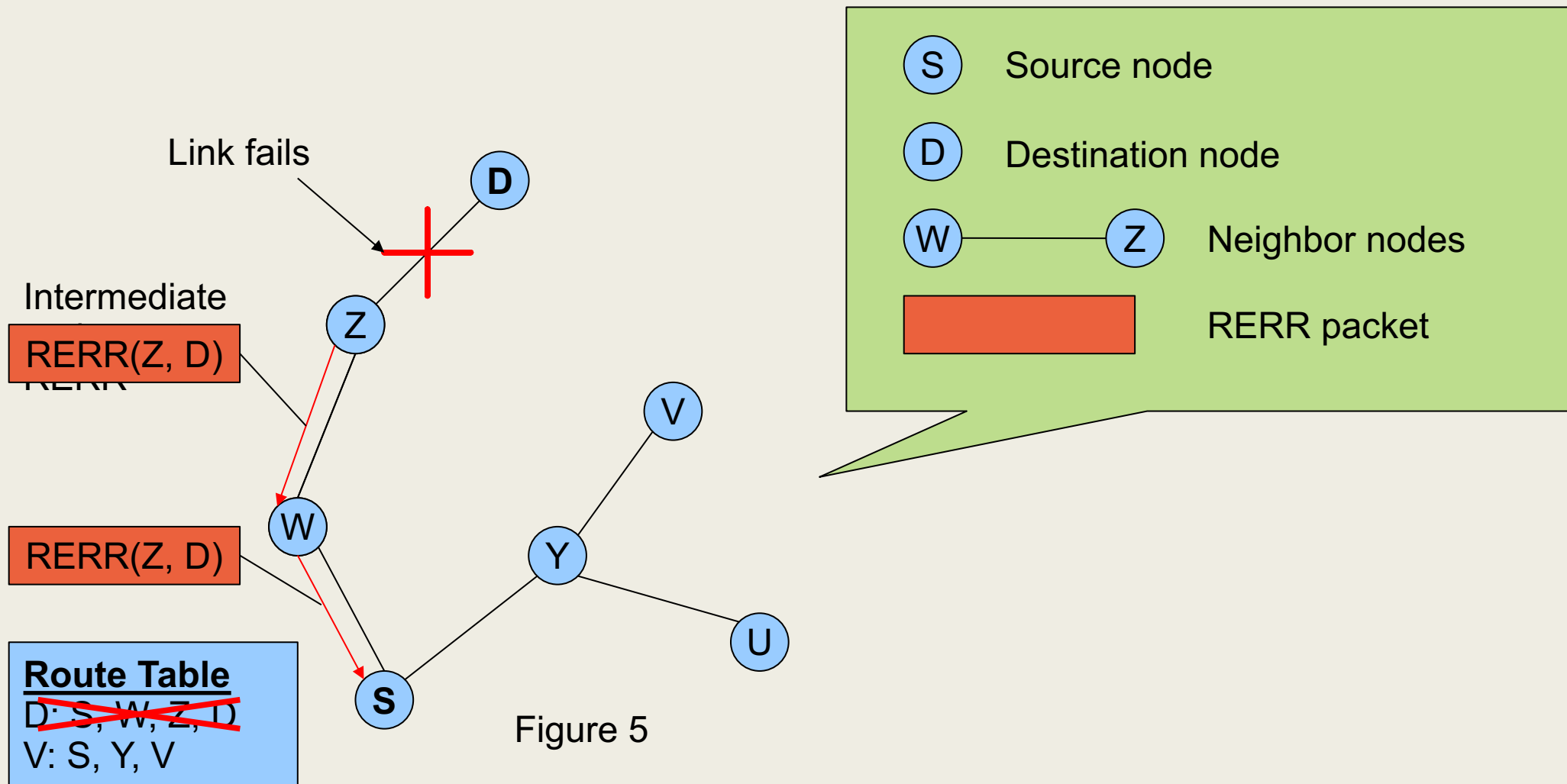
- is responsible for confirming that the packet has been received by the next hop along the source route
- The confirmation it is done with a standard part of MAC layer (e.g. Link-level ACKs in 802.11)
- If none exists, a DSR-specific software takes the responsibility to sent back an ACK
- When retransmissions of a packet in a node reach a maximum number, a Route Error Packet (RERR) is sent from the node back to the source, identifying the broken link

# (DSR) Basic Route Maintenance

## The source:

- ❑ Removes from the routing table the broken route
- ❑ Retransmission of the original packet is a function of upper layers (e.g. TCP)
- ❑ It searches the routing table for another route, or start a new Route Discovery process

# (DSR) Basic Route Maintenance





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# (DSR) Conclusion

- ❑ Excellent performance for routing in multi-hop wireless ad hoc networks
- ❑ Very low routing overhead even with continuous rapid motion, which scales to :
  1. zero when nodes are stationary
  2. the affected routes when nodes are moving
- ❑ Completely self-organized & self-configuring network
- ❑ Entirely on-demand operation. No periodic activity of any kind at any level

# DSR



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# Comparison of AODV and DSR

**Main common features:**

- ❑ On-demand route requesting
- ❑ Route discovery based on requesting and replying control packets
- ❑ Broadcast route discovery mechanism

# Comparison of AODV and DSR

**Main common features: (continue)**

- ❑ Route information is stored in all intermediate nodes along the established path
- ❑ Inform source node for a broken links
- ❑ Loop-free routing

# Comparison of AODV and DSR

## Main differences:

- ❑ DSR can handle uni and bi-directional links, AODV uses only bi-directional
- ❑ In DSR, using a single RREQ - RREP cycle, source and intermediate nodes can learn routes to other nodes on the route
- ❑ DSR maintains many alternate routes to the destination, instead of AODV that maintains at most one entry per destination

# Comparison of AODV and DSR

## Main differences: (continue)

- ❑ DSR doesn't contain any explicit mechanism to expire stale routes in the cache , In AODV if a routing table entry is not recently used , the entry is expired
- ❑ DSR can't prefer "fresher" routes when faced multiple choices for routes. In contrast, AODV always choose the fresher route (based on destination sequence numbers)



# Comparison of AODV and DSR

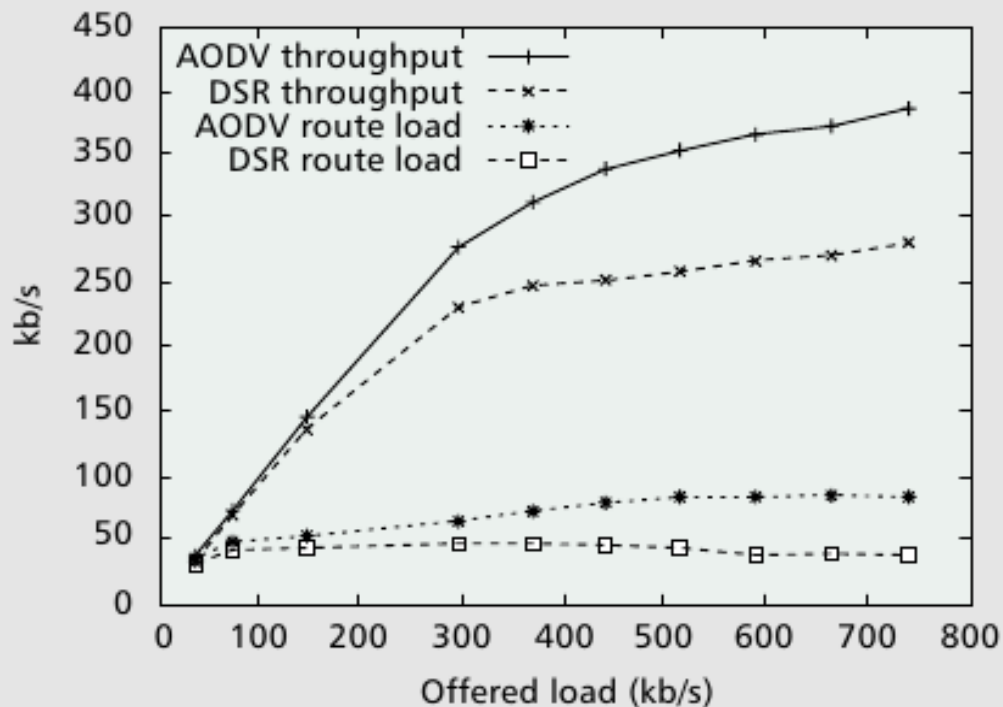
## Main differences: (continue)

- ❑ DSR's RREQ has variable length depending on the nodes that the packet has traveled. AODV's RREQ size is constant
- ❑ As a result DSR's header overhead may increase as more nodes become active, so we expect that AODV throughput in those scenarios to be better

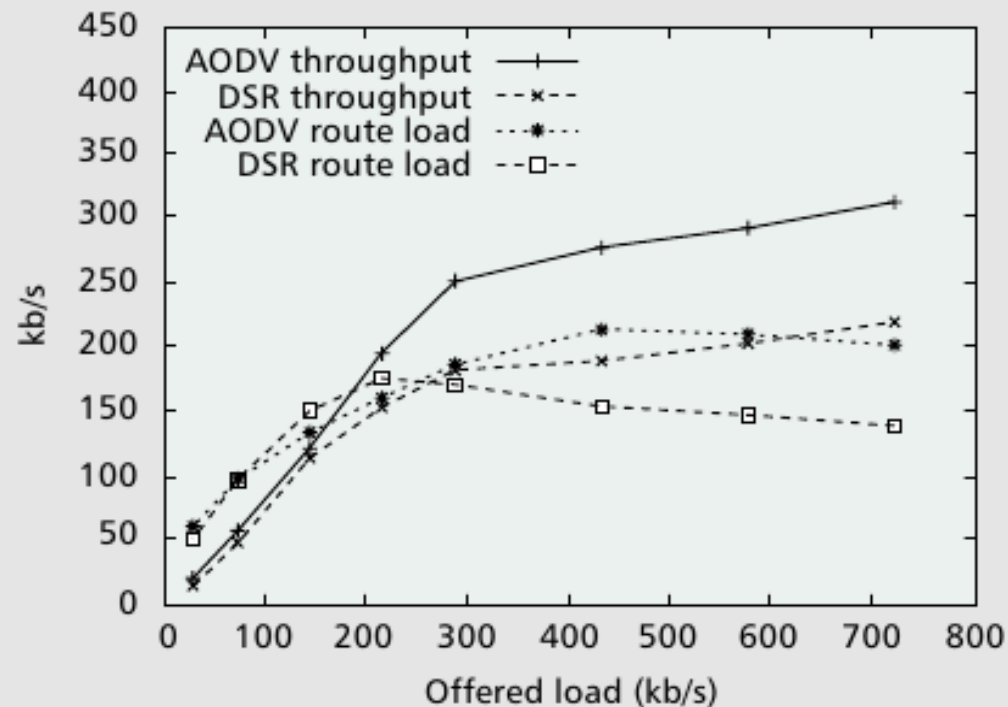
# Comparison of AODV and DSR

Test bench set up:

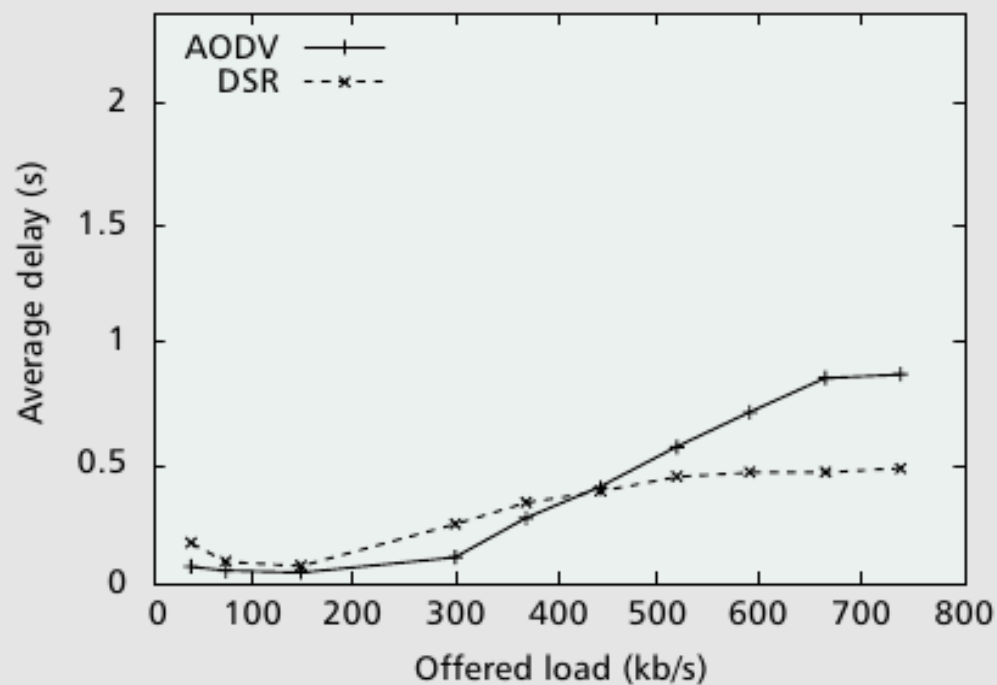
- ❑ 100 nodes, some of them as sources
- ❑ Nominal bit rate of 2 Mb/s
- ❑ Nominal node range of 250 m
- ❑ Continuously moving nodes



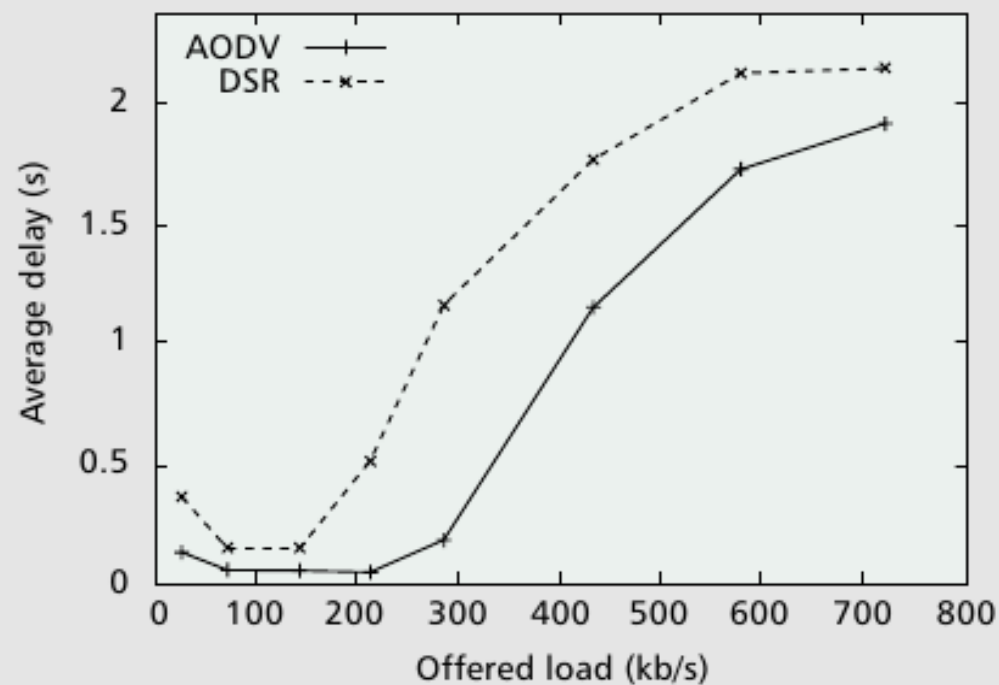
(a) 10 sources - throughput, routing load



(b) 40 sources - throughput, routing load



(c) 10 sources - delay



(d) 40 sources - delay

# Comparison of AODV and DSR

Performance metrics	DSR	AODV
Packets delivered /Packets sent (%)	56.88	83.66
Average delay (s)	1.36	0.26

Routing Packets	DSR	AODV
Route requests	37774	228094
Route replies	82710	17753
Route errors	26591	9808
Total	147075	255655

*Application and routing statistics for an example scenario for a network of 100 nodes with continuous mobility and 40 sources*

# Conclusion

- ❑ DSR outperforms AODV in less stressful situations (i.e., smaller number of nodes and lower load and/or mobility)
- ❑ AODV outperforms DSR in more stressful situations (e.g., more load, higher mobility)
- ❑ DSR commonly generates less routing load than AODV
- ❑ Poor delay and throughput of DSR due to lack of any mechanism to expire stale routes or determine the freshness of routes