

# **Software Verification & Validation**

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Graph Coverage

# Acknowledgement

- Some slides in this lecture are adapted from
  - **Paul Ammann and Jeff Offutt's** slides for their textbook **“Introduction to Software Testing”**. Cambridge University Press, 2008.
  - **Lee Copeland's “A Practitioner's Guide to Software Test Design”**. Artech House, 2004.

# Graphs for Test Design

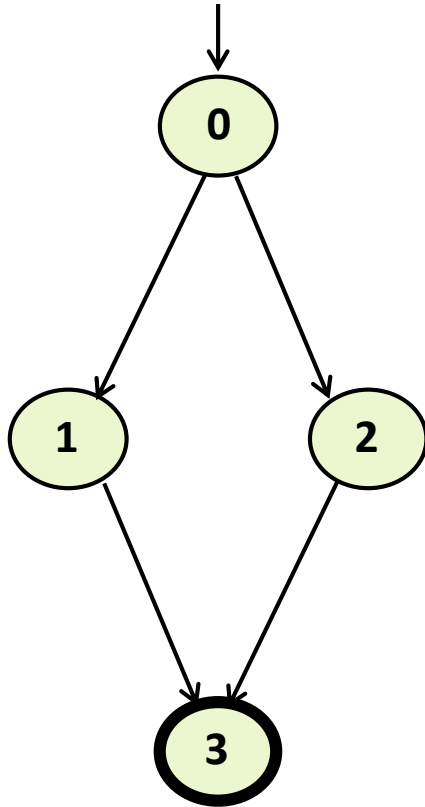
- Graphs are the most commonly used structure for testing
- Graphs can come from many sources
  - Control flow graphs
  - Design structure
  - State machines
  - Use cases
- Tests usually are intended to “cover” the graph in some way

# Definition of a Graph

For our purpose, we will be using a directed graph with one or more initial nodes and one or more final nodes. Precisely, we define a graph to be a structure  $(N, N_o, N_f, E)$  where

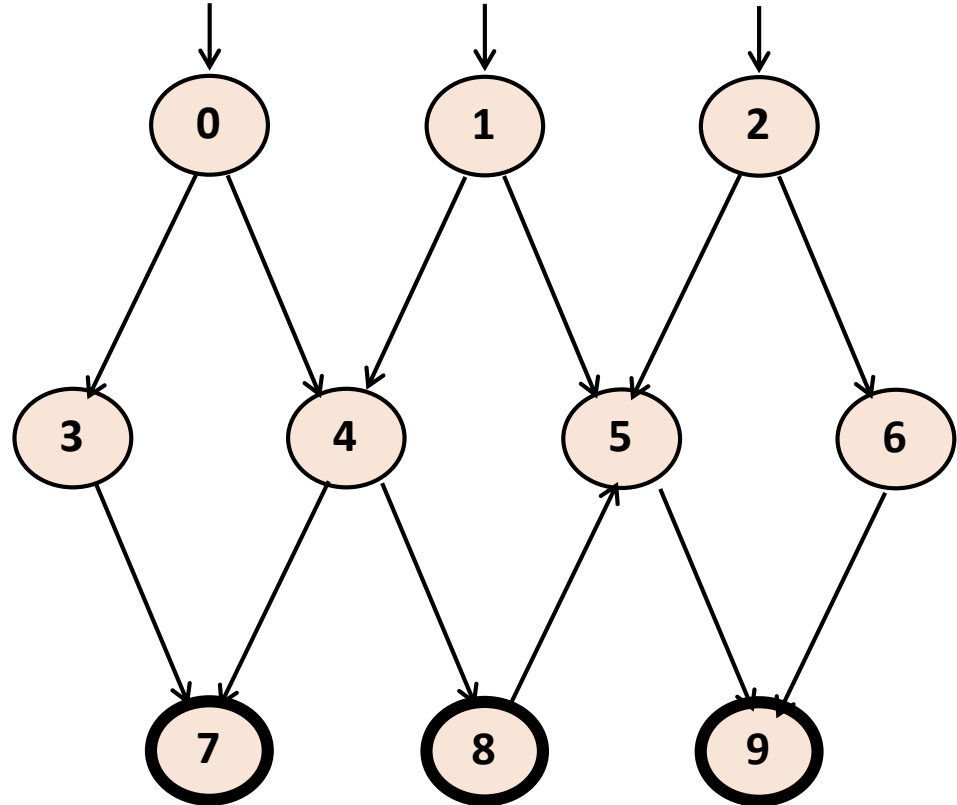
- $N$  is a non-empty set of nodes,
- $N_o \subseteq N$  is a non-empty set of initial nodes,
- $N_f \subseteq N$  is a non-empty set of final nodes,
- $E \subseteq N \times N$  is a set of edges, each edge from one node to another
  - If  $(n_i, n_j)$  is in  $E$ , then  $n_i$  is called a predecessor of  $n_j$  and  $n_j$  is called a successor of  $n_i$

# Three Example Graphs



$$N_0 = \{ 0 \}$$

$$N_f = \{ 3 \}$$

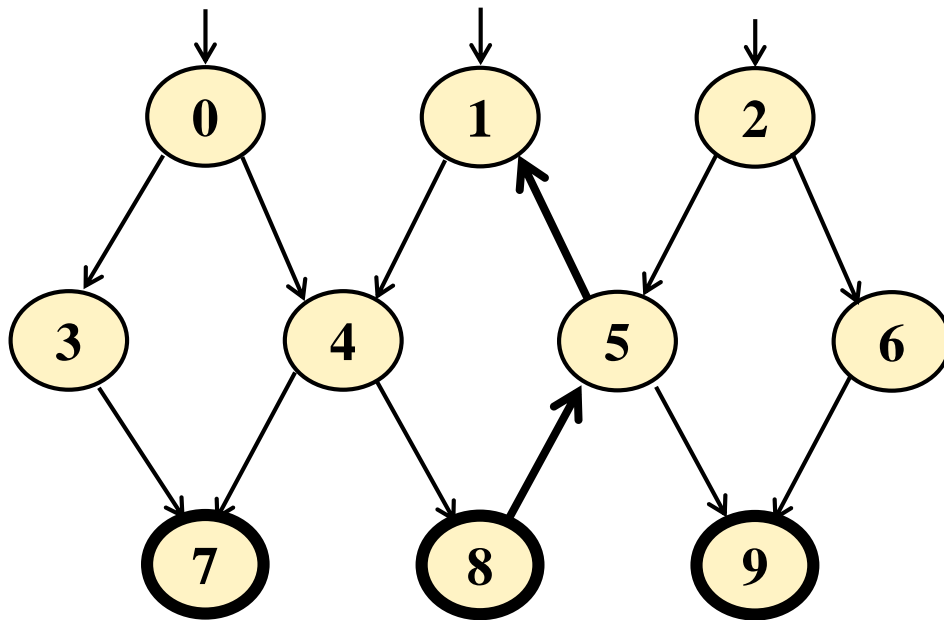


$$N_0 = \{ 0, 1, 2 \}$$

$$N_f = \{ 7, 8, 9 \}$$

# Paths in Graphs

- Path : A sequence of nodes –  $[n_1, n_2, \dots, n_M]$  such that each pair of adjacent nodes  $(n_i, n_{i+1})$  is an edge.
- Length : The number of edges
  - A single node is a path of length 0
- Subpath : A subsequence of consecutive nodes in  $p$  is a subpath of  $p$
- Reach( $n$ ) : The set of nodes that can be reached from  $n$
- Reach( $A$ ) : The set of nodes that can be reached from some node in set  $A$



## Paths

[ 0, 3, 7 ]

[ 1, 4, 8, 5, 1 ]

[ 2, 6, 9 ]

Reach (0) = { 0, 3, 4, 7, 8, 5, 1, 9 }

Reach ({0, 2}) = G

# Visiting and Touring

- Visit : A path  $p$  visits node  $n$  iff  $n$  is in  $p$   
A path  $p$  visits edge  $e$  iff  $e$  is in  $p$
- Tour : A path  $p$  tours path  $q$  iff  $q$  is a subpath of  $p$

**For example, the path [ 0, 1, 3, 4, 6 ]**

**visits nodes 0, 1, 3, 4, 6**

**visits edges (0, 1), (1, 3), (3, 4), (4, 6)**

**And tours paths [0, 1, 3], [1, 3, 4], [3, 4, 6], [0, 1, 3, 4], [1, 3, 4, 6]**

# Tests and Test Paths

- A test path in a graph is a path in the graph following the execution of some test case.
- path ( $t$ ) : The test path executed by test  $t$
- path ( $T$ ) : The set of test paths executed by the set of tests  $T$
- In a deterministic program, each test executes **one and only one** test path
- A location in a graph (node or edge) can be reached from another location if there is a sequence of edges from the first location to the second
  - Syntactic reach : A subpath exists in the graph
  - Semantic reach : A test exists that can execute that subpath



# Testing and Covering Graphs

- We use graphs in testing as follows :
  - Developing a model of the software as a graph
  - Requiring tests to visit or tour specific sets of nodes, edges or subpaths
- Structural Coverage Criteria : Defined on a graph just in terms of nodes and edges
- Data Flow Coverage Criteria : Requires a graph to be annotated with references to variables (*this will be studied in subsequent lectures*)

# Node and Edge Coverage

- The first (and simplest) two criteria require that each node and edge in a graph be executed

**Node Coverage Criterion (NC)**: Test suite  $T$  satisfies *node coverage* on graph  $G$  iff, for every node  $n$ , some path in  $path(T)$  visits  $n$ .

Given a test suite  $T$ , the degree of node coverage of  $T$  can be calculated form:

$$C_{NC} = \frac{\text{Number of nodes visited by some test path}}{\text{Number of nodes}} \times 100\%$$

# Node and Edge Coverage

- Edge coverage is slightly stronger than node coverage

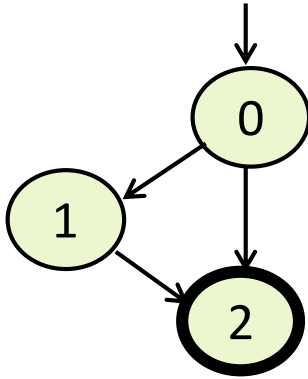
**Edge Coverage Criterion (EC)** : Test suite  $T$  satisfies *edge coverage* on graph  $G$  iff, for every edge  $(m, n)$ , some path in  $path(T)$  visits  $(m, n)$ .

Given a test suite  $T$ , the degree of edge coverage of  $T$  can be calculated form:

$$C_{EC} = \frac{\text{Number of edges visited by some test path}}{\text{Number of edges}} \times 100\%$$

# Node and Edge Coverage

- NC and EC are only different when there is an edge and another subpath between a pair of nodes



**Node Coverage** : Nodes = { 0, 1, 2 }

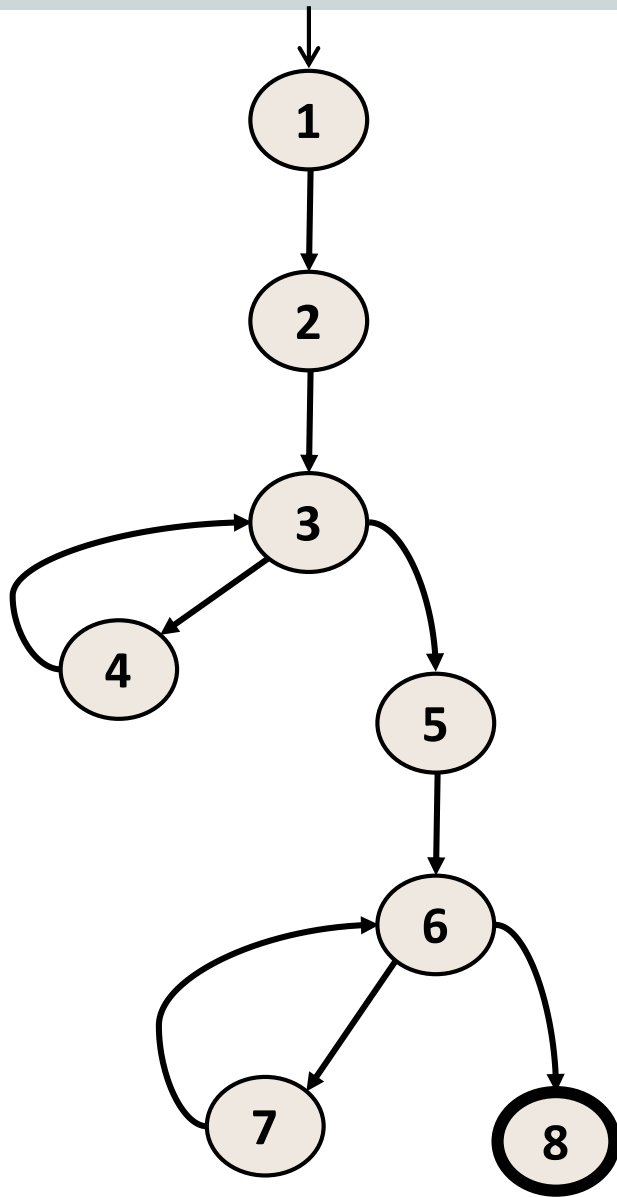
Test Path = [ 0, 1, 2 ]

**Edge Coverage** : Edges = { (0,1), (0, 2), (1, 2) }

Test Paths = [ 0, 1, 2 ]

[ 0, 2 ]

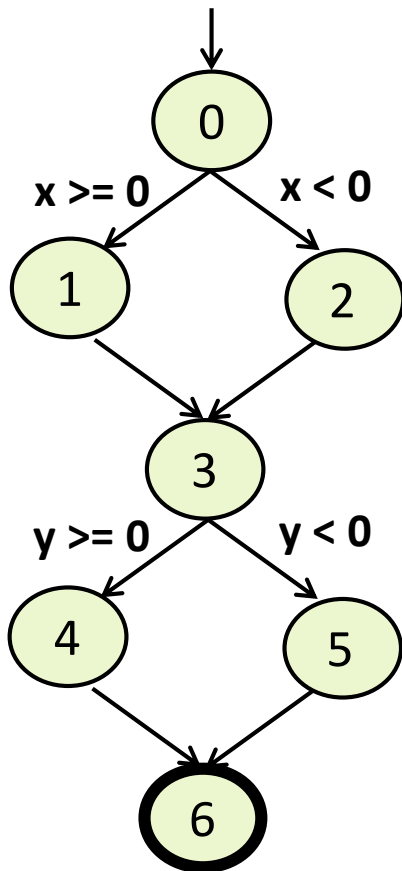
# Control Flow TRs and Test Paths – EC



Edge Coverage	
Edges	Test Path
A. (1, 2)	[ 1, 2, 3, 4, 3, 5, 6, 7, 6, 8 ]
B. (2, 3)	
C. (3, 4)	
D. (3, 5)	
E. (4, 3)	
F. (5, 6)	
G. (6, 7)	
H. (6, 8)	
I. (7, 6)	

# Limitation of Edge Coverage

- **Edge coverage** requires that all alternatives in each choice node in the graph must be chosen in some test path.
- However, in graph with consecutive choice nodes, edge coverage criteria does **not** require a test path for each **combination of alternatives** of the choice nodes.



- For example, in this graph, nodes **0** and **3** are choice nodes.
- Only two paths are required to cover every edge. For example, **[0, 1, 3, 4, 6]** and **[0, 2, 3, 5, 6]**.
- The first path covers the case where  $x \geq 0$  and  $y \geq 0$ . The second path covers the case where  $x < 0$  and  $y < 0$ .
- But no paths cover the case where  $x \geq 0$  and  $y < 0$  nor the case where  $x < 0$  and  $y \geq 0$ .
- To cover these cases, we need to add more paths. For example, **[0, 1, 3, 5, 6]** and **[0, 2, 3, 4, 6]**.
- We will look at some coverage criteria which are more comprehensive than edge coverage.

# Edge-Pair Coverage

- Edge-pair coverage is a stronger version than edge coverage

**Edge-Pair Coverage Criterion (EC)** : Test suite  $T$  satisfies *edge-pair coverage* on graph  $G$  iff, for every path  $p$  of length up to 2, some path in  $path(T)$  tours  $p$ .

Given a test suite  $T$ , the degree of edge coverage of  $T$  can be calculated form:

$$C_{EC} = \frac{\text{Number of paths of length up to 2 toured by some test path}}{\text{Number of paths of length up to 2}} \times 100\%$$

# Complete Path Coverage

- The Complete Path Coverage requires that every possible path in the graph is toured by some test path.

**Complete Path Coverage (CPC)** : Test suite  $T$  satisfies *complete path coverage* on graph  $G$  iff, for every path  $q$ , some path in  $path(T)$  tours  $q$ .

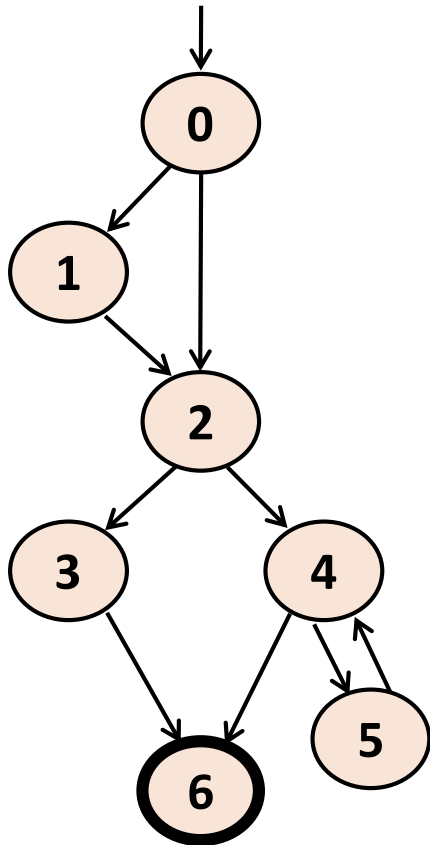
Given a test suite  $T$ , the degree of complete path coverage of  $T$  can be calculated form:

$$C_{CPC} = \frac{\text{Number of paths toured by some test path}}{\text{Number of paths}} \times 100\%$$

- Unfortunately, this is **impossible** if the graph has a loop because such graph will have infinitely many paths.



# Structural Coverage Example



## Node Coverage

**Nodes** = { 0, 1, 2, 3, 4, 5, 6 }

**Test Paths:** [ 0, 1, 2, 3, 6 ] [ 0, 1, 2, 4, 5, 4, 6 ]

## Edge Coverage

**Edges** = { (0,1), (0,2), (1,2), (2,3), (2,4), (3,6), (4,5), (4,6), (5,4) }

**Test Paths:** [ 0, 1, 2, 3, 6 ] [ 0, 2, 4, 5, 4, 6 ]

## Complete Path Coverage

**Paths** = { [ 0 ], [ 0,1 ], [ 0,2 ], [ 0,1,2 ], [ 1,2 ], [ 1,2,3 ], ... }

**Test Paths:** [ 0, 1, 2, 3, 6 ] [ 0, 1, 2, 4, 6 ] [ 0, 1, 2, 4, 5, 4, 6 ]  
[ 0, 1, 2, 4, 5, 4, 5, 4, 6 ] [ 0, 1, 2, 4, 5, 4, 5, 4, 5, 4, 6 ] ...

# Loops in Graphs

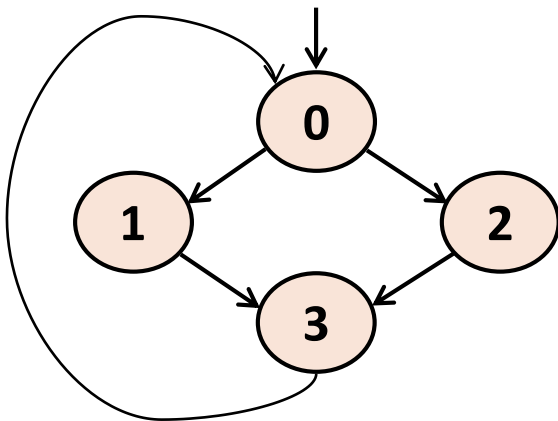
- If a graph contains a loop, it has an infinite number of paths
- Thus, Complete Path Coverage is not feasible
- History of attempts to “deal with” loops:
  - Years 1980s : Execute each loop, exactly once ([4, 5, 4] in previous example)
  - Years 1990s : Execute loops 0 times, once, twice, ..., up to specified limit
  - Years 2000s : *Prime paths*

**Optimus  
Prime**



# Simple Paths and Prime Paths

- Simple Path : A path from node ***m*** to ***n*** is simple if no node appears more than once, except possibly the first and last nodes are the same
  - No internal loops,
  - But the path it self may be a loop.
- Prime Path : A simple path that does not appear as a proper subpath of any other simple path



Simple Paths : [ 0, 1, 3, 0 ], [ 0, 2, 3, 0 ], [ 1, 3, 0, 1 ],  
[ 2, 3, 0, 2 ], [ 3, 0, 1, 3 ], [ 3, 0, 2, 3 ], [ 1, 3, 0, 2 ],  
[ 2, 3, 0, 1 ], [ 0, 1, 3 ], [ 0, 2, 3 ], [ 1, 3, 0 ], [ 2, 3, 0 ],  
[ 3, 0, 1 ], [ 3, 0, 2 ], [ 0, 1 ], [ 0, 2 ], [ 1, 3 ], [ 2, 3 ], [ 3, 0 ],  
[ 0 ], [ 1 ], [ 2 ], [ 3 ]

Prime Paths : [ 0, 1, 3, 0 ], [ 0, 2, 3, 0 ], [ 1, 3, 0, 1 ],  
[ 2, 3, 0, 2 ], [ 3, 0, 1, 3 ], [ 3, 0, 2, 3 ], [ 1, 3, 0, 2 ],  
[ 2, 3, 0, 1 ]

# Prime Path Coverage

- A simple, elegant and finite criterion that requires loops to be executed as well as skipped

**Prime Path Coverage (PPC)** : Test suite  $T$  satisfies *prime path coverage* on graph  $G$  iff, for every prime path  $q$ , some path in  $path(T)$  tours  $q$  with sidetrips allowed.

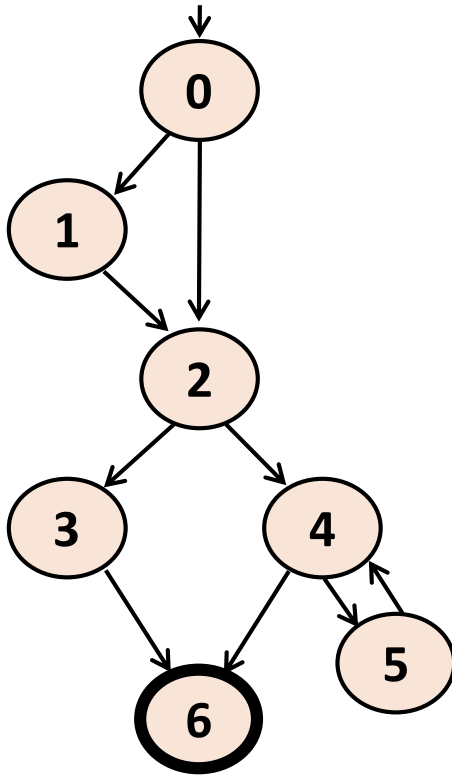
- Prime Path Coverage *subsumes* node and edge coverage, i.e. if a test set satisfies prime path coverage, it also satisfies node and edge coverage.

Given a test suite  $T$ , the degree of prime path coverage of  $T$  can be calculated form:

$$C_{PPC} = \frac{\text{Number of prime paths toured by some test path}}{\text{Number of prime paths}} \times 100\%$$

# Prime Path Example

- The following example has 38 **simple** paths
- Only **nine prime paths**



## Prime Paths

[ 0, 1, 2, 3, 6 ]

[ 0, 1, 2, 4, 5 ]

[ 0, 1, 2, 4, 6 ]

[ 0, 2, 3, 6 ]

[ 0, 2, 4, 5 ]

[ 0, 2, 4, 6 ]

[ 5, 4, 6 ]

[ 4, 5, 4 ]

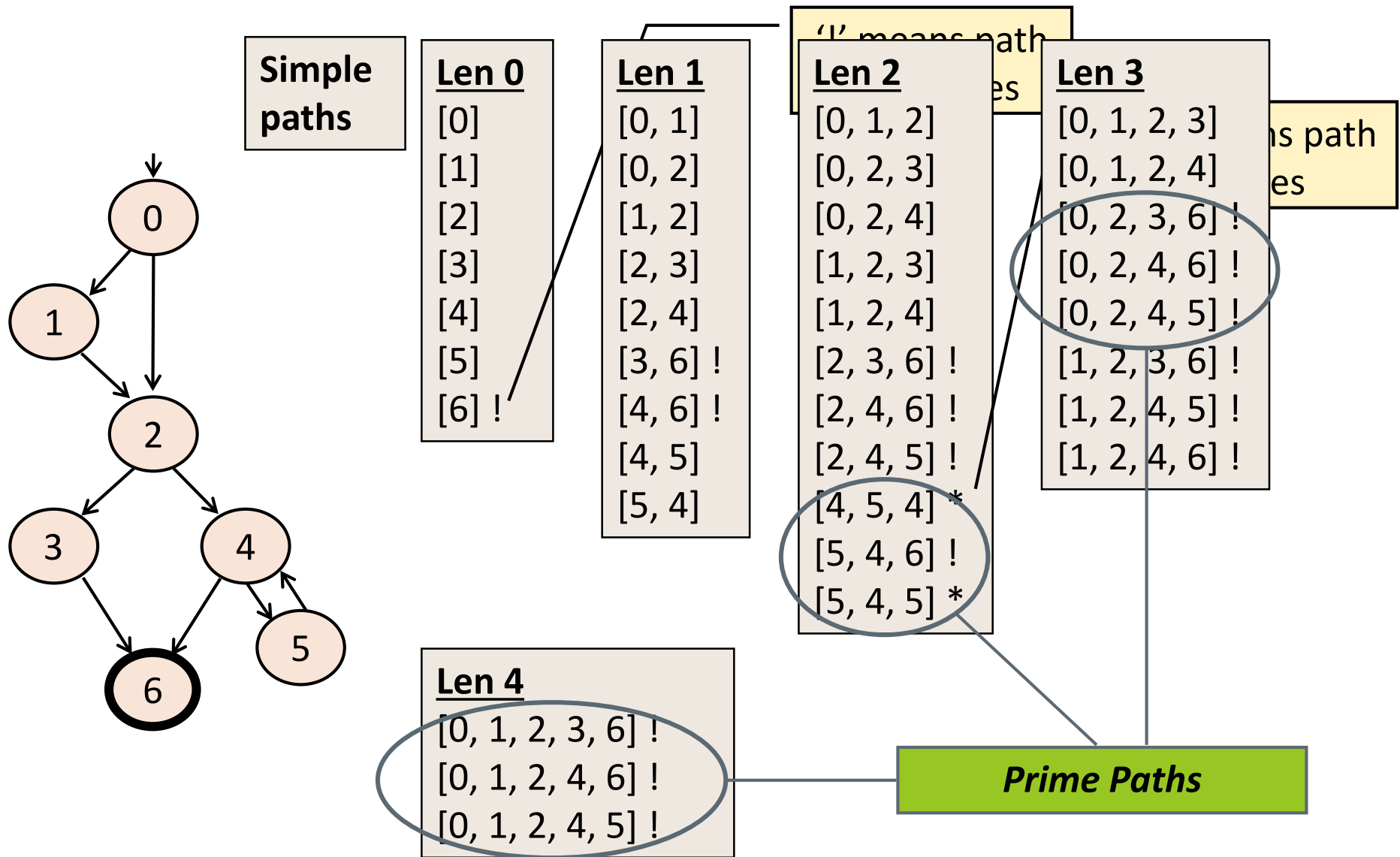
[ 5, 4, 5 ]

Execute loop  
0 times

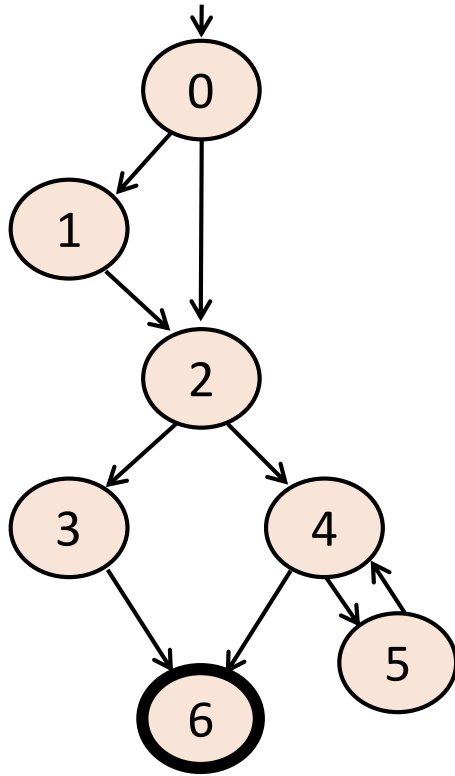
Execute loop  
once

Execute loop  
more than once

# How to Find Prime Paths



# Prime Path Coverage



## Prime Paths

[ 0, 1, 2, 3, 6 ]

[ 0, 1, 2, 4, 5 ]

[ 0, 1, 2, 4, 6 ]

[ 0, 2, 3, 6 ]

[ 0, 2, 4, 5 ]

[ 0, 2, 4, 6 ]

[ 5, 4, 6 ]

[ 4, 5, 4 ]

[ 5, 4, 5 ]

- An example of a set of test paths that covers all prime paths:

{ [ 0, 1, 2, 3, 6 ], [ 0, 1, 2, 4, 6 ],  
[ 0, 2, 3, 6 ], [ 0, 2, 4, 6 ],  
[ 0, 1, 2, 4, 5, 4, 6 ],  
[ 0, 2, 4, 5, 4, 5, 4, 6 ] }