

Ad Auctions

Internet Analytics (COM-308)

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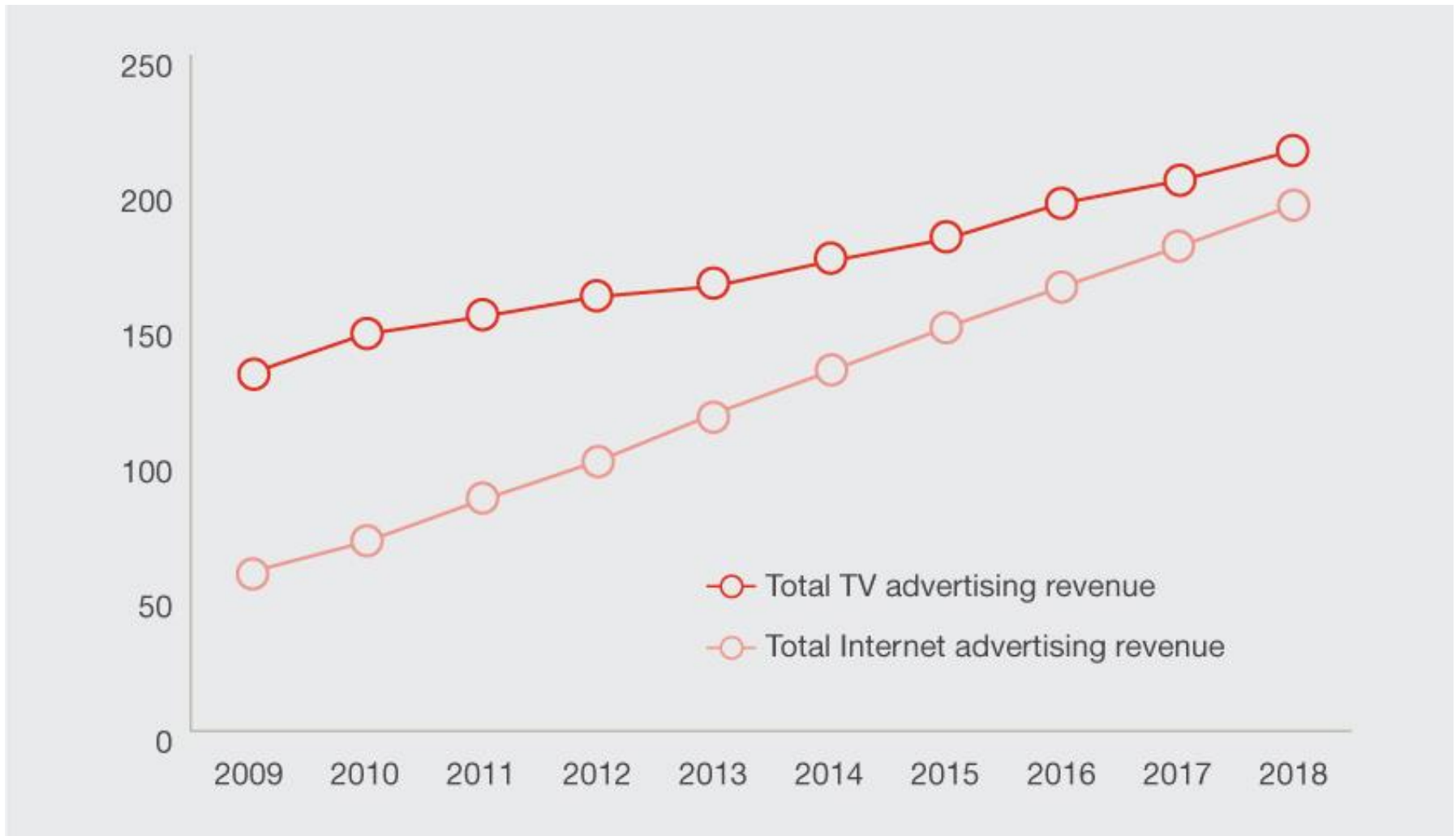
Overview

- Advertisement is **the** Internet business model
 - Plus a bit of subscription, freemium, etc.
- Most ad impressions are auctioned
 - Google AdSense: banner ads on third-party websites
 - Google AdWords: search ads
- Classical auctions: single item
 - First vs second price
- Ad auctions: multiple items = several “positions” on web page
 - Generalized Second Price auction
 - VCG auction

Online advertisement market

- Business models for the Internet:
 - (a) Paying for it: subscription, per-use, taxes, ...
 - (b) Providing data and consuming advertisement
- “If you’re not paying for it, you’re the product”
 - Willingness to sacrifice privacy and “paying through attention”
 - Not just about the money, but ease of use
- Ad market:
 - Globally ~ 700bn\$/year
 - Online: ~ 150bn\$/year (2015)
 - Web: maturing; mobile: big challenge

Online vs TV advertisement market

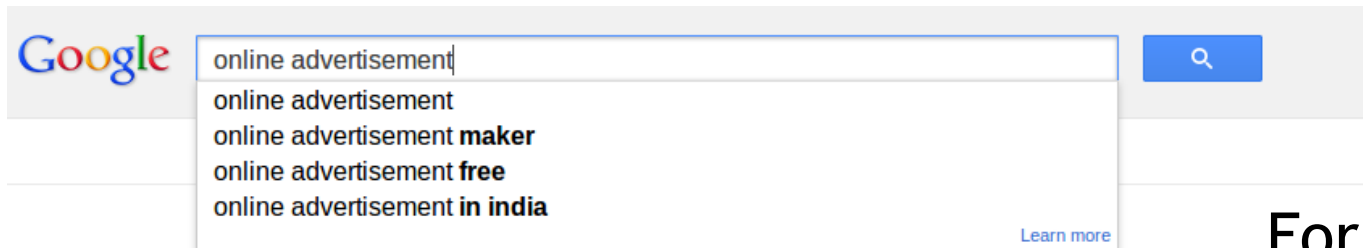


[source: PwC 2014]

Online advertisement market

- Banner ads:
 - Pay-per-impression (user sees ad)
 - CPM (cost-per-mille)
 - Pay-per-click (PPC) (user takes action)
 - CTR (clickthrough rate = clicks per impression)
 - CPC (cost per click)
- Search ads / sponsored search
 - Google AdWords: bidding for keywords plus other constraints (geographic, max cost,...)
 - Google shows your ads in response to searches, charges for traffic

Ad auctions



Ads related to **online advertisement** ⓘ

Online Email Advertising - VerticalResponse.com
www.verticalresponse.com/FreeTrial
Trusted by 500k+ for Reliable Email Marketing Since 2001. Try it Free!
238 people +1'd or follow VerticalResponse
Flexible Pricing Options - Features - Pay As You Go - Start Free Trial Now

Online Advertising - Create your online marketing strat.
www.agencevirtuelle.com/OnlineMarketing
AdWords, SEO, Mobile, Social Media.

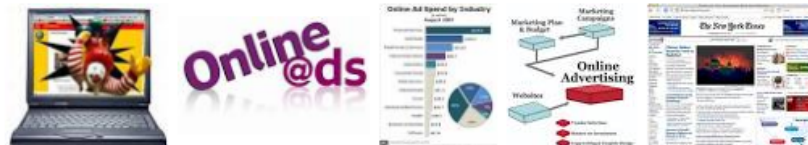
For each search, this table of “sponsored search results” is the result of an online auction

Online advertising - Wikipedia, the free encyclopedia
en.wikipedia.org/wiki/Online_advertising

Online advertising, also known as **online advertisement**, internet marketing, online marketing or e-marketing, is the marketing and promotion of products or ...
[History of online advertising - Competitive advantage over ... - Online advertisement](#)

Online Advertising: How to Do It Right | Small Business Trends
smallbiztrends.com/2010/11/online-advertising-how-to-do.html
Nov 4, 2010 – Helpful tips on using **online advertising** for small businesses.

Images for online advertisement - Report images



Internet advertising: The ultimate marketing machine | The Economist
www.economist.com/node/7138905
Jul 6, 2006 – IN TERMS of efficiency, if not size, the **advertising** industry is only now

Search ads: how to allocate?

- Resource allocation problem:
 - One seller: google, amazon, bing, etc.
 - Many bidders: advertisers buying visibility
 - Several items: multiple locations on results panel
- Auction:
 - Google runs an auction for every search to sell the ad space!
 - Several billion searches per day
 - Very heavy tailed CPC distribution:
 - “mesothelioma” had CPC up to 79 US\$ a few years ago
 - asbestos lawyers

Auctions

- Open vs sealed envelope
 - Public vs private bids
- Open: ascending (English) vs descending (Dutch)
 - Ascending: increase the price until single bidder left
 - Descending: decrease the price until a bidder calls out
- First-price vs second-price
 - First: winning bidder pays highest bid
 - Second: winning bidder pays second-highest bid
- Equivalences:
 - Descending equivalent to sealed first-price
 - Winning bidder calls out when willing to pay the price
 - Ascending equivalent to sealed second-price
 - Winning bidder stops bidding when second-highest drops out

Single-item auctions

- One item for sale, many potential buyers
 - N bidders
 - Bid of bidder i : b_i
 - Internal valuation of bidder i : v_i
 - Advertisement: revenue generated by ad
 - Art: monetary measure for viewing pleasure/pride/envy of friends/...
 - Price paid by bidder i : p_i
 - 0 if lost, price determined by auction mechanism if won
- Payoff (or utility):
 - $$U_i = \begin{cases} 0 & \text{if lost} \\ v_i - p_i & \text{if won} \end{cases}$$

Single-item auction

- Valuation, price, payoff
 - Valuation v_i : depends only on bidder (personal preference, business case, etc.)
 - Price $p_i = p_i(b_1, b_2, \dots, b_N)$: depends on everybody's bid through the auction mechanism
- Strategy:
 - Each bidder selects bid b_i that maximizes $U_i(b_1, b_2, \dots, b_N)$
 - b_i too low: risk not winning the auction
 - b_i too high: risk paying too much

Why second-price auction?

- Intuition:

- First price seems natural and reasonable:
 - Bid what you are willing to pay = value
- Second-price seems manipulable:
 - Bid very high to win, pay only 2nd price

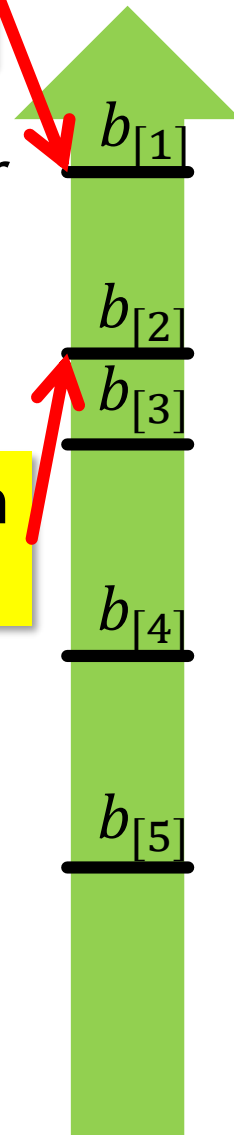
- Theory:

- First price: bidding value means zero payoff → must bid less than value
- Second-price: bidding too high is a bad strategy, if others follow the same strategy → will pay above value, negative payoff

Price with
1st -price

winner

Price with
2nd -price



Truthful bidding in second-price

- Theorem: in a second-price auction, truthful bidding ($b = v$) is a dominating strategy
 - Dominating: regardless of what strategy other players use, best strategy for myself
- Proof: assume I bid b' instead of $b = v$
 - Case $b' < v$ (“under-bidding”):
 - Affects outcome only if 2nd-highest bid $b_{[2]}$ is $b' < b_{[2]} < v \rightarrow$ auction lost, $U = 0$ instead of $U = v - p = v - b_{[2]} \geq 0$
 - Case $b' > v$ (“over-bidding”):
 - Affects outcome only if highest bid $b_{[1]}$ is $v < b_{[1]} < b' \rightarrow$ auction won, $U = v - p = v - b_{[1]} < 0$ instead of $U = 0$

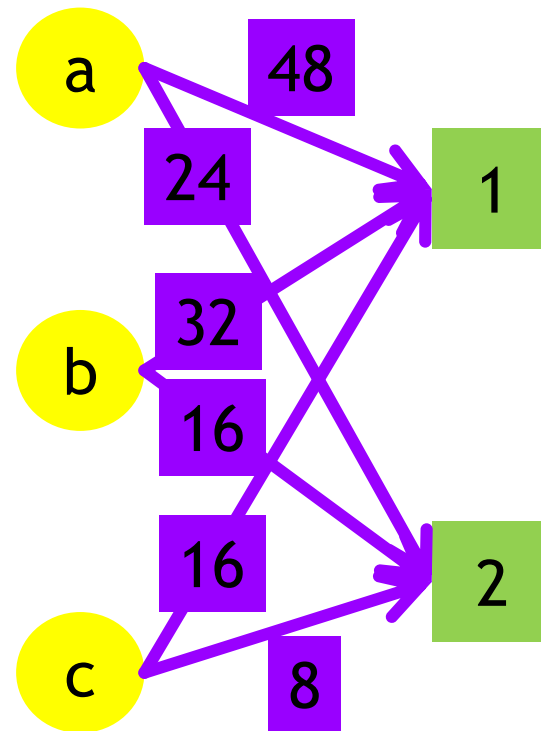
Second price below b'

Multiple-item auctions

- K spaces for sale (decreasing value)
- Bid vector: $b_i = (b_{i1}, b_{i2}, \dots, b_{iK})$
- Allocation:
 - Maximum matching: M maximizes $\sum_{(i,j) \in M} b_{ij}$
- Generalized Second Price (GSP):
 - Winner of k th item pays winning bid for $(k + 1)$ st item
 - Simple, used in Google AdWords
- Vickrey-Clarke-Groves (VCG):
 - Bidder i pays its “damage” (externality in economics-speak) on everybody else
 - More on this later...

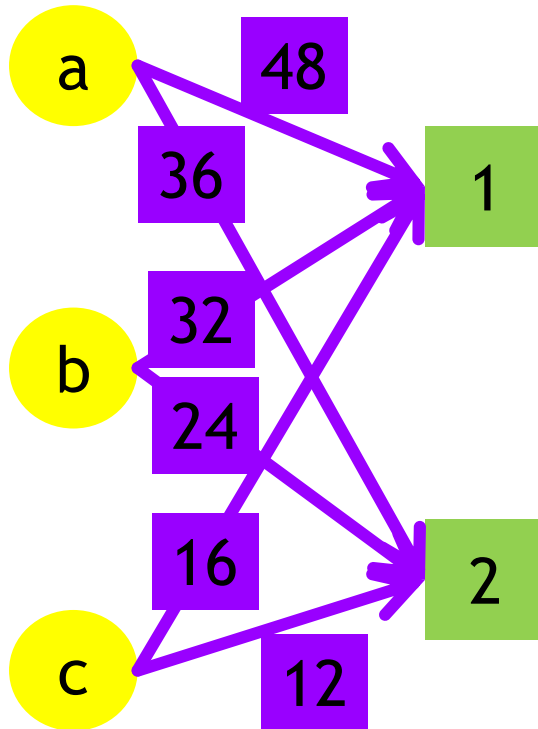
Generalized Second-Price (GSP)

- Ad auctions:
 - Value = CTR x value per click
 - Normally CTR decreases with list, value-per-click assumed independent of position
 - Maximum matching = {(highest bidder, 1st position), (2nd highest bidder, 2nd position), ...}
- Example:
 - Assume CTR of (4%,2%) and valuations per click of (12,8,4)
 - Prices (for truthful bids $b = v$):
 - a pays 32
 - b pays 8



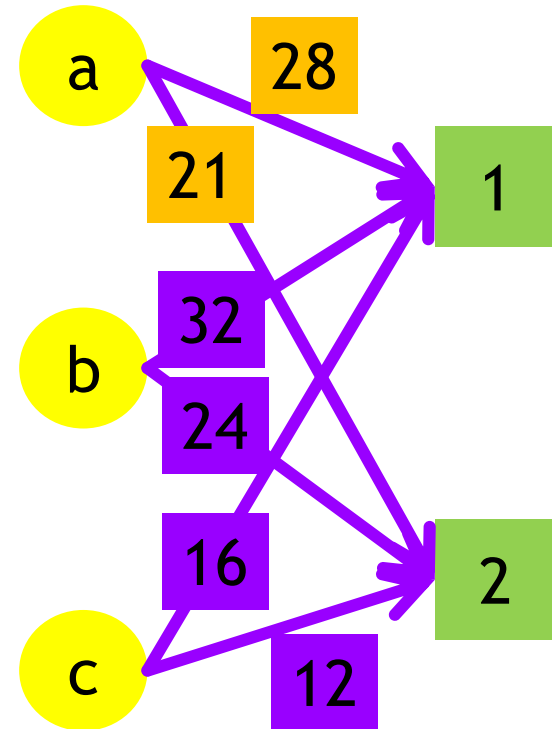
Non-truthful bidding in GSP

- GSP is not incentive-compatible - example:



Truthful bidding:

$$\begin{aligned} \text{a: } u_a &= v_a - p_a = \\ &= 48 - 32 = 16 \end{aligned}$$



Tactical bidding:

$$\begin{aligned} \text{a: } u_a &= v_a - p_a = \\ &= 28 - 24 = 4 \end{aligned}$$

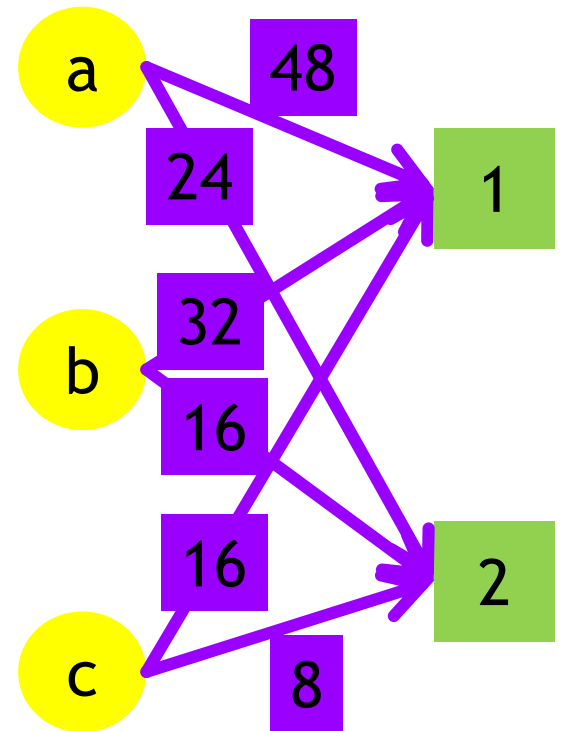
VCG (Vickrey-Clarke-Groves) auction

- Suppose everybody bids truthfully: $b_{ij} = v_{ij}$
- Total valuation (for best matching): $V = \sum_M v_{ij}$
- Def: if bidder i gets item j : $V_{i \leftarrow j} = \sum_{M-(i,j)} v_{ij}$
 - i.e., the value of all bidder-item pairs except for (i, j)
 - Total valuation $V = v_{ij} + V_{i \leftarrow j}$
- Def: valuation without i : V_{-i}
 - Best total valuation if bidder i is completely removed (different matching M')
- Price:
 - Compute matching M
 - $p_i = V_{-i} - V_{i \leftarrow j}$
 - Interpretation: price for i is the decrease in total valuation for everybody else due to i 's participation

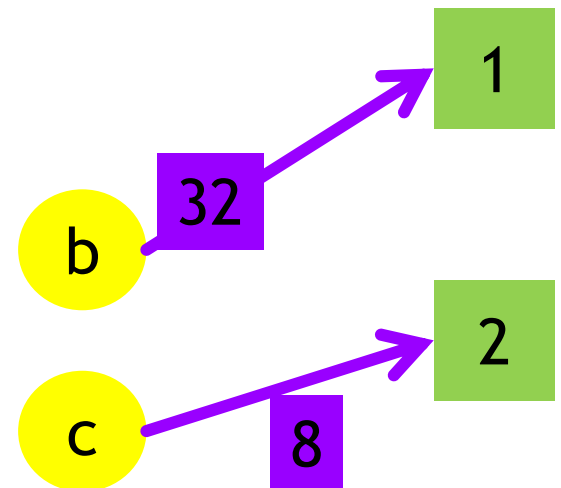
VCG pricing: example

- Allocation identical
- Prices:
 - $p_a = (32 + 8) - (16 + 0) = 24$
 - 16 for reducing b from 32 \rightarrow 16
 - 8 for reducing c from 8 \rightarrow 0
 - $p_b = 8$:
 - 8 for reducing c from 8 \rightarrow 0
- Note: VCG with single item is equivalent to second-price
 - $V_{i \leftarrow 1} = 0$
 - $V_{-i} = b_2$

$V = 64$:




$V_{-a} = 40$:




Truthful bidding in VCG

- Theorem: in a VCG auction, truthful bidding ($b = v$) is a dominating strategy
- Proof:
 - Suppose i bids non-truthfully for some item j
 - $V = v_{ij} + V_{i \leftarrow j} \geq v_{ix} + V_{i \leftarrow x}$ for all $x \neq j$, because V is max valuation (best matching)
 - Subtract V_{-i} from both sides:
$$v_{ij} + V_{i \leftarrow j} - V_{-i} \geq v_{ix} + V_{i \leftarrow x} - V_{-i}$$
 - $v_{ij} - p_{ij} \geq v_{ix} - p_{ix}$



payoff when truthful



payoff when not truthful

Summary

- Online advertisement: scarce resources + competing interests → auction
- Classification: open/sealed; ascending/descending; 1st/2nd price
- Truthful bidding: bid = value
- Second-price: truthful bidding is dominant strategy
- Multi-item:
 - GSP: very simple rule, but not incentive-compatible (truthful not dominant); used by Google AdWords, very high frequency
 - VCG: incentive-compatible

References

- [M. Chiang: Networked Life, Cambridge 2012 (chapter 2)]

Internet Analytics: conclusion

- Types of data: all about people / user-generated
 - Social and info networks, likes & preferences, text & language
 - Real & diverse datasets for labs
- Learning outcomes:
 - Key **models** to characterize data in social web, social media, and mobile apps
 - E.g.: $G(n, p)$; latent factors; vector space; bidding...
 - Key **methods**: prediction, filtering, ranking, searching, selling
 - E.g.: link prediction; graph sampling; PCA; topic models; auction mechanisms; ...
 - **Tools** of the trade: distributed (non-sql) data processing
 - Spark (make sure to put on resume! ;-))

Internet Analytics: conclusion

- Fields:
 - Data Mining: dim reduction; streaming
 - Machine Learning: learning, prediction; regularization; Bayesian networks
 - Network Science: net structure, evolution, processes
 - Graph theory & probability: random graph models; MCMC
- Breadth & straddling fields:
 - Depth was often limited → probe further!
 - No textbook with full coverage → will be built over time
- Where to go from here?
 - Master Specialization in Data Science
 - Master in Data Science
- **Final exam: Wed, 08:15-11:00**

Thanks & good luck!

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