

# Software Engineering

- Exercises 2023 -

# References

- [LL05] T. Lethbridge, R. Laganriere. *Object Oriented Software Engineering Practical Software Development using UML and Java* (1<sup>st</sup>, 2<sup>nd</sup> editions). McGraw Hill, (2001, 2005).
- [FP20] A. Fox, D. Patterson. *Engineering Software as a Service: An Agile Approach using Cloud Computing*. (1<sup>st</sup>, 2<sup>nd</sup> editions) Strawberry Canyon LLC, (2016, 2020). <http://www.saasbook.info/>
- [BRJ99] G. Booch, J. Rumbaugh, I. Jacobson. *The Unified Modeling Language User Guide*. Addison-Wesley, 1999.
- [JBR99] I. Jacobson, G. Booch, J. Rumbaugh. *The Unified Software Development Process*. Addison-Wesley, 1999.
- [BW90] J.C. Baeten, W.P. Wijland. *Process algebra*. Cambridge Univ. Press, 1990.
- [BB02] W. Boggs, M. Boggs. *Mastering UML with Rational Rose* (2<sup>nd</sup> edition). Sybex, 2002.
- [BK08] C. Baier, J.-P. Katoen, *Principles of Model Checking*, MIT Press, 2008.
- [CHVB18] E.M. Clarke, T.A. Henzinger, H. Veith and R. Bloem, Editors, *Handbook of Model Checking*, Springer, 2018.
- [GoF94] E. Gamma, R. Helm, R. Johnson and J. Vlissides, *Design Patterns: Elements of Reusable Object-Oriented Software*, Addison-Wesley, 1994.
- [Nik16] I. Nikolov, *Scala Design Patterns - Write efficient, clean, and reusable code with Scala*, Packt Publishing, 2016.
- [Pie23] B. Pierce et al, *Software Foundations*, 2023. <https://softwarefoundations.cis.upenn.edu/>
- [RS13] D. Rosenberg, M. Stephens. *Use Case Driven Object Modeling with UML: Theory and Practice*, 2<sup>nd</sup> edition, Apress, 2013.
- [Som16] I. Sommerville, *Software Engineering*, (10th edition). Addison-Wesley, 2016.
- [IBM07] <http://www-306.ibm.com/software/rational/>
- [Haskell] Haskell web page, [www.haskell.org](http://www.haskell.org)
- [PRISM] PRISM model checker, <http://www.prismmodelchecker.org>
- [PRISM-tut] PRISM tutorial, available from: <http://www.prismmodelchecker.org/tutorial/>  
<http://www.prismmodelchecker.org/courses/pmc1112/>
- [PRISM-sem] The PRISM language –semantics, available from: <http://www.prismmodelchecker.org/doc/semantics.pdf>

# Developing requirements - Exercise 1

- (Source: [LL05], ch. 4) Draw an UML use case diagram to specify the requirements for the library application described in [LL05] Example 4.9

# PRISM case study - Exercise 2

- Source: Dynamic Power Management (DPM) case study available from PRISM tutorial <http://www.prismmodelchecker.org/tutorial/>  
<http://www.prismmodelchecker.org/courses/pmc1112/>
- Study the (P, S and R) operators of the PRISM property specification language (described in the course lecture slides and the PRISM manual <http://www.prismmodelchecker.org/manual/>).
  - Analyze the behavior of DPM system by following the steps in the tutorial (see part 3 of the PRISM tutorial). This involves creating PRISM properties and may involve developing appropriate reward structures.

# PRISM semantics - Exercise 3

The following PRISM code describes a DTMC:

```
dtmc
module M1
  v1 : [0..1] init 0;
  [] v1=0 & v2=0 -> 0.9:(v1'=0) + 0.1:(v1'=1);
  [a] v1=0 & v2=1 -> 1:(v1'=1);
  [b] v1=1 -> 1:true;
endmodule
module M2
  v2 : [0..1] init 0;
  [] v1=0 & v2=0 -> 0.7:(v2'=0) + 0.3:(v2'=1);
  [a] v1=0 & v2=1 -> 1:true;
  [b] v1=1 -> 1:true;
endmodule
```

[1] The PRISM language semantics, available from:  
<http://www.prismmodelchecker.org/doc/semantics.pdf>

Define the semantics of this PRISM model following [1]

- a) Construct the system module from the component modules
  - b) Give the semantics of the system module as a transition probability matrix  $P:S \times S \rightarrow [0,1]$
- We recall that a (discrete) probability distribution over a countable set  $S$  is a function  $\mu:S \rightarrow [0,1]$  satisfying
$$\sum_{s \in S} \mu(s) = 1$$
  - We use  $[s_0 \rightarrow p_0, \dots, s_n \rightarrow p_n]$  to denote the distribution that chooses  $s_i$  with probability  $p_i$  (for  $1 \leq i \leq n$ ) and  $\text{Dist}(S)$  for the set of distributions over  $S$

# PRISM semantics - Exercise 4

The following PRISM code describes a CTMC:

```
ctmc
module M1
v1 : [0..1] init 0;
[] v1=0 & v2=0 -> 4.5:(v1'=0) + 0.5:(v1'=1);
[a] v1=0 & v2=1 -> 1:(v1'=1);
[b] v1=1 -> 1:true;
endmodule
module M2
v2 : [0..1] init 0;
[] v1=0 & v2=0 -> 3.5:(v2'=0) + 1.5:(v2'=1);
[a] v1=0 & v2=1 -> 1:true;
[b] v1=1 -> 2:true;
endmodule
```

Define the semantics of this PRISM model following [1]

- a) Construct the system module from the component modules
- b) Give the semantics of the system module as a transition rate matrix  $R: S \times S \rightarrow \mathbb{R}_{\geq 0}$

[1] The PRISM language semantics, available from:  
<http://www.prismmodelchecker.org/doc/semantics.pdf>

# PRISM semantics - Exercise 5

The following PRISM code describes a MDP:

```
mdp
module M1
  v1 : [0..1] init 0;
  [] v1=0 & v2=0 -> 0.9:(v1'=0) + 0.1:(v1'=1);
  [a] v1=0 & v2=1 -> 1:(v1'=1);
  [b] v1=1 -> 1:true;
endmodule
module M2
  v2 : [0..1] init 0;
  [] v1=0 & v2=0 -> 0.7:(v2'=0) + 0.3:(v2'=1);
  [a] v1=0 & v2=1 -> 1:true;
  [b] v1=1 -> 1:true;
endmodule
```

[1] The PRISM language semantics, available from:  
<http://www.prismmodelchecker.org/doc/semantics.pdf>

Define the semantics of this PRISM model following [1]

- Construct the system module from the component modules
- Give the semantics of the system module as the function (given in [1]):

$$\text{Steps} : S \rightarrow 2^{\text{Dist}(S)}$$

- $S$  is the set of states,  $2^{\text{Dist}(S)}$  = power set of  $\text{Dist}(S)$ , i.e., the set of all subsets of  $\text{Dist}(S)$
- Give the semantics of the system module when  $\text{Steps} : S \rightarrow 2^{(\text{Act} \cup \{\wedge\}) \times \text{Dist}(S)}$  is (re)defined as follows:

$$\text{Steps}(s) = \{(\alpha, \mu_{c,s}) \mid c \in C, s \in S\}$$

- $(\alpha \in) \text{Act} \cup \{\wedge\}$ ;  $\text{Act}$  = set of action labels,  $\wedge \notin \text{Act}$  ( $\wedge$  is not an element of  $\text{Act}$ )
- $\alpha = a$ , if  $c = [a] g \rightarrow \lambda_1:u_1 + \dots + \lambda_n:u_n$
- $\alpha = \wedge$ , if  $c = [] g \rightarrow \lambda_1:u_1 + \dots + \lambda_n:u_n$
- $\mu_{c,s} : \text{Dist}(S)$  -  $\mu_{c,s}$  are given in [1]

# Modeling with classes - Exercise 6

- (Source: [LL05], E91) Draw a class diagram corresponding to the object diagram shown in [LL05] Fig. 5.18

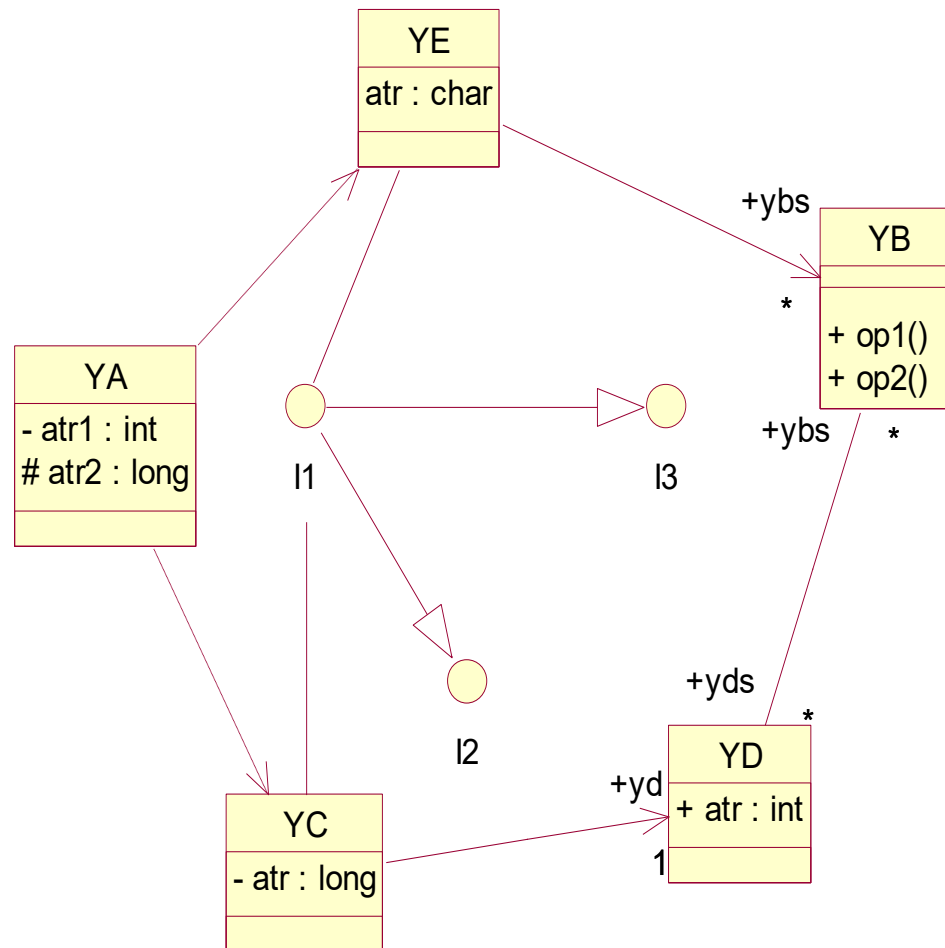


# Modeling with classes - Exercise 7

- (Source: [LL05], ch.5) Consider the description of the Airline system described in [LL05] ch.5 (and the course lecture slides)
- Analyze and design the system using UML class diagrams (apply the methodology given in [LL05] ch.5 "Modeling with classes")

# Forward and reverse engineering

## Forward Engineering (UML -> Java) – Exercise 8



# Forward and reverse engineering

## Reverse Engineering (Java -> UML) – Exercise 9

```
class F extends G implements I3 {  
    private K[] ks;  
    int i;  
    public K k;  
    protected H h;  
    F() {}  
}
```

```
class G {  
    public G() {  
    }  
    public void f(H h) {}  
}
```

```
class H {  
    public H() {  
    }  
}
```

```
class K {  
    private F f;  
    public K() {  
    }  
    void g() { H h = new H(); }  
}
```

```
class L {  
    public static void main(String args[]) {  
        H h = new H();  
        G g = new G();  
        F f = new F();  
        System.out.println("Test");  
    }  
}
```

```
interface I1 {  
}
```

```
interface I2 extends I1,I3 {  
}
```

```
interface I3 {  
}
```

# Design patterns – Exercise 10

- Exercise: Use the General Hierarchy (or Composite) design pattern to model the Lisp data type, which can be specified using Haskell algebraic data types as follows:

```
data Lisp = Nil | Number Int | Symbol String | Cons Lisp Lisp
```

# Design patterns – Exercise 11

- [Source: [LL05], ch.6, Exercise E113] By using the Abstraction-Occurrence design pattern develop a class diagram describing:
  - The issues of a periodical
  - The copies of the issues of a periodical
- [Source: [LL05], ch. 6] Consider the list of design patterns presented in [LL05] Chapter 6 (“Using Design Patterns”) : Abstraction-Occurrence, Player-Role, General Hierarchy, Singleton, Observer, Delegation, Façade, Immutable, Read-Only Interface and Proxy. Solve Exercise E128 from [LL05] by determining the design patterns (from the above list) that would apply in specific circumstances.

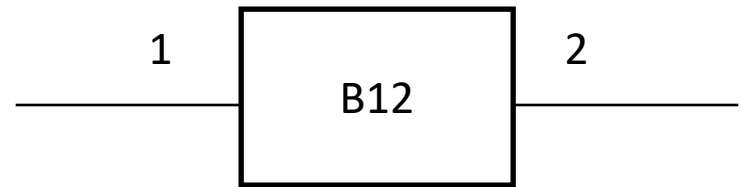
# Design patterns – Exercise 12

- Using the General Hierarchy (or the Composite) pattern create a class diagram to represent PCTL (Probabilistic Computation Tree Logic) expressions (PCTL is a probabilistic extension of the temporal logic CTL). The syntax of PCTL is given below in BNF (p = probability, a = atomic proposition):

$$\Phi ::= \text{true} \mid a \mid \neg \Phi \mid \Phi \wedge \Phi \mid P_{\sim p}[\Psi]$$
$$\Psi ::= X \Phi \mid \Phi \cup_{\leq k} \Phi$$

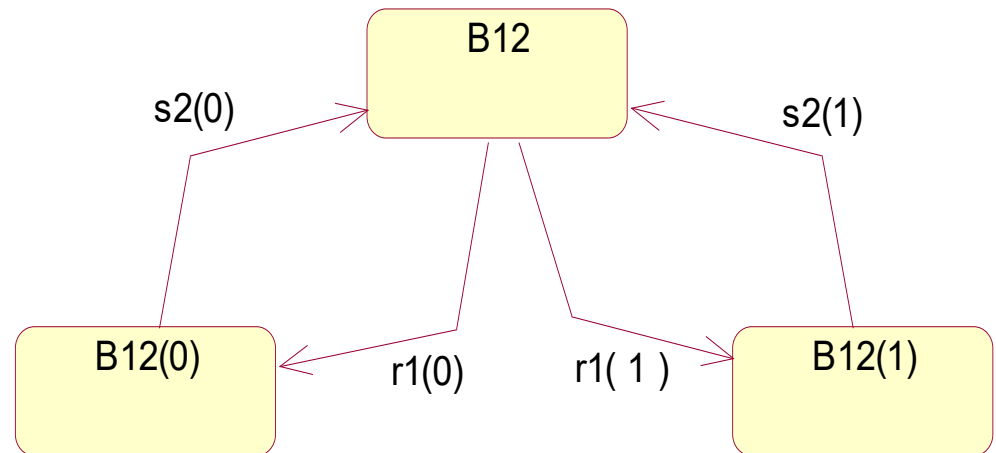
# Modeling interaction and behavior – Exercise 13

- Consider a **bit buffer** B12 of **capacity one** with ports 1 (input) and 2 (output):



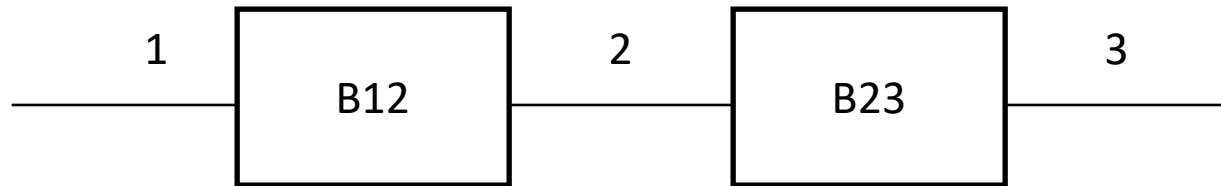
- The behavior of B12 is described by the following STD, where:

- $rj(b)$  = receive bit  $b$  at port  $j$
- $sj(b)$  = send bit  $b$  at port  $j$
- $B12$  = buffer empty
- $B12(b)$  = buffer full ( $b = 0,1$ )



# Modeling interaction and behavior – Exercise 13

- Draw a STD to model the behavior of two buffers B12 and B23 connected at port 2: B12 can send data at port 2, this data is received by B23.
  - By convention, the simultaneous execution of a pair of actions  $s_j(b)$  (send bit  $b$  at port  $j$ ) and  $r_j(b)$  (receive bit  $b$  at port  $j$ ) results in a transmission of  $b$  by a communication at port  $j$ :  $c_j(b)$



- Use the resulted STD to show that two one bit buffers connected in this way behave as a buffer of capacity two [BW90].



# Modeling interaction and behavior – Exercise 14

(Source: [LL05], ch. 8) Implement the class `CourseSection` in Java based on the class diagram given in [LL05] Fig. 8.2 and the state transition diagram given in [LL05] Fig. 8.14

# Architecting & Designing Software – Exercise 15

- (Source: [LL05], ch. 9) A software design should lead to modules that exhibit high cohesion and low coupling. Solve exercises E172 and E176 (from book [LL05]), which require you to categorize certain design aspects according to the types of cohesion or coupling of the modules involved.

# Testing & Inspecting to Ensure High Quality –

## Exercise 16

- (Source: [LL05], ch. 10) Solve exercise E193 (from book [LL05]) which refers to the use of terms *failure*, *defect* and *error* (according to the definitions given in the book [LL05] and the course lecture slides)
- (Source: [LL05], ch. 10, excerpt from Exercise E195)  
Describe a good set of equivalence class test cases for an information form that asks for the current *day* using the format *dd* (assuming that the current year and the current month are known)

# Use case driven development – Exercise 17

Employ the use-case driven methodology in the Unified Process [JBR99] in developing a software solution for an automated teller machine system providing the following services: Withdraw Money, Deposit Money, Transfer between Accounts, View Balance. Develop complete UML models (specification, analysis, design, implementation and test) for one of the following services:

1. Withdraw Money
2. Deposit Money
3. View Balance

# Declarative prototyping – Exercise 18

Consider the following Haskell prototype (quicksort with difference lists):

```
qsd :: [Int] -> [Int] -> [Int]
qsd []      ys = ys
qsd (x : xs) ys =
    let (ls,gs) = partit x xs
    in qsd ls (x : qsd gs ys)

partit :: Int -> [Int] -> ([Int],[Int])
partit p []      = ([],[ ])
partit p (x:xs) =
    let (ls,gs) = partit p xs
    in  if (x < p) then (x : ls,gs) else (ls,x:gs)

Main> qsd [1,7,3,12,4,2,5] [0,0,0]
[1,2,3,4,5,7,12,0,0,0]
```

1. Prove the correctness of the prototype by mathematical induction
2. Implement the prototype by using an imperative language (e.g. C)

# Declarative prototyping – Exercise 19

Consider the following Haskell prototype (binary search tree merging using difference lists):

```
data Tree = NIL | T Tree Int Tree
flatt :: Tree -> [Int] -> [Int]
flatt NIL      ys = ys
flatt (T l n r) ys = flatt l (n : flatt r ys)

merge :: Tree -> Tree -> [Int] -> [Int]
merge NIL t ys = flatt t ys
merge t NIL ys = flatt t ys
merge (T NIL n1 r1) (T NIL n2 r2) ys =
    if (n1 < n2) then n1:merge r1 (T NIL n2 r2) ys
    else n2:merge (T NIL n1 r1) r2 ys
merge (T (T ll1 ln1 lr1) n1 r1) t2 ys =
    merge (T ll1 ln1 (T lr1 n1 r1)) t2 ys
merge t1 (T (T ll2 ln2 lr2) n2 r2) ys =
    merge t1 (T ll2 ln2 (T lr2 n2 r2)) ys

t1 = T (T NIL 1 (T NIL 3 NIL)) 5 (T NIL 7 NIL)
t2 = T (T NIL 2 NIL) 4 (T (T NIL 6 NIL) 8 NIL)
Main> merge t1 t2 [0,0]
[1,2,3,4,5,6,7,8,0,0]
```

1. Prove the correctness of the prototype by mathematical induction
2. Implement the prototype by using an imperative language (e.g. C)