

# **Trabajo Práctico III**

*Análisis de Lenguajes de Programación*

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## Ejercicio 1

Árbol de derivación de tipos de  $S = \lambda x : B \rightarrow B \rightarrow B. \lambda y : B \rightarrow B. \lambda z : B. (xz)(yz)$  en página 4.

El resultado del árbol queda comprobado usando el intérprete:

```
ST> :type S
(B -> B -> B) -> (B -> B) -> B -> B
```

## Ejercicio 2

infer retorna un valor de tipo Either para poder devolver mensajes de error cuando no se pueda inferir el tipo de la expresión.

$\gg=$  toma un valor  $v$  de tipo Either String Type y una función  $f$  de tipo  $Type \rightarrow Either String Type$  y retorna Left  $v$  si es una cadena (Left) o la función  $f$  aplicada a  $v$  si es un tipo (Right). Sirve para propagar errores.

## Ejercicio 5

Árbol de tipado para  $(let\ z = ((\lambda x : B.x)\ as\ B \rightarrow B)\ in\ z)\ as\ B \rightarrow B$ :

$$\frac{\frac{\frac{x : B \in x : B}{x : B \vdash x : B} T-VAR}{\vdash (\lambda x : B.x) : B \rightarrow B} T-ABS \quad \frac{\frac{\vdash (\lambda x : B.x) : B \rightarrow B}{\vdash (\lambda x : B.x)\ as\ B \rightarrow B : B \rightarrow B} T-ASCRIBE \quad \frac{z : B \rightarrow B \in z : B \rightarrow B}{z : B \rightarrow B \vdash z : B \rightarrow B} T-VAR}{\vdash (let\ z = ((\lambda x : B.x)\ as\ B \rightarrow B)\ in\ z) : B \rightarrow B} T-LET} \vdash (let\ z = ((\lambda x : B.x)\ as\ B \rightarrow B)\ in\ z)\ as\ B \rightarrow B : B \rightarrow B T-ASCRIBE$$

El resultado del árbol queda comprobado usando el intérprete:

```
ST> :type (let z = ((\x:B.x) as B -> B) in z) as B->B
B -> B
```

## Ejercicio 7

Reglas de evaluación para pares:

$$\begin{array}{c}
\frac{t_1 \rightarrow t'_1}{(t_1, t_2) \rightarrow (t'_1, t_2)} \text{ (E-PAIR1)} \\
\\
\frac{t \rightarrow t'}{FST\ t \rightarrow FST\ t'} \text{ (E-FST0)} \\
\\
\frac{}{FST\ (v_1, v_2) \rightarrow v_1} \text{ (E-FST1)} \\
\\
\frac{t_2 \rightarrow t'_2}{(v, t_2) \rightarrow (v, t'_2)} \text{ (E-PAIR2)} \\
\\
\frac{t \rightarrow t'}{SND\ t \rightarrow SND\ t'} \text{ (E-SND0)} \\
\\
\frac{}{SND\ (v_1, v_2) \rightarrow v_2} \text{ (E-SND1)}
\end{array}$$

## Ejercicio 9

Árbol de tipado para  $fst\ (unit\ as\ Unit, \lambda x : (B, B).snd\ x)$ :

$$\frac{\frac{\frac{}{\vdash unit : Unit} T-UNIT}{\vdash unit\ as\ Unit : Unit} T-ASCRIBE \quad \frac{\frac{\frac{x : (B, B) \in x : (B, B)}{x : (B, B) \vdash x : (B, B)} T-VAR}{x : (B, B) \vdash snd\ x : B} T-SND}{\vdash \lambda x : (B, B).snd\ x : (B \rightarrow B) \rightarrow B} T-ABS}{\vdash (unit\ as\ Unit, \lambda x : (B, B).snd\ x) : (Unit, (B \rightarrow B) \rightarrow B)} T-PAIR} \vdash fst\ (unit\ as\ Unit, \lambda x : (B, B).snd\ x) : Unit \text{ T-FST}$$

El resultado del árbol queda comprobado usando el intérprete:

```
ST> :type fst (unit as Unit, \x:(B,B). snd x)
Unit
```

## Ejercicio 11

La función Ack queda definida como:

```
def ack = \m:Nat. R (\n:Nat. suc n) (\x:Nat->Nat. \y:Nat. \n:Nat. R (x (suc 0))
                                     (\p:Nat. \q:Nat. x p) n) m
```



Los ejercicios 3, 4, 6, 8 y 10 se encuentran resueltos en el código escrito a continuación:

Código de Common.hs:

```
module Common where

-- Comandos interactivos o de archivos
data Stmt i = Def String i  -- Declarar un nuevo identificador x, let x = t
             | Eval i        -- Evaluar el término
             deriving (Show)

instance Functor Stmt where
  fmap f (Def s i) = Def s (f i)
  fmap f (Eval i)  = Eval (f i)

-- Tipos de los nombres
data Name
  = Global String
  | Quote  Int
  deriving (Show, Eq)

-- Entornos
type NameEnv v t = [(Name, (v, t))]

-- Tipo de los tipos
data Type = Base
          | Unit
          | Nat
          | Fun Type Type
          | Tup Type Type
          deriving (Show, Eq)

-- Términos con nombres
data LamTerm = LUnit
             | LZero
             | LSuc LamTerm
             | LR LamTerm LamTerm LamTerm
             | LVar String
             | Abs String Type LamTerm
             | App LamTerm LamTerm
```

```

        | LLet String LamTerm LamTerm
        | LAs LamTerm Type
        | LTup LamTerm LamTerm
        | LFst LamTerm
        | LSnd LamTerm
    deriving (Show, Eq)

```

-- Términos localmente sin nombres

```

data Term = TUnit
    | Zero
    | Bound Int
    | Free Name
    | Suc Term
    | R Term Term Term
    | Term :@: Term
    | Lam Type Term
    | Let Term Term
    | As Term Type
    | TTup Term Term
    | Fst Term
    | Snd Term
    deriving (Show, Eq)

```

-- Valores

```

data Value = VUnit
    | VZero
    | VSuc Value
    | VLam Type Term
    | VTup Value Value
    deriving Show

```

-- Contextos del tipado

```

type Context = [Type]

```

Código de Symplytyped.hs

```

module Simplytyped (
    conversion,    -- conversion a terminos localmente sin nombre
    eval,          -- evaluador
    infer,         -- inferidor de tipos
    quote         -- valores -> terminos
)
where

```

```

import Data.List
import Data.Maybe
import Prelude hiding ((>=))
import Text.PrettyPrint.HughesPJ (render)
import PrettyPrinter
import Common

-- conversion a términos localmente sin nombres
conversion :: LamTerm -> Term
conversion = conversion' []

conversion' :: [String] -> LamTerm -> Term
conversion' b LUnit      = TUnit
conversion' b LZero      = Zero
conversion' b (LVar n)    = maybe (Free (Global n)) Bound (n `elemIndex` b)
conversion' b (LFst t)    = Fst (conversion' b t)
conversion' b (LSnd t)    = Snd (conversion' b t)
conversion' b (App t u)   = conversion' b t :@: conversion' b u
conversion' b (Abs n t u) = Lam t (conversion' (n:b) u)
conversion' b (LLet n t1 t2) = Let (conversion' b t1) (conversion' (n:b) t2)
conversion' b (LAs u t)   = As (conversion' b u) t
conversion' b (LTup t1 t2) = TTup (conversion' b t1) (conversion' b t2)
conversion' b (LSuc t)     = Suc (conversion' b t)
conversion' b (LR t0 t1 t2) = R (conversion' b t0) (conversion' b t1)
                                (conversion' b t2)

---- eval -----

sub :: Int -> Term -> Term -> Term
sub i t (Bound j) | i == j    = t
sub _ _ (Bound j) | otherwise = Bound j
sub _ _ (Free n)              = Free n
sub _ _ TUnit                 = TUnit
sub _ _ Zero                  = Zero
sub i t (Fst t')              = Fst (sub i t t')
sub i t (Snd t')              = Snd (sub i t t')
sub i t (u :@: v)              = sub i t u :@: sub i t v
sub i t (Lam t' u)             = Lam t' (sub (i+1) t u)
sub i t (Let t0 t1)           = Let (sub i t t0) (sub (i+1) t t1)
sub i t (As u t')             = As (sub i t u) t'
sub i t (Suc t')              = Suc (sub i t t')
sub i t (TTup t0 t1)          = TTup (sub i t t0) (sub i t t1)
sub i t (R t0 t1 t2)          = R (sub i t t0) (sub i t t1) (sub i t t2)

```

```

-- evaluador de términos
eval :: NameEnv Value Type -> Term -> Value
eval _ TUnit      = VUnit
eval _ Zero       = VZero
eval e (Suc t)    = VSuc (eval e t)
eval _ (Bound _)  = error "variable ligada inesperada en eval"
eval e (Free n)   = fst $ fromJust $ lookup n e
eval _ (Lam t u)  = VLam t u
eval e (Lam _ u :@: Lam s v) = eval e (sub 0 (Lam s v) u)
eval e (Lam t u :@: v) = case eval e v of
    VLam t' u' -> eval e (Lam t u :@: Lam t' u')
    VUnit      -> eval e (sub 0 TUnit u)
    VZero      -> eval e (sub 0 Zero u)
    VSuc x     -> eval e (sub 0 (quote (VSuc x)) u)
    _         -> error "Error de tipo en run-time,
                        verificar type checker"
eval e (u :@: v) = case eval e u of
    VLam t u' -> eval e (Lam t u' :@: v)
    _         -> error "Error de tipo en run-time,
                        verificar type checker"
eval e (Let t0 t1) = eval e (sub 0 (quote (eval e t0)) t1)
eval e (As u t)    = eval e u
eval e (Fst t)     = case eval e t of
    VTup t0 t1 -> t0
    _          -> error "Error de tipo en run-time,
                        el argumento debe ser una tupla"
eval e (Snd t)     = case eval e t of
    VTup t0 t1 -> t1
    _          -> error "Error de tipo en run-time,
                        el argumento debe ser una tupla"
eval e (TTup t0 t1) = VTup (eval e t0) (eval e t1)
eval e (R t0 t1 t2) = case eval e t2 of
    VZero -> eval e t0
    VSuc t -> eval e ((t1 :@: (R t0 t1 (quote t))) :@: (quote t))
    _     -> error "Error de tipo en run-time, el
                    tercer argumento de R tiene que ser Nat"

-----
----- quoting
-----

quote :: Value -> Term
quote VUnit      = TUnit
quote VZero = Zero

```



```

quote (VSuc t) = Suc (quote t)
quote (VLam t f) = Lam t f
quote (VTup t0 t1) = TTup (quote t0) (quote t1)

-----
--- type checker
-----

-- type checker
infer :: NameEnv Value Type -> Term -> Either String Type
infer = infer' []

-- definiciones auxiliares
ret :: Type -> Either String Type
ret = Right

err :: String -> Either String Type
err = Left

(>>=) :: Either String Type -> (Type -> Either String Type) -> Either String Type
(>>=) v f = either Left f v

-- fcs. de error

matchError :: Type -> Type -> Either String Type
matchError t1 t2 = err $ "se esperaba " ++
    render (printType t1) ++
    ", pero " ++
    render (printType t2) ++
    " fue inferido."

notfunError :: Type -> Either String Type
notfunError t1 = err $ render (printType t1) ++ " no puede ser aplicado."

notfoundError :: Name -> Either String Type
notfoundError n = err $ show n ++ " no está definida."

infer' :: Context -> NameEnv Value Type -> Term -> Either String Type
infer' _ _ TUnit      = ret Unit
infer' _ _ Zero        = ret Nat
infer' c e (Suc t)     = infer' c e t >>= \tt->
    case tt of
        Nat -> ret Nat
        _    -> matchError Nat tt
infer' c _ (Bound i) = ret (c !! i)

```

```

infer' _ e (Free n) = case lookup n e of
    Nothing -> notfoundError n
    Just (_,t) -> ret t
infer' c e (t :@: u) = infer' c e t >=> \tt ->
    infer' c e u >=> \tu ->
    case tt of
        Fun t1 t2 -> if (tu == t1)
            then ret t2
            else matchError t1 tu
        _ -> notfunError tt
infer' c e (Lam t u) = infer' (t:c) e u >=> \tu ->
    ret $ Fun t tu
infer' c e (Let t0 t1) = infer' c e t0 >=> \tt0 ->
    infer' (tt0:c) e t1
infer' c e (As u t) = infer' c e u >=> \tu ->
    if tu == t then ret t else matchError t tu
infer' c e (TTup t0 t1) = infer' c e t0 >=> \tt0 ->
    infer' c e t1 >=> \tt1 ->
    ret (Tup tt0 tt1)
infer' c e (Fst t) = infer' c e t >=> \tt ->
    case tt of
        Tup tt0 tt1 -> ret tt0
        _ -> notfunError tt
infer' c e (Snd t) = infer' c e t >=> \tt ->
    case tt of
        Tup tt0 tt1 -> ret tt1
        _ -> notfunError tt
infer' c e (R t0 t1 t2) = infer' c e t0 >=> \tt0 ->
    infer' c e t1 >=> \tt1 ->
    infer' c e t2 >=> \tt2 ->
    case tt1 of
        Fun t t' -> if t == tt0 && t' == Fun Nat t && tt2 == Nat
            then ret t
            else matchError (Fun t t') tt1
        _ -> notfunError tt1

```

Código de Parse.y:

```

{
module Parse where
import Common
import Data.Maybe
import Data.Char

}

```

```

%monad { P } { thenP } { returnP }
%name parseStmt Def
%name parseStmts Defs
%name term Exp

```

```

%tokentype { Token }
%lexer {lexer} {TEOF}

```

```

%token
    '='      { TEquals }
    ':'      { TColon }
    '\\\''   { TAbs }
    '.'      { TDot }
    '('      { TOpen }
    ')'      { TClose }
    '->'     { TArrow }
    ','      { TComma }
    LET      { TLet }
    IN       { TIn }
    VAR      { TVar $$ }
    TYPE     { TType }
    DEF      { TDef }
    TYPEUNIT { TTypeUnit }
    NAT      { TNat }
    ZERO     { TZero }
    SUC      { TSuc }
    REC      { TR }
    AS       { TAs }
    UNIT     { TokenUnit }
    FST      { TFst }
    SND      { TSnd }

```

```

%right VAR
%left '='
%right '->'
%right '\\\'' '.' LET IN
%left AS
%right REC
%right SUC
%right SND FST

```

```

%%

```

```

Def      : Defexp                               { $1 }
          | Exp                                 { Eval $1 }
Defexp   : DEF VAR '=' Exp                     { Def $2 $4 }

Exp      :: { LamTerm }
          : '\\\ VAR ':' Type '.' Exp         { Abs $2 $4 $6 }
          | LET VAR '=' Exp IN Exp             { LLet $2 $4 $6 }
          | Exp AS Type                        { LAs $1 $3 }
          | REC Atom Atom Exp                  { LR $2 $3 $4 }
          | FST Atom                           { LFst $2 }
          | SND Atom                           { LSnd $2 }
          | NAbs                               { $1 }

NAbs     :: { LamTerm }
          : NAbs Atom                          { App $1 $2 }
          | Atom                               { $1 }

Atom     :: { LamTerm }
          : VAR                                { LVar $1 }
          | SUC Atom                           { LSuc $2 }
          | '(' Exp ',' Exp ')'                { LTup $2 $4 }
          | '(' Exp ')'                        { $2 }
          | UNIT                               { LUnit }
          | ZERO                               { LZero }

Type     : TYPE                                { Base }
          | TYPEUNIT                           { Unit }
          | NAT                                { Nat }
          | '(' Type ',' Type ')'              { Tup $2 $4 }
          | Type '->' Type                      { Fun $1 $3 }
          | '(' Type ')'                       { $2 }

Defs     : Defexp Defs                        { $1 : $2 }
          |                                     { [] }

```

```
{
```

```

data ParseResult a = Ok a | Failed String
                    deriving Show
type LineNumber = Int
type P a = String -> LineNumber -> ParseResult a

getLineNo :: P LineNumber
getLineNo = \s l -> Ok l

```

```

thenP :: P a -> (a -> P b) -> P b
m 'thenP' k = \s l-> case m s l of
    Ok a      -> k a s l
    Failed e -> Failed e

returnP :: a -> P a
returnP a = \s l-> Ok a

failP :: String -> P a
failP err = \s l -> Failed err

catchP :: P a -> (String -> P a) -> P a
catchP m k = \s l -> case m s l of
    Ok a      -> Ok a
    Failed e -> k e s l

happyError :: P a
happyError = \ s i -> Failed $ "Línea "++(show (i::LineNumber))++":
                                Error de parseo\n"++(s)

data Token = TVar String
            | TType
            | TDef
            | TTypeUnit
            | TNat
            | TAbs
            | TDot
            | TOpen
            | TClose
            | TColon
            | TArrow
            | TComma
            | TEquals
            | TLet
            | TAs
            | TIn
            | TokenUnit
            | TFst
            | TSnd
            | TZero
            | TSuc
            | TR
            | TEOF
            deriving Show

```

```

lexer cont s = case s of
  [] -> cont TEOF []
  ('\\n':s) -> \\line -> lexer cont s (line + 1)
  (c:cs)
    | isSpace c -> lexer cont cs
    | isAlpha c -> lexVar (c:cs)
  ('-':('-':cs)) -> lexer cont $ dropWhile ((/=) '\\n') cs
  ('{':('{':cs)) -> consumirBK 0 0 cont cs
  ('-':('}':cs)) -> \\ line -> Failed $ "Línea "++(show line)++":
                                     Comentario no abierto"

  ('-':('>':cs)) -> cont TArrow cs
  ('\\':cs)-> cont TAbs cs
  ('.':cs) -> cont TDot cs
  ('(':cs) -> cont TOpen cs
  ('-':('>':cs)) -> cont TArrow cs
  (')':cs) -> cont TClose cs
  (':':cs) -> cont TColon cs
  ('=':cs) -> cont TEquals cs
  (',':cs) -> cont TComma cs
  ('0':cs) -> cont TZero cs
  unknown -> \\line -> Failed $ "Línea "++(show line)++":
    No se puede reconocer "++(show $ take 10 unknown)++ "...
where lexVar cs = case span isAlpha cs of
  ("B",rest) -> cont TType rest
  ("Unit", rest) -> cont TTypeUnit rest
  ("Nat", rest) -> cont TNat rest
  ("def",rest) -> cont TDef rest
  ("let",rest) -> cont TLet rest
  ("in",rest) -> cont TIn rest
  ("as", rest) -> cont TAs rest
  ("unit", rest) -> cont TokenUnit rest
  ("fst", rest) -> cont TFst rest
  ("snd", rest) -> cont TSnd rest
  ("suc", rest) -> cont TSuc rest
  ("R", rest) -> cont TR rest
  (var,rest) -> cont (TVar var) rest
  consumirBK anidado cl cont s = case s of
    ('-':('-':cs)) ->
      consumirBK anidado cl
      cont $ dropWhile
        ((/=) '\\n') cs
    ('{':('{':cs)) ->
      consumirBK (anidado+1)
      cl cont cs

```

```

('-' : ('}' : cs)) ->
  case anidado of
    0 -> \line -> lexer
          cont cs (line+cl)
    _ -> consumirBK
          (anidado-1) cl
          cont cs
  ('\n' : cs) ->
    consumirBK anidado
    (cl+1) cont cs
  (_ : cs) -> consumirBK
    anidado cl cont cs

```

```

stmts_parse s = parseStmts s 1
stmt_parse s = parseStmt s 1
term_parse s = term s 1
}

```

Código de PrettyPrinter.hs:

```

module PrettyPrinter (
  printTerm,    -- pretty printer para terminos
  printType,    -- pretty printer para tipos
)
where

import Common
import Text.PrettyPrint.HughesPJ

-- lista de posibles nombres para variables
vars :: [String]
vars = [ c : n | n <- "" : map show [1..], c <- ['x','y','z'] ++ ['a'..'w'] ]

parensIf :: Bool -> Doc -> Doc
parensIf True  = parens
parensIf False = id

-- pretty-printer de términos

pp :: Int -> [String] -> Term -> Doc
pp ii vs TUnit      = text "unit"
pp ii vs Zero       = text "0"
pp ii vs (Suc t)     = text "suc " <>
  parensIf (not (isAtm t)) (pp ii vs t)
pp ii vs (Bound k)  = text (vs !! (ii - k - 1))

```

```

pp _ vs (Free (Global s)) = text s
pp ii vs (i :@: c)        = sep [parensIf (isLam i) (pp ii vs i),
    nest 1 (parensIf (isLam c || isApp c) (pp ii vs c))]
pp ii vs (Lam t c)        = text "\\ " <>
    text (vs !! ii) <>
    text ":" <>
    printType t <>
    text ". " <>
    pp (ii+1) vs c
pp ii vs (Let t0 t1)      = text "let " <>
    text (vs !! ii) <>
    text " = " <>
    pp ii vs t0 <>
    text " in " <>
    pp (ii+1) vs t1
pp ii vs (As u t)         = pp ii vs u <>
    text " as " <>
    printType t
pp ii vs (TTup t0 t1)     = text "(" <>
    pp ii vs t0 <>
    text "," <>
    pp ii vs t1 <>
    text ")"
pp ii vs (Fst t)          = text "fst " <>
    pp ii vs t
pp ii vs (Snd t)          = text "snd" <>
    pp ii vs t
pp ii vs (R t0 t1 t2)     = text "R" <>
    pp ii vs t0 <>
    pp ii vs t1 <>
    pp ii vs t2

isLam (Lam _ _) = True
isLam _         = False

isApp (_ :@: _) = True
isApp _         = False

isAtm TUnit      = True
isAtm Zero       = True
isAtm (TTup t0 t1) = True
isAtm _          = False

```



```

-- pretty-printer de tipos
printType :: Type -> Doc
printType Base      = text "B"
printType Unit      = text "Unit"
printType Nat       = text "Nat"
printType (Fun t1 t2) = sep [ parensIf (isFun t1) (printType t1),
                             text "->",
                             printType t2]

printType (Tup t1 t2) = text "(" <>
                        printType t1 <>
                        text "," <>
                        printType t2 <>
                        text ")"

isFun (Fun _ _)      = True
isFun _              = False

fv :: Term -> [String]
fv TUnit             = []
fv Zero              = []
fv (Suc t)           = fv t
fv (Bound _)         = []
fv (Free (Global n)) = [n]
fv (Free _)          = []
fv (t :@: u)         = fv t ++ fv u
fv (Lam _ u)         = fv u
fv (Let t0 t1)       = fv t0 ++ fv t1
fv (As u t)          = fv u
fv (Fst t)           = fv t
fv (Snd t)           = fv t
fv (TTup t0 t1)      = fv t0 ++ fv t1
fv (R t0 t1 t2)      = fv t0 ++ fv t1 ++ fv t2

---
printTerm :: Term -> Doc
printTerm t = pp 0 (filter (\v -> not $ elem v (fv t)) vars) t

```