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# 通訊系統 (II)

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Prof. Tsai

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## Chapter 1 Signal-Space Analysis

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# Introduction

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- We discuss some basic issues that relate to signal transmission over an **additive white Gaussian noise (AWGN)** channel
  - **Geometric representation** of signals with finite energy
  - **Maximum likelihood (ML)** procedure for signal detection in an AWGN channel
  - Derivation of the **correlation receiver** that is equivalent to the **matched filter receiver**
  - Probability of symbol error and the **union bound** approximation

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## Digital Communication Systems

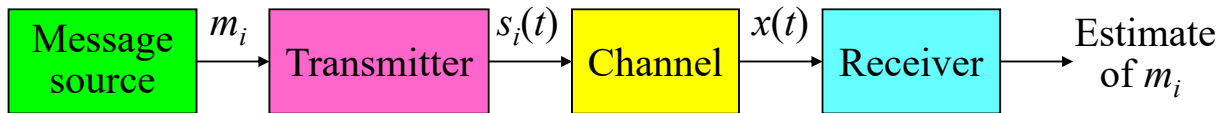
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# Message

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- Consider the most basic form of a **digital** communication system
  - A message source emits one symbol every  $T$  seconds, with the symbols belonging to an alphabet of  $M$  symbols denoted by  $m_1, m_2, \dots, m_M$
  - The ***a priori* probabilities**  $p_1, p_2, \dots, p_M$  specify the **message source** output
    - It is customary to assume that the  $M$  symbols of the alphabet are **equally likely**

$$p_i = P(m_i) = \frac{1}{M} \quad \text{for } i = 1, 2, \dots, M$$



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# Transmitter

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- The transmitter takes the message source output  $m_i$  and codes it into a distinct signal  $s_i(t)$  **suitable for transmission** over the **analog channel**.
- The signal  $s_i(t)$  is a **real-valued energy signal**
$$E_i = \int_0^T s_i^2(t) dt, \quad i = 1, 2, \dots, M$$
  - A signal with **finite energy**
- The design of the signal  $s_i(t)$  is a key issue in communication systems

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# Channel

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- Passing through the channel, the received signal  $x(t)$  is
$$x(t) = \alpha s_i(t) + w(t), \quad 0 \leq t \leq T \text{ and } i = 1, 2, \dots, M$$
  - $\alpha$  is the **complex-valued** channel gain
- The channel is assumed to have two characteristics:
  - The channel is **linear**, with a bandwidth that is **wide enough** to accommodate the transmission of signal  $s_i(t)$ 
    - with **negligible or no distortion**
    - $\alpha$  includes the **attenuation** and **phase rotation**
  - The channel noise  $w(t)$  is the sample function of a zero-mean **white Gaussian noise process**



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# Receiver

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- The receiver has the task of
  - **Observing** the received signal  $x(t)$  for a duration of  $T$
  - Making a best **estimate** of the transmitted signal  $s_i(t)$  (or  $m_i$ )
- However, owing to the presence of channel noise, the receiver will make occasional **errors**
  - To design the receiver so as to **minimize** the average **probability of symbol error**

$$P_e = \sum_{i=1}^M p_i P(\hat{m} \neq m_i | m_i)$$

- where  $m_i$  is the transmitted symbol;  $\hat{m}$  is the estimate produced by the receiver;  $P(\hat{m} \neq m_i | m_i)$  is the **conditional** error probability given that the  $i$ -th symbol was sent

⇒ **Optimum in the minimum probability of error sense**

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# Geometric Representation of Signals

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## Linear Vector Space

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- In signal analysis, we can represent signals as vectors
  - To remove some **redundancy** in the signals
  - To provide a more compact form for the signals
- The signal space could be constructed by **amplitude, phase, frequency** and/or **time**
- A vector space is called a **linear vector space** if it satisfies the following conditions:
  - 1:  $\mathbf{x} + \mathbf{y} = \mathbf{y} + \mathbf{x}$
  - 2:  $\mathbf{x} + (\mathbf{y} + \mathbf{z}) = (\mathbf{x} + \mathbf{y}) + \mathbf{z}$
  - 3:  $\alpha (\mathbf{x} + \mathbf{y}) = \alpha \mathbf{x} + \alpha \mathbf{y}$
  - 4:  $(\alpha + \beta) \mathbf{x} = \alpha \mathbf{x} + \beta \mathbf{x}$
  - where  $\mathbf{x}$  and  $\mathbf{y}$  are arbitrary vectors and  $\alpha$  and  $\beta$  are scalars

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## Linear Vector Space (Cont.)

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- In an  $N$ -dimensional linear vector space, we define a **inner product** as  $\mathbf{x} \cdot \mathbf{y} \triangleq \sum_{i=1}^N x_i y_i$

- where  $x_i$  and  $y_i$  are the elements of  $\mathbf{x}$  and  $\mathbf{y}$ , respectively

- Two vectors  $\mathbf{x}$  and  $\mathbf{y}$  are said to be orthogonal if  $\mathbf{x} \cdot \mathbf{y} = 0$ .

- The **norm** (or the length) of a vector  $\mathbf{x}$  is denoted by  $\|\mathbf{x}\|$

$$\|\mathbf{x}\| \triangleq \sqrt{\mathbf{x} \cdot \mathbf{x}} = \sqrt{\sum_{i=1}^N x_i^2}$$

- This norm has the following properties:

- 5:  $\|\mathbf{x}\| \geq 0$

- 6:  $\|\mathbf{x}\| = 0 \Leftrightarrow \mathbf{x} = \mathbf{0}$

- 7:  $\|\mathbf{x} + \mathbf{y}\| \leq \|\mathbf{x}\| + \|\mathbf{y}\|$

- 8:  $\|\alpha \mathbf{x}\| = |\alpha| \cdot \|\mathbf{x}\|$

- The Schwarz inequality:  $|\mathbf{x} \cdot \mathbf{y}| \leq \|\mathbf{x}\| \cdot \|\mathbf{y}\|$

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## Orthonormal Basis Functions

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- The signal space is assumed to be an  $N$ -dimensional space
  - Constructed by  $N$  **orthonormal** basis functions
- The goal of Geometric Representation of Signals is to represent any set of  $M$  energy signals  $\{s_i(t), i = 1, 2, \dots, M\}$  as **linear combinations** of  $N$  **orthonormal** basis functions,  $N \leq M$

$\Rightarrow$  Given a set of real-valued energy signals  $s_1(t), s_2(t), \dots, s_M(t)$

$$s_i(t) = \sum_{j=1}^N s_{ij} \phi_j(t), \quad 0 \leq t \leq T \text{ and } i = 1, 2, \dots, M$$

$$s_{ij} = \int_0^T s_i(t) \phi_j(t) dt$$

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## Orthonormal Basis Functions (Cont.)

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- The real-valued basis functions  $\phi_1(t), \dots, \phi_N(t)$  are **orthonormal**

$$\int_0^T \phi_i(t) \phi_j(t) dt = \underline{\delta_{ij}} = \begin{cases} 1 & \text{if } i = j \\ 0 & \text{if } i \neq j \end{cases}$$

**Kronecker delta function**

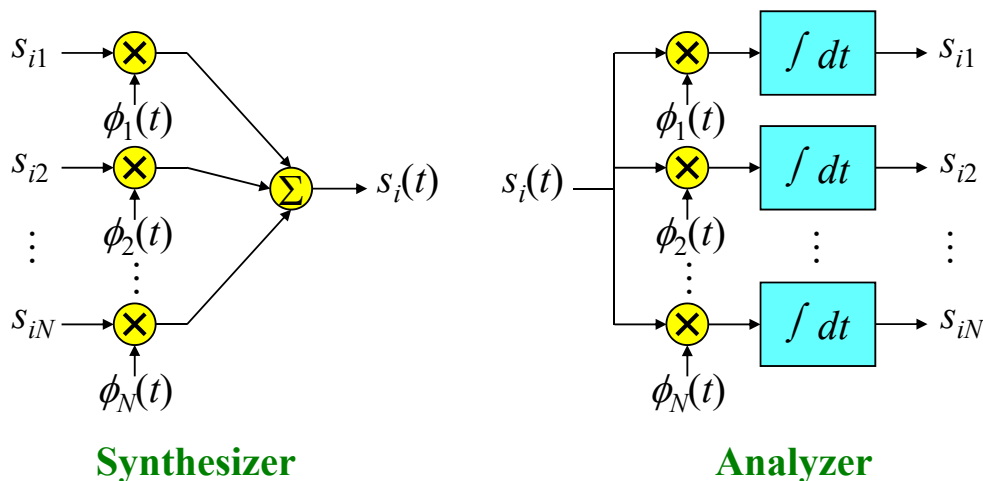
- Each basis function is normalized to have **unit energy**
- The basis functions  $\phi_1(t), \phi_2(t), \dots, \phi_N(t)$  are **orthogonal** with respect to each other over the interval  $0 \leq t \leq T$
- The set of  $\{s_{ij}\}_{j=1}^N$  may be viewed as an  $N$ -dimensional vector  $\mathbf{s}_i$ 
  - A **one-to-one** relationship with the transmitted signal  $s_i(t)$

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## Signal Synthesizer and Analyzer

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- Synthesizer:** given the  $N$  elements of the vectors  $\mathbf{s}_i$  (i.e.,  $s_{i1}, s_{i2}, \dots, s_{iN}$ ) as input to generate the signal  $s_i(t)$
- Analyzer:** given the signals  $s_i(t), i = 1, 2, \dots, M$ , as input to calculate the coefficients  $s_{i1}, s_{i2}, \dots, s_{iN}$



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# Signal Vector

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- Each signal in the set  $\{s_i(t)\}$  is **completely determined** by the vector of its coefficients

$$\mathbf{s}_i = \begin{bmatrix} s_{i1} \\ s_{i2} \\ \vdots \\ s_{iN} \end{bmatrix}, \quad i = 1, 2, \dots, M$$

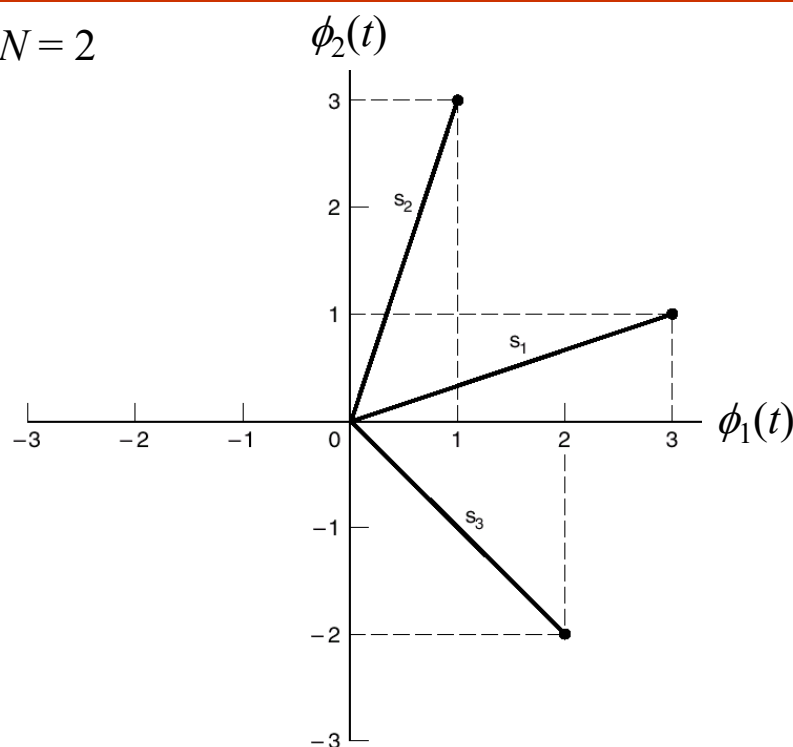
- The vector  $\mathbf{s}_i$  is called a **signal vector**
- Consider an  **$N$ -dimensional** Euclidean space
  - There are  $N$  mutually perpendicular axes labeled  $\phi_1(t)$ ,  $\phi_2(t)$ , ...,  $\phi_N(t)$
  - The set of signal vectors  $\{\mathbf{s}_i \mid i = 1, 2, \dots, M\}$  defines a corresponding set of  **$M$  points** in the Euclidean space
  - This Euclidean space is called the **signal space**

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# Signal Space

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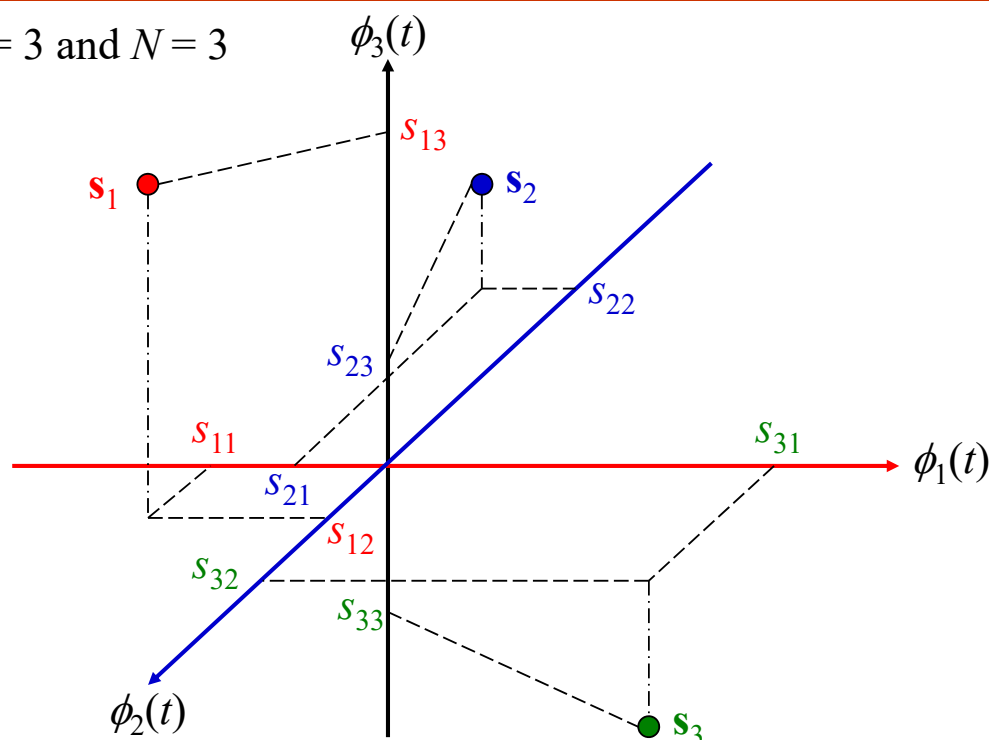
- $M = 3$  and  $N = 2$





# Signal Space

- $M = 3$  and  $N = 3$



# Signal Energy

- The squared-length of any signal vector  $\mathbf{s}_i$  is defined as the **inner product** or **dot product** of  $\mathbf{s}_i$  with itself

$$\|\mathbf{s}_i\|^2 = \mathbf{s}_i^T \mathbf{s}_i = \sum_{j=1}^N s_{ij}^2, \quad i = 1, 2, \dots, M$$

- The **energy** of a signal  $s_i(t)$  of duration  $T$  seconds is defined as

$$\begin{aligned} E_i &= \int_0^T s_i^2(t) dt = \int_0^T \left[ \sum_{j=1}^N s_{ij} \phi_j(t) \right] \left[ \sum_{k=1}^N s_{ik} \phi_k(t) \right] dt \\ &= \sum_{j=1}^N \sum_{k=1}^N s_{ij} s_{ik} \int_0^T \phi_j(t) \phi_k(t) dt \end{aligned}$$

$s_i(t) = \sum_{j=1}^N s_{ij} \phi_j(t)$

- Since the  $\phi_j(t), j = 1, 2, \dots, N$ , form an **orthonormal set**

$$E_i = \sum_{j=1}^N s_{ij}^2 = \|\mathbf{s}_i\|^2$$

$\int_0^T \phi_j(t) \phi_k(t) dt = \delta_{jk}$

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## Distance and Angle

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- For a pair of signals  $s_i(t)$  and  $s_k(t)$ , the inner product of the signals over the interval  $[0, T]$  is

$$\int_0^T s_i(t)s_k(t) dt = \mathbf{s}_i^T \mathbf{s}_k$$

- For a specific pair of signals  $s_i(t)$  and  $s_k(t)$ , the inner product is **invariant** to the choice of basis functions  $\{\phi_j(t)\}_{j=1}^N$ 
  - The **rotation** of the coordinate system does not change the locations of signal points

- The **Euclidean distance** of two vectors  $\mathbf{s}_i$  and  $\mathbf{s}_k$  is

$$d_{ik} = \|\mathbf{s}_i - \mathbf{s}_k\|^2 = \sum_{j=1}^N (s_{ij} - s_{kj})^2 = \int_0^T [s_i(t) - s_k(t)]^2 dt$$

- The **angle**  $\theta_{ik}$  between two signal vectors  $\mathbf{s}_i$  and  $\mathbf{s}_k$  follows

$$\cos \theta_{ik} = \frac{\mathbf{s}_i^T \mathbf{s}_k}{\|\mathbf{s}_i\| \|\mathbf{s}_k\|}$$

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## Example 1: Schwarz Inequality

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- Consider any pair of energy signals  $s_1(t)$  and  $s_2(t)$ . The **Schwarz inequality** state that

$$\left( \int_{-\infty}^{\infty} s_1(t)s_2(t) dt \right)^2 \leq \left( \int_{-\infty}^{\infty} s_1^2(t) dt \right) \left( \int_{-\infty}^{\infty} s_2^2(t) dt \right)$$

- The equality holds if and only if  $s_2(t) = cs_1(t)$ , where  $c$  is a constant

- Proof:** Let  $s_1(t)$  and  $s_2(t)$  be expressed in terms of the pair of orthonormal basis functions  $\phi_1(t)$  and  $\phi_2(t)$  as follows:

$$s_1(t) = s_{11}\phi_1(t) + s_{12}\phi_2(t); \quad s_2(t) = s_{21}\phi_1(t) + s_{22}\phi_2(t)$$

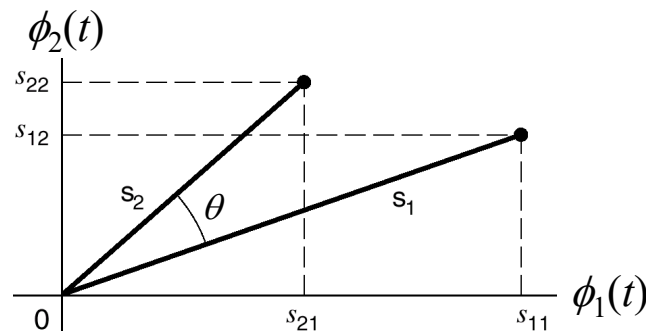
- The angle  $\theta$  between two signal vectors  $\mathbf{s}_1$  and  $\mathbf{s}_2$  is

$$\cos \theta = \frac{\mathbf{s}_1^T \mathbf{s}_2}{\|\mathbf{s}_1\| \|\mathbf{s}_2\|} = \frac{\int_{-\infty}^{\infty} s_1(t)s_2(t) dt}{\left( \int_{-\infty}^{\infty} s_1^2(t) dt \right)^{1/2} \left( \int_{-\infty}^{\infty} s_2^2(t) dt \right)^{1/2}}$$

## Example 1: Schwarz Inequality (Cont.)

- Since  $|\cos\theta| \leq 1 \Rightarrow$  the Schwarz inequality holds
- $|\cos\theta| = 1$  if and only if  $\theta = 0 \Rightarrow \mathbf{s}_2 = c\mathbf{s}_1$ 
  - In other words,  $s_2(t) = cs_1(t)$
- If  $s_1(t)$  and  $s_2(t)$  are **complex-valued** signals, the Schwarz inequality becomes

$$\left| \int_{-\infty}^{\infty} s_1(t) s_2^*(t) dt \right|^2 \leq \left( \int_{-\infty}^{\infty} |s_1(t)|^2 dt \right) \left( \int_{-\infty}^{\infty} |s_2(t)|^2 dt \right)$$



## Gram-Schmidt Orthogonalization Procedure

- Suppose we have a set of  $M$  energy signals:  $s_1(t), s_2(t), \dots, s_M(t)$  with signal energy  $E_1, E_2, \dots, E_M$
- We need to generate a complete orthonormal set of basis functions  $\Rightarrow$  **Gram-Schmidt Orthogonalization Procedure**
- Starting with  $s_1(t)$  chosen from this set **arbitrarily**,
  - The first basis function is defined by

$$\phi_1(t) = s_1(t) / \sqrt{E_1}, \quad s_1(t) = \sqrt{E_1} \phi_1(t) = s_{11} \phi_1(t)$$

- Next, using the signal  $s_2(t)$ , we define the coefficient  $s_{21}$  as

$$s_{21} = \int_0^T s_2(t) \phi_1(t) dt$$

- We obtain a new function  $g_2(t)$  orthogonal to  $\phi_1(t)$  over the interval  $0 \leq t \leq T$

$$g_2(t) = s_2(t) - \underline{s_{21} \phi_1(t)}$$

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## G-S Orthogonalization Procedure (Cont.)

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- Define the second basis function as

$$\phi_2(t) = \frac{g_2(t)}{\sqrt{\int_0^T g_2^2(t) dt}} = \frac{s_2(t) - s_{21}\phi_1(t)}{\sqrt{E_2 - s_{21}^2}}$$

$$s_2(t) = s_{21}\phi_1(t) + \sqrt{E_2 - s_{21}^2}\phi_2(t) = s_{21}\phi_1(t) + s_{22}\phi_2(t)$$

- Continuing in this fashion, we have

$$g_i(t) = s_i(t) - \sum_{j=1}^{i-1} s_{ij}\phi_j(t)$$

$$s_{ij} = \int_0^T s_i(t)\phi_j(t) dt, \quad j = 1, 2, \dots, i-1$$

- Define the set of basis functions (an orthonormal set)

$$\phi_i(t) = \frac{g_i(t)}{\sqrt{\int_0^T g_i^2(t) dt}}, \quad i = 1, 2, \dots, N$$

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## G-S Orthogonalization Procedure (Cont.)

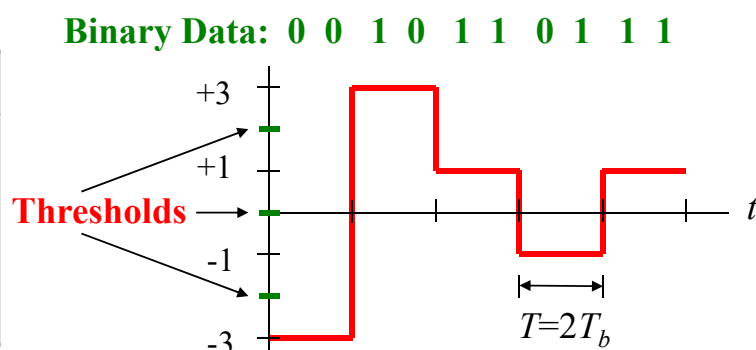
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- The dimension  $N$  is less than or equal to the number  $M$ 
  - If the signals  $s_1(t), s_2(t), \dots, s_M(t)$  form a **linearly independent** set, we have  $N = M$
  - If the signals  $s_1(t), s_2(t), \dots, s_M(t)$  are **not linearly independent**, we have  $N < M$  and the function  $g_i(t)$  is zero for  $i > N$  ( $s_i(t), i > N$ , is fully expanded by  $\phi_i(t), i = 1, \dots, M$ )
- The form of the basis functions  $\phi_1(t), \phi_2(t), \dots, \phi_N(t)$  has **not been specified**
  - $\phi_i(t)$  is **not restricted to** be either sinusoidal or sinc functions of time
- The expansion of the signal  $s_i(t)$  is an **exact expression** where  $N$  and only  $N$  terms are significant

## Example 2: 2B1Q Code

- For **baseband transmission**, the **Gray-encoded 2B1Q** code is used as the quaternary PAM signal
  - Gray encode**: any symbol differs from an adjacent symbol in a **single bit position**

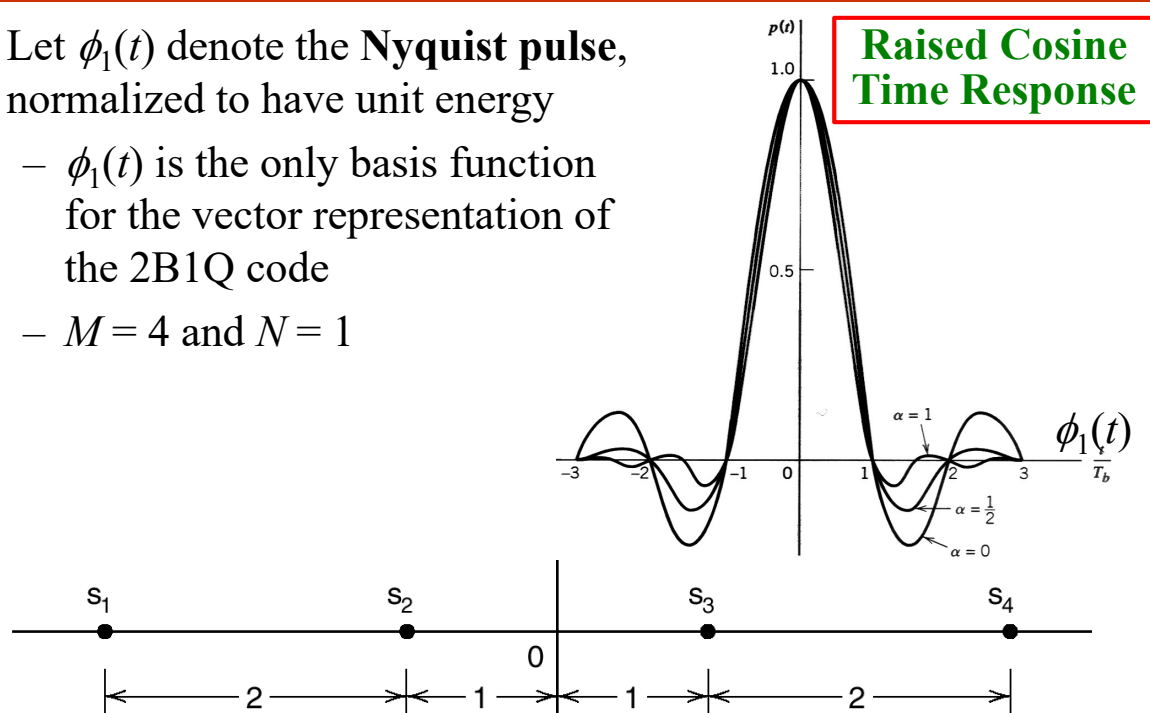
Symbol	Amplitude
10	+3
11	+1
01	-1
00	-3



- The four possible signals,  $s_1(t)$ ,  $s_2(t)$ ,  $s_3(t)$  and  $s_4(t)$ , are amplitude-scaled versions of a **Nyquist pulse**

## Example 2: 2B1Q Code (Cont.)

- Let  $\phi_1(t)$  denote the **Nyquist pulse**, normalized to have unit energy
  - $\phi_1(t)$  is the only basis function for the vector representation of the 2B1Q code
  - $M = 4$  and  $N = 1$



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# Conversion of the Continuous AWGN Channel into a Vector Channel

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## Correlator Outputs of the Received Signal

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- Let the input to the bank of  $N$  product correlators (**analyzer**) be the received signal  $x(t) = s_i(t) + w(t)$ 
  - where  $w(t)$  is a sample function of a white Gaussian noise process  $W(t)$  of **zero mean** and **power spectral density**  $N_0/2$
- The output of correlator  $j$  is the sample value of a random variable  $X_j$

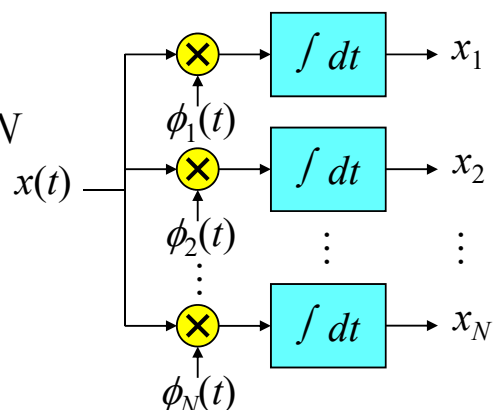
$$\begin{aligned}x_j &= \int_0^T x(t)\phi_j(t) dt \\ &= s_{ij} + w_j, \quad j = 1, 2, \dots, N\end{aligned}$$

– where

$$s_{ij} = \int_0^T s_i(t)\phi_j(t) dt$$

– The sample value of noise  $W_j$

$$w_j = \int_0^T w(t)\phi_j(t) dt$$



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## Correlator Outputs of the Received Signal(Cont.)

- Consider a new random process  $\tilde{X}(t)$  whose sample function is

$$\begin{aligned}\tilde{x}(t) &= x(t) - \sum_{j=1}^N x_j \phi_j(t) \\ &= \underline{s_i(t)} + w(t) - \sum_{j=1}^N (\underline{s_{ij}} + w_j) \phi_j(t) \\ &= w(t) - \sum_{j=1}^N w_j \phi_j(t) \triangleq w'(t)\end{aligned}$$

- which depends solely on the channel noise  $w(t)$
- The received signal can be expressed as
$$x(t) = \sum_{j=1}^N x_j \phi_j(t) + \tilde{x}(t) \triangleq \sum_{j=1}^N x_j \phi_j(t) + w'(t)$$
  - $w'(t)$  is the **remainder term** that **cannot be expanded** by the selected basis functions

## Statistical Characterization

- The random process  $X(t)$  is a **Gaussian process**
  - $X_j$  is a Gaussian random variable for all  $j$

- The **mean** of  $X_j$  depends only on  $s_{ij}$ ,

$$\mu_{X_j} = E[X_j] = E[s_{ij} + W_j] = s_{ij} + E[W_j] = s_{ij}$$

- The **variance** of  $X_j$  is

$$\sigma_{X_j}^2 = \text{var}[X_j] = E[(X_j - s_{ij})^2] = E[W_j^2]$$

- Note that the random variable  $W_j$  is defined as

$$W_j = \int_0^T W(t) \phi_j(t) dt$$

- Therefore, we have

$$\begin{aligned}\sigma_{X_j}^2 &= E\left[\int_0^T W(t) \phi_j(t) dt \int_0^T W(u) \phi_j(u) du\right] \\ &= E\left[\int_0^T \int_0^T \phi_j(t) \phi_j(u) W(t) W(u) dt du\right]\end{aligned}$$

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## Statistical Characterization (Cont.)

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- Interchanging the order of integration and expectation:

$$\begin{aligned}\sigma_{X_j}^2 &= \int_0^T \int_0^T \phi_j(t) \phi_j(u) E[W(t)W(u)] dt du \\ &= \int_0^T \int_0^T \phi_j(t) \phi_j(u) R_W(t, u) dt du\end{aligned}$$

- where  $R_W(t, u)$  is the **autocorrelation function** of  $W(t)$

$$R_W(t, u) = \frac{N_0}{2} \delta(t - u)$$

- Thus, we obtain

$$\sigma_{X_j}^2 = \frac{N_0}{2} \int_0^T \phi_j^2(t) dt$$

- Since the basis functions  $\phi_j(t)$  have **unit energy**, we get

$$\sigma_{X_j}^2 = \frac{N_0}{2}, \quad \forall j$$

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## Statistical Characterization (Cont.)

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- Since the basis functions  $\phi_j(t)$  form an orthogonal set,  $X_j$  are **mutually uncorrelated**

$$\begin{aligned}\text{cov}[X_j X_k] &= E[(X_j - \mu_{X_j})(X_k - \mu_{X_k})] \\ &= E[(X_j - s_{ij})(X_k - s_{ik})] = E[W_j W_k] \\ &= \int_0^T \int_0^T \phi_j(t) \phi_k(u) R_W(t, u) dt du \\ &= \frac{N_0}{2} \int_0^T \int_0^T \phi_j(t) \phi_k(u) \delta(t - u) dt du \\ &= \frac{N_0}{2} \int_0^T \phi_j(t) \phi_k(t) dt \\ &= 0, \quad \text{for } j \neq k\end{aligned}$$

- Since the  $X_j$  are Gaussian random variables, they are also **statistically independent** (Property of a Gaussian Process)



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## Statistical Characterization (Cont.)

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- Define the **observation vector** of  $N$  random variables as

$$\mathbf{X} = [X_1 \ X_2 \ \cdots \ X_N]^T$$

- The conditional probability density function (pdf) of  $\mathbf{X}$ , given that  $s_i(t)$  or correspondingly the symbol  $m_i$  was transmitted, is

$$f_{\mathbf{X}}(\mathbf{x}|m_i) = \prod_{j=1}^N f_{X_j}(x_j|m_i), \quad i = 1, 2, \dots, M$$

**Statistically independent**

- Since each  $X_j$  is a Gaussian random variable

$$f_{X_j}(x_j|m_i) = \frac{1}{\sqrt{\pi N_0}} \exp \left[ -\frac{1}{N_0} (x_j - s_{ij})^2 \right]$$

- The conditional pdf of  $\mathbf{X}$  is

$$f_{\mathbf{X}}(\mathbf{x}|m_i) = (\pi N_0)^{-N/2} \exp \left[ -\frac{1}{N_0} \sum_{j=1}^N (x_j - s_{ij})^2 \right]$$

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## Statistical Characterization (Cont.)

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- Note that the observation vector  $\mathbf{X}$  completely characterizes the received signal  $x(t)$ , except for the **remaining noise term**  $w'(t)$

$$x(t) = \sum_{j=1}^N x_j \phi_j(t) + w'(t)$$

- Since the noise process  $W(t)$  is Gaussian with zero mean, the noise process  $W'(t)$ , with the sample function  $w'(t)$ , is also a **zero-mean Gaussian** process
- Any random variable  $W'(t_k)$ , derived from  $W'(t)$ , is **statistically independent** of the set of random variables  $\{X_j\}$ , i.e.,

$$E[X_j W'(t_k)] = 0, \quad j = 1, 2, \dots, N$$

- $W'(t_k)$  is **irrelevant to** the message decision
  - $W'(t_k)$  is outside the  $N$ -dimensional signal space
  - The  $N$  correlator outputs are used for **decision-making**

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## Theorem of Irrelevance

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- **Theorem of irrelevance:** For signal detection in **additive white Gaussian noise**, only the projections of the noise onto the basis functions of the signal set  $\{s_i(t)\}$  affects the **sufficient statistics** (i.e., the statistics of  $\mathbf{X}$ ) of the detection problem; **the remainder of the noise is irrelevant**.
- The **AWGN channel** is equivalent to an  **$N$ -dimensional vector channel** described by the observation vector

$$\mathbf{x} = \mathbf{s}_i + \mathbf{w}, \quad i = 1, 2, \dots, M$$

$$\mathbf{w} = [w_1 \ w_2 \ \cdots \ w_N]$$

- where the dimension  $N$  is **the number of basis functions** involved in formulating the signal vector  $\mathbf{s}_i$

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## Likelihood Functions

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## Likelihood Functions

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- **At the receiver**, we are given the observation vector  $\mathbf{x}$ 
  - The requirement is **to estimate the message symbol**  $m_i$  that is responsible for generating  $\mathbf{x}$
- The definition of the **likelihood function**
$$L(m_i) = f_{\mathbf{x}}(\mathbf{x}|m_i), \quad i = 1, 2, \dots, M$$
  - The **possibility** that the message symbol  $m_i$  was transmitted when the observation vector is  $\mathbf{x}$
- In practice, we generally use the **log-likelihood function**
$$l(m_i) = \log L(m_i), \quad i = 1, 2, \dots, M$$
- The log-likelihood function bears a **one-to-one mapping** to the likelihood function
  - A probability density function is always **nonnegative**
  - The logarithmic function is **monotonically increasing**

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## Likelihood Functions

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- For an **AWGN channel**, the log-likelihood function is

$$l(m_i) = -\frac{1}{N_0} \sum_{j=1}^N (x_j - s_{ij})^2, \quad i = 1, 2, \dots, M$$

- where the constant term  $-(N/2)\log(\pi N_0)$  is ignored
- The constant term is the same for different values of  $m_i$
- It bears no relation whatsoever to the message symbol  $m_i$

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# Coherent Detection of Signals in Noise: Maximum Likelihood Decoding

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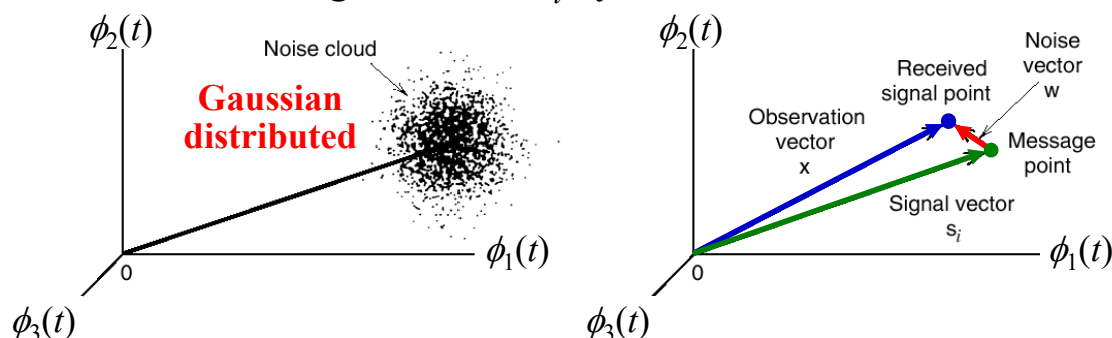
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## Signal Points

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- The transmitted signal  $s_i(t)$  can be represented as a point in a **Euclidean space** of dimension  $N \leq M$ 
  - The **transmitted signal point** or **message point** of  $s_i(t)$
- The set of message points corresponding to the set of transmitted signals is called as a **signal constellation**
- The observation vector  $\mathbf{x}$  (**received signal point**) differs from the transmitted signal vector  $\mathbf{s}_i$  by a **random noise vector**  $\mathbf{w}$



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## Signal Detection – MAP Decision Rule

- Given the observation vector  $\mathbf{x}$ , perform a mapping from  $\mathbf{x}$  to an **estimate**  $\hat{m}$  of the transmitted symbol  $m_i$ 
  - In a way that would **minimize the probability of error** in the decision-making process

$$P_e(m_i|\mathbf{x}) = P(m_i \text{ not sent}|\mathbf{x}) = 1 - P(m_i \text{ sent}|\mathbf{x})$$

- We can state the **optimum decision rule** as:

$$\text{Set } \hat{m} = m_i \text{ if } P(m_i \text{ sent}|\mathbf{x}) \geq P(m_k \text{ sent}|\mathbf{x}) \text{ for all } k \neq i$$

- This decision rule is referred to as the **maximum a posteriori probability (MAP)** rule

## Signal Detection – MAP Decision Rule (Cont.)

- Using **Bayes' rule**, we may restate the MAP rule as follows:

$$\text{Set } \hat{m} = m_i \text{ if } \frac{p_k f_{\mathbf{X}}(\mathbf{x}|m_k)}{f_{\mathbf{X}}(\mathbf{x})} \text{ is maximum for } k = i$$

**The *a priori* probabilities are required**

$$P(A|B) = \frac{P(A \cap B)}{P(B)} = \frac{P(A)P(A \cap B)}{P(A)P(B)} = \frac{P(A)P(B|A)}{P(B)}$$

- $p_k$  is the ***a priori* probability** of transmitting symbol  $m_k$
- $f_{\mathbf{X}}(\mathbf{x}|m_k)$  is the **conditional pdf** of  $\mathbf{X}$  given the transmission of  $m_k$ , and  $f_{\mathbf{X}}(\mathbf{x})$  is the unconditional pdf of  $\mathbf{X}$
- The denominator  $f_{\mathbf{X}}(\mathbf{x})$  is independent of the transmitted symbol
- If all the symbols are **equally likely**,  $p_k = p_i$  for all  $i$  and  $k$ 
  - The conditional pdf  $f_{\mathbf{X}}(\mathbf{x}|m_k)$  bears a **one-to-one mapping** to the **log-likelihood function**  $l(m_k)$

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# Maximum Likelihood Decision Rule

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- Accordingly, we can restate the decision rule as follows:

Set  $\hat{m} = m_i$  if  
 $l(m_k)$  is maximum for  $k = i$

- This decision rule is referred to as the **maximum likelihood (ML) rule** and the device for its implementation is the **maximum likelihood decoder**
- Based on the observation vector  $\mathbf{x}$ , the decoder **computes** the log-likelihood functions as metrics for all the  $M$  possible message symbols, **compares** them, and then **decides** in favor of the maximum
- The ML decoder differs from the MAP decoder in that **it assumes equally likely message symbols**
  - The ML decoder **does not require** the *a priori* probabilities

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## Comparison: MAP and ML Decision Rules

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- If the transmitting message symbols are **equally likely**, the MAP and ML decision rules have **the same** performance.
- If the transmitting message symbols are **not equally likely**, the MAP decision rule is **superior to** the ML decision rule
  - Since the information of *a priori* probabilities is available
- For example, there are two message symbols with the *a priori* probabilities  $p_0 = 0.9999$  and  $p_1 = 0.0001$ 
  - Considering the *a priori* probabilities is very important
- However, in general, the transmitting message symbols are **equally likely** for practical systems
  - The ML decision rule is commonly used

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## Maximum Likelihood Decision Rule (Cont.)

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- Let  $Z$  denote the  $N$ -dimensional space (**observation space**) of all possible observation vectors  $\mathbf{x}$
- The observation space  $Z$  is partitioned into  $M$ -**decision regions**,  $Z_1, Z_2, \dots, Z_M$  (which are **non-overlapping**)
  - Accordingly, we can restate the ML decision rule as follows:  
Observation vector  $\mathbf{x}$  lies in region  $Z_i$  if  
$$l(m_k) \text{ is maximum for } k = i$$
- For an AWGN channel, the LLF  $l(m_k)$  attains its maximum value when  $\sum_{j=1}^N (x_j - s_{kj})^2$  is **minimized by the choice**  $k = i$ 
  - Accordingly, we can restate the ML decision rule as follows:  
Observation vector  $\mathbf{x}$  lies in region  $Z_i$  if  
$$\sum_{j=1}^N (x_j - s_{kj})^2 \text{ is minimum for } k = i$$

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## Maximum Likelihood Decision Rule (Cont.)

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- $\sum_{j=1}^N (x_j - s_{kj})^2$  is the Euclidean distance square between  $\mathbf{x}$  and  $\mathbf{s}_k$ 
  - Accordingly, we can restate the ML decision rule as follows:  
Observation vector  $\mathbf{x}$  lies in region  $Z_i$  if  
the Euclidean distance  $\|\mathbf{x} - \mathbf{s}_k\|$  is minimum for  $k = i$
- The ML decision rule is simply **to choose the message point closest to the received signal point**
- Note that 
$$\sum_{j=1}^N (x_j - s_{kj})^2 = \sum_{j=1}^N x_j^2 - 2 \sum_{j=1}^N x_j s_{kj} + \sum_{j=1}^N s_{kj}^2$$

**Irrelevant to  $k$**  

  - Accordingly, we can restate the ML decision rule as follows:  
Observation vector  $\mathbf{x}$  lies in region  $Z_i$  if  
$$\sum_{j=1}^N x_j s_{kj} - E_k/2 \text{ is maximum for } k = i$$

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## Maximum Likelihood Decision Rule (Cont.)

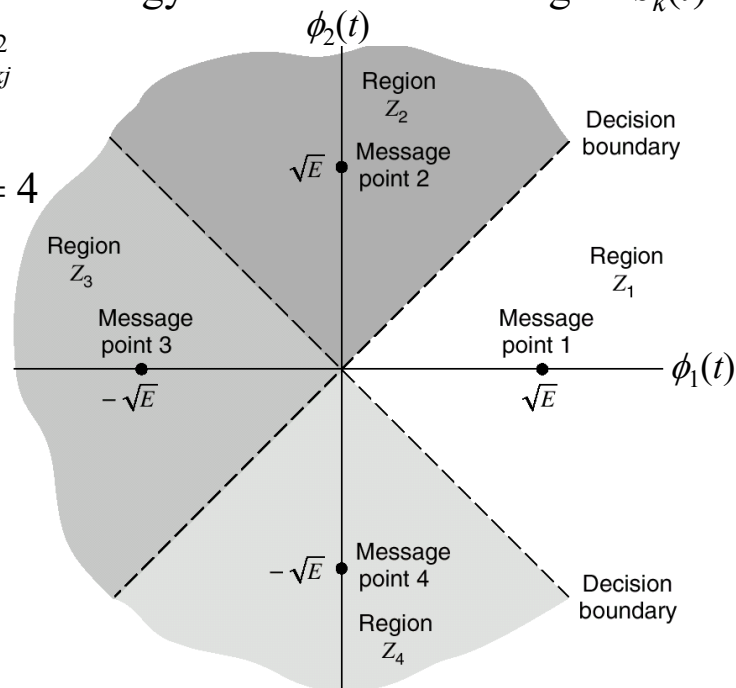
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– where  $E_k$  is the energy of the transmitted signal  $s_k(t)$

$$E_k = \sum_{j=1}^N s_{kj}^2$$

• Example:

$N = 2$  and  $M = 4$



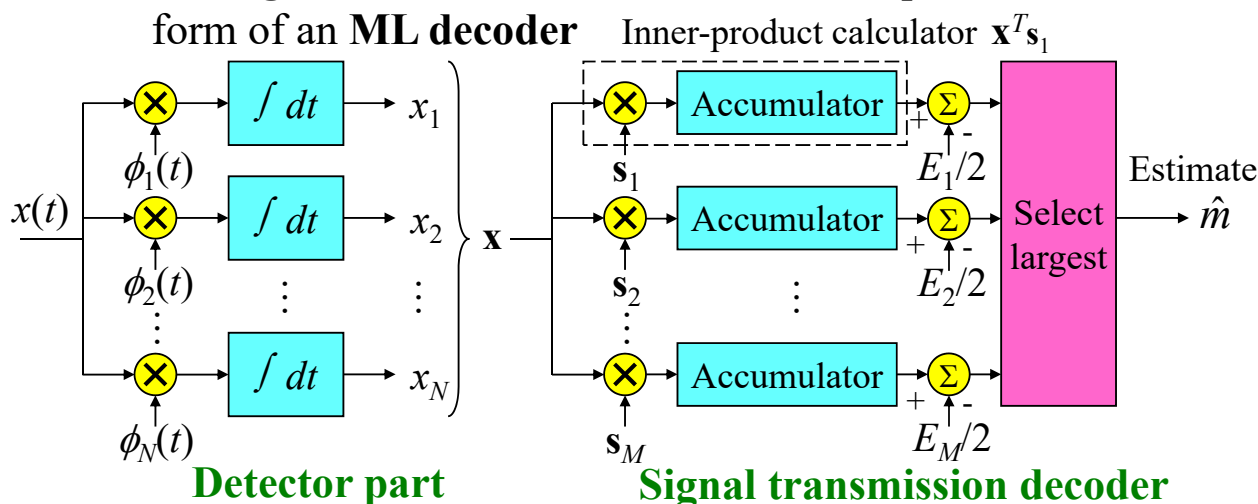
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## Correlation Receiver



# Optimum Receiver

- The optimum receiver consists of two subsystems:
  - The **detector part**: It consists of a bank of  $M$  product-integrators or **correlators**
  - The **signal transmission decoder**: It is implemented in the form of an **ML decoder**



## Equivalence: Correlation – Matched Filter

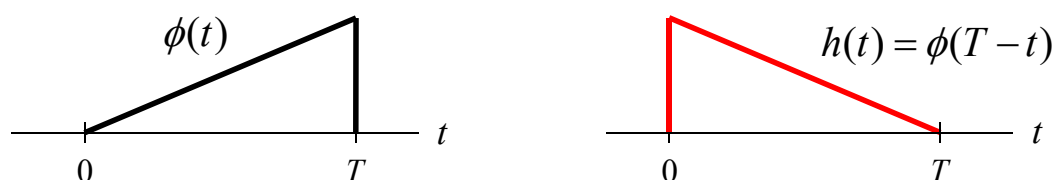
- We can use a corresponding set of **matched filters** to build the detector
- Consider a linear time-invariant filter with impulse response  $h_j(t)$
- When the received signal  $x(t)$  is used as the filter input, the resulting filter output  $y_j(t)$  is

$$y_j(t) = \int_{-\infty}^{\infty} x(\tau) h_j(t - \tau) d\tau$$

$$\begin{aligned} x(t) &= s_i(t) + w(t) \\ s_i(t) &= \sum_{j=1}^N s_{ij} \phi_j(t) \end{aligned}$$

- The impulse response  $h_j(t)$  **matched to** an input signal  $\phi_j(t)$  is

$$h_j(t) = \phi_j(T - t)$$



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## Equivalence: Correlation – Matched Filter(Cont.)

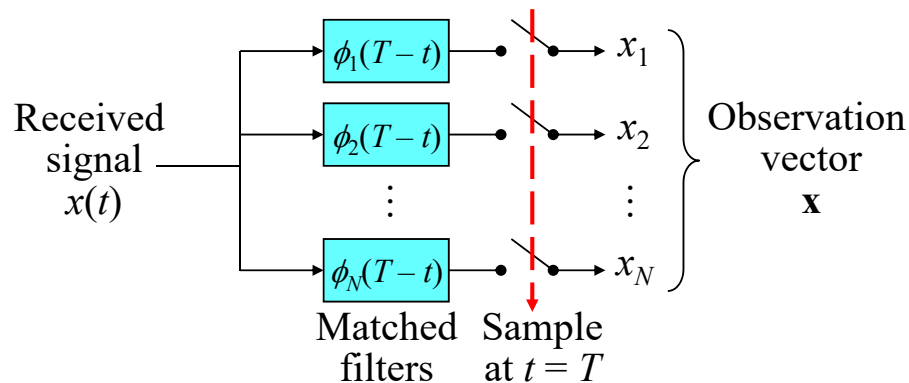
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- Then the filter output is  $y_j(t) = \int_{-\infty}^{\infty} x(\tau)\phi_j(T-t+\tau) d\tau$
- Sampling the output at time  $t = T$ , we get

$$y_j(T) = \int_{-\infty}^{\infty} x(\tau)\phi_j(\tau) d\tau$$

- Since  $\phi_j(t)$  is zero outside the interval  $0 \leq t \leq T$

$$y_j(T) = \int_0^T x(\tau)\phi_j(\tau) d\tau$$



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## Probability of Error

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## Noise Performance

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- Suppose that symbol  $m_i$  is transmitted and an observation vector  $\mathbf{x}$  is received
  - An **error** occurs whenever the received signal point does not fall inside region  $Z_i$
- The average probability of symbol error is

$$\begin{aligned} P_e &= \sum_{i=1}^M p_i P(\mathbf{x} \text{ does not lie in } Z_i | m_i \text{ sent}) \\ &= 1 - \frac{1}{M} \sum_{i=1}^M P(\mathbf{x} \text{ lies in } Z_i | m_i \text{ sent}) \end{aligned} \quad \boxed{p_i = \frac{1}{M}, \quad \forall i}$$

- Since  $\mathbf{x}$  is the sample value of random vector  $\mathbf{X}$ ,  $P_e$  can be expressed in terms of the **likelihood function** as follows:

$$P_e = 1 - \frac{1}{M} \sum_{i=1}^M \int_{Z_i} f_{\mathbf{x}}(\mathbf{x} | m_i) d\mathbf{x}$$

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## Invariance to Rotation and Translation

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- For the ML detection of a signal in AWGN, **changes in the orientation of the signal constellation** with respect to both the **coordinate axes** and **origin** of the signal space **do not affect** the probability of symbol error  $P_e$ 
  - In ML detection,  $P_e$  depends solely on the **relative Euclidean distances** between the message points
  - The AWGN is **spherically symmetric** in all directions
- The effect of a **rotation** applied to all the message points is equivalent to multiplying the signal vector  $\mathbf{s}_i$  by an  **$N$ -by- $N$  orthonormal matrix  $\mathbf{Q}$**  for all  $i$ 
  - where the matrix  $\mathbf{Q}$  satisfies

$$\mathbf{Q}\mathbf{Q}^T = \mathbf{I} \longleftarrow \text{Identity matrix}$$

## Invariance to Rotation and Translation (Cont.)

- The signal (noise) vector  $\mathbf{s}_i$  ( $\mathbf{w}$ ) is replaced by the rotated version

$$\begin{aligned}\mathbf{s}_{i,\text{rotate}} &= \mathbf{Q}\mathbf{s}_i, \quad i = 1, 2, \dots, M \\ \mathbf{w}_{\text{rotate}} &= \mathbf{Q}\mathbf{w}\end{aligned}$$

- The statistical characteristics of the noise vector are **unaffected**

$$E[\mathbf{w}_{\text{rotate}}] = \mathbf{0}; \quad E[\mathbf{w}_{\text{rotate}} \mathbf{w}_{\text{rotate}}^T] = \frac{N_0}{2} \mathbf{I}$$

- The observation vector for the **rotated signal constellation** is

$$\mathbf{x}_{\text{rotate}} = \mathbf{Q}\mathbf{s}_i + \mathbf{w}_{\text{rotate}} = \mathbf{Q}\mathbf{s}_i + \mathbf{w}, \quad i = 1, 2, \dots, M$$

- The Euclidean distance between  $\mathbf{x}_{\text{rotate}}$  and  $\mathbf{s}_{i,\text{rotate}}$  is

$$\|\mathbf{x}_{\text{rotate}} - \mathbf{s}_{i,\text{rotate}}\| = \|\mathbf{x} - \mathbf{s}_i\|, \quad \text{for all } i$$

- Also note that

$$\mathbf{x} = \mathbf{Q}^T \mathbf{x}_{\text{rotate}} = \mathbf{Q}^T \mathbf{Q} \mathbf{s}_i + \mathbf{Q}^T \mathbf{w} = \mathbf{s}_i + \mathbf{w}'_{\text{rotate}} = \mathbf{s}_i + \mathbf{w}$$

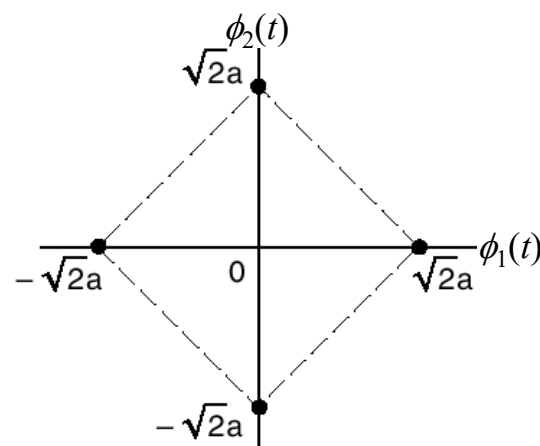
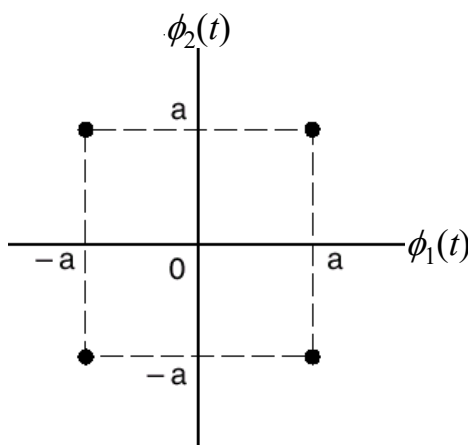
## Invariance to Rotation and Translation (Cont.)

- For translation, we have

$$\mathbf{s}_{i,\text{translate}} = \mathbf{s}_i - \mathbf{a}, \quad i = 1, 2, \dots, M$$

$$\mathbf{x}_{\text{translate}} = \mathbf{x} - \mathbf{a},$$

$$\|\mathbf{x}_{\text{translate}} - \mathbf{s}_{i,\text{translate}}\| = \|\mathbf{x} - \mathbf{s}_i\|, \quad \text{for all } i$$



# Minimum Energy Signals

- Based on the principle of **translational invariance**, we can translate the signal constellation to **minimize** the average energy

- The average energy of the signal constellation translated by a vector  $\mathbf{a}$  is  $\mathcal{E}_{\text{translate}} = \sum_{i=1}^M \|\mathbf{s}_i - \mathbf{a}\|^2 p_i$

$$\|\mathbf{s}_i - \mathbf{a}\|^2 = \|\mathbf{s}_i\|^2 - 2\mathbf{a}^T \mathbf{s}_i + \|\mathbf{a}\|^2$$

$$\mathcal{E}_{\text{translate}} = \sum_{i=1}^M \|\mathbf{s}_i\|^2 p_i - 2 \sum_{i=1}^M \mathbf{a}^T \mathbf{s}_i p_i + \sum_{i=1}^M \|\mathbf{a}\|^2 p_i = \mathcal{E} - 2\mathbf{a}^T E[\mathbf{s}] + \|\mathbf{a}\|^2$$

– where  $\mathcal{E}$  is the average energy of the **original** signal

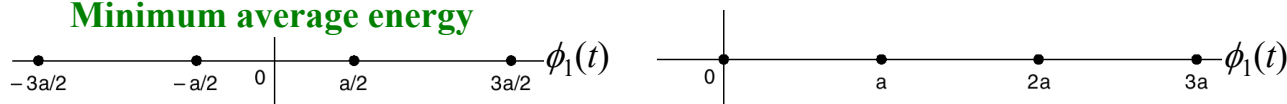
constellation and  $E[\mathbf{s}] = \sum_{i=1}^M \mathbf{s}_i p_i$

- Differentiating  $\mathcal{E}_{\text{translate}}$  with respect to  $\mathbf{a}$  and setting it to zero

– The minimizing translate is  $\mathbf{a}_{\min} = E[\mathbf{s}]$

– The minimum average energy  $\mathcal{E}_{\text{translate,min}} = \mathcal{E} - \|\mathbf{a}_{\min}\|^2 = E[\mathbf{s}]$

**Minimum average energy**



# Union Bound on the Probability of Error

- The average probability of symbol error  $P_e$  is

$$P_e = 1 - \frac{1}{M} \sum_{i=1}^M \int_{Z_i} f_{\mathbf{x}}(\mathbf{x} | m_i) d\mathbf{x}$$

- Numerical computation of the integral may be **impractical**

– We can approximate  $P_e$  by **simplifying the integral** or **simplifying the region of integration**

- Union bound:** a simple upper bound that bases on simplifying the region of integration

**$A_{ik}$  and  $A_{ij}$  may be overlapped**

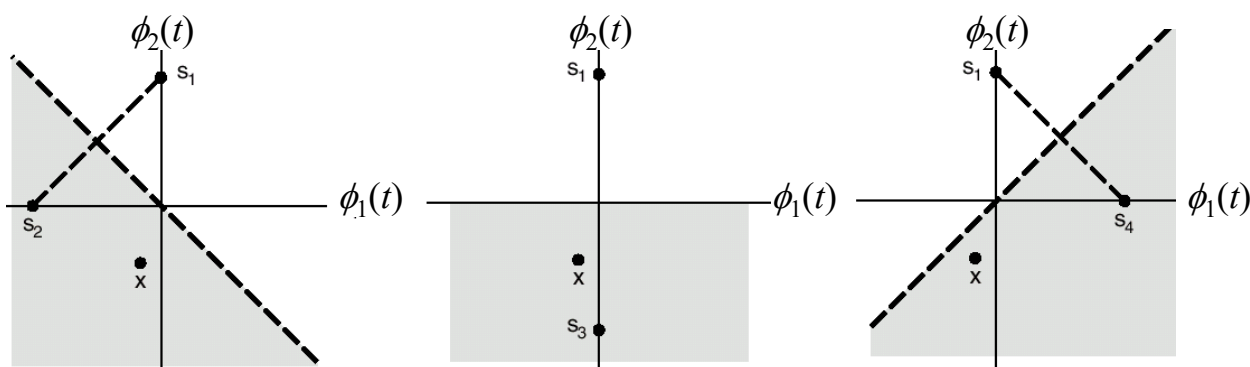
- Let  $A_{ik}$ ,  $i, k \in \{1, 2, \dots, M\}$ , denote the event that the observation vector  $\mathbf{x}$  is **closer to  $\mathbf{s}_k$  than to  $\mathbf{s}_i$**  when the symbol  $m_i$  is sent

- The conditional probability of symbol error  $P_e(m_i)$  is equal to the probability of the **union events**  $A_{i1}, A_{i2}, \dots, A_{i,i-1}, A_{i,i+1}, \dots, A_{iM}$

## Union Bound on the Probability of Error (Cont.)

- The probability of a finite **union of events** is overbounded by the **sum of the probabilities** of the constituent events

$$P_e(m_i) \leq \sum_{k=1, k \neq i}^M P(A_{ik}), \quad i = 1, 2, \dots, M$$



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## Union Bound on the Probability of Error (Cont.)

- Note that the probability  $P(A_{ik})$  is **different from** the probability  $P(\hat{m} = m_k | m_i)$
- $P(A_{ik}) \triangleq P_2(\mathbf{s}_i, \mathbf{s}_k)$  is the **pairwise error probability** in that the system uses only a pair of signals  $\mathbf{s}_i$  and  $\mathbf{s}_k$

$$\begin{aligned} P_2(\mathbf{s}_i, \mathbf{s}_k) &= P(\mathbf{x} \text{ is closer to } \mathbf{s}_k \text{ than } \mathbf{s}_i, \text{ when } \mathbf{s}_i \text{ is sent}) \\ &= \int_{d_{ik}/2}^{\infty} \frac{1}{\sqrt{\pi N_0}} \exp(-v^2/N_0) dv \end{aligned}$$

–  $d_{ik} = \|\mathbf{s}_i - \mathbf{s}_k\|$  is the Euclidean distance between  $\mathbf{s}_i$  and  $\mathbf{s}_k$

- Setting  $z = v/\sqrt{N_0}$ ,

$$P_2(\mathbf{s}_i, \mathbf{s}_k) = \frac{1}{2} \operatorname{erfc}(d_{ik}/2\sqrt{N_0})$$

$$\operatorname{erfc}(u) = \frac{2}{\sqrt{\pi}} \int_u^{\infty} \exp(-z^2) dz$$

$$\Rightarrow P_e(m_i) \leq \frac{1}{2} \sum_{k=1, k \neq i}^M \operatorname{erfc}(d_{ik}/2\sqrt{N_0}), \quad i = 1, 2, \dots, M$$

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## Union Bound on the Probability of Error (Cont.)

- The probability of symbol error is overbounded as follows:

$$P_e = \sum_{i=1}^M p_i P_e(m_i) \leq \frac{1}{2} \sum_{i=1}^M \sum_{k=1, k \neq i}^M p_i \operatorname{erfc}\left(d_{ik}/2\sqrt{N_0}\right)$$

- Suppose that the signal constellation is **circularly symmetric** about the origin  $\Rightarrow P_e(m_i)$  is the same for all  $i$

$$P_e \leq \frac{1}{2} \sum_{k=1, k \neq i}^M \operatorname{erfc}\left(d_{ik}/2\sqrt{N_0}\right)$$

- Define  $d_{\min}$  as the **minimum distance** between any two signals

$$P_e \leq \frac{M-1}{2} \operatorname{erfc}\left(d_{\min}/2\sqrt{N_0}\right) \quad d_{\min} = \min_{k \neq i} d_{ik}, \text{ for all } i \text{ and } k$$

- We can also further simplify the union bound on  $P_e$  as

$$P_e \leq \frac{M-1}{2\sqrt{\pi}} \exp\left(-\underline{d_{\min}^2}/4N_0\right) \quad \operatorname{erfc}\left(\frac{d_{\min}}{2\sqrt{N_0}}\right) \leq \frac{1}{\sqrt{\pi}} \exp\left(-\frac{d_{\min}^2}{4N_0}\right)$$

## Bit versus Symbol Error Probabilities

- For binary data transmission, it is more meaningful to consider the **bit error rate (BER)**: there are  $K = \log_2 M$  bits per symbol
- Case 1: Gray encode** is applied
  - Any two adjacent symbols differ in only **one bit** position
  - Given a symbol error, the most probable number of bit errors is **1**  $\Rightarrow$  The BER is bounded as follows:

**Only 1 bit is in error**  $\rightarrow P_e/\log_2 M \leq \text{BER} \leq P_e \leftarrow$  **All bits are in error**

- Case 2:** All symbol errors occurs **equally likely**
  - The occurrence probability is  $P_e/(M-1) = P_e/(2^K - 1)$
  - There are  $2^{K-1}$  error symbols that the  $i$ -th bit is in error
  - The bit error rate is  $\text{BER} = \frac{2^{K-1}}{2^K - 1} P_e = \frac{M/2}{M-1} P_e \approx \underline{P_e/2}$