

# Python ICPC Cheatsheet

Comprehensive Reference for Competitive Programming

## Contents

### 1 Input/Output

**Description:** Efficient input/output is crucial in competitive programming especially for problems with large datasets. Using `sys.stdin.readline` is significantly faster than the default `input()` function.

```

1 # Fast I/O - Essential for large inputs
2 import sys
3 input = sys.stdin.readline
4
5 # Read single integer
6 n = int(input())
7
8 # Read multiple integers on one line
9 a, b = map(int, input().split())
10
11 # Read array of integers
12 arr = list(map(int, input().split()))
13
14 # Read strings (strip to remove trailing newline)
15 s = input().strip()
16 words = input().split()
17
18 # Multiple test cases pattern
19 t = int(input())
20 for _ in range(t):
21     # process each test case
22
23 # Print without newline
24 print(x, end=' ')
25
26 # Formatted output with precision
27 print(f"{x:.6f}") # 6 decimal places

```

### 2 Basic Data Structures

#### 2.1 List Operations

**Description:** Python lists are dynamic arrays with O(1) amortized append and O(n) insert/delete at arbitrary positions.

```

1 # Initialize lists
2 arr = [0] * n # n zeros
3 matrix = [[0] * m for _ in range(n)] # Correct way!
4
5 # List comprehension - concise and efficient
6 squares = [x**2 for x in range(n)]
7 evens = [x for x in arr if x % 2 == 0]
8
9 # Sorting - O(n log n)
10 arr.sort() # in-place, modifies arr
11 arr.sort(reverse=True) # descending
12 arr.sort(key=lambda x: (x[0], -x[1])) # custom
13 sorted_arr = sorted(arr) # returns new list
14
15 # Binary search in sorted array
16 from bisect import bisect_left, bisect_right
17 idx = bisect_left(arr, x) # leftmost position
18 idx = bisect_right(arr, x) # rightmost position
19
20 # Common operations
21 arr.append(x) # O(1) amortized
22 arr.pop() # O(1) - remove last
23 arr.pop(0) # O(n) - remove first (slow!)
24 arr.reverse() # O(n) - in-place
25 arr.count(x) # O(n) - count occurrences
26 arr.index(x) # O(n) - first occurrence

```

#### 2.2 Deque (Double-ended Queue)

**Description:** Deque (pronounced "deck") provides O(1) append and pop operations from both ends, unlike lists which have O(n) for operations at the front. Essential for BFS, sliding window problems, implementing efficient queues/stacks, and maintaining monotonic queues. Use when you need fast insertions/deletions at both ends.

```

1 from collections import deque
2 dq = deque()
3
4 # O(1) operations on both ends
5 dq.append(x) # add to right
6 dq.appendleft(x) # add to left
7 dq.pop() # remove from right
8 dq.popleft() # remove from left
9
10 # Sliding window maximum - O(n)

```

```

11 # Maintains decreasing order of elements
12 def sliding_max(arr, k):
13     dq = deque() # stores indices
14     result = []
15
16     for i in range(len(arr)):
17         # Remove indices outside window
18         while dq and dq[0] < i - k + 1:
19             dq.popleft()
20
21         # Remove smaller elements (not useful)
22         while dq and arr[dq[-1]] < arr[i]:
23             dq.pop()
24
25         dq.append(i)
26         if i >= k - 1:
27             result.append(arr[dq[0]])
28
29     return result

```

#### 2.3 Heap (Priority Queue)

**Description:** Python's `heapq` module implements a min-heap (smallest element always at index 0). Provides O(log n) insert and extract-min operations, O(n) heapify, and O(1) peek. For max-heap, negate values before insertion. Critical for Dijkstra's algorithm, Prim's MST, k-th largest/smallest problems, merge k sorted lists, and any problem requiring repeated access to minimum-/maximum elements. More efficient than sorting when you only need partial ordering.

```

1 import heapq
2
3 # Min heap (default)
4 heap = []
5 heapq.heappush(heap, x) # O(log n)
6 min_val = heapq.heappop(heap) # O(log n)
7 min_val = heap[0] # O(1) peek
8
9 # Max heap - negate values
10 heapq.heappush(heap, -x)
11 max_val = -heapq.heappop(heap)
12
13 # Convert list to heap in-place - O(n)
14 heapq.heapify(arr)
15
16 # K largest/smallest - O(n log k)
17 k_largest = heapq.nlargest(k, arr)
18 k_smallest = heapq.nsmallest(k, arr)
19
20 # Custom comparator using tuples
21 # Compares first element, then second, etc.
22 heapq.heappush(heap, (priority, item))

```

#### 2.4 Dictionary & Counter

**Description:** Hash maps with O(1) average case insert/lookup. Counter is specialized for counting occurrences.

```

1 from collections import defaultdict, Counter
2
3 # defaultdict - provides default value
4 graph = defaultdict(list) # empty list default
5 count = defaultdict(int) # 0 default
6
7 # Counter - count elements efficiently
8 cnt = Counter(arr)
9 cnt['x'] += 1
10 most_common = cnt.most_common(k) # k most frequent
11
12 # Dictionary operations
13 d = {}
14 d.get(key, default_val)
15 d.setdefault(key, default_val)
16 for k, v in d.items():
17     pass

```

#### 2.5 Set Operations

**Description:** Hash sets provide O(1) average-case membership testing, insertion, and deletion. Unlike lists, sets store only unique elements (no duplicates) and are unordered. Essential for removing duplicates, fast membership queries, and mathematical set operations (union, intersection, difference). Use when element uniqueness matters or you need fast lookups without caring about order. For sorted sets, consider using sorted containers or maintaining a sorted list separately.

```

1 s = set()
2 s.add(x)          # O(1)
3 s.remove(x)       # O(1), KeyError if not exists
4 s.discard(x)     # O(1), no error if not exists
5
6 # Set operations - all O(n)
7 a | b            # union
8 a & b            # intersection
9 a - b            # difference
10 a ^ b           # symmetric difference
11
12 # Ordered set workaround
13 from collections import OrderedDict
14 oset = OrderedDict.fromkeys([])

```

## 2.6 Skip List

**Description:** Skip lists are probabilistic balanced data structures that support  $O(\log n)$  expected time for search, insert, and delete. They are an alternative to balanced BSTs and are simple to implement. Good when you want ordered set/map operations with expected logarithmic time.

**How it works:** A skip list is composed of multiple levels of sorted linked lists; level 0 contains all elements and higher levels act as “express lanes” with a subset of elements (each element appears in a random number of levels, forming a “tower”). To search, start at the head on the highest level, move right while the next key is less than the target, and drop down one level when you can no longer move right; repeat until level 0. New node levels are chosen randomly with probability  $p$  (common choice  $p = 0.5$ ). For contests, use  $p = 0.5$  and  $L_{\max} \approx \lceil \log_2 N \rceil + 2$  for safety.

```

1 import random
2
3 class SkipNode:
4     def __init__(self, key, value=None, level=0):
5         self.key = key
6         self.value = value
7         # forward pointers for levels 0..level
8         self.forward = [None] * (level + 1)
9
10 class SkipList:
11     def __init__(self, max_level=16, p=0.5):
12         self.max_level = max_level
13         self.p = p
14         self.level = 0
15         self.head = SkipNode(None, level=max_level)
16
17     def random_level(self):
18         lvl = 0
19         while random.random() < self.p and lvl < self.max_level:
20             lvl += 1
21         return lvl
22
23     def search(self, key):
24         node = self.head
25         # start from highest level
26         for i in range(self.level, -1, -1):
27             while node.forward[i] and node.forward[i].key < key:
28                 node = node.forward[i]
29         node = node.forward[0]
30         if node and node.key == key:
31             return node.value
32         return None
33
34     def insert(self, key, value=None):
35         update = [None] * (self.max_level + 1)
36         node = self.head
37         for i in range(self.level, -1, -1):
38             while node.forward[i] and node.forward[i].key < key:
39                 node = node.forward[i]
40             update[i] = node
41         node = node.forward[0]
42
43         # If key exists, update value
44         if node and node.key == key:
45             node.value = value
46             return
47
48         lvl = self.random_level()
49         if lvl > self.level:
50             for i in range(self.level + 1, lvl + 1):
51                 update[i] = self.head
52             self.level = lvl
53
54         new_node = SkipNode(key, value, lvl)
55         for i in range(lvl + 1):
56             new_node.forward[i] = update[i].forward[i]
57             update[i].forward[i] = new_node
58
59     def delete(self, key):
60         update = [None] * (self.max_level + 1)
61         node = self.head
62         for i in range(self.level, -1, -1):
63             while node.forward[i] and node.forward[i].key < key:
64                 node = node.forward[i]

```

```

65         update[i] = node
66         node = node.forward[0]
67
68         if not node or node.key != key:
69             return False # not found
70
71         for i in range(self.level + 1):
72             if update[i].forward[i] != node:
73                 break
74             update[i].forward[i] = node.forward[i]
75
76         # adjust current level
77         while self.level > 0 and self.head.forward[self.level] is None:
78             self.level -= 1
79         return True
80
81 # Example usage:
82 # s = SkipList()
83 # s.insert(10, 'ten')
84 # print(s.search(10))
85 # s.delete(10)

```

## 3 String Operations

**Description:** Strings in Python are immutable. For building strings, use list and join for  $O(n)$  complexity instead of repeated concatenation which is  $O(n^2)$ .

```

1 # Common string methods
2 s.lower(), s.upper()
3 s.strip() # remove whitespace both ends
4 s.lstrip() # remove left whitespace
5 s.rstrip() # remove right whitespace
6 s.split(delimiter)
7 delimiter.join(list)
8 s.replace(old, new)
9 s.startswith(prefix)
10 s.endswith(suffix)
11 s.isdigit(), s.isalpha(), s.isalnum()
12
13 # String building - EFFICIENT O(n)
14 result = []
15 for x in data:
16     result.append(str(x))
17 s = ''.join(result)
18
19 # String concatenation - SLOW O(n^2)
20 # s = ""
21 # for x in data:
22 #     s += str(x) # Don't do this!
23
24 # ASCII values
25 ord('a') # 97
26 chr(97) # 'a'
27
28 # String to character array (for mutations)
29 chars = list(s)
30 chars[0] = 'x'
31 s = ''.join(chars)

```

### 3.1 KMP Pattern Matching

**Description:** Find all occurrences of pattern in text. Time:  $O(n+m)$ .

```

1 def kmp_search(text, pattern):
2     # Build LPS (Longest Proper Prefix which is Suffix)
3     def build_lps(pattern):
4         m = len(pattern)
5         lps = [0] * m
6         length = 0 # Length of previous longest prefix
7         i = 1
8
9         while i < m:
10             if pattern[i] == pattern[length]:
11                 length += 1
12                 lps[i] = length
13                 i += 1
14             else:
15                 if length != 0:
16                     length = lps[length - 1]
17                 else:
18                     lps[i] = 0
19                     i += 1
20
21         return lps
22
23 n, m = len(text), len(pattern)
24 lps = build_lps(pattern)
25
26 matches = []
27 i = j = 0 # Indices for text and pattern
28
29 while i < n:
30     if text[i] == pattern[j]:
31         i += 1

```

```

32     j += 1
33
34     if j == m:
35         matches.append(i - j)
36         j = lps[j - 1]
37     elif i < n and text[i] != pattern[j]:
38         if j != 0:
39             j = lps[j - 1]
40         else:
41             i += 1
42
43 return matches

```

## 3.2 Z-Algorithm

**Description:** Compute Z-array where  $Z[i]$  = length of longest substring starting from  $i$  that matches prefix. Time:  $O(n)$ .

```

1 def z_algorithm(s):
2     n = len(s)
3     z = [0] * n
4     l, r = 0, 0
5
6     for i in range(1, n):
7         if i <= r:
8             z[i] = min(r - i + 1, z[i - 1])
9
10    while i + z[i] < n and s[z[i]] == s[i + z[i]]:
11        z[i] += 1
12
13    if i + z[i] - 1 > r:
14        l, r = i, i + z[i] - 1
15
16 return z
17
# Pattern matching using Z-algorithm
18 def z_search(text, pattern):
19     # Concatenate pattern + $ + text
20     s = pattern + '$' + text
21     z = z_algorithm(s)
22
23     matches = []
24     m = len(pattern)
25
26     for i in range(m + 1, len(s)):
27         if z[i] == m:
28             matches.append(i - m - 1)
29
30     return matches
31

```

## 3.3 Rabin-Karp (Rolling Hash)

**Description:** Fast pattern matching using hashing. Average:  $O(n+m)$   
Worst:  $O(nm)$ .

```

1 def rabin_karp(text, pattern):
2     MOD = 10**9 + 7
3     BASE = 31 # Prime base for hashing
4
5     n, m = len(text), len(pattern)
6     if m > n:
7         return []
8
9     # Compute hash of pattern
10    pattern_hash = 0
11    power = 1
12    for i in range(m):
13        pattern_hash = (pattern_hash * BASE +
14                        ord(pattern[i])) % MOD
15    if i < m - 1:
16        power = (power * BASE) % MOD
17
18    # Rolling hash
19    text_hash = 0
20    matches = []
21
22    for i in range(n):
23        # Add new character
24        text_hash = (text_hash * BASE +
25                      ord(text[i])) % MOD
26
27        # Remove old character if window full
28        if i >= m:
29            text_hash = (text_hash -
30                          ord(text[i - m])) * power) % MOD
31            text_hash = (text_hash + MOD) % MOD
32
33        # Check match
34        if i >= m - 1 and text_hash == pattern_hash:
35            # Verify actual match (avoid hash collision)
36            if text[i - m + 1:i + 1] == pattern:
37                matches.append(i - m + 1)
38
39 return matches

```

# 4 Mathematics

## 4.1 Basic Math Operations

```

1 import math
2
3 # Common functions
4 math.ceil(x), math.floor(x)
5 math.gcd(a, b)      # Greatest common divisor
6 math.lcm(a, b)      # Python 3.9+
7 math.sqrt(x)
8 math.log(x), math.log2(x), math.log10(x)
9
10 # Powers
11 x ** y
12 pow(x, y, mod)    # (x^y) % mod - efficient modular exp
13
14 # Infinity
15 float('inf'), float('-inf')
16
17 # Custom GCD using Euclidean algorithm - O(log min(a,b))
18 def gcd(a, b):
19     while b:
20         a, b = b, a % b
21
22     return a
23
24 def lcm(a, b):
25     return a * b // gcd(a, b)

```

## 4.2 Combinatorics

**Description:** Compute combinations and permutations. For modular arithmetic, compute factorial arrays and use modular inverse.

```

1 from math import factorial, comb, perm
2
3 # nCr (combinations) - "n choose r"
4 comb(n, r) # Built-in Python 3.8+
5
6 # nPr (permutations)
7 perm(n, r) # Built-in Python 3.8+
8
9 # Manual nCr implementation
10 def ncr(n, r):
11     if r > n: return 0
12     r = min(r, n - r) # Optimization: C(n,r) = C(n,n-r)
13     num = den = 1
14     for i in range(r):
15         num *= (n - i)
16         den *= (i + 1)
17     return num // den
18
# Precompute factorials with modulo
19 MOD = 10**9 + 7
20 def modfact(n):
21     fact = [1] * (n + 1)
22     for i in range(1, n + 1):
23         fact[i] = fact[i-1] * i % MOD
24     return fact
25
26 # Modular combination using precomputed factorials
27 # First precompute inverse factorials
28 # def compute_inv_factorials(n, mod):
29 #     fact = modfact(n)
30 #     inv_fact = [1] * (n + 1)
31 #     inv_fact[n] = pow(fact[n], mod - 2, mod)
32 #     for i in range(n - 1, -1, -1):
33 #         inv_fact[i] = inv_fact[i + 1] * (i + 1) % mod
34 #     return fact, inv_fact
35
36 def modcomb(n, r, fact, inv_fact, mod):
37     if r > n or r < 0: return 0
38     return fact[n] * inv_fact[r] % mod * inv_fact[n-r] % mod

```

## 4.3 Number Theory (Advanced)

**Description:** Utilities for modular arithmetic, prime generation, Diophantine equations and CRT commonly needed in ICPC. Each helper includes when to use it, complexity and important edge-cases to remember.

```

1 # --- Assumes standard imports already present (e.g. 'import math') ---
2
3 # Sieve of Eratosthenes - O(n log log n)
4 # Returns (is_prime, primes).
5 # Use when you need primality answers for all values up to n or a list of
6 # primes for factorization.
7 # Memory: O(n).
8 def sieve(n):
9     is_prime = [True] * (n + 1)
10    if n >= 0:
11        is_prime[0] = False
12    if n >= 1:
13        is_prime[1] = False
14    for i in range(2, int(n**0.5) + 1):
15        if is_prime[i]:

```

```

15     step = i
16     start = i * i
17     is_prime[start:n+1:step] = [False] * (((n - start) // step) +
18         1)
19   primes = [i for i, v in enumerate(is_prime) if v]
20   return is_prime, primes
21
# Extended GCD: returns (g, x, y) such that a*x + b*y = g = gcd(a,b)
22 # Complexity: O(log(min(a,b))).
23 # Use to compute modular inverses when mod is not prime, or to solve a*x +
24 # b*y = c.
25 def extgcd(a, b):
26   if b == 0:
27     return a, 1, 0
28   g, x1, y1 = extgcd(b, a % b)
29   return g, y1, x1 - (a // b) * y1
30
# Modular inverse using extended gcd (works when mod is not prime)
31 # Returns integer inverse in [0, mod) or None if not invertible.
32 def modinv(a, mod):
33   g, x, _ = extgcd(a % mod, mod)
34   if g != 1:
35     return None # inverse doesn't exist (a and mod not coprime)
36   return x % mod
37
# CRT for two congruences: x = a1 (mod n1), x = a2 (mod n2)
38 # Handles non-coprime moduli by checking consistency.
39 # Returns (x, lcm) where x is the smallest non-negative solution modulo
40 # lcm(n1,n2), or (None, None) if none.
41 def crt_pair(a1, n1, a2, n2):
42   g, m1, m2 = extgcd(n1, n2)
43   if (a2 - a1) % g != 0:
44     return None, None
45   l = n1 // g * n2
46   t = ((a2 - a1) // g) * m1 % (n2 // g)
47   x = (a1 + n1 * t) % l
48   return x, l
49
# Fast integer floor-sum (AtCoder library pattern)
50 # sum_{i=0}^{n-1} floor((a*i + b) / m)
51 # Use for counting lattice points or evaluating floor-sum formulas
52 # efficiently.
53 def floor_sum(n, m, a, b):
54   res = 0
55   while True:
56     if a >= m:
57       res += (n - 1) * n * (a // m) // 2
58       a %= m
59     if b >= m:
60       res += n * (b // m)
61       b %= m
62     y = a * n + b
63     if y < m:
64       break
65     n, m, a, b = y // m, a, m, y % m
66   return res
67
# Modular exponentiation and inverse shortcuts
68 # pow(a, e, mod) is fast and should be preferred for a**e % mod.
69 # When mod is prime, use pow(a, -1, mod) for inverse (Python 3.8+),
70 # otherwise use modinv.
71
# Simple small-prime filter for primality checks. Add Miller-Rabin for
72 # large values.
73 def is_prime_simple(n):
74   if n < 2:
75     return False
76   small_primes = [2,3,5,7,11,13,17,19,23,29]
77   for p in small_primes:
78     if n == p:
79       return True
80     if n % p == 0:
81       return False
82   # For contest use: if n is large, apply deterministic Miller-Rabin
83   # bases for 64-bit.
84   return True
85
# Usage notes:
86 # - Prefer pow(a, e, mod) and pow(a, -1, mod) (when applicable) over
87 # custom loops.
88 # - Use sieve when many primality tests or prime lists are needed; use
89 # trial division or Miller-Rabin when testing a few large numbers.

```

```

7   def is_close(a, b, eps=EPS):
8     """Return True when |a-b| <= eps. Good when values are around 1.
9     Use relative comparison for widely varying magnitudes.
10    """
11    return abs(a - b) <= eps
12
13  # Relative tolerance comparison:
14  def is_close_rel(a, b, rel_eps=1e-12, abs_eps=EPS):
15    """Return True when |a-b| <= max(rel_eps * max(|a|,|b|), abs_eps).
16    Use this to avoid false negatives when values are very large or very
17    small.
18    """
19    return abs(a - b) <= max(rel_eps * max(abs(a), abs(b)), abs_eps)
20
21  # Summation: use math.fsum for numerically stable summation of many floats
22  :
23  # total = math.fsum(list_of_floats)
24
25  # Avoid repeated sqrt for comparisons: compare squared distances instead.
26  # Example: to check if d1 < d2 compare d1_sq < d2_sq.
27
28  # Fraction: exact rational arithmetic. Use when geometric intersections
29  # must be compared exactly.
30  # Example: p = Fraction(3,7)
31  # Caution: Fraction is slower and uses more memory; prefer integers or
32  # controlled eps comparisons in contests.
33
34  # safe_div: return infinities instead of ZeroDivisionError for convenience
35  # in some algorithms.
36  def safe_div(a, b):
37    if b == 0:
38      return float('inf') if a >= 0 else float('-inf')
39    return a / b
40
41  # Quick tips summary:
42  # - Keep computations integer when possible (use area2, orientations,
43  # squared distances).
44  # - Use EPS and relative comparisons consistently in a problem.
45  # - Use Fraction only when correctness requires exact rationals.
46  # - Use math.fsum when summing many floating values to reduce error.

```

## 5 Number Theory

**Description:** Essential algorithms for problems involving primes, modular arithmetic, and divisibility.

### 5.1 Modular Arithmetic

```

1  # Modular inverse using Fermat's Little Theorem
2  # Only works when mod is prime
3  # a^{(-1)} = a^{(mod-2)} (mod p)
4  def modinv(a, mod):
5    return pow(a, mod - 2, mod)
6
7  # Extended Euclidean Algorithm
8  # Returns (gcd, x, y) where ax + by = gcd(a,b)
9  # Can find modular inverse for any coprime a,mod
10 def extgcd(a, b):
11   if b == 0:
12     return a, 1, 0
13   g, x1, y1 = extgcd(b, a % b)
14   x = y1
15   y = x1 - (a // b) * y1
16   return g, x, y

```

### 5.2 Sieve of Eratosthenes

**Description:** Find all primes up to n in  $O(n \log \log n)$  time. Memory:  $O(n)$ .

```

1  def sieve(n):
2    is_prime = [True] * (n + 1)
3    is_prime[0] = is_prime[1] = False
4
5    for i in range(2, int(n**0.5) + 1):
6      if is_prime[i]:
7        # Mark multiples as composite
8        for j in range(i*i, n + 1, i):
9          is_prime[j] = False
10
11    return is_prime
12
13  # Get list of primes
14  primes = [i for i in range(n+1) if is_prime[i]]

```

### 5.3 Prime Factorization

**Description:** Decompose n into prime factors in  $O(\sqrt{n})$  time.

```

1  def factorize(n):
2    factors = []
3    d = 2
4
5    # Check divisors up to sqrt(n)
6    while d * d <= n:

```

## 4.4 Floating point and precision tips

**Description:** Practical guidance to avoid precision bugs in math and geometry. Prefer integer arithmetic when possible; use rational arithmetic (Fraction) only when exactness is required.

```

1  # --- Assumes 'import math' and 'from fractions import Fraction' exist at
2  # the top-level of the cheatsheet.
3
4  # EPS: absolute epsilon useful for values around magnitude 1.
5  EPS = 1e-9
6
# Absolute tolerance comparison:

```

```

7     while n % d == 0:
8         factors.append(d)
9         n //= d
10    d += 1
11
12    # If n > 1, it's a prime factor
13    if n > 1:
14        factors.append(n)
15
16    return factors
17
18 # Get prime factors with counts
19 from collections import Counter
20 def prime_factor_counts(n):
21     return Counter(factorize(n))
22
23 # Count divisors
24 def count_divisors(n):
25     count = 0
26     i = 1
27     while i * i <= n:
28         if n % i == 0:
29             count += 1 if i * i == n else 2
30         i += 1
31     return count
32
33 # Sum of divisors
34 def sum_divisors(n):
35     total = 0
36     i = 1
37     while i * i <= n:
38         if n % i == 0:
39             total += i
40             if i != n // i:
41                 total += n // i
42         i += 1
43     return total

```

## 5.4 Chinese Remainder Theorem

**Description:** Solve system of congruences  $x \equiv a_1 \pmod{m_1}$ ,  $x \equiv a_2 \pmod{m_2}$ , ... Time:  $O(n \log M)$  where  $M$  is product of moduli.

```

1 def chinese_remainder(remainders, moduli):
2     # Solve  $x = remainders[i] \pmod{moduli[i]}$ 
3     # Assumes moduli are pairwise coprime
4
5     def extgcd(a, b):
6         if b == 0:
7             return a, 1, 0
8         g, x1, y1 = extgcd(b, a % b)
9         return g, y1, x1 - (a // b) * y1
10
11    total = 0
12    prod = 1
13    for m in moduli:
14        prod *= m
15
16    for r, m in zip(remainders, moduli):
17        p = prod // m
18        g, inv, _ = extgcd(p, m)
19        # inv may be negative, normalize it
20        inv = (inv % m + m) % m
21        total += r * inv * p
22
23    return total % prod

```

## 5.5 Euler's Totient Function

**Description:**  $\phi(n)$  = count of numbers  $\leq n$  coprime to  $n$ . Time:  $O(\sqrt{n})$ .

```

1 def euler_phi(n):
2     result = n
3     p = 2
4
5     while p * p <= n:
6         if n % p == 0:
7             # Remove factor p
8             while n % p == 0:
9                 n /= p
10            # Multiply by  $(1 - 1/p)$ 
11            result -= result // p
12        p += 1
13
14    if n > 1:
15        result -= result // n
16
17    return result
18
19 # Phi for range [1, n] using sieve
20 def phi_sieve(n):
21     phi = list(range(n + 1)) # phi[i] = i initially
22
23     for i in range(2, n + 1):
24         if phi[i] == i: # i is prime

```

```

25         for j in range(i, n + 1, i):
26             phi[j] = phi[j] // i * (i - 1)
27
28     return phi

```

## 5.6 Fast Exponentiation with Matrix

**Description:** Already covered in matrix section, but useful pattern.

```

1 # Modular exponentiation
2 def mod_exp(base, exp, mod):
3     result = 1
4     base %= mod
5
6     while exp > 0:
7         if exp & 1:
8             result = (result * base) % mod
9             base = (base * base) % mod
10            exp >>= 1
11
12     return result

```

# 6 Graph Algorithms

## 6.1 Constructing Graphs

**Description:** Common patterns for building graphs from problem input. Always check whether the input is 0-indexed or 1-indexed (subtract 1 for 1-indexed inputs when using 0-based arrays). Choose representation based on density and operations you need (adjacency list for sparse graphs, adjacency matrix for dense/small graphs).

```

1 # Unweighted undirected (defaultdict) - common and flexible
2 from collections import defaultdict
3 graph = defaultdict(list)
4 for _ in range(m):
5     u, v = map(int, input().split())
6     u -= 1 # remove if input already 0-indexed
7     v -= 1
8     graph[u].append(v)
9     graph[v].append(u)
10
11 # Directed: append only one direction
12 # graph[u].append(v)
13
14 # Weighted edges: store (neighbor, weight)
15 graph = defaultdict(list)
16 for _ in range(m):
17     u, v, w = map(int, input().split())
18     u -= 1; v -= 1
19     graph[u].append((v, w))
20     # graph[v].append((u, w)) # add for undirected
21
22 # If n is known, list-of-lists is slightly faster
23 graph = [[] for _ in range(n)]
24 graph[u].append(v)
25
26 # Avoid duplicate edges when needed
27 graph_sets = [set() for _ in range(n)]
28 graph_sets[u].add(v)
29
30 # Adjacency matrix for dense graphs / small n
31 adj = [[0] * n for _ in range(n)]
32 adj[u][v] = 1 # or assign weight
33
34 # Fast parsing for large input (useful in contests)
35 import sys
36 data = list(map(int, sys.stdin.buffer.read().split()))
37 it = iter(data)
38 n, m = next(it), next(it)
39 for _ in range(m):
40     u, v = next(it), next(it)
41     u -= 1; v -= 1
42     graph[u].append(v)
43     graph[v].append(u)

```

## Constructing a graph from a grid

**Description:** Many problems give a grid (maze, weights, terrain). Convert grid cells to nodes ( $id = x * m + y$ ) or work with implicit neighbours on the fly.

```

1 # Example grid (list of strings) where '.' is open and '#' is wall
2 grid = [
3     ".#.",
4     "#.#.",
5     ".#..",
6     "....",
7 ]
8 n = len(grid)
9 m = len(grid[0])
10
11 # Map (x,y) to node id and back
12 def id_of(x, y):
13     return x * m + y

```

```

13 def coord_of(id):
14     return divmod(id, m) # (x, y)
15
16 # 4-direction neighbours
17 dirs4 = [(0,1), (1,0), (0,-1), (-1,0)]
18
19 # Build adjacency list for traversable cells
20 graph = [[] for _ in range(n * m)]
21 for x in range(n):
22     for y in range(m):
23         if grid[x][y] == '#':
24             continue
25         u = id_of(x, y)
26         for dx, dy in dirs4:
27             nx, ny = x + dx, y + dy
28             if 0 <= nx < n and 0 <= ny < m and grid[nx][ny] != '#':
29                 v = id_of(nx, ny)
30                 graph[u].append(v)
31
32
33 # Use BFS on the implicit/constructed graph
34 from collections import deque
35 start = id_of(0, 0)
36 if grid[0][0] == '#':
37     start = None
38
39 if start is not None:
40     dist = [-1] * (n * m)
41     q = deque([start])
42     dist[start] = 0
43     while q:
44         u = q.popleft()
45         for v in graph[u]:
46             if dist[v] == -1:
47                 dist[v] = dist[u] + 1
48                 q.append(v)
49
50 # Diagonals (8-direction)
51 dirs8 = dirs4 + [(-1,-1), (-1,1), (1,-1), (1,1)]
52
53 # Weighted grid: movement cost stored separately
54 costs = [[1] * m for _ in range(n)] # example uniform cost
55 wg = [[0] for _ in range(n*m)]
56 for x in range(n):
57     for y in range(m):
58         if grid[x][y] == '#':
59             continue
60         u = id_of(x,y)
61         for dx, dy in dirs4:
62             nx, ny = x + dx, y + dy
63             if 0 <= nx < n and 0 <= ny < m and grid[nx][ny] != '#':
64                 v = id_of(nx, ny)
65                 wg[u].append((v, costs[nx][ny]))
66
67 # Implicit neighbour generator (avoid building edges when memory matters)
68 def neighbours(x, y):
69     for dx, dy in dirs4:
70         nx, ny = x + dx, y + dy
71         if 0 <= nx < n and 0 <= ny < m and grid[nx][ny] != '#':
72             yield nx, ny
73
74 # BFS using implicit neighbours
75 dist2 = [[-1] * m for _ in range(n)]
76 from collections import deque
77 q = deque()
78 if grid[0][0] != '#':
79     dist2[0][0] = 0
80     q.append((0,0))
81
82 while q:
83     x, y = q.popleft()
84     for nx, ny in neighbours(x, y):
85         if dist2[nx][ny] == -1:
86             dist2[nx][ny] = dist2[x][y] + 1
87             q.append((nx, ny))

```

## 6.2 Graph Representation

**Description:** Adjacency list is most common for sparse graphs. Use defaultdict for convenience.

```

1 from collections import defaultdict, deque
2
3 # Unweighted graph
4 graph = defaultdict(list)
5 for _ in range(m):
6     u, v = map(int, input().split())
7     graph[u].append(v)
8     graph[v].append(u) # for undirected
9
10 # Weighted graph - store (neighbor, weight) tuples
11 graph[u].append((v, weight))

```

## 6.3 BFS (Breadth-First Search)

**Description:** Explores graph level by level. Finds shortest path in unweighted graphs. Time: O(V+E), Space: O(V).

```

1 def bfs(graph, start):
2     visited = set([start])
3     queue = deque([start])
4     dist = {start: 0}
5
6     while queue:
7         node = queue.popleft()
8
9         for neighbor in graph[node]:
10            if neighbor not in visited:
11                visited.add(neighbor)
12                queue.append(neighbor)
13                dist[neighbor] = dist[node] + 1
14
15    return dist
16
17 # Grid BFS - common in maze/path problems
18 def grid_bfs(grid, start):
19     n, m = len(grid), len(grid[0])
20     visited = [[False] * m for _ in range(n)]
21     queue = deque([start])
22     visited[start[0]][start[1]] = True
23
24     # 4 directions: right, down, left, up
25     dirs = [(0,1), (1,0), (0,-1), (-1,0)]
26
27     while queue:
28         x, y = queue.popleft()
29
30         for dx, dy in dirs:
31             nx, ny = x + dx, y + dy
32
33             # Check bounds and validity
34             if (0 <= nx < n and 0 <= ny < m
35                 and not visited[nx][ny]
36                 and grid[nx][ny] != '#'):
37
38                 visited[nx][ny] = True
39                 queue.append((nx, ny))

```

## 6.4 DFS (Depth-First Search)

**Description:** Explores as far as possible along each branch. Used for connectivity, cycles, topological sort. Time: O(V+E), Space: O(V).

```

1 # Recursive DFS
2 def dfs(graph, node, visited):
3     visited.add(node)
4
5     for neighbor in graph[node]:
6         if neighbor not in visited:
7             dfs(graph, neighbor, visited)
8
9 # Iterative DFS using stack
10 def dfs_iterative(graph, start):
11     visited = set()
12     stack = [start]
13
14     while stack:
15         node = stack.pop()
16
17         if node not in visited:
18             visited.add(node)
19
20             for neighbor in graph[node]:
21                 if neighbor not in visited:
22                     stack.append(neighbor)
23
24 # Cycle detection in undirected graph
25 def has_cycle(graph, n):
26     visited = [False] * n
27
28     def dfs(node, parent):
29         visited[node] = True
30
31         for neighbor in graph[node]:
32             if not visited[neighbor]:
33                 if dfs(neighbor, node):
34                     return True
35
36             # Back edge to non-parent = cycle
37             elif neighbor != parent:
38                 return True
39
40         return False
41
42 # Check all components
43 for i in range(n):
44     if not visited[i]:
45         if dfs(i, -1):
46             return True

```

```

47     return False
48
49 # Cycle detection in directed graph
50 def has_cycle_directed(graph, n):
51     WHITE, GRAY, BLACK = 0, 1, 2
52     color = [WHITE] * n
53
54     def dfs(node):
55         color[node] = GRAY
56
57         for neighbor in graph[node]:
58             if color[neighbor] == GRAY:
59                 return True # Back edge = cycle
60             if color[neighbor] == WHITE:
61                 if dfs(neighbor):
62                     return True
63
64             color[node] = BLACK
65
66     return False
67
68 for i in range(n):
69     if color[i] == WHITE:
70         if dfs(i):
71             return True
72
73 return False
74
75 # Connected components count
76 def count_components(graph, n):
77     visited = [False] * n
78     count = 0
79
80     def dfs(node):
81         visited[node] = True
82         for neighbor in graph[node]:
83             if not visited[neighbor]:
84                 dfs(neighbor)
85
86     for i in range(n):
87         if not visited[i]:
88             dfs(i)
89             count += 1
90
91     return count
92
93 # Bipartite check (2-coloring)
94 def is_bipartite(graph, n):
95     color = [-1] * n
96
97     def bfs(start):
98         from collections import deque
99         queue = deque([start])
100        color[start] = 0
101
102        while queue:
103            node = queue.popleft()
104
105            for neighbor in graph[node]:
106                if color[neighbor] == -1:
107                    color[neighbor] = 1 - color[node]
108                    queue.append(neighbor)
109                elif color[neighbor] == color[node]:
110                    return False
111
112    return True
113
114    for i in range(n):
115        if color[i] == -1:
116            if not bfs(i):
117                return False
118
119    return True

```

## 6.5 Strongly Connected Components (SCC)

Description: Find all SCCs in directed graph using Tarjan's algorithm.  
Time: O(V+E).

```

1 def tarjan_scc(graph, n):
2     index_counter = [0]
3     stack = []
4     lowlink = [0] * n
5     index = [0] * n
6     on_stack = [False] * n
7     index_initialized = [False] * n
8     sccs = []
9
10    def strongconnect(v):
11        index[v] = index_counter[0]
12        lowlink[v] = index_counter[0]
13        index_counter[0] += 1
14        index_initialized[v] = True
15        stack.append(v)
16        on_stack[v] = True
17
18        for w in graph[v]:
19            if not index_initialized[w]:

```

```

20            strongconnect(w)
21            lowlink[v] = min(lowlink[v], lowlink[w])
22
23    elif on_stack[w]:
24        lowlink[v] = min(lowlink[v], index[w])
25
26    if lowlink[v] == index[v]:
27        scc = []
28        while True:
29            w = stack.pop()
30            on_stack[w] = False
31            scc.append(w)
32            if w == v:
33                break
34        sccs.append(scc)
35
36    for v in range(n):
37        if not index_initialized[v]:
38            strongconnect(v)
39
40    return sccs

```

## 6.6 Bridges and Articulation Points

Description: Find critical edges (bridges) and vertices (articulation points).  
Time: O(V+E).

```

1 def find_bridges(graph, n):
2     visited = [False] * n
3     disc = [0] * n
4     low = [0] * n
5     parent = [-1] * n
6     time = [0]
7     bridges = []
8
9     def dfs(u):
10        visited[u] = True
11        disc[u] = low[u] = time[0]
12        time[0] += 1
13
14        for v in graph[u]:
15            if not visited[v]:
16                parent[v] = u
17                dfs(v)
18                low[u] = min(low[u], low[v])
19
20            # Bridge condition
21            if low[v] > disc[u]:
22                bridges.append((u, v))
23
24            elif v != parent[u]:
25                low[u] = min(low[u], disc[v])
26
27        for i in range(n):
28            if not visited[i]:
29                dfs(i)
30
31    return bridges
32
33 def find_articulation_points(graph, n):
34     visited = [False] * n
35     disc = [0] * n
36     low = [0] * n
37     parent = [-1] * n
38     time = [0]
39     ap = set()
40
41     def dfs(u):
42         children = 0
43         visited[u] = True
44         disc[u] = low[u] = time[0]
45         time[0] += 1
46
47         for v in graph[u]:
48             if not visited[v]:
49                 children += 1
50                 parent[v] = u
51                 dfs(v)
52                 low[u] = min(low[u], low[v])
53
54             # Articulation point conditions
55             if parent[u] == -1 and children > 1:
56                 ap.add(u)
57             if parent[u] != -1 and low[v] >= disc[u]:
58                 ap.add(u)
59
60             elif v != parent[u]:
61                 low[u] = min(low[u], disc[v])
62
63        for i in range(n):
64            if not visited[i]:
65                dfs(i)
66
67    return list(ap)

```

## 6.7 Lowest Common Ancestor (LCA)

**Description:** Find LCA of two nodes in a tree. Binary lifting preprocessing:  $O(n \log n)$ , Query:  $O(\log n)$ .

```

1 class LCA:
2     def __init__(self, graph, root, n):
3         self.n = n
4         self.LOG = 20 # log2(n) + 1
5         self.parent = [[-1] * self.LOG for _ in range(n)]
6         self.depth = [0] * n
7
8         # DFS to set parent and depth
9         visited = [False] * n
10
11        def dfs(node, par, d):
12            visited[node] = True
13            self.parent[node][0] = par
14            self.depth[node] = d
15
16            for neighbor in graph[node]:
17                if not visited[neighbor]:
18                    dfs(neighbor, node, d + 1)
19
20        dfs(root, -1, 0)
21
22        # Binary lifting preprocessing
23        for j in range(1, self.LOG):
24            for i in range(n):
25                if self.parent[i][j-1] != -1:
26                    self.parent[i][j] = self.parent[
27                        self.parent[i][j-1]][j-1]
28
29    def lca(self, u, v):
30        # Make u deeper
31        if self.depth[u] < self.depth[v]:
32            u, v = v, u
33
34        # Bring u to same level as v
35        diff = self.depth[u] - self.depth[v]
36        for i in range(self.LOG):
37            if (diff >> i) & 1:
38                u = self.parent[u][i]
39
40        if u == v:
41            return u
42
43        # Binary search for LCA
44        for i in range(self.LOG - 1, -1, -1):
45            if self.parent[u][i] != self.parent[v][i]:
46                u = self.parent[u][i]
47                v = self.parent[v][i]
48
49        return self.parent[u][0]
50
51    def dist(self, u, v):
52        # Distance between two nodes
53        l = self.lca(u, v)
54        return self.depth[u] + self.depth[v] - 2 * self.depth[l]
```

```

29 def dijkstra_with_path(graph, start, n):
30     dist = [float('inf')] * n
31     parent = [-1] * n
32     dist[start] = 0
33     heap = [(0, start)]
34
35     while heap:
36         d, node = heapq.heappop(heap)
37         if d > dist[node]:
38             continue
39
40         for neighbor, weight in graph[node]:
41             new_dist = dist[node] + weight
42             if new_dist < dist[neighbor]:
43                 dist[neighbor] = new_dist
44                 parent[neighbor] = node
45                 heapq.heappush(heap, (new_dist, neighbor))
46
47     return dist, parent
48
49 def reconstruct_path(parent, target):
50     path = []
51     while target != -1:
52         path.append(target)
53         target = parent[target]
54     return path[::-1]
```

## 7.2 Bellman-Ford Algorithm

**Description:** Finds shortest paths with negative edges. Detects negative cycles. Time:  $O(VE)$ .

```

1 def bellman_ford(edges, n, start):
2     # edges = [(u, v, weight), ...]
3     dist = [float('inf')] * n
4     dist[start] = 0
5
6     # Relax edges n-1 times
7     for _ in range(n - 1):
8         for u, v, w in edges:
9             if dist[u] != float('inf') and \
10                 dist[u] + w < dist[v]:
11                     dist[v] = dist[u] + w
12
13     # Check for negative cycles
14     for u, v, w in edges:
15         if dist[u] != float('inf') and \
16             dist[u] + w < dist[v]:
17                 return None # Negative cycle exists
18
19     return dist
```

## 7.3 Floyd-Warshall Algorithm

**Description:** All-pairs shortest paths. Works with negative edges (no negative cycles). Time:  $O(V^3)$ .

```

1 def floyd_marshall(n, edges):
2     # Initialize distance matrix
3     dist = [[float('inf')]] * n for _ in range(n)
4
5     for i in range(n):
6         dist[i][i] = 0
7
8     for u, v, w in edges:
9         dist[u][v] = min(dist[u][v], w)
10
11    # Dynamic programming
12    for k in range(n): # Intermediate vertex
13        for i in range(n):
14            for j in range(n):
15                dist[i][j] = min(dist[i][j],
16                                 dist[i][k] + dist[k][j])
17
18    return dist
19
20    # Check for negative cycle
21    def has_negative_cycle(dist, n):
22        for i in range(n):
23            if dist[i][i] < 0:
24                return True
25        return False
```

## 7.4 Minimum Spanning Tree

### 7.4.1 Kruskal's Algorithm

**Description:** MST using Union-Find. Sort edges by weight. Time:  $O(E \log E)$ .

```

1 def kruskal(n, edges):
2     # edges = [(weight, u, v), ...]
3     edges.sort() # Sort by weight
4
```

```

5   uf = UnionFind(n)
6   mst_weight = 0
7   mst_edges = []
8
9   for weight, u, v in edges:
10    if uf.union(u, v):
11      mst_weight += weight
12      mst_edges.append((u, v, weight))
13
14  return mst_weight, mst_edges
15
16 class UnionFind:
17     def __init__(self, n):
18         self.parent = list(range(n))
19         self.rank = [0] * n
20
21     def find(self, x):
22         if self.parent[x] != x:
23             self.parent[x] = self.find(self.parent[x])
24         return self.parent[x]
25
26     def union(self, x, y):
27         px, py = self.find(x), self.find(y)
28         if px == py:
29             return False
30         if self.rank[px] < self.rank[py]:
31             px, py = py, px
32             self.parent[py] = px
33         if self.rank[px] == self.rank[py]:
34             self.rank[px] += 1
35         return True

```

```

25         queue.append(neighbor)
26
27     # If not all nodes processed, cycle exists
28     return result if len(result) == n else []

```

## 8.2 DFS-based Topological Sort

**Advantages:** Simpler code, uses less space.

```

1  def topo_dfs(graph, n):
2      visited = [False] * n
3      stack = []
4
5      def dfs(node):
6          visited[node] = True
7
8          # Visit all neighbors first
9          for neighbor in graph[node]:
10             if not visited[neighbor]:
11                 dfs(neighbor)
12
13         # Add to stack after visiting all descendants
14         stack.append(node)
15
16     # Process all components
17     for i in range(n):
18         if not visited[i]:
19             dfs(i)
20
21     # Reverse stack gives topological order
22     return stack[::-1]

```

### 7.4.2 Prim's Algorithm

**Description:** MST using heap. Good for dense graphs. Time:  $O(E \log V)$ .

```

1 import heapq
2
3 def prim(graph, n):
4     # graph[u] = [(v, weight), ...]
5     visited = [False] * n
6     min_heap = [(0, 0)] # (weight, node)
7     mst_weight = 0
8
9     while min_heap:
10        weight, u = heapq.heappop(min_heap)
11
12        if visited[u]:
13            continue
14
15        visited[u] = True
16        mst_weight += weight
17
18        for v, w in graph[u]:
19            if not visited[v]:
20                heapq.heappush(min_heap, (w, v))
21
22    return mst_weight

```

## 9 Union-Find (Disjoint Set Union)

**Description:** Efficiently tracks disjoint sets and supports union and find operations. Used for Kruskal's MST, connected components, cycle detection. Time:  $O(\alpha(n)) \approx O(1)$  per operation with path compression and union by rank.

### Applications:

- Kruskal's minimum spanning tree
- Detecting cycles in undirected graphs
- Finding connected components
- Network connectivity problems

```

1  class UnionFind:
2      def __init__(self, n):
3          # Each node is its own parent initially
4          self.parent = list(range(n))
5          # Rank for union by rank optimization
6          self.rank = [0] * n
7
8      def find(self, x):
9          # Path compression: point directly to root
10         if self.parent[x] != x:
11             self.parent[x] = self.find(self.parent[x])
12         return self.parent[x]
13
14      def union(self, x, y):
15          # Find roots
16          px, py = self.find(x), self.find(y)
17
18          # Already in same set
19          if px == py:
20              return False
21
22          # Union by rank: attach smaller tree under larger
23          if self.rank[px] < self.rank[py]:
24              px, py = py, px
25
26          self.parent[py] = px
27
28          # Increase rank if trees had equal rank
29          if self.rank[px] == self.rank[py]:
30              self.rank[px] += 1
31
32          return True
33
34      def connected(self, x, y):
35          return self.find(x) == self.find(y)
36
37      # Count number of disjoint sets
38      def count_sets(self):
39          return len(set(self.find(i)))
40          for i in range(len(self.parent)))
41
42      # Example: Detect cycle in undirected graph
43      def has_cycle_uf(edges, n):
44          uf = UnionFind(n)
45          for u, v in edges:
46              if uf.connected(u, v):
47                  return True # Cycle found
48                  uf.union(u, v)
49

```

## 8 Topological Sort

**Description:** Linear ordering of vertices in a DAG (Directed Acyclic Graph) such that for every edge  $u \rightarrow v$ ,  $u$  comes before  $v$ . Used for task scheduling, course prerequisites, build systems. Time:  $O(V+E)$ .

### 8.1 Kahn's Algorithm (BFS-based)

**Advantages:** Detects cycles, can process nodes level by level.

```

1 from collections import deque
2
3 def topo_sort(graph, n):
4     # Count incoming edges for each node
5     indegree = [0] * n
6     for u in range(n):
7         for v in graph[u]:
8             indegree[v] += 1
9
10    # Start with nodes having no dependencies
11    queue = deque([i for i in range(n)
12                  if indegree[i] == 0])
13    result = []
14
15    while queue:
16        node = queue.popleft()
17        result.append(node)
18
19        # Remove this node from graph
20        for neighbor in graph[node]:
21            indegree[neighbor] -= 1
22
23        # If neighbor has no more dependencies
24        if indegree[neighbor] == 0:

```

## 10 Binary Search

**Description:** Search in  $O(\log n)$  time. Works on monotonic functions where condition changes from false to true (or vice versa) exactly once.

### 10.1 Basic Templates

```

1 # Standard binary search (find exact element)
2 def binary_search(arr, target):
3     left, right = 0, len(arr) - 1
4
5     while left <= right:
6         mid = (left + right) // 2
7
8         if arr[mid] == target:
9             return mid
10        elif arr[mid] < target:
11            left = mid + 1
12        else:
13            right = mid - 1
14
15    return -1 # Not found
16
17 # Find FIRST position where condition is True
18 def lower_bound(arr, target):
19     left, right = 0, len(arr)
20
21     while left < right:
22         mid = (left + right) // 2
23
24         if arr[mid] < target:
25             left = mid + 1
26         else:
27             right = mid
28
29     return left
30
31 # Find LAST position where condition is True + 1
32 def upper_bound(arr, target):
33     left, right = 0, len(arr)
34
35     while left < right:
36         mid = (left + right) // 2
37
38         if arr[mid] <= target:
39             left = mid + 1
40         else:
41             right = mid
42
43     return left
44
45 # 2. Find range of target in sorted array
46 def search_range(nums, target):
47     left = lower_bound(nums, target)
48     if left >= len(nums) or nums[left] != target:
49         return [-1, -1]
50     right = upper_bound(nums, target) - 1
51     return [left, right]
52
53 # 3. Peak element (not strictly increasing/decreasing)
54 def find_peak_element(nums):
55     left, right = 0, len(nums) - 1
56
57     while left < right:
58         mid = (left + right) // 2
59
60         if nums[mid] > nums[mid + 1]:
61             right = mid # Peak is on left or mid
62         else:
63             left = mid + 1 # Peak is on right
64
65     return left
66
67 # 4. Rotated sorted array
68 def search_rotated(nums, target):
69     left, right = 0, len(nums) - 1
70
71     while left <= right:
72         mid = (left + right) // 2
73
74         if nums[mid] == target:
75             return mid
76
77         # Left half is sorted
78         if nums[left] <= nums[mid]:
79             if nums[left] <= target < nums[mid]:
80                 right = mid - 1
81             else:
82                 left = mid + 1
83
84         # Right half is sorted
85         else:
86             if nums[mid] < target <= nums[right]:
87                 left = mid + 1
88             else:
89                 right = mid - 1
90
91     return -1
92
93 # 5. Square root (integer)
94 def sqrt_int(x):
95     if x < 2:
96         return x
97
98     left, right = 1, x // 2
99
100    while left <= right:
101        mid = (left + right) // 2
102        square = mid * mid
103
104        if square == x:
105            return mid
106        elif square < x:
107            left = mid + 1
108        else:
109            right = mid - 1
110
111    return right # Largest integer whose square <= x

```

### 10.2 Advanced Templates

```

1 # Find first True in boolean array
2 def find_first_true(check_func, left, right):
3     while left < right:
4         mid = (left + right) // 2
5
6         if check_func(mid):
7             right = mid # Could be answer, search left
8         else:
9             left = mid + 1 # Not answer, search right
10
11    return left
12
13 # Find last True in boolean array
14 def find_last_true(check_func, left, right):
15     while left < right:
16         mid = (left + right + 1) // 2 # Round up!
17
18         if check_func(mid):
19             left = mid # Could be answer, search right
20         else:
21             right = mid - 1 # Not answer, search left
22
23    return left
24
25 # Binary search on floating point
26 def binary_search_float(check_func, left, right, eps=1e-9):
27     while right - left > eps:
28         mid = (left + right) / 2
29
30         if check_func(mid):
31             right = mid
32         else:
33             left = mid
34
35    return left

```

### 10.3 Common Use Cases

```

1 # 1. Find insertion point
2 def search_insert_position(nums, target):
3     return lower_bound(nums, target)

```

```

1 # Template for optimization problems
2 def minimize_max_distance(stations, k):
3     # Can we place k stations with max distance <= max_dist?
4     def possible(max_dist):
5         stations_needed = 0
6         for i in range(len(stations) - 1):
7             gap = stations[i + 1] - stations[i]
8             stations_needed += int(gap / max_dist)
9         return stations_needed <= k
10
11     left, right = 0.0, stations[-1] - stations[0]
12
13     while right - left > 1e-6:
14         mid = (left + right) / 2
15
16         if possible(mid):
17             right = mid
18         else:
19             left = mid
20
21     return left
22
23 # Capacity allocation problem
24 def split_array_largest_sum(nums, m):
25     # Can we split into m subarrays with max sum <= max_sum?

```

```

26     def can_split(max_sum):
27         count = 1
28         current_sum = 0
29
30         for num in nums:
31             if current_sum + num > max_sum:
32                 count += 1
33                 current_sum = num
34                 if count > m:
35                     return False
36             else:
37                 current_sum += num
38
39         return True
40
41 left, right = max(nums), sum(nums)
42
43 while left < right:
44     mid = (left + right) // 2
45
46     if can_split(mid):
47         right = mid
48     else:
49         left = mid + 1
50
51 return left
52
53 # Resource allocation with greedy check
54 def minimum_speed_to_eat_bananas(piles, h):
55     # Can we eat all bananas in h hours at speed k?
56     def can_finish(k):
57         return sum((pile + k - 1) // k for pile in piles) <= h
58
59     left, right = 1, max(piles)
60
61     while left < right:
62         mid = (left + right) // 2
63
64         if can_finish(mid):
65             right = mid
66         else:
67             left = mid + 1
68
69 return left

```

**Key Points:**

- Always clarify what you're searching for (first/last occurrence, exact match, etc.)
- For "find first True": use `right = mid`
- For "find last True": use `left = mid` and `mid = (left + right + 1) // 2`
- Binary search on answer: define a monotonic check function
- Handle edge cases: empty arrays, single elements, all same elements

## 10.5 Common Pitfalls

**Watch out for:**

- **Infinite loops:** Wrong midpoint calculation for "find last" (missing +1)
- **Off-by-one errors:** Mixing `len(arr)` vs `len(arr)-1` in initial bounds
- **Integer overflow:** Use `left + (right - left) // 2` instead of `(left + right) // 2`
- **Wrong condition:** `<` vs `<=` in loop condition (`left < right` vs `left <= right`)
- **Non-monotonic function:** Binary search only works if condition changes at most once
- **Empty result handling:** Check bounds before accessing `arr[result]`
- **Floating point precision:** Use appropriate epsilon for convergence
- **Wrong search space:** Ensure your `left` and `right` bounds contain the answer

# 11 Dynamic Programming

**Description:** Solve problems by breaking them into overlapping subproblems. Store results to avoid recomputation.

## 11.1 Longest Increasing Subsequence

**Description:** Find length of longest strictly increasing subsequence. Time:  $O(n \log n)$  using binary search.

```

1 def lis(arr):
2     from bisect import bisect_left
3
4     # dp[i] = smallest ending value of LIS of length i+1
5     dp = []
6
7     for x in arr:
8         # Find position to place x
9         idx = bisect_left(dp, x)
10
11        if idx == len(dp):
12            dp.append(x) # Extend LIS
13        else:
14            dp[idx] = x # Better ending for this length

```

```

15     return len(dp)
16
17
18 # LIS with actual sequence
19 def lis_with_sequence(arr):
20     from bisect import bisect_left
21
22     n = len(arr)
23     dp = []
24     parent = [-1] * n
25     dp_idx = [] # indices in dp
26
27     for i, x in enumerate(arr):
28         idx = bisect_left(dp, x)
29
30         if idx == len(dp):
31             dp.append(x)
32             dp_idx.append(i)
33         else:
34             dp[idx] = x
35             dp_idx[idx] = i
36
37         if idx > 0:
38             parent[i] = dp_idx[idx - 1]
39
40     # Reconstruct sequence
41     result = []
42     idx = dp_idx[-1]
43     while idx != -1:
44         result.append(arr[idx])
45         idx = parent[idx]
46
47     return result[::-1]

```

## 11.2 0/1 Knapsack

**Description:** Maximum value with weight capacity. Each item can be taken 0 or 1 time. Time:  $O(n \times \text{capacity})$ , Space:  $O(n \times \text{capacity})$ .

```

1 def knapsack(weights, values, capacity):
2     n = len(weights)
3     # dp[i][w] = max value using first i items,
4     #           weight <= w
5     dp = [[0] * (capacity + 1) for _ in range(n + 1)]
6
7     for i in range(1, n + 1):
8         for w in range(capacity + 1):
9             # Don't take item i-1
10                dp[i][w] = dp[i-1][w]
11
12             # Take item i-1 if it fits
13             if weights[i-1] <= w:
14                 dp[i][w] = max(
15                     dp[i][w],
16                     dp[i-1][w - weights[i-1]] + values[i-1]
17                 )
18
19     return dp[n][capacity]
20
21 # Space-optimized O(capacity)
22 def knapsack_optimized(weights, values, capacity):
23     dp = [0] * (capacity + 1)
24
25     for i in range(len(weights)):
26         # Iterate backwards to avoid using updated values
27         for w in range(capacity, weights[i] - 1, -1):
28             dp[w] = max(dp[w],
29                         dp[w - weights[i]] + values[i])
30
31     return dp[capacity]

```

## 11.3 Edit Distance (Levenshtein Distance)

**Description:** Minimum operations (insert, delete, replace) to transform s1 to s2. Time:  $O(m \times n)$ , Space:  $O(m \times n)$ .

```

1 def edit_dist(s1, s2):
2     m, n = len(s1), len(s2)
3     # dp[i][j] = edit distance of s1[:i] and s2[:j]
4     dp = [[0] * (n + 1) for _ in range(m + 1)]
5
6     # Base cases: empty string transformations
7     for i in range(m + 1):
8         dp[i][0] = i # Delete all
9     for j in range(n + 1):
10        dp[0][j] = j # Insert all
11
12    for i in range(1, m + 1):
13        for j in range(1, n + 1):
14            if s1[i-1] == s2[j-1]:
15                # Characters match, no operation needed
16                dp[i][j] = dp[i-1][j-1]
17            else:
18                dp[i][j] = 1 + min(
19                    dp[i-1][j], # Delete from s1

```

```

20     dp[i][j-1],      # Insert into s1
21     dp[i-1][j-1]    # Replace in s1
22
23
24 return dp[m][n]

```

## 11.4 Longest Common Subsequence (LCS)

Description: Longest subsequence common to two sequences. Time: O( $m \times n$ ).

```

1 def lcs(s1, s2):
2     m, n = len(s1), len(s2)
3     dp = [[0] * (n + 1) for _ in range(m + 1)]
4
5     for i in range(1, m + 1):
6         for j in range(1, n + 1):
7             if s1[i-1] == s2[j-1]:
8                 dp[i][j] = dp[i-1][j-1] + 1
9             else:
10                dp[i][j] = max(dp[i-1][j], dp[i][j-1])
11
12 return dp[m][n]
13
14 # Reconstruct LCS
15 def lcs_string(s1, s2):
16     m, n = len(s1), len(s2)
17     dp = [[0] * (n + 1) for _ in range(m + 1)]
18
19     for i in range(1, m + 1):
20         for j in range(1, n + 1):
21             if s1[i-1] == s2[j-1]:
22                 dp[i][j] = dp[i-1][j-1] + 1
23             else:
24                 dp[i][j] = max(dp[i-1][j], dp[i][j-1])
25
26 # Backtrack
27 result = []
28 i, j = m, n
29 while i > 0 and j > 0:
30     if s1[i-1] == s2[j-1]:
31         result.append(s1[i-1])
32         i -= 1
33         j -= 1
34     elif dp[i-1][j] > dp[i][j-1]:
35         i -= 1
36     else:
37         j -= 1
38
39 return ''.join(reversed(result))

```

## 11.5 Coin Change

Description: Minimum coins to make amount, or count ways. Time: O( $n \times \text{amount}$ ).

```

1 # Minimum coins
2 def coin_change_min(coins, amount):
3     dp = [float('inf')] * (amount + 1)
4     dp[0] = 0
5
6     for coin in coins:
7         for i in range(coin, amount + 1):
8             dp[i] = min(dp[i], dp[i - coin] + 1)
9
10    return dp[amount] if dp[amount] != float('inf') else -1
11
12 # Count ways
13 def coin_change_ways(coins, amount):
14     dp = [0] * (amount + 1)
15     dp[0] = 1
16
17     for coin in coins:
18         for i in range(coin, amount + 1):
19             dp[i] += dp[i - coin]
20
21    return dp[amount]

```

## 11.6 Palindrome Partitioning

Description: Minimum cuts to partition string into palindromes. Time: O( $n^2$ ).

```

1 def min_palindrome_partition(s):
2     n = len(s)
3
4     # is_pal[i][j] = True if s[i:j+1] is palindrome
5     is_pal = [[False] * n for _ in range(n)]
6
7     # Every single character is palindrome
8     for i in range(n):
9         is_pal[i][i] = True
10
11    # Check all substrings
12    for length in range(2, n + 1):

```

```

13     for i in range(n - length + 1):
14         j = i + length - 1
15         if s[i] == s[j]:
16             is_pal[i][j] = (length == 2 or
17                             is_pal[i+1][j-1])
18
19     # dp[i] = min cuts for s[0:i+1]
20     dp = [float('inf')] * n
21
22     for i in range(n):
23         if is_pal[0][i]:
24             dp[i] = 0
25         else:
26             for j in range(i):
27                 if is_pal[j+1][i]:
28                     dp[i] = min(dp[i], dp[j] + 1)
29
30    return dp[n-1]

```

## 11.7 Subset Sum

Description: Check if subset sums to target. Time: O( $n \times \text{sum}$ ).

```

1 def subset_sum(arr, target):
2     n = len(arr)
3     dp = [[False] * (target + 1) for _ in range(n + 1)]
4
5     # Base case: sum 0 is always achievable
6     for i in range(n + 1):
7         dp[i][0] = True
8
9     for i in range(1, n + 1):
10        for s in range(target + 1):
11            # Don't take arr[i-1]
12            dp[i][s] = dp[i-1][s]
13
14            # Take arr[i-1] if possible
15            if s >= arr[i-1]:
16                dp[i][s] = dp[i][s] or dp[i-1][s - arr[i-1]]
17
18    return dp[n][target]
19
20 # Space optimized
21 def subset_sum_optimized(arr, target):
22     dp = [False] * (target + 1)
23     dp[0] = True
24
25     for num in arr:
26         for s in range(target, num - 1, -1):
27             dp[s] = dp[s] or dp[s - num]
28
29    return dp[target]

```

## 12 Array Techniques

### 12.1 Prefix Sum

Description: Precompute cumulative sums for O(1) range queries. Time: O( $n$ ) preprocessing, O(1) query.

```

1 # 1D prefix sum
2 prefix = [0] * (n + 1)
3 for i in range(n):
4     prefix[i + 1] = prefix[i] + arr[i]
5
6 # Range sum query [l, r] inclusive
7 range_sum = prefix[r + 1] - prefix[l]
8
9 # 2D prefix sum - for rectangle sum queries
10 def build_2d_prefix(matrix):
11     n, m = len(matrix), len(matrix[0])
12     prefix = [[0] * (m + 1) for _ in range(n + 1)]
13
14     for i in range(1, n + 1):
15         for j in range(1, m + 1):
16             prefix[i][j] = (matrix[i-1][j-1] +
17                             prefix[i-1][j] +
18                             prefix[i][j-1] -
19                             prefix[i-1][j-1])
20
21    return prefix
22
23 # Rectangle sum from (x1, y1) to (x2, y2) inclusive
24 def rect_sum(prefix, x1, y1, x2, y2):
25     return (prefix[x2+1][y2+1] -
26             prefix[x2+1][y2+1] -
27             prefix[x2+1][y1] +
28             prefix[x1][y1])

```

### 12.2 Difference Array

Description: Efficiently perform range updates. O(1) per update, O( $n$ ) to reconstruct.

```

1 # Initialize difference array
2 diff = [0] * (n + 1)
3
4 # Add 'val' to range [l, r]
5 def range_update(diff, l, r, val):
6     diff[l] += val
7     diff[r + 1] -= val
8
9 # After all updates, reconstruct array
10 def reconstruct(diff):
11     result = []
12     current = 0
13     for i in range(len(diff) - 1):
14         current += diff[i]
15         result.append(current)
16     return result
17
18 # Example: Multiple range updates
19 diff = [0] * (n + 1)
20 for l, r, val in updates:
21     range_update(diff, l, r, val)
22 final_array = reconstruct(diff)

```

## 12.3 Sliding Window

**Description:** Maintain a window of elements while traversing. Time:  $O(n)$ .

```

1 # Fixed size window
2 def max_sum_window(arr, k):
3     window_sum = sum(arr[:k])
4     max_sum = window_sum
5
6     # Slide window: add right, remove left
7     for i in range(k, len(arr)):
8         window_sum += arr[i] - arr[i - k]
9         max_sum = max(max_sum, window_sum)
10
11    return max_sum
12
13 # Variable size window - two pointers
14 def min_subarray_sum_geq_target(arr, target):
15     left = 0
16     current_sum = 0
17     min_len = float('inf')
18
19     for right in range(len(arr)):
20         current_sum += arr[right]
21
22         # Shrink window while condition holds
23         while current_sum >= target:
24             min_len = min(min_len, right - left + 1)
25             current_sum -= arr[left]
26             left += 1
27
28     return min_len if min_len != float('inf') else 0
29
30 # Longest substring with at most k distinct chars
31 def longest_k_distinct(s, k):
32     from collections import defaultdict
33
34     left = 0
35     char_count = defaultdict(int)
36     max_len = 0
37
38     for right in range(len(s)):
39         char_count[s[right]] += 1
40
41         # Shrink if too many distinct
42         while len(char_count) > k:
43             char_count[s[left]] -= 1
44             if char_count[s[left]] == 0:
45                 del char_count[s[left]]
46             left += 1
47
48     max_len = max(max_len, right - left + 1)
49
50    return max_len

```

## 13 Advanced Data Structures

### 13.1 Segment Tree

**Description:** Supports range queries and point updates in  $O(\log n)$ . Can be modified for range updates with lazy propagation.

```

1 class SegmentTree:
2     def __init__(self, arr):
3         self.n = len(arr)
4         # Tree size: 4n is safe upper bound
5         self.tree = [0] * (4 * self.n)
6         self.build(arr, 0, 0, self.n - 1)
7
8     def build(self, arr, node, start, end):
9         if start == end:

```

```

10        # Leaf node
11        self.tree[node] = arr[start]
12
13    else:
14        mid = (start + end) // 2
15        # Build left and right subtrees
16        self.build(arr, 2*node+1, start, mid)
17        self.build(arr, 2*node+2, mid+1, end)
18        # Combine results (sum in this case)
19        self.tree[node] = (self.tree[2*node+1] +
20                           self.tree[2*node+2])
21
22    def update(self, node, start, end, idx, val):
23        if start == end:
24            # Leaf node - update value
25            self.tree[node] = val
26        else:
27            mid = (start + end) // 2
28            if idx <= mid:
29                # Update left subtree
30                self.update(2*node+1, start, mid, idx, val)
31            else:
32                # Update right subtree
33                self.update(2*node+2, mid+1, end, idx, val)
34            # Recompute parent
35            self.tree[node] = (self.tree[2*node+1] +
36                               self.tree[2*node+2])
37
38    def query(self, node, start, end, l, r):
39        # No overlap
40        if r < start or end < l:
41            return 0
42
43        # Complete overlap
44        if l <= start and end <= r:
45            return self.tree[node]
46
47        # Partial overlap
48        mid = (start + end) // 2
49        left_sum = self.query(2*node+1, start, mid, l, r)
50        right_sum = self.query(2*node+2, mid+1, end, l, r)
51        return left_sum + right_sum
52
53    def update_val(self, idx, val):
54        self.update(0, 0, self.n-1, idx, val)
55
56    def range_sum(self, l, r):
57        return self.query(0, 0, self.n-1, l, r)

```

### 13.2 Fenwick Tree (Binary Indexed Tree)

**Description:** Simpler than segment tree, supports prefix sum and point updates in  $O(\log n)$ . More space efficient.

```

1 class FenwickTree:
2     def __init__(self, n):
3         self.n = n
4         # 1-indexed for easier implementation
5         self.tree = [0] * (n + 1)
6
7     def update(self, i, delta):
8         # Add delta to position i (1-indexed)
9         while i <= self.n:
10             self.tree[i] += delta
11             # Move to next node: add LSB
12             i += i & (-i)
13
14     def query(self, i):
15         # Get prefix sum up to i (1-indexed)
16         s = 0
17         while i > 0:
18             s += self.tree[i]
19             # Move to parent: remove LSB
20             i -= i & (-i)
21         return s
22
23     def range_query(self, l, r):
24         # Sum from l to r (1-indexed)
25         return self.query(r) - self.query(l - 1)
26
27 # Usage example
28 bit = FenwickTree(n)
29 for i, val in enumerate(arr, 1):
30     bit.update(i, val)
31
32 # Range sum [l, r] (1-indexed)
33 result = bit.range_query(l, r)

```

### 13.3 Trie (Prefix Tree)

**Description:** Tree for storing strings, enables fast prefix searches. Time:  $O(m)$  for operations where  $m$  is string length.

```

1 class TrieNode:
2     def __init__(self):

```

```

3     self.children = {} # char -> TrieNode
4     self.is_end = False # End of word marker
5
6 class Trie:
7     def __init__(self):
8         self.root = TrieNode()
9
10    def insert(self, word):
11        # Insert word - O(len(word))
12        node = self.root
13        for char in word:
14            if char not in node.children:
15                node.children[char] = TrieNode()
16            node = node.children[char]
17        node.is_end = True
18
19    def search(self, word):
20        # Exact word search - O(len(word))
21        node = self.root
22        for char in word:
23            if char not in node.children:
24                return False
25            node = node.children[char]
26        return node.is_end
27
28    def starts_with(self, prefix):
29        # Prefix search - O(len(prefix))
30        node = self.root
31        for char in prefix:
32            if char not in node.children:
33                return False
34            node = node.children[char]
35        return True
36
37    # Find all words with given prefix
38    def words_with_prefix(self, prefix):
39        node = self.root
40        for char in prefix:
41            if char not in node.children:
42                return []
43            node = node.children[char]
44
45        # DFS to collect all words
46        words = []
47        def dfs(n, path):
48            if n.is_end:
49                words.append(prefix + path)
50            for char, child in n.children.items():
51                dfs(child, path + char)
52
53        dfs(node, "")
54        return words

```

```

38     _upd(t)
39     return (l, t)
40
41 else:
42     t.right, r = _split(t.right, key)
43     _upd(t)
44     return (t, r)
45
46 def _erase(t, key):
47     if not t:
48         return None
49     if key == t.key:
50         return _merge(t.left, t.right)
51     if key < t.key:
52         t.left = _erase(t.left, key)
53     else:
54         t.right = _erase(t.right, key)
55     _upd(t)
56     return t
57
58 class Treap:
59     def __init__(self):
60         self.root = None
61
62     def __len__(self):
63         return _sz(self.root)
64
65     def contains(self, key):
66         t = self.root
67         while t:
68             if key == t.key:
69                 return True
70             t = t.left if key < t.key else t.right
71         return False
72
73     def insert(self, key):
74         if self.contains(key):
75             return
76         node = TreapNode(key)
77         l, r = _split(self.root, key)
78         self.root = _merge(_merge(l, node), r)
79
80     def remove(self, key):
81         self.root = _erase(self.root, key)
82
83     def kth_smallest(self, k):
84         # 0-indexed k
85         t = self.root
86         while t:
87             ls = _sz(t.left)
88             if k < ls:
89                 t = t.left
90             elif k == ls:
91                 return t.key
92             else:
93                 k -= ls + 1
94                 t = t.right
95         return None # k out of range
96
97     def count_less_than(self, key):
98         # number of keys < key
99         t, cnt = self.root, 0
100        while t:
101            if key <= t.key:
102                t = t.left
103            else:
104                cnt += 1 + _sz(t.left)
105                t = t.right
106        return cnt
107
108    def lower_bound(self, key):
109        # smallest key >= key; returns None if none
110        t, ans = self.root, None
111        while t:
112            if t.key >= key:
113                ans = t.key
114                t = t.left
115            else:
116                t = t.right
117        return ans
118
119    # Usage example
120    T = Treap()
121    for x in [5, 1, 7, 3]:
122        T.insert(x)
123    T.contains(3) # True
124    T.kth_smallest(1) # 3 (0-indexed)
125    T.count_less_than(6) # 3 (1,3,5)
126    T.remove(5)
127    len(T) # 3

```

## 13.4 Treap (Randomized Balanced BST)

**Description:** Ordered set/map with expected  $O(\log n)$  insert, erase, search, k-th, and rank. Combines a BST by key and a heap by random priority. Stores unique keys; for multiset, store (key, uid) or maintain a count.

```

1 import random
2
3 class TreapNode:
4     __slots__ = ("key", "prio", "left", "right", "size")
5     def __init__(self, key):
6         self.key = key
7         self.prio = random.randint(1, 1 << 30)
8         self.left = None
9         self.right = None
10        self.size = 1
11
12    def _sz(t):
13        return t.size if t else 0
14
15    def _upd(t):
16        if t:
17            t.size = 1 + _sz(t.left) + _sz(t.right)
18
19    def _merge(a, b):
20        # assumes all keys in a < all keys in b
21        if not a or not b:
22            return a or b
23        if a.prio > b.prio:
24            a.right = _merge(a.right, b)
25            _upd(a)
26            return a
27        else:
28            b.left = _merge(a, b.left)
29            _upd(b)
30            return b
31
32    def _split(t, key):
33        # returns (l, r): l has keys < key, r has keys >= key
34        if not t:
35            return (None, None)
36        if key <= t.key:
37            l, t.left = _split(t.left, key)

```

## 14 Bit Manipulation

**Description:** Efficient operations using bitwise operators. Useful for sets, flags, and optimization.

```

1 # Check if i-th bit (0-indexed) is set
2 is_set = (n >> i) & 1
3
4 # Set i-th bit to 1
5 n |= (1 << i)
6
7 # Clear i-th bit (set to 0)
8 n &= ~(1 << i)
9
10 # Toggle i-th bit
11 n ^= (1 << i)
12
13 # Count set bits (popcount)
14 count = bin(n).count('1')
15 count = n.bit_count() # Python 3.10+
16
17 # Get lowest set bit
18 lsb = n & -n # Also n & (~n + 1)
19
20 # Remove lowest set bit
21 n &= (n - 1)
22
23 # Check if power of 2
24 is_pow2 = n > 0 and (n & (n - 1)) == 0
25
26 # Check if power of 4
27 is_pow4 = n > 0 and (n & (n-1)) == 0 and (n & 0x55555555) != 0
28
29 # Iterate over all subsets of set represented by mask
30 mask = (1 << n) - 1 # All bits set
31 submask = mask
32 while submask > 0:
33     # Process submask
34     submask = (submask - 1) & mask
35
36 # Iterate through all k-bit masks
37 def iterate_k_bits(n, k):
38     mask = (1 << k) - 1
39     while mask < (1 << n):
40         # Process mask
41         yield mask
42         # Gosper's hack
43         c = mask & ~mask
44         r = mask + c
45         mask = (((r ^ mask) >> 2) // c) | r
46
47 # XOR properties
48 # a ^ a = 0 (number XOR itself is 0)
49 # a ^ 0 = a (number XOR 0 is itself)
50 # XOR is commutative and associative
51 # Find unique element when all others appear twice:
52 def find_unique(arr):
53     result = 0
54     for x in arr:
55         result ^= x
56     return result
57
58 # Subset enumeration
59 n = 5 # Number of elements
60 for mask in range(1 << n):
61     subset = [i for i in range(n) if mask & (1 << i)]
62     # Process subset
63
64 # Check parity (odd/even number of 1s)
65 def parity(n):
66     count = 0
67     while n:
68         count ^= 1
69         n &= n - 1
70     return count # 1 if odd, 0 if even
71
72 # Swap two numbers without temp variable
73 a, b = 5, 10
74 a ^= b
75 b ^= a
76 a ^= b
77 # Now a=10, b=5

```

```

10
11     return C
12
13 # With modulo
14 def matmul_mod(A, B, mod):
15     n = len(A)
16     C = [[0] * n for _ in range(n)]
17
18     for i in range(n):
19         for j in range(n):
20             for k in range(n):
21                 C[i][j] = (C[i][j] +
22                             A[i][k] * B[k][j]) % mod
23
24     return C

```

## 15.2 Matrix Exponentiation

**Description:** Compute  $M^n$  in  $O(k^3 \log n)$  where  $k$  is matrix dimension. Used for solving linear recurrences efficiently.

```

1 def matpow(M, n, mod):
2     size = len(M)
3
4     # Identity matrix
5     result = [[1 if i==j else 0
6                for j in range(size)]
7               for i in range(size)]
8
9     # Binary exponentiation
10    while n > 0:
11        if n & 1:
12            result = matmul_mod(result, M, mod)
13        M = matmul_mod(M, M, mod)
14        n >>= 1
15
16    return result
17
18 # Example: Fibonacci using matrix exponentiation
19 # F(n) = [[1, 1], [1, 0]]^n
20 def fibonacci(n, mod):
21     if n == 0: return 0
22     if n == 1: return 1
23
24     M = [[1, 1], [1, 0]]
25     result = matpow(M, n - 1, mod)
26     return result[0][0]
27
28 # Linear recurrence: a(n) = c1*a(n-1) + c2*a(n-2) + ...
29 # Build transition matrix and use matrix exponentiation
30 def linear_recurrence(coeffs, init, n, mod):
31     k = len(coeffs)
32
33     if n < k:
34         return init[n]
35
36     # Transition matrix
37     # [a(n), a(n-1), ..., a(n-k+1)]
38     M = [[0] * k for _ in range(k)]
39     M[0] = coeffs # First row
40     for i in range(1, k):
41         M[i][i-1] = 1 # Identity for shifting
42
43     # Initial state vector [a(k-1), a(k-2), ..., a(0)]
44     state = init[k-1::-1]
45
46     #  $M^{(n-k+1)}$ 
47     result_matrix = matpow(M, n - k + 1, mod)
48
49     # Multiply with initial state
50     result = 0
51     for i in range(k):
52         result = (result + result_matrix[0][i] * state[i]) % mod
53
54     return result
55
56 # Example: Tribonacci T(n) = T(n-1) + T(n-2) + T(n-3)
57 def tribonacci(n, mod):
58     if n == 0: return 0
59     if n == 1 or n == 2: return 1
60
61     coeffs = [1, 1, 1]
62     init = [0, 1, 1]
63     return linear_recurrence(coeffs, init, n, mod)

```

# 15 Matrix Operations

**Description:** Matrix operations for DP optimization, graph algorithms, and recurrence relations.

## 15.1 Matrix Multiplication

```

1 # Standard matrix multiplication - O(n^3)
2 def matmul(A, B):
3     n, m, p = len(A), len(A[0]), len(B[0])
4     C = [[0] * p for _ in range(n)]
5
6     for i in range(n):
7         for j in range(p):
8             for k in range(m):
9                 C[i][j] += A[i][k] * B[k][j]

```

## 16 Miscellaneous Tips

### 16.1 Python-Specific Optimizations

```

1 # Fast input for large datasets
2 import sys
3 input = sys.stdin.readline
4
5 # Increase recursion limit for deep DFS/DP

```

```

6 sys.setrecursionlimit(10**6)
7
8 # Threading for higher stack limit (CAUTION: use carefully)
9 import threading
10 threading.stack_size(2**26) # 64MB
11 sys.setrecursionlimit(2**20)
12
13 # Deep copy (be careful with performance)
14 from copy import deepcopy
15 new_list = deepcopy(old_list)
16
17 # Fast output (for printing large results)
18 import sys
19 print = sys.stdout.write # Only use for string output

```

## 16.2 Useful Libraries

```

1 # Iterator tools - powerful combinations
2 from itertools import *
3
4 # permutations(iterable, r) - all r-length permutations
5 perms = list(permutations([1,2,3], 2))
6 # [(1,2), (1,3), (2,1), (2,3), (3,1), (3,2)]
7
8 # combinations(iterable, r) - r-length combinations
9 combs = list(combinations([1,2,3], 2))
10 # [(1,2), (1,3), (2,3)]
11
12 # product - cartesian product
13 prod = list(product([1,2], ['a','b']))
14 # [(1,'a'), (1,'b'), (2,'a'), (2,'b')]
15
16 # accumulate - running totals
17 acc = list(accumulate([1,2,3,4]))
18 # [1, 3, 6, 10]
19
20 # chain - flatten iterables
21 chained = list(chain([1,2], [3,4]))
22 # [1, 2, 3, 4]

```

## 16.3 Common Patterns

```

1 # Lambda sorting with multiple keys
2 arr.sort(key=lambda x: (-x[0], x[1]))
3 # Sort by first desc, then second asc
4
5 # All/Any - short-circuit evaluation
6 all(x > 0 for x in arr) # True if all positive
7 any(x > 0 for x in arr) # True if any positive
8
9 # Zip - parallel iteration
10 for a, b in zip(list1, list2):
11     pass
12
13 # Enumerate - index and value
14 for i, val in enumerate(arr):
15     print(f"arr[{i}] = {val}")
16
17 # Custom comparison function
18 from functools import cmp_to_key
19
20 def compare(a, b):
21     # Return -1 if a < b, 0 if equal, 1 if a > b
22     if a + b > b + a:
23         return -1
24     return 1
25
26 arr.sort(key=cmp_to_key(compare))
27
28 # DefaultDict with lambda
29 from collections import defaultdict
30 d = defaultdict(lambda: float('inf'))
31
32 # Multiple assignment
33 a, b = b, a # Swap
34 a, *rest, b = [1,2,3,4,5] # a=1, rest=[2,3,4], b=5

```

## 16.4 Common Pitfalls

```

1 # Integer division - floors toward negative infinity
2 print(7 // 3) # 2
3 print(-7 // 3) # -3 (not -2!)
4
5 # For ceiling division toward zero:
6 def div_ceil(a, b):
7     return -(a // b)
8
9 # Modulo with negative numbers
10 print((-5) % 3) # 1 (not -2!)
11 print(5 % -3) # -1
12
13 # List multiplication creates references!
14 matrix = [[0] * m] * n # WRONG! All rows same object
15 matrix[0][0] = 1 # Changes all rows!

```

```

16 # Correct way
17 matrix = [[0] * m for _ in range(n)]
18
19 # Float comparison - don't use ==
20 a, b = 0.1 + 0.2, 0.3
21 print(a == b) # False!
22
23 # Use epsilon comparison
24 eps = 1e-9
25 print(abs(a - b) < eps) # True
26
27 # String immutability
28 s = "abc"
29 # s[0] = 'd' # ERROR!
30 s = 'd' + s[1:] # OK
31
32 # For many string mutations, use list
33 chars = list(s)
34 chars[0] = 'd'
35 s = ''.join(chars)
36
37 # Mutable default arguments - dangerous!
38 def func(arr=[]): # WRONG!
39     arr.append(1)
40     return arr
41
42 # Each call modifies same list
43 print(func()) # [1]
44 print(func()) # [1, 1]
45
46 # Correct way
47 def func(arr=None):
48     if arr is None:
49         arr = []
50     arr.append(1)
51     return arr
52
53
54 # Generator expressions save memory
55 sum(xx for x in range(10**6)) # Memory efficient
56 # vs
57 sum([xx for x in range(10**6)]) # Creates full list
58
59 # Ternary operator
60 x = a if condition else b
61
62 # Dictionary get with default
63 count = d.get(key, 0) + 1
64
65 # Matrix rotation 90 degrees clockwise
66 def rotate_90(matrix):
67     return [list(row) for row in zip(*matrix[::-1])]
68
69 # Matrix transpose
70 def transpose(matrix):
71     return [list(row) for row in zip(*matrix)]

```

## 16.5 Time Complexity Reference

Common time complexities (Python, rough guides for 1–2s limits):

- $O(1)$ ,  $O(\log n)$ : instant
- $O(n)$ : usually fine up to  $\sim 10^7$  operations ( $\sim 1s$ )
- $O(n \log n)$ : OK for  $n$  up to several  $10^5$  depending on constants
- $O(n\sqrt{n})$ : risky in Python (may be OK for  $n$  up to a few  $10^4$  with low constants)
- $O(n^2)$ : often TLE for  $n > 10^4$
- $O(2^n)$ : TLE for  $n > 20$  (unless heavy pruning/memoization)
- $O(n!)$ : TLE for  $n > 11$

Input size guidelines (Python-focused):

- $n \leq 12$ :  $O(n!)$  (brute-force permutations)
- $n \leq 20$ :  $O(2^n)$  (subset DP / bitmask DP)
- $n \leq 500$ :  $O(n^3)$  may sometimes pass for small constants
- $n \leq 5000$ :  $O(n^2)$  borderline; optimize heavily
- $n \leq 10^6$ :  $O(n \log n)$  common;  $O(n)$  preferred when possible
- $n \leq 10^7$ :  $O(n)$  may be OK for tight loops
- $n > 10^7$ : aim for  $O(n)$  with very low constants, or  $O(\log n)/O(1)$

Complexity examples (Python implementations)

- $O(1)$ : array access, dictionary lookup, push/pop from list end.
- $O(\log n)$ : binary search (bisect), heap push/pop (heappq), operations in sortedcontainers.
- $O(n)$ : single-pass scans, two-pointers, prefix sums, counting frequencies (Counter).
- $O(n \log n)$ : sorting (Timsort via sorted()/list.sort()), heap construction, divide-and-conquer merges.
- $O(n\sqrt{n})$ : sqrt-decomposition queries, some Mo's algorithm variants (constant-sensitive).
- $O(n^2)$ : nested loops for pairwise checks, naive DP on pairs (be cautious for  $n > 10,000$ ).
- $O(n^3)$ : triple loops (Floyd–Warshall), usually too slow unless  $n \leq 200$ .
- $O(2^n)$ : bitmask DP, subset enumerations, recursion over subsets (recom-

mended for  $n \leq 20$ ).

- $O(n!)$ : full permutations, exhaustive search over orderings (recommended for  $n \leq 10$ ; occasionally up to 11).

**How to use:** This quick reference maps input size  $n$  (left) to typical feasible time complexities (right) for contest time limits (1–2s) targeting Python implementations. Use it to pick algorithmic approaches and to decide when to optimize or change strategy.

#### Notes on filling the table:

- Start by checking the problem's time limit and target language. These guidelines are Python-focused (assume roughly  $\approx 10^7$  simple operations/s; actual throughput depends on implementation details and input shapes). 17  
18  
19
- Convert algorithm cost to operation count: roughly cost =  $c \cdot f(n)$ . If cost  $>$  time\_limit  $\times$  ops\_per\_sec, it will TLE. 20  
21  
22
- When in doubt, aim one complexity class lower (e.g. prefer  $O(n \log n)$  over  $O(n^2)$  for  $n$  around  $10^5$ ). 23  
24
- Consider memory limits—some faster algorithms use more memory (e.g. segment trees vs. Fenwick tree). 25  
26
- For multivariate inputs, replace  $n$  with the product/dominant parameter (e.g.  $n \cdot m$ ) and apply the same rules. 27  
28
- If an algorithm theoretically fits but is close to the limit, try to reduce constant factors: use local variables, avoid heavy Python objects in inner loops; use built-in functions, or move hot code to PyPy/Cython if allowed. 29  
30  
31  
32  
33  
34  
35  
36  
37  
38

## 17 Computational Geometry

### 17.1 Basic Geometry

**Description:** Fundamental geometric operations for 2D points.

```

1 # Assumes 'math' and basic constants (e.g. EPS) are available see
  Mathematics section for shared imports.
2
3 # Point operations
4 def dist(p1, p2):
5     # Euclidean distance
6     return math.sqrt((p1[0] - p2[0])**2 + (p1[1] - p2[1])**2)
7
8 def cross_product(O, A, B):
9     # Cross product of vectors OA and OB
10    # Positive: counter-clockwise
11    # Negative: clockwise
12    # Zero: collinear
13    return (A[0] - O[0]) * (B[1] - O[1]) - \
14        (A[1] - O[1]) * (B[0] - O[0])
15
16 def dot_product(A, B, C, D):
17     # Dot product of vectors AB and CD
18     return (B[0] - A[0]) * (D[0] - C[0]) + \
19         (B[1] - A[1]) * (D[1] - C[1])
20
21 # Check if point is on segment
22 def on_segment(p, q, r):
23     # Check if q lies on segment pr
24     return (q[0] <= max(p[0], r[0]) and
25             q[0] >= min(p[0], r[0]) and
26             q[1] <= max(p[1], r[1]) and
27             q[1] >= min(p[1], r[1]))
28
29 # Segment intersection
30 def segments_intersect(p1, q1, p2, q2):
31     o1 = cross_product(p1, q1, p2)
32     o2 = cross_product(p1, q1, q2)
33     o3 = cross_product(p2, q2, p1)
34     o4 = cross_product(p2, q2, q1)
35
36     # General case
37     if o1 * o2 < 0 and o3 * o4 < 0:
38         return True
39
40     # Special cases (collinear)
41     if o1 == 0 and on_segment(p1, p2, q1):
42         return True
43     if o2 == 0 and on_segment(p1, q2, q1):
44         return True
45     if o3 == 0 and on_segment(p2, p1, q2):
46         return True
47     if o4 == 0 and on_segment(p2, q1, q2):
48         return True
49
50     return False

```

### 17.2 Convex Hull

**Description:** Find convex hull using Graham's scan. Time:  $O(n \log n)$ .

```

1 def convex_hull(points):
2     # Graham's scan algorithm
3     points = sorted(points) # Sort by x, then y
4
5     if len(points) <= 2:
6         return points

```

```

7
8     # Build lower hull
9     lower = []
10    for p in points:
11        while (len(lower) >= 2 and
12               cross_product(lower[-2], lower[-1], p) <= 0):
13            lower.pop()
14        lower.append(p)
15
16    # Build upper hull
17    upper = []
18    for p in reversed(points):
19        while (len(upper) >= 2 and
20               cross_product(upper[-2], upper[-1], p) <= 0):
21            upper.pop()
22        upper.append(p)
23
24    # Remove last point (duplicate of first)
25    return lower[:-1] + upper[:-1]
26
27 # Convex hull area
28 def polygon_area(points):
29     # Shoelace formula
30     n = len(points)
31     area = 0
32
33     for i in range(n):
34         j = (i + 1) % n
35         area += points[i][0] * points[j][1]
36         area -= points[j][0] * points[i][1]
37
38     return abs(area) / 2

```

### 17.3 Point in Polygon

**Description:** Check if point is inside polygon. Time:  $O(n)$ .

```

1 def point_in_polygon(point, polygon):
2     # Ray casting algorithm
3     x, y = point
4     n = len(polygon)
5     inside = False
6
7     p1x, p1y = polygon[0]
8     for i in range(1, n + 1):
9         p2x, p2y = polygon[i % n]
10
11    if y > min(p1y, p2y):
12        if y <= max(p1y, p2y):
13            if x <= max(p1x, p2x):
14                if p1y != p2y:
15                    xinters = (y - p1y) * (p2x - p1x) / \
16                               (p2y - p1y) + p1x
17
18                if p1x == p2x or x <= xinters:
19                    inside = not inside
20
21    p1x, p1y = p2x, p2y
22
23    return inside

```

### 17.4 Closest Pair of Points

**Description:** Find closest pair using divide and conquer. Time:  $O(n \log n)$ .

```

1 def closest_pair(points):
2     points_sorted_x = sorted(points, key=lambda p: p[0])
3     points_sorted_y = sorted(points, key=lambda p: p[1])
4
5     def closest_recursive(px, py):
6         n = len(px)
7
8         # Base case: brute force
9         if n <= 3:
10            min_dist = float('inf')
11            for i in range(n):
12                for j in range(i + 1, n):
13                    min_dist = min(min_dist, dist(px[i], px[j]))
14            return min_dist
15
16         # Divide
17         mid = n // 2
18         midpoint = px[mid]
19
20         pyl = [p for p in py if p[0] <= midpoint[0]]
21         pyr = [p for p in py if p[0] > midpoint[0]]
22
23         # Conquer
24         dl = closest_recursive(px[:mid], pyl)
25         dr = closest_recursive(px[mid:], pyr)
26         d = min(dl, dr)
27
28         # Combine: check strip
29         strip = [p for p in py if abs(p[0] - midpoint[0]) < d]
30
31         for i in range(len(strip)):

```

```
32     j = i + 1
33     while j < len(strip) and strip[j][1] - strip[i][1] < d:
34         d = min(d, dist(strip[i], strip[j]))
35         j += 1
36
37     return d
38
39
return closest_recursive(points_sorted_x, points_sorted_y)
```

```

dx = convex[i][0] - convex[j][0]
dy = convex[i][1] - convex[j][1]
return dx*dx + dy*dy

t = 0
i in range(n):
    ni = (i + 1) % n
    while True:
        nj = (j + 1) % n
        cur = abs(cross_vec(vec(convex[i], convex[ni]), vec(convex[j], convex[nj])))
        nex = abs(cross_vec(vec(convex[i], convex[ni]), vec(convex[nj], convex[(nj+1)%n])))
        if False else None
        # move j while area increases (alternate robust check below)
        if dot_vec(vec(convex[i], convex[ni]), vec(convex[j], convex[nj])) > 0:
            j = nj
        else:
            break
    best = max(best, sqdist(i, j))
urn math.sqrt(best)

fer integer arithmetic (area2, orientation) whenever possible.
: Fraction for exact intersections if you must compare intersection
ordinates.
: did atan2 in tight loops; use cross/dot for comparisons when
possible.

```

## 17.5 Robust techniques and helpers for ICPC

**Description:** Practical tips and helper functions that make geometric code safer and easier to reason about in contests: prefer integer arithmetic when possible, use exact rationals for intersections, and common utilities.

```

1 # Assumes 'math', 'Fraction' and shared constants (EPS) are available see
2     Mathematics section for shared imports.
3
4 # Vector helpers
5 def vec(a, b):
6     return (b[0] - a[0], b[1] - a[1])
7
8 def cross_vec(u, v):
9     return u[0] * v[1] - u[1] * v[0]
10
11 def dot_vec(u, v):
12     return u[0] * v[0] + u[1] * v[1]
13
14 # Signed orientation: >0 = ccw, <0 = cw, 0 = collinear
15 def orient(a, b, c):
16     return cross_vec(vec(a, b), vec(a, c))
17
18 # Twice signed polygon area (integer if points are integer)
19 def polygon_area2(poly):
20     s = 0
21     n = len(poly)
22     for i in range(n):
23         j = (i + 1) % n
24         s += cross_vec(poly[i], poly[j])
25     return s
26
27 # Exact intersection point of two lines (a-b) and (c-d)
28 # Returns (Fraction, Fraction) or None when parallel/collinear
29 def line_intersection(a, b, c, d):
30     r = vec(a, b)
31     s = vec(c, d)
32     rxs = cross_vec(r, s)
33     if rxs == 0:
34         return None
35     t_num = cross_vec(vec(a, c), s)
36     t = Fraction(t_num, rxs)
37     return (Fraction(a[0]) + t * Fraction(r[0]), Fraction(a[1]) + t * Fraction(r[1]))
38
39 # Project point p onto line a-b, returns a point (float)
40 def project_point_on_line(a, b, p):
41     r = vec(a, b)
42     if r == (0, 0):
43         return a
44     t = dot_vec(vec(a, p), r) / dot_vec(r, r)
45     return (a[0] + t * r[0], a[1] + t * r[1])
46
47 # Distance from point p to line a-b (float)
48 def point_line_distance(a, b, p):
49     r = vec(a, b)
50     if r == (0, 0):
51         return dist(a, p)
52     return abs(cross_vec(vec(a, b), vec(a, p))) / math.hypot(r[0], r[1])
53
54 # Rotate point p around origin by theta (radians)
55 def rotate(p, theta):
56     x, y = p
57     c, s = math.cos(theta), math.sin(theta)
58     return (x * c - y * s, x * s + y * c)
59
60 # Polar sort key around origin (or given center)
61 def polar_sort_key(origin=(0,0)):
62     def key(p):
63         return (math.atan2(p[1]-origin[1], p[0]-origin[0]), (p[0]-origin[0])**2 + (p[1]-origin[1])**2)
64     return key
65
66 # Rotating calipers to compute diameter (farthest pair) of convex polygon
67 # convex_hull is expected to return points in CCW order without duplicate
68 # endpoints
69 def diameter(convex):
70     n = len(convex)
71     if n == 0:
72         return 0
73     if n == 1:
74         return 0
75     if n == 2:
76         return math.hypot(convex[0][0]-convex[1][0], convex[0][1]-convex[1][1])
77
78     def sadist(i, j):
79
80

```

18 Network Flow

## 18.1 Maximum Flow - Edmonds-Karp (BFS-based Ford-Fulkerson)

**Description:** Find maximum flow from source to sink. Time:  $O(VE^2)$ .

```

from collections import deque, defaultdict

def max_flow(graph, source, sink, n):
    # graph[u][v] = capacity from u to v
    # Build residual graph
    residual = defaultdict(lambda: defaultdict(int))
    for u in graph:
        for v in graph[u]:
            residual[u][v] = graph[u][v]

    def bfs_path():
        # Find augmenting path using BFS
        parent = {source: None}
        visited = {source}
        queue = deque([source])

        while queue:
            u = queue.popleft()

            if u == sink:
                # Reconstruct path
                path = []
                while parent[u] is not None:
                    path.append((parent[u], u))
                    u = parent[u]
                return path[::-1]

            for v in range(n):
                if v not in visited and residual[u][v] > 0:
                    visited.add(v)
                    parent[v] = u
                    queue.append(v)

        return None

    max_flow_value = 0

    # Find augmenting paths
    while True:
        path = bfs_path()
        if path is None:
            break

        # Find minimum capacity along path
        flow = min(residual[u][v] for u, v in path)

        # Update residual graph
        for u, v in path:
            residual[u][v] -= flow
            residual[v][u] += flow

        max_flow_value += flow

    return max_flow_value

# Example usage
# graph[u][v] = capacity

```

```

58 graph = defaultdict(lambda: defaultdict(int))
59 graph[0][1] = 10
60 graph[0][2] = 10
61 graph[1][3] = 4
62 graph[1][4] = 8
63 graph[2][4] = 9
64 graph[3][5] = 10
65 graph[4][3] = 6
66 graph[4][5] = 10
67
68 n = 6 # Number of nodes
69 result = max_flow(graph, 0, 5, n)

```

## 18.2 Dinic's Algorithm (Faster)

**Description:** Faster max flow using level graph and blocking flow. Time:  $O(V^2E)$ .

```

6 visited[source] = True
7
8 while queue:
9     u = queue.popleft()
10    for v in range(n):
11        if not visited[v] and residual[u][v] > 0:
12            visited[v] = True
13            queue.append(v)
14
15 # Cut edges
16 cut_edges = []
17 for u in range(n):
18    if visited[u]:
19        for v in range(n):
20            if not visited[v] and graph[u][v] > 0:
21                cut_edges.append((u, v))
22
23 return cut_edges

```

```

1 from collections import deque, defaultdict
2
3 class Dinic:
4     def __init__(self, n):
5         self.n = n
6         self.graph = defaultdict(lambda: defaultdict(int))
7
8     def add_edge(self, u, v, cap):
9         self.graph[u][v] += cap
10
11     def bfs(self, source, sink):
12         # Build level graph
13         level = [-1] * self.n
14         level[source] = 0
15         queue = deque([source])
16
17         while queue:
18             u = queue.popleft()
19
20             for v in range(self.n):
21                 if level[v] == -1 and self.graph[u][v] > 0:
22                     level[v] = level[u] + 1
23                     queue.append(v)
24
25         return level if level[sink] != -1 else None
26
27     def dfs(self, u, sink, pushed, level, start):
28         if u == sink:
29             return pushed
30
31         while start[u] < self.n:
32             v = start[u]
33
34             if (level[v] == level[u] + 1 and
35                 self.graph[u][v] > 0):
36
37                 flow = self.dfs(v, sink,
38                                   min(pushed, self.graph[u][v]),
39                                   level, start)
40
41                 if flow > 0:
42                     self.graph[u][v] -= flow
43                     self.graph[v][u] += flow
44                     return flow
45
46             start[u] += 1
47
48         return 0
49
50     def max_flow(self, source, sink):
51         flow = 0
52
53         while True:
54             level = self.bfs(source, sink)
55             if level is None:
56                 break
57
58             start = [0] * self.n
59
60             while True:
61                 pushed = self.dfs(source, sink, float('inf'),
62                                   level, start)
63                 if pushed == 0:
64                     break
65                 flow += pushed
66
67         return flow

```

## 18.3 Min Cut

**Description:** Find minimum cut after computing max flow.

```

1 def min_cut(graph, source, n, residual):
2     # After running max_flow, residual graph is available
3     # Min cut = set of reachable nodes from source
4     visited = [False] * n
5     queue = deque([source])

```

## 18.4 Bipartite Matching

**Description:** Maximum matching in bipartite graph using flow.

```

1 def max_bipartite_matching(left_size, right_size, edges):
2     # edges = [(left_node, right_node), ...]
3     # Add source (0) and sink (left_size + right_size + 1)
4
5     n = left_size + right_size + 2
6     source = 0
7     sink = n - 1
8
9     graph = defaultdict(lambda: defaultdict(int))
10
11     # Source to left nodes
12     for i in range(1, left_size + 1):
13         graph[source][i] = 1
14
15     # Left to right edges
16     for l, r in edges:
17         graph[l + 1][left_size + r + 1] = 1
18
19     # Right nodes to sink
20     for i in range(1, right_size + 1):
21         graph[left_size + i][sink] = 1
22
23     return max_flow(graph, source, sink, n)

```

## 19 Advanced Python Syntax

**Description:** Beyond basic loops - powerful Python constructs for competitive programming.

### 19.1 List Comprehensions and Generators

```

1 # Basic list comprehension
2 squares = [x**2 for x in range(10)]
3 evens = [x for x in range(20) if x % 2 == 0]
4
5 # Nested comprehensions
6 matrix = [[i*j for j in range(5)] for i in range(3)]
7 flattened = [x for row in matrix for x in row]
8
9 # Dictionary comprehensions
10 char_count = {char: text.count(char) for char in set(text)}
11 squares_dict = {x: x**2 for x in range(5)}
12
13 # Set comprehensions
14 unique_lengths = {len(word) for word in words}
15
16 # Generator expressions (memory efficient)
17 sum_squares = sum(x**2 for x in range(1000000))
18 any_even = any(x % 2 == 0 for x in numbers)
19
20 # Conditional expressions in comprehensions
21 processed = [x if x > 0 else 0 for x in numbers]
22 filtered = [x for x in numbers if x > 0 and x < 100]

```

### 19.2 Advanced Iteration Patterns

```

1 # Zip - parallel iteration
2 names = ['Alice', 'Bob', 'Charlie']
3 scores = [85, 92, 78]
4 for name, score in zip(names, scores):
5     print(f'{name}: {score}')
6
7 # Enumerate with custom start
8 for i, item in enumerate(items, 1): # Start from 1
9     print(f'Item {i}: {item}')
10
11 # Zip longest (from itertools)
12 from itertools import zip_longest
13 for a, b in zip_longest(list1, list2, fillvalue=0):
14     # Continues until longest list is exhausted

```

```

15 # Unpacking with star operator
16 first, *middle, last = [1, 2, 3, 4, 5]
17 # first=1, middle=[2,3,4], last=5
18
19 # Multiple assignment
20 a, b = b, a # Swap variables
21 x, y, z = input().split() # Parse multiple inputs
22
23 # Walking with indices and slicing
24 for i in range(len(arr) - 1):
25     current, next_item = arr[i], arr[i + 1]
26
27 # Sliding window with enumerate
28 for i, val in enumerate(arr[:-k+1]):
29     window = arr[i:i+k] # k-length sliding window
30

```

## 19.3 Advanced Data Structure Operations

```

1 # Dictionary operations
2 from collections import defaultdict, Counter
3
4 # DefaultDict with different types
5 adj_list = defaultdict(list) # Adjacency list
6 counts = defaultdict(int) # Frequency counter
7 groups = defaultdict(set) # Group sets
8
9 # Counter - frequency counting
10 text = "hello world"
11 freq = Counter(text)
12 print(freq['l']) # 3
13 most_common = freq.most_common(3) # Top 3 frequent
14
15 # Dictionary merging (Python 3.9+)
16 dict1 = {'a': 1, 'b': 2}
17 dict2 = {'c': 3, 'd': 4}
18 merged = dict1 | dict2
19
20 # Dictionary comprehension with conditions
21 filtered_dict = {k: v for k, v in original.items() if v > 0}
22
23 # Nested dictionary access
24 from collections import defaultdict
25 nested = defaultdict(lambda: defaultdict(int))
26 nested[key1][key2] += 1
27
28 # Set operations
29 set1 & set2 # Intersection
30 set1 | set2 # Union
31 set1 - set2 # Difference
32 set1 ^ set2 # Symmetric difference
33
34 # List slicing tricks
35 arr[::-1] # Reverse
36 arr[::2] # Every 2nd element
37 arr[1::2] # Every 2nd starting from index 1
38 arr[-3:] # Last 3 elements
39 arr[:-2] # All except last 2

```

## 19.4 Functional Programming Features

```

1 # Map, filter, reduce
2 numbers = [1, 2, 3, 4, 5]
3
4 # Map - apply function to each element
5 doubled = list(map(lambda x: x * 2, numbers))
6 strings = list(map(str, numbers))
7
8 # Filter - select elements meeting condition
9 evens = list(filter(lambda x: x % 2 == 0, numbers))
10
11 # Reduce - accumulate values (from functools)
12 from functools import reduce
13 product = reduce(lambda x, y: x * y, numbers)
14 max_val = reduce(max, numbers)
15
16 # Lambda functions for sorting
17 points = [(1, 3), (2, 1), (0, 5)]
18 sorted_by_y = sorted(points, key=lambda p: p[1])
19 sorted_by_dist = sorted(points, key=lambda p: p[0]**2 + p[1]**2)
20
21 # Multiple sort keys
22 students = [('Alice', 85, 'A'), ('Bob', 92, 'B'), ('Charlie', 85, 'A')]
23 # Sort by grade (desc), then score (desc), then name (asc)
24 sorted_students = sorted(students,
25                         key=lambda x: (-ord(x[2]), -x[1], x[0]))

```

## 19.5 Advanced String Operations

```

1 # String formatting
2 name, score = "Alice", 95
3 formatted = f"{name} scored {score:>5} points" # Right align
4 binary = f"{42:08b}" # 00101010 (8-digit binary)
5 hex_val = f"{255:02x}" # ff (2-digit hex)

```

```

6 # String methods for parsing
7 text = " hello,world "
8 words = text.strip().split(',')
9 joined = ' | '.join(words)
10
11 # String manipulation
12 s = "programming"
13 s.startswith("prog") # True
14 s.endswith("ming") # True
15 s.find("gram") # 3 (index, -1 if not found)
16 s.replace("ram", "ROM") # "progROMming"
17
18 # Character operations
19 char = 'A'
20 ord(char) - ord('A') # 0 (offset from 'A')
21 chr(ord('a') + 1) # 'b'
22
23 # Regular expressions (when needed)
24 import re
25 pattern = r'\d+' # One or more digits
26 matches = re.findall(pattern, "abc123def456") # ['123', '456']
27

```

## 19.6 Control Flow and Error Handling

```

1 # Try-except for parsing
2 def safe_int(s):
3     try:
4         return int(s)
5     except ValueError:
6         return 0
7
8 # Else clause with loops
9 for item in items:
10     if condition(item):
11         break
12 else:
13     # Executed if loop completed without break
14     print("No item found")
15
16 # While-else similar pattern
17 while condition:
18     if found:
19         break
20 else:
21     print("Condition became false, not broken")
22
23 # Ternary operator (conditional expression)
24 result = "positive" if x > 0 else "non-positive"
25 max_val = a if a > b else b
26
27 # Chained comparisons
28 if 0 <= x < len(arr): # Bounds check
29     if a < b < c: # Three-way comparison
30
31 # Walrus operator (Python 3.8+) - assignment in expression
32 while (line := input()) != "END":
33     process(line)
34
35 if (n := len(items)) > 10:
36     print(f"Many items: {n}")

```

## 19.7 Advanced Input/Output Patterns

```

1 # Multiple assignment from input
2 a, b, c = map(int, input().split())
3 arr = list(map(int, input().split()))
4
5 # Reading matrix
6 n, m = map(int, input().split())
7 matrix = []
8 for _ in range(n):
9     row = list(map(int, input().split()))
10    matrix.append(row)
11
12 # Or with comprehension
13 matrix = [list(map(int, input().split())) for _ in range(n)]
14
15 # Reading until EOF
16 import sys
17 lines = sys.stdin.read().strip().split('\n')
18
19 # Fast I/O: see the dedicated Input/Output section for recommended setup
20 using
21 # 'input' = sys.stdin.readline' and other fast I/O patterns (Section 01_io)
22
23 # Avoid redefining 'input'/'print' repeatedly across sections; set them
24 once
25 # in your submission template when needed.
26
27 # Multiple test cases pattern
28 t = int(input())
29 for _ in range(t):
30     # Process each test case
31     pass

```

```

29 # Output formatting
30 print(*arr)           # Space-separated
31 print(*arr, sep='\n')  # Newline-separated
32 print(f"{'x':.6f}")   # 6 decimal places

```

## 19.8 Type conversions and casting

```

1 # Common conversions
2 s = str(123)          # '123'
3 i = int('42')         # 42
4 f = float('3.14')     # 3.14
5 b = bool(0)           # False
6
7 # Container conversions
8 lst = list([1,2])     # [1,2]
9 tup = tuple([1,2])    # (1,2)
10 st = set([1,2,2])    # {1,2}
11 d = dict([('a',1), ('b',2)]) # {'a':1, 'b':2}
12
13 # Characters and bytes
14 c = chr(65)           # 'A'
15 o = ord('A')          # 65
16 by = bytes('abc', 'utf-8') # b'abc'
17 s2 = by.decode('utf-8') # 'abc'
18
19 # Number formatting to hex/bin/oct
20 h = hex(255)           # '0xff'
21 bstr = bin(10)          # '0b1010'
22 octr = oct(8)           # '0o10'
23 # Formatting without prefixes
24 h2 = format(255, 'x')  # 'ff'
25 b2 = format(10, 'b')   # '1010'
26
27 # Safe literal evaluation (use instead of eval when parsing literals)
28 import ast
29 data = ast.literal_eval("[1, 2, {'a':3}]")
30
31 # Common contest patterns
32 # Convert split input to ints
33 arr = list(map(int, input().split()))
34 # Join values back into a string
35 s = ' '.join(map(str, arr))

```

## ASCII conversions (quick reference)

Char	Code	Char	Code
'A'	65	'a'	97
'Z'	90	'z'	122
'0'	48	'9'	57
' '	(space)	'\n'	10
'!'	33	'?'	63

Use `ord(char)` to get the numeric code and `chr(code)` to convert back.

### Key Benefits:

- **Readability:** More concise and expressive code
- **Performance:** Built-in functions are often faster than manual loops
- **Memory:** Generators save memory for large datasets
- **Debugging:** Functional style often easier to debug
- **Contest speed:** Write solutions faster with fewer lines

## 20 Problem procedure and contest workflow

This section summarizes a concise, repeatable procedure to approach each contest problem efficiently.

1. **Quick read (1–2 minutes):** Read statement carefully. Identify input/output formats and required output exactly.
2. **Classify and estimate (1–3 minutes):** Determine problem type (graphs, math, DP, greedy, flow, geometry, data structures). Estimate difficulty and a rough solution approach and complexity.
3. **Decide whether to attempt now**
  - Solve immediately if it looks easy or you already know the pattern.
  - Skip and mark for later if it looks hard or needs long derivation.
4. **Concrete examples and edge cases (2–5 minutes):** Create small examples, corner cases, and verify understanding (empty inputs, single element, max/min values).
5. **Design outline (5–15 minutes):** Write a clear plan, data structures, and algorithm steps. Record complexity and memory. If proof/argument needed, sketch correctness and invariants.
6. **Implement carefully**
  - Start from a short, tested template (fast I/O, common helpers).
  - Implement core logic in readable chunks and name variables clearly.
  - Avoid clever one-liners that are hard to debug under time pressure.
7. **Test locally (2–8 minutes):** Run sample tests, your handcrafted examples, and stress small random cases if possible. Check edge cases (overflows, boundary indices, integer division, floating precision, modulo behavior).
8. **Optimize only if necessary**
  - If complexity is borderline, identify hotspots and optimize (use faster data structures, avoid repeated work, precompute).
  - Prefer simple, robust fixes (change algorithmic approach if needed).
9. **Verify and submit**
  - Re-check input parsing and exact output formatting.
  - Comment out debug prints and run final tests.
  - Submit when confident; if WA, read feedback and debug quickly (reproduce failing case locally if possible).
10. **Time management and team strategy**
  - Track time per problem; don't spend excessive time on a single problem early.
  - Coordinate with teammates (divide/parallelize work, let one implement while another tests or reads other problems).
  - Keep a visible board of which problems are attempted, solved, or to revisit.

## Quick checklist before submitting

- Correct handling of limits and types (int vs float, long long, big integers).
- Off-by-one and index orientation.
- Exact output formatting (extra spaces/newlines).
- Performance within time/memory limits.
- No remaining debug prints.

## Common patterns and hints

- Try greedy first for constructive problems; if it fails, look for DP/graph framing.
- Reduce to shortest paths / flow for routing and matching problems.
- For counting/number theory, consider modular arithmetic and combinatorics shortcuts.
- Use Union-Find for connectivity under merges; use binary search on answer when monotone.
- When stuck, simplify constraints or solve a special case to gain insight.