

MARKOV PROCESS

$(\Omega, \mathcal{F}, \mathbb{P})$ probability space
 (E, \mathcal{E}) measurable space (E typically is finite/countable / $\mathbb{R}/\mathbb{R}^n/A \subseteq \mathbb{R}^n$)

Def. STOCHASTIC PROCESS: A stochastic process with values in E is a collection of random variables $X_t: \Omega \rightarrow E$ and it's denoted by $(X_t)_t$

Def. MARKOV PROCESS: $(X_t)_t$ is a Markov process if $(X_t)_t$ is a stochastic process with Markov property:

$$\mathbb{P}(X_{t_{m+1}} \in E_{m+1} | X_{t_m} \in E_m, \dots, X_{t_1} \in E_1) = \mathbb{P}(X_{t_{m+1}} \in E_{m+1} | X_{t_m} \in E_m)$$

$$\forall t_1 < t_2 < \dots < t_m < t_{m+1}, \quad \forall E_1, \dots, E_{m+1} \in \mathcal{E}, \quad \forall m \in \mathbb{N}$$

$t_{m+1} = \text{future}$

$t_m = \text{present}$

$t_1, \dots, t_{m-1} = \text{past}$

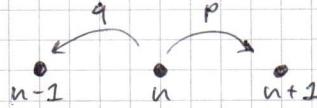
Example: Discrete Markov Chain

E finite/countable, $\mathcal{E} = \mathcal{P}(E)$, $t \in \mathbb{N} = \{0, 1, 2, \dots\}$

$$\begin{aligned} \text{Markov property} \Leftrightarrow \mathbb{P}(X_{t_{m+1}} = e_{m+1} | X_{t_m} = e_m, \dots, X_{t_1} = e_1) &= \\ &= \mathbb{P}(X_{t_{m+1}} = e_{m+1} | X_{t_m} = e_m) \end{aligned}$$

(with discrete Markov chains we can replace sets with elements of the set)

Example: Discrete time, \mathbb{Z} -valued Markov chain: Random walk on \mathbb{Z}



Starting from n , in a unit of time we jump to:

- $n+1$ with probability p
- $n-1$ with probability q

Defining the realization of $(Y_k)_{k \geq 1}$ independent random variables with values $\{-1, 1\}$, we can define the position at time n :

$$X_n = x_0 + \sum_{k=1}^n Y_k \quad x_0 \in \mathbb{Z}$$

$(X_n)_{n \geq 1}$ is a Markov chain.

(Discrete time $t \in \mathbb{N} = \{0, 1, 2, \dots\}$ ≠ discrete values (state space))

Def. TRANSITION PROBABILITIES:

$$p_{ij}(t, s) := \mathbb{P}(X_t = j | X_s = i) \quad s < t$$

Def. TIME HOMOGENEOUS PROCESS: The process $(X_t)_t$ is time homogeneous if $p_{ij}(t, s)$ depends only on $t-s$. In this case we can write $p_{ij}(t-s)$ or $p_{ij}^{(t-s)}$. For $t-s=1$ we simply write p_{ij} or $p_{ij}^{(1)}$ and we call $(p_{ij})_{i,j \in E}$ the TRANSITION MATRIX.

Lemma: $\mathbb{P}(X_{t+z} = j | X_t = i) = p_{ij}^{(z)}$

$$= \sum_{k \in E} p_{ik} p_{kj}$$

$\xrightarrow{p_{ij}^{(z)}}$
 $s \quad t$
power z of the transition matrix

proof.

$$\mathbb{P}(X_{t+z} = j | X_t = i) = \frac{\mathbb{P}(X_{t+z} = j, X_t = i)}{\mathbb{P}(X_t = i)}$$

$$\begin{aligned}
 \text{IP}(X_{t+2} = j \mid X_t = i) &= \frac{\sum_{k \in E} \text{IP}(X_{t+2} = j, X_{t+1} = k, X_t = i)}{\text{IP}(X_t = i)} \\
 &\stackrel{\text{Markov property}}{=} \sum_{k \in E} \text{IP}(X_{t+2} = j \mid X_{t+1} = k, X_t = i) \frac{\text{IP}(X_{t+1} = k, X_t = i)}{\text{IP}(X_t = i)} \\
 &= \sum_{k \in E} \text{IP}(X_{t+2} = j \mid X_{t+1} = k) \text{IP}(X_{t+1} = k \mid X_t = i) \\
 &= \sum_{k \in E} p_{kj} p_{ik} \leftarrow \text{elements of } (p_{ij}) \text{ power 2}
 \end{aligned}$$

More generally:

$$\text{IP}(X_{t_m} = j_m, X_{t_{m-1}} = j_{m-1}, \dots, X_0 = j_0) = \text{IP}(X_0 = j_0) p_{j_0 j_1}^{(t_1)} p_{j_1 j_2}^{(t_2 - t_1)} \dots p_{j_{m-1} j_m}^{(t_m - t_{m-1})}$$

with $(p_{ij})_{ij}$ transition matrix.

Properties:

- $0 \leq p_{ij} \leq 1$
- $\sum_{j \in E} p_{ij} = 1 \quad (\sum_{j \in E} \text{IP}(X_1 = j \mid X_0 = i) = 1)$

Suppose $(X_n)_{n \geq 0}$, $n \in \mathbb{N}$, E countable:

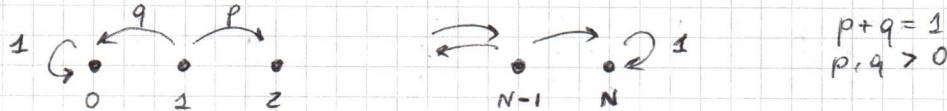
Def.: ACCESSIBLE STATE: $i, j \in E$, j accessible from i if $\exists n \geq 0$ s.t.:
 $\text{IP}(X_n = j \mid X_0 = i) = p_{ij}^{(n)} > 0$.

COMMUNICATIVE STATES: $i, j \in E$, j and i communicate if each one is accessible from the other one.

CLASS OF STATES: $C \subseteq E$ is a class of states if all states in C communicate and they do not communicate with states in $E \setminus C$.

Example: Gambler's ruin problem

States $E = \{0, 1, \dots, N\}$



Transition matrix:

$$\begin{aligned}
 p_{ij} : i \rightarrow j : & \quad \begin{matrix} 0 & 1 & 2 & \dots & N-1 & N \end{matrix} \quad \text{(going to)} \\
 & \begin{bmatrix} 0 & 1 & 2 & \dots & N-1 & N \\ 1 & 0 & 0 & 0 & 0 & 0 \\ 2 & q & 0 & p & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ N-1 & 0 & 0 & 0 & 0 & q \\ N & 0 & 0 & 0 & 0 & p \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & \dots & 0 \\ 0 & q & 0 & p & 0 \\ 0 & 0 & q & 0 & p \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \quad \text{(going from)} \\
 & \begin{matrix} 0 & 1 & 2 & \dots & N-1 & N \end{matrix} \\
 & \begin{bmatrix} p & q & 0 & \dots & 0 & 0 \\ q & p & 0 & \dots & 0 & 0 \\ 0 & 0 & p & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & p & q \\ 0 & 0 & 0 & \dots & q & p \end{bmatrix} \\
 & \begin{matrix} 0 & 1 & 2 & \dots & N-1 & N \end{matrix} \\
 & \begin{bmatrix} p & q & 0 & \dots & 0 & 0 \\ q & p & 0 & \dots & 0 & 0 \\ 0 & 0 & p & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & p & q \\ 0 & 0 & 0 & \dots & q & p \end{bmatrix}
 \end{aligned}$$

$\{0\}$ is a class of state since 0 communicates only with itself.

The classes are $\{0\}$, $\{N\}$, $\{1, 2, 3, \dots, N-1\}$

Example: Random queue at a counter:

- a counter processes 1 customer per unit of time
- the number of customers arriving in the system per unit of time is a random variable with density $(a_n)_{n \geq 0}$:

A_n = number of arrivals at time n : $A_1, \dots, A_n \perp \!\!\! \perp$

$\text{IP}(A_n = k) = a_k$ = at the time n there comes A_n people with probability a_k

X_n = number of customers in the queue at time n

$$X_{n+1} = \begin{cases} + A_{n+1} - 1 & \text{if } X_n > 0 \\ A_{n+1} & \text{if } X_n = 0 \end{cases}$$

$$= (X_n - 1)^+ + A_{n+1}$$

it appears
only if $X_n > 0$

Transition matrix:

$$\begin{matrix} & 0 & 1 & 2 & \dots \\ 0 & a_0 & a_1 & a_2 & a_3 & \dots \\ 1 & a_0 & a_1 & a_2 & a_3 & \dots \\ 2 & 0 & a_0 & a_1 & a_2 & \dots \\ \vdots & 0 & 0 & a_0 & a_1 & \dots \\ \vdots & \vdots & \vdots & \vdots & \ddots \end{matrix}$$

If a_k are sufficiently "generic", namely $a_0 > 0$, $a_1 > 0$, and $\exists \bar{n} \geq 2$ such that $a_{\bar{n}} > 0$, what are the classes of the states?

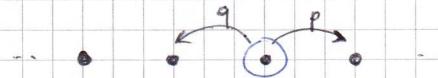
Starting from $i > 0$ one can reach 0 in i steps with probability $a_0^i > 0$.

Starting from $i > 0$ one can reach any state $j + (\bar{n}-1)n$ and then only $j > i$ for $i < j < 1 + (\bar{n}-1)n$

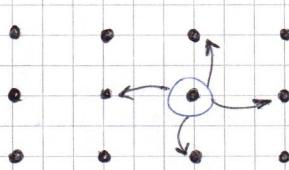
Def. **IRREDUCIBLE MC**: A Markov chain is irreducible if all states communicate.

(Examples:)

- Random walk on \mathbb{Z} :



- Random walk on \mathbb{Z}^2 :



in this case we need a 4 dimensional matrix for the representation

Def. **RECURRENT STATE**: $(X_n)_{n \geq 0}$ discrete time Markov chain. A state is called recurrent if:

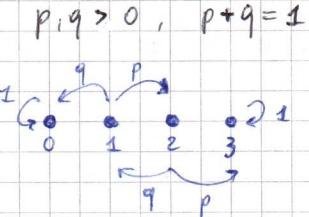
$$\Pr \left(\bigcup_{1 \leq n \leq \infty} \{X_n = i\} \mid X_0 = i \right) = 1$$

Namely, the probability of return to i in finite time (starting from i) is 1. We're sure to come back to i . Otherwise the state i is called TRANSIENT.

Example: Gambler's ruin with $N = 3$:

Transition matrix:

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ q & 0 & p & 0 \\ 0 & q & 0 & p \\ 0 & 0 & 0 & 1 \end{bmatrix}$$



0, 3 are recurrent (absorbing states)
1, 2 are transient

Def. **FIRST ENTRANCE TIME**: Let i be a state.

$$T_i := \begin{cases} \min \{n \mid X_n = i\} & \text{if } \{n \geq 1 \mid X_n = i\} \neq \emptyset \\ +\infty & \text{if } \{n \geq 1 \mid X_n = i\} = \emptyset \end{cases}$$

T_i is called first entrance time (or first visit time) to i .

Notation: $f_{ji}^{(n)} = \Pr(T_i = n \mid X_0 = j)$

Prop. Renewal equations

$$\forall i, j \in E : p_{ij}^{(n)} = \underbrace{\sum_{j=1}^n f_{ij}^{(n)}}_{\substack{\text{probability of} \\ \text{going from } i \text{ to } j \\ \text{in } n \text{ steps}}} \underbrace{p_{jj}^{(n-j)}}_{\substack{\sum \text{ outer in } j \text{ after } j \\ \text{steps and then remain} \\ \text{in } j \text{ (the } \sum \text{ represents} \\ \text{all the possible cases)}}}$$

proof.

$$p_{ij}^{(n)} = \mathbb{P}(X_n = j \mid X_0 = i)$$

$$\{X_n = j\} = \bigcup_{j=1}^n \{X_n = j, T_j = \infty\} = \bigcup_{j=1}^n \{X_n = j, \dots, X_{j-1} \neq j, \dots\}$$

disjoint events

$$\begin{aligned} p_{ij}^{(n)} &= \sum_{j=1}^n \mathbb{P}(X_n = j, \dots, X_{j-1} \neq j, \dots, X_0 = i) \\ &= \sum_{j=1}^n \mathbb{P}(X_n = j, \dots \mid X_{j-1} = j, X_{j-1} \neq j, \dots) \mathbb{P}(X_{j-1} = j, \dots, X_0 = i) \\ &= \sum_{j=1}^n \mathbb{P}(X_n = j, \dots \mid X_{j-1} = j) \mathbb{P}(X_{j-1} = j, \dots, X_0 = i) \\ &= \sum_{j=1}^n p_{jj}^{(n-j)} \mathbb{P}(T_j = \infty \mid X_0 = i) \\ &= \sum_{j=1}^n p_{jj}^{(n-j)} f_{ij}^{(n)} \end{aligned}$$

Theorem: For any state i of a discrete Markov chain $(X_n)_n$ the following are equivalent:

(a) i is recurrent

(b) $\sum_{n \geq 0} p_{ii}^{(n)} = +\infty$ = the average time spent in i is $+\infty$

If the series in (b) is convergent ($\sum_{n \geq 0} p_{ii}^{(n)} < +\infty$) then:

$$\sum_{n \geq 0} p_{ii}^{(n)} = \frac{1}{1 - \mathbb{P}_i(T_i < \infty)}$$

where $\mathbb{P}_i(\bullet) = \mathbb{P}(\bullet \mid X_0 = i)$

= probability of never returning in i

Meaning of (b):

$$p_{ii}^{(n)} = \mathbb{P}(X_n = i \mid X_0 = i) = \underbrace{\mathbb{E}[\mathbb{1}_{\{X_n = i\}}]}_{\substack{\text{expectation with respect} \\ \text{to } \mathbb{P}(\bullet \mid X_0 = i)}}$$

$$\Rightarrow \sum_{n \geq 0} p_{ii}^{(n)} = \sum_{n \geq 0} \mathbb{E}[\mathbb{1}_{\{X_n = i\}}] = \mathbb{E} \left[\sum_{n \geq 0} \mathbb{1}_{\{X_n = i\}} \right]$$

$\underbrace{\text{total time spent}}_{\text{in } i} \underbrace{\text{average time}}_{\text{spent in } i}$

$\Rightarrow [i \text{ recurrent} \Leftrightarrow (\text{a}): \text{for sure we'll be back to } i \Leftrightarrow (\text{b}): \text{the (average)time spent in } i \text{ is } +\infty]$

proof.

Recall: $f_{ii}^{(n)} = \mathbb{P}_i(T_i = n) = \text{probability of first return to } i \text{ after } n \text{ steps}$

We define: $P_{ii}(s) = \sum_{n \geq 0} s^n p_{ii}^{(n)}$ } power series, absolutely convergent if $|s| < 1$
 $F_{ii}(s) = \sum_{n \geq 1} s^n f_{ii}^{(n)}$

Renewal equations: $p_{ii}^{(n)} = \sum_{k=1}^n f_{ii}^{(k)} p_{ii}^{(n-k)}$

 $\Rightarrow s^n p_{ii}^{(n)} = \sum_{k=1}^n (s^k f_{ii}^{(k)}) (s^{n-k} p_{ii}^{(n-k)})$
 $\Rightarrow \sum_{n=1}^{\infty} s^n p_{ii}^{(n)} = \sum_{n=1}^{\infty} \sum_{k=1}^n (s^k f_{ii}^{(k)}) (s^{n-k} p_{ii}^{(n-k)})$
 $= \sum_{k=1}^{\infty} \sum_{n=k}^{\infty} (s^k f_{ii}^{(k)}) (s^{n-k} p_{ii}^{(n-k)})$
 $= \sum_{k=1}^{\infty} (s^k f_{ii}^{(k)}) \sum_{n=k}^{\infty} (s^{n-k} p_{ii}^{(n-k)})$
 $= \sum_{k=1}^{\infty} (s^k f_{ii}^{(k)}) \sum_{m=0}^{\infty} (s^m p_{ii}^{(m)})$
 $= F_{ii}(s) P_{ii}(s)$

$\Rightarrow P_{ii}(s) - 1 = F_{ii}(s) P_{ii}(s)$

$\Rightarrow \text{For } |s| < 1 : P_{ii}(s) = \frac{1}{1 - F_{ii}(s)} \quad (*)$

(a) \Rightarrow (b): i recurrent means $\mathbb{P}_i(\tau_i < \infty) = 1$, namely:

$\sum_{n=1}^{\infty} f_{ii}^{(n)} = 1 \quad (f_{ii}^{(n)} = \mathbb{P}_i(\tau_i = n))$

$\Rightarrow \lim_{s \rightarrow 1^-} F_{ii}(s) = \lim_{s \rightarrow 1^-} \sum_{n=1}^{\infty} s^n f_{ii}^{(n)} = 1$

$$\underbrace{\lim_{s \rightarrow 1^-} P_{ii}(s)}_{\sum_{n=0}^{\infty} p_{ii}^{(n)}} = \lim_{s \rightarrow 1^-} \frac{1}{1 - F_{ii}(s)} = +\infty$$

(b) \Rightarrow (a): (b) means that $\lim_{s \rightarrow 1^-} P_{ii}(s) = +\infty$

From (*) : $F_{ii}(s) = 1 - \frac{1}{P_{ii}(s)}$, then:

$\lim_{s \rightarrow 1^-} F_{ii}(s) = 1 \quad (\text{i.e. } \sum_{n=1}^{\infty} f_{ii}^{(n)} = 1) \Leftrightarrow \mathbb{P}_i(\tau_i < \infty) = 1$

If the series is convergent: $\sum_{n=0}^{\infty} p_{ii}^{(n)} < \infty$:

$P_{ii}(s) = \frac{1}{1 - F_{ii}(s)} \xrightarrow{s \rightarrow 1^-} \sum_{n=0}^{\infty} p_{ii}^{(n)} = \frac{1}{1 - \mathbb{P}_i(\tau_i < \infty)}$

Remark: $\sum_{n \geq 0} p_{ii}^{(n)} = \mathbb{E}_i \left[\underbrace{\sum_{n \geq 0} \mathbb{1}_{\{X_n=i\}}}_{\text{number of visits in } i} \right]$

If i is transient both of the terms are $< \infty$, and so:

$i \text{ transient} \Rightarrow \mathbb{P} \left(\sum_{n \geq 0} \mathbb{1}_{\{X_n=i\}} = +\infty \right) = 0$

Conclusion: or MC visits a transient state only a finite number of times with probability = 1

Corollary: Two communicating states i, j are both recurrent or both transient
 (This implies that: let J be a class of states, then if $j \in J$ is recurrent (or transient) all the elements of J are recurrent (or transient))

proof.

Suppose that i is recurrent $\Rightarrow \sum_{n \geq 0} p_{ii}^{(n)} = +\infty$

since i and j communicate $\exists s, t \in \mathbb{N} : p_{ji}^{(s)} > 0, p_{ij}^{(t)} > 0$

$\Rightarrow p_{jj}^{(s+t+n)} = \sum_{k_1, k_2} p_{jk_1}^{(s)} p_{k_1 k_2}^{(n)} p_{k_2 j}^{(t)} \geq p_{ji}^{(s)} p_{ii}^{(n)} p_{ij}^{(t)}$

$$\Rightarrow \sum_{n \geq 0} p_{jj}^{(s+t+n)} \geq p_{ji}^{(s)} \left(\underbrace{\sum_{n \geq 0} p_{ii}^{(n)}}_{> 0} \right) \underbrace{p_{ij}^{(t)}}_{> 0} = +\infty$$

since i is recurrent

$$\Rightarrow \sum_{n \geq 0} p_{jj}^{(s+t+n)} = +\infty \Rightarrow j \text{ is recurrent.}$$

This is enough since every state is either recurrent or transient. ■

Corollary:

1. If j is transient $\Rightarrow \forall i \sum_{n \geq 0} p_{ij}^{(n)} < +\infty$
2. In particular (j transient): $\lim_{n \rightarrow \infty} P_i(X_n = j) = 0$

3. If the set of states E is finite then there exists at least a recurrent state (this is not necessary true for a transient state).
(Notice: E is the set of states, not class of states)

if j is transient, starting from any i the Markov Chain goes through j almost certainly only a finite number of times

proof.

$$\begin{aligned} 1. \text{ Renewal equation: } \sum_{n \geq 1} p_{ij}^{(n)} &= \sum_{n \geq 1} \sum_{k=1}^n f_{ij}^{(k)} p_{jj}^{(n-k)} \\ &= \sum_{k=1}^{\infty} \sum_{n=k}^{\infty} f_{ij}^{(k)} p_{jj}^{(n-k)} \\ &= \sum_{k=1}^{\infty} f_{ij}^{(k)} \sum_{n=k}^{\infty} p_{jj}^{(n-k)} \\ &= \sum_{k=1}^{\infty} f_{ij}^{(k)} \underbrace{\sum_{m=0}^{\infty} p_{jj}^{(m)}}_{P_i(T_j < \infty)} \underbrace{\text{convergent } (< \infty)}_{(\text{since } j \text{ is transient})} \end{aligned}$$

$$\Rightarrow \sum_{n \geq 1} p_{ij}^{(n)} < \infty$$

$$2. \sum_{n \geq 1} p_{ij}^{(n)} < \infty \Rightarrow \lim_{n \rightarrow \infty} p_{ij}^{(n)} = 0 \quad \equiv \text{the probability of visiting } j \text{ goes to 0 as } n \rightarrow \infty$$

$$3. E \text{ finite. Suppose that all the states are transient} \Rightarrow \lim_{n \rightarrow \infty} p_{ij}^{(n)} = 0 \quad \forall i, j$$

$$\text{but } \forall n : \sum_{j \in E} p_{ij}^{(n)} = 1$$

$$\Rightarrow (\text{E finite}) \quad n \rightarrow \infty : \quad 0 = \lim_{n \rightarrow \infty} \sum_{j \in E} p_{ij}^{(n)} = 1, \text{ contradiction}$$

$\Rightarrow \exists$ at least 1 recurrent state. ■

Example: Random walk on \mathbb{Z} :



$$p, q > 0, \quad p + q = 1$$

$$X_{n+1} = X_n + Y_{n+1} \quad \text{where } (Y_n)_{n \in \mathbb{N}} \text{ are iid:} \quad \begin{aligned} P(Y_n = 1) &= p \\ P(Y_n = -1) &= q \end{aligned}$$

$$X_n = x_0 + \sum_{k=1}^n Y_k \quad \left(\frac{1}{n} \sum_{k=1}^n Y_k \xrightarrow{\text{a.s.}} \mathbb{E}[Y_1] = p - q \right)$$

If we start from 0 we can find a moment (a time) for which after that moment we can't go back to 0 $\Rightarrow 0$ is transient.

LLN \rightarrow all states are transient if $p \neq q$
Law of large numbers

Application of
the theorem
(recurrent and
transient
states)

The MC is irreducible since all the states communicate.
 \Rightarrow all the states are either transient or recurrent)

Let's consider the state 0 and study the series:

$$\sum_{n \geq 0} p_{00}^{(n)} = \begin{cases} +\infty \\ < \infty \end{cases}$$

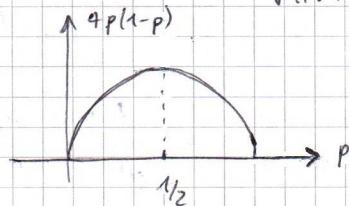
The probability of starting from 0 and come back to 0 after n steps:

$$p_{00}^{(n)} = \begin{cases} 0 & n = \text{odd} \\ \binom{2m}{m} p^m q^m & n = 2m \end{cases}$$

$$p_{00}^{(2m)} = \frac{(2m)!}{(m!)^2} (p(1-p))^m \sim \frac{(2m)^{2m} e^{-2m}}{m^{2m} e^{-2m} 2^{2m}} \frac{\sqrt{4\pi m}}{(p(1-p))^m}$$

$$\text{Stirling: } \lim_{n \rightarrow \infty} \frac{n^n e^{-n} \sqrt{2\pi n}}{n!} = 1$$

$$\Rightarrow p_{00}^{(2m)} \sim \frac{2^{2m}}{\sqrt{\pi m}} (p(1-p))^m = \frac{(4p(1-p))^m}{\sqrt{\pi m}}$$



$$p_{00}^{(2m)} = \begin{cases} \frac{1}{\sqrt{\pi m}} & p = \frac{1}{2} \quad (p = q) \\ \frac{(4p(1-p))^m}{\sqrt{\pi m}} & p \neq \frac{1}{2} \quad (p \neq q) \end{cases}$$

$$\Rightarrow \begin{cases} \bullet \quad p = q = \frac{1}{2} : \quad \sum_{n \geq 0} p_{00}^{(n)} = +\infty \quad (\text{state is recurrent}) \\ \bullet \quad p \neq q \left(\neq \frac{1}{2}\right) : \quad \sum_{n \geq 0} p_{00}^{(n)} < +\infty \quad (\text{state is transient}) \end{cases}$$

Moreover: (for $p=q=\frac{1}{2}$)

$$p_{00}^{(2m)} = \binom{2m}{m} (pq)^m \Rightarrow P_{00}(s) = \sum_{n \geq 0} s^n p_{00}^{(n)}$$

$$\Rightarrow P_{00}(s) = \sum_{n \geq 0} s^{2m} p_{00}^{(2m)}$$

$$= \sum_{n \geq 0} s^{2m} \binom{2m}{m} (pq)^m$$

$$= \frac{1}{\sqrt{1 - 4pq s^2}} \quad |s| < 1$$

$$\Rightarrow F_{00}(s) = 1 - \frac{1}{P_{00}(s)} = 1 - \sqrt{1 - 4pq s^2} = 1 - \sqrt{1 - s^2}$$

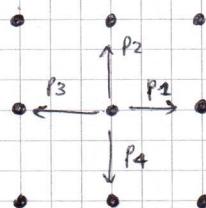
$$\Rightarrow F_{00}(s) = \sum_{n \geq 1} s^n P_0(T_0=n) = 1 - \sqrt{1 - s^2}$$

$$\frac{\partial}{\partial s} \Rightarrow \sum_{n \geq 1} n s^{n-1} P_0(T_0=n) = \frac{s}{(1-s^2)^{3/2}}$$

$$\lim_{s \rightarrow 1^-} \Rightarrow \sum_{n \geq 1} n P_0(T_0=n) = E_0[T_0] = +\infty$$

so even if the return to 0 is almost certain,
 the return time is ∞

Example: Random walk on \mathbb{Z}^2 :



$$p_1 + p_2 + p_3 + p_4 = 1$$

- p_i all equal $\Rightarrow p_i = \frac{1}{4}$ $\forall i$, all states are recurrent
- p_i all different \Rightarrow all states are transient

$$p_{00}^{(2m)} = \sum_{h+k=m} \frac{(2m)!}{(h!)^2 (k!)^2} \left(\frac{1}{4}\right)^{2m} \quad \leftrightarrow \quad \begin{matrix} h \\ k \end{matrix}$$

$$\begin{aligned}
 p_{00}^{(2n)} &= \sum_{k=0}^m \frac{(2n)!}{(k!)^2 ((m-k)!)^2} \left(\frac{1}{4}\right)^{2n} \\
 &= \binom{2n}{m} \left(\frac{1}{4}\right)^{2n} \sum_{k=0}^m \binom{m}{k} \binom{m}{m-k} \\
 &\stackrel{?}{=} \binom{2n}{m} \left(\frac{1}{4}\right)^{2n} \binom{2n}{m} \\
 &\stackrel{?}{=} \left(\binom{2n}{m} \left(\frac{1}{4}\right)^m\right)^2 \sim \left(\frac{1}{\sqrt{\pi m}}\right)^2 = \frac{1}{\pi m}
 \end{aligned}$$

$\Rightarrow \sum_{n \geq 0} p_{00}^{(n)} = +\infty \Rightarrow (0,0)$ is recurrent (and so all the states are)

Example: Random walk on \mathbb{Z}^d :

$$P(X_{n+1} = (j_1, \dots, j_d) | X_n = (i_1, \dots, i_d)) = \begin{cases} \frac{1}{2d} & \text{if } \sum_{k=1}^d |i_k - j_k| = 1 \\ 0 & \text{otherwise} \end{cases}$$

Considering the MC irreducible: are the states transient or recurrent?

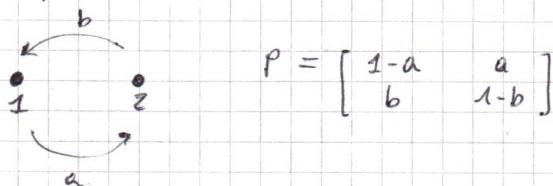
$$\begin{aligned}
 P(X_{2n} = 0 | X_0 = 0) &= p_{00}^{(2n)} \\
 &= \sum_{k_1 + \dots + k_d = n} \frac{(2n)!}{(k_1!)^2 \dots (k_d!)^2} \left(\frac{1}{2d}\right)^{2n} \\
 &= \binom{2n}{n} \left(\frac{1}{2d}\right)^{2n} \sum_{k_1 + \dots + k_d = n} \frac{(n!)^2}{(k_1!)^2 \dots (k_d!)^2} \\
 &\leq \max_{k_1 + \dots + k_d = n} \left\{ \frac{n!}{k_1! \dots k_d!} \right\} \cdot \boxed{\sum_{k_1 + \dots + k_d = n} \frac{n!}{k_1! \dots k_d!}} \\
 &= \max_{k_1 + \dots + k_d = n} \left\{ \frac{n!}{k_1! \dots k_d!} \right\} \cdot d^n
 \end{aligned}$$

$$(a_1 + \dots + a_d)^n = \sum_{\sum k_j = n} \frac{n!}{(k_1!) \dots (k_d!)} (a_1^{k_1} \dots a_d^{k_d})$$

$$\begin{aligned}
 \Rightarrow p_{00}^{(2n)} &\leq d^n \max_{k_1 + \dots + k_d = n} \left\{ \frac{n!}{k_1! \dots k_d!} \right\} \cdot \binom{2n}{n} \left(\frac{1}{2d}\right)^{2n} = \\
 &= d^n \frac{n!}{\left(\frac{n}{d}\right)!} \binom{2n}{n} \left(\frac{1}{2d}\right)^{2n} \\
 &\sim \frac{d^{d/2}}{\pi^{d/2} 2^{(d-1)/2} n^{d/2}} = \frac{\text{constant}(d)}{n^{d/2}}
 \end{aligned}$$

\Rightarrow if $d \geq 3$ $\sum_{n \geq 0} p_{00}^{(n)} < +\infty \Rightarrow$ all the states are transient

We want now to focus on the invariant distribution.
for example:



We would like to know how much we stay in a state. The invariant distribution answers to this problem:

$$P(X_n = 1) \stackrel{n \rightarrow \infty}{\approx} \pi_1 \text{ of an invariant distribution}$$

↑
asymptotical distribution

Def. INARIANT DISTRIBUTION: $(X_n)_{n \geq 0}$ MC with transition matrix $(p_{ij})_{i,j}$ and $\pi = (\pi_i)_{i \in E}$ probability density of E ($0 \leq \pi_i \leq 1$, $\sum_i \pi_i = 1$). π is an invariant density if $X_n \sim \pi \forall n \geq 0$.

We can write it as:

π is an invariant density if $X_{n+1} \sim \pi$ whenever $X_n \sim \pi \forall n$.

$$\text{Prop. } \text{IP}(X_{n+1} = j) = \sum_{i \in E} \text{IP}(X_{n+1} = j | X_n = i) \text{IP}(X_n = i)$$

$$\text{If } \text{IP}(X_n = i) = \pi_i \implies \text{IP}(X_{n+1} = j) = \sum_{i \in E} p_{ij} \pi_i$$

$$\implies \text{If } X_n \sim \pi, X_{n+1} \sim \pi \iff \pi_j = \sum_{i \in E} p_{ij} \pi_i = \text{of } P \text{ with eigenvalue } 1$$

$$\text{Example: } P = \begin{bmatrix} 1-a & a \\ b & 1-b \end{bmatrix}$$

$$\pi P = \pi \iff \begin{cases} \pi_1(1-a) + \pi_2 b = \pi_1 \\ \pi_1 a + (1-b)\pi_2 = \pi_2 \end{cases} \iff \pi_1 a = \pi_2 b$$

$$\implies \begin{cases} \pi_1 + \pi_2 = 1 \\ \pi_1 a = \pi_2 b \end{cases} \iff \pi = \left[\frac{b}{a+b}, \frac{a}{a+b} \right]$$

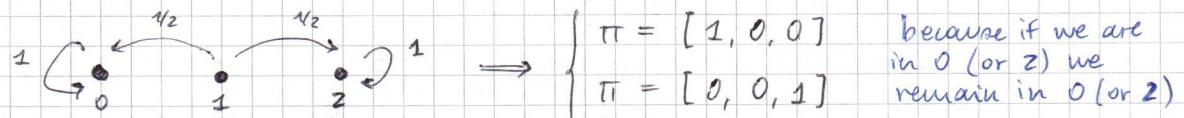
$$\implies \lim_{n \rightarrow \infty} \text{IP}(X_n = 1) = \pi_1 = \frac{b}{a+b} \quad \text{if we know the invariant distribution we can quantity the frequency of arriving}$$

Theorem: IF E is finite then there exists at least one invariant density.

Remark: The invariant is not necessarily unique:

$$P = \begin{bmatrix} 1 & 0 & 0 \\ 1/2 & 0 & 1/2 \\ 0 & 0 & 1 \end{bmatrix} \quad (\text{Gambler's ruin with } p=q=\frac{1}{2}, N=2)$$

$E = \{0, 1, 2\}$: there are two absorbing states: $\{0\}, \{2\}$



$$\text{IP}(X_n = 0) = 1 \implies \pi P = [1, 0, 0]P$$

$$\text{IP}(X_n = 2) = 1 \implies \pi P = [0, 0, 1]P$$

Convex combination:

$$\lambda [1, 0, 0] + (1-\lambda) [0, 0, 1] = [\underbrace{\lambda, 0, 1-\lambda}_{\text{they're all invariant distributions}}] \quad \lambda \in [0, 1]$$

$$\begin{cases} u = \mu P \\ v = \nu P \end{cases} \implies \lambda u + (1-\lambda)v = (\lambda u + (1-\lambda)v)P$$

Remark: Infinite MC may not have invariant distributions

$$\text{Random walk on } \mathbb{Z}: \text{IP}(X_n = 0) = \frac{1}{\sqrt{n}} \rightarrow 0$$

Rigorous computation:

$$\pi P = \pi : \quad p_{ij} = \begin{cases} 1/2 & j = i \pm 1 \\ 0 & \text{otherwise} \end{cases}$$

$$\implies \pi_i = \frac{1}{2} \pi_{i-1} + \frac{1}{2} \pi_{i+1}$$

$$\implies \pi_{i+1} - \pi_i = \pi_i - \pi_{i-1}$$

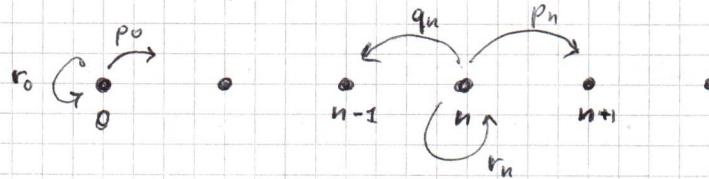
$$\implies \pi_i = \pi_0 + \sum_{k=1}^i (\pi_k - \pi_{k-1}) = \pi_0 + (\pi_1 - \pi_0)i$$

• $\pi_1 \neq \pi_0 \implies \pi_i$ can be negative for $i \rightarrow -\infty \implies$ no solutions 5

• $\pi_1 = \pi_0 \implies \pi_i = \pi_0 \quad \forall i \implies \sum_{i \in \mathbb{Z}} \pi_i = +\infty \implies$ no solutions

Example: Discrete birth and death process
 (only the name of the model, it actually has a lot of applications)

Status $N = \{0, 1, 2, \dots\}$



$$P = \begin{bmatrix} r_0 p_0 & \cdots & & \\ q_1 r_1 p_1 & 0 & 0 & 0 & \cdots \\ 0 & q_2 r_2 p_2 & 0 & 0 & \cdots \\ 0 & 0 & q_3 r_3 p_3 & 0 & \cdots \end{bmatrix}$$

$$\begin{aligned} r_i, q_i, p_i &\geq 0 \\ r_0 + p_0 &= 1 \\ r_i + p_i + q_i &= 1 \quad \forall i \neq 0 \end{aligned}$$

r_i = no change

p_i = P(birth)

q_i = P(death)

This MC is irreducible $\Rightarrow p_i > 0, q_i > 0$.

Invariant distribution π :

$$\begin{aligned} \pi_i = \pi P &\iff \begin{cases} \pi_0 = r_0 \pi_0 + q_1 \pi_1 \\ \pi_n = p_{n-1} \pi_{n-1} + r_n \pi_n + q_{n+1} \pi_{n+1} \end{cases} \\ \implies \pi_1 &= \frac{1 - r_0}{q_1} \pi_0 = \frac{p_0}{q_1} \pi_0 \end{aligned}$$

$$\pi_2 = \pi_1 \frac{1 - r_1}{q_2} - \frac{p_0}{q_2} \pi_0 = \frac{p_0 p_1}{q_1 q_2} \pi_0$$

:

$$\pi_n = \frac{p_0 p_1 \cdots p_{n-1}}{q_1 q_2 \cdots q_n} \pi_0$$

$(\pi_n)_{n \geq 0}$ are positive numbers if $\pi_0 > 0$.

$$\sum_{n \geq 0} \pi_n = 1 \iff \pi_0 + \pi_0 \sum_{n \geq 1} \frac{p_0 p_1 \cdots p_{n-1}}{q_1 q_2 \cdots q_n} = 1$$

If the series converges we can define: $Z := 1 + \sum_{n \geq 1} \frac{p_0 p_1 \cdots p_{n-1}}{q_1 q_2 \cdots q_n}$

$$\implies \pi_n = \frac{1}{Z} \cdot \frac{p_0 p_1 \cdots p_{n-1}}{q_1 q_2 \cdots q_n}$$

(unique)

Invariant distribution

If the series diverges there exists no invariant distribution.

$$\begin{aligned} \text{Remark: } \pi_n P_{n,n} &= \frac{1}{Z} \frac{p_0 \cdots p_{n-1}}{q_1 \cdots q_n} p_n \\ \pi_{n+1} P_{(n+1),n} &= \frac{1}{Z} \frac{p_0 \cdots p_n}{q_1 \cdots q_{n+1}} q_{n+1} \end{aligned} \quad \left. \begin{array}{l} \\ \end{array} \right\} \text{they're =}$$

Given π_n , if all the variables have the same invariant distribution, going on or going back is the same

Def. REVERSABLE PROBABILITY DENSITY: A probability density $(\pi_i)_{i \in E}$ is reversible with respect to a stochastic matrix $(p_{ij})_{i,j \in E}$ if:

$$\pi_i p_{ij} = \pi_j p_{ji} \quad \forall i, j \in E$$

Remark: Not every invariant density is a reversible density, but if a density is reversible then it's also invariant

reversible $(\pi_i)_i \iff (\pi_i)_i$ is invariant

proof.

$$\sum_{i \in E} \pi_i p_{ij} = \sum_{j \in E} \pi_j p_{ij} = \pi_i \sum_{j \in E} p_{ij} = \pi_i$$

(TP); by reversibility

Remark: If $p_{ij}^{(n)} \xrightarrow{n} \mu_j$ (converges) and $\lim_{n \rightarrow \infty} p_{ij}^{(n)} = \mu_j \quad \forall i$
 (so it depends only on j) and $\sum_{j \in E} \mu_j$ (which is ≤ 1) satisfies
 $\sum_{j \in E} \mu_j = 1 \implies \mu_j = \sum_i \mu_i p_{ij}$ is an invariant density
 (because $p_{ij}^{(n+1)} \rightarrow \mu_j \implies \sum_{k \in E} p_{ik}^{(n)} p_{kj} \xrightarrow{n \rightarrow \infty} \sum_{k \in E} \mu_k p_{kj}$)

Theorem: If limits $\lim_{n \rightarrow \infty} p_{ij}^{(n)}$ exist and are strictly positive and depend only on $j \implies$ MC has a unique invariant distribution π :

$$\lim_{n \rightarrow \infty} \mathbb{E}[f(X_n)] = \int f d\pi = \sum_{j \in E} f(j) \pi_j$$

for all f bounded.

In this situation: $X_n \rightarrow \pi$ in distribution ($P(X_n=j) \rightarrow \pi_j \quad \forall j$)

Def. PERIOD OF A STATE: The period of the state $i \in E$ is defined as:

$$\text{MCD} \{ n \geq 1 \mid p_{ii}^{(n)} > 0 \}$$

maximum common divisor

Ex. random walk: $p_{ii}^{(n)} = \begin{cases} 0 & n \text{ odd} \\ \neq 0 & n \text{ even} \end{cases} \implies \text{period} = 2$

Def. APERIODIC STATE: If the period is 1 the state is called aperiodic.

Prop: States of the same class have the same period.
 (for irreducible MC all states have the same period and we can talk about periodic/aperiodic Markov chains)

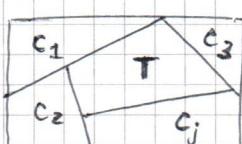
Theorem: Limits of subsequences drawn from $\frac{1}{n} \sum_{k=1}^n p_{ij}^{(k)} \xrightarrow{n \rightarrow \infty} L$
 determine invariant densities (also for periodic MC)

$$\frac{1}{n} \sum_{k=1}^n p_{ij}^{(k)} = \frac{1}{n} \sum_{k=1}^n \mathbb{E}_i [\mathbb{1}_{\{X_k=j\}}] = \mathbb{E}_i \left[\frac{1}{n} \sum_{k=1}^n \mathbb{1}_{\{X_k=j\}} \right]$$

mean frequency of visits in j frequency of visit in j
mean frequency of visit in j

$$\frac{1}{n} \sum_{k=1}^n \mathbb{1}_{\{X_k=j\}} \xrightarrow{\text{as.}} \pi_j$$

Reminder (Absorption probabilities / Hitting probabilities)



C_i = recurrent classes
 T = transient class

$$h_i = P_i(\cup_{n \geq 1} \{X_n \in C_i\})$$

= probability of absorption in C_1
 starting from a transient state

$$h_i = \sum_{j \in C_1} p_{ij} + \sum_{j \in T} p_{ij} h_j \quad (*)$$

the probabilities $(h_i)_{i \in T}$ are the smallest solution with values in $[0,1]$ of $(*)$. If T is finite $\Rightarrow (h_i)_{i \in T}$ are the only solutions with values in $[0,1]$ of $(*)$.

(Recall) Markov property for discrete MC:

$$\mathbb{P}(X_{n+1} = j \mid X_n = i_n, X_{n-1} = i_{n-1}, \dots) = \mathbb{P}(X_{n+1} = j \mid X_n = i_n) \quad \forall n, j, i_1, \dots, i_n$$

Notice that here we have fixed times ($n+1, n, n-1, \dots$).

Def. STOPPING TIME: The random variable $T: \Omega \rightarrow \mathbb{N}$ is a stopping time of the Markov chain $(X_n)_{n \geq 0}$ if $\forall n$ the event $\{T \leq n\}$ belongs to the σ -algebra generated by $\{X_n = i_n, X_{n-1} = i_{n-1}, \dots, X_0 = i_0\}$. (with i_1, \dots, i_0 arbitrary states in the countable set of states)

The stopping time T is \mathbb{U} from the future.

If (E, \mathcal{E}) and (F, \mathcal{F}) are measurable spaces and $f: E \rightarrow F$ is a function, then the σ -algebra generated by f (denoted by $f^{-1}(\mathcal{F})$) is the set of $f^{-1}(B) \quad B \in \mathcal{F}$.

The set of the events $\{X_n = i_n, X_{n-1} = i_{n-1}, \dots, X_0 = i_0\}$ represent the history of observations till time n .

Example: First entrance time T_j :

$$T_j = \inf_{n \geq 1} \{n \mid X_n = j\} \quad (\text{if } \{n \mid X_n = j\} = \emptyset \Rightarrow T_j = +\infty)$$

The time of the first visit in the state j is a stopping time, because:

$$\underbrace{\{T_j \leq n\}}_{\substack{\text{the first entrance in } j \\ \text{is at the step } n \text{ at most}}} = \underbrace{\bigcup_{m=1}^n \{X_m = j\}}_{\substack{\exists m \leq n \text{ for which} \\ X_m = j}}$$

$$\{T_j \leq n\} = \bigcup_{m=1}^n \{X_m = j\} = \bigcup_{m=1}^n \bigcup_{i_1, \dots, i_m} \{X_n = i_n, \dots, X_{m+1} = i_{m+1}, X_m = i_m, \dots\}$$

$$\{T_j = n\} = \{X_n = j, X_{n-1} \neq j, X_{n-2} \neq j, \dots, X_1 \neq j\}$$

Example: First exit time U_j :

$$U_j = \inf_{n \geq 1} \{n \mid X_n \neq j\} \quad (\text{if } \{n \mid X_n \neq j\} = \emptyset \Rightarrow U_j = +\infty)$$

$$\{U_j \leq n\} = \{U_j > n\}^c = \{X_n = j, \dots, X_1 = j\}^c$$

Prop. Law of the first exit time: $(X_n)_{n \geq 0}$ MC with transition matrix $(p_{ij})_{i,j}$. If $p_{jj} \in (0, 1)$ then the stopping time U_j has a geometric* density with parameter $1 - p_{jj}$ and the random variable X_{U_j} :

$$(X_{U_j})(\omega) = \sum_{n \geq 1} X_n(\omega) \mathbb{1}_{\{U_j=n\}}(\omega)$$

$$\text{satisfies: } \mathbb{P}_j(X_{U_j} = k) = \frac{p_{jk}}{1 - p_{jj}} \quad \forall k \neq j$$

* the Geometric distribution is the only one without memory, that's why it's correlated with Markov's property.

proof.

Geometric distribution: $\mathbb{P}(U_j = n) = ?$

$$\mathbb{P}(U_j > n) = \mathbb{P}(X_n = j, \dots, X_1 = j) = \sum_k p_{jj} p_{jj} \dots p_{jk} \mathbb{P}(X_0 = k)$$

$$\text{If we consider } \mathbb{P}(\cdot) = \mathbb{P}_j(\cdot) = \mathbb{P}(\cdot \mid X_0 = j) \Rightarrow \mathbb{P}_j(U_j > n) = (p_{jj})^n$$

$$\Rightarrow (n \geq 0): \mathbb{P}_j(U_j = n) = \mathbb{P}_j(U_j > n-1) - \mathbb{P}_j(U_j > n) \\ = (p_{jj})^{n-1} - (p_{jj})^n = (p_{jj})^{n-1} (1 - p_{jj})$$

So the parameter is $1 - p_{jj}$.

$$Y \sim \text{Geom}(\theta) \iff P(Y=n) = (1-\theta)^{n-1} \theta$$

Formula's proof:

$$\begin{aligned} P_j(X_{U_j}=k) &= \sum_{n=1}^{\infty} P_j(X_n=k, U_j=n) \\ &= \sum_{n=1}^{\infty} P_j(X_n=k, X_n \neq j, X_{n-1}=j, \dots, X_1=j) \end{aligned}$$

$$\text{If } k=j \implies P_j(X_n=k, X_n \neq j, X_{n-1}=j, \dots, X_1=j) = 0$$

$$\begin{aligned} \text{If } k \neq j \implies \sum_{n=1}^{\infty} P_j(X_n=k, X_n \neq j, X_{n-1}=j, \dots, X_1=j) &= \\ &= \sum_{n=1}^{\infty} \underbrace{p_{jj} p_{jj} \cdots p_{jj}}_{n-1 \text{ factors}} p_{jk} \end{aligned}$$

$$= p_{jk} \sum_{n \geq 1} p_{jj}^{n-1} = p_{jk} \frac{1}{1-p_{jj}}$$

Theorem (Restarted Markov chain): $(X_n)_{n \geq 0}$ MC, T stopping time (a.s. finite)

$$Y_n(w) := \begin{cases} X_{T(w)+n}(w) & \text{if } T(w) < +\infty \\ \text{arbitrary} & \text{if } T(w) = +\infty \end{cases}$$

$(Y_n)_{n \geq 0}$ is a Markov chain with the same transition matrix of the MC $(X_n)_{n \geq 0}$.

restart the process after the stopping time

proof.

We must check that Y_n is a random variable $\forall n$ and that $\forall n \forall j, i_1, \dots, i_n \in E$:

$$\begin{aligned} P(Y_{n+1}=j | Y_n=i_n, \dots, Y_1=i_1) &= P(Y_{n+1}=j | Y_n=i_n) && \text{(Markov property)} \\ &= p_{i_n j} && \text{(trans. matrix } p) \end{aligned}$$

namely:

$$\begin{aligned} P(X_{T+(n+1)}=j | X_{T+n}=i_n, \dots, X_{T+1}=i_1) &\stackrel{*}{=} P(X_{T+n+1}=i_{n+1} | X_{T+n}=i_n) = p_{i_n j} \\ &\stackrel{*}{=} \frac{P(X_{T+n+1}=j, X_{T+n}=i_n, \dots)}{P(X_{T+n}=i_n, \dots)} \\ &= \frac{\sum_{m=1}^{\infty} P(X_{m+n+1}=j, X_{m+n}=i_n, \dots, T=m)}{P(X_{T+n}=i_n, \dots)} \\ &= \frac{\sum_{m=1}^{\infty} P(X_{m+n+1}=j | X_{m+n}=i_n, \dots, T=m) P(X_{m+n}=i_n, \dots, T=m)}{P(X_{T+n}=i_n, \dots)} \\ &= \frac{\sum_{m=1}^{\infty} p_{i_n j} P(X_{m+n}=i_n, \dots, T=m)}{P(X_{T+n}=i_n, \dots, X_{T+1}=i_1)} \\ &= p_{i_n j} \end{aligned}$$

Same computation for $P(X_{T+n+1}=j | X_{T+n}=i_n) = p_{i_n j}$

Theorem (Strong Markov property): $(X_n)_{n \geq 0}$ MC, $(p_{ij})_{i,j \in E}$ transition matrix, T stopping time

$$(X_{T \wedge n})_{n \geq 0} \quad (T \wedge n)(w) = \min \{T(w), n\}$$

$$(Y_n)_{n \geq 0} = (X_{T+n})_{n \geq 0}$$

MC stopped at T
MC restarted at T

The stopped MC and restarted MC are $\perp \!\!\! \perp$ w.r.t. $P(\cdot | X_T=i, T < \infty)$ 7

proof. (Hint)

We have to check that all pairs of events:

$$\{Y_m = j_m, \dots, Y_0 = j_0\}, \quad \{X_{T \wedge n} = i_n, \dots, X_T = i_0\} \quad \forall m, n \quad \forall j_k, i_k$$

are independent.

" ■ "

Theorem: The number of visits of a recurrent state is infinite almost surely.
proof.

T
transience,
recurrency and
absorption

We define the number of visits in $i := N_i = \sum_{n \geq 1} \mathbb{1}_{\{X_n=i\}}$.

Times of visit: $\begin{cases} T_i^{(1)} = \inf_{n \geq 1} \{n \mid X_n = i\} & \text{time of the 1st visit} \\ \vdots \\ T_i^{(k+1)} = \inf_{n \geq k} T_i^{(k)} \{n \mid X_n = i\} & \text{time of the } (k+1)\text{th visit} \end{cases}$

$$i \text{ recurrent} \Leftrightarrow P_i(T_i^{(1)} < +\infty) = 1 \Leftrightarrow P_i(N_i \geq 1) = 1$$

(def.)

$T_i^{(2)}$ is the first visit time for the MC restarted from time $T^{(1)}$.
This MC has the same transition matrix

$$\Rightarrow P_i(T_i^{(2)} < \infty) = 1$$

$$\begin{aligned} \text{By induction: } & P_i(T_i^{(k)} < +\infty) = 1 \quad \forall k \\ \Leftrightarrow & P_i(N_i \geq k) = 1 \quad \forall k \\ \Leftrightarrow & P_i(N_i = +\infty) = 1 \end{aligned}$$

■

Theorem (Probability of staying forever in transient states):

Let T be the set of transient states.

The probability of remain forever in transient states is defined as:

$$U_i = P_i(\bigcap_{n=1}^{+\infty} \{X_n \in T\}) \quad \text{if } i \in T$$

Then, $(U_i)_{i \in T}$ is the biggest solution with $0 \leq U_i \leq 1 \quad \forall i \in T$ of the system of equations:

$$U_i = \sum_{j \in T} p_{ij} U_j$$

Remark: Generally there is no unique solution, but if T is finite then the only solution is $U_i = 0$.

proof. (Thm.)

We define $U_i^{(n)} = P_i(X_n \in T, \dots, X_1 \in T) = \text{probability of staying in transient states from 1 to } n \text{ (starting from 0)}$

$(U_i^{(n)})_{n \geq 1}$ is a non increasing sequence (since the probability of staying longer in transient states is lower than the probability of staying shorter)

$$\begin{aligned} U_i &= P_i(\bigcap_{n \geq 1} \{X_n \in T, \dots, X_1 \in T\}) \\ &= \lim_{n \rightarrow \infty} U_i^{(n)} = \lim_{n \rightarrow \infty} P_i(X_n \in T, \dots, X_1 \in T) \end{aligned}$$

$$\begin{aligned} U_i^{(n+1)} &= P_i(X_{n+1} \in T, \dots, X_2 \in T, X_1 \in T) \\ &= \underbrace{P_i(X_{n+1} \in T, \dots, X_2 \in T \mid X_1 \in T)}_{P_i(X_1 \in T)} \underbrace{P_i(X_1 \in T)}_{P_i(X_1 \in T)} \end{aligned}$$

$\underbrace{\quad}_{\text{IP(remain } n \text{ steps in } T \text{ starting from } X_1 \in T)}$

More precisely, we have:

$$\begin{aligned}
 U_i^{(n+1)} &= \sum_{j \in T} P_i(X_{n+1} \in T, \dots, X_2 \in T, X_1 = j) \\
 &= \underbrace{\sum_{j \in T} P_i(X_{n+1} \in T, \dots, X_2 \in T | X_1 = j)}_{\substack{\text{probability of remaining} \\ n \text{ steps in } T \text{ starting from} \\ X_1 = j \text{ (and not } X_0 = i\text{)}}} \underbrace{P_i(X_1 = j)}_{p_{ij}} \\
 &= \sum_{j \in T} U_j^{(n)} p_{ij}
 \end{aligned}$$

Applying the limits on both sides we get: $U_i = \sum_{j \in T} p_{ij} U_j \quad \forall i \in T$

Now we have to prove that is the biggest $[0,1]$ -valued solution.

Suppose we have a bigger solution $V_i = \sum_{j \in T} p_{ij} V_j \quad \forall i \in T, 0 \leq V_i \leq 1$

$$\Rightarrow U_i^{(1)} = P_i(X_1 \in T) = \sum_{j \in T} p_{ij} \geq \sum_{j \in T} p_{ij} V_j = V_i \Rightarrow U_i^{(1)} \geq V_i \quad \forall i \in T$$

↑
since $0 \leq V_j \leq 1$

By induction: $U_i^{(n)} \geq V_i \quad \forall i \in T$, then:

$$U_i^{(n+1)} = \sum_{j \in T} p_{ij} U_j^{(n)} \geq \sum_{j \in T} p_{ij} V_j = V_i$$

And so, by considering the limit for $n \rightarrow \infty$ we have that $U_i^{(n)}$ increases to U_i and so $U_i \geq V_i \quad \forall i \in T$, which is a contradiction.

We also prove the remark: T finite $\Rightarrow U_i = 0 \quad \forall i \in T$

$$\begin{aligned}
 U_i &= \sum_{j \in T} p_{ij} U_j = \sum_{j \in T} p_{ij} \sum_{k \in T} p_{jk} U_k \\
 &= \sum_{j \in T} \sum_{k \in T} p_{ij} p_{jk} U_k \\
 &\leq \sum_{k \in T} \left(\sum_{j \in T} p_{ij} p_{jk} \right) U_k = \sum_{k \in T} p_{ik}^{(2)} U_k
 \end{aligned}$$

By iterating n times we get:

$$\begin{aligned}
 U_i &= \sum_{k_1, \dots, k_n \in T} p_{ik_1} \cdots p_{ik_{n-1} k_n} U_{k_n} \\
 &\leq \sum_{\substack{k_1, \dots, k_{n-1} \in E \\ k_n \in T}} (\dots) U_{k_n} = \sum_{k_n \in T} p_{ik_n}^{(n)} U_{k_n} \\
 &\Rightarrow \forall n: U_i \leq \sum_{j \in T} p_{ij}^{(n)} U_j \leq \sum_{j \in T} p_{ij}^{(n)} \leq 1
 \end{aligned}$$

$U_j \in [0,1]$

Since $\forall j \in T \quad p_{ij}^{(n)} \xrightarrow{n \rightarrow \infty} 0$ (transient) $\Rightarrow U_i = 0 \quad \forall i \in T$ ■

Example: Gambler's ruin problem (against a bank)

$(X_n)_{n \geq 0}$ MC, X_n = player's capital at time n , $E \subset \mathbb{N}$

$$\begin{bmatrix} 1 & 0 & 0 & 0 & 0 & \dots \\ q & 0 & p & 0 & 0 & \dots \\ 0 & q & 0 & p & 0 & \dots \end{bmatrix} \quad \begin{array}{l} p+q=1 \\ p, q > 0 \end{array}$$

Classes: $\{0\}$ recurrent

$\{1, 2, 3, \dots\}$ transient

$$P_{00}^{(n)} = 1, \quad \sum_{n \geq 1} P_{00}^{(n)} = +\infty$$

$$\forall i \geq 1 \quad P_i(\bigcap_{i \geq 1} (X_n = i)) < 1 - q^i < 1$$

We want to calculate the probability that the gambler is never ruined.
 (This is equivalent to calculate the probability of remain forever in the transient states $T = \{1, 2, 3, \dots\} = N \setminus \{0\}$)

$$U_i = \sum_{j \geq 1} p_{ij} U_j \quad : \quad \begin{cases} U_1 = p U_2 & (U_0 = 0) \\ U_i = p U_{i+1} + q U_{i-1} & i > 1 \end{cases}$$

General difference equation:

$$\begin{array}{l} a y_{i+1} + b y_i + c y_{i-1} = 0 \\ a, b, c \in \mathbb{R}, a \neq 0 \end{array} \quad \Rightarrow \quad \begin{cases} y_i = k_1 \lambda_1^i + k_2 \lambda_2^i \\ \lambda_1, \lambda_2: a\lambda^2 + b\lambda + c = 0 \\ (\text{if } \lambda_1 \neq \lambda_2) \\ (\text{if } \lambda_1 = \lambda_2: y_i = (k_1 + i k_2) \lambda_1^i) \end{cases}$$

In our case (defining $U_0 = 0$):

$$p U_{i+1} - U_i + q U_{i-1} = 0 \quad \Rightarrow \quad U_i = p U_{i+1} + q U_{i-1} \quad \forall i \geq 1$$

$$\text{characteristic equation: } p\lambda^2 - \lambda + q = 0 \quad \Rightarrow \quad \lambda_1 = 1, \lambda_2 = \frac{q}{p}$$

- $p \neq q$: Generic solution: $U_i = k_1 1^i + k_2 \left(\frac{q}{p}\right)^i$

$$U_0 = 0 \iff k_2 = -k_1 \quad \text{and so:}$$

$$U_i = k_1 \left(1 - \left(\frac{q}{p}\right)^i\right)$$

- $q > p$: $U_i \in [0, 1] \iff k_1 = 0$, and so:

$$U_i = 0 \quad \forall i \iff \text{Ruin, i.e. absorption to 0}$$

- $q < p$: $\lim_{i \rightarrow \infty} U_i = k_1$, in order to get the bigger

solution we get $k_1 = 1$, and so:

$$U_i = 1 - \left(\frac{q}{p}\right)^i = \text{probability that the gambler becomes \infty-rich}$$

$$\left(\frac{q}{p}\right)^i = \text{probability of the ruin} \\ (\Leftarrow \text{probability of absorption in 0})$$

- $p = q = \frac{1}{2}$: $U_i = k_1 + k_2 i$: $U_0 = 0 \Rightarrow c k_2 = -k_1$
 $U_0 = 0 \Rightarrow k_1 = 0 \Rightarrow k_2 = 0$

Theorem (Absorption probability in a recurrent class):

The probability of absorption in a recurrent class C is defined as $(V_i)_{i \in T}$.
 V_i is the smallest $[0, 1]$ -valued solution of:

$$V_i = \sum_{j \in C} p_{ij} + \sum_{j \in T} p_{ij} V_j$$

Moreover, if T is finite, the solution $((V_i)_i)$ is unique.

Proof.

We define $V_i^{(n)} = \mathbb{P}_i(X_n \in C, X_{n-1} \notin C, \dots, X_1 \notin C) = \text{probability of absorption in } C \text{ at time } n$

$$V_i^{(1)} = \sum_{j \in C} p_{ij}$$

$$V_i^{(n)} = \mathbb{P}_i(X_n \in C, X_{n-1} \notin C, \dots, X_1 \notin C) \quad (n > 1)$$

$$= \sum_{j \in T} \mathbb{P}_i(X_n \in C, X_{n-1} \notin C, \dots, X_2 \notin C, X_1 = j)$$

$$= \sum_{j \in T} \mathbb{P}_i(X_n \in C, \dots, X_2 \notin C \mid X_1 = j) \quad (\mathbb{P}(X_1 = j))$$

$$= \sum_{j \in T} V_j^{(n-1)} p_{ij}$$

The event "absorption in C " can be written as: $\bigcup_{n=1}^{\infty} \{X_n \in C\}$

Moreover: $\bigcup_{n \geq 1} \{X_n \in C\} = \bigcup_{n \geq 1} \{X_n \in C, X_{n-1} \notin C, \dots, X_1 \notin C\}$
 ↑
 disjoint union

$$V_i^{(n)} = \sum_{j \in T} p_{ij} V_j^{(n-1)}$$

$$\begin{aligned} \implies V_i &= P_i(\bigcup_{n \geq 1} \{X_n \in C\}) \\ &= P_i(\bigcup_{n \geq 1} \{X_n \in C, X_{n-1} \notin C, \dots, X_1 \notin C\}) \\ &= \sum_{n \geq 1} P_i(\{X_n \in C, X_{n-1} \notin C, \dots, X_1 \notin C\}) \\ &= \sum_{n \geq 1} V_i^{(n)} \\ &= V_i^{(1)} + \sum_{n \geq 1} V_i^{(n)} \\ &= \sum_{j \in C} p_{ij} + \sum_{n \geq 2} \sum_{j \in T} p_{ij} V_j^{(n-1)} \\ &= \sum_{j \in C} p_{ij} + \sum_{j \in T} \sum_{n \geq 1} p_{ij} V_j^{(n)} \\ &= \sum_{j \in C} p_{ij} + \sum_{j \in T} p_{ij} V_j \end{aligned}$$

Moreover, it's the smallest $[0,1]$ -valued solution:

let $(X_i)_{i \in T}$ be another $[0,1]$ -valued solution of $X_i = \sum_{j \in C} p_{ij} + \sum_{j \in T} p_{ij} X_j$

since $V_i^{(1)} = \sum_{j \in C} p_{ij}$ we have:

$$X_i = V_i^{(1)} + \sum_{j \in T} p_{ij} X_j > V_i^{(1)} \quad \forall i \in T.$$

Suppose $X_i \geq \sum_{k=1}^n V_i^{(k)}$

$$\begin{aligned} X_i &= \sum_{j \in C} p_{ij} + \sum_{j \in T} p_{ij} X_j \geq \underbrace{\sum_{j \in C} p_{ij}}_{V_j^{(1)}} + \underbrace{\sum_{j \in T} p_{ij} \sum_{k=1}^n V_j^{(k)}}_{\sum_{k=1}^n (\sum_{j \in T} p_{ij} V_j^{(k)})} = \sum_{k=1}^{n+1} V_i^{(k)} \\ &\quad V_j^{(1)} + \sum_{k=1}^n (\sum_{j \in T} p_{ij} V_j^{(k)}) \\ &\quad V_j^{(1)} + \sum_{k=1}^n V_j^{(k+1)} \end{aligned}$$

so, by induction, we showed that $X_i \geq \sum_{k=1}^n V_i^{(k)} \quad \forall n$

\implies taking the limit $n \rightarrow \infty$; $X_i \geq V_i \quad \forall i \in T$

Example: Gambler's ruin problem

$$\begin{bmatrix} 1 & 0 & 0 & 0 & \dots \\ 0 & p & 0 & 0 & \dots \\ 0 & 0 & p & 0 & \dots \\ 0 & 0 & 0 & q & \dots \end{bmatrix} \quad \text{States} = \{0, 1, 2, 3, \dots\}$$

$p > q$ (since if $q \geq p$ the probability of ruin is 1)

{0} absorbing, {1, 2, 3, ...} transient

Ruin \iff absorption in 0 : $C = \{0\}, T = \{1, 2, 3, \dots\}$

$$\begin{cases} V_1 = pV_2 + q \\ V_i = pV_{i+1} + qV_{i-1}, \quad i > 1 \end{cases}$$

= q

Or, with the convention $V_0 = 1$:

$$\begin{cases} V_0 = 1 \\ V_i = pV_{i+1} + qV_{i-1} \quad i \geq 1 \end{cases}$$

General solution:

$$V_i = k_1 + k_2 \left(\frac{q}{p}\right)^i$$

$$\text{and since } V_0 = 1 \Rightarrow k_1 + k_2 = 1 \Rightarrow V_i = k_1 + (1-k_1) \left(\frac{q}{p}\right)^i$$

We want the smallest $[0,1]$ -valued solution:

$$\begin{aligned} i \rightarrow \infty \quad V_i &\rightarrow k_1 \in [0,1] \Rightarrow k_1 = 0 \text{ is the smallest} \\ \Rightarrow V_i &= \left(\frac{q}{p}\right)^i \end{aligned}$$

Theorem (Mean absorption time in recurrent classes):

Suppose that the set of states E is finite and there is an unique recurrent class C . Then, the mean absorption time w_i in C is finite and satisfies:

$$w_i = 1 + \sum_{j \in T} p_{ij} w_j \quad i \in T \quad (w_i \text{ is the unique solution})$$

Interpretation: the mean time for absorption in C (w_i) is the time for the 1-step absorption ($=1$) + mean time for absorption in C after moving to a different transient state j

Lemma: Let V be a \mathbb{N} -valued random variable $\Rightarrow E[V] = \sum_{n=0}^{\infty} P(V > n)$

proof. (lemma)

$$\begin{aligned} E[V] &= \sum_{m=1}^{\infty} m P(V=m) = \sum_{m=1}^{\infty} \sum_{k=1}^{\infty} 1 \cdot P(V=m) \\ &= \sum_{k=1}^{\infty} \sum_{m=k}^{\infty} P(V=m) \\ &= \sum_{k=1}^{\infty} P(V \geq k) \\ &= \sum_{n=0}^{\infty} P(V > n) \quad \blacksquare \end{aligned}$$

proof. (Theorem)

Step 1. Check that w_i are finite by: $\exists c \in (0,1) : P_i(V > n) \leq c^n$, where V is the absorption time in C .

$$P_i(V > n) = P(X_n \in T, \dots, X_1 \in T) := U_i^{(n)} = \text{probability of staying in transient states from 1 to } n \text{ steps}$$

$$U_i^{(n+1)} = \sum_{j \in T} p_{ij} U_j^{(n)}$$

$$\begin{aligned} U_i^{(n+2)} &= \sum_{\substack{j \in T \\ k \in T}} p_{ij} p_{jk} U_k^{(n)} \leq \sum_{k \in T} \left(\sum_{j \in E} p_{ij} p_{jk} \right) U_k^{(n)} = \\ &= \sum_{k \in T} p_{ik}^{(2)} U_k^{(n)} \end{aligned}$$

$$\text{Iterating: } U_i^{(n+m)} \leq \sum_{j \in T} p_{ij}^{(m)} U_j^{(n)}$$

$$\text{If } i, j \text{ are transient } (\sum_{n \geq 1} p_{ij}^{(n)} < \infty) \Rightarrow \lim_{m \rightarrow \infty} p_{ij}^{(m)} = 0$$

\Rightarrow If T is finite we can choose m s.t. $\sum_{j \in T} p_{ij}^{(m)} < 1 \quad \forall i$ and

$$\text{we set } M = \max_{i \in T} \left\{ \sum_{j \in T} p_{ij}^{(m)} \right\}, \text{ that will be } < 1.$$

$$\begin{aligned}
 \text{Then: } U_i^{(n)} &\leq \sum_{j \in T} p_{ij}^{(m)} U_j^{(n-m)} \\
 &\leq \max_{j \in T} U_j^{(n-m)} \sum_{j \in T} p_{ij}^{(m)} \\
 &\leq \max_{j \in T} U_j^{(n-m)} M \quad \rightarrow \text{the probability of remaining forever in a transient state is exponentially decreasing}
 \end{aligned}$$

$$\text{Iterating: } \max_{i \in T} U_i^{(n)} \leq M^2 \max_{i \in T} U_i^{(n-2m)} \leq \dots$$

$$\leq M^{\frac{n}{m}} \max_{i \in T} U_i^{(n - (\frac{n}{m})m)}$$

≤ 1 since they're probabilities

$$\Rightarrow P_i(V > n) \leq \max_{i \in T} U_i^{(n)} \leq M^{\frac{n}{m}} \leq M^{\frac{n}{m}-1}$$

$$\leq M^{-1} (M^{\frac{1}{m}})^n$$

$$\Rightarrow E_i[V] = \sum_{n=0}^{\infty} P_i(V > n) \leq M^{-1} \sum_{n=0}^{\infty} (M^{\frac{1}{m}})^n = \\ \leq (M(1-M^{\frac{1}{m}}))^{-1} < \infty$$

Step 2. Check the formula: $w_i = 1 + \sum_{j \in T} p_{ij} w_j$

$$\begin{aligned}
 E_i[V] &= \sum_{n=1}^{\infty} n P_i(V=n) \\
 &= \sum_{n=1}^{\infty} n P_i(V=n, X_1 \in T) + \sum_{n=1}^{\infty} n P_i(V=n, X_1 \in C) \\
 &\quad \text{if } X_1 \in T \text{ then } V \geq 2 \\
 &\quad \text{since we're starting from } X_0 \in T, X_1 \in T \Rightarrow \text{we can be in } C \text{ not before than } X_2. \\
 &= \left[\sum_{n=2}^{+\infty} n \sum_{j \in T} P_i(V=n, X_1=j) \right] + \sum_{j \in C} p_{ij} \\
 &= \left[\sum_{n=2}^{\infty} n \sum_{j \in T} P_i(V=n | X_1=j) p_{ij} \right] + \sum_{j \in C} p_{ij} \\
 &= P(V=n | X_1=j, X_0=i) \stackrel{M}{=} P(V=n | X_1=j) \\
 &= P(V=n-1 | X_0=j) = P_j(V=n-1) \\
 &= \left[\sum_{n=1}^{\infty} (n+1) \sum_{j \in T} P_j(V=n) p_{ij} \right] + \sum_{j \in C} p_{ij} \\
 &= \left[\sum_{j \in T} \underbrace{\sum_{n=1}^{\infty} n P_j(V=n) p_{ij}}_{E_j[V]} \right] + \left[\sum_{j \in T} \underbrace{\sum_{n=1}^{\infty} p_{ij} \underbrace{\sum_{n=1}^{\infty} P_j(V=n)}_1} \right] + \left[\sum_{j \in C} p_{ij} \right] \\
 &\quad \boxed{1}
 \end{aligned}$$

$$\Rightarrow w_i = \sum_{j \in T} w_j p_{ij} + 1$$

Remark: If there is a situation with two (or more) recurrent classes C_1, C_2 and there is a strictly positive probability of ending in C_1 and a strictly positive prob. of ending in C_2 , then the mean absorption time of ending in C_1 is ∞ (since we can end in C_2). To evaluate the "true mean" for, for instance, C_1 we have to neglect all C_i $i \neq 1$.

Example: Coupon collector problem.

The collection is made of N pictures. At time n we buy an envelope containing a random picture. How many pictures do we have to buy to complete the collection?

$(X_n)_{n \geq 0}$ MC : $X_n = \#$ different pictures collected at time n
 $(X_0 = 0)$

Transitions probabilities :

$$P(X_{n+1} = k+1 | X_n = k) = \frac{N-k}{N}$$

$$P(X_{n+1} = k | X_n = k) = \frac{k}{N}$$

$$P(X_{n+1} = j | X_n = k) = 0 \quad \forall j \neq k, k+1$$

Transition matrix :

$$\begin{matrix} 0 & 0 & 1 & 0 & 0 & 0 & \dots \\ 1 & 0 & \frac{1}{N} & \frac{N-1}{N} & 0 & 0 & \dots \\ 2 & 0 & 0 & \frac{2}{N} & \frac{N-2}{N} & 0 & \dots \\ \vdots & \ddots & & & & & \\ N-1 & \ddots & & & 0 & \frac{N-1}{N} & \frac{1}{N} \\ N & \ddots & & & 0 & 0 & 1 \end{matrix} \quad T = \{0, 1, \dots, N-1\} \quad C = \{N\}$$

w_i = mean time for reaching N starting from i :

$$\Rightarrow w_i = 1 + \frac{i}{N} w_i + \frac{N-i}{N} w_{i+1}$$

$$\Rightarrow \dots \Rightarrow w_i = \frac{N}{N-i} + \frac{N}{N-i-1} + \dots + \frac{N}{N-(N-1)} \quad (i=0, \dots, N-1)$$

$$\Rightarrow w_0 = \sum_{k=0}^{N-1} \frac{N}{N-k} = N \sum_{k=1}^N \frac{1}{k} \underset{N \text{ large}}{\approx} N \log(N)$$

If a collector has $N \log(N)$ pictures (generally) :

- probability that there is no second copy of a given copy?

$$\left(\frac{N-1}{N}\right)^{N \log(N)} = \left(\left(\frac{N-1}{N}\right)^N\right) \log(N) \approx e^{-\log(N)} = \frac{1}{N}$$

- probability that k fixed copies have no second copy:

$$\left(\frac{N-k}{N}\right)^{N \log(N)} = \left(1 - \frac{k}{N}\right)^{N \log(N)} \approx e^{-k \log(N)} = \left(\frac{1}{N}\right)^k$$

Example: Gambler's ruin against a bank

We consider the fair case of the game, namely $p=q=\frac{1}{2}$

$$\begin{bmatrix} 1 & 0 & 0 & 0 & 0 & \dots \\ \frac{1}{2} & 0 & \frac{1}{2} & 0 & 0 & \dots \\ 0 & \frac{1}{2} & 0 & \frac{1}{2} & 0 & \dots \\ 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} & 0 & \dots \\ 0 & 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} & \dots \\ 0 & 0 & 0 & 0 & \frac{1}{2} & 0 & \dots \\ 0 & 0 & 0 & 0 & 0 & \frac{1}{2} & \dots \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

States = $\{0, 1, 2, \dots, N\}$

Absorbing classes : $\{0\}, \{N\}$

For $\{1, \dots, N-1\}$ we are interested in computing the mean win time w.r.t. $P_i(\cdot | \bigcap_{n=1}^{\infty} (X_n \neq N))$

win occurs with probability = 1

We have to exclude the possibility of going from $N-1$ to N because $\{N\}$ is a recurrent class and if we consider it too then the mean of absorption in $C_1 = \infty$ will be ∞ .
 (since we would consider the case in which we end up in $C_2 = \{N\}$)

$$\begin{cases} w_0 = 0 \\ w_i = 1 + \frac{1}{2} w_{i-1} + \frac{1}{2} w_{i+1} & 1 \leq i \leq N-1 \\ w_{N-1} = 1 + w_{N-2} \end{cases}$$

This is a non-homogeneous finite diff. equation.
 The associated homogeneous eq. and charact. eq.:

$$\begin{aligned} y_i &= \frac{1}{2} y_{i+1} + \frac{1}{2} y_{i-1} \\ \lambda^2 - 2\lambda + 1 &= 0 \quad (\lambda=1) \end{aligned} \quad \Rightarrow \quad y_i = c_1 + i c_2 \quad (\text{general solution})$$

We try with: $w_i = a i^2 + b i + c$ (*)

$$\begin{aligned} \Rightarrow 1 + \frac{1}{2} w_{i-1} + \frac{1}{2} w_{i+1} &= 1 + \frac{1}{2} a(i-1)^2 + \frac{1}{2} b(i-1) + c \frac{1}{2} + \\ &\quad + \frac{1}{2} a(i+1)^2 + \frac{1}{2} b(i+1) + c \frac{1}{2} \\ &= 1 + a i^2 + a + b i + c \end{aligned}$$

We compare $1 + a i^2 + a + b i + c$ with (*) and we obtain:

$$a + c + 1 = 0 \quad \Rightarrow \quad \text{we consider the simpler: } a = -1, c = 0$$

$$\Rightarrow \text{special solution: } w_i = -i^2$$

$$\text{general solution: } y_i = c_1 + i c_2$$

$$\Rightarrow \text{boundary conditions } (w_i = -i^2 + c_1 + i c_2) :$$

$$\begin{cases} w_0 = 0 \\ w_1 = 1 + \frac{1}{2} w_2 \end{cases} \quad \Rightarrow \quad c_1 = 0$$

$$w_{N-1} = 1 + w_{N-2} \quad \Rightarrow \quad c_2 = 1 - (N-2)^2 + (N-1)^2 = 2N-2$$

\Rightarrow Ultimate solution:

$$w_i = -i^2 + 2(N-1)i \quad = \text{mean duration of the game if you're sure you'll be visited}$$

Theorem (Transience criterion):

$(X_n)_{n \geq 0}$ irreducible MC (with countable space E) is transient if and only if there exists a bounded non-constant solution of:

$$\sum_{k \in E} p_{jk} y_k = y_j \quad \text{for all } j \in E \quad \underline{\text{but one.}}$$

((*))

for all $j \in E$ except for
(at most) one.

proof. (⊗)

1. This condition is necessary:

Suppose $(X_n)_{n \geq 0}$ transient. We denote by e the (unique) state for which ((*)) does not hold.

We consider the transformed MC in which e is absorbing. Its transition matrix is:

$$\tilde{p}_{ij} = \begin{cases} \delta_{ij} & j = e \\ p_{ij} & j \neq e \end{cases}$$

Since the given MC is transient (the transformed is not), there exists $i \in E$ such that:

$$\tilde{V}_i := P_i(T_e < +\infty) < 1 \quad (T_e = \text{first entrance time in } e) \quad 11$$

< 1 because if $\tilde{V}_i = P_i(T_e < \infty) = 1 \quad \forall i$ then:

$$\begin{aligned} P_e(T_e < +\infty) &= p_{ee} + \sum_{k \neq e} p_{ek} P_k(T_e < +\infty) \\ &= p_{ee} + \sum_{k \neq e} p_{ek} = 1 \end{aligned}$$

which contradicts the transiency of $(X_n)_{n \geq 0}$.

$$\Rightarrow \tilde{V}_i < 1 \text{ for some } i \text{ and } \tilde{V}_e = 1.$$

$\Rightarrow (V_i)_{i \in E}$ is $[0, 1]$ -valued (bounded) and non-constant.

Moreover, $(V_i)_{i \in E}$ are the absorption probabilities in e for the MC with $(\tilde{p}_{ij})_{i, j \in E}$, therefore:

$$\tilde{V}_i = \tilde{p}_{ie} + \sum_{k \neq e} \tilde{p}_{ik} \tilde{V}_k \quad (\star)$$

$$\text{but } \tilde{p}_{ik} = p_{ik} \quad \forall k \neq e \quad (\text{if } i \neq e), \quad \tilde{V}_e = 1 \Rightarrow \tilde{V}_i = \sum_{k \in E} p_{ik} \tilde{V}_k$$

$\Rightarrow (\star)$ and $((*)$) are equivalent and satisfied

2. This condition is sufficient:

$$\text{Suppose: } \sum_{k \in E} p_{ik} y_k = y_i \quad \forall i \in E \setminus \{e\} \quad (\Delta)$$

where y_i is bounded and non-constant. $((y_i)_{i \in E})$.

Then, if we consider $(\tilde{p}_{ij})_{i, j \in E}$ as before, from (Δ) we get:

$$\sum_{k \in E} \tilde{p}_{ik} y_k = y_i \quad \forall i \in E$$

$$\begin{aligned} \text{Iterating: } \underbrace{\sum_{k \in E, h \in E} \tilde{p}_{ik} \tilde{p}_{kh} y_h}_{\sum_{h \in E} \left(\sum_{k \in E} \tilde{p}_{ik} \tilde{p}_{kh} \right) y_h} &= y_i \\ &= \sum_{h \in E} \tilde{p}_{ih}^{(2)} y_h \end{aligned}$$

Iterating n times:

$$\sum_{k \in E} \tilde{p}_{ik}^{(n)} y_k = y_i$$

If the MC $(X_n)_{n \geq 0}$ was recurrent then the transformed MC $(\tilde{p}_{ij})_{i, j}$ would converge to the state e , namely:

$$\lim_{n \rightarrow \infty} \tilde{p}_{ie}^{(n)} = 1 \quad \forall i \in E$$

$$(\lim_{n \rightarrow \infty} \tilde{p}_{ie}^{(n)} = \lim_{n \rightarrow \infty} P_j(X_n = e) = \lim_{n \rightarrow \infty} P_j(T_e \leq n) = 1 \Leftrightarrow \text{recurrence})$$

$$\text{Then } \forall j \neq e: |y_j - y_e| = \lim_n |y_j - \tilde{p}_{je}^{(n)} y_e|$$

$$\begin{aligned} \textcircled{*} \quad \sum_{k \in E} \tilde{p}_{jk}^{(n)} y_k &= \sum_{k \neq e} \tilde{p}_{jk}^{(n)} y_k + \tilde{p}_{je}^{(n)} y_e \\ &\downarrow y_j \quad \forall j \in E \\ &= \lim_n |y_j - (\tilde{p}_{je}^{(n)} y_e - \sum_{k \neq e} \tilde{p}_{jk}^{(n)} y_k)| \\ &= \lim_n | \sum_{k \neq e} \tilde{p}_{jk}^{(n)} y_k | \\ &\leq \sup_k |y_k| \lim_n \sum_{k \neq e} \tilde{p}_{jk}^{(n)} \\ &\leq \sup_k |y_k| \lim_n (1 - \tilde{p}_{je}^{(n)}) = 0 \end{aligned}$$

$$\Rightarrow y_j = y_e \quad \forall j \neq e$$

$$\Rightarrow (y_i)_{i \in E} \text{ constant} \Rightarrow \text{contradiction}$$

Theorem (Recurrence criterion):

$(X_n)_{n \geq 0}$ irreducible MC. If there exists a family $(y_j)_{j \in E}$ such that:

- $\sum_{k \in E} p_{jk} y_k \leq y_j \quad \forall j \in E$ but (at most) one
- $\lim_{k \rightarrow \infty} y_k = +\infty \quad (\forall M > 0 \quad \exists F \subseteq E \text{ finite s.t. } y_k > M \quad \forall k \notin F)$

then the Markov chain $(X_n)_{n \geq 0}$ is recurrent.

Example: Application of this criteria to a queue model.

- the counter serves one customer per unit of time
- the number of customers arriving in the system at each unit of time is a random variable $A_n = \text{arrivals at } n$ s.t.

$$P(A_n = k) = a_k \quad 0 \leq a_k \leq 1, \quad \sum_k a_k = 1$$

- the number of customers arriving at different times are **II**

Transition matrix:

$$\begin{bmatrix} a_0 & a_1 & a_2 & a_3 & \dots \\ a_0 & a_1 & a_2 & a_3 & \dots \\ 0 & a_0 & a_1 & a_2 & \dots \\ \dots & & & & \end{bmatrix} \quad \left. \begin{array}{l} a_0 > 0 \\ 0 < a_0 + a_1 < 1 \end{array} \right\} \begin{array}{l} \text{there is always the} \\ \text{possibility of 0 arrivals} \end{array}$$

→ the Markov chain is irreducible

We want to find out if this MC is transient or recurrent

(transient = the queue explodes (at finite time))

recurrent = the queue can be very long but it'll go back to 0).

The average number of customers arriving at each time:

$$\lambda = \sum_{k \geq 1} a_k k \quad \text{intuitively: } \lambda > 1 \Rightarrow \text{transient} \quad \lambda \leq 1 \Rightarrow \text{recurrent}$$

- $\lambda > 1$

We want to prove that is transient → first criteria.

We look for solutions of:

$$\sum_{k \geq 0} p_{jk} y_k = y_j \quad \forall j > 0 \quad (j \neq 0)$$

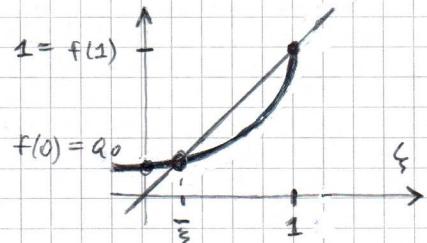
of the form $y_k = \xi^k$ with $0 < \xi < 1$

$$\Rightarrow \sum_{k \geq j+1} a_{k-j+1} \xi^k = \xi^j$$

$$\Leftrightarrow \sum_{k \geq j+1} a_{k-j+1} \xi^{k-j+1} = \xi^j$$

$$\Leftrightarrow \sum_{h \geq 0} a_h \xi^h = \xi^j$$

$$f(\xi) := \sum_{h \geq 0} a_h \xi^h \quad |\xi| \leq 1 :$$



$a_h \geq 0 \Rightarrow \xi \mapsto f(\xi)$ is increasing in $[0, 1]$

$$f(1) = \sum_{h \geq 1} h a_h = \lambda$$

$$f'(1) = \lambda > 1$$

⇒ ∃ξ ∈ (0, 1) such that $f(\xi) = 1$

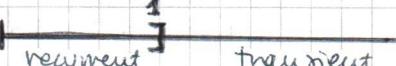
• $\lambda \leq 1$

We want to prove that if recurrent \Rightarrow second criteria.
We look for the unbounded solution of:

$$\sum_{k \geq 0} p_{jk} y_k = y_j \quad \forall j > 0$$

Choose $y_n = k$, for $j > 0$:

$$\begin{aligned} \sum_{k \geq 0} p_{jk} y_k &= \sum_{k \geq j-1} a_{k-j+1} k \\ &= \sum_{h \geq 0} a_h (h + (j-1)) \\ &= \sum_{h \geq 0} a_h h + (j-1) \sum_{h \geq 0} a_h \\ &= \lambda + j-1 = (\lambda - 1) + j = (\lambda - 1) + y_j \leq 1 \quad (\lambda \leq 1) \end{aligned}$$

So: 

The smallest is $\lambda \Rightarrow$ the smallest the number of arrivals

Indeed for $0 < \lambda < 1$ it is fast recurrent $\Leftrightarrow \exists$ invariant distribution
(the system has a stationary state)

Theorem (Sufficient condition for the existence of invariant distributions)

$(x_n)_{n \geq 0}$ irreducible MC.

If we find $(y_j)_j, (x_j)_j$ both unbounded ($\lim_{j \rightarrow \infty} y_j = \lim_{j \rightarrow \infty} x_j = +\infty$) s.t.:

$$\sum_{k \geq 0} p_{jk} y_k \leq y_j - x_j \quad \forall j$$

\Rightarrow the MC admits a unique invariant density

Example: Application at the queue model.

We assume $\sum_{k \geq 0} k^2 a_k < +\infty$ (= the arrivals have a finite second order moment, we call it m_2)

Consider $y_j = j^2$

$$\begin{aligned} j > 0 : \sum_{k \geq 0} p_{jk} y_k &= \sum_{k \geq j-1} a_{k-j+1} k^2 \\ &= \sum_{h \geq 0} a_h (h + (j-1))^2 \\ &= \sum_{h \geq 0} h^2 a_h + 2(j-1) \sum_{h \geq 0} h a_h + (j-1)^2 \sum_{h \geq 0} a_h \\ &= m_2 + 2(j-1)\lambda + (j-1)^2 \\ &= j^2 - 2j + 1 + 2j\lambda - 2\lambda + m_2 \\ &= j^2 - 2(1-\lambda)j + (1-2\lambda + m_2) \\ &= y_j - (2(1-\lambda)j - (1-2\lambda + m_2)) \\ &:= y_j - x_j \quad (x_j \rightarrow \infty \text{ as } j \rightarrow \infty) \quad \text{if } \lambda < 1 \end{aligned}$$

\Rightarrow The MC admits a unique invariant density if $\lambda < 1$.

Def. SOJOURN TIME: We call $T = \text{set of transient states}$. Consider the subset $S \subseteq T$.
We define the sojourn time in S the total time spent in S :

$$T_S = \sum_{n \geq 0} 1_{\{X_n \in S\}}$$

Def. MOMENT GENERATING FUNCTION: The moment gen. fun. of a random var. T is:

$$m_i = E_i[z^T] \quad \forall i \in E$$

Notice that: if S ($\subseteq T$ = transient states) is finite, then the moment generator function:

$$m_i = E_i [z^T] \quad \text{is well defined}$$

Alternative notation: $E[z^T] = E[e^{Tt}] (= \phi(t), t: e^t = z)$

Properties :

(of the moment generator function of $T = T_S$)

- $m_i(z) = \sum_{k \geq 0} z^k P_i(T=k)$

it can be seen as Taylor series, where $P_i(T=0)$ is the constant, $P_i(T=1)$ is the coeff. multiplying the first order term, etc..

- $m_i'(z) = E_i \left[\frac{d}{dz} (e^{T \log(z)}) \right] = E_i \left[\frac{1}{z} z^T \right] = E_i [z^{T-1}]$

- $m_i(z) = 1$ if i is recurrent

- $m_i(z) = z \sum_j p_{ij} m_j(z)$ if $i \in S$

- $m_i(z) = \sum_j p_{ij} m_j(z)$ if $i \in T \setminus S = (\text{transient states}) \setminus S$

Proof.

- i recurrent $\Rightarrow T=0 \Rightarrow E_i [z^T] = 1$

- $i \in T = \text{transient states}$

$$T_S = \mathbb{1}_{\{X_0 \in S\}} + \sum_{j \neq 1} \mathbb{1}_{\{X_0 \in S\}} := \mathbb{1}_{\{X_0 \in S\}} + \tilde{T}$$

$$E_i [z^T] = \sum_{j \in E} E_i [z^T | X_1=j] P_i(X_1=j)$$

$$= \sum_{j \in E} E_i [z^{\mathbb{1}_{\{X_0 \in S\}}} z^{\tilde{T}} | X_1=j] P_i(X_1=j)$$

$$\begin{aligned} z^{\mathbb{1}_{\{X_0 \in S\}}} &= \\ &= \begin{cases} z^0 = 1 & i \in S \\ z^1 = z & i \notin S \end{cases} \end{aligned}$$

$$= z^{\mathbb{1}_{\{i \in S\}}} \sum_{j \in E} E_i [z^{\tilde{T}} | X_1=j, X_0=i] p_{ij}$$

Markov property

$$= z^{\mathbb{1}_{\{i \in S\}}} \sum_{j \in E} E_i [z^{\tilde{T}} | X_1=j] p_{ij}$$

$$= z^{\mathbb{1}_{\{i \in S\}}} \sum_{j \in E} E_i [z^{\tilde{T}}] p_{ij}$$

$$\Rightarrow \bullet i \in S: m_i = z \sum_{j \in E} E_i [z^{\tilde{T}}] p_{ij}$$

$$\bullet i \notin S: m_i = \sum_{j \in E} E_i [z^{\tilde{T}}] p_{ij}$$

Example: Gambler's ruin

symmetric case ($p=q=\frac{1}{2}$)

finite state space := $\{0, 1, 2, 3\}$:

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ \frac{1}{2} & 0 & \frac{1}{2} & 0 \\ 0 & \frac{1}{2} & 0 & \frac{1}{2} \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

We want to compute the moment generator function of the total time spent in $S=\{1\}$ starting from the state 1.

$$\left\{ \begin{array}{l} m_0(z) = m_3(z) = 1 \\ m_1(z) = z \left(\frac{1}{2} m_2(z) + \frac{1}{2} m_0(z) \right) \end{array} \right.$$

$$m_2(z) = \frac{1}{2} m_1(z) + \frac{1}{2}$$

$$\Rightarrow m_1(z) = \frac{3/4 z}{1 - z/4}, \quad m_2(z) = \frac{1 + z/2}{2(1 - z/4)}$$

