

CSIT 495/595 - Introduction to Cryptography Introduction

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Need to Learn Cryptography

- Data (or information) is the most crucial thing for many companies and government agencies
- Strong need to protect data while at rest, being transmitted and shared with others
- **Cryptography** offers very useful tools to protect data
- Examples:
 - All your emails on Gmail are encrypted by Google
 - Where is your password stored and how it is your protected?

What is Cryptography?

- Cryptography is the art of writing or solving codes
- Cryptography deals with *the scientific study of techniques for securing digital information, transactions, and distributed computations* (Textbook)

Importance of Cryptography

- Historically, major consumers of cryptography are military and intelligence organizations
 - Breaking Japanese Naval code in the middle of the second world war
 - Breaking Enigma machine
- Breaking Blue-ray and DVD encryption
- Today, cryptography is everywhere!!
 - Security mechanisms that rely on cryptography are an integral part of almost any computer system

Importance of Cryptography

"Windtalkers" (2002) - about the secret code used by American military during world war II which baffled the Japanese.

"Breaking the Code" (1996) - about British mathematician, Alan Turing, "Father of computer science", who worked on breaking German military code during World War II.
<http://www.turing.org.uk/turing/>

"A Beautiful Mind" (2001) - Oscar-winning movie about US mathematician and Nobel laureate John Nash. Nash worked for the US military on secret codes.

Others: **"U-571" (2000)**, **"Swordfish" (2001)**, **"Enigma" (2001)**, **"Hackers" (1995)**

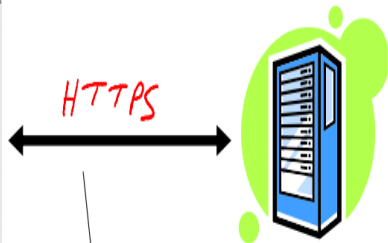
Check www.imdb.com for details.



Some well-known Applications

- Secure communication (e.g., HTTPS)
 - Email, e-commerce, ATM machines or cellular phones
- Encrypting files on Disk (EFS, TrueCrypt)
- Content Protection (e.g., Blue-Ray, DVD): CSS, AACS
- User Authentication (e.g., verifying user password/identity)
- and many other electronic applications over Internet

Secure Communication



no eavesdropping
no tampering

Data Confidentiality

Objectives

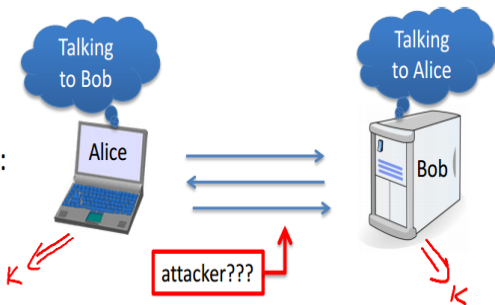
- Ensure data is accessible only to intended parties
- Data, whether at rest or being transmitted, need to remain confidential
- Often related to privacy

How this is achieved

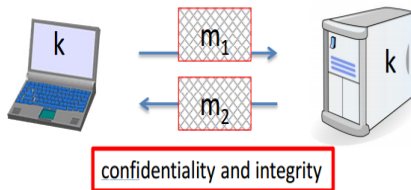
- Typically achieved by encryption the data
- Example: user passwords and bio-metric data need to be protected from other users/attackers

Data Confidentiality

Secret key establishment:



Secure communication:



Data Integrity

Objectives

- Maintain the accuracy and trustworthiness of data
- Data at rest or in transit should not be changed by unauthorized people (e.g., in the case of hacking)

How this is achieved

- Checksums or digital signatures
- Example: when accessing your bank account details, how can the user ensure that his/her account information is accurate and trustworthy?

Data Authentication

- Only authenticated people can receive/send messages in communication
- User has to show proper identity to be verified by the verifier - data authentication allows a receiver to verify that the data really was sent by the claimed sender
- Examples: PIN, password, biometrics
- **Key question:** Can the user authenticate him/herself to the server without revealing his/her identity?
 - YES!!! using proper cryptographic techniques

Private-Key Encryption: Introduction

- Classical cryptosystems are mostly based on private-key encryption where the security depends on a secret, commonly called **key**
- **General setting**: two parties sharing a key can communicate and exchange messages using the key
- The eavesdropper can monitor all the communication between the two parties
- **Basic Steps**:
 - Party 1 encrypts a message (known as **plaintext**) using the shared **key**
 - Party 1 sends the encrypted message (known as **ciphertext**) to party 2
 - Party 2 decrypts ciphertext to get the actual message
- Also referred to as **symmetric-key** setting

Private-Key Encryption: Applications

Secure Communication

- Two different parties separated in space (e.g., a CIA agent in USA communicating with his colleague in Europe)
- Parties can exchange messages after establishing key
- **Key Question:** How the parties can share the key securely?

Disk Encryption

- Same party communicating with itself over time
- User encrypts his files and stores them on Disk
- At a later time, he uses the same key to decrypt the files using the same key
- **Key Question:** How the user can remember/store the key securely and reliably?

Private-Key Encryption: Formal Definition

- Encryption schemes are defined over the message space \mathcal{M} , the key space \mathcal{K} , and the ciphertext \mathcal{C}

Key Generation (Gen)

- A probabilistic algorithm that selects a key $k \in \mathcal{K}$

Encryption (Enc)

- Input: message $m \in \mathcal{M}$ and k
- Output: Ciphertext $c \leftarrow \text{Enc}_k(m)$, where $c \in \mathcal{C}$

Decryption (Dec)

- Input: c and k
- Output: $m \leftarrow \text{Dec}_k(c)$

Kerckhoff's Principle (1)

- An adversary (or eavesdropper) can decrypt a ciphertext if he/she knows the decryption function (Dec) and the key used (k)
- Therefore, k has to be shared securely between the participating parties
- Do we also need to protect Dec from the adversary?
- Hide all the information related to the encryption scheme -
Is this a better idea??

Kerckhoff's Principle (2)

- Auguste Kerckhoff, a Dutch Cryptographer, argues the opposite in a paper he wrote in the late 19th century

***The cipher method must not be required to be secret,
and it must be able to fall into hands of the enemy
without inconvenience***

- Commonly known as **Kerckhoff's Principle**
- Encryption should be designed to be secure even if all the details of the encryption scheme are revealed to the adversary, except the key
- In short, security relies solely on the secrecy of the key

Kerckhoff's Principle (3)

- But.. Why this is correct?
- Three reasons:
 - 1 It is often unrealistic to assume encryption/decryption function will remain secret
 - 2 Easier to replace a key than to replace the entire encryption system
 - 3 Feasible for large-scale deployment: it is better to use same encryption software/algorithm + users might use different keys

Caesar's Cipher

- Named after Julius Caesar, who described it in 110 CE
- He encrypted the messages by shifting each alphabet 3 places forward

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C

Alphabet shifted by 3 spaces.

Caesar's Cipher - Example

Plaintext (the message)

a	t	t	a	c	k	a	t	d	a	w	n
---	---	---	---	---	---	---	---	---	---	---	---

e.g., 'c' is 'a' shifted by 2

e.g., 'f' is 'd' shifted by 2

c	v	v	c	e	m	c	v	f	c	y	p
---	---	---	---	---	---	---	---	---	---	---	---

Ciphertext (encrypted message)

- **Note:** punctuation, spaces, and numbers are removed
- What is the ciphertext when Caesar's cipher is used?

Caesar's Cipher - Security

- Encryption method is fixed and there is **no key**
- It is very easy for anyone to decrypt ciphertexts
- **ROT-13:**
 - A variant of Caesar's cipher with 13 places shift instead of 3
 - Still used in various online forums

Shift Cipher

- A keyed variant of Caesar's cipher
- Key k is a number between 0 and 25
- **Encryption**: letters in plaintext are shifted by k places forward
- **Decryption**: letters in ciphertext are shifted by k places backward

Shift Cipher - Formal Definition

- Suppose English alphabet set is mapped to the set $\{0, \dots, 25\}$ (i.e., $a = 0, b = 1$, and so on)
- Given a key k , encryption of a message $m = m_1 \dots m_\ell$ is

$$\text{Enc}_k(m) = c_1 \dots c_\ell, \text{ where } c_i = (m_i + k) \bmod 26$$

- Decryption function is given by

$$\text{Dec}_k(c_1 \dots c_\ell) = m_1 \dots m_\ell, \text{ where } m_i = (c_i - k) \bmod 26$$

Shift Cipher - Example

Let $k = 20$

Original	M	A	T	H	R	U	L	E	S
Number-fied	12	0	19	7	17	20	11	4	18
+key	32	20	39	27	37	40	31	24	38
mod 26	6	20	13	1	11	14	5	24	12
Letter-fied	G	U	N	B	L	O	F	Y	M

Shift Cipher - Attack Scenarios

- Is it possible to decrypt a ciphertext without knowing k ?
- YES!! It is trivial
- Brute-force or exhaustive-search attack:
 - There are only 26 possible keys
 - Given a ciphertext, use every possible key and compute 26 possible plaintexts (one of them certainly being the correct one)
- In general, when the ciphertext is long enough, then the only one on the list that “makes sense” is the correct plaintext

Key Space - Important Observations

- An encryption scheme should not be vulnerable to brute-force attacks
- **sufficient key-space principle**: Any secure encryption scheme must have a key space that is sufficiently large to avoid brute-force attacks in-feasible
- Example: key size of at least 2^{80}
- Nevertheless, the sufficient key space is a necessary condition, but not the only requirement, to ensure proper security

Substitution Ciphers

- Mono-alphabetic
- Poly-alphabetic

Mono-alphabetic Substitution Cipher (1)

- The key defines a fixed substitution for individual letters in the plaintext
- Each letter is mapped to one of the remaining letters
- **key space**: consists of all the bijections or permutations

③ Substitution Cipher

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
																									
Q	W	E	R	T	Y	U	I	O	P	A	S	D	F	G	H	J	K	L	Z	X	C	V	B	N	M

GRAY	FOX	HAS	ARRIVED
UKQN	YGB	IQL	QKKOCTR

Mono-alphabetic Substitution Cipher (2)

- The key space is of size $26!$ ($\approx 2^{88}$)
- Brute-force attack becomes difficult
- Does it mean this cipher is secure? ... **NO**
- It is still easy to break this scheme using frequency cryptanalysis - the frequency of encrypted characters in the ciphertext is used to derive the plaintext

Mono-alphabetic Substitution Cipher (3)

- The attack can be carried based on the following facts (assuming English language text)
 - 1 Given any key, the mapping for each letter is fixed. For example, if a is mapped to q , then every a in the plaintext is mapped to q
 - 2 The frequency distribution of individual letters in English is known
- Count the frequencies of each letter in the ciphertext and compare to the English letter frequency distribution
- Additional knowledge can also be used (e.g., h is likely to appear between t and e)

Poly-alphabetic Substitution Cipher

- Key is defined on blocks of plaintext characters
- **Example:** for plaintext *abac* and 2-character block, the ciphertext might be *DZTY*
- Here we have two blocks *ab* and *ac*
- Observe that *a* is not mapped to a fixed ciphertext character

Vigenere Cipher

- A poly-alphabetic shift cipher
- Applies several independent instances of shift cipher in sequence
- **Example 1:**
Plaintext: attackatdawn
Key: lemonlemonle
Ciphertext: lxfopefrnhr
- **Example 2:**
Plaintext: tellhimaboutme
Key: cafecafecafeca
Ciphertext: ??

How can we define security

Two Important components:

- **Security Guarantee**: Defines what the scheme is intended to protect from the attacker
- **Threat Model**: Defines the attacker's capability

Security Guarantee - Some Thoughts

What is a good security guarantee??

- Impossible for an attacker to recover the key
- Impossible for an attacker to recover the entire plaintext
- Impossible for an attacker to recover any character of the plaintext
- **Right Answer:** Attacker should not know any information about the plaintext other than what he has already known

Threat Models

In order of increasing power of attacker:

- Ciphertext-only attack
- Known-plaintext attack
- Chosen-plaintext attack
- Chosen-ciphertext attack

Some Tips

- Never use your own or non-standardized crypto algorithms
- Understand the application requirements and use the existing schemes that are mathematically strong, provably secure and widely used

Summary

- Importance of Cryptography
- Data Confidentiality + Integrity + Authentication
- Kerckhoff's principle and its broad impact
- Historical Ciphers and their cryptanalysis
- Security goals and threat models

Useful References

- Chapter 1, Introduction to Modern Cryptography by Jonathan Katz and Yehuda Lindell, 2nd Edition, CRC Press, 2015.
- `http://cseweb.ucsd.edu/~mihir/crypto-links.html`