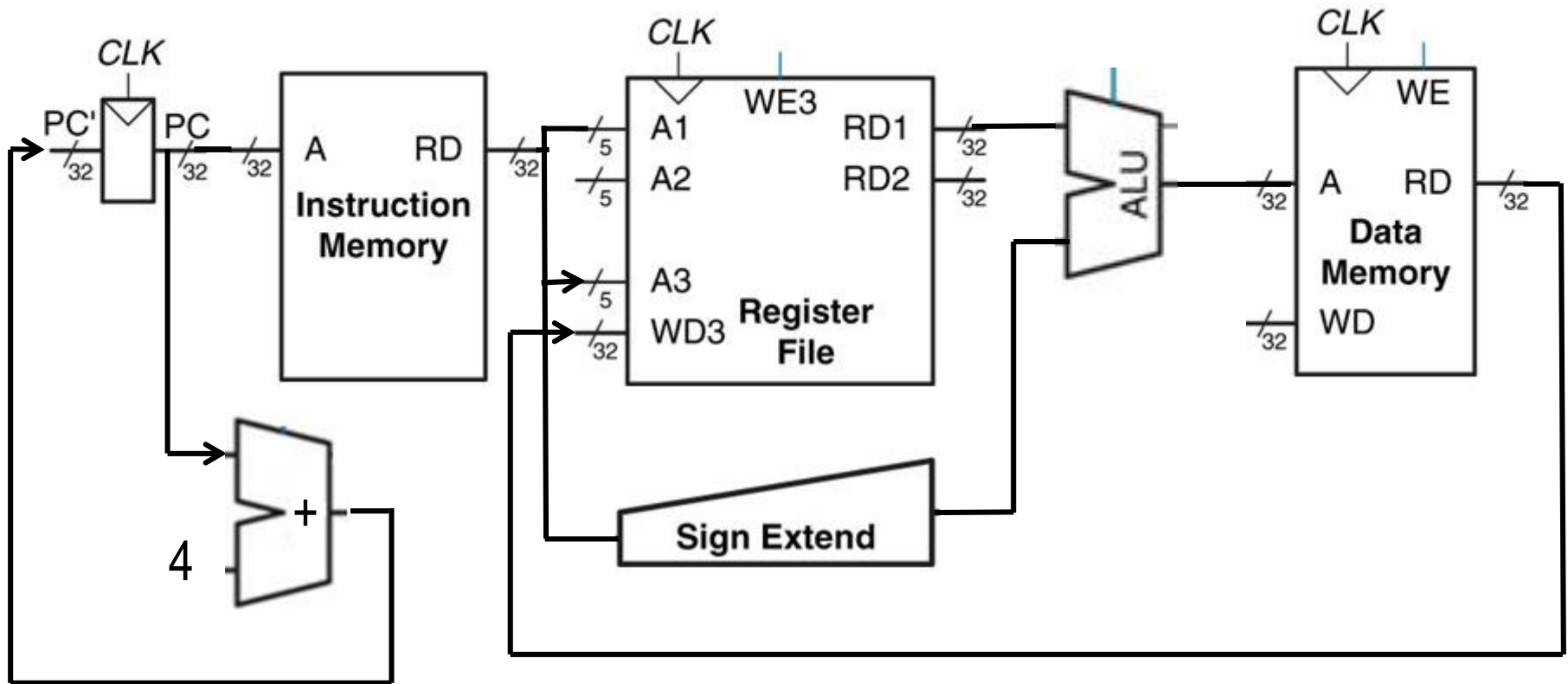


MIPS Assembly Programming

**The Language of the
Electronic Computer**

32-bit RegisterFile + ALU

MIPS

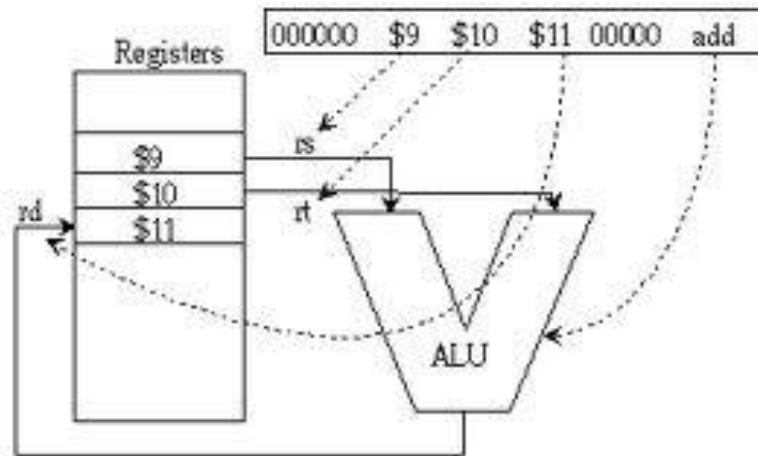


Why Assembly?

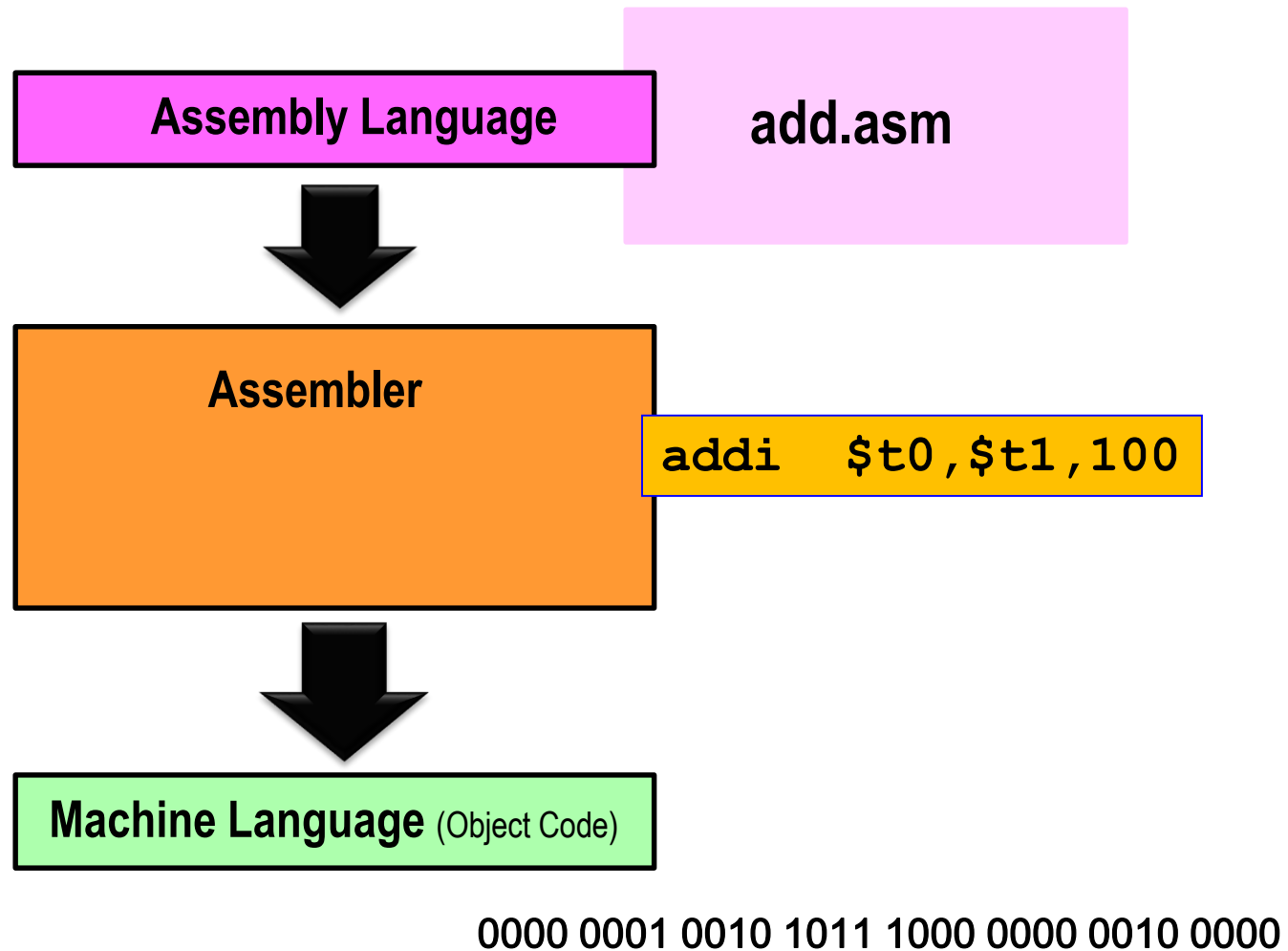
- **Assembly is widely used in industry:**
 - Embedded systems
 - Real time systems
 - Low level and direct access to hardware
- **Assembly is widely used not in industry:**
 - Cracking software protections: patching, patch-loaders and emulators
 - Hacking into computer systems: buffer under/overflows, worms and Trojans.

Assembly-Machine Language

- Each assembly language is specific to a particular computer architecture
- Each computer architecture has its own machine language.



Assembly-Machine Language



MIPS Architecture

- MIPS is a register-to-register, or load/store, architecture
- The destination and sources must **all** be registers
- Special instructions are needed to access the main memory.

MIPS: Register File

MIPS processors have 32-registers, each of which holds a 32-bit value

- Register addresses are 5-bits ($2^5 = 32$ -bits) long
- The data inputs and outputs are 32-bits wide.

MIPS register names convention

- MIPS register names begin with a dollar sign → \$

1. By number:

\$0, \$1, ..., \$31

2. By a letter and a number:

\$a0-\$a2 ...

MIPS registers

Registers	Coproc 1	Coproc 0
Name	Number	Value
\$zero	0	0
\$at	1	0
\$v0	2	0
\$v1	3	0
\$a0	4	0
\$a1	5	0
\$a2	6	0
\$a3	7	0
\$t0	8	0
\$t1	9	0
\$t2	10	0
\$t3	11	0
\$t4	12	0
\$t5	13	0
\$t6	14	0
\$t7	15	0
\$s0	16	0
\$s1	17	0
\$s2	18	0
\$s3	19	0
\$s4	20	0
\$s5	21	0
\$s6	22	0
\$s7	23	0
\$t8	24	0
\$t9	25	0
\$k0	26	0
\$k1	27	0
\$gp	28	268468224
\$sp	29	2147479548
\$fp	30	0
\$ra	31	0
pc		4194304
hi		0
lo		0

MIPS registers

Name	Number	Use	Preserved across a call?
\$zero	0	The constant value 0	N.A.
\$at	1	Assembler temporary	No
\$v0–\$v1	2–3	Values for function results and expression evaluation	No
\$a0–\$a3	4–7	Arguments	No
\$t0–\$t7	8–15	Temporaries	No
\$s0–\$s7	16–23	Saved temporaries	Yes
\$t8–\$t9	24–25	Temporaries	No
\$k0–\$k1	26–27	Reserved for OS kernel	No
\$gp	28	Global pointer	Yes
\$sp	29	Stack pointer	Yes
\$fp	30	Frame pointer	Yes
\$ra	31	Return address	Yes

Figure 1.4 MIPS registers and usage conventions. In addition to the 32 general-purpose registers (R0–R31), MIPS has 32 floating-point registers (F0–F31) that can hold either a 32-bit single-precision number or a 64-bit double-precision number.

Our first assembly demo program

```
1.asm
1 # Folder L1/1.asm
2
3     .text
4     .globl main
5 main:
6     li      $t0, 2
7     li      $t1, 3
8     add     $t5, $t1, $t0
9     li      $v0, 10
10    syscall
```

System Calls: 10

Registers	Coproc 1	Coproc 0
Name	Number	Value
\$zero	0	0
\$at	1	0
\$v0	2	0
\$v1	3	0
\$a0	4	0
\$a1	5	0
\$a2	6	0
\$a3	7	0
\$t0	8	0
\$t1	9	0
\$t2	10	0
\$t3	11	0
\$t4	12	0
\$t5	13	0
\$t6	14	0
\$t7	15	0
\$s0	16	0
\$s1	17	0
\$s2	18	0
\$s3	19	0
\$s4	20	0
\$s5	21	0
\$s6	22	0
\$s7	23	0
\$t8	24	0
\$t9	25	0
\$k0	26	0
\$k1	27	0
\$gp	28	268468224
\$sp	29	2147479548
\$fp	30	0
\$ra	31	0
pc		4194304
hi		0
lo		0

Our first assembly demo program

The diagram shows an assembly program with the following code and annotations:

```
1 # Folder L1/1.asm
2
3 .text
4 .globl main
5 main:
6     li    $t0, 2
7     li    $t1, 3
8     add   $t5, $t1, $t0
9     li    $v0, 10
10    syscall
```

Annotations:

- Comments:** Points to the green comment lines starting with `#`.
- Assembly directives:** Points to the pink directives `.text` and `.globl main`.
- Label:** Points to the label `main:` on line 5.
- Opcode:** Points to the instruction `li` on line 6.
- Operand:** Points to the register and value `$t0, 2` on line 6.

Additional comments on the right side of the code:

- `#` Informs the assembler that instructions follow
- `#` Declare as global the label main
- `#` Execution starts at main:
- `#` \$t0 = 2
- `#` \$t1 = 3
- `#` \$t5 = \$t1 + \$t0
- `#` System call for exit (Load code 10)
- `#` Call operating system to perform operation (exit)

`li` is a pseudo-instruction; will talk about it in the next lecture

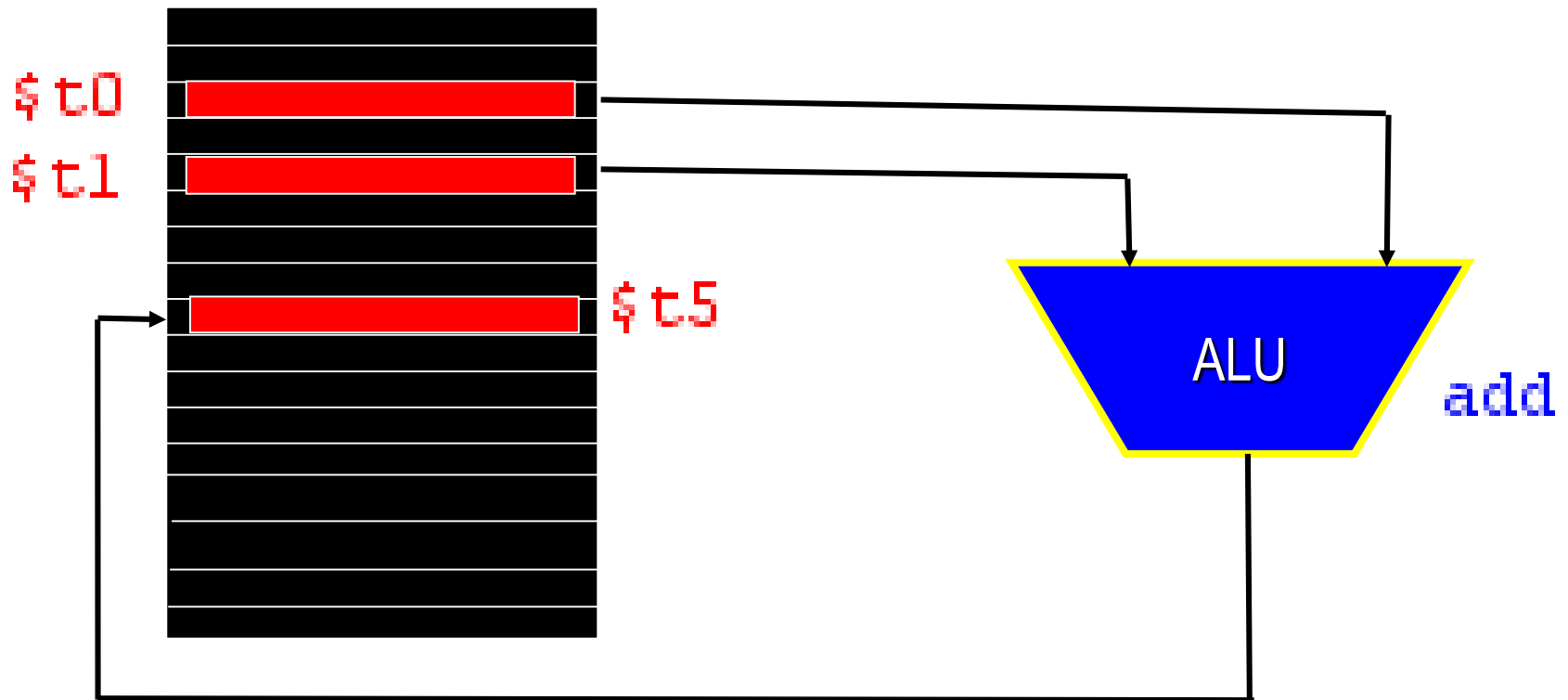
Assemble ... GO

\$t0	8	2	←
\$t1	9	3	←
\$t2	10	0	
\$t3	11	0	
\$t4	12	0	
\$t5	13	5	←
\$t6	14	0	
\$t7	15	0	

Registers	Coproc 1	Coproc 0
Name	Number	Value
\$zero	0	0
\$at	1	0
\$v0	2	10
\$v1	3	0
\$a0	4	0
\$a1	5	0
\$a2	6	0
\$a3	7	0
\$t0	8	2
\$t1	9	3
\$t2	10	0
\$t3	11	0
\$t4	12	0
\$t5	13	5
\$t6	14	0
\$t7	15	0
\$s0	16	0
\$s1	17	0
\$s2	18	0
\$s3	19	0
\$s4	20	0
\$s5	21	0
\$s6	22	0
\$s7	23	0
\$t8	24	0
\$t9	25	0
\$k0	26	0
\$k1	27	0
\$gp	28	268468224
\$sp	29	2147479548
\$fp	30	0
\$ra	31	0
pc		4194324
hi		0
lo		0

[32] 32-bit RegisterFile + ALU

add \$t5, \$t1, \$t0



Assembler COMMENTS and LABELS

- **Comments:** Text following a '#' (sharp) to the end of the line is ignored
- **Labels:** Are symbols that represent memory addresses
 - Labels take on the values of the address where they are declared
 - Labels declarations appear at the beginning of a line, and are terminated by a colon.

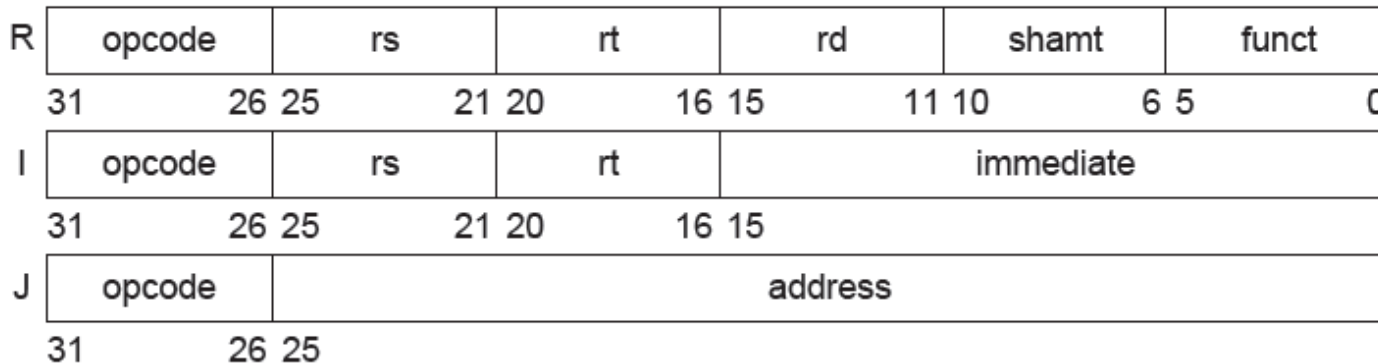
Instructions are divided into three kinds of format

- (**R**, **I** and **J** format)
 - **Register** arithmetic instructions (**R**-format)
 - Memory **Immediate** load and store (**I**-format)
 - Branching and **Jump** instructions (**J**-format).

MIPS instruction format



Basic instruction formats



Floating-point instruction formats

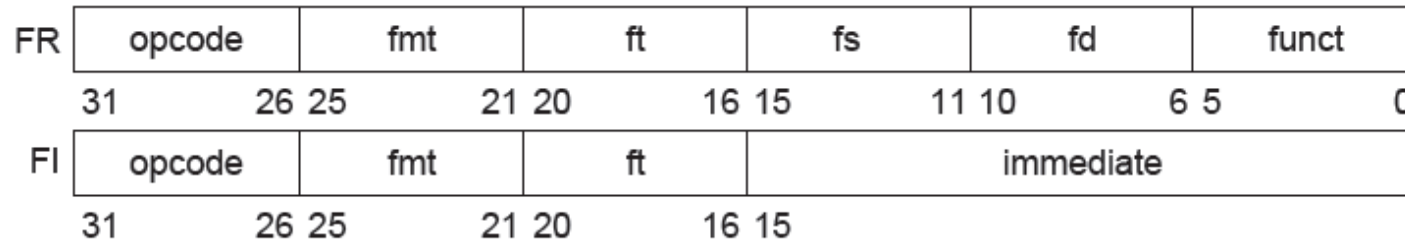


Figure 1.6 MIPS64 instruction set architecture formats. All instructions are 32 bits long. The R format is for integer register-to-register operations, such as DADDU, DSUBU, and so on. The I format is for data transfers, branches, and immediate instructions, such as LD, SD, BEQZ, and DADDIs. The J format is for jumps, the FR format for floating-point operations, and the FI format for floating-point branches.

MIPS instruction format

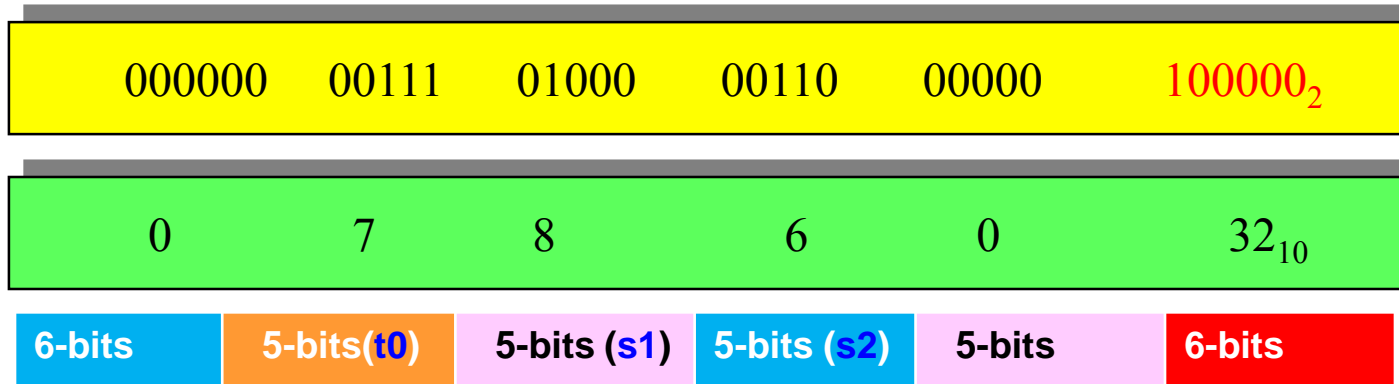
Format	6 bits	5 bits	5 bits	5 bits	5 bits	6 bits	Comments
R	op	rs	rt	rd	shamt	funct	Arithmetic
I	op	rs	rt	address/immediate			Transfer, branch,immediate
J	op	target address					Jump

- **op**: basic operation of instruction
- **funct**: variant of instruction
- **rs**: first register source operand
- **rt**: second register source operand
- **rd**: register destination operand
- **shamt**: shift amount

R-Format example

```
add $t0, $s1, $s2
```

add \$t0, \$s1, \$s2 < R-Format



- **rs** = 8₁₀: first source operand is: **\$s1**
- **rt** = 6₁₀: second source operand is: **\$s2**
- **rd** = 7₁₀: register destination is: **\$t0**
- **funct** = 32 (**add**)

MIPS **Instruction** set

- Arithmetic, Logic, and Shifting Instructions
- Conditional Branch Instructions
- Load and Store Instructions
- Function Call Instructions

Arithmetic

```
add $t5, $t1, $t0
```

Our first assembly demo program

```
1 # Folder L1/1.asm
2
3     .text
4     .globl main
5 main:
6     li      $t0, 2
7     li      $t1, 3
8     add     $t5, $t1, $t0
9     li      $v0, 10
10    syscall
```

li

add

li




syscall

`.text` and `.globl` directives

- `.text` directive
 - Defines the section of a program containing instructions
- `.globl main`
 - Declares `main` as global
- `main:` label that represents a memory address.

```
main:      .text
           .globl main
           li      $t0, 2
           li      $t1, 3
           add     $t5, $t1, $t0
           li      $v0, 10
           syscall
```


A Breakdown of Segment and Linker Directives

Name	Parameters	Description
 .data	<i>addr</i>	The following items are to be assembled into the data segment. By default, begin at the next available address in the data segment. If the optional argument <i>addr</i> is present, then begin at <i>addr</i> .
 .text	<i>addr</i>	The following items are to be assembled into the text segment. By default, begin at the next available address in the text segment. If the optional argument <i>addr</i> is present, then begin at <i>addr</i> . In SPIM, the only items that can be assembled into the text segment are instructions and words (via the <code>.word</code> directive).
.kdata	<i>addr</i>	The kernel data segment. Like the data segment, but used by the Operating System.
.ktext	<i>addr</i>	The kernel text segment. Like the text segment, but used by the Operating System.
.extern	<i>sym size</i>	Declare as global the label <i>sym</i> , and declare that it is <i>size</i> bytes in length (this information can be used by the assembler).
 .globl	<i>sym</i>	Declare as global the label <i>sym</i> .

li *des, const* # load the constant *const* into *des*

li *\$t0, 2*

li *\$t1, 3*

Op	Operands	Description
○ la	<i>des, addr</i>	Load the address of a label.
lb(u)	<i>des, addr</i>	Load the byte at <i>addr</i> into <i>des</i> .
lh(u)	<i>des, addr</i>	Load the halfword at <i>addr</i> into <i>des</i> .
→ ○ li	<i>des, const</i>	Load the constant <i>const</i> into <i>des</i> .
lui	<i>des, const</i>	Load the constant <i>const</i> into the upper halfword of <i>des</i> , and set the lower halfword of <i>des</i> to 0.
lw	<i>des, addr</i>	Load the word at <i>addr</i> into <i>des</i> .
lwl	<i>des, addr</i>	
lwr	<i>des, addr</i>	
○ ulh(u)	<i>des, addr</i>	Load the halfword starting at the (possibly unaligned) address <i>addr</i> into <i>des</i> .
○ ulw	<i>des, addr</i>	Load the word starting at the (possibly unaligned) address <i>addr</i> into <i>des</i> .

addu ← new instruction

add instruction

add/addu instructions

- **add**; signed addition
- **addu**; **u**nsigned addition.


4.4.1 Arithmetic Instructions

Op	Operands	Description
◦ <code>abs</code>	<i>des, src1</i>	<i>des</i> gets the absolute value of <i>src1</i> .
→ <code>add(u)</code>	<i>des, src1, src2</i>	<i>des</i> gets <i>src1</i> + <i>src2</i> .
<code>and</code>	<i>des, src1, src2</i>	<i>des</i> gets the bitwise and of <i>src1</i> and <i>src2</i> .
<code>div(u)</code>	<i>src1, reg2</i>	Divide <i>src1</i> by <i>reg2</i> , leaving the quotient in register <i>lo</i> and the remainder in register <i>hi</i> .
◦ <code>div(u)</code>	<i>des, src1, src2</i>	<i>des</i> gets <i>src1</i> / <i>src2</i> .
◦ <code>mul</code>	<i>des, src1, src2</i>	<i>des</i> gets <i>src1</i> × <i>src2</i> .
◦ <code>mulo</code>	<i>des, src1, src2</i>	<i>des</i> gets <i>src1</i> × <i>src2</i> , with overflow.
<code>mult(u)</code>	<i>src1, reg2</i>	Multiply <i>src1</i> and <i>reg2</i> , leaving the low-order word in register <i>lo</i> and the high-order word in register <i>hi</i> .
◦ <code>neg(u)</code>	<i>des, src1</i>	<i>des</i> gets the negative of <i>src1</i> .
<code>nor</code>	<i>des, src1, src2</i>	<i>des</i> gets the bitwise logical nor of <i>src1</i> and <i>src2</i> .
◦ <code>not</code>	<i>des, src1</i>	<i>des</i> gets the bitwise logical negation of <i>src1</i> .
<code>or</code>	<i>des, src1, src2</i>	<i>des</i> gets the bitwise logical or of <i>src1</i> and <i>src2</i> .
◦ <code>rem(u)</code>	<i>des, src1, src2</i>	<i>des</i> gets the remainder of dividing <i>src1</i> by <i>src2</i> .
◦ <code>rol</code>	<i>des, src1, src2</i>	<i>des</i> gets the result of rotating left the contents of <i>src1</i> by <i>src2</i> bits.
◦ <code>ror</code>	<i>des, src1, src2</i>	<i>des</i> gets the result of rotating right the contents of <i>src1</i> by <i>src2</i> bits.
<code>sll</code>	<i>des, src1, src2</i>	<i>des</i> gets <i>src1</i> shifted left by <i>src2</i> bits.
<code>sra</code>	<i>des, src1, src2</i>	Right shift arithmetic.
<code>srl</code>	<i>des, src1, src2</i>	Right shift logical.
<code>sub(u)</code>	<i>des, src1, src2</i>	<i>des</i> gets <i>src1</i> - <i>src2</i> .
<code>xor</code>	<i>des, src1, src2</i>	<i>des</i> gets the bitwise exclusive or of <i>src1</i> and <i>src2</i> .

MIPS: Three-address instructions

- MIPS uses three-address instructions for data manipulation
- Each ALU instruction contains a destination and two sources
- For example, an addition instruction ($a = b + c$) has the form.

add **\$t5** , **\$t1** , **\$t0**

The diagram shows the instruction format for the MIPS 'add' instruction. The word 'add' is in blue. The destination register '\$t5' is in red and is positioned above a green box containing the number '1'. The first source register '\$t1' is in red and is positioned above a blue box containing the number '2'. The second source register '\$t0' is in red and is positioned above a blue box containing the number '3'. Commas separate the three components.

add signed addition

add \$t5, \$t1, \$t0

- Performs the Binary Addition algorithm on two 32-bits;
 - Signed Binary
- add *des, src1, src2* # des gets src1 + src2
- add \$t5, \$t1, \$t0 # \$t5: 4 = 1 + 3
- Three registers (\$t5, \$t1, \$t0) are involved
- Overflow **trap** is possible

$$\begin{array}{r} 1001 \ 1001 \\ + 0111 \ 0001 \\ \hline 10001 \ 1010 \end{array}$$

- A **trap** is an interruption in the normal machine cycle.

addu; **u**nsigned addition

- Performs the Binary Addition Algorithm on two 32-bits ;
 - Unsigned Binary
 - Two's Complement
- The destination register can be the same as one of the source registers
- **add(u)** *des, src1, src2* # *des* gets *src1* + *src2*
- **add(u)** *\$t0, \$t0, \$t1* # *\$t0* = *\$t0* + *\$t1*
- **addu** ... ignores overflow trap.

add ... addu

- The **add** instruction it is used in the cases that overflow is an important factor. Otherwise we use the **addu** instruction
 - For **signed** numbers, use **add**
 - For **unsigned** numbers, use **addu**
- ... and the **addi** instruction ...

addi ← new instruction

`addi` (add immediate) > `addi $r1, $r2, 4`

MIPS32; `addi` ... add immediate

`addi`

`$r1`

`$r2`

`4`

addi (add immediate) > **addi** \$r1, \$r2, 4

MIPS32; addi ... add immediate

addi	\$r1	\$r2	4
001000	00001	00010	0000000000000000100

addi (add immediate) > **addi** \$r1, \$r2, 4

MIPS32; addi ... add immediate			
addi	\$r1	\$r2	4
001000	00001	00010	000000000000000100
Op Code	rt	rs	Immediate value 4

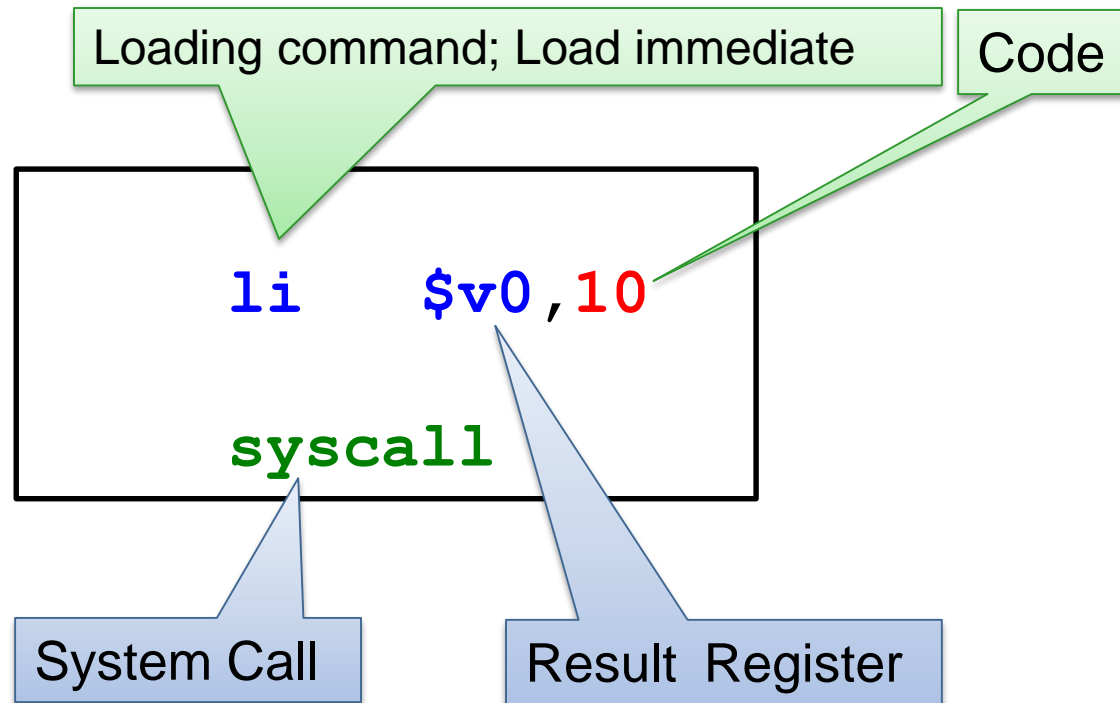
li (**l**oad **i**mmEDIATE)

Another immediate instruction

li (load immediate)

- `li $v0, 10` # load immediate `$v0` = 10
- `Syscall` # instruction; action depends on
Code loaded in the register: `$v0`

Load 10 into \$v0; ... terminate



- A system call starts off by loading a specific code into the Result Register
- Then, the `syscall` instruction is called. The final result depends on the code loaded into the Result register
- The above example is the exit `syscall`, since loading the code "10" and calling the "`syscall`" instruction terminates the program.

addi

2.asm*

```
1  # Folder L1/2.asm
2
3      .text
4      .globl main
5 main:
6      li      $t0, 2
7      addi    $t5, $t0, 3
8      li      $v0, 10
9      syscall
10
```


addi

2.asm*

```
1 # Folder L1/2.asm
2
3     .text                # Informs the assembler that instructions follow
4     .globl main          # Declare as global the label main
5 main:                    # Execution starts at main:
6     li    $t0, 2         # $t0 = 2
7     addi   $t5, $t0, 3    # $t5 = $t0 + 3
8     li    $v0, 10        # System call for exit (Load code 10)
9     syscall              # Call operating system to perform operation (exit)
10
```

Assemble ... GO

```
1 # Folder L1/2.asm
2
3     .text
4     .globl main
5 main:
6     li      $t0, 2
7     addi    $t5, $t0, 3
8     li      $v0, 10
9     syscall
10
```



\$t0	8	2
\$t1	9	0
\$t2	10	0
\$t3	11	0
\$t4	12	0
\$t5	13	5
\$t6	14	0
\$t7	15	0

addi ...add *i*mmEDIATE *u*nsigned

- `addi des, src1, const` # *des* gets *src1* + *const*
- `addi $t0, $t0, 1` # *\$t0* = *\$t0* + 1

(Overflow trap)

- Sign Extension (done)

addiu

← new instruction

Addiu ← add immediate unsigned

- `addiu des, src1, const` # *des* gets *src1* + *const*
- `addiu $t0, $t0, 1` # *\$t0* = *\$t0* + 1

(No overflow trap)

Add three numbers ...

- $1 + 3 + 4$
- To add the numbers, just use the instruction: **add**



In Class

5 minutes

Program/Solution


3.asm*

```
1 # Folder L1/3.asm
2 # Andrew De Stefano and Chris Crockett, 2012
3 # This program adds three numbers.
4
5     .text
6     .globl main
7 main:
8     li    $t0,1           # $t0 = 1
9     li    $t1,3           # $t1 = 3
10    li    $t2,4           # $t2 = 4
11    add    $t6,$t1,$t0     # $t6 = 4
12    add    $t7,$t6,$t2     # $t7 = 8
13    li    $v0,10          # System call for exit (Load code 10)
14    syscall               # Call operating system to perform operation (exit)
```


Assemble ... GO



Registers	Coproc 1	Coproc 0
Name	Number	Value
\$zero	0	0
\$at	1	0
\$v0	2	10
\$v1	3	0
\$a0	4	0
\$a1	5	0
\$a2	6	0
\$a3	7	0
\$t0	8	1
\$t1	9	3
\$t2	10	4
\$t3	11	0
\$t4	12	0
\$t5	13	0
\$t6	14	4
\$t7	15	8



“More ...”

System **C**alls

System Calls (`syscall`)

- MIPS programs can make system calls **by placing parameters in specified registers**, depending on the call, and executing a code instruction.
- Returned results are made available in other specified registers, also depending on the call.

Syscalls

Service	Code	Arguments	Result
print_int	1	\$a0	<i>none</i>
print_float	2	\$f12	<i>none</i>
print_double	3	\$f12	<i>none</i>
print_string	4	\$a0	<i>none</i>
read_int	5	<i>none</i>	\$v0
read_float	6	<i>none</i>	\$f0
read_double	7	<i>none</i>	\$f0
read_string	8	\$a0 (address), \$a1 (length)	<i>none</i>
sbrk	9	\$a0 (length)	\$v0
exit	10	<i>none</i>	<i>none</i>

- "Service" explains the function of the `syscall` code
- "Code" is the number to be loaded
- "Arguments" states the arguments used and where specifically they'd be located
- "Result" explains the output

How to use `syscall` system services

1. Load the service number in register `$v0`
2. Load argument values, if any, in registers:
`$a0, $a1, $a2`)
3. Issue the `syscall` instruction
4. Retrieve return values, if any, from the result registers.

la ← new instruction

load the **a**ddress

Syscall → 4

Print a string of text

Hello world

System Calls: 4,10

folder:

Assembly directive

start:

.text

la \$a0, msg

li \$v0, 4

syscall

li \$v0, 10

syscall

Assembly directive

.data

.asciiz "hello world! "

Assembly directive

msg:

label

The following are to be assembled in to text segment

Load the address of the message text

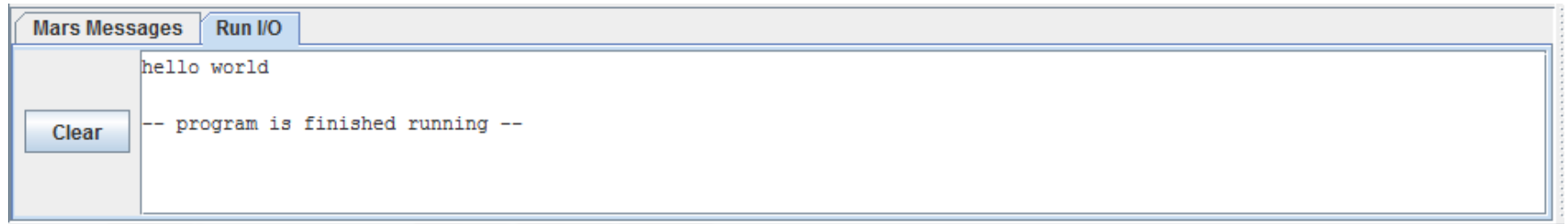
Load the syscall (4) code for printing the string of text

Load the syscall (10) code for exiting



Informs the assembler that data needed within instructions follows

la is a pseudo-instruction; will talk about it in the next lecture


Assemble ... GO



Data Directives

Name	Parameters	Description
 .data	<i>addr</i>	The following items are to be assembled into the data segment. By default, begin at the next available address in the data segment. If the optional argument <i>addr</i> is present, then begin at <i>addr</i> .
 .text	<i>addr</i>	The following items are to be assembled into the text segment. By default, begin at the next available address in the text segment. If the optional argument <i>addr</i> is present, then begin at <i>addr</i> . In SPIM, the only items that can be assembled into the text segment are instructions and words (via the <code>.word</code> directive).
.kdata	<i>addr</i>	The kernel data segment. Like the data segment, but used by the Operating System.
.ktext	<i>addr</i>	The kernel text segment. Like the text segment, but used by the Operating System.
.extern	<i>sym size</i>	Declare as global the label <i>sym</i> , and declare that it is <i>size</i> bytes in length (this information can be used by the assembler).
.globl	<i>sym</i>	Declare as global the label <i>sym</i> .

Data Directives



Name	Parameters	Description
<code>.align</code>	<i>n</i>	Align the next item on the next 2^n -byte boundary. <code>.align 0</code> turns off automatic alignment.
<code>.ascii</code>	<i>str</i>	Assemble the given string in memory. Do not null-terminate.
<code>.asciiiz</code>	<i>str</i>	Assemble the given string in memory. Do null-terminate.
<code>.byte</code>	<i>byte1</i> ... <i>byteN</i>	Assemble the given bytes (8-bit integers).
<code>.half</code>	<i>half1</i> ... <i>halfN</i>	Assemble the given halfwords (16-bit integers).
<code>.space</code>	<i>size</i>	Allocate <i>n</i> bytes of space in the current segment. In SPIM, this is only permitted in the data segment.
<code>.word</code>	<i>word1</i> ... <i>wordN</i>	Assemble the given words (32-bit integers).

la (load address)

- `la $a0, msg` # load address of string to be printed into `$a0`

```
start:  .text                                # The following are to be assembled in to text segment
        la      $a0, msg                    # Load the address of the message text
        li      $v0, 4                      # Load the syscall (4) code for printing the string of text
        syscall
        li      $v0, 10                     # Load the syscall (10) code for exiting
        syscall

msg:    .data                                # Informs the assembler that data needed within instructions follows
        .ascii  "hello world! "
```

Syscall → 1

Print-out Integer

Prints “out” the result

System Calls: 4,1,10

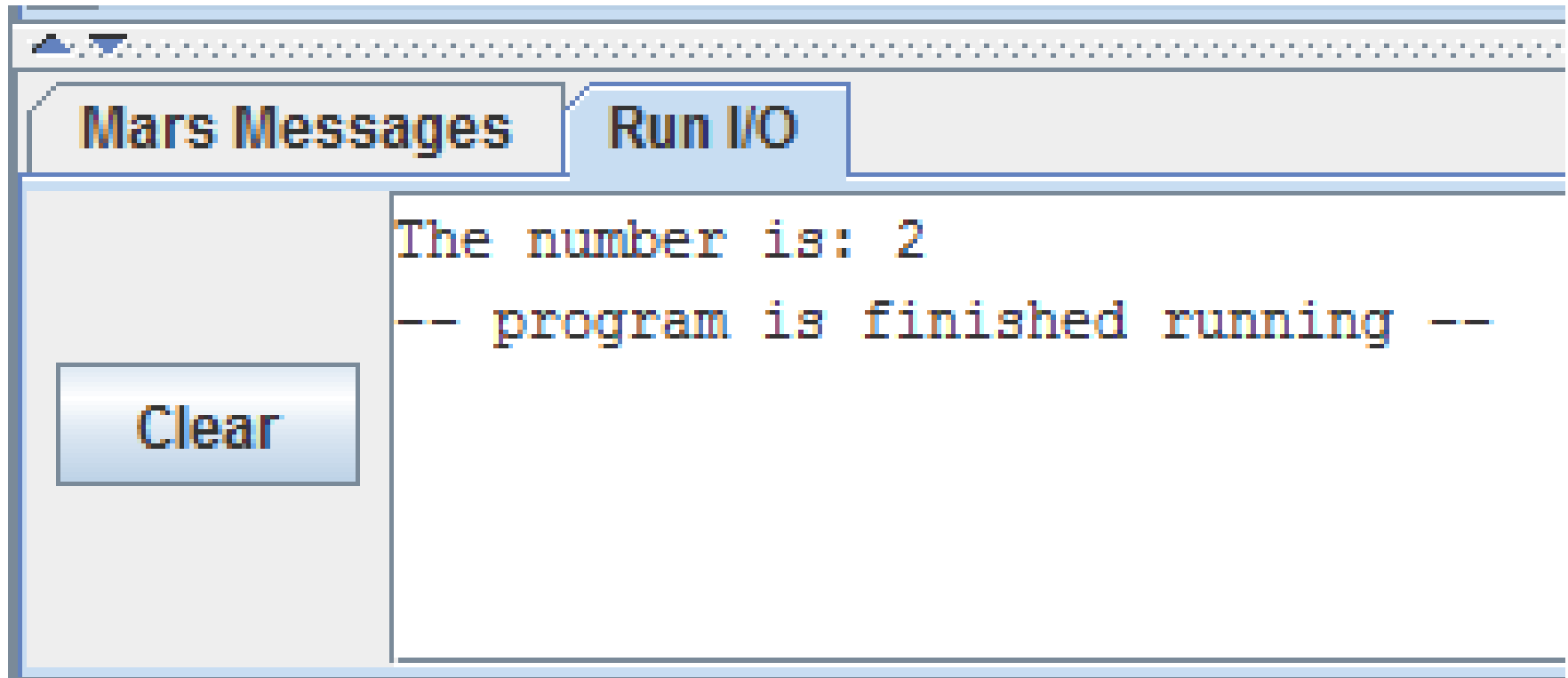
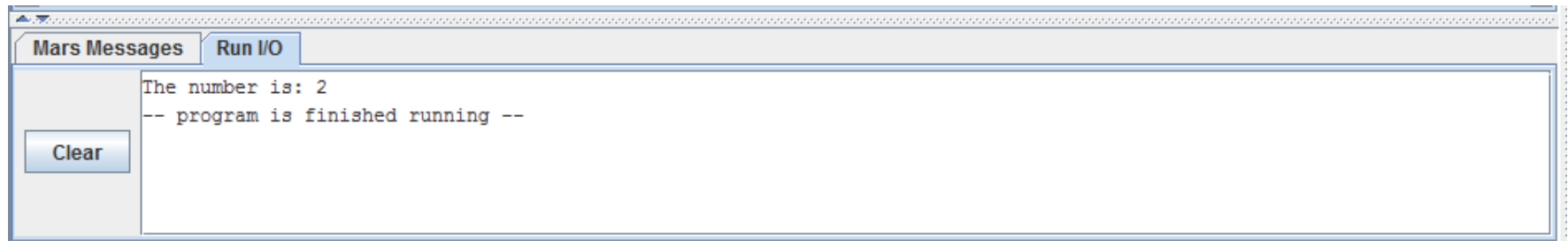
```
1 # Folder L1\4.asm
2 # Prints-"out"
3
4     .text
5     .globl main
6 main:
7     la      $a0, output
8     li      $v0, 4
9     syscall
10
11     li      $t0, 2
12     move    $a0, $t0  # The contents of $t0 are to be copied into register $a0
13     li      $v0, 1
14     syscall
15
16     li      $v0, 10
17     syscall
18
19     .data
20 output:
21     .asciiz "The number is: "
22
```

Prints “out” the result

```
1 # Folder L1\4.asm
2 # Prints-"out"
3
4     .text
5     .globl main
6 main:
7     la    $a0, output      # load address of string to be printed into $a0
8     li    $v0,4            # System call for printing string (code = 4)
9     syscall               # Call operating system to perform operation (Print string)
10
11     li    $t0,2           # $t0 = 2
12     move  $a0,$t0         # The contents of $t0 are to be copied into register $a0
13     li    $v0,1          # System call for printing integer (code = 1)
14     syscall               # Call operating system to perform operation (Print integer)
15
16     li    $v0,10          # System call for exit (code = 10)
17     syscall               # Call operating system to perform operation exit
18
19     .data                 # Directive; Informs the assembler that data needed within instructions follows
20 output:                  # Label (output)
21     .ascii "The number is: " # Declaration for string variable (directive makes string null terminated)
22
```

move is a pseudo-instruction; will talk about it in the next lecture

Assemble ... GO; **The number is: 2**



Example

System Calls: 4,4,1,10

5.asm

```
1  # Folder: L1/5.asm
2  # Andrew De Stefano and Chris Crockett, 2012
3
4      .text
5      .globl main
6  main:
7      la      $a0,prompt1      # Print string
8      li      $v0,4            #
9      syscall                  #
10     li      $t0,3            # $t0=3
11     li      $t1,5            # $t1=5
12     li      $t2,2            # $t2=2
13     addu    $t0,$t0,$t1      # $t0=
14     addu    $t0,$t0,$t2      # $t0=
15     la      $a0,prompt2      # Print string
16     li      $v0,4            #
17     syscall                  #
18     move    $a0,$t0          #
19     li      $v0,1            #
20     syscall                  #
21     li      $v0,10           #
22     syscall
23
24     .data
25 prompt1:
26     .asciiz "Sum of three numbers. "
27 prompt2:
28     .asciiz "Sum =  "
29
```

Assemble ...

D:\Courses\ALL COURSES\CMPT280-2014\LECTURES+HW\L7-ExamSolution(HW3)\DemoCode\L2-arithmetic\5.asm - MARS 4.4

File Edit Run Settings Tools Help

Run speed at max (no interaction)

Edit Execute

Text Segment

Bkpt	Address	Code	Basic	Source
	4194304	0x3c011001	lui \$1,4097	7: la \$a0,prompt1 # Print string
	4194308	0x34240000	ori \$4,\$1,0	
	4194312	0x24020004	addiu \$2,\$0,4	8: li \$v0,4 #
	4194316	0x0000000c	syscall	9: syscall #
	4194320	0x24080003	addiu \$8,\$0,3	10: li \$t0,3 # \$t0=3
	4194324	0x24090005	addiu \$9,\$0,5	11: li \$t1,5 # \$t1=5
	4194328	0x240a0002	addiu \$10,\$0,2	12: li \$t2,2 # \$t2=3
	4194332	0x01094021	addu \$8,\$8,\$9	13: addu \$t0,\$t0,\$t1 # \$t0=8
	4194336	0x010a4021	addu \$8,\$8,\$10	14: addu \$t0,\$t0,\$t2 # \$t0=10
	4194340	0x3c011001	lui \$1,4097	15: la \$a0,prompt2 # Print string
	4194344	0x34240017	ori \$4,\$1,23	

Data Segment

Address	Value (+0)	Value (+4)	Value (+8)	Value (+12)	Value (+16)	Value (+20)	Value (+24)	Value (+28)
268500992	544044371	1948280431	1701147240	1836412448	1936876898	1392511278	1025535349	2105376
268501024	0	0	0	0	0	0	0	0
268501056	0	0	0	0	0	0	0	0
268501088	0	0	0	0	0	0	0	0
268501120	0	0	0	0	0	0	0	0
268501152	0	0	0	0	0	0	0	0
268501184	0	0	0	0	0	0	0	0
268501216	0	0	0	0	0	0	0	0
268501248	0	0	0	0	0	0	0	0

0x10010000 (.data) Hexadecimal Addresses Hexadecimal Values ASCII

Mars Messages Run I/O

Assemble: assembling D:\Courses\ALL COURSES\CMPT280-2014\LECTURES+HW\L7-ExamSolution(HW3)\DemoCode\L2-arithmetic\5.asm

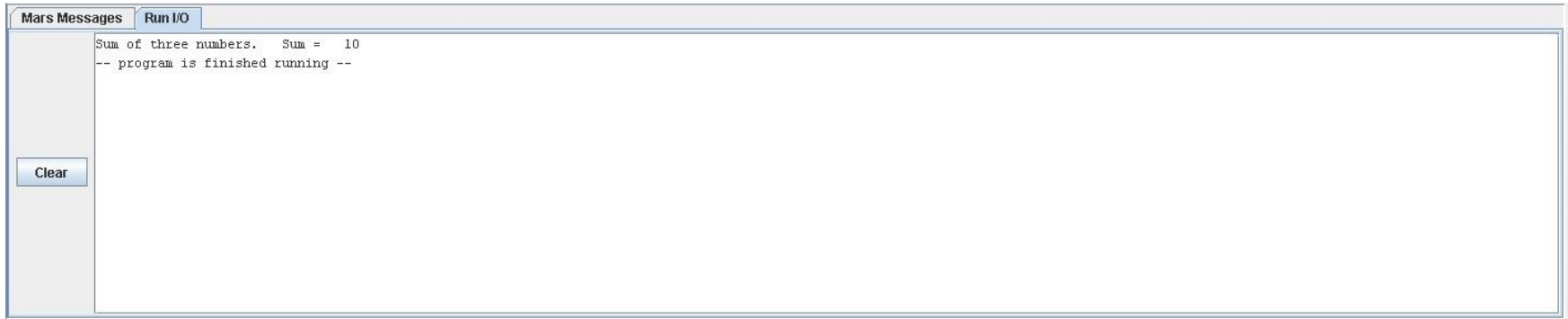
Clear Assemble: operation completed successfully.

Registers Coproc 1 Coproc 0

Name	Number	Value
\$zero	0	0
\$at	1	0
\$v0	2	0
\$v1	3	0
\$a0	4	0
\$a1	5	0
\$a2	6	0
\$a3	7	0
\$t0	8	0
\$t1	9	0
\$t2	10	0
\$t3	11	0
\$t4	12	0
\$t5	13	0
\$t6	14	0
\$t7	15	0
\$s0	16	0
\$s1	17	0
\$s2	18	0
\$s3	19	0
\$s4	20	0
\$s5	21	0
\$s6	22	0
\$s7	23	0
\$t8	24	0
\$t9	25	0
\$k0	26	0
\$k1	27	0
\$gp	28	268468224
\$sp	29	2147479548
\$fp	30	0
\$ra	31	0
pc		4194304
hi		0
lo		0

EN 9:22 AM 11/1/2014

GO (result)



```
Sum of three numbers.    Sum =    10
-- program is finished running --
```

Syscall → 5

Read Integer from the Command line

Read and Print integer from command line

System Calls: 5,1,10

```
1 # Folder: L1/6.asm
2 # Read and Print integer f
3
4
5     .text
6     .globl main
7 main:
8     li     $v0, 5
9     syscall
10
11     move   $a0, $v0
12     li     $v0, 1
13     syscall
14
15     li     $v0, 10
16     syscall
```

Read and Print integer from command line

```
1 # Folder: L1/6.asm
2 # Read and Print integer from command line
3
4
5     .text
6     .globl main
7 main:
8     li    $v0, 5          # syscall for reading integer from the command line (code = 5)
9     syscall
10
11     move  $a0, $v0        # Move integer from $v0 to $a0
12     li    $v0, 1          # syscall for printing integer to command line (code = 1)
13     syscall
14
15     li    $v0, 10         # Call operating system to perform operation exit
16     syscall
```

Assemble ... GO;

Mars Messages

Run I/O

Clear

25
25
-- program is finished running --

Registers	Coproc 1	Coproc 0
Name	Number	Value
\$zero	0	0
\$at	1	0
\$v0	2	10
\$v1	3	0
\$a0	4	25
\$a1	5	0
\$a2	6	0

System Services (**syscall**) in MIPS

To print an integer to the screen:

Set **\$v0** to **1**

syscall

System Services (**syscall**) in MIPS

To print an integer to the screen:

Set **\$v0** to **1**

syscall

To print a string to the screen:

Set **\$v0** to **4**

syscall

System Services (`syscall`) in MIPS

To print an integer to the screen:

Set `$v0` to `1`

`syscall`

To print a string to the screen:

Set `$v0` to `4`

`syscall`

To read an integer from the keyboard:

Set `$v0` to `5`

`syscall`

System Services (**syscall**) in MIPS

To print an integer to the screen:

Set **\$v0** to **1**

syscall

To print a string to the screen:

Set **\$v0** to **4**

syscall

To read an integer from the keyboard:

Set **\$v0** to **5**

syscall

To exit:

Set **\$v0** to **10**

syscall