

Relations

Chapter 9

Chapter Summary

- Relations and Their Properties
- n -ary Relations and Their Applications (*not currently included in overheads*)
- Representing Relations
- Closures of Relations (*not currently included in overheads*)
- Equivalence Relations
- Partial Orderings

Relations and Their Properties

Section 9.1

Section Summary

- Relations and Functions
- Properties of Relations
 - Reflexive Relations
 - Symmetric and Antisymmetric Relations
 - Transitive Relations
- Combining Relations

Announcements

- ◆ HW7 is due now
- ◆ HW8 has been posted

Relational Databases

- ◆ Relational databases standard organizing structure for large databases
 - Simple design
 - Powerful functionality
 - Allows for efficient algorithms
- ◆ Not all databases are relational
 - Ancient database systems
 - XML –tree based data structure
 - Modern database must: easy conversion to relational

Example 1

A relational database with *schema* :

1	<i>Name</i>
2	<i>Favorite Soap</i>
3	<i>Favorite Color</i>
4	<i>Occupation</i>

1	Kate Winslet	Leonardo DiCaprio
2	Dove	Dial
3	Purple	Green
4	Movie star	Movie star

...etc.

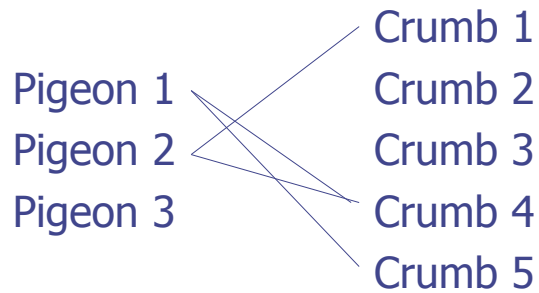
Example 2

The table for **mod** 2 addition:

+	0	1
0	0	1
1	1	0

Example 3

Example of a pigeon to crumb pairing
where pigeons may share a crumb:



Example 4

The concept of "siblinghood".

Relations: Generalizing Functions

Some of the examples were function-like (e.g. **mod** 2 addition, or crumbs to pigeons) but violations of definition of function were allowed (not well-defined, or multiple values defined).

All of the 4 examples had a common thread: They relate elements or properties with each other.

Relations: Represented as Subsets of Cartesian Products

In more rigorous terms, all 4 examples could be represented as subsets of certain Cartesian products.

Q: How is this done for examples 1, 2, 3 and 4?

Relations: Represented as Subsets of Cartesian Products

The 4 examples:

- 1) Database \subseteq
- 2) **mod** 2 addition \subseteq
- 3) Pigeon-Crumb feeding \subseteq
- 4) Siblinghood \subseteq

Relations: Represented as Subsets of Cartesian Products

A:

- 1) Database \subseteq
 $\{\text{Names}\} \times \{\text{Soaps}\} \times \{\text{Colors}\} \times \{\text{Jobs}\}$
- 2) **mod** 2 addition \subseteq
 $\{0,1\} \times \{0,1\} \times \{0,1\}$
- 3) Pigeon-Crumb feeding \subseteq
 $\{\text{pigeons}\} \times \{\text{crumbs}\}$
- 4) Siblinghood \subseteq
 $\{\text{people}\} \times \{\text{people}\}$

Q: What is the actual subset for **mod** 2 addition?

Relations as Subsets of Cartesian Products

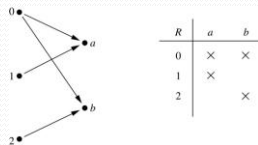
A: The subset for **mod 2** addition:
 $\{ (0,0,0), (0,1,1), (1,0,1), (1,1,0) \}$

Binary Relations

Definition: A *binary relation* R from a set A to a set B is a subset $R \subseteq A \times B$.

Example:

- Let $A = \{0,1,2\}$ and $B = \{a,b\}$
- $\{(0, a), (0, b), (1, a), (2, b)\}$ is a relation from A to B .
- We can represent relations from a set A to a set B graphically or using a table:



Relations are more general than functions. A function is a relation where exactly one element of B is related to each element of A .

Binary Relation on a Set

Definition: A binary relation R on a set A is a subset of $A \times A$ or a relation from A to A .

Example:

- Suppose that $A = \{a, b, c\}$. Then $R = \{(a, a), (a, b), (a, c)\}$ is a relation on A .
- Let $A = \{1, 2, 3, 4\}$. The ordered pairs in the relation $R = \{(a, b) \mid a \text{ divides } b\}$ are $(1, 1), (1, 2), (1, 3), (1, 4), (2, 2), (2, 4), (3, 3)$, and $(4, 4)$.

Binary Relation on a Set (cont.)

Question: How many relations are there on a set A ?

Solution: Because a relation on A is the same thing as a subset of $A \times A$, we count the subsets of $A \times A$. Since $A \times A$ has n^2 elements when A has n elements, and a set with m elements has 2^m subsets, there are $2^{|A|^2}$ subsets of $A \times A$. Therefore, there are $2^{|A|^2}$ relations on a set A .

Binary Relations on a Set (*cont.*)

Example: Consider these relations on the set of integers:

$$\begin{aligned} R_1 &= \{(a,b) \mid a \leq b\}, & R_4 &= \{(a,b) \mid a = b\}, \\ R_2 &= \{(a,b) \mid a > b\}, & R_5 &= \{(a,b) \mid a = b + 1\}, \\ R_3 &= \{(a,b) \mid a = b \text{ or } a = -b\}, & R_6 &= \{(a,b) \mid a + b \leq 3\}. \end{aligned}$$

Note that these relations are on an infinite set and each of these relations is an infinite set.

Which of these relations contain each of the pairs

$(1,1)$, $(1, 2)$, $(2, 1)$, $(1, -1)$, and $(2, 2)$?

Solution: Checking the conditions that define each relation, we see that the pair $(1,1)$ is in R_1 , R_3 , R_4 , and R_6 ; $(1,2)$ is in R_1 and R_6 ; $(2,1)$ is in R_2 , R_5 , and R_6 ; $(1, -1)$ is in R_2 , R_3 , and R_6 ; $(2,2)$ is in R_1 , R_3 , and R_4 .

Reflexive Relations

Definition: R is *reflexive* iff $(a,a) \in R$ for every element $a \in A$. Written symbolically, R is reflexive if and only if

$$\forall x[x \in U \rightarrow (x,x) \in R]$$

Example: The following relations on the integers are reflexive:

$$\begin{aligned} R_1 &= \{(a,b) \mid a \leq b\}, \\ R_3 &= \{(a,b) \mid a = b \text{ or } a = -b\}, \\ R_4 &= \{(a,b) \mid a = b\}. \end{aligned}$$

If $A = \emptyset$ then the empty relation is reflexive vacuously. That is the empty relation on an empty set is reflexive!

The following relations are not reflexive:

$$\begin{aligned} R_2 &= \{(a,b) \mid a > b\} \text{ (note that } 3 \not> 3), \\ R_5 &= \{(a,b) \mid a = b + 1\} \text{ (note that } 3 \neq 3 + 1), \\ R_6 &= \{(a,b) \mid a + b \leq 3\} \text{ (note that } 4 + 4 \not\leq 3). \end{aligned}$$

Symmetric Relations

Definition: R is *symmetric* iff $(b,a) \in R$ whenever $(a,b) \in R$ for all $a,b \in A$. Written symbolically, R is symmetric if and only if

$$\forall x \forall y [(x,y) \in R \rightarrow (y,x) \in R]$$

Example: The following relations on the integers are symmetric:

$$R_3 = \{(a,b) \mid a = b \text{ or } a = -b\},$$

$$R_4 = \{(a,b) \mid a = b\},$$

$$R_6 = \{(a,b) \mid a + b \leq 3\}.$$

The following are not symmetric:

$$R_1 = \{(a,b) \mid a \leq b\} \text{ (note that } 3 \leq 4, \text{ but } 4 \not\leq 3),$$

$$R_2 = \{(a,b) \mid a > b\} \text{ (note that } 4 > 3, \text{ but } 3 \not> 4),$$

$$R_5 = \{(a,b) \mid a = b + 1\} \text{ (note that } 4 = 3 + 1, \text{ but } 3 \neq 4 + 1).$$

Antisymmetric Relations

Definition: A relation R on a set A such that for all $a,b \in A$ if $(a,b) \in R$ and $(b,a) \in R$, then $a = b$ is called *antisymmetric*. Written symbolically, R is antisymmetric if and only if

$$\forall x \forall y [(x,y) \in R \wedge (y,x) \in R \rightarrow x = y]$$

- Example:** The following relations on the integers are antisymmetric:

$$R_1 = \{(a,b) \mid a \leq b\},$$

$$R_2 = \{(a,b) \mid a > b\},$$

$$R_4 = \{(a,b) \mid a = b\},$$

$$R_5 = \{(a,b) \mid a = b + 1\}.$$

The following relations are not antisymmetric:

$$R_3 = \{(a,b) \mid a = b \text{ or } a = -b\}$$

(note that both $(1,-1)$ and $(-1,1)$ belong to R_3),

$$R_6 = \{(a,b) \mid a + b \leq 3\} \text{ (note that both } (1,2) \text{ and } (2,1) \text{ belong to } R_6).$$

For any integer, if $a \leq b$ and $a \leq b$, then $a = b$.

Transitive Relations

Definition: A relation R on a set A is called transitive if whenever $(a,b) \in R$ and $(b,c) \in R$, then $(a,c) \in R$, for all $a,b,c \in A$. Written symbolically, R is transitive if and only if

$$\forall x \forall y \forall z [(x,y) \in R \wedge (y,z) \in R \rightarrow (x,z) \in R]$$

- Example:** The following relations on the integers are transitive:

$$R_1 = \{(a,b) \mid a \leq b\},$$

$$R_2 = \{(a,b) \mid a > b\},$$

$$R_3 = \{(a,b) \mid a = b \text{ or } a = -b\},$$

$$R_4 = \{(a,b) \mid a = b\}.$$

For every integer, $a \leq b$
and $b \leq c$, then $b \leq c$.

The following are not transitive:

$$R_5 = \{(a,b) \mid a = b + 1\} \text{ (note that both } (4,3) \text{ and } (3,2) \text{ belong to } R_5, \text{ but not } (4,2)),$$

$$R_6 = \{(a,b) \mid a + b \leq 3\} \text{ (note that both } (2,1) \text{ and } (1,2) \text{ belong to } R_6, \text{ but not } (2,2)).$$

Combining Relations

- Given two relations R_1 and R_2 , we can combine them using basic set operations to form new relations such as $R_1 \cup R_2$, $R_1 \cap R_2$, $R_1 - R_2$, and $R_2 - R_1$.
- Example:** Let $A = \{1,2,3\}$ and $B = \{1,2,3,4\}$. The relations $R_1 = \{(1,1), (2,2), (3,3)\}$ and $R_2 = \{(1,1), (1,2), (1,3), (1,4)\}$ can be combined using basic set operations to form new relations:

$$R_1 \cup R_2 = \{(1,1), (1,2), (1,3), (1,4), (2,2), (3,3)\}$$

$$R_1 \cap R_2 = \{(1,1)\} \quad R_1 - R_2 = \{(2,2), (3,3)\}$$

$$R_2 - R_1 = \{(1,2), (1,3), (1,4)\}$$

Composition

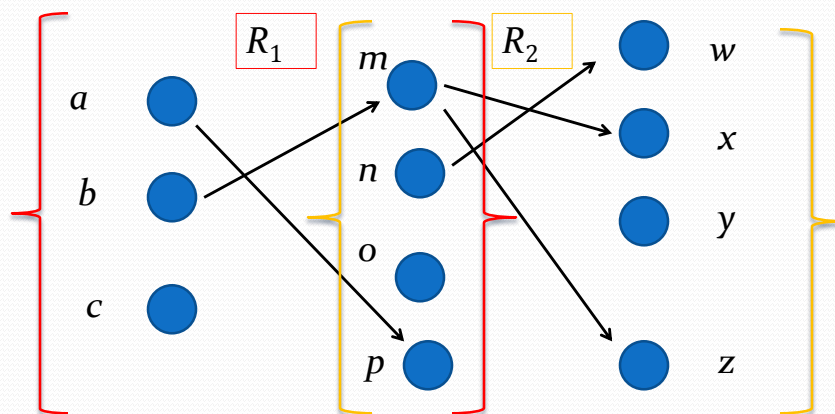
Definition: Suppose

- R_1 is a relation from a set A to a set B .
- R_2 is a relation from B to a set C .

Then the *composition* (or *composite*) of R_2 with R_1 , is a relation from A to C where

- if (x,y) is a member of R_1 and (y,z) is a member of R_2 , then (x,z) is a member of $R_2 \circ R_1$.

Representing the Composition of a Relation



$$R_2 \circ R_1 = \{(b,x), (b,z)\}$$

Powers of a Relation

Definition: Let R be a binary relation on A . Then the powers R^n of the relation R can be defined inductively by:

- Basis Step: $R^1 = R$
- Inductive Step: $R^{n+1} = R^n \circ R$

(see the slides for Section 9.3 for further insights)

The powers of a transitive relation are subsets of the relation. This is established by the following theorem:

Theorem 1: The relation R on a set A is transitive iff $R^n \subseteq R$ for $n = 1, 2, 3, \dots$

(see the text for a proof via mathematical induction)

Representing Relations

Section 9.3

Section Summary

- Representing Relations using Matrices
- Representing Relations using Digraphs

Representing Relations Using Matrices

- A relation between finite sets can be represented using a zero-one matrix.
- Suppose R is a relation from $A = \{a_1, a_2, \dots, a_m\}$ to $B = \{b_1, b_2, \dots, b_n\}$.
 - The elements of the two sets can be listed in any particular arbitrary order. When $A = B$, we use the same ordering.
- The relation R is represented by the matrix $M_R = [m_{ij}]$, where

$$m_{ij} = \begin{cases} 1 & \text{if } (a_i, b_j) \in R, \\ 0 & \text{if } (a_i, b_j) \notin R. \end{cases}$$

- The matrix representing R has a 1 as its (i,j) entry when a_i is related to b_j and a 0 if a_i is not related to b_j .

Examples of Representing Relations Using Matrices

Example 1: Suppose that $A = \{1,2,3\}$ and $B = \{1,2\}$. Let R be the relation from A to B containing (a,b) if $a \in A$, $b \in B$, and $a > b$. What is the matrix representing R (assuming the ordering of elements is the same as the increasing numerical order)?

Solution: Because $R = \{(2,1), (3,1), (3,2)\}$, the matrix is

$$M_R = \begin{bmatrix} 0 & 0 \\ 1 & 0 \\ 1 & 1 \end{bmatrix}.$$

Examples of Representing Relations Using Matrices (*cont.*)

Example 2: Let $A = \{a_1, a_2, a_3\}$ and $B = \{b_1, b_2, b_3, b_4, b_5\}$. Which ordered pairs are in the relation R represented by the matrix

$$M_R = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 \end{bmatrix}?$$

Solution: Because R consists of those ordered pairs (a_i, b_j) with $m_{ij} = 1$, it follows that:

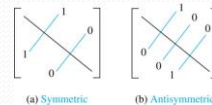
$$R = \{(a_1, b_2), (a_2, b_1), (a_2, b_3), (a_2, b_4), (a_3, b_1), (a_3, b_3), (a_3, b_5)\}.$$

Matrices of Relations on Sets

- If R is a reflexive relation, all the elements on the main diagonal of M_R are equal to 1.



- R is a symmetric relation, if and only if $m_{ij} = 1$ whenever $m_{ji} = 1$. R is an antisymmetric relation, if and only if $m_{ij} = 0$ or $m_{ji} = 0$ when $i \neq j$.



Example of a Relation on a Set

Example 3: Suppose that the relation R on a set is represented by the matrix

$$M_R = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 1 \end{bmatrix}.$$

Is R reflexive, symmetric, and/or antisymmetric?

Solution: Because all the diagonal elements are equal to 1, R is reflexive. Because M_R is symmetric, R is symmetric and not antisymmetric because both $m_{1,2}$ and $m_{2,1}$ are 1.