Buffer Overflow Exploitation

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The goal of this session is to analyse and exploit buffer-overflow vulnerabilities. Buffer-overflow vulnerabilities usually occur when someone is allowed to write and/or to execute code in areas that one should not, and usually derives from the usage of unsafe function like gets.

Modern OS's and compilers already incorporate some security features that prevent these attacks such as canaries, Executable space protection (XP)/Data Execution Prevention (DEP), and Address Space Layout Randomisation (ASLR).

Canaries are known values that are placed between a buffer and control data on the stack to monitor buffer overflows. When the buffer overflows, the first data to be corrupted will be the canary, and a failed verification of the canary data is therefore an alert of an overflow; Executable space protection (XP) prevents certain memory sectors, e.g. the stack, from being executed; and Address Space Layout Randomisation (ASLR) is a technology used to help prevent shell-code from being successful. It does this by randomly offsetting the location of modules and certain in-memory structures.

Combining these techniques makes it very difficult to exploit vulnerabilities in applications using shell-code or return-oriented programming (ROP) techniques. In order for our attacks to succeed we'll thus need to disable them. For that we'll need to compile our programs using the following flags (we use here vuln.c as the program to be compiled, and vuln as the generated binary).

Disable canaries: gcc vuln.c -o vuln -fno-stack-protector

Disable XP: gcc vuln.c -o vuln -z execstack

To disable ASLR we do it for the whole system:

Disable ASLR: echo 0 | sudo tee /proc/sys/kernel/randomize_va_space

Disable ASLR now as we need it for all our exercises.

1 Setup

These exercises will run on the machine downloaded for the first lab. You may want to install some of the auxiliary tools by running ./basic-install.sh.

Task 1 – Exercises with No Protection

0-Simple Overflow

Let us start with program O-simple.c. Compile this file with no-canaries. The goal of this attack is to print the message "YOU WIN!!!" in the screen. How can we do it?

- Recall how variables are recorded in the stack; can variable buffer interfere with variable control?
- You might want to use GDB to see where buffer and control are stored in memory.

1-Match an Exact Value

Now that you know how to overflow a buffer, can you do it with an exact value?. Compile 1-match.c file with no-canaries. The goal of this attack is to print the "Congratulations" message in the screen. Can you do it?

- Notice that the argument is input in the command line, ie, ./1-match <string to write in buffer>;
- Recall that 0x61626364 is the string abcd;
- Have you heard of little-endian and big-endian?

2-Calling Functions

Ok, we already know how to overflow a buffer in a controlled way and change variables. But can I call a function that is not called anywhere in our code? Compile 2-functions.c file with no-canaries. The goal of this attack is to call function win and print the "Congratulations" message in the screen. Can you do it?

- Recall that the name of a function in C is the address where this function is written in the memory;
- Can fp be win?

3-Return Address

We now know everything about the stack and how to change its values, change the functions that are called and so on. But can I call a function even if NO function is called anywhere in our code? Compile 3-return.c file with no-canaries. The goal of this attack is to call function win and print the "Congratulations" message in the screen. Can you do it?

- Recall that the name of a function in C is the address where this function is written in the memory;
- Is it true that no function is called in our program? How can I call win? Recall how the stack is organised and what are the values stored in the stack.

4-Unhauthorized Access

The goal of this attack is to gain access to the system without introducing the correct password *V3RY53cr37*. Can you do it? Compile 4-auth.c file with no-canaries and elaborate on Exercise 1.

5-Unhauthorized Access

The goal of this attack is again to gain access to the system without introducing the correct password V3RY53cr37. But now, we have changed lines 11 and 12 of the previous program. Can you still do it? Compile 5-auth.c file with no-canaries and elaborate on Exercise 3.

6-Running Arbitrary Code

Congratulations, You know now how to change the values of your stack variables, and to modify the expected flow of execution of a program. But can we do more? Can we exploit buffer-overflows to run arbitrary code?

The goal of this attack is to run several pieces of code that do nasty things. Compile 6-code.c file with no-canaries and no XP. You should do it in root mode. Type

sudo gcc 6-code.c -o 6-code -fno-stack-protector -z execstack

To run the program with some attack-string that is stored in file run

```
cat <file> | sudo ./6-code
```

Shell-code The pieces of code we want to execute are available in the file test-shellcode.c. Without getting into much details the 4 pieces of code in this file are the low-level instructions of programs that:

- 1. Print "Hello world";
- 2. Print the file /etc/passwd on the screen;
- 3. Reboot our machine;
- 4. Open a connection that allows one to control this shell remotely.

If you want to try these pieces in isolation just uncomment one at a time and run

gcc test-shellcode.c -o test-shellcode -fno-stack-protector -z execstack; ./test-shellcode
Some important notes

2. Change the permissions of /etc/passwd so that only root can read and write to this file.

```
sudo chmod 600 /etc/passwd
```

Can you still run this code and get the output? No?

Good! We'll explore this later in our examples. For now let's revert it otherwise you might not be able to login again :-)

```
sudo chmod 644 /etc/passwd
```

- 3. Be carefull and save all your open files before running this one! Also, when you boot your machine up again, do not forget to disable ASLR again.
- 4. Read the instructions and follow the steps indicated in the test-shellcode.c file for this case. What happens when you do ls? And ls /root? And if you do sudo ls /root?

Injecting Shell-code Now that you have experimented the 4-pieces of code and that you know that they do what they are supposed to do (in your machine) when you have the correct privileges, can you do it in another machine/process? Saying it differently, can you execute them without running ./test-shellcode?

- 2. Notice that 6-code is running as root. Can you make it show the /etc/passwd file when you change its permissions to 600? Never forget to revert the permissions of /etc/passwd back to 644 afterwards!!!
- 4. Supposing that 6-code is running on a different machine (as root), can you read the content of its directories from the second terminal? Can you ls /root now from the second terminal? And from the first terminal?

7-Environment Variables

Good, we already know how to run arbitrary code. but what happens when the buffer we can inject is to small? The goal of this attack is to run the same pieces of code as in the previous exercise but when the buffer is much smaller (128 to 16). Compile 7-environment.c file with no-canaries and no XP.

- Have you heard off environment variables?
- export MYVAR =; echo \$MYVAR?
- Try to find this variable's address in order to run the attacks.

References

[SEEDLabs] http://www.cis.syr.edu/~wedu/seed/Labs_12.04/Software/Buffer_Overflow/Buffer_Overflow.pdf

Extras – Strings with non-printable chars

To perform these attacks you may need to input chars that are non-printable. The easier way to do it is to write such input string to a file, and then use this file as input to the program. For this, do the following:

```
python -c 'print string1 + string2 + ... + stringn' > input-file.txt
./program < input-file.txt
Example:
python -c 'print "\x48\x31\xc0\x48" + "\x90" * 30 + "\x7f\x00\x00"' > input-exerciseX.txt
./file < input-exerciseX.txt</pre>
```

Extras – Some GDB commands

```
gdb ./file — open GDB to debug file
disas function — disassemble function function
disas address — disassemble function at this address
b address — inserts a breakpoint at address address
r < file — run the program with input file
s — execute next step of program
c — execute until the next breakpoint of the program
stack n — show the n registers after the stack pointer
x/nx $rsp — show the n registers after the register $rsp
x/nx address — show the n registers after the address
```