

INFINITESIMAL CALCULUS

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Ultrafilters

Let S be a set. If $\mathcal{F} \subseteq \mathcal{P}(S)$, we say that \mathcal{F} is an *ultrafilter* on S if:

- For any $A, B \in \mathcal{F}$, we have $A \cap B \in \mathcal{F}$
- If $A \in \mathcal{F}$ and $A \subseteq B \in \mathcal{P}(S)$, then $B \in \mathcal{F}$
- For any $A \subseteq S$, either $A \in \mathcal{F}$ or $A^c \in \mathcal{F}$
- \mathcal{F} is a proper subset of $\mathcal{P}(S)$

We will use \mathcal{F} to denote an arbitrary ultrafilter on \mathbb{N} that is *non-principle*, meaning (in this context) it does not contain any finite sets. Since \mathcal{F} is an ultrafilter, this means any *cofinite set*, i.e. any set whose complement is finite, is in \mathcal{F} .

Ultrapower of \mathbb{R}

Let $\mathbb{R}^{\mathbb{N}}$ denote the set of sequences of real numbers. We will denote a member $r = \langle r_1, r_2, r_3, \dots \rangle$ of $\mathbb{R}^{\mathbb{N}}$ by $\langle r_n \rangle$. We can define operations \oplus and \odot on $\mathbb{R}^{\mathbb{N}}$ by $\langle r_n \rangle \oplus \langle s_n \rangle = \langle r_n + s_n \rangle$ and $\langle r_n \rangle \odot \langle s_n \rangle = \langle r_n \cdot s_n \rangle$, giving us a ring $(\mathbb{R}^{\mathbb{N}}, \oplus, \odot)$. We define an equivalence relation \equiv by $\langle r_n \rangle \equiv \langle s_n \rangle$ if $\{n \in \mathbb{N} \mid r_n = s_n\} \in \mathcal{F}$. Write $[[r_n = s_n]] = \{n \in \mathbb{N} \mid r_n = s_n\}$. Now, \equiv is reflexive, as $[[r_n = r_n]] = \mathbb{N}$, and $\emptyset \notin \mathcal{F}$ implies $\mathbb{N} \in \mathcal{F}$. \equiv is symmetric, since $[[r_n = s_n]] = [[s_n = r_n]]$. And \equiv is transitive, since if $[[r_n = s_n]] \in \mathcal{F}$ and $[[s_n = t_n]] \in \mathcal{F}$, we have both $[[r_n = s_n]] \cap [[s_n = t_n]] \in \mathcal{F}$ and $[[r_n = s_n]] \cap [[s_n = t_n]] \subseteq [[r_n = t_n]]$. We let $[r]$ denote the equivalence class of r under \equiv . This can also be written as $[\langle r_n \rangle]$, or abbreviated as $[r_n]$. We then define the *hyperreals* ${}^*\mathbb{R} = \{[r] \mid r \in \mathbb{R}^{\mathbb{N}}\}$. Note that ${}^*\mathbb{R}$ in some sense contains \mathbb{R} : we can identify any real number a with he hyperreal $[\langle a, a, a, \dots \rangle]$.

Hyperreal extensions

We define addition on the hyperreals $[r] + [s] = [r \oplus s]$, and we define multiplication similarly. Furthermore, we can extend any function $f : \mathbb{R} \rightarrow \mathbb{R}$ to the hyperreals by letting $*f([r_n]) = [f(r_n)]$. For instance, if $f(x) = x^2$ and $[r] = [\langle 1, 1/2, 1/3, \dots \rangle]$, then $*f([r]) = [\langle 1, 1/4, 1/9, 1/16, \dots \rangle]$. Better yet, we can extend any k -ary relation R_k by saying that for any $[r^1], [r^2], [r^3], \dots, [r^k] \in {}^*\mathbb{R}$ (that's an upper index, not an exponent), we have $*R_k([r^1], [r^2], \dots, [r^k])$ if $[[R_k(r_n^1, r_n^2, \dots, r_n^k)]] \in \mathcal{F}$. In particular, $[r] < [s]$ if $[[r_n < s_n]] \in \mathcal{F}$. Doing this for 1-ary relations lets us extend any set to the hyperreals. For instance, $[r] \in {}^*\mathbb{N}$ if $[[r_n \in \mathbb{N}]] \in \mathcal{F}$. We must, of course, prove that all of these extensions are well-defined and do not depend on the representative we choose for $[\langle r_n \rangle]$, but this follows from the properties of \mathcal{F} .

Transfer

We construct a mathematical language with the following symbols:

- \wedge (and), \vee (or), \rightarrow (if-then), and \neg (not)
- a symbol c for every $c \in \mathbb{R}$
- a symbol f for every function $f : \mathbb{R} \rightarrow \mathbb{R}$
- a symbol R_k for every k -ary relation on the \mathbb{R} (including $=$)
- \forall (for all) and \exists (there exists)

Any sentence in this language about the reals we can “reinterpret” as being about the hyperreals. For instance, say $f : \mathbb{R} \rightarrow \mathbb{R}$. If we want to say that for every real number x , there is a natural number larger than $f(x)$, we write:

$$(\forall x \in \mathbb{R})(\exists y \in \mathbb{N})(f(x) < y)$$

To “reinterpret” this about the hyperreals, we replace objects with their extensions where possible:

$$(\forall x \in {}^*\mathbb{R})(\exists y \in {}^*\mathbb{N})(*f(x) < y)$$

Transfer Principle: Any sentence we can write in this language is true in \mathbb{R} if and only if its “reinterpretation” in ${}^*\mathbb{R}$ is true.

Structure of ${}^*\mathbb{R}$

We call an element $x \in {}^*\mathbb{R}$ *infinitesimal* if $|x| < r$ for any $r \in \mathbb{R}^+$ (here treating $[\langle r, r, \dots \rangle]$ as r). The only infinitesimal real number is 0, but consider $x = [\langle 1, 1/2, 1/3, 1/4 \dots \rangle]$. For any $r \in \mathbb{R}^+$, by the archimedean property there is some N such that $1/N < r$, and so $[[x_n < r]] \subseteq \{n \in \mathbb{N} \mid N \leq n\}$. This set is cofinite, and so is in \mathcal{F} . Hence, $x < r$. However, $[[0 < x_n]] = \mathbb{N} \in \mathcal{F}$, so $0 < x$.

We call an element $x \in {}^*\mathbb{R}$ *unbounded* if $|x| > r$ for any $r \in \mathbb{R}^+$. For example, $[\langle 1, 2, 3, 4, \dots \rangle]$ is unbounded. An element that is not unbounded is *bounded*. An element that is bounded but not infitesimal is *appreciable*.

If $|x - y|$ is infinitesimal, we say that x and y are *infinitely close* and write $x \simeq y$. So x is infinitesimal iff $x \simeq 0$. If x is bounded, then there is a unique $r \in \mathbb{R}$ such that $x \simeq r$ called the *standard part* of x , denoted $\text{st}(x)$. Note \simeq is a transitive relation, so if $\text{st}(x) = \text{st}(y)$ we have $x \simeq \text{st}(x) = \text{st}(y) \simeq y$ and so $x \simeq y$. Similarly, if $x \simeq y$ then $\text{st}(x) \simeq x \simeq y \simeq \text{st}(y)$, so $\text{st}(x) \simeq \text{st}(y)$. But $\text{st}(x)$ and $\text{st}(y)$ are both real, and so their difference is real. Since $\text{st}(x) \simeq \text{st}(y)$ implies their difference is infinitesimal, and the only real infinitesimal is 0, their difference must be 0 and so $\text{st}(x) = \text{st}(y)$. Hence $\text{st}(x) = \text{st}(y)$ iff $x \simeq y$.

Derivatives

We want to denote by $f'(x)$ an infinitesimal change in $f(x)$ over an infinitesimal change in x . Let Δx be a nonzero infinitesimal change in x . The corresponding change in $f(x)$ is $f(x + \Delta x) - f(x)$. So we define

$$f'(x) = \text{st} \left(\frac{f(x + \Delta x) - f(x)}{\Delta x} \right)$$

Since Δx can be any nonzero infinitesimal, this might not be well-defined. If it is, we say that f is *differentiable* at x . Otherwise, we say that $f'(x)$ is undefined. This is equivalent to the standard definition

$$f'(x) = \lim_{h \rightarrow 0} \frac{f(x + h) - f(x)}{h}$$

Simple Proofs with Infinitesimals

The primary motivation behind infinitesimal calculus is to allow us to prove reuslts in intuitive ways. Take for instance, the Chain Rule: standardly, the proof has an intuitive outline but runs into several techincal issues. Nonstandardly, a rigorous proof is much simpler:

Theorem (Chain Rule). *If $f, g : \mathbb{R} \rightarrow \mathbb{R}$ are differentiable, then $(f \circ g)'(x) = f'(g(x)) \cdot g'(x)$.*

Proof. Let Δx be any nonzero infinitesimal, and $\Delta g = g(x + \Delta x) - g(x)$. Since $g'(x) = \text{st}(\Delta g / \Delta x)$ is defined, $\Delta g / \Delta x$ is bounded, and so Δg must be infinitesimal (as a non-infinitesimal divided by an infinitesimal is unbounded). If $\Delta g \neq 0$, then

$$\begin{aligned} (f \circ g)'(x) &\simeq \frac{f(g(x + \Delta x)) - f(g(x))}{\Delta x} \\ &= \frac{f(g(x) + \Delta g) - f(g(x))}{\Delta g} \cdot \frac{\Delta g}{\Delta x} \\ &\simeq f'(g(x)) \cdot g'(x) \end{aligned}$$

So $(f \circ g)'(x) \simeq f'(g(x)) \cdot g'(x)$. But since both of these numbers are real, they must be identical. In the case where $\Delta g = 0$, we clearly have $f(g(x) + \Delta g) - f(g(x)) = 0$ and so we still get $(f \circ g)'(x) = 0$. \square