

# Infinitesimal Calculus

Paul Schulze

# History

- ▶ Newton & Leibniz formulated calculus using the idea of *infinitesimals*
- ▶ Infinitesimals are really really small, but not 0
- ▶ Considered nonsensical, replaced with  $\delta$ - $\epsilon$
- ▶ Early 1960's: Abraham Robinson formalizes Nonstandard Analysis
- ▶ Our formulation of infinitesimals is based off work by Jerzy Łoś



(a) Abraham Robinson



(b) Jerzy Łoś

# Basic Idea

- ▶ We construct a set of *hyperreals*, denoted  ${}^*\mathbb{R}$
- ▶  ${}^*\mathbb{R}$  includes  $\mathbb{R}$ , along with (a lot of) hyperreals
- ▶ We construct a language of mathematical logic, using symbols like  $\neg$ ,  $\wedge$ ,  $\vee$ ,  $\forall$ , etc.
- ▶ We show that any sentence of that language is true in  $\mathbb{R}$  iff it is true in  ${}^*\mathbb{R}$  (transfer principle)
- ▶ We use transfer to prove things about  $\mathbb{R}$

# Constructing ${}^*\mathbb{R}$

- ▶ We start with the ring  $\mathbb{R}^\infty$
- ▶ Identify sequences that are the same “almost everywhere,” like  $\langle 0, 1, 1, \dots \rangle \sim \langle 1, 1, 1, \dots \rangle$
- ▶ Sequences are the same almost everywhere if the set of indices at which they are the same is “big”
- ▶ The set of “big” sets of natural numbers is an *ultrafilter*  $\mathcal{F} \subseteq \mathcal{P}(\mathbb{N})$

# Constructing ${}^*\mathbb{R}$

- ▶ Now we can define our equivalence relation  $\sim$  by saying that  $\langle r_1, r_2, r_3, \dots \rangle \sim \langle s_1, s_2, s_3, \dots \rangle$  iff  $\{n \in \mathbb{N} \mid r_n = s_n\} \in \mathcal{F}$
- ▶ Write the equivalence class  $[\langle r_1, r_2, r_3, \dots \rangle]$
- ▶ Define  ${}^*\mathbb{R}$  as the quotient ring of  $\mathbb{R}^\infty$  under  $\sim$ , i.e.  
$${}^*\mathbb{R} = \{[\langle r_1, r_2, r_3, \dots \rangle] \mid \langle r_1, r_2, \dots \rangle \in \mathbb{R}^\infty\}$$
- ▶ Extend any function  $f : \mathbb{R} \rightarrow \mathbb{R}$  to a new  ${}^*f : {}^*\mathbb{R} \rightarrow {}^*\mathbb{R}$   
$${}^*f([\langle r_1, r_2, r_3, \dots \rangle]) = [\langle f(r_1), f(r_2), f(r_3), \dots \rangle]$$
- ▶ We can also extend relations, like “ $\leq$ ”  
$$\langle r_1, r_2, r_3, \dots \rangle \leq \langle s_1, s_2, s_3, \dots \rangle \text{ iff } \{n \in \mathbb{N} \mid r_n \leq s_n\} \in \mathcal{F}$$
- ▶ For any  $x \in \mathbb{R}$ , we can take  $x \in {}^*\mathbb{R}$  to mean  $[\langle x, x, x, \dots \rangle]$

# Transfer Principle—Language

- ▶ Our language is made up of:
  - ▶ logical connectives  $\wedge$ ,  $\vee$ ,  $\rightarrow$ ,  $\leftrightarrow$ , and  $\neg$
  - ▶ quantifiers  $\forall$ ,  $\exists$
  - ▶ parenthesis ( and )
  - ▶ variables  $v_1, v_2, v_3, \dots$
  - ▶ symbols for every element of  $\mathbb{R}$ , every relation on  $\mathbb{R}$ , and every function on  $\mathbb{R}$

exp

# Infinitesimals

- ▶ We construct a set of *hyperreals*  ${}^*\mathbb{R} \supseteq \mathbb{R}$ .
- ▶  ${}^*\mathbb{R}$  is “like”  $\mathbb{R}$ , but it includes *infinitesimals*, elements  $\delta$  such that  $\delta \neq 0$  but  $|\delta| < r$  for every  $r \in \mathbb{R}^+$ .
- ▶ We can add these infinitesimals to other numbers to get things like  $1 + \delta$ , a number that is “infinitely close to” 1 but not 1.
- ▶ If  $|x - y|$  is infinitesimal or 0, we say  $x \simeq y$
- ▶ If  $x \in {}^*\mathbb{R}$ , we denote by  $\text{st}(x)$  the *standard part of*  $x$ , the unique real number that is infinitely close to  $x$ .  $\text{st}(1 + \delta) = 1$ .
- ▶ We can also take the reciprocals of these infinitesimals to get *unbounded* hyperreals, like  $\frac{1}{\delta}$ . These have no standard part.
- ▶ We can of course combine all these elements however we’d like. If  $\delta$  and  $\gamma$  are infinitesimals, we have  $\frac{1}{\delta} + 4 + \pi + \gamma \in {}^*\mathbb{R}$ .



# Derivatives, the way Leibniz intended

- ▶ Say  $f : \mathbb{R} \rightarrow \mathbb{R}$ . We “extend”  $f$  to  $*f : *\mathbb{R} \rightarrow *\mathbb{R}$ .
- ▶ Fix  $b \in \mathbb{R}$ . Let  $\Delta x$  be infinitesimal, and let  $\Delta f$  be  $*f(b + \Delta x) - *f(b)$ .
- ▶ Then define  $f'(b) = \text{st} \left( \frac{\Delta f}{\Delta x} \right)$ . So  $f'(b) \simeq \frac{\Delta f}{\Delta x}$ .
- ▶ **Example:** Say  $f(x) = x^2$ . Then we have

$$\begin{aligned} f'(3) &\simeq \frac{(3 + \Delta x)^2 - 3^2}{\Delta x} = \frac{9 + 6\Delta x + (\Delta x)^2 - 9}{\Delta x} \\ &= \frac{6\Delta x + (\Delta x)^2}{\Delta x} = 6 + \Delta x \simeq 6 \end{aligned}$$

- ▶ So  $f'(3) \simeq 6$ . But these are both real numbers, so their difference can't be infinitesimal. Hence  $f'(3) = 6$ .

## Proof: Chain Rule

Let  $f, g : \mathbb{R} \rightarrow \mathbb{R}$  be differentiable. Let  $\Delta x$  be any infinitesimal, and  $\Delta g = g(x + \Delta x) - g(x)$ . Since  $g'(x) = \text{st}(\Delta g / \Delta x)$  is defined,  $\Delta g$  must be infinitesimal. Then

$$\begin{aligned}(f \circ g)'(x) &\simeq \frac{f(g(x + \Delta x)) - f(g(x))}{\Delta x} \\&= \frac{f(g(x) + \Delta g) - f(g(x))}{\Delta g} \cdot \frac{\Delta g}{\Delta x} \\&\simeq f'(g(x)) \cdot g'(x)\end{aligned}$$

So  $(f \circ g)'(x) \simeq f'(g(x)) \cdot g'(x)$ . But since both of these numbers are real, they must be identical.

In the case where  $\Delta g = 0$ , we clearly have  $f(g(x) + \Delta g) - f(g(x)) = 0$  and so  $(f \circ g)'(x) = 0$ .