Number Theory Notebook

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1 Chapter 1

Divisibility and congruence

Theorem 1.1. Let a, b, and c be integers. If $a \mid b$ and $a \mid c$, then $a \mid (b+c)$.

Proof.

$$\begin{array}{lll} (1.1.1) & \exists \ d_b \in \mathbb{Z} \ni b = a \cdot d_b & \text{because } a \mid b \\ (1.1.2) & \exists \ d_c \in \mathbb{Z} \ni c = a \cdot d_c & \text{because } a \mid c \\ (1.1.3) & b + c = a \cdot d_b + a \cdot d_c & \text{by (1.1.1) and (1.1.2)} \\ (1.1.4) & b + c = a \cdot (d_b + d_c) & \text{by distributive property} \\ (1.1.5) & d_b + d_c \in \mathbb{Z} & \text{because } d_b \in \mathbb{Z} \text{ and } d_c \in \mathbb{Z} \\ & a \mid (b + c) & \text{by def'n of divides} \end{array}$$

Theorem 1.2. Let a, b, and c be integers. If $a \mid b$ and $a \mid c$, then $a \mid (b-c)$.

Proof.

$$\begin{array}{lll} (1.2.1) & \exists \ d_b \in \mathbb{Z} \ni b = a \cdot d_b & \text{because } a \mid b \\ (1.2.2) & \exists \ d_c \in \mathbb{Z} \ni c = a \cdot d_c & \text{because } a \mid c \\ (1.2.3) & b - c = a \cdot d_b - a \cdot d_c & \text{by (1.2.1) and (1.2.2)} \\ (1.2.4) & b - c = a \cdot (d_b - d_c) & \text{by distributive property} \\ (1.2.5) & d_b - d_c \in \mathbb{Z} & \text{because } d_b \in \mathbb{Z} \text{ and } d_c \in \mathbb{Z} \\ & a \mid (b - c) & \text{by def'n of divides} \end{array}$$

Theorem 1.3. Let a, b, and c be integers. If $a \mid b$ and $a \mid c$, then $a \mid bc$.

Proof.

$$\begin{array}{lll} (1.3.1) & \exists \ d_b \in \mathbb{Z} \ni b = a \cdot d_b & \text{because } a \mid b \\ (1.3.2) & \exists \ d_c \in \mathbb{Z} \ni c = a \cdot d_c & \text{because } a \mid c \\ (1.3.3) & bc = (a \cdot d_b) \cdot (a \cdot d_c) & \text{by } (1.3.1) \text{ and } (1.3.2) \\ (1.3.4) & bc = a \cdot (a \cdot d_b \cdot d_c) & \text{by associativity and commutativity} \\ (1.3.5) & a \cdot d_b \cdot d_c \in \mathbb{Z} & \text{because } d_b \in \mathbb{Z} \text{ and } d_c \in \mathbb{Z} \\ & a \mid bc & \text{by def'n of divides} & \Box \end{array}$$

Question 1.4. Can you weaken the hypothesis of the previous theorem and still prove the conclusion? Can you keep the same hypothesis, but replace the conclusion by the stronger conclusion that $a^2|bc$ and still prove the theorem?

Yes. You can remove the a|c condition to weaken the hypothesis, or with both a|b and a|c you can show $a^2|bc$.

Question 1.5. Can you formulate your own conjecture along the lines of the above theorems and then prove it to make it your theorem?

Yes.

Paul's Conjecture. Let a, b, and c be integers. If a|b and a|c, then $a^2|bc$.

Proof. First, take lines (1.3.1) through (1.3.4) of the proof of Theorem 1.3. Then,

$$d_b \cdot d_c \in \mathbb{Z}$$
 because $d_b \in \mathbb{Z}$ and $d_c \in \mathbb{Z}$ by def'n of divides

Theorem 1.6. Let a, b, and c be integers. If a|b, then a|bc.

Proof.

$$(1.6.1) \exists d \in \mathbb{Z} \ni ad = b because a|b$$

(1.6.2)
$$bc = adc$$
 by (1.6.1)

Exercise 1.7. Answer each of the following questions, and prove that your answer is correct.

1. Is
$$45 \equiv 9 \pmod{4}$$
?
Yes. $4 \cdot 9 = 36 = 45 - 9$.

2. Is
$$37 \equiv 2 \pmod{5}$$
?
Yes. $5 \cdot 7 = 35 = 37 - 2$.

3. Is
$$37 \equiv 3 \pmod{5}$$
?
No. $37 - 3 = 34$ which is not a multiple of 5.

4. Is
$$37 \equiv -3 \pmod{5}$$
?
Yes. $5 \cdot 8 = 40 = 37 - (-3)$.

Exercise 1.8. For each of the following congruences, characterize all the integers m that satisfy that congruence.

1.
$$m \equiv 0 \pmod{3}$$

 $m \in \{3z \mid z \in \mathbb{Z}\}$

2.
$$m \equiv 1 \pmod{3}$$

 $m \in \{3z + 1 \mid z \in \mathbb{Z}\}$

3.
$$m \equiv 2 \pmod{3}$$

 $m \in \{3z + 2 \mid z \in \mathbb{Z}\}$

4.
$$m \equiv 3 \pmod{3}$$

 $m \in \{3z \mid z \in \mathbb{Z}\}$

5.
$$m \equiv 4 \pmod{3}$$

 $m \in \{3z + 1 \mid z \in \mathbb{Z}\}$

Theorem 1.9. Let a and n be integers with n > 0. Then $a \equiv a \pmod{n}$.

Proof.

$$(1.9.1) 0 \in \mathbb{Z}$$

$$(1.9.2) n \cdot 0 = 0$$

(1.9.3)
$$n|0$$
 By def'n of divides

$$(1.9.4) a - a = 0$$

(1.9.5)
$$n|(a-a)$$
 By (1.9.3) and (1.9.4) $a \equiv a \pmod{n}$ By def'n of modular congruence

Theorem 1.10. Let a, b, and n be integers with n > 0. If $a \equiv b \pmod{n}$, then $b \equiv a \pmod{n}$.

Proof.

$$(1.10.1) a \equiv b \pmod{n} Given$$

(1.10.2)
$$\exists d \in \mathbb{Z} \ni nd = a - b$$
 By def'n of modular congruence

(1.10.3)
$$-1nd = -1 \cdot (a - b)$$
 By multiplicative property of equality

(1.10.4)
$$n \cdot (-d) = b - a$$
 By various algebra

(1.10.5)
$$-d \in \mathbb{Z}$$
 By multiplicative closure of \mathbb{Z}

(1.10.6)
$$n|(b-a)$$
 By (1.10.4), (1.10.5)

$$b \equiv a \pmod{n}$$
 By def'n of modular congruence

Theorem 1.11. Let a, b, and n be integers with n > 0. If $a \equiv b \pmod{n}$ and $b \equiv c \pmod{n}$, then $a \equiv c \pmod{n}$.

Proof.

$$(1.11.1) n|a-b By a \equiv b \pmod{n}$$

$$(1.11.2) n|b-c \text{By } b \equiv c \pmod{n}$$

(1.11.3)
$$\exists d_1 \in \mathbb{Z} \ni nd_1 = a - b$$
 By (1.11.1)

(1.11.4)
$$\exists d_2 \in \mathbb{Z} \ni nd_2 = b - c$$
 By (1.11.2)

(1.11.5)
$$nd_1 + nd_2 = (a-b) + (b-c)$$
 By additive property of equality

$$(1.11.6) n(d_1 + d_2) = a - c By various algebra$$

(1.11.7)
$$d_1 + d_2 \in \mathbb{Z}$$
 By closure of integers under addition

(1.11.8)
$$n|(a-c)$$
 By def'n of divides $a \equiv c \pmod{n}$ By def'n of modular congruence

Theorem 1.12. Let a, b, c, d, and n be integers with n > 0. If $a \equiv b \pmod{n}$ and $c \equiv d \pmod{n}$, then $a + c \equiv b + d \pmod{n}$.

Proof.

$$(1.12.1) n|(a-b) By a \equiv b \pmod{n}$$

(1.12.2)
$$\exists d_1 \in \mathbb{Z} \ni nd_1 = a - b$$
 By def'n divides

$$(1.12.3) n|(c-d) By c \equiv d \pmod{n}$$

$$(1.12.4) \exists d_2 \in \mathbb{Z} \ni nd_2 = c - d By def'n divides$$

$$(1.12.5) nd_1 + nd_2 = (a-b) + (c-d) By additive property of equality$$

(1.12.6)
$$n \cdot (d_1 + d_2) = (a+c) - (b+d)$$
 By various algebra

$$(1.12.7) d_1 + d_2 \in \mathbb{Z} By additive closure of \mathbb{Z}$$

$$(1.12.8)$$
 $n|((a+c)-(b+d))$ By def'n of divides

$$a + c \equiv b + d \pmod{n}$$
 By def'n of modular congruence

Theorem 1.13. Let a, b, c, d, and n be integers with n > 0. If $a \equiv b \pmod{n}$ and $c \equiv d \pmod{n}$, then $a - c \equiv b - d \pmod{n}$.

Proof. Notice -c and -d are integers, and $-c \equiv -d \pmod{n}$ (glossing over the proof of that for now). Then simply cite 1.12 and we're done.

Theorem 1.14. Let a, b, c, d, and n be integers with n > 0. If $a \equiv b \pmod{n}$ and $c \equiv d \pmod{n}$, then $ac \equiv bd \pmod{n}$.

Proof.

$$(1.14.1) n|(a-b) By a \equiv b \pmod{n}$$

$$(1.14.2) \exists k_1 \in \mathbb{Z} \ni a - b = nk_1$$

$$(1.14.3) a = nk_1 + b$$

(1.14.4)
$$n|(c-d)$$
 By $c \equiv d \pmod{n}$

$$(1.14.5) \exists k_2 \in \mathbb{Z} \ni c - d = nk_2$$

$$(1.14.6) c = nk_2 + d$$

(1.14.7)
$$ac = (nk_1 + b)(nk_2 + d)$$
 By (1.14.3) and (1.14.6)

$$(1.14.8) ac = n^2 k_1 k_2 + nk_1 d + nk_2 b + bd$$

$$(1.14.9) ac - bd = n \cdot (nk_1k_2 + k_1d + k_2b)$$

(1.14.10)
$$n|(ac-bd)$$
 Since $nk_1k_2 + k_1d + k_2b \in \mathbb{Z}$
$$ac \equiv bd \pmod{n}$$

Exercise 1.15. Let a, b, and n be integers with n > 0. Show that if $a \equiv b \pmod{n}$, then $a^2 \equiv b^2 \pmod{n}$.

Proof.

(1.15.1)
$$a \equiv b \pmod{n}$$
 Given
(1.15.2) $a \cdot a \equiv b \cdot b \pmod{n}$ 1.14 $a^2 \equiv b^2 \pmod{n}$

Exercise 1.16. Let a, b, and n be integers with n > 0. Show that if $a \equiv b \pmod{n}$, then $a^3 \equiv b^3 \pmod{n}$.

Proof.

$$(1.16.1) a \equiv b \pmod{n} Given$$

(1.16.2)
$$a^2 \equiv b^2 \; (\bmod \; n)$$
 1.15

(1.16.3)
$$a \cdot a^2 \equiv b \cdot b^2 \pmod{n}$$
 By 1.14 on (1.16.1) and (1.16.2)
$$a^3 \equiv b^3 \pmod{n}$$

Exercise 1.17. Let a, b, k, and n be integers with n > 0 and k > 1. Show that if $a \equiv b \pmod{n}$ and $a^{k-1} \equiv b^{k-1} \pmod{n}$, then $a^k \equiv b^k \pmod{n}$.

Proof.

$$(1.17.1) a \equiv b \pmod{n} Given$$

$$(1.17.2) a^{k-1} \equiv b^{k-1} \pmod{n} 1.15$$

(1.17.3)
$$a \cdot a^{k-1} \equiv b \cdot b^{k-1} \pmod{n}$$
 By 1.14 on (1.17.1) and (1.17.2)
$$a^k \equiv b^k \pmod{n}$$

Theorem 1.18. Let a, b, k, and n be integers with n > 0 and k > 0. If $a \equiv b \pmod{n}$, then $a^k \equiv b^k \pmod{n}$

Proof. Our base case is 1.9. Our induction hypothesis is "a, b, k, and n are integers with n > 0 and k > 1 such that $\forall j \ni 0 < j < k$, we find $a^j \equiv b^j \pmod{n}$. Notice our induction hypothesis fulfills the criteria for 1.17, and in fact 1.17 covers our induction step.

Exercise 1.19. Illustrate each of Theorems 1.12 - 1.18 with an example using actual numbers

1.12 $2 \equiv 12 \pmod{1}0$ and $5 \equiv 15 \pmod{1}0$ imply $7 \equiv 27 \pmod{1}0$.

1.13 $7 \equiv 27 \pmod{1}0$ and $12 \equiv 2 \pmod{1}0$ imply that $-5 \equiv 25 \pmod{1}0$.

1.14 $2 \equiv 7 \pmod{5}$ and $3 \equiv 8 \pmod{5}$ imply that $6 \equiv 56 \pmod{5}$.

1.15 $2 \equiv 7 \pmod{5}$ implies that $4 \equiv 49 \pmod{5}$.

1.16 $1 \equiv 3 \pmod{2}$ implies that $1 \equiv 27 \pmod{2}$.

1.17 $1 \equiv 3 \pmod{2}$ and $1 \equiv 27 \pmod{2}$ imply that $1 \equiv 81 \pmod{2}$.

1.18 $1 \equiv 3 \; (\bmod \; 2) \; \text{implies that} \; 1 \equiv 81 \; (\bmod \; 2).$

Question 1.20. Let a, b, c, and n be integers for which $ac \equiv bc \pmod{n}$. Can we conclude that $a \equiv b \pmod{n}$? If you answer "yes", try to give a proof. If you answer "no", try to give a counterexample.

No. Notice $1 \cdot 0 \equiv 2 \cdot 0 \pmod{5}$ and yet $1 \not\equiv 2 \pmod{5}$.

Theorem 1.21. Let a natural number n be expressed in base 10 as

$$n = a_k a_{k-1} \dots a_1 a_0$$

If $m = a_k + a_{k-1} + \dots + a_1 + a_0$ then $n \equiv m \pmod{3}$.

First, a Lemma that will help us later.

Lemma 1.21.1. Let a be an integer and j a natural number. Then $a \equiv a \cdot 10^{j} \pmod{3}$.

Proof. Notice that $1 \equiv 10 \pmod{3}$. Then, by 1.18, we find $1^j \equiv 10^j \pmod{3}$ and thus that $1 \equiv 10^j \pmod{3}$. Then, since $a \equiv a \pmod{3}$ (by 1.9), we invoke 1.14 to find $a \cdot 1 \equiv a \cdot 10^j \pmod{3}$, implying that $a \equiv a \cdot 10^j \pmod{3}$.

Now we begin our proof of the theorem in full.

Proof. Notice that n can be written as $a_k \cdot 10^k + a_{k-1} \cdot 10^{k-1} + \cdots + a_1 \cdot 10 + a_0$, or more easily as

$$n = \sum_{i=0}^{k} a_i \cdot 10^i$$

Now notice that

$$m = \sum_{i=0}^{k} a_i$$

By 1.21.1, we notice that $\forall i \ a_i \equiv a_i \cdot 10^i \pmod{3}$. Thus, n and m are sums of terms that are congruent modulo 3. By repeatedly invoking 1.12, we eventually find that the two strings of congruent sums are themselves are congruent, i.e. that $n \equiv m \pmod{3}$.

Theorem 1.22. If a natural number is divisible by 3, then, when expressed in base 10, the sum of its digits is divisible by 3.

Proof. Let the natural number be n, and the sum of its digits m. We're given by the theorem $n \equiv 0 \pmod{3}$, and by 1.21 we know $n \equiv m \pmod{3}$, so we can cite 1.11 and conclude $m \equiv 0 \pmod{3}$, i.e. m is divisible by 3.

Theorem 1.23. If the sum of the digits of a natural number expressed in base 10 is divisible by 3, then the number is divisible by 3 as well.

Proof. Let the natural number be n, and the sum of its digits m. We're given by the theorem $m \equiv 0 \pmod{3}$, and by 1.21 we know $n \equiv m \pmod{3}$, so we can cite 1.11 and conclude $n \equiv 0 \pmod{3}$, i.e. n is divisible by 3.

Exercise 1.24. Devise and prove other divisibility criteria similar to the preceding one.

A number is divisible by 2 if and only if its last digit is divisible by 2, because any (base 10) number $n = a_k a_{k-1} \dots a_1 a_0 = a_k a_{k-1} \dots a_1 \cdot 10 + a_0$, and 2|10 so $2|\dots \cdot 10$. Thus, $2|\dots \cdot 10 + a_0$ iff $2|a_0$.

Similar proofs can be done for 5 and the last digit, 4 and the last 2 digits, 8 and the last 3 digits, 16 and the last 4 digits, 32 and the last 5 digits, etc.

The Division Algorithm

Exercise 1.25. Illustrate the division algorithm for:

1.
$$m = 25$$
, $n = 7$.
 $25 = 7 \cdot 3 + 4$.

2.
$$m = 277$$
, $n = 4$. $277 = 4 \cdot 69 + 1$.

3.
$$m = 33$$
, $n = 11$. $33 = 11 \cdot 3 + 0$.

4.
$$m = 33$$
, $n = 45$. $33 = 44 \cdot 0 + 33$.

Theorem 1.26. Prove the existence part of the Division Algorithm. In other words, given natural numbers n and m, show there exist integers q and r such that m = nq + r and $0 \le r \le n - 1$.

Proof. Let $S = \{x \in \mathbb{Z} \mid nx > m\}$. By the Well-Ordering Axiom, S has a smallest element: call it s. Let q = s - 1. This definition gives us two important properties:

- 1. $nq \le m$, for if nq > m then $q \in S$ with q < s, which is impossible since s is the smallest element of S.
- 2. m < n(q+1) = nq + n, for q+1 = s and sx > m because $s \in S$.

Now, we define r = m - nq, so that by definition m = nq + r. Since $nq \le m$, we know $r \ge 0$. Since m < nq + n, and yet m = nq + r, implying $nq + r < nq + n \implies r < n \implies r \le n - 1$.

Thus, we have found q, r such that m = nq + r and $0 \le r \le n - 1$.

Theorem 1.27. Prove the uniqueness part of the Division Algorithm. In other words, given natual numbers n and m, if there are 4 integers q, q', r, and r', such that m = nq + r = nq' + r' with $0 \le r, r' \le n - 1$ then q = q' and r = r'.

Proof. Notice that nq + r = nq' + r' implies that $nq - nq' = r' - r \implies n(q - q') = r' - r$.

Since $0 \le r, r' \le n-1$, we conclude that $-n+1 \le r'-r \le n-1$. By our previous equality, then, $-n+1 \le n(q-q') \le n-1 \implies -n < n(q-q') < n$. Since n is a natural number, we can divide by n to get -1 < q-q' < 1. Since q and q' are integers, q-q' must also be an integer. The only integer between -1 and 1 is 0, so we conclude $q-q'=0 \implies q=q'$.

Once we have q = q', we see that $nq + r = nq' + r' \implies nq + r = nq + r' \implies r = r'$.

Theorem 1.28. Let a, b, and n be integers with n > 0. Then $a \equiv b \pmod{n}$ if and only if a and b have the same remainder when divided by n. Equivalently, $a \equiv b \pmod{n}$ if and only if when $a = nq_1 + r_1 \pmod{n}$ and $b = nq_2 + r_2 \pmod{n}$ then $r_1 = r_2$.

First, we will show that $a \equiv b \pmod{n} \implies r_1 = r_2$.

Proof. Notice by the definition of modular congruence that $a \equiv b \pmod{n}$ implies that n|(b-a), or $\exists d \in \mathbb{Z} \ni nd = b-a$. Using $a = nq_1 + r_1$ and $b = nq_2 + r_2$ we get $nd = nq_1 + r_1 - nq_2 - r_2 = n(q_1 - q_2) + r_1 - r_2$. Then we get $nd - n(q_1 - q_2) = r_1 - r_2$ or $n(d - q_1 + q_2) = r_1 - r_2$.

Since $0 \le r_1, r_2 \le n-1$ we find that $-n+1 \le r_1-r_2 \le n-1 \implies -n < r_1-r_2 < n$. Using our previous equation with r_1-r_2 we get that $-n < n(d-q_1+q_2) < n$, and dividing by n (which we can do because n > 0) we get $-1 < d-q_1+q_2 < 1$. Since d, q_1 , and q_2 are all integers, $d-q_1+q_2$ is also an integer, and the only integer between -1 and 1 is 0 so we find $d-q_1+q_2=0$.

Plugging this back in to $n(d-q_1+q_2)=r_1-r_2$, we find $n\cdot 0=r_1-r_2$, which implies $0=r_1-r_2$, or $r_1=r_2$. \square

Second, we will show that $r_1 = r_2 \implies a \equiv b \pmod{n}$.

Proof. Notice $a - b = nq_1 + r_1 - (nq_2 + r_2)$. With some simple rearranging, we obtain $a - b = n(q_1 - q_2) + r_1 - r_2$. Since we know $r_1 = r_2$, we know $r_1 - r_2 = 0$, and plugging this in we obtain $a - b = n(q_1 - q_2)$.

Since q_1 and q_2 are integers, $q_1 - q_2$ is also an integer. Thus, n times some integer is a - b: in other words, n|(a - b).

Then, by the definition of modular congruence, we obtain $a \equiv b \pmod{n}$.

Greatest common divisors and linear Diophantine equations

Question 1.29. Do every two integers have at least one common divisor?

Yes. For any two integers a and b, $1 \cdot a = a$ and $1 \cdot b = b$ so 1|a and 1|b, making 1 a common divsor of a and b.

Question 1.30. Can two integers have infinitely many common divisors?

No, if the two integers are distinct. Any nonzero integer n can only have finitely many divisors, as any integer d such that d < -|n| or d > |n| cannot be a divisor (since 1d and -1d have a greater absolute value than n, and $0d = 0 \neq n$). In other words, only the numbers f such that $-n \leq f \leq n$ are "eligibile" to be divisors of n, so there can only be finitely many divisors of n.

Exercise 1.31. Find the following greatest common divisors. Which pairs are relatively prime?

- 1. (36, 22) 2
- 2. (45, -15)15
- 3. (-296, -88)
- 4. (0,256)256
- 5. (15, 28) 1 (relatively prime)
- 6. (1, -2436) 1 (relatively prime)

Theorem 1.32. Let a, n, b, r, and k be integers. If a = nb + r and k|a and k|b, then k|r.

Proof. Let $a = d_a k$ and $b = d_b k$, where d_a and d_b are the integers guaranteed by the facts that k|a and k|b. Then, we have $d_a k = n d_b k + r$. Isolating r, we get $r = d_a k - n d_b k = k(d_a - n d_b)$. Since n, d_a , and d_b are all integers, we know $d_a - n d_b$ is an integer. Thus, we've found r is equal to k times some integer, so k|r.

Theorem 1.33. Let $a, b, n_1, and r_1$ be integers with a and b not both b. If $a = n_1b + r_1$, then $(a, b) = (b, r_1)$.

Proof. We will show that the common divisors of a and b are the same as the common divisors of b and r_1 , and thus conclude that the greatest element of S is also the greatest element of T.

Let S be the set of common divisors of a and b, and let T be the set of common divisors of b and r_1 . We will show S = T by double inclusion.

First, let's show $S \subset T$. Take an arbitrary $s \in S$. Since s|a and s|b, we conclude $\exists d_a, d_b \in \mathbb{Z} \ni a = sd_a, b = sd_b$. We can then rearrange $a = n_1b + r_1$ to read $r_1 = a - n_1b$, and then plug in our previous two equations to get $r_1 = sd_a - n_1sd_b \implies r_1 = s(d_a - n_1d_b)$. Since d_a , d_b , and n_1 are all integers, we know $d_a - n_1d_b$ is an integer, thus implying that $s|r_1$. Since we know s|b since $s \in S$, we conclude $s \in T$. Thus, any arbitrary $s \in S$ is an element of T, so $S \subset T$.

Showing that $T \subset S$ proceeds in much the same way. Take $t \in T$, conclude since t|b and $t|r_1$ we find $\exists d_b d_r \in \mathbb{Z} \ni b = td_b, r_1 = td_r$, and then plug those in to $a = n_1b + r_1$ to get $a = n_1td_b + td_r \implies a = t(n_1d_b + d_r)$. Since n_1, d_b , and d_r are integers, we find t|a, and since t|b because $t \in T$, we thus conclude $t \in S$. Thus any arbitrary $t \in T$ is an element of S, so $T \subset S$.

Thus, by double inclusion, S = T. This implies that the greates element of S, i.e. (a,b), is equal to the greatest element of T, i.e. (b,r_1) .

Exercise 1.34. Use the preceding theorem to show that if a = 51 and b = 15, then (51, 15) = (6, 3) = 3.

Proof. Since $51 = 3 \cdot 15 + 6$, we find (51, 15), we cite 1.33 to see (51, 15) = (15, 6). Then, since $15 = 2 \cdot 6 + 3$, we again cite 1.33 to find (15, 6) = (6, 3). We see that (6, 3) = 3 by inspection. Then, since equality is transitive, we conclude (51, 15) = (6, 3) = 3.

Exercise 1.35. Using the previous theorem and the Division Algorithm successively, devise a procedure for finding the greatest common divisor of two integers.

Well you kind of gave the game away when you said to use 1.33 and the division algorithm successively huh. If you're trying to find (a, b), you simply invoke the division algorithm to get a = nb + r (assuming WLOG that $a \ge b$), and then rewrite (a, b) as (b, r). Then, you use the divion algorithm to get b = nr + r', simplifying to (r, r'), etc. etc., until at some point you have (x, 0), which by inspection is equal to x.

You will always reach (x,0) because the divison algorithm produces a remainder r that is strictly less than the smaller input b, so (informally) the smaller of the two numbers you're working with always gets smaller while never going negative.

Exercise 1.36. Use the Euclidean Algorithm to find the following.

- 1. (96, 112)112 = 1.96 + 16, simplifying the problem to (96, 16). Then 96 = 5.16 + 0, so we get (16, 0) = 16
- 2. (162,31) $162 = 5 \cdot 31 + 7 \implies (31,7) \implies 31 = 4 \cdot 7 + 3 \implies (7,3) \implies 7 = 2 \cdot 3 + 1 \implies (3,1) = 1.$
- 3. (0, 256) Since everything divides 0, this is trivially 256.
- - 5. (1, -2436)Since the only integers that divide 1 are -1, 0, and 1, we trivially find 1.

Exercise 1.37. Find integers x and y such that 162x + 31y = 1.

By division algorithm, $162 = 5 \cdot 31 + 7 \implies 7 = 1 \cdot 162 + (-5) \cdot 31$. By division algorithm, $31 = 4 \cdot 7 + 3 \implies 3 = 1 \cdot 31 + (-4) \cdot 7 = 1 \cdot 31 + (-4) \cdot (1 \cdot 162 + (-5) \cdot 31) = (-4) \cdot 162 + 21 \cdot 31$. By division algorithm, $7 = 2 \cdot 3 + 1 \implies 1 = 1 \cdot 7 + (-2) \cdot 3 = 1 \cdot (1 \cdot 162 + (-5) \cdot 31) + (-2) \cdot ((-4) \cdot 162 + 21 \cdot 31) = 9 \cdot 162 + (-47) \cdot 31$.

Thus, we've found our solution x = 9 and y = -47.

Theorem 1.38. Let a and b be integers. If (a,b) = 1, then there exist integers x and y such that ax + by = 1.

Proof. If either a or b is negative, replace it with -a or -b for the rest of this proof. At then end, you can replace either x or y with -x or -y to get an answer; for instance, if a = -3, we can replace a = 3, do the proof to obtain x_0 and y_0 such that $3x_0 + by_0 = 1$ and then realize that $(-3)(-x_0) + by_0 = 1$, which since $-x_0$ is still an integer still suffices. Now we will only be worrying about non-negative a and bs.

We will demonstrate an algorithm to find x and y. WLOG, assume $a \ge b$. Invoke the division algorithm to get $a = n_1b + r_1$. Then invoke it again to get $b = n_2r_1 + r_2$. Then invoke it again to get $r_1 = n_3r_2 + r_3$. Etc. etc. etc.

We will show that the series "remainder" generated by this algorithm eventually has to hit 0: in other words, $\exists i \in \mathbb{N} \ni r_i = 0$. To do this, we must notice that for any index j, since r_j is generated by calling the division algorithm on r_{j-2} and r_{j-1} , we find that $r_j \leq r_{j-1} - 1$. Notice, then, that we can apply this to r_{j-1} to obtain $r_{j-1} \leq r_{j-2} - 1$, and then plug that in to our previous inequality to get $r_j \leq r_{j-1} - 1 \leq r_{j-2} - 2$.

By inspection (i.e. I'm lazy and don't want to formalize this), we notice we can continually apply this. We will apply this to r_b , and notice that $r_b \le r_{b-1} - 1 \le r_{b-2} - 2 \le \cdots \le r_1 - (b-1) \le b-b$. Since b-b=0, we find $r_b \le 0$, but since r_b is a remainder from the division algorithm we know $r_b \ge 0$, so we conclude $r_b = 0$.

Notice we have not proven that r_b is the first 0, only that the remainders must eventually reach 0 at some point. Now, keep invoking the division algorithm until the "remainder" generated by the algorithm is 0: we will label that step k+1, so that we find $r_{k-1} = n_{k+1}r_k + 0$. We will show that r_k is 1.

By invoking 1.33 repeatedly, we find that $(a,b) = (b,r_1) = (r_1,r_2) = \cdots = (r_{k-1},r_k) = (r_k,r_{k+1})$. Since $r_{k+1} = 0$, we conclude $(a,b) = (r_k,0)$. Since 0 divides everything, $(r_k,0) = r_k$, so $(a,b) = r_k$, and since a and b are relatively prime we conclude $1 = r_k$.

Now, we take all of our equations and rewrite them to solve for the remainder. For example, $a = n_1b + r_1$ becomes $r_1 = a + (-n_1)b$, and $b = n_2r_1 + r_2$ becomes $r_2 = b + (-n_2)r_1$.

This gives us a bunch of equations of the form $r_j = \delta_j r_{j-2} + \gamma_j r_{j-1}$. This includes one for r_k , namely $r_k = \delta_k r_{k-2} + \gamma_k r_{k-1}$. We can then substitute in lower indices of r for r_{k-2} and r_{k-1} , using the generic equation, to get something like $r_k = \delta_k (\delta_{k-2} r_{k-4} + \gamma_{k-2} r_{k-3}) + \gamma_k (\delta_{k-1} r_{k-3} + \gamma_{k-1} r_{k-2})$.

That looks horrifying, but the important bit is that we notice if we simplify it we get $r_k = Ar_{k-4} + Br_{k-3} + Cr_{k-2}$ with $A, B, C \in \mathbb{Z}$. That is, by replacing all r_j 's with their respective equations, we have reduced the highest index on an r in the right hand side by 1. Previously, the highest index was k-1, but now it's k-2, because we had an equation to represent r_{k-1} in terms of r_{k-3} and r_{k-4} .

Notice, though, that not all r's satisfy this property: namely, r_1 and r_2 simplify down to a and b, which then don't have equations of their own. So, we apply the equations for r_k through r_1 in "reverse" order, pairing down the maximum index of k each time, until we're left with only r_1 's and r_2 's on the left hand side and can apply those equations to get a linear expression in a and b on the right hand side.

We've been talking a lot about the right hand side, but remember, the left hand side is r_k , and we've shown $r_k = 1$, so we've just found a linear expression in a and b that is equal to 1. In other words, 1 = ax + by for some $x, y \in \mathbb{Z}$.

Theorem 1.39. Let a and b be integers. If there exist integers x and y with ax + by = 1, then (a, b) = 1.

Proof. Readers of the last proof will be glad to hear this one is much simpler.

By definition, (a,b)|a and (a,b)|b. Then, (a,b)|ax and (a,b)|by by 1.6. Then, (a,b)|ax+by by 1.1. Then, since ax+by=1, we find (a,b)|1. We know 1|a and 1|b, so $(a,b) \ge 1$. The only number ≥ 1 that divides 1 is 1, so since $(a,b) \ge 1$ and (a,b)|1 we conclude (a,b)=1.

Theorem 1.40. For any integers a and b not both 0, there are integers x and y such that ax + by = (a, b).

Proof. Let c = a/(a,b) and d = b/(a,b). Notice that since (c,d)|c and $a = c \cdot (a,b)$ we find $((c,d) \cdot (a,b))|a$, and similarly since (c,d)|d and $b = d \cdot (a,b)$ we find $((c,d) \cdot (a,b))|b$.

Since c and d are integers not both 0, (c,d) must be a positive integer. Since $(c,d) \cdot (a,b)$ is a common factor of a and b, and (a,b) is the *greatest* common factor of a and b, we find $(c,d) \cdot (a,b) \leq (a,b) \implies (c,d) = 1$.

Thus, we invoke 1.38 to find integers x and y such that cx + dy = 1. Then, we multiply both sides by (a, b) to find that $(a, b) \cdot cx + (a, b) \cdot dy = (a, b)$. Since $a = (a, b) \cdot c$ and $b = (a, b) \cdot d$ we conclude ax + by = (a, b).