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1 Problem

Maximise $Z = -x + 2y$, subject to the constraints:
 $x \geq 3$, $x + y \geq 5$, $x + 2y \geq 6$, $y \geq 0$.

2 Solution

Objective function:

$$z = \max_x (-1 \ 2) x \quad (1)$$

constraints:

$$\begin{aligned} x+y &\geq 5 \\ x+2y &\geq 6 \\ x &\geq 3 \\ y &\geq 0 \end{aligned}$$

Writing all constraints in the matrix form

$$px = q \quad (2)$$

$$\begin{pmatrix} 1 & 1 \\ 1 & 2 \\ 1 & 0 \\ 0 & 1 \end{pmatrix} x = \begin{pmatrix} 5 \\ 6 \\ 3 \\ 0 \end{pmatrix} \quad (3)$$

By providing the objective function and constraints to cvxpy, the optimal value gives infinity as result and the problem is unbounded.

Reason: Unbounded means if there exists some direction within the feasible region along which the objective function value can increase (maximization case) or decrease (minimization case) without bound. In such a formulation, the optimal value is negative infinity for a minimization problem, and conversely, positive infinity for a maximization problem.

An unbounded solution is something that typically does not arise in practical applications. When it does occur, it's usually because the formulation is ill-posed, i.e., incorrect in some way, and/or missing some necessary constraints for properly modeling the dynamics of the system under consideration.

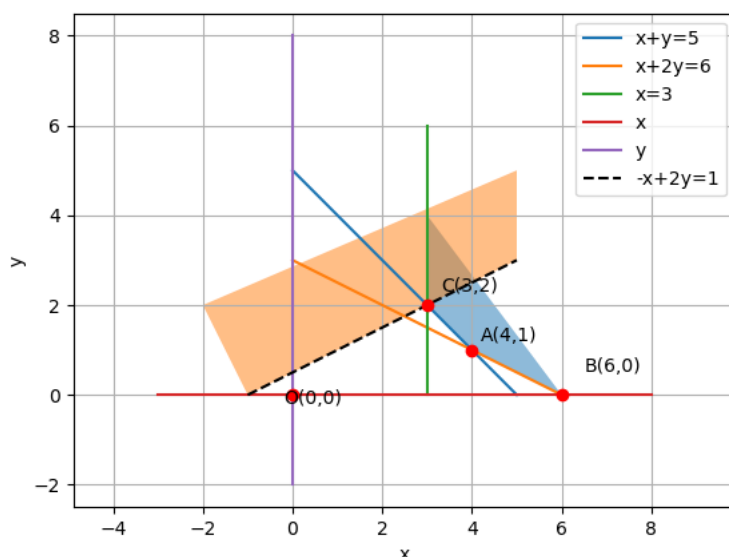
termux commands :

```
bash basic.sh.....using shell command
```

cvxpy code:

```
https://github.com/FWC\_module1/blob/main/optimization/basic.py
```

Graphical method:



From Graph,

corner points	z
(3,2)	1
(4,1)	-2
(6,0)	-6

Clearly the corner point (3,2) has maximum value for the given objective function.

But, as the feasible region is unbounded, '1' may or may not be the maximum value. So, we need to graph the inequality $-x+2y > 1$.

\therefore Feasible region of $-x+2y > 1$ has some points in common with the given constraints. So, there is no maximum value for z subject to given constraints.