Array Element Assignments

CS 536: Science of Programming, Fall 2022

A. Why?

Array assignments aren't like assignments to plain variables because the actual item to change
can't be determined until runtime. We can handle this by extending our notion of assignment
and/or substitution.

B. Outcomes

At the end of this work you should:

- Be able to perform textual substitution to replace an array element.
- Be able to calculate the wp of an array element assignment.

C. Questions

For each of the questions below, calculate the given weakest precondition. Then try logically simplifying it to something easier to read.

- 1. Calculate wp(b[0] := 9, x > b[k]).
- 2. Calculate wp(b[k] := b[j], b[m] = 0).
- 3. Calculate wp(b[k] := b[j], b[j] = z).
- **4.** Calculate wp(b[k] := 1, b[k] = b[j]).
- 5. Calculate $wp(b[k] := x; b[j] := y, b[k] \neq b[j])$.
- 6. Calculate $wp(b[k] := x, b[b[k]] \neq b[k])$.
- 7. Is the triple $\{k < b[k] < b[i]\}$ b[b[k]] := b[i] $\{b[k] \neq b[i]\}$ valid?
- 8. Define a predicate function swapped(b_1 , b_2 , k, j) that yields true iff b_1 and b_2 are equal except their values at k and j are swapped.
- 9. Let the function search(b, m, n, x) = k where if $m \le k \le n$ then b[k] = x and if $k < m \lor k > n$, then no index for x exists in the range m ... n. (I.e., for all j where $m \le j \le n$, we have $b[j] \ne x$.) Complete the predicate $\exists k$ search(b, m, n, x) = $k \rightarrow ...$ where the ellipsis formalizes this description of what search does.
- 10. In Example 5 in the notes, we had (b[b[k]])[5/b[0]]

```
= if (if k = 0 then 5 else b[k] fi) = 0 then 5 else b[ if k = 0 then 5 else b[k] fi ] fi
```

Show how to optimize this to *if* k = 0 *then* b[5] *else if* b[k] = 0 *then* 5 *else* b[b[k]] *fi fi.*

Solution to Practice 21 (Array Element Assignments)

```
1. wp(b[0] := 9, x > b[k]) = (x > b[k])[9/b[0]]
          = x > if k = 0 then 9 else b[k] fi
          \Leftrightarrow if k = 0 then x > 9 else x > b[k] fi
          \Leftrightarrow (k = 0 \land x > 9) \lor (k \neq 0 \land x > b[k])
2. wp(b[k] := b[i], b[m] = 0) = (b[m] = 0)[b[i]/b[k]]
          = if m = k then b[i] else b[m] fi = 0
          \Leftrightarrow if m = k then b[i] = 0 else b[m] = 0 fi
                                                                     (one possible alternative rewriting)
          \Leftrightarrow (m = k \rightarrow b[i] = 0) \land (m \neq k \rightarrow b[m] = 0)
                                                                     (another possible alternative rewriting)
3. wp(b[k] := b[i], b[i] = z) = (b[i] = z)[b[i]/b[k]]
          = if | = k then b[i] else b[i] fi = z
          \Leftrightarrow b[i] = z
4. wp(b[k] := 1, b[k] = b[i]) = (b[k] = b[i])[1/b[k]]
          = (b[k])[1/b[k]] = (b[i])[1/b[k]]
          = 1 = (if \mid = k then 1 else b[i] fi)
          \Leftrightarrow k = j \vee b[j] = 1 (\Leftrightarrow k \neq j \rightarrow b[j] = 1 if you prefer \rightarrow to \vee)
5. wp(b[k] := x; b[j] := y, b[k] \neq b[j])
          = wp(b[k] := x, wp(b[j] := y, b[k] \neq b[j])
     For the embedded wp, wp(b[i] := y, b[k] \neq b[i])
          = (b[k] \neq b[i])[y/b[i]]
          = b[k] [y/b[j]] \neq b[j] [y/b[j]]
          = if k = i then y else b[k] fi \neq y
          \Leftrightarrow k \neq i \land b[k] \neq y
     So wp(b[k] := x, wp(b[j] := y, b[k] \neq b[j])
          \Leftrightarrow wp(b[k] := x, k \neq j \land b[k] \neq y)
          = (k \neq j \land b[k] \neq y)[x/b[k]]
          \equiv k \neq j \land b[k][x/b[k]] \neq y
          = k \neq i \land x \neq y
```

Intuitively, *if* k = j, *then* we can't have $b[k] \neq b[j]$. Even *if* $k \neq j$, we need $x \neq y$ to ensure that the assignments b[k] := x; b[j] := y make $b[k] \neq b[j]$.

6. For $wp(b[k] := x, b[b[k]] \neq b[k])$, let's first look at substituting x for b[k] in b[b[k]]. It's complicated because we have to recursively substitute in the index of the outer b[...]. The general rule is

```
(b[e_2])[e_1/b[e_0]] \equiv (if e_2' = e_0 then e_1 else b[e_2'] fi) where e_2' \equiv (e_2)[e_1/b[e_0]]
```

```
So let e' = (b[k])[x/b[k]] = x, so
(b[b[k]])[x/b[k]]
= if e' = k then \times else b[e'] fi
= if \times = k then \times else b[x] fi
Then
wp(b[k] := x, b[b[k]] \neq b[k])
= (b[b[k]] \neq b[k])[x/b[k]]
= (b[b[k]])[x/b[k]] \neq (b[k])[x/b[k]]
= if \times = k then \times else b[x] fi \neq x
\Leftrightarrow x \neq k \wedge b[x] \neq x
```

7. For the triple to be valid, it's sufficient to show that its precondition implies the wp of the assignment and postcondition. I.e., $k < b[k] < b[j] \rightarrow wp(b[b[k]] := b[j], b[k] \neq b[j])$ First let's calculate $wp(b[b[k]] := b[j], b[k] \neq b[j])$

```
= (b[k] \neq b[j])[b[j] / b[b[k]]]
= (b[k])[b[j] / b[b[k]]] \neq (b[j])[b[j] / b[b[k]]]
= if k = b[k] then b[j] else b[k] fi \neq if j = b[k] then b[j] else b[j] fi
\Leftrightarrow if k = b[k] then b[j] else b[k] fi \neq b[j]
\Leftrightarrow if k = b[k] then b[j] \neq b[j] else b[k] \neq b[j] fi
\Leftrightarrow k \neq b[k] \land b[k] \neq b[j]
```

Going back to our original question, if the implication below is valid, then our triple is valid (because we have the precondition of the triple implying the wp of its body and postcondition).

```
k < b[k] < b[j] \rightarrow wp(b[b[k]] := b[j], b[k] \neq b[j])

\Leftrightarrow k < b[k] < b[j] \rightarrow k \neq b[k] \land b[k] \neq b[j]

\Leftrightarrow T
```

So our original triple is indeed valid.

- 8. swapped(b₁, b₂, k, j) = b₁[k] = b₂[j] \wedge b₁[j] = b₂[k] \wedge (\forall m.(m \neq k \wedge m \neq j \rightarrow b₁[m] = b₂[m]))
- 9. \exists k.search(b, m, n, x) = k \rightarrow (m ≤ k ≤ n \rightarrow b[k] = x) \land (m > k \lor k > n \rightarrow (\forall j.m \le j \le n \rightarrow b[j] \neq x))

(Extra parentheses added for readability)

- **10.** (Optimize (b[b[k]])[5/b[0]])
 - 1. We start with if (if k = 0 then 5 else b[k] fi) = 0 then 5 else b[if k = 0 then 5 else b[k] fi] fi

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2. First, let's optimize the test for the outer *if*:

(if
$$k = 0$$
 then 5 else $b[k]$ fi) = 0
 \mapsto if $k = 0$ then 5 = 0 else $b[k] = 0$ fi
 \mapsto if $k = 0$ then F else $b[k] = 0$ f
 \mapsto $k \neq 0 \land b[k] = 0$

3. Then let's optimize the false branch of the outer *if*:

$$b[if k = 0 then 5 else b[k] fi]$$

 $\mapsto if k = 0 then b[5] else b[b[k]] fi$

4. If we substitute (2) and (3) back into (1), we get

if (if
$$k = 0$$
 then 5 else $b[k]$ fi) = 0 then 5 else $b[$ if $k = 0$ then 5 else $b[k]$ fi f if f if

- 5. A little case analysis tells us that we want
 - b[5] when k = 0 (regardless of $b[k] = or \neq 0$)
 - 5 when $k \neq 0 \land b[k] = 0$
 - b[b[k]] when $k \neq 0 \land b[k] \neq 0$
- 6. This gives us

if
$$k = 0$$
 then $b[5]$
else /* $k \neq 0$ */ if $b[k] = 0$ then 5
else /* $k \neq 0 \land b[k] \neq 0$ */ $b[b[k]]$ fi fi