

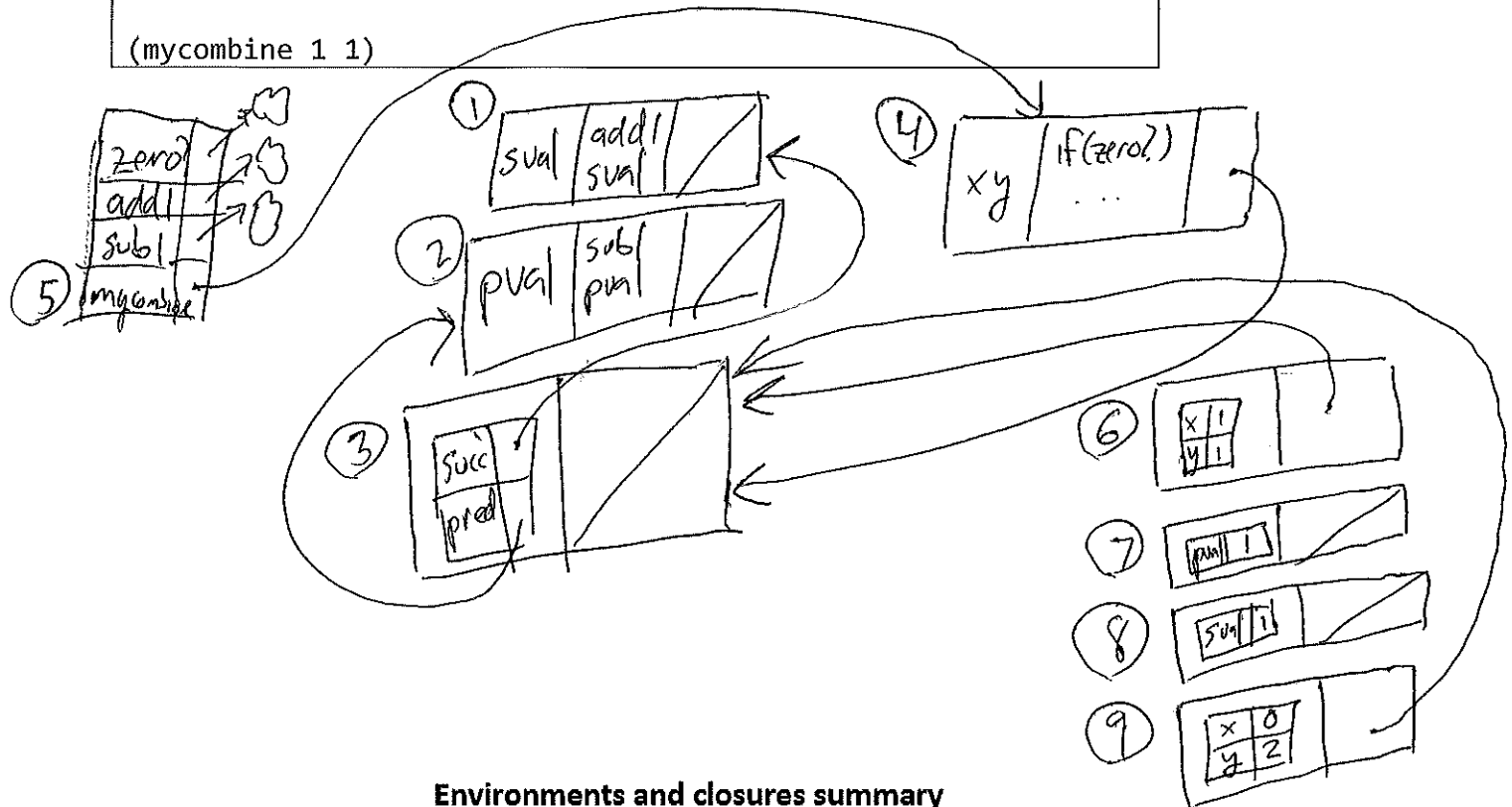
(20 points) Draw the environments and closures created during the execution of the following code segment. Include a sequence number for each new environment, new closure, or global environment change.

If you need to start over, use the back, and put a big X through what you wrote here.

Value returned by the expression: _____

```
(define mycombine
  (let ([succ (lambda (sval) (add1 sval))]
        [pred (lambda (pval) (sub1 pval))])
    (lambda (x y)
      (if (zero? x) y (mycombine (pred x) (succ y))))))

(mycombine 1 1)
```



Environments and closures summary

Environment:

env	env	Points to enclosing environment
env	env	
env	env	

Closure:

list of formal argument names	code (body of the procedure)	local environment that existed when the procedure was created
-------------------------------	------------------------------	---

A procedure (a.k.a. closure) is created when a `lambda` expression is evaluated. The body of the closure is not evaluated at this time.

Application of a closure (user-defined procedure):

- The expressions for the procedure and its arguments are evaluated.
- A new local environment is created.
 - Each variable from the procedure's formal parameter list is bound to the corresponding value in the actual argument list.
 - The new environment's "pointer to an enclosing environment" is set to be a copy of the local environment pointer that is the third part of the closure.
- The body of the procedure is evaluated, using this new local environment. If a variable is not found in this local environment or something it points to, look in the global environment. If not in global environment either, it is an error.

Evaluate a `let` expression:

- Evaluate (in the current environment) the expressions to get the values to be assigned to the `let` variables.
- Create a new local environment with bindings for the `let` variables. The "enclosing environment" pointer points to the current environment.
- Evaluate the body of the `let` in this new environment, as in 3 above.

Evaluate a `letrec` expression:

- Create a new local environment, similar to a `let` environment, except that:
 - The "saved environment" pointers of any closures that are bound to the `letrec` variables point to the new `letrec` environment, not the enclosing environment.
- Evaluate the body of the `letrec` in this new environment, as in 3 above.