

FOOD → MONEY

FOOD

↓

MOOD → MOND

↓

MONEY → MONEY

Operations.
 - edit a letter
 - delete a letter
 - add a letter

→ FOOD → MONEY } Edit operation (3).

FOOD → MOOD → MOO

↓

MONEY ← MONE ← MON

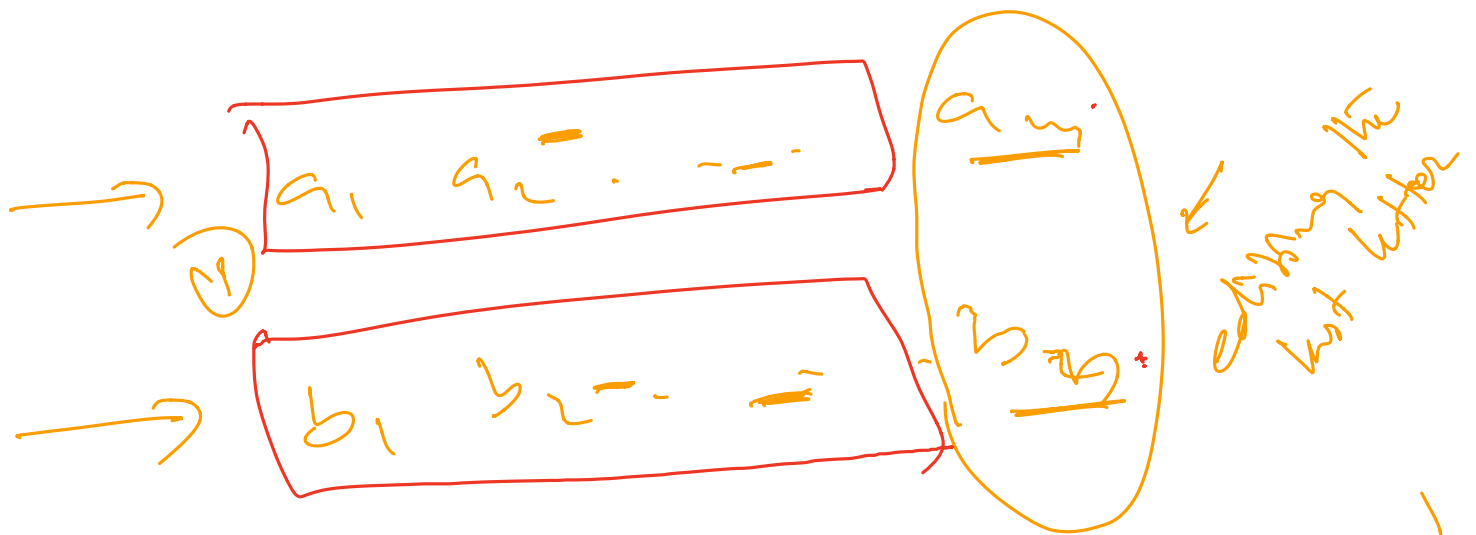
FOOD
 MON

Edit operation (3)

FOOD - -
MON - - -

Input: Two sequences of letters
 $A[1..n]$ and $B[1..m]$

Output: The number of minimum
edit operation required
to convert $A[1..n]$ to $B[1..m]$



$a_m \neq b_m$

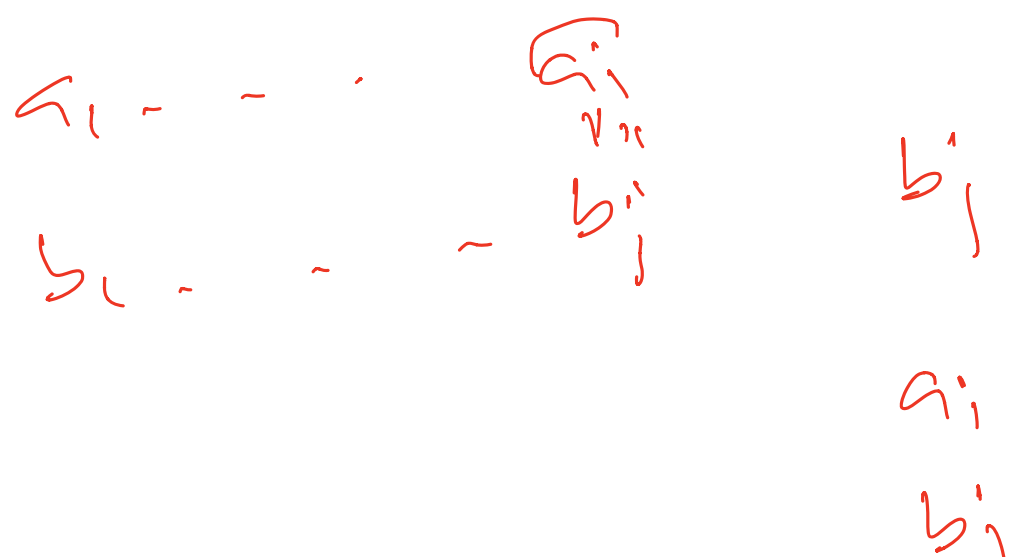


$$[b_1 - b_2 - \dots - b_n] =$$

$Edit(i, j)$ = min number of operations required to convert $[a_1 \dots i]$ to $[b_1 \dots j]$.

$Edit(m, n)$ — Our answer.

$$Edit(i, j) = \begin{cases} \min \begin{cases} (1 + Edit(i-1, j-1)) & \text{if } a_i \neq b_j \\ 1 + Edit(i, j-1) \\ 1 + Edit(i-1, j) \end{cases} & \text{if } i > 0 \text{ and } j > 0 \\ Edit(i-1, j-1) & \text{if } i=0 \text{ and } j=0 \\ Edit(i, j-1) + 1 & \text{if } i=0 \text{ and } j > 0 \\ Edit(i-1, j) + 1 & \text{if } i > 0 \text{ and } j=0 \end{cases}$$



$a_1 \dots$

$a_i -$

~~b_j~~ ~~b_j~~

$a \leq c \leq d \rightarrow \underline{b} \leq \underline{c} \leq \underline{d}$

delete a to $b \leq c \leq d$.

① Come-up with rec. formula.
[Exc. proof].

② $i \rightarrow 1$ to ∞
 $j \rightarrow 1$ to ∞

$\text{edit}(i, j)$

③ DS. $\text{EDIT}[n+1][n+1]$
 $\text{EDIT}[i, j] \leftarrow \text{edit}(i, j)$

EDITDISTANCE($A[1..m], B[1..n]$):

for $j \leftarrow 0$ to n

$Edit[0, j] \leftarrow j$

base case

for $i \leftarrow 1$ to m

$Edit[i, 0] \leftarrow i$

base case

for $j \leftarrow 1$ to n

$ins \leftarrow Edit[i, j-1] + 1$

$del \leftarrow Edit[i-1, j] + 1$

if $A[i] = B[j]$

$rep \leftarrow Edit[i-1, j-1]$

else

$rep \leftarrow Edit[i-1, j-1] + 1$

$Edit[i, j] \leftarrow \min\{ins, del, rep\}$

return $Edit[m, n]$

Time complexity

$$T(n) = O(n \cdot m)$$

Space complexity

$$S(n) = O(n \cdot m)$$

D.P

①

Write recur. formula
[proof]

②

Find out evaluation order

③

Find out suitable D.S

④

compute the values
as prescribed

in the system
and store/use in the
D.S.



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