Introduction to Singly Linked List

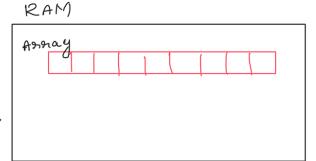
Day 35 14/08/2025

As we know the Asseaux are fixed In size, if any new data to be added to the array then the New assay with more size is onested and Elements are Copped to 9t.

To address this issue (using Memory in Efficiently) The Linkedlist is designed.

What are linked list? RAM

In cash of Annay Continous Memory Block in the Memory is allocated



address limitations of armay

- (1) fixed size (Not dynamic)
- (3) Mostage of the mostly Space (3) avoid copyring of the elements

Linked List

Me allocate Memosy for 1 Element at a fine In case of Java & Python Objects are created In Heap Membery

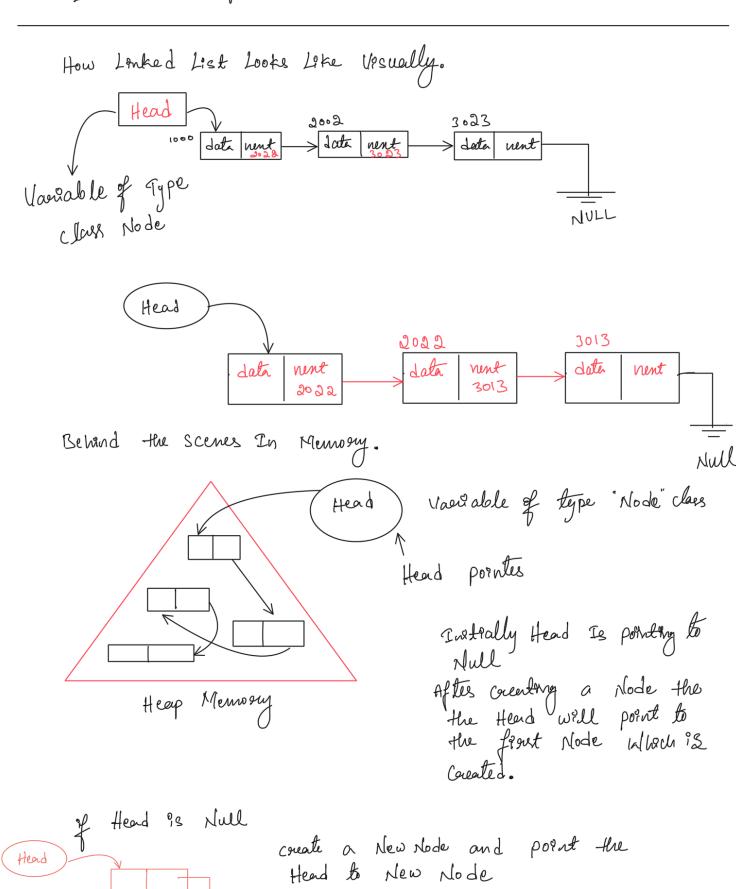
When We Great objects of NODE'S then the Memory get allocated In the Heap

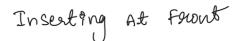
Node: 1 Element for Which Memory allocated dynamically In the RAM

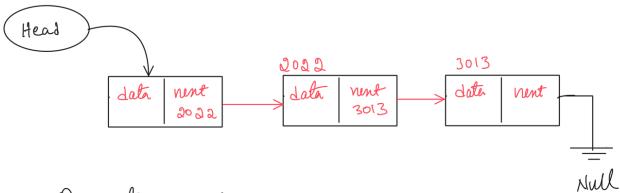
7 Contains different types of Node contains -> Reference, address of Next Mode

also Have address of pregrous No de

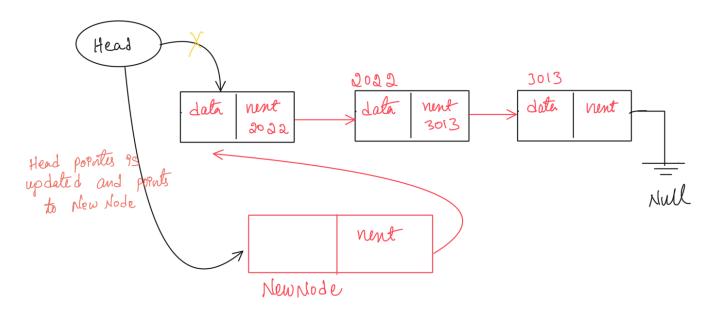
NOTE : 4 Data Depends on Requisiements.







- (1) caeate New Node
- point New Mode. nent to Heard update Heard to New Mode



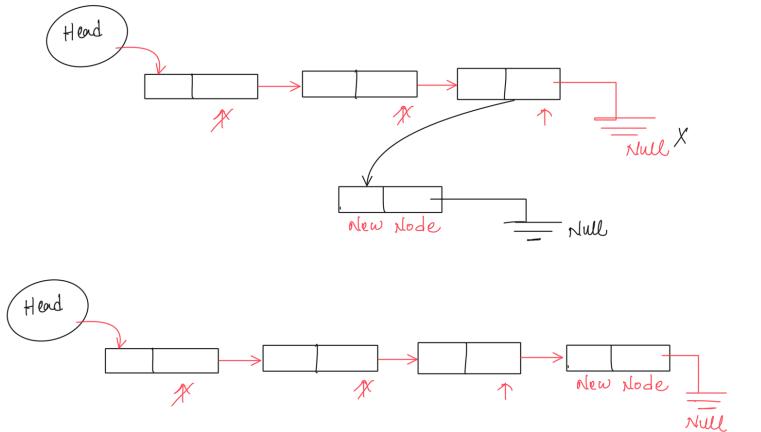
The Node 95 successfully added to the buked lest

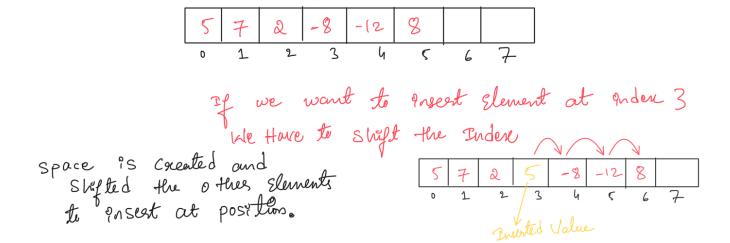
New Lanked List Looks Lake Head

Insenting At End

Traverse from Head Node 1911 we get a last Node. Last Node -> Have nent address Null, last Node doesn't Have Any Nerghbour.

ASSign New Mode Address to the Last Mode's Hent.

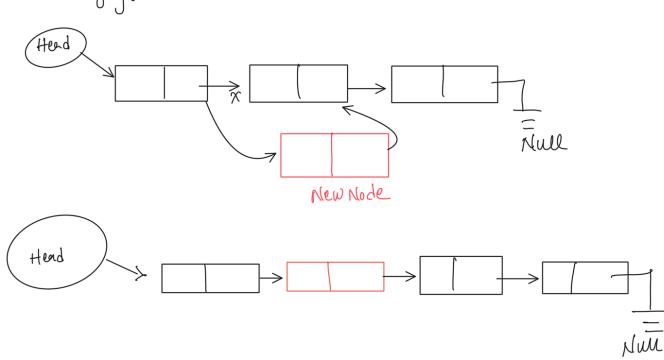




In Case of Annays

Inserting In the Middle

Inserting Elements at the Middle of Linked List 95 More fine Efficient Its Turk Cherrying quescumies accordingly.



| The Control of the Co | | |
|--|-------|-------------------------|
| Feature | Array | Linked List |
| Access time | O(1) | O(n) |
| Insertion/Deletion | O(n) | O(1) (if pointer known) |
| Memory usage | Low | Higher (extra pointers) |
| Size flexibility | Fixed | Dynamic |
| Cache friendliness | High | Low |
| | | |

Annual Annual Annual Annual Access (Random Access) Access any element in O(1) time using index. Example: arr[5] is instant. Memory Locality Elements are stored contiguously, improving CPU cache performance. Ease of Use Simple to implement and use for fixed-size collections. Low Memory Overhead Only stores the actual data (no extra pointers like in linked lists).

✓ Advantages 1. Dynamic Size • Can grow/shrink at runtime without resizing. 2. Efficient Insertion/Deletion • Inserting/removing at the beginning or middle is O(1) (if you already have the pointer). 3. No Wasted Space for Fixed Size • Memory is allocated as needed. ✓ Disadvantages 1. No Random Access • Must traverse nodes one by one (O(n)) to access element.



2. Extra Memory Overhead

• Each node stores extra pointer(s) (next, sometimes prev).

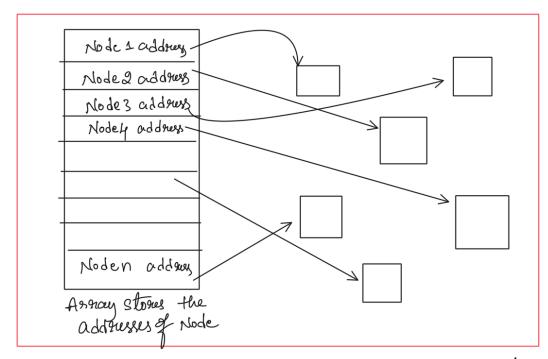
3. Poor Cache Performance

• Elements may be scattered in memory, reducing CPU cache efficiency.

4. More Complex Implementation

• Harder to manage (need careful handling of pointers to avoid memory leaks).

Annay List



To overcome the desadvantages of Lenbedlest Assemblest Come ente pecture.

Array list -> Combination of Annay and Linked List