Encoding Higher Inductive Types Without Boilerplate

A Study in Agda Metaprogramming

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Abstract

Higher inductive types are inductive types that include non-trivial higher-dimensional structure, represented as identifications that are not reflexivity. While work proceeds on type theories with a computational interpretation of univalence and higher inductive types, it is convenient to encode these structures in more traditional type theories with mature implementations. However, these encodings involve a great deal of error-prone additional syntax. We present a library that uses Agda's metaprogramming facilities to automate this process, allowing higher inductive types to be specified with minimal additional syntax.

Keywords Higher inductive type, Elaboration, Elimination rules, Computation rules

1 Introduction

Type theory unites programming and mathematics in a delightful synthesis, in which we can write programs and proofs in the same language. Work on higher-dimensional type theory has revealed a beautiful higher-dimensional structure, lurking just beyond reach. In particular, higher inductive types provide a natural encoding of many otherwise-difficult mathematical concepts, and univalence lets us work in our type theory the way we do on paper: up to isomorphism. Homotopy type theory, however, is not yet done. We do not yet have a mature theory or a mature implementation.

While work proceeds on prototype implementations of higher-dimensional type theories [4, 10], much work remains before they will be as convenient for experimentation with new ideas as Coq, Agda, or Idris is today. In the meantime, it is useful to be able to experiment with ideas from higher-dimensional type theory in our existing systems. If one is willing to put up with some boilerplate code, it is possible to encode higher inductive types and univalence using a mixture of postulated identities and traditional datatypes. We use a technique developed by Licata [17].

Boilerplate postulates, however, are not just inconvenient, they are also an opportunity to make mistakes. Luckily, this boilerplate code can be mechanically generated using Agda's recent support for *elaborator reflection* [7], a paradigm for metaprogramming in an implementation of type theory. An elaborator is the part of the implementation that translates a convenient language designed for humans into a much simpler, more explicit, verbose language designed to be easy for a machine to process. Elaborator reflection directly exposes the primitive components of the elaborator to metaprograms written in the language being elaborated, allowing them to put these components to new uses.

Using Agda's elaborator reflection, we automatically generate the support code for many useful higher inductive types, specifically those that include additional paths between constructors, but not paths between paths. We automate the production of the datatype definitions, postulated paths, induction principles, and their computational behavior. Angiuli et al.'s encoding of patch theory as a higher inductive type [5] requires approximately 1500 lines of code when represented using Licata's technique. Using our library, the encoding can be expressed in just 70 lines.

This paper makes the following contributions:

- We describe the design and implementation of a metaprogram that automates an encoding of higher inductive types with one path dimension using Agda's new metaprogramming system.
- We demonstrate applications of this metaprogram to examples from the literature, including both standard textbook examples of higher inductive types as well as larger systems, including both patch theory and specifying cryptographic schemes.
- This metaprogram serves as an example of the additional power available in Agda's elaborator reflection relative to earlier metaprogramming APIs.

2 Background

2.1 Agda Reflection

Agda's reflection library enables compile-time metaprogramming. This reflection library directly exposes parts of the implementation of Agda's type checker and elaborator for use by metaprograms, in a manner that is similar to Idris's elaborator reflection [6, 7] and Lean's tactic metaprogramming [14]. The type checker's implementation is exposed as effects in a monad called `TC`.

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```
macro
  mc1 : Term → Term → TC ⊤
  mc1 exp hole =
    do exp' ← quoteTC exp
       unify hole exp'

sampleTerm : Term
sampleTerm = mc1 (λ (n : Nat) → n)
```

Figure 1. A macro that quotes its argument

Agda exposes a representation of its syntax to metaprograms, including datatypes for expressions (called `Term`) and definitions (called `Definition`). The primitives exposed in `TC` include declaring new metavariables, unifying two `Terms`, declaring new definitions, adding new postulates, computing the normal form or weak head normal form of a `Term`, inspecting the current context, and constructing fresh names. This section describes the primitives that are used in our code generation library; more information on the reflection library can be found in the Agda documentation [1].

`TC` computations can be invoked in three ways: by macros, which work in expression positions, using the `unquoteDecl` operator in a declaration position, which can bring new names into scope, and using the `unquoteDef` operator in a declaration position, which can automate constructions using names that are already in scope. This preserves the principle in Agda's design that the system never invents a name.

An Agda *macro* is a function of type $t_1 \to t_2 \to \ldots \to$ `Term` $\to$ `TC` $\top$ that is defined inside a `macro` block. Macros are special: their last argument is automatically supplied by the type checker, and consists of a `Term` that represents the metavariable to be solved by the macro. If the remaining arguments are quoted names or `Terms`, then the type checker will automatically quote the arguments at the macro's use site. At some point, the macro is expected to unify the provided metavariable with some other term, thus solving it.

Figure 1 demonstrates a macro that quotes its argument. The first step is to quote the quoted expression argument again, using `quoteTC`, yielding a quotation of a quotation. This doubly-quoted expression is passed, using Agda's new support for Haskell-style do-notation, into a function that unifies it with the hole. Because unification removes one layer of quotation, `unify` inserts the original quoted term into the hole. The value of `sampleTerm` is

```
lam visible (abs "n" (var 0 []))
```

The constructor `lam` represents a lambda, and its body is formed by the abstraction constructor `abs` that represents a scope in which a new name `"n"` is bound. The body of the abstraction is a reference back to the abstracted name using de Bruijn index `0`.

The `unquoteTC` primitive removes one level of quotation. Figure 2 demonstrates the use of `unquoteTC`. The macro

```
macro
  mc2 : Term → Term → TC ⊤
  mc2 exp hole =
    do exp' ← unquoteTC exp
       unify hole exp'

sampleSyntax : Nat → Nat
sampleSyntax =
  mc2 (lam visible (abs "n" (var 0 [])))
```

Figure 2. A macro that unquotes its argument

`mc2` expects a quotation of a quotation, and substitutes its unquotation for the current metavariable.

The `unquoteDecl` and `unquoteDef` primitives, which run `TC` computations in a declaration context, will typically introduce new declarations by side effect. A function of a given type is declared using `declareDef`, and it can be given a definition using `defineFun`. Similarly, a postulate of a given type is defined using `declarePostulate`. Figure 3 shows an Agda implementation of addition on natural numbers, while figure 4 demonstrates an equivalent metaprogram that adds the same definition to the context.

```
plus : Nat → Nat → Nat
plus zero b = b
plus (suc n) b = suc (plus n b)
```

Figure 3. Addition on natural numbers

In Figure 4, `declareDef` declares the type of `plus`. The constructor `pi` represents dependent function types, but a pattern synonym is used to make it shorter. Similarly, `def` constructs references to defined names, and the pattern synonym `` `Nat `` abbreviates references to the defined name `Nat`, and `vArg` represents the desired visibility and relevance settings of the arguments. Once declared, `plus` is defined

```
pattern vArg x = arg (arg-info visible relevant) x
pattern _`⇒_ a b = pi (vArg a) (abs "_" b)
pattern `Nat = def (quote Nat) []

unquoteDecl plus =
  do declareDef (vArg plus) (`Nat `⇒ `Nat `⇒ `Nat)
     defineFun plus
       (clause (vArg (con (quote zero) []) ::
                vArg (var "y") ::
                [])
         (var 0 []) ::
        clause (vArg (con (quote suc)
                          (vArg (var "x") :: [])) ::
                vArg (var "y") ::
                [])
          (con (quote suc)
            (vArg (def plus
                    (vArg (var 1 []) ::
                     vArg (var 0 []) :: [])) ::
             [])) ::
        [])
```

Figure 4. Addition, defined by metaprogramming

using `defineFun`, which takes a name and a list of clauses, defining the function by side effect. Each clause consists of a pattern and a right-hand side. Patterns have their own datatype, while right-hand sides are `Terms`. The name `con` is overloaded: in patterns, it denotes a pattern that matches a particular constructor, while in `Terms`, it denotes a reference to a constructor.

2.2 Higher Inductive Types

Homotopy type theory [25] is a research program that aims to develop univalent, higher-dimensional type theories. A type theory is *univalent* when equivalences between types are considered equivalent to identifications between types; it is *higher-dimensional* when we allow non-trivial identifications that every structure in the theory must nevertheless respect. Identifications between elements of a type are considered to be at the lowest dimension, while identifications between identifications at dimension n are at dimension $n + 1$. Voevodsky added univalence to type theories as an axiom, asserting new identifications without providing a means to compute with them. While more recent work arranges the computational mechanisms of the type theory such that univalence can be derived, as is done in cubical type theories [4, 10], we are concerned with modeling concepts from homotopy type theory in existing, mature implementations of type theory, so we follow Univalent Foundations Program [25] in modeling paths using Martin-Löf's identity type. Higher-dimensional structure can arise from univalence, but it can also be introduced by defining new type formers that introduce not only introduction and elimination principles, but also new non-trivial identifications.

In homotopy type theories, one tends to think of types not as collections of distinct elements, but rather through the metaphor of topological spaces. The individual elements of the type correspond with points in the topological space, and identifications correspond to paths in this space.

While there is not yet a general schematic characterization of a broad class of higher inductive types along the lines of Dybjer's inductive families [13], it is convenient to syntactically represent the higher inductive types that we know are acceptable as if we had such a syntax. Thus, we sometimes specify a higher inductive type similarly to a traditional inductive type by providing its constructors (*i.e.* its points); we additionally specify the higher-dimensional structure by providing additional constructors for paths. For example, figure 5 describes `Circle`, which is a higher inductive type with one point constructor `base` and one non-trivial path constructor `loop`.

```
data Circle : Set where
  base : Circle
  loop : base ≡ base
```

Figure 5. A specification of a higher inductive type

Agda [19] is a functional programming language with full dependent types and dependent pattern matching. Agda's type theory has gained a number of new features over the years, among them the ability to restrict pattern matching to that subset that does not imply Streicher's Axiom K [9], which is inconsistent with univalence. The convenience of programming in Agda, combined with the ability to avoid axiom K, makes it a good laboratory for experimenting with the idioms and techniques of univalent programming while more practical implementations of univalent type theories are under development.

In Agda, we don't have built-in primitives to support the definition of higher inductive type such as `Circle`. One approach is to use Agda's rewrite rules [8] mechanism to define higher inductive types. In this approach, we define the dependent and non-dependent eliminators of a higher inductive type as paramterized modules inside which we declare the computation rules for points as rewrite rules using `{-# REWRITE , ...#-}` pragma. However, Agda's reflection library do not have interfaces to support introducing new pragmas and defining new modules. Another approach to define higher inductive types is to use Licata's method [17]. According to this method, a higher inductive type is defined using type abstraction inside a module. The module consists of a boiler-plate code segment which defines the higher inductive type using a private base type. Inside the module, the recursion and the induction principles acts on the constructors of the private base type. The abstract type is then exported allowing the reduction rules for point constructors to hold definitionally. For example, `Circle` is defined using Licata's method as follows.

Inside the module `Circle`, the type `S` is defined using a private datatype `S*`. The constructor `base` is defined using `base*` and the path `loop` is given as a postulated identification. The recursion and induction principles are defined by pattern matching on the constructor `base*` of the type `S*`, and thus compute as expected. The clients of `Circle` will not have access to the constructor `base*` of the private type `S*`, as it is not visible outside the module, which prevents them from writing functions that distinguish between multiple constructors of a higher inductive type that may be identified by additional path constructors. The client's *only* access to the constructor is through the provided elimination rules. The following code gives the non-dependent eliminator (sometimes called the *recursion rule*) `recS`.

`recS` ignores the path argument and simply computes to the appropriate answer for the point constructor. The computational behavior for the path constructor loop is postulated using reduction rule β`recS`. The operator `apPerhaps we should move the discussion of ap earlier, so that we don't need the digression` is frequently referred to as `cong`, because it expresses that propositional equality is a congruence. However, when viewed through a homotopy

```
module Circle where
  private
    data S* : Set where
      base* : S*

  S : Set
  S = S*

  base : S
  base = base*

  postulate
    loop : base ≡ base

  recS : {C : Set} →
    (cbase : C) →
    (cloop : cbase ≡ cbase) →
    S → C
  recS cbase cloop base* = cbase

  postulate
    βrecS : {C : Set} →
      (cbase : C) →
      (cloop : cbase ≡ cbase) →
      ap (recS cbase cloop) loop ≡ cloop

  indS : {C : S → Set} →
    (cbase : C base) →
    (cloop : transport C loop cbase ≡ cbase) →
    (circle : S) → C circle
  indS cbase cloop base* = cbase

  postulate
    βindS : {C : S → Set} →
      (cbase : C base) →
      (cloop : transport C loop cbase ≡ cbase) →
      apd (indS {C} cbase cloop) loop ≡ cloop
```

Figure 6. A HIT encoded using Licata's method

type theory lens, it is often called ap, as it describes the action of a function on paths. In a higher inductive type, ap should compute new paths from old ones.

```
ap : {A B : Set} {x y : A}
     (f : A → B) (p : x ≡ y) → f x ≡ f y
```

In addition to describing the constructors of the points and paths of S, figure 6 additionally demonstrates the dependent eliminator (that is, the induction rule) indS and its computational meaning. The dependent eliminator relies on another operation on identifications, called transport, that coerces an inhabitant of a family of types at a particular index into an inhabitant at another index. Outside of homotopy type theory, transport is typically called subst or replace, because it also expresses that substituting equal elements for equal elements is acceptable.

```
transport : {A : Set} {x y : A} →
  (P : A → Set) → (p : x ≡ y) → P x → P y
```

In the postulated computation rule for indS, the function apd is the dependent version of ap: it expresses the action of dependent functions on paths.

```
apd : {A : Set} {B : A → Set} {x y : A} →
      (f : (a : A) → B a) → (p : x ≡ y) →
      transport B p (f x) ≡ f y
```

The next section introduces the necessary automation features by describing the automatic generation of eliminators for a variant on Dybjer's inductive families. Section 4 then generalizes this feature to automate the production of eliminators for higher inductive types using Licata's technique. Section 5 revisits Angiuli et al.'s encoding of Darcs's patch theory [5] and demonstrates that the higher inductive types employed in that paper can be generated succinctly using our library.

3 Code Generation for Inductive Types

An inductive type D is a type that is freely generated by a finite collection of constructors. The constructors of D accept zero or more arguments, and result in an D. The constructors can also take an element of type D itself as an argument, but only *strictly positively*: any occurrences of the type constructor D in the type of an argument to a constructor of D must not be to the left of any arrows. Type constructors can have a number of *parameters*, which may not vary between the constructors, as well as *indices*, which may vary.

In Agda, constructors are given a function type. In Agda's reflection library, the constructor `data-type` of the datatype `Definition` stores the constructors of an inductive type as a list of `Names`. The type of a constructor can be retrieved by giving its `Name` as an input to the `getType` primitive. In this section, we discuss how to use the list of constructors and their types to generate code for the elimination rules of an inductive type.

3.1 Non-dependent Eliminators

In Agda, we define an inductive type using `data` keyword. A definition of an inductive datatype declares its type and specifies its constructors. While Agda supports a variety of ways to define new datatypes, we will restrict our attention to the subset that correspond closely to Dyber's inductive families. In general, the definition of an inductive datatype D with constructors $c_1 \ldots c_n$ has the following form:

$$\begin{array}{l}\textbf{data}\ D\ (a_1 : A_1) \ldots (a_n : A_n) : (i_1 : I_1) \to \ldots \to (i_m : I_m) \to \text{Set}\ \textbf{where} \\ \quad c_1 : \Delta_1 \to D\ a_1 \ldots a_n\ e_{11} \ldots e_{1m} \\ \qquad\qquad \vdots \\ \quad c_r : \Delta_n \to D\ a_1 \ldots a_n\ e_{r1} \ldots e_{rm}\end{array}$$

where the index instantiations $e_{k1} \ldots e_{km}$ are expressions in the scope induced by the telescope Δ_k. Every expression

```
data Vec (A : Set) : Nat → Set where
  []   : Vec A zero
  _::_ : {n : Nat} →
         (x : A) → (xs : Vec A n) →
         Vec A (suc n)
```

Figure 7. Length-indexed lists

in the definition must also be well-typed according to the provided declarations.

As an example, the datatype `Vec` represents lists of a known length. It is defined in figure 7. There is one parameter, namely `(A : Set)`, and one index, namely `Nat`. The second constructor, `_::_`, has a recursive instance of `Vec` as an argument.

While inductive datatypes are essentially characterized by their constructors, it must also be possible to eliminate their inhabitants, exposing the information in the constructors. This section describes an Agda metaprogram that generates a non-dependent recursion principle for an inductive type; section 3.2 generalizes this technique to fully dependent induction principles.

For `Vec`, the recursion principle says that, in order to eliminate a `Vec A n`, one must provide a result for the empty `Vec` and a means for transforming the head and tail of a non-empty `Vec` combined with the result of recursion onto a tail into the desired answer for the entire `Vec`. Concretely, the type of the recursor `recVec` is:

```
recVec : (A : Set) →
         {n : Nat} → Vec A n →
         (C : Set) →
         (base : C) →
         (step : {n : Nat} →
                 (x : A) →
                 (xs : Vec A n) → C →
                 C) →
         C
```

The recursor `recVec` maps the constructor `[]`, which takes zero arguments, to `base`. It maps `(x :: xs)` to `(step x xs (recVec xs C base step))`. Because `step` is applied to a recursive call to the recursor, it takes one more argument than the constructor `_::_`.

Based on the schematic presentation of inductive types D earlier in this section, we can define a schematic representation for their non-dependent eliminators D_{rec}.

$$
\begin{aligned}
D_{rec} : {} & (a_1 : A_1) \to \ldots \to (a_n : A_n) \to \\
& (i_1 : I_1) \to \ldots \to (i_m : I_m) \to \\
& (tgt : D\ a_1 \ldots a_n\ i_1\ \ldots\ i_n) \to \\
& (C : \mathsf{Set}) \to \\
& (f_1 : \Delta_1' \to C) \to \ldots \to (f_r : \Delta_r' \to C) \to \\
& C
\end{aligned}
$$

```
pattern _[_v]⇒_ a s b = pi (vArg a) (abs s b)
pattern _[_h]⇒_ a s b = pi (hArg a) (abs s b)

(agda-sort (lit 0) [ "A" h]⇒          -- A
 (def (quote Nat) [] [ "n" h]⇒        -- n
  (var 1 [] [ "x" v]⇒                 -- x
   (def (quote Vec)                   -- xs : Vec A n
        (vArg (var 2 []) ::
         vArg (var 1 []) :: [])
         [ "xs" v]⇒
    def (quote Vec)                  -- Vec A (suc n)
        (vArg (var 3 []) ::
         vArg (con (quote suc)
                (vArg (var 2 []) :: []))
         :: [])))))
```

Figure 8. Abstract syntax tree for the type of `_::_`

The type of f_i, which is the method for fulfilling the desired type C when eliminating the constructor c_i, is determined by the type of c_i. The telescope Δ_i' is the same as Δ_i for non-recursive constructor arguments. However, Δ_i' binds additional variables when there are recursive occurrences of D in the arguments. If Δ_i has an argument $(y : B)$, where B is not an application of D or a function returning such an application, Δ_i' binds $(y : B)$ directly. If B is an application of D, then an additional binding $(y' : C)$ is inserted following y. Finally, if B is a function type $\Psi \to D$, the additional binding is $(y' : \Psi \to C)$.

To construct the type of `recVec`, we need to build the types of `base` and `step`. These are derived from the corresponding types of `base` and `_::_`, which can be discovered using reflection primitives. Since `[]` requires no arguments, its corresponding method is `(base : C)`. The constructor `pi` of type `Term` encodes the abstract syntax tree (AST) representation of `_::_` (figure 8). We can retrieve and traverse the AST of `_::_`, and add new type information into it to build a new type representing `step`.

During the traversal of abstract syntax tree of the type of `_::_`, when the type `Vec` occurs directly as an argument, the result type `C` is added next to it. For example, in figure 8, a new function is built from the argument `(xs : Vec A n)` by modifying it to `(Vec A n) → C` (figure 9). Arguments other than `Vec` require no modifications. Therefore, `(x : A)` is copied into the new type without any changes. Finally, the codomain `Vec A (suc n)` of `_::_`'s type is replaced with `C`, resulting in an AST for the type of `step`.

In general, when automating the production of D_{rec}, all the information that is needed to produce the type signature is available in the `TC` monad by looking up D's definition. The constructor `data-type` contains the number of parameters occurring in a defined type. It also encodes the constructors of the type as a list of `Names`. Metaprograms can retrieve the index count by finding the difference between the number of parameters and the length of the constructor list. The constructors of D refer to the parameter and the index using de Bruijn indices.

```
(agda-sort (lit 0) [ "A" h]⇒              -- A
 (def (quote Nat) [] [ "n" h]⇒            -- n
  (def (quote Vec) (vArg (var 1 []) ::    -- Vec A n
     vArg (var 0 []) :: []) [ "_" v]⇒
   (agda-sort (lit 0) [ "C" v]⇒           -- C
    (var 0 [] [ "_" v]⇒                   -- base
     ((def (quote Nat) [] [ "n" h]⇒       -- step
       (var 5 [] [ "x" v]⇒                -- x
        (def (quote Vec)
          (vArg (var 6 []) ::
           vArg (var 1 []) ::
           [])
          [ "xs" v]⇒                      -- xs
          (var 4 [] [ "_" v]⇒             -- → C
           var 5 []))))                   -- C
      [ "_" v]⇒ var 2 [])))))             -- C
```

Figure 9. Abstract syntax tree of `recVec`'s type

The method `step` for the constructor `_::_` in `Vec`, refers to the parameter and the index using de Bruijn indices. During the construction of the type of `step`, the recursor generator updates the de Bruijn indices accordingly. Note that not all indices occur as arguments to a constructor: in `Vec`, the constructor `[]` instantiates the index with a constant.

Once the AST for `step`'s type has been found, it is possible to build the type of `recVec` in figure 9. To quantify over the return type `(C : Set)`, the `Term` constructor `agda-sort` refers to `Set)`.

The general schema for the computation rules corresponding to D_{rec} and constructors $c_1, \ldots, c_n$ follows:

$$
\begin{aligned}
&D_{rec}\ a_1\ \ldots\ a_n\ i_1\ \ldots\ i_m\ (c_1\ \Delta_1)\ C\ f_1 \ldots f_r = \\
&\quad \mathsf{RHS}\,(f_1, \Delta_1') \\
&\vdots \\
&D_{rec}\ a_1\ \ldots\ a_n\ i_1\ \ldots\ i_m\ (c_r\ \Delta_r)\ C\ f_1 \ldots f_r = \\
&\quad \mathsf{RHS}\,(f_r, \Delta_r')
\end{aligned}
$$

Here, $\overline{\Delta_j}$ is the sequence of variables bound in Δ_j. RHS constructs the application of the method f_j to the arguments of c_j, such that C is satisfied. It is defined by recursion on Δ_j. $\mathsf{RHS}\,(f_j, \cdot)$ is f_j, because all arguments have been accounted for. $\mathsf{RHS}\,(f_j, (y : B)\Delta_k)$ is $\mathsf{RHS}\,(f_j\ y, \Delta_k)$ when B does not mention D. $\mathsf{RHS}\,(f_j, (y : D)(y' : C)\Delta_k)$ is $\mathsf{RHS}\,(f_j\ y\ \ (D_{rec}\ \ldots y \ldots), \Delta_k)$, where the recursive use of D_{rec} is applied to the recursive constructor argument as well as the appropriate indices, and the parameters, result type, and methods remain constant. Higher-order recursive arguments are a generalization of first-order arguments. Finally,

$$\mathsf{RHS}\,(f_j, (y : \Psi \to D)(y' : \Psi \to C)\Delta_k)$$

is

$$\mathsf{RHS}\left(f_j\ y\ \left(\lambda\overline{\Psi}.D_{rec}\ \ldots\left(y\ \overline{\Psi}\right)\ldots\right), \Delta_k\right)$$

where the recursive use of D_{rec} is as before.

After declaring `recVec`'s type using `declareDef`, it is time to define its computational meaning. The computation rule

```
(clause
  (vArg (con (quote _::_)           -- _::_
         (vArg (var "x") ::         -- x
          vArg (var "xs") :: []))   -- xs
   vArg (var "C") ::                -- C
   vArg (var "base") ::             -- base
   vArg (var "step") :: [])         -- step
  (var 0
    (vArg (var 4 []) ::             -- x
     vArg (var 3 []) ::             -- xs
     vArg
      (def recVec                   -- recursion
       (vArg (var 3 [])             -- xs
        vArg (var 2 []) ::          -- C
        vArg (var 1 []) ::          -- base
        vArg (var 0 []) :: [])))))  -- step
```

Figure 10. Clause definition for the computation rule of `_::_`

```
data W (A : Set) (B : A → Set) : Set where
  sup : (a : A) → (B a → W A B) → W A B
```

Figure 11. W-Type

representing the action of function `recVec` on `[]` and `_::_` using `clause` (figure 10). The first argument to `clause` encodes variables corresponding to the above type, and it also includes the abstract representation of the constructors `[]` and `_::_` on which the pattern matching should occur. The second argument to `clause`, which is of type `Term`, refers to the variables in the first argument using de Bruijn indices, and it encodes the output of `recVec` when the pattern matches. The constructor `var` of `Pattern` is used to introduce new variables in the `clause` definition. The type `Pattern` also has another constructor `con` that represents patterns that match specific constructors. The type `Term` has similar constructors `var` and `con`, that respectively represent variable references and constructor invocations. The computation rules for `recVec` are

```
recVec []        C base step =
  base
recVec (x :: xs) C base step =
  step x xs (f xs C base step)
```

Figure 10 presents the AST for the clause that matches `_::_`. The de Bruijn index reference increments right-to-left, starting from the last argument. Definitions by pattern matching are added to the global context using the `defineFun` primitive.

Figure 11 presents an Agda encoding of Martin-Löf's well-orderings, the so-called *W-type*. `W` is interesting because its constructor exhibits a *higher-order* recursive instance of the type being defined. Following the recipe yields the recursor in figure 12, in which the recursive call occurs under a function representing arbitrary choices of tag.

The type of `step` is built by traversing the AST of `sup`'s type. The first argument to `sup`, which is a constant type `A`,

```
recW : ∀ {A B}
       (tgt : W A B) →
       (C : Set) →
       (step : (x : A) →
               (f : B x → W A B) →
               (f' : B x → C) →
               C) →
       C
recW (sup x f) C step =
  step x f (λ b → recW (f b) C step)
```

Figure 12. The non-dependent eliminator for `W`

```
generateRec : Arg Name → Name → TC ⊤
generateRec (arg i f) t =
  do cns ← getConstructors t
     lcons ← getLength cns
     cls ← getClause lcons zero t f cns
     RTy ← getType t
     funType ← getRtype t zero RTy
     declareDef (arg i f) funType
     defineFun f cls
```

Figure 13. Implementation for `generateRec`

is copied directly into `step`'s type. The second argument is a `(B x → W A B)`, which is a function whose codomain is a recursive instance of `W`. The resulting arguments must account for the recursion and are thus `(B x → W A B)` and `(B x → C)`. Finally, the codomain `W A B` of `sup` is replaced by `C`.

In the above computation rule, the third argument to `step` is a function that works for any choice of tag `b`. The arguments to `lam` are referenced using de Bruijn indices inside the lambda body. Thus, the de Bruijn indices for referring variables outside the lambda body must be updated accordingly.

Figure 13 demonstrates the implementation of `generateRec`, which constructs recursors. `generateRec` uses `getClause` and `getRtype` to build the computation and elimination rules respectively. It takes two arguments: the name of the function to be defined (represented by an element of type `Arg Name`), the quoted `Name` of the datatype. `generateRec` can be used to automate the generation of recursion rules for inductive types having the general schema given at the beginning of this section. The recursion rule generated by `generateRec` is brought into scope using `unquoteDecl` as follows.

```
unquoteDecl f = generateRec (vArg f)
                            (quote Vec)
```

3.2 Dependent Eliminators

The dependent eliminator for a datatype, also known as the *induction principle*, is used to eliminate elements of a datatype when the type resulting from the elimination mentions the very element being eliminated. The type of the

```
indW : {A : Set} → {B : A → Set} →
       (tgt : W A B) →
       (mot : W A B → Set) →
       (step : (x : A) →
               (f : B x → W A B) →
               (f' : (b : B x) → mot (f b)) →
               mot (sup x f)) →
       mot tgt
indW (sup x f) mot step =
  step x f (λ b -> indW (f b) mot step)
```

Figure 14. The induction principle for `W`

induction principle for D is:

$$
\begin{aligned}
D_{ind} : {} & (a_1 : A_1) \to \ldots \to (a_n : A_n) \to \\
& (i_1 : I_1) \to \ldots \to (i_m : I_m) \to \\
& (tgt : D\ a_1 \ldots a_n\ i_1\ \ldots\ i_m) \to \\
& (C : (i_1 : I_1) \to \ldots \to (i_m : I_m) \to \\
& \quad D\ a_1 \ldots a_n\ i_1\ \ldots\ i_n \to \\
& \quad \mathsf{Set}) \to \\
& (f_1 : \Delta_1' \to C\ e_{11} \ldots e_{1p}\ (c_1\ \overline{\Delta_1})) \to \\
& (f_r : \Delta_r' \to C\ e_{r1} \ldots e_{rp}\ (c_r\ \overline{\Delta_r})) \to \\
& C\ i_1\ \ldots\ i_n\ tgt
\end{aligned}
$$

Unlike the non-dependent recursion principle D_{rec}, the result type is now computed from the target and its indices. Because it expresses the reason that the target must be eliminated, the function C is often referred to as the *motive*. Similarly to D_{rec}, the type of each method f_i is derived from the type of the constructor c_i—the method argument telescope Δ_k' is similar, except the arguments that represent the result of recursion now apply the motive C to appropriate arguments. If Δ_i has an argument $(y : B)$, where B is not an application of D or a function returning such an application, Δ_i' still binds $(y : B)$ directly. If B is an application of D to parameters $a \ldots$ and indices $e \ldots$, then an additional binding $(y' : C\ e \ldots\ y)$ is inserted following y. Finally, if B is a function type $\Psi \to D\ a \ldots\ e \ldots$, the additional binding is $(y' : \Psi \to C\ e \ldots (y\ \overline{\Psi}))$.

The computation rules for the induction principle are the same as for the recursion principle. Following these rules, the induction principle for `W` can be seen in figure 14.

Automating the production of the dependent eliminator is an extension of the procedure for automating the production of the non-dependent eliminator.

We can construct the AST of `d` using the static type information obtained from `sup`. To construct `indW`, during the traversal of the AST of `sup`'s type, the argument `(a : A)` is copied without any changes, just as it is in the case of the non-dependent eliminator. The next argument, however, is a function that returns a `W`. An additional argument, representing the induction hypothesis, is needed. The induction hypothesis `(f' : (b : B x) → mot (f b))` takes the

```
(agda-sort (lit 0) [ "A" h]⇒                -- A
 ((var 0 [] [ "_" v]⇒                       -- B : A → Set
    agda-sort (lit 0)) [ "B" h]⇒
  (def (quote W)
       (vArg (var 1 []) ::                 -- tgt : W A B
        vArg (var 0 []) :: []) [ "tgt" v]⇒
   ((def (quote W)
      (vArg (var 2 []) ::                  -- mot : W A B → Set
       vArg (var 1 []) :: []) [ "_" v]⇒
     agda-sort (lit 0)) [ "mot" v]⇒
    ((var 3 [] [ "x" v]⇒                    -- x : A
      ((var 3
        (vArg (var 0 []) :: [])             -- f : B x → W A B
        [ "_" v]⇒
        def (quote W)
         (vArg (var 5 []) ::
          vArg (var 4 []) :: [])) [ "f" v]⇒
       ((var 4 (vArg (var 1 []) ::          -- f' : B x → mot (f x)
                [])
         [ "f'" v]⇒
         var 3
          (vArg (var 1 (vArg (var 0 []) ::
                        [])) ::
                []))
        [ "z" v]⇒
        var 3                               -- mot (sup x f)
        (vArg
         (con (quote sup)
          (vArg (var 2 []) ::
           vArg (var 1 []) :: []))
         :: [])))))
    [ "_" v]⇒ var 1 (vArg (var 2 []) ::     -- mot tgt
                     [])))))))
```

Figure 15. Abstract syntax tree for the dependent eliminator of `W`

```
generateInd : Arg Name → Name → TC ⊤
generateInd (arg i f) t =
  do cns ← getConstructors t
     lcons ← getLength cns
     cls ← getClauseDep lcons zero t f cns
     RTy ← getType t
     funType ← getRtypeInd t zero RTy
     declareDef (arg i f) funType
     defineFun f cls
```

Figure 16. Implementation for `generateInd`

same arguments as `f`, but it returns the motive instantiated at the application of `f`. The final return type is found by applying the motive to the target and its indices (figure 15).

We can construct the type of the induction principle `f` using `d`. The type `C` in the mapping `f` depends on the element of the input type `W A B`. Operationally, the induction principle computes just like the recursion principle. It is constructed using `clause` definitions following the same approach. The generation of induction principles is carried out using `generateInd`, in figure 16.

`generateInd` uses `getClauseDep` to generate the `clause` definitions representing the computation rules. The abstract representation of the type is provided by `getRtypeInd`. A version of the induction principle called `indW'` generated by `generateInd` is brought into scope by `unquoteDecl` as follows:

```
unquoteDecl indW' = generateInd (vArg indW')
                                (quote W)
```

4 Code Generation for Higher Inductive Types

In Agda, there are no built-in primitives to support the definition of higher inductive types. However, we can still define a higher inductive type with a base type using Licata's [17] method, as described in section 2.1. In this section, we discuss the automation of code generation for the boiler-plate code segments defining the higher inductive types.

4.1 Defining Higher Inductive Types

Our library defines a higher inductive type G as a top-level definition using a base type D, similar to the module `Circle` in section 2.1. The first step is to generate the private base type. After this, public definitions are created to provide access to the constructors, and the paths are postulated. Finally, the elimination principles are generated, and their actions on paths are postulated.

The first step is to copy the type former D's type signature, re-using it for the declaration of G. Next, the library copies the type of each constructor c_i to a definition g_i, except each reference to D is replaced by a reference to G.

Unlike the type former and the point constructors, which can be copied from the base type, the paths in G must be provided explicitly. The code generator takes the path constructor types as input and postulates them as identifications using the reflection primitive `declarePostulate`. The data type `ArgPath` in figure 17 represents paths as their types.

HITs are defined by invoking the `data-hit` metaprogram with a syntax reminiscent of ordinary data declarations and the putative HIT syntax from figure 5.

```
data-hit (quote D) G
  Gpoints (g₁ :: ... :: gᵣ :: [])
  Gpaths (p₁ :: ... :: pₛ :: [])
  (argPath
    ({a₁ : A₁} → ... → {aₙ : Aₙ} →
     Δ₁ →
     (cᵢ e ...) ≡ (cⱼ e ...)) ::
                  ⋮
   argPath
    ({x₁ : P₁} → ... → {xₙ : Pₙ} →
     Δₛ →
     (cᵢ e ...) ≡ (cⱼ e ...)) :: [])
```

```
data ArgPath {ℓ₁} : Set (lsuc ℓ₁) where
  argPath : Set ℓ₁ → ArgPath
```

Figure 17. Definition of ArgPath

```
data-hit : ∀{ℓ₁}
  (baseType : Name) → (indType : Name) →
  (pointHolder : Name) → (lcons : List Name) →
  (pathHolder : Name) → (lpaths : List Name) →
  (lpathTypes : (List (ArgPath {ℓ₁}))) → TC ⊤
data-hit base ind h1 lcons h2 lpaths pTy =
  do defineHindType base ind
     cns ← getConstructors base
     defineHitCons base ind cns lcons
     pTy' ← getPathTypes base ind cns lcons pTy
     defineHitPathCons lpaths pTy'
     definePointHolder h1 lcons
     definePathHolder h2 lpaths
```

Figure 18. The implementation of `data-hit`

`data-hit` defines top-level definitions `Gpoints` for point constructors and `Gpaths` for path constructors as part of the definition of the encoded higher inductive type `G`. This is because Agda does not natively support these definitions, so there is no way to retrieve the constructors of the G using the ordinary reflection mechanisms, so an additional registry is needed. This parallels the use of compile-time bindings in Racket [15] to store structure type metadata.

The elements of the `argPath` list represent the type of the path constructors $p_1 \ldots p_q$. The constructor g_i is not in scope when used in the path type passed to `argPath`, so the base types' constructors are used instead. Figure 18 contains the implementation of `data-hit`.

In the implementation of `data-hit` in figure 18, the helper `defineHindType` defines the higher inductive type as a top-level definition using the base type, and `defineHitCons` defines specifies the point constructors of the higher inductive type as references to the base type's constructors. `defineHitPathCons` builds the path constructors of the higher inductive type using the `argPath` list. For example, `Circle` from section 2.1 is defined thusly:

```
unquoteDecl S Spoints base Spaths loop =
  data-hit (quote S*) S
    Spoints (base :: []) -- point constructors
    Spaths (loop :: [])  -- path constructors
    (argPath (base* ≡ base*) :: [])
    -- base replaces base*
```

The identity type input (`base* ≡ base*`) to argPath represents the type of the path `loop`. The constructor `base` comes into scope only during the execution of `unquoteDecl`, and so cannot be used in the identity type reference in `argPath`. The references to `base*` in the type of the path `loop` are replaced by `base` once it is defined.

4.2 Non-dependent Eliminators for HITs

The recursion principle of a higher inductive type G maps the points and paths of G to points and paths in an output type C. We extend the general schema of the recursion principle given in section 3.1 by adding methods for path constructors as follows.

$$
\begin{aligned}
G_{rec} : {} & (a_1 : A_1) \to \ldots \to (a_n : A_n) \to \\
& (i_1 : I_1) \to \ldots \to (i_m : I_m) \to \\
& (tgt : G\ a_1 \ldots a_n\ i_1\ \ldots\ i_n) \to \\
& (C : \mathrm{Set}) \to \\
& (f_1 : \Delta'_1 \to C) \ldots (f_r : \Delta'_r \to C) \to \\
& (k_1 : \Delta'_1 \to (f_i \ldots) \equiv (f_j \ldots)) \to \\
& \vdots \\
& (k_q : \Delta'_q \to (f_i \ldots) \equiv (f_j \ldots)) \to \\
& C
\end{aligned}
$$

The schematic definition of G_{rec} supports only we have given only one-dimensional paths. Our metaprogram currently supports only one-dimensional paths, but we are planning to improve the tool to support higher-dimensional paths in the future.

The type of the method f_i for the point constructor g_i in G_{rec} is built the same way as for the normal inductive type D, as described in section 3.1. The code generator builds the type of k_i, method for path constructor p_i in G_{rec}, by traversing the AST of p_i. The arguments of k_i are handled the same way as for the point constructor's method f_i. During the traversal, the code generator uses the base type recursor D_{rec} to map the point constructors g_i of G in the codomain of p_i to f_i. Determining the computation rules corresponding to points g_i is similar to the computation rules corresponding to constructors c_i of the inductive type D, except that there are additional methods to handle paths. Paths compute new paths; the computation rules that govern the interaction of recursors and paths p_i are named and postulated. They identify the action of the recursor on the path with the corresponding method.

$$
\begin{aligned}
\beta G_{rec} : {} & (a_1 : A_1) \to \ldots \to (a_n : A_n) \to \\
& (C : \mathrm{Set}) \to \\
& (f_1 : \Delta'_1 \to C) \ldots (f_r : \Delta'_r \to C) \to \\
& (k_1 : \Delta'_1 \to (f_i \ldots) \equiv (f_j \ldots)) \to \\
& \vdots \\
& (k_q : \Delta'_q \to (f_i \ldots) \equiv (f_j \ldots)) \to \\
& ap\ (\lambda\ x.G_{rec}\ x\ C\ f_1\ \ldots\ f_r\ k_1\ \ldots\ k_q) \\
& \quad (p_i \ldots) \equiv (k_i \ldots)
\end{aligned}
$$

As an example, if code for the circle HIT from section 2.1 has been generated, and the type is called `S`, then the recursor

```
(def (quote S) [] [ "_" v]⇒     -- S
 (agda-sort (lit 0) [ "C" v]⇒  -- C : Set
  (var 0 [] [ "cbase" v]⇒      -- cbase
   (def (quote _≡_)           -- cloop
    (vArg (var 0 []) ::
     vArg (var 0 []) :: [])
    [ "cloop" v]⇒ var 2 []))))
```

Figure 19. Abstract syntax tree for recS's type

needs a method for base and one for loop. The method for base should be an inhabitant of C. If it is called cbase, then the method for loop should be a path cbase ≡ cbase. The types of the path methods depend on the values of the point methods. The code generator builds the type of loop's method by traversing the AST of loop's type, replacing references to point constructors with the result of applying the base type's recursor to the point methods.

The recursion rule recS, corresponding to figure 19, follows this pattern.

```
recS : S →
  (C : Set) →
  (cbase : C) →
  (cloop : cbase ≡ cbase) →
  C
```

The code generator builds the computation rule for the point constructor base using the same approach as described in section 3.1, as if it were for the base type. Additionally, it includes variables in the clause definition for the path constructor loop. The code generator postulates the following computation rule βrecS for the path constructor loop:

```
βrecS : (C : Set) → (cbase : C) →
  (cloop : cbase ≡ cbase) →
  ap (λ x → recS x C cbase cloop) loop
  ≡ cloop
```

The application of function recS to the path loop substitutes the point base for the argument x, and it evaluates to the path cloop in the output type C. The definition of generateRecHit can be found in figure 20.

generateRecHit takes the base types' recursion rule as input and uses that to eliminate the points during the construction of the path methods in the recursor G_{rec}. The second argument argD is a list of terms representing the computation rules for the path constructors. The generateβRecHit interface takes argD as input and builds the computation rules for the path constructors. Other inputs to generateRecHit are the point and path collections declared during the definition of the HIT.

```
generateRecHit : Arg Name → List (Arg Name) →
  (baseType : Name) → (baseRec : Name) →
  (indType : Name) → (points : List Name) →
  (paths : List Name) → TC ⊤
generateRecHit (arg i f) argD b br i pts paths =
  do lcons ← getConstructors b
     lpts ← getLength pts
     lpaths ← getLength paths
     clauses ← getPathClause lpts lpaths br
     RTy ← getType b
     fTy ← getRtypePath b i br paths zero RTy
     declareDef (arg i f) fTy
     defineFun f clauses
     generateβRecHit argD b br i f pts paths
```

Figure 20. Implementation for generateRecHit

4.3 Dependent Eliminators for HITs

The dependent eliminator for a higher inductive type G is a dependent function that maps an element g of G to an output type $C\,g$. The general schema for the induction principle of G is given as follows.

$$
\begin{aligned}
G_{ind} : {} & (a_1 : A_1) \to \ldots \to (a_n : A_n) \to \\
& (i_1 : I_1) \to \ldots \to (i_m : I_m) \to \\
& (tgt : G\ a_1 \ldots a_n\ i_1\ \ldots\ i_n) \to \\
& (C : (i_1 : I_1) \to \ldots \to (i_m : I_m) \to \\
& \quad G\ a_1 \ldots a_n\ i_1\ \ldots\ i_n \to \\
& \quad \mathrm{Set}) \to \\
& (f_1 : \Delta_1' \to C\ j_{11} \ldots j_{1p}\ (c_1\ \overline{\Delta_1})) \to \\
& \vdots \\
& (f_r : \Delta_r' \to C\ j_{r1} \ldots j_{rp}\ (c_r\ \overline{\Delta_r})) \to \\
& (k_1 : \Delta_1' \to \mathtt{transport}\ C\, p_1\ (f_i \ldots) \equiv (f_j \ldots)) \to \\
& \vdots \\
& (k_q : \Delta_q' \to \mathtt{transport}\ C\, p_q\ (f_i \ldots) \equiv (f_j \ldots)) \to \\
& C\ i_1\ \ldots\ i_n\ tgt
\end{aligned}
$$

Similar to G_{rec}, the type of f_i is built the same way as for the normal inductive type D. The code generator builds the type of the method for path constructor p_i, called k_i, in G_{ind}, by traversing the AST of p_i. During the traversal, the code generator uses the base eliminator D_{ind} to map the point constructors g_i of G in the codomain of p_i to f_i. In the first argument to the identity type in the codomain of k_i, the code generator adds an application of transport to the motive C and the path p_i. The arguments of k_i are handled the same way as for f_i. The computation rules corresponding to paths

p_i are postulated as follows:

$$
\begin{aligned}
&\beta G_i : (a_1 : A_1) \to \ldots \to (a_n : A_n) \to \\
&\quad (C : (i_1 : I_1) \to \ldots \to (i_m : I_m) \to \\
&\qquad G\ a_1 \ldots a_n\ i_1\ \ldots\ i_n \to \\
&\qquad \mathsf{Set}) \to \\
&\quad (f_1 : \Delta_1' \to C\ j_{11} \ldots j_{1p}\ (c_1\ \overline{\Delta_1})) \to \\
&\quad (f_r : \Delta_r' \to C\ j_{r1} \ldots j_{rp}\ (c_r\ \overline{\Delta_r})) \to \\
&\quad (k_1 : \Delta_1' \to \mathsf{transport}\ C\ p_1\ (f_i \ldots) \equiv (f_j \ldots)) \to \\
&\quad (k_r : \Delta_r' \to \mathsf{transport}\ C\ p_r\ (f_i \ldots) \equiv (f_j \ldots)) \to \\
&\quad \mathsf{apd}\ (\lambda\ x\ .\ G_{ind}\ x\ C\ f_1\ \ldots\ f_r\ k_1\ \ldots\ k_r) \\
&\qquad (p_i \ldots) \equiv \\
&\quad (k_i \ldots)
\end{aligned}
$$

For the type `S` with point constructor `base` and path constructor `loop`, to define a mapping `indS : (x : S) → C x`, we need `cbase : C base` and `cloop : transport C loop cbase ≡ cbase`, where `cloop` is a heterogeneous path transported over `loop`. The code generator builds the type of `cloop` by adding relevant type information to the type of `loop`. The type of the method for the path constructor `cloop` is derived by inserting a call to `transport` with arguments `C`, `loop` and `cbase`. The code generator applies the eliminator of base type `S*` to map `base` to `cbase` during the construction of the codomain of `cloop`. The following declaration gives the type of `indS`.

```
indS : (circle : S) →
  (C : S → Set) →
  (cbase : C base) →
  (cloop : transport C loop cbase ≡ cbase) →
  C circle
```

The computation rule for `base`, which defines the action of `indS` on `base`, is built using the same approach as for the non-dependent eliminator `recS`. The postulated computation rule `βindS` for the path `loop` uses `apd` which gives the action of dependent function `indS` on the path `loop`.

```
βindS : (C : S → Set) →
  (cbase : C base) →
  (cloop : transport C loop cbase ≡ cbase) →
  apd (λ x → indS x C cbase cloop) loop ≡ cloop
```

Figure 21 gives the implementation of `generateIndHit`. The helper `generateβIndHit` builds the computation rules for the path constructors.

```
generateIndHit : Arg Name → List (Arg Name) →
  (baseType : Name) → (baseElm : Name) →
  (indType : Name) → (points : List Name) →
  (paths : List Name) → TC ⊤
generateIndHit (arg i f) argD b br i pts paths =
  do lcons ← getConstructors b
     lp1 ← getLength pts
     lp2 ← getLength paths
     cls ← getPathClauseDep lp1 lp2 b br lcons
     RTy ← getType b
     fTy ← getRtypePathDep b i br pts paths
             zero RTy
     declareDef (arg i f) fTy
     defineFun f cls
     generateβIndHit argD b br i f pts paths
```

Figure 21. Implementation for `generateIndHit`

5 Applications

Homotopy type theory has thus far primarily been applied to the encoding of mathematics, rather than to programming. Nevertheless, there are a few applications of homotopy type theory to programming. Applications such as homotopical patch theory [5] discuss a model of the core of the of Darcs [23] version control system using patch theory [12, 18] encoded as a HIT. Containers in homotopy type theory [2, 3] implement data structures such as multisets and cycles. Automating the HIT boilerplate code allows more programmers to begin experimenting with programming with HITs.

5.1 Patch Theory Revisited

A patch is a syntactic representation of a function that modifies a repository when applied. For example, a patch $(s_1 \leftrightarrow s_2 \mathbin{@} l)$, which replaces string s_1 with s_2 at line l, when applied to a repository containing with string s_1 at line l results in a repository with string s_2 at line l. In homotopical patch theory [5], the patches are modeled as paths in a higher inductive type. Because patches are paths, they automatically satisfy groupoid laws, such as that the composition of patches is associative, that all patches have an inverse, and that inverse patches compose to the identity. Domain-specific laws related to the patches, such as that two swaps at independent lines commute, are implemented as higher dimensional paths. The computational content of the patches is extracted by mapping them to bijections in the universe with the help of univalence. Due to the functoriality of mappings in type theory, the functions preserve the path structures and thus the desired axioms.

We reimplemented Angiuli et al.'s patch theory in Agda. We implemented basic patches such as the insertion of a string as line l_1 in a file and deletion of a line l_2 from a file. The functions implementing insertion and deletion in the universe are not bijective. So, to map the paths representing the patches insert and delete into the universe, we used Angiuli et al.'s patch history approach. According to this approach, we developed a separate higher inductive type `History` which serves as the types of patches. In addition to basic patches, we also implemented patches involving

encryption or decryption with cryptosystems like RSA [22] and Paillier [20].

Having implemented patch theory, we then reimplemented `History` and the encrypted repository type `cryptR` it using our code generator. We also automated the code generation for the elimination and the computation rules for the higher inductive types `History` and `cryptR`. In addition to easing the implementation difficulties of higher inductive types, the code generator greatly reduced the code size. The type definitions shrank from around 1500 to around 70 lines, resulting in a 60% decrease in the overall number of lines of code in the development.

5.2 Cryptographic Protocols

Vivekanandan [28] models certain cryptographic protocols using homotopy type theory, introducing a new approach to formally specificying cryptographic schemes using types. The work discusses modeling cryptDB [21] using a framework similar to Angiuli et al.'s patch theory. CryptDB employs layered encryption techniques and homomorphic encryption. We can implement cryptDB by modeling the database queries as paths in a higher inductive type and mapping the paths to the universe using singleton types [5]. The code generator can be applied to generate code for the higher inductive type representing cryptDB and its corresponding elimination and computation rules. By using the code generator, we can decrease the length and increase the readability of the definitions, hopefully making it more accessible to the broad cryptographic community.

6 Related Work

Kokke and Swierstra [16] implemented a library for proof search using Agda's old reflection primitives, from before it had elaborator reflection. They describe a Prolog interpreter in the style of Stutterheim et al. [24]. It employs a hint database and a customizable depth-first traversal, with lemmas to assist in the proof search.

van der Walt and Swierstra [27] and van der Walt [26] discuss automating specific categories of proofs using proof by reflection. A key component of this proof technique is a means for converting an expression into a quoted representation. They automate this process, giving a user-defined datatype. van der Walt gives an overview of Agda's old metaprogramming tools.

Ongoing work on cubical type theories [4, 10, 11] provides a computational interpretation of univalence and HITs. We strenuously hope that these systems quickly reach maturity, rendering our code generator obsolete. In the meantime, however, these systems are not yet as mature as Agda.

7 Conclusion and Future Work

We presented a code generator that generates the encodings of higher inductive types, developed using Agda's new support for Idris-style elaborator reflection. In particular, the tool generates formation, introduction, path, and elimination rules for 1-dimensional higher inductive types. This syntax is greatly simplified with respect to writing the encoding by hand. We demonstrated an extensive reduction in code size by employing our tool. Next, we intend to extend the tool to support higher-dimensional paths in the definition of HITs, bringing its benefits to a wider class of problems.

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