

CSE-5331-001-DBMS MODELS AND IMPLEMENTATION

Project 1 Phase 1

Rigorous 2PL with Wound-Wait protocol

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Data Structure:

The data structure used for rigorous 2PL with wound-wait protocol is as follows:

- Hash Map for Transaction Table and Lock Table
- Queue for List of blocked transactions operations

Additional Information in the lock table:

A separate column to keep track on the number of read operation in the lock table is added.

Information stored in transaction Table:

- transaction_id
- transaction_timestamp
- transaction_state
- list_of_items_locked

Information stored in Lock Table:

- item_name
- lock_state
- transaction_holding_lock
- transaction_id_waiting_for_lock
- lock_status
- no_of reads

Pseudo Code for rigorous 2PL with wound-wait protocol:

This program simulates the behavior of rigorous 2 phase locking

Create a hash map for Transaction table

Create a hash map for Lock table

declare timestamp_counter to 0

declare transaction_id, data_item

read the input file

get the operation, transaction id and data item from the file

if operation is not equal to "b"

 check if transaction_id is already blocked or aborted in the transaction table

 if transaction_state is equal to "blocked":

 add the operation to the list of blocked operations

 else if transaction_state is equal to "aborted":

 Ignore the operation and fetch the next operation and proceed

Iterate until all the operations are read from the file

if operation is equal to "b"

 increment timestamp_counter

 set transaction_id

 goto: begin_transaction(transaction_id, timestamp_counter)

else if operation is equal "r"

 set transaction_id and data_item;

 goto: read_lock(transaction_id, data_item)

else if operation is equal to "w"

 set transaction_id and data_item;

 goto: write_lock(transaction_id, data_item)

else if operation is equal to "e"

 set transaction_id;

 goto: end_transaction(transaction_id)

Fetch next operation and proceed

// Pseudo code for every new transaction begin

begin_transaction(transaction_id, timestamp_counter):

```
set transaction_id in the transaction table
set the counter value to the transaction_timestamp in the transaction table
set transaction_state to active in the transaction table
set null in the list_of_items_locked
print "A record is created for the transaction_id in the transaction table"
```

```
// Pseudo code for read lock(data_item)
```

```
read_lock(transaction_id, data_item):
```

```
    traverse the lock table for the requested data_item
    check if the data_item is already locked by any other transaction in lock table
    if the data_item is not in the table
        set data_item to item_name in the lock table
        set read_lock to lock_state in the lock table
        set the transaction_id to the transaction_holding_lock in the lock table
        set lock_status to locked
        set data_item to the list_of_items_locked in the transaction table
        increment no_of_reads by 1 in the lock table for that data_item
        print "transaction_id acquires read_lock on data_item"
```

```
else if lock_state is "write_locked" and transaction_holding_lock is equal to
    transaction_id
        downgrade the lock_state to read_locked in lock table
        print "transaction_id downgrades from write_lock to read_lock on data_item"
```

```
else if lock_state is "read_locked"
    add the transaction_id to the transaction_holding_lock
    add the data_item to the list_of_items_locked for the current transaction
    increment no_of_reads by 1 in the lock table for that data_item
    print "transaction_id acquires read_lock on data_item"
```

```
else if lock_state is "write_locked and transaction_id is not equal to
    transactionID_holding_lock
        // prevent deadlock using wound-wait protocol
        if timestamp of requesting transaction is less than transaction holding the lock
            abort the transaction holding the lock and release all the locks held by
            transaction holding the lock
            set lock_status to unlock to all the items held by transaction_holding_lock
            change the transaction_state of the transaction_holding_lock to abort in
            transaction table
            print "transaction holding the lock is aborted"
```

```
else
```

```
    add transaction_id of transaction_requesting_write_lock to
    transactionID_waiting_for_lock in lock table
    change the transaction_state of the transaction requesting write lock to
    waiting in transaction table
    print "requesting transaction is blocked"
```

```
// Pseudo code for write_lock(data_item)
```

```
write_lock(transaction_id,data_item)
```

```
    traverse the lock table for the requested data_item
    check if data_item is already write_locked by any other transaction in lock table
    if the data_item is not in the table
        set data_item to item_name
        set write_locked to lock_state
        set the transaction_id to the transaction_holding_write_lock
        set lock_status to lock
        print "transaction_id acquires write_lock on data_item"
```

```
    else if transactionID_holding_lock is equal to transactionID_requesting_write_lock and
        lock_state is equal to "read_locked" and no_of_reads < 2
        upgrade the lock_state to write_lock in lock table
        print "transaction_id upgraded from read_lock to write_lock on data_item"
```

```
    else if lock_state is "write_locked" or "read_locked" and transaction_id is not equal to
    transactionID_holding_lock
```

```
        // prevent deadlock using wound-wait protocol
```

```
        if timestamp of requesting transaction is less than transaction holding the lock
            abort the transaction holding the lock and release all the locks held by
            transaction holding the lock
            update the no_of_reads in the lock table
            set lock_status to unlock to all the items held by transaction
            change the transaction_state of the transaction_holding_write_lock to
            abort in transaction table
```

```
        print "transaction holding the lock is aborted"
```

```
    else
```

```
        add transaction_id of transaction_requesting_write_lock to
        transaction_id_waiting_for_lock in lock table
        change the transaction_state of the transaction_requesting_write_lock to
        waiting in transaction table
        print "requesting transaction is blocked"
```

```

// Pseudo code for end transaction
end_transaction(transaction_id):
    check the transaction state of the current transaction to determine whether to commit or
    abort the transaction
    if transaction_state is equal to "active"
        goto: commit
    else if transaction_state is equal to "blocked"
        goto:abort_transaction(transaction_id)
    else if transaction_state is equal to "aborted"
        ignore the end operation

// Pseudo code for abort transaction
abort_transaction(transaction_id):
    if transaction_id is equal to transactionID_holding_lock
        goto: unlock_transaction(transaction_id)
        change the status of the transaction_state to aborted in the transaction table
        print "transaction_id aborted"

// Pseudo code for commit
commit(transaction_id):
    change the status of the transaction_state to committed in the transaction table
    goto:unlock_transaction(transaction_id)
    print "transaction_id committed"

// Pseudo code for unlock transaction
unlock_transaction(transaction_id):
    release all locks held by the transaction_id
    change the lock_status to unlock in the lock table
    update the no_of_reads in the lock table if it is read lock
    traverse the list of blocked operation
    if there are any waiting operation in the list of blocked operations, process list one by one
        iterate each item in the list of blocked operations
        if operation is equal to "r"
            goto:read_lock(transaction_id, data_item)
        else if operation is equal to "w"
            goto:write_lock(transaction_id, data_item)
    else
        return

```

