

Functional Programming (FPR)

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Introduction

Getting started

This document is an essay for Functional Programming course at Software Engineering Programme. I have been given two tasks, solve weighting puzzle, and explain haskell. To separate those two tasks and to improve clarity of this essay I have introduced two sections `Puzzle` and `Haskell`.

`Puzzle` section in an narrative around a problem and setps that leads to solution. `Haskell` section contains explanation to syntax and functional programming knowledge.

Haskell

In order to use some functions, I imported two modules. `Import` keyword imports modules, and makes its content aviable for us.

```
import Data.List import Data.Ord import Data.Function
```

State and Test Algebraic datatypes

State

Puzzle

Intro

Haskell

First we shall define two algebraic datatypes. `State` and `Test`. `State` datatype has two contructors `Pair` and `Triple`. It should be noted here that everything in Haskell is a function therefore constructor is also a function.

`Pair` constructor is a type of `Int -> Int -> State`. It means that we can looks at this two ways. It takes two arguments and returns `State`, or it takes one argument and returns a function that takes one argument and returns `State`. Which is called partial application.

```
data State = Pair Int Int | Triple Int Int Int
deriving (Eq, Show)
```

Test

`Test` data type has two constructors, `TPair` and `TTrip`. `TPair` constructor takes two tuples. Tuple represents cartisian product. As you can see there are two diffrent ways of

representing a arguments. As cartesian product or as a series of functions in the case of `State` constructors.

In haskell we have ability to transpose cartisian product to series of functions by using currying, and vice versa.

```
data Test = TPair (Int, Int) (Int, Int) | TTrip (Int,Int,Int) (Int,Int,Int)
    deriving (Eq, Show)
```

Another piece that needs an explanation is keyword `deriving`. Keyword deriving allows us to make a instance of type classes (`Eq` and `Show`).

Otherwise we would need to make manually an instance of desired type class by denoting

```
instance Eq Test where
```

By using pattern matching this guards that `TPair` test will be only conducted in a `Pair` state, and a `TTrip` test in a `Triple` state. In addition to that there are predicates that checks validity of test against state. if the number of coins is the same in each pan of the scale if there is sufficiently many coins in the variousse piles.

```
valid :: State -> Test -> Bool
valid (Pair u g) (TPair (a, b) (c ,d)) =
    (a+b) == (c+d) &&
    (a+c) <= u &&
    (a+b+c+d) <= (u+g)
valid (Triple l h g) (TTrip (a, b, c) (d, e, f)) =
    (a+b+c) == (d+e+f) &&
    (a+d) <= l &&
    (b+e) <= h &&
    (c+f) <= g
```

Choosing and conducting a test

OUTCOMES TBD

```
outcomes :: State -> Test -> [State]
outcomes (Pair u g) (TPair (a, b) (c, d)) == True =
  | valid (Pair u g) (TPair (a, b) (c, d)) == True =
    [Pair un gc] ++
    [Triple l h gcc] ++
    [Triple l h gcc]
  | otherwise = error ("Invalid state or test" ++ (show (Pair u g)))
  where
    un = (u - (a + c))
    gcc = (u - (a + c)) + g
    gc = g + a + c
    l = a
    h = c
outcomes (Triple l h g) (TTrip (a, b, c) (d, e, f))
  | valid (Triple l h g) (TTrip (a, b, c) (d, e, f)) == True =
    [Triple (a+d) (b+e) (g+(1-(a+d))+(h-(b+e)))] ++ -- 1 1 10
    [Triple (a+d) (b+e) (g+(1-(a+d))+(h-(b+e)))] ++ -- 1 1 10
    [Triple (1-(a+d)) (h-(b+e)) (g+a+b+d+e)] -- 2 2 8
  | otherwise = error ("Invalid state or test:" ++ " "
    ++ (show (Triple l h g)) ++ " "
    ++ (show (TTrip (a, b, c) (d, e, f))))
```

Weighings

Weighings function has different implementation for each of State.

```
weighings :: State -> [Test]
weighings (Pair u g) = [TPair (a,b) (a+b, 0) | a<-[0..u], b<-[0..g],
  (a+b) > 0,
  ((2*a)+b) <= u,
  b <= g]
weighings (Triple l h g) = [TTrip (a, b, c) (d, e, f) | k1<-[1..k],
  (a, b, c) <- choices k1 (l, h, g),
  (d, e, f) <- choices k1 (l, h, g),
  c == 0 || f == 0, (a,b,c) <= (d,e,f), (c+f) <= g, (b+e) <= h, (a+d) <= l, (a+b+
  where
    k = (l+h+g) `div` 2
```

Choices function uses set comprehension with predicates

```
choices :: Int -> (Int, Int, Int) -> [(Int, Int, Int)]
choices k (l, h, g) = [(i,j,k-i-j) | i<-[0..l], j<-[0..h],
  (k-i-j) <= g,
  (k-i-j) >= 0]
```

This is a case of manually set up an instance of type class.

```
instance Ord State where
  (Pair _ _) < (Triple _ _ _) = False
  (Pair _ g1) < (Pair _ g2) = g2 < g1
  (Triple _ _ g1) < (Triple _ _ g2) = g2 < g1

  (Pair _ _) <= (Triple _ _ _) = False
  (Pair _ g1) <= (Pair _ g2) = g2 <= g1
  (Triple _ _ g1) <= (Triple _ _ g2) = g2 <= g1
```

Test function uses two other functions. `weighings`, `productive` and `filter`. First two are defined by us. Filter is part of lib. Which filters out an collection by provided predicate

```
productive :: State -> Test -> Bool
productive s t = all (s > ) (outcomes s t)

tests :: State -> [Test]
tests s = filter (productive s) (weighings s)
```

Decision tree

Now we can introduce `Tree` data type that represents weighting process. It's a ternary tree. This data type has two constructors, `Stop` that represents final state as a leaf of the tree. 'Node that represents weighting as a node of the tree.

```
data Tree = Stop State | Node Test [Tree]
  deriving (Show)
```

Final is a predicate that determine whether State is final.

```
final :: State -> Bool
final (Pair u g)
  | u == 0 = True
  | otherwise = False
final (Triple l h g)
  | l == 1 && h == 0 = True
  | l == 0 && h == 1 = True
  | otherwise = False
```

In height function, case for `(Stop s)` is not really interesting. But `(Node _ xs)` recursively calculates height of a tree and then selects maximum value.

```
height :: Tree -> Int
height (Stop s) = 0
height (Node _ xs) = 1 + maximum (map height xs)
```

`minHeight` is a partial function that throws an error in case of empty collection. Then it calculates height on elements and returns a tuple of `Tree` and its height. Then collection is sorted and selected first element.

```
minHeight :: [Tree] -> Tree
minHeight [] = error "Tree cannot be empty"
minHeight xs = snd
  $ head
  $ sortBy (compare `on` (\(y,_) -> y))
  $ map (\x -> (height x, x)) xs
```

`mktree` function generate tree recursively for each outcome `s t`.

```
mktree :: State -> Tree
mktree s
  | (final s) == True = Stop s
  | otherwise = minHeight
    $ map subTree
    $ productiveTests
  where
    productiveTests = tests s
    subTree = (\t -> (Node t (map mktree (getOutcomes s t))))
    getOutcomes = (\s t -> (outcomes s t))
```

Caching heights

Program in the previous sections works but it is rather slow. One of the problems is that it is recomputing the heights of trees. Now we will try to introduce height on each Node so that we can compute it in constant time.

We define new datatype `TreeH`. As you can see it's very similar to previous `Tree` with one change. We introduce notion of height in the `NodeH` constructor.

```
data TreeH = StopH State | NodeH Int Test [TreeH]
  deriving Show
```

Function that extracts height out of `TreeH` type.

```
heightH :: TreeH -> Int
heightH (StopH s) = 0
heightH (NodeH h t ts) = h
```

`treeH2tree` that will map `TreeH` type to `Tree`. In this case `TreeH` loses direct access to its height.

```
treeH2tree :: TreeH -> Tree
treeH2tree (StopH s) = (Stop s)
treeH2tree (NodeH h t ts) = (Node t (map treeH2tree ts))
```

`nodeH` is a function that for given `Test` and list of tries, it will construct new `TreeH`.

```
nodeH :: Test -> [TreeH] -> TreeH
nodeH t ths = NodeH ((+) 1 $ maximum $ map heightH ths) t ths
```

`tree2treeH` is just an inverse of `treeH2tree`.

```
tree2treeH :: Tree -> TreeH
tree2treeH (Stop s) = (StopH s)
tree2treeH (Node t ts) = nodeH t (map tree2treeH ts)
```

As you can notice `heightH . tree2treeH = height`. This equality holds due to function composition law. Let's do quick type check.

`heightH` has a type `TreeH -> Int`, `tree2treeH` has a type `Tree -> TreeH`. Composition of those two functions has type `Tree -> TreeH -> TreeH -> Int` it can be simplified to `Tree -> Int` (composition law), which is a type signature of our `height` function.

We could also prove it by induction.

Finally we can implement function that constructs tree for given state.

```
mktreeH :: State -> TreeH
mktreeH s
  | (final s) == True = (StopH s)
  | otherwise = head $ sortBy (compare `on` heightH) $ subTree $ productiveTests
    where
      productiveTests = tests s
      subTree = map (\t -> nodeH t (map mktreeH (getOutcomes s t)))
      getOutcomes = (\s t -> (outcomes s t))
```

As it was stated in an assignment. This approach does not massively improve performance.

Greedy solution

This function was copied from assignment description.

```
optimal :: State -> Test -> Bool
optimal (Pair u g) (TPair (a,b) (ab,0)) =
    (2 * a + b <= p) && (u - 2 * a - b <= q)
    where
        p = 3 ^ (t - 1)
        q = (p - 1) `div` 2
        t = ceiling (logBase 3 (fromIntegral (2 * u + k)))
        k = if g == 0 then 2 else 1
optimal (Triple l h g) (TTrip (a,b,c) (d,e,f)) =
    (a+e) `max` (b+d) `max` (l-a-d+h-b-e) <= p
    where
        p = 3 ^ (t - 1)
        t = ceiling (logBase 3 (fromIntegral (l+h)))
```

Function that filters unoptimal weighings, leaving just optimals.

```
bestTests :: State -> [Test]
bestTests s = filter (optimal s) (weighings s)
```

Function that builds tree out of the first optimal test.

```
mktreeG :: State -> TreeH
mktreeG s
    | (final s) == True = (StopH s)
    | otherwise = subTree $ bestTest
    where
        bestTest = head $ bestTests s
        subTree = (\t -> nodeH t (map mktreeG (getOutcomes s t)))
        getOutcomes = (\s t -> (outcomes s t))
```

Function that builds tries based on all optimal tests

```
mktreesG :: State -> [TreeH]
mktreesG s
    | (final s) == True = [StopH s]
    | otherwise = concat $ makeTree
    where
        optimalTests = bestTests s
        makeTree = map (\t -> map mktreeG (outcomes s t)) $ optimalTests
```

Conclusion