

Programming Project 1: Boolean Expression Evaluator

You will be writing a C++ program that reads a string, determines whether it is a syntactically correct Boolean expression, and, if so, determines its value. For a grade of 100, This program will be due on Thursday, Sept 28. Each week or part of a week thereafter counts off 10%. The program must pass all of the test cases created by the grader to receive credit. Programs that do not pass all test cases may be resubmitted the following week.

1. Data model

A **symbol** is one of the following nine strings: "T", "F", "^", "v", "~", "=>", "<=>", "(", ")". Intuitively, these are interpreted as *true*, *false*, *and*, *or*, *not*, *implies*, *if-and-only-if*, and left and right parentheses.

A **symbol string** is the concatenation of zero or more symbols and/or spaces. For example, the following are symbol strings:

- "T "
- "=> T v ~ F ((F"
- "T => F ^ (F v F) "
- ""
- " "

and the following are not symbols strings

- "T X "
- " F= > T"
- "TF p * ^) "

A **Boolean expression** is a C++ vector of strings satisfying certain conditions. We will write C++ vectors by writing their elements, separated in commas and enclosed in brackets; for example, we will write [10,20] for the vector of length 2 whose first element is 10 and whose second element is 20.. "Followed by" means *concatenated with*. For example, ["F"] followed by ["^", "T"] is the vector ["F","^","T"]. The rules for forming Boolean expressions of various sorts are given below.

1. A **Boolean constant** is ["T"] or ["F"].
2. An **unbreakable expression** is either a Boolean constant, or ["("] followed by a Boolean expression followed by [")"].
3. A **negation** is either an unbreakable expression, or ["~"] followed by a negation.
4. A **conjunction** is either a negation, or a conjunction followed by ["^"] followed by a negation.
5. A **disjunction** is either a conjunction, or a disjunction followed by ["v"] followed by a conjunction.

6. An **implication** is either a disjunction, or a disjunction followed by ["=>"] followed by an implication.
7. A **Boolean expression** is either an implication, or an implication followed by ["<=>"] followed by a Boolean expression.

This grammar can be formalized in [BNF notation](#) as follows:

```

Const  →   "T" |  "F"
U  →   Const |  "(" B ")" // note, this rule has been corrected
N  →   U |  "~" N
C  →   N |  C "^" N
D  →   C |  D "v" C      // note, this rule has been corrected
I  →   D |  D "=>" I
B  →   I |  I "<=>" B

```

An *AST* (short for *abstract syntax tree*, the standard name for what the book calls an *expression tree*) is defined below. This follows the pattern for defining trees given in the Aho & Ullman book in Chapter 5, p. 232. We will use AST's as a data structure to store the semantic structure of Boolean expressions.

```

typedef struct AST* pNODE;
struct AST {string info; pNODE children[2];};

```

The *info* member of an AST is a symbol, as defined above. If *info* is "T" or "F" then both children are NULL. If *info* is "~" then *children*[1] is NULL. Otherwise, both children are non-NULL. Sample code illustrating the use of this data structure can be found [here](#).

A *tokRslt* is a struct with two fields:

- *success*, a bool
- *syms*, a C++ vector of strings.

A *parseRslt* is a struct with two fields:

- *success*, a bool
- *ast*, an AST

A *TPERslt* is a struct with two fields:

- *val*, a bool
- *msg*, a string

2. Functions

Implement the following five functions in a single C++ file:

`tokRslt tokenize(string s)`

1. If s is a string, $tokenize(s).success$ is *true* if s is a string of symbols, and *false* otherwise.
2. If s is a string of symbols, then $tokenize(s).syms$ is a vector of the symbols occurring in s , in order. For example, if $s = "T \vee v \Rightarrow"$ then $tokenize(s).syms = ["T", "v", "v", "\Rightarrow"]$

`parseRslt parse(vector<string> V)`

1. If V is a Boolean expression, then $parse(V).success$ is *true* and $parse(V).ast$ is the abstract syntax tree of V according to the [standard grammar of Boolean expressions](#).
2. Otherwise, $parse(V).success$ is *false*.

`bool eval (AST T)`

1. $eval(T)$ is the value of T according to the [standard semantics of Boolean expressions](#).

`TPERslt TPE(string s) (tokenize, parse, and evaluate)`

1. If s is a string of symbols whose tokenization is a Boolean expression, then $TPE.msg$ is "success" and $TPE(s).val$ is the value of that Boolean expression.
2. If s is a string of symbols whose tokenization is not a Boolean expression, then $TPE.msg$ is "grammar error".
3. If s is not a string of symbols, then $TPE.msg$ is "symbol error".

`string TPEOut(string s)`

1. If s is a string of symbols whose tokenization is a Boolean expression, then $TPEOut(s)$ is the value of that expression, converted to a string, which is either "true" or "false".
2. If s is a string of symbols whose tokenization is not a Boolean expression, then $TPEOut(s)$ is "grammar error".
3. If s is not a string of symbols, then $TPEOut(s)$ is "symbol error".

For example,

- If $s = "T \vee F \wedge T"$, then $TPEOut(s)$ is "true"
- If $s = "T \Rightarrow (F \times T)"$, then $TPEOut(s)$ is "symbol error"
- If $s = "T \ T \ (F \ \& \ T \Rightarrow F) "$, then $TPEOut(s)$ is "grammar error"

Academic Integrity

Students are allowed to discuss this assignment verbally with each other, search the Web for useful source code, and download and use code if it helps you (though I doubt one can find code that will use the same data structures we use). The following are prohibited:

1. Looking at another student's code
2. Sharing code with another student
3. Asking another person to write code for you or to see their code.

Violations of 1-3 will result in a request to withdraw from the class with a *W*, if done before the last day to drop. After this date, it will result in an *F*. If those options are not acceptable, the result will be a university-level academic honesty proceeding, with the aim of expulsion from the university.