ASSIGNMENT: HW3

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1. Considering the potential function as mentioned below,

$$\phi_i = \begin{cases} 2 \cdot num_i - sire_i & \text{if } d_i \ge 1/2 \\ \text{sire}_i - 2 \cdot num_i & \text{if } d_i \le 1/2 \end{cases}$$

we will calculate the amortized costs for insentions and deletion costs for call possible combinations.

Amortized costs for insentions:

(9) For all the blobdow cases m= num; and sis=sire;

(a) di-1 = 1: Here m=5

(b) \frac{1}{2} \le d_{i-1} < 1:

$$\frac{C_{i}}{1} \left| \frac{\phi_{i}}{2(m+1)-2s} \right| \frac{\phi_{i-1}}{2m-s} \left| \frac{C_{i}}{1+(2(m+1)-s)-(2m-s)} \right|$$

$$= 1+2m+2-s+-2m+s$$

$$= 2$$

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The amortized cast for insertion operation in all cases is a constant and so it is O(1).

Now, we consider the deletion operation.

$$\frac{C_{i}}{1} \left| \frac{\phi_{i}}{2(m-1)-S} \right| \frac{\phi_{i-1}}{2m-S} \left| \frac{C_{i}}{=1+(2(m-1)-S)-(2m-S)} \right|$$

$$= 1+2m-2-S-2m+S$$

$$= -\frac{1}{2}$$

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(6)
$$\alpha_{i-1} = 1/2$$
: Here 2m = s

$$\frac{C_{i}}{1} \left| \frac{\phi_{i-1}}{s-2(m-1)} \right| \frac{C_{i}}{s-2m} \left| \frac{C_{i}}{z_{1}+(s-2(m-1))-(s-2m)} \right| \\
= 1 + (s-2m+2-s+2m) \\
= 2 + (s-2m+2-s+2m) \\
= 2 + (s-2m+2-s+2m) \\
= 3 + (s-2m$$

(d)
$$d_{s-1} = \frac{1}{4}$$
: $m = \frac{5}{4}$ and $d_i = \frac{1}{2}$

$$\frac{\epsilon_{i}}{m} = \frac{\beta_{i}}{\frac{5}{2} - 2(m-1)} = \frac{\beta_{i-1}}{5-2m} = \frac{\zeta_{i}}{m+\frac{5}{2}-2m+2-5+2m} = \frac{2+m-\frac{5}{2}}{\frac{2-\frac{5}{4}}{4}}$$

The deletion operation is not of O(1)

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- 2 (a) The amobired cost of ORDER is O(1) because, the Operation will either return TRUE if label(1) < label(1), otherwise it returns false. There is no change in the label number of entires (a) the values of the labels and so it takes only a constant time. Hence the amortized cost is O(1).
- 2 (b) The amortized cost of DELETE is O(1) because, the operation will only reduced delete on element and so the overall potential decreases only by the label of the deleted element and that will be some contant number only Hence the amortized cost is O(1):

2.

2-(0)

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Considers the scenario, where in during the INSERT operation rebelling is not required. So the actual cost will be just only for labelling the new inserted element. Now; if we look at change in potential, lets say we have inserted element y after x-

Similarly , &i- will be

$$\frac{9}{5-1} = \underbrace{\xi - clog\left(\frac{9}{m}\right)}_{0 \le k \in \mathbb{N}-1}$$

$$= \underbrace{\xi - clog\left(\frac{9}{m}\right) - clog\left(\frac{9}{m}\right) - clog\left(\frac{9}{m}\right) - clog\left(\frac{9}{m}\right) - \underbrace{\xi - clog\left(\frac{9}{m}\right)}_{242 \le k \le n-1}$$

The change in potential would be at element 'y' and premous successor of 'x', which is

$$\psi_{j} - \psi_{x-1} = -c \left[log(\frac{g_{y}}{n}) + log(\frac{g_{y+1}}{n}) - log(\frac{g_{x+1}}{n}) - log(\frac{g_{x+1}}{n}) \right]$$

$$= c \left[log M - log g_{y} + log M - log g_{y+1} + log g_{x+1} + log g_{x+1} + log g_{x+1} + log g_{x+1} \right]$$

$$= c \left[log M - log g_{y} - log g_{y+1} + log g_{x+1} \right]$$

gs, gg+1 and gg+1 are the differences in the valences of label and they will be can be treated of contints and so, the order is amortized lat in this cyc is O(log 17)