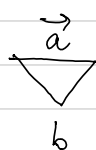


A multicategory \mathcal{C} is

- a set of objects
- for each object list \vec{a} and object b ,
a set of morphisms

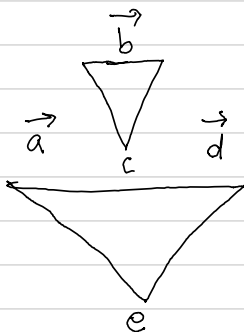


written $\mathcal{C}(\vec{a}; b)$

- for each object a , an identity



- Composition



(upper & lower in our pictures)

satisfying four equations: left & right unitality

associativity

interchange

$$\begin{array}{c} \vec{b} \\ \triangle \\ \vec{a} \end{array} = \begin{array}{c} \vec{b} \\ \triangle \\ \vec{a} \end{array}$$

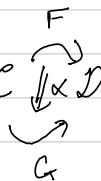
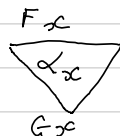
$$\begin{array}{c} \vec{a} \\ \triangle \\ \vec{a} \end{array} = \begin{array}{c} \vec{a} \\ \triangle \\ \vec{a} \end{array}$$

Multicategories form a 2-category (ahem, size)

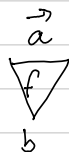
Multifunctor $F: \mathcal{C} \rightarrow \mathcal{D}$ is obvious

Given $F, G: \mathcal{C} \rightarrow \mathcal{D}$,

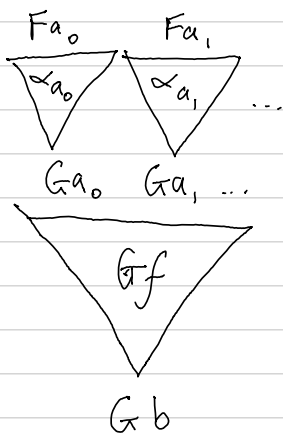
sends each $x \in \mathcal{C}$ to unary



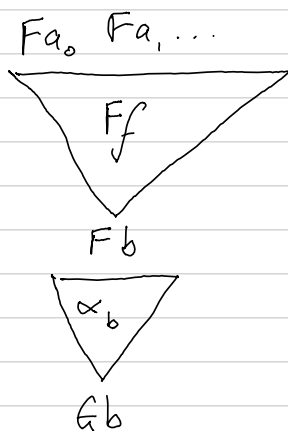
s.t. for any \mathcal{C} -morphism



we have



\cong



Left-skew multicategory \mathcal{C} (Bonnet & Lack) is

- a set of objects

no post by will afford

- for each object list \vec{a} and object b ,

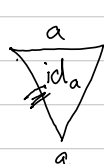
a set of loose morphisms $\triangleleft_{\vec{a}}^b$ written $\mathcal{C}(\vec{a}, b)$

- for each object a , object list \vec{b} and object c

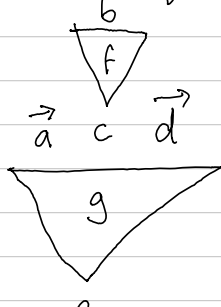
a set of tight morphisms $\triangleleft_a^{\vec{b}}^c$ written $\mathcal{C}(a, \vec{b}; c)$

- a loosening operation $\mathcal{C}(a, \vec{b}; c) \rightarrow \mathcal{C}(\vec{a}, b)$
 $f \mapsto \hat{f}$

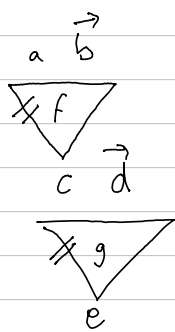
- for each object a , a tight identity



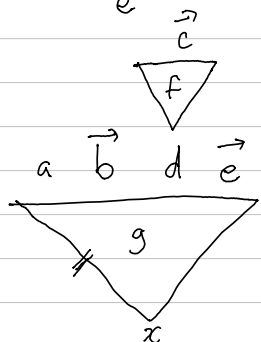
- loose composition



- tight composition

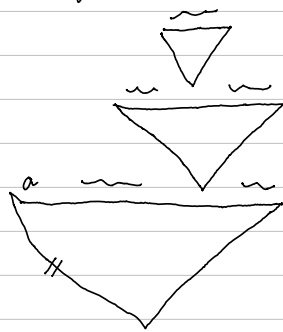


- mixed composition

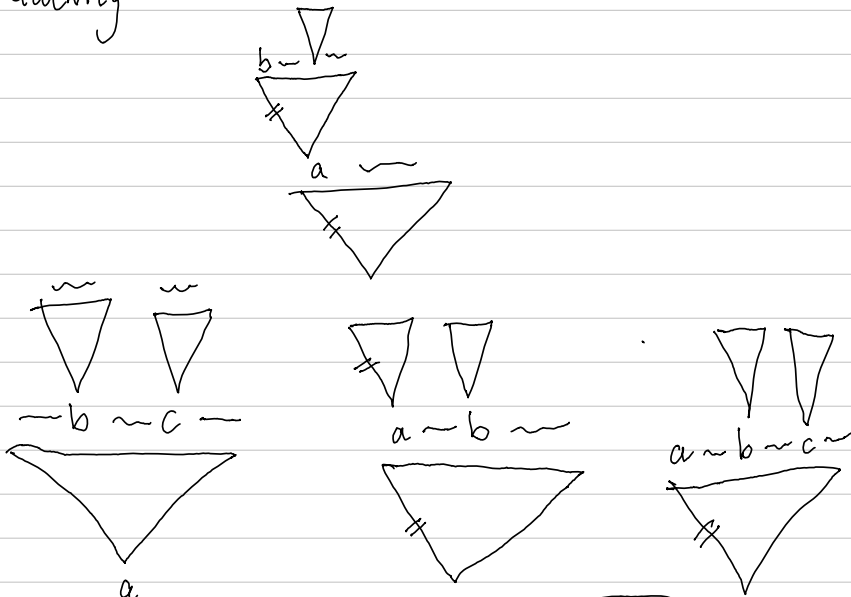


satisfying

- mixed composition commutes with loosening
- tight composition commutes with loosening
- left unitality for each kind of composition
- right unitality for loose & tight composition
- associativity for loose
- associativity for tight
- associativity



- associativity

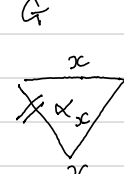


14 equations

These form a 2-category.

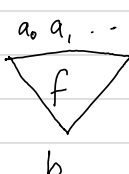
Given $F, G: \mathcal{C} \rightarrow \mathcal{D}$, a natural transformation $\alpha: F \rightarrow G$

sends each \mathcal{C} -object x to a tight unary morphism

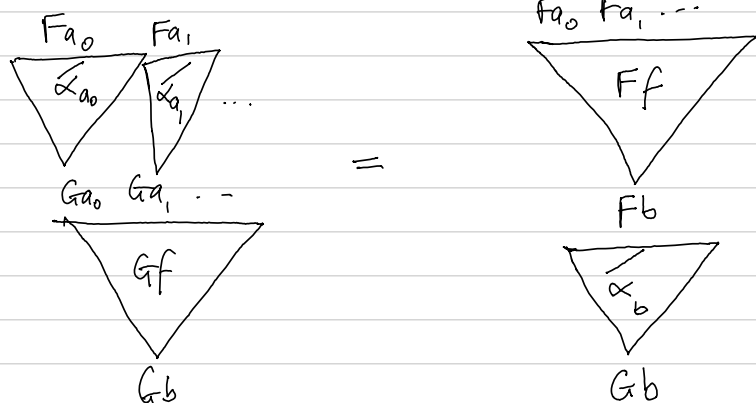


such that

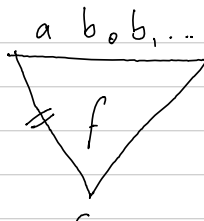
- for each loose morphism



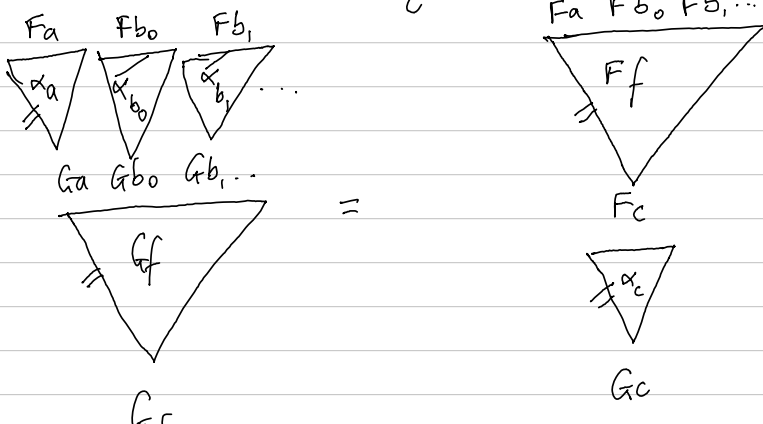
we have



- for each tight morphism



we have

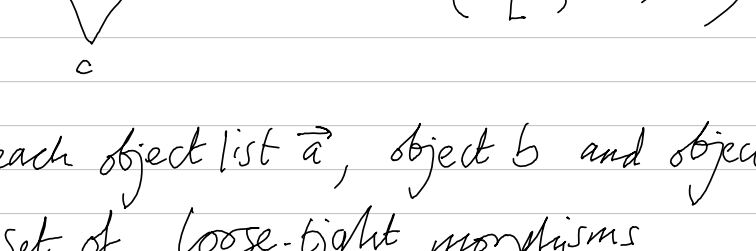


A *bi*kew multicategory \mathcal{C} consists of

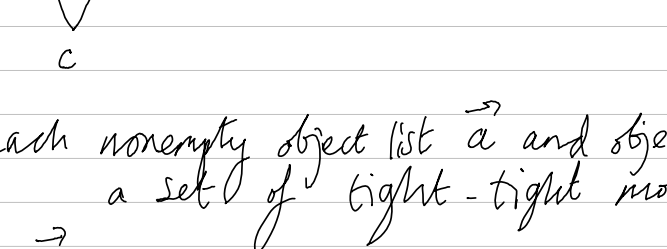
- a set of objects \mathcal{C}
- for each object list \vec{a} and object b , a set of loose-loose morphisms



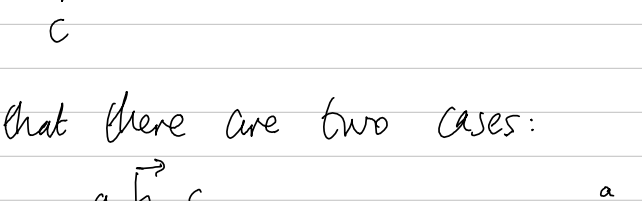
- for each object a , object list \vec{b} and object c , a set of tight-loose morphisms



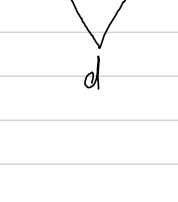
- for each object list \vec{a} , object b and object c , a set of loose-tight morphisms



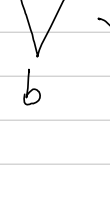
- for each nonempty object list \vec{a} and object b , a set of tight-tight morphisms



Note that there are two cases:



and



omitted in the CT 2019 slides

- left-loosening operations

$$\mathcal{C}([\vec{a}]; b) \rightarrow \mathcal{C}(\vec{a}; b) \quad f \mapsto \hat{f}$$

$$\mathcal{C}([\vec{a}]; b) \rightarrow \mathcal{C}(\vec{a}; b)$$

- right-loosening operations

$$\mathcal{C}(\vec{a}; b) \rightarrow \mathcal{C}(\vec{a}; b) \quad f \mapsto \hat{f}$$

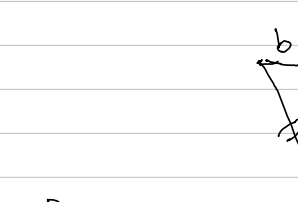
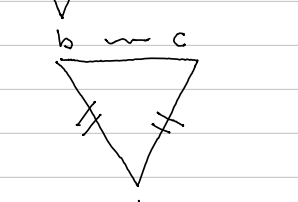
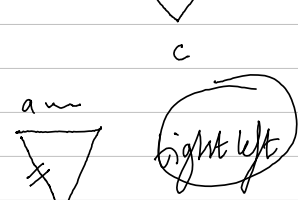
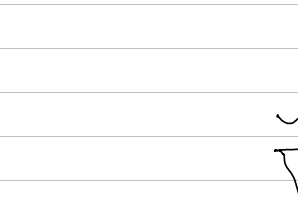
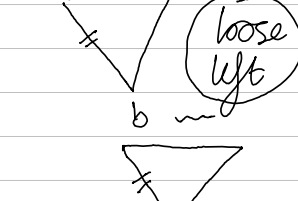
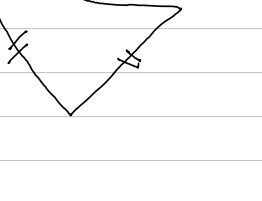
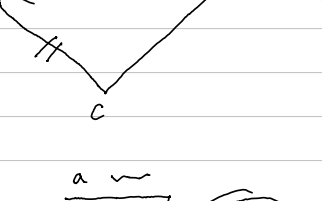
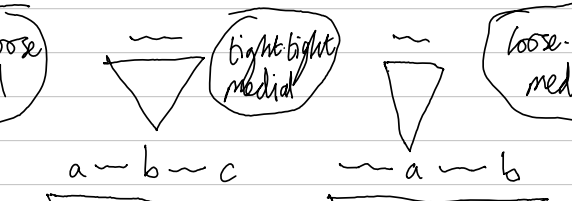
$$\mathcal{C}(\vec{a}; b) \rightarrow \mathcal{C}(\vec{a}; b)$$

(\vec{a} is nonempty in the above)

- for each object a , a tight-tight identity



- nine kinds of composition



satisfying the following equations.

- left-loosening & right-loosening commute

$$\mathcal{C}([\vec{a}]; b) \xrightarrow{\text{written } f \mapsto \hat{f}} \mathcal{C}(\vec{a}; b) \xrightarrow{\vec{a} \text{ nonempty}} \mathcal{C}(\vec{a}; b)$$

- each kind of composition resulting in a tight-loose or tight-tight morphism

- viz. {
- tight-loose medial
 - tight-tight medial
 - loose left
 - tight left
 - bi-tight

commutes with left-loosening

- each kind of composition resulting in a loose-tight or tight-tight morphism

(5 kinds)

commutes with right-loosening

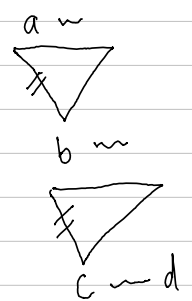
- each kind of composition is left-unital

- each kind of morphism is right-unital

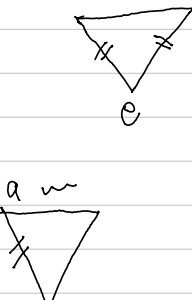
- sixteen associativity laws:

nine consisting of medial composition above and any kind of composition below

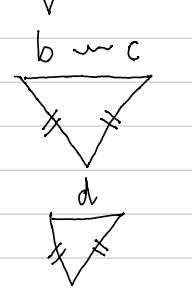
seven more:



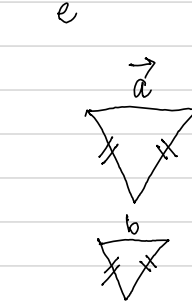
and mirror image



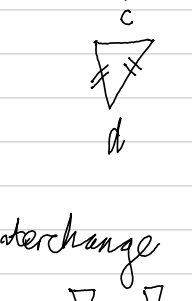
and mirror image



and mirror image



and mirror image



and mirror image

nine interchange laws

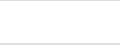
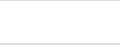
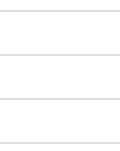
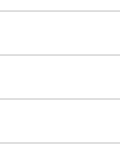
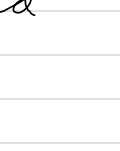
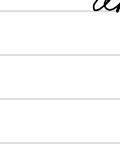
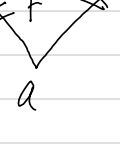
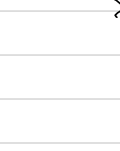
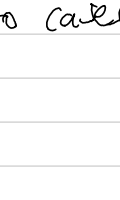
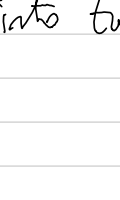
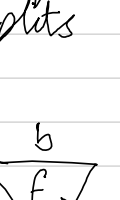
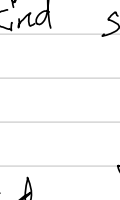
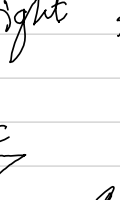
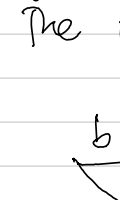
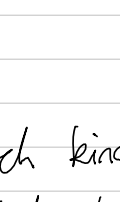
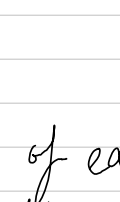
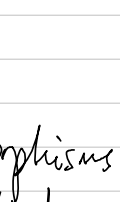
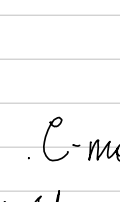
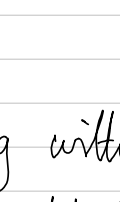
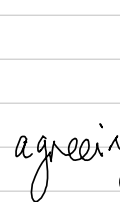
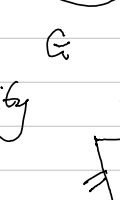
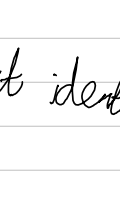
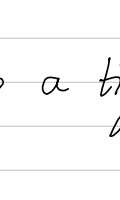
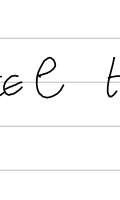
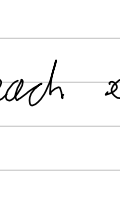
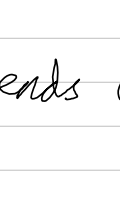
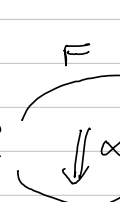
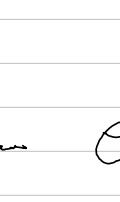
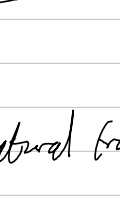
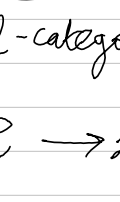
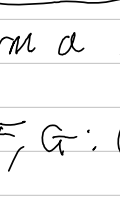
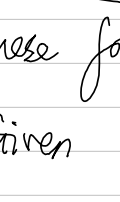
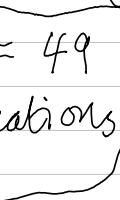
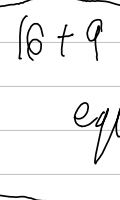
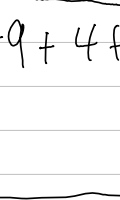
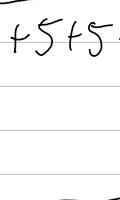
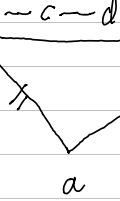
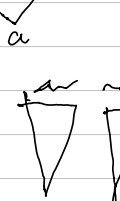
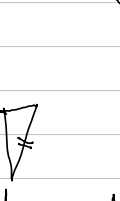
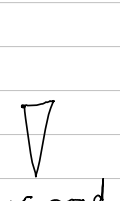
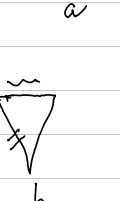
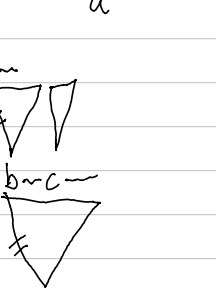


Diagram of 2-categories

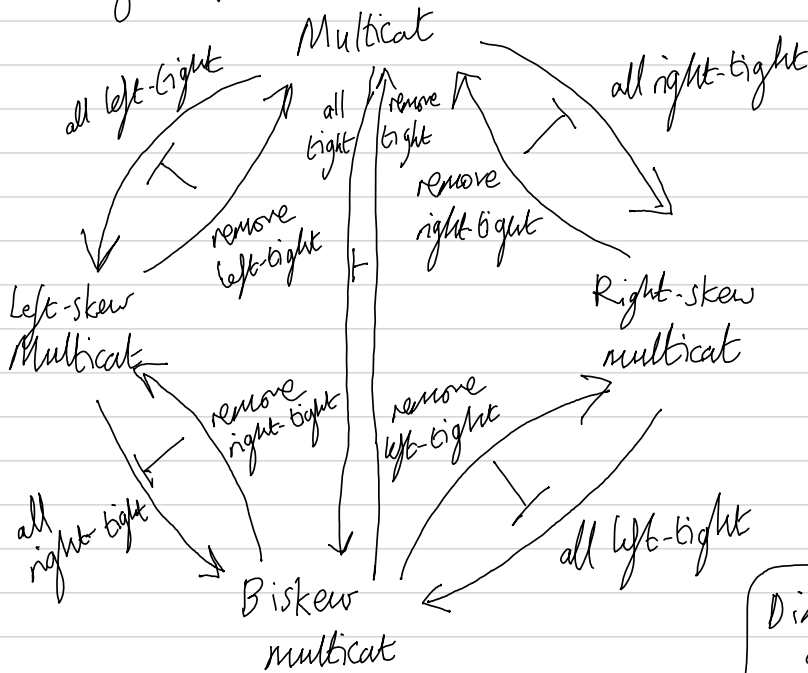


Diagram commutes

Units are identity

- Given a multicategory \mathcal{C}

$$\text{All left-right}(\mathcal{C})(\vec{a}; b) \stackrel{\text{def}}{=} \mathcal{C}(\vec{a}; b)$$

$$\text{All left-right}(\mathcal{C})([a, \vec{b}; c) \stackrel{\text{def}}{=} \mathcal{C}(a, \vec{b}; c)$$

Loosening is identity.
Identity on a is id_a

- Given a left skew-multicategory \mathcal{D}

$$\text{Remove left-right}(\mathcal{D})(\vec{a}; b) \stackrel{\text{def}}{=} \mathcal{D}(\vec{a}; b)$$

Identity on a is id_a

- Given a left-skew multicategory \mathcal{D} and category \mathcal{C} we have

$$\text{Multicat}(\text{Remove left-right}(\mathcal{C}), \mathcal{D}) \cong \text{Left-skew multicat}(\mathcal{C}, \text{All left-right}(\mathcal{D}))$$

In a biskew multicategory \mathcal{C} ,

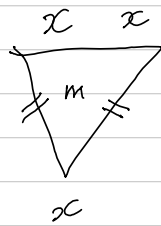
a monoid consists of

- an object x

- a loose-loose morphism



- a tight-tight morphism



satisfying equations:

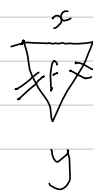
$$- \quad \begin{array}{c} \text{triangle } u \\ \downarrow x \quad x \\ \text{triangle } m \end{array} = \begin{array}{c} x \\ \downarrow \text{id} \\ \text{triangle } m \end{array}$$

$$- \quad \begin{array}{c} \text{triangle } u \\ \downarrow x \quad x \\ \text{triangle } m \end{array} = \begin{array}{c} x \\ \downarrow \text{id} \\ \text{triangle } m \end{array}$$

$$\cdot \quad \begin{array}{c} x \quad x \\ \downarrow m \\ \text{triangle } m \end{array} = \begin{array}{c} x \quad x \\ \downarrow m \\ \text{triangle } m \end{array}$$

A homomorphism $(x, u, m) \longrightarrow (y, v, n)$

is a tight-tight morphism



such that

$$\begin{array}{c} \text{triangle } u \\ \downarrow x \\ \text{triangle } \hat{f} \end{array} = \begin{array}{c} \text{triangle } v \\ \downarrow y \end{array}$$

$$\begin{array}{c} x \quad x \\ \downarrow m \\ \text{triangle } m \end{array} = \begin{array}{c} x \quad x \\ \downarrow \hat{f} \quad \hat{f} \\ \text{triangle } n \end{array}$$

Hence we obtain a notion of monoid & homomorphism
in a multicategory
left-skew multi-
right-skew multi-

Let \mathcal{C} be a multicategory

A left \mathcal{C} -multicategory \mathcal{D} consists of

- a set of objects

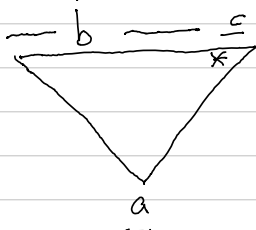
- for each \mathcal{C} -object list \vec{a} and \mathcal{D} -objects \underline{b} and \underline{c} ,
a set of morphisms



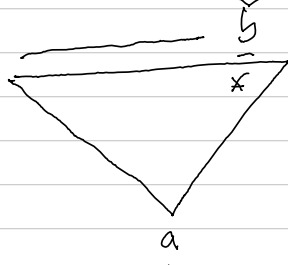
- for each \mathcal{D} -object \underline{a} , an identity morphism



- two kinds of composition



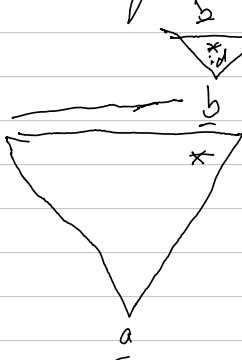
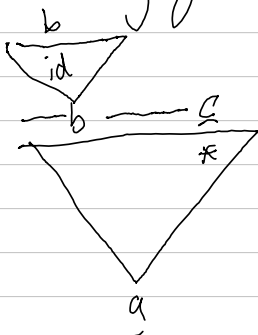
outer



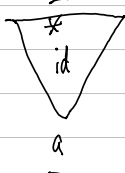
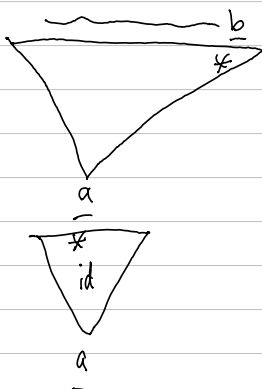
inner

satisfying the following equations:

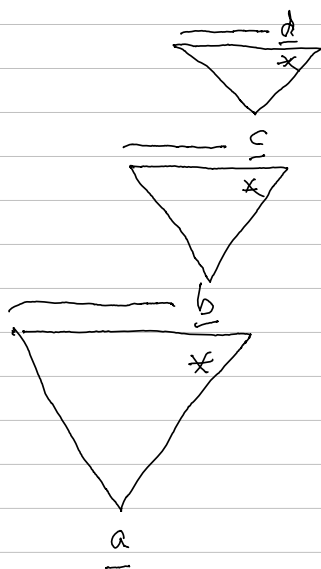
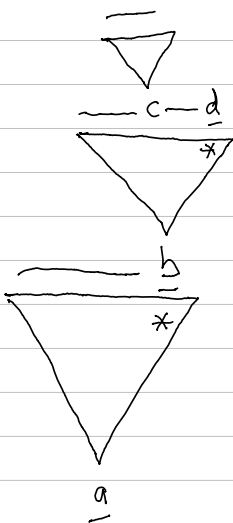
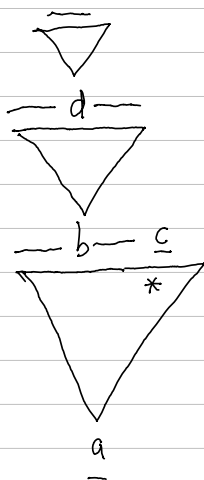
left-identity for each kind of composition



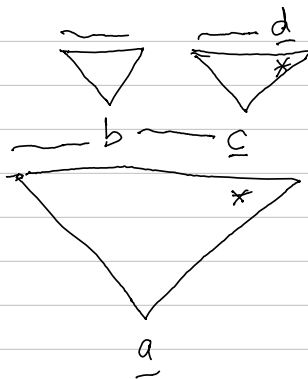
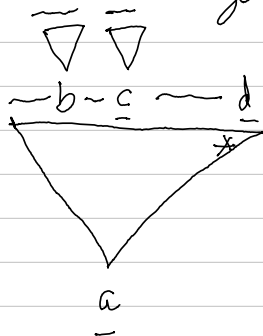
right identity



three associativity laws



two interchange laws

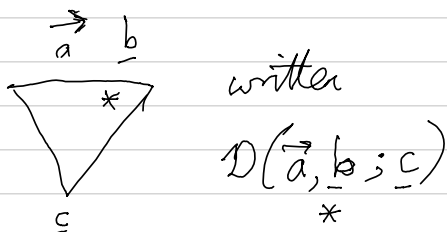


8 equations

Let \mathcal{C} be a left-skew multicategory.

A left-skew \mathcal{C} -multicategory \mathcal{D} is

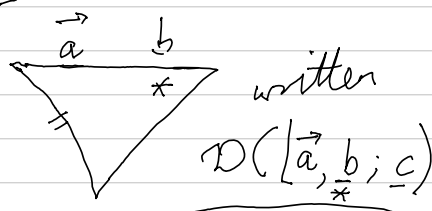
- a set of objects
- for each \mathcal{C} -object list \vec{a} and \mathcal{D} -objects $\underline{b}, \underline{c}$ a set of loose morphisms



written

$$D(\vec{a}, \underline{b}; \underline{c})$$

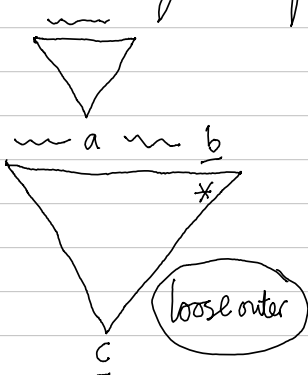
and a set of tight morphisms



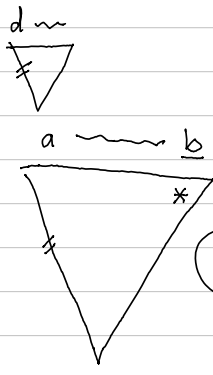
written

$$D([\vec{a}, \underline{b}]; \underline{c})$$

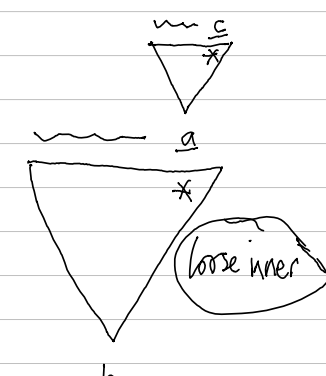
- a loosening operation $D([\vec{a}, \underline{b}]; \underline{c}) \rightarrow D(\vec{a}, \underline{b}; \underline{c})$ written $f \mapsto f^{\sim}$
- four kinds of composition



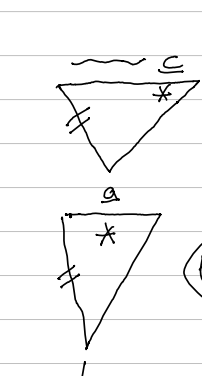
loose outer



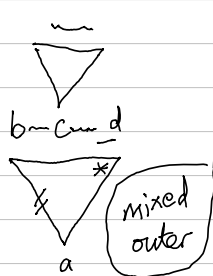
tight outer



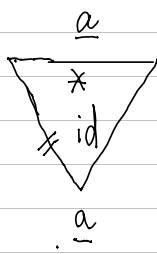
loose inner



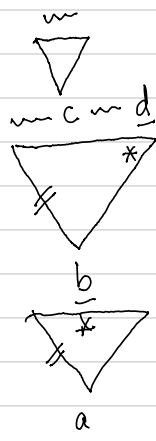
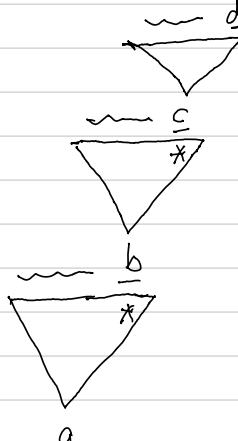
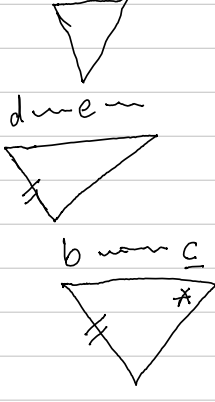
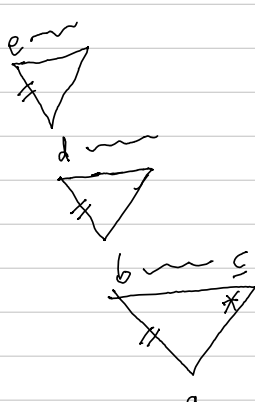
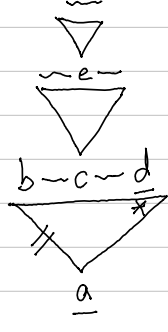
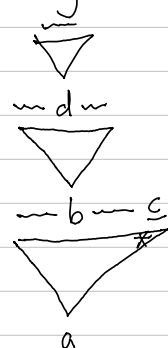
tight inner



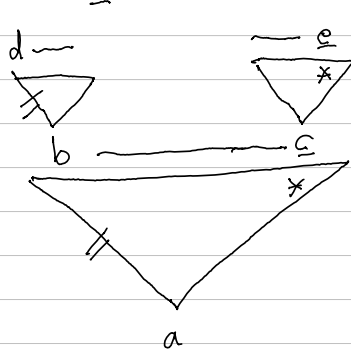
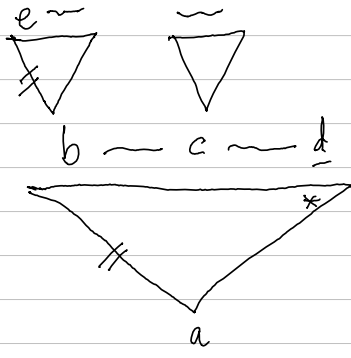
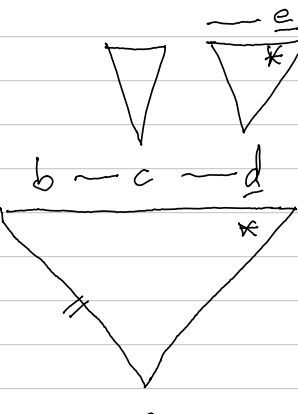
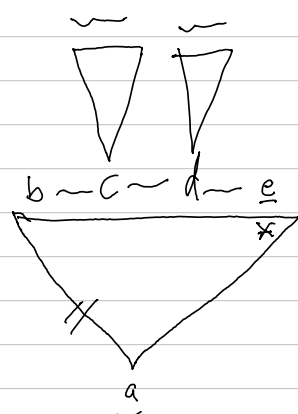
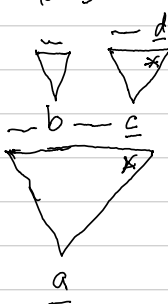
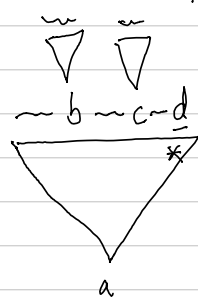
mixed outer



- for each \mathcal{D} -object \underline{a} , a tight identity satisfying the following equations:
- tight outer composition commutes with loosening
- tight inner composition commutes with loosening
- mixed outer composition commutes with loosening
- left unitality for each kind of composition
- right unitality for each kind of morphism (with inner composition)
- nine associativity laws



six interchange laws



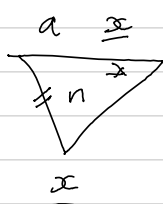
25 equations

Given a left-skew multicategory \mathcal{C}
 and a left-skew \mathcal{C} -multicategory \mathcal{D}
 and a monoid $M = (a, u, m)$ in \mathcal{C}

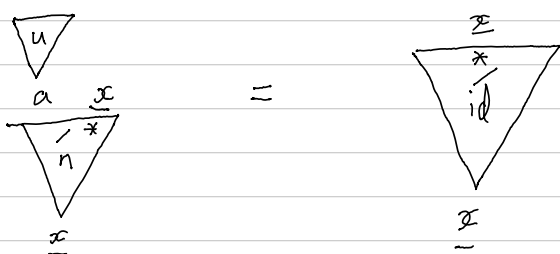
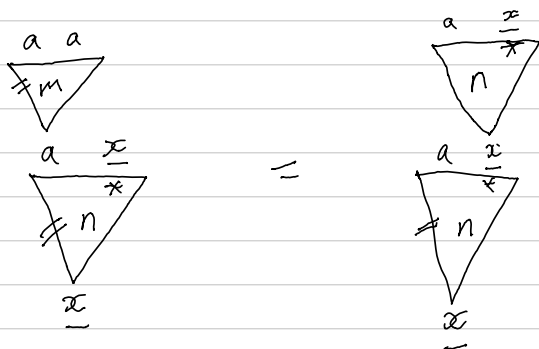
a left M -set is

- a \mathcal{D} -object \underline{x}

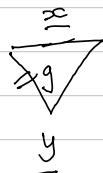
- a tight morphism



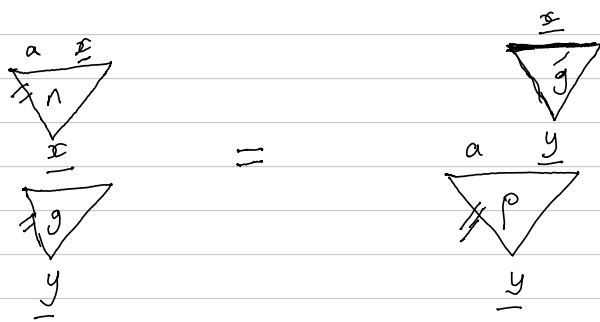
satisfying the following equations



A homomorphism $(\underline{x}, n) \rightarrow (y, p)$ is
 a tight morphism



satisfying the following equation



A monoid homomorphism $M \xrightarrow{f} M'$
 $(a, m, u) \rightarrow (a', m', u')$

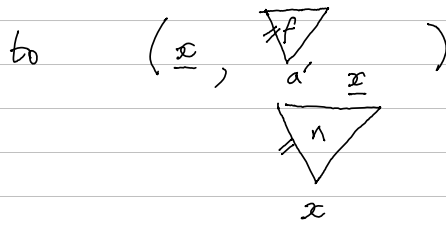
gives a faithful functor

Left M' -set

\longrightarrow

Left M -set

It sends (\underline{x}, n) to

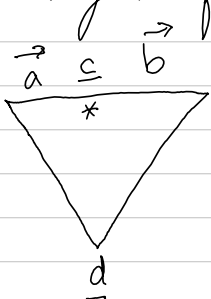


and $f: (\underline{x}, n) \rightarrow (y, p)$ to f .

Let \mathcal{C} and \mathcal{D} be multicategories.

A $(\mathcal{C}, \mathcal{D})$ -multicategory \mathcal{E} is

- a set of objects
- for each \mathcal{C} -object list \vec{a} , \mathcal{D} -object list \vec{b} and \mathcal{E} -objects \underline{c} and \underline{d} , a set of morphisms

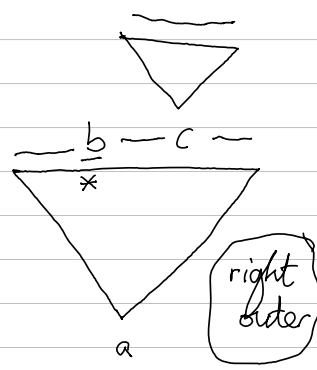
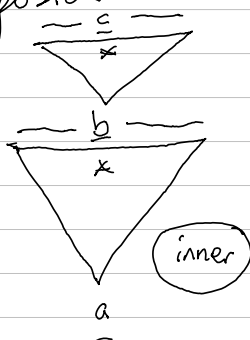
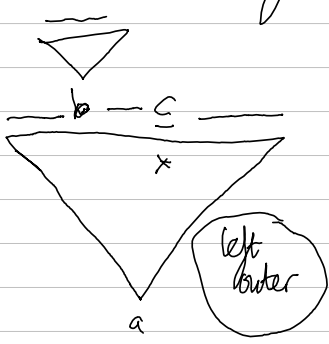


written $\mathcal{E}(\vec{a}, \underline{c}, \vec{b}; \underline{d})$

- for each \mathcal{E} -object \underline{a} , an identity morphism

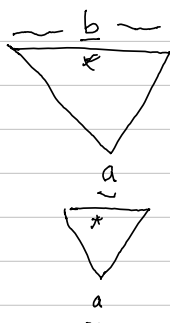


- three kinds of composition

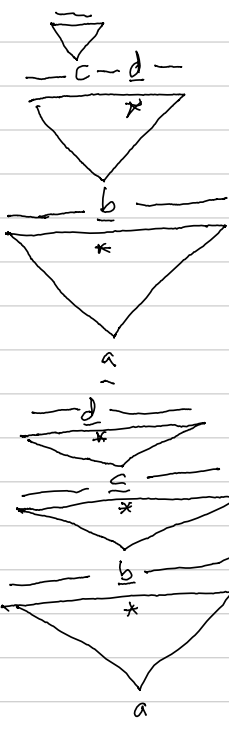
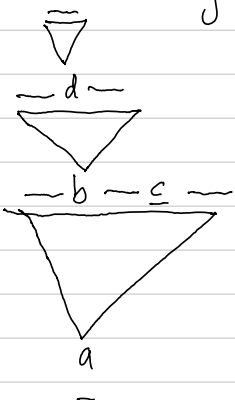


satisfying the following equations:

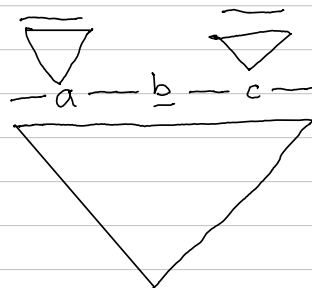
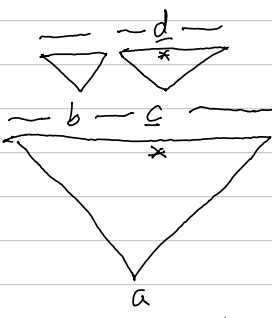
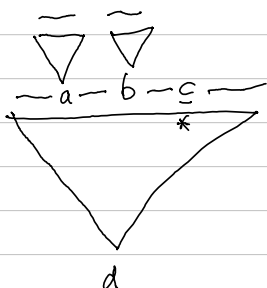
- left-unitality for each kind of composition
- right-unitality



- five associativity laws



- five interchange laws



and mirror

14 equations

Let \mathcal{C} be a left-skew multicategory and \mathcal{D} a multicategory.

A left-skew $(\mathcal{C}, \mathcal{D})$ -multicategory \mathcal{E} is

- a set of objects
- for each \mathcal{C} -object list \vec{a} , \mathcal{D} -object list \vec{b} and \mathcal{E} -objects \underline{c} and \underline{d}

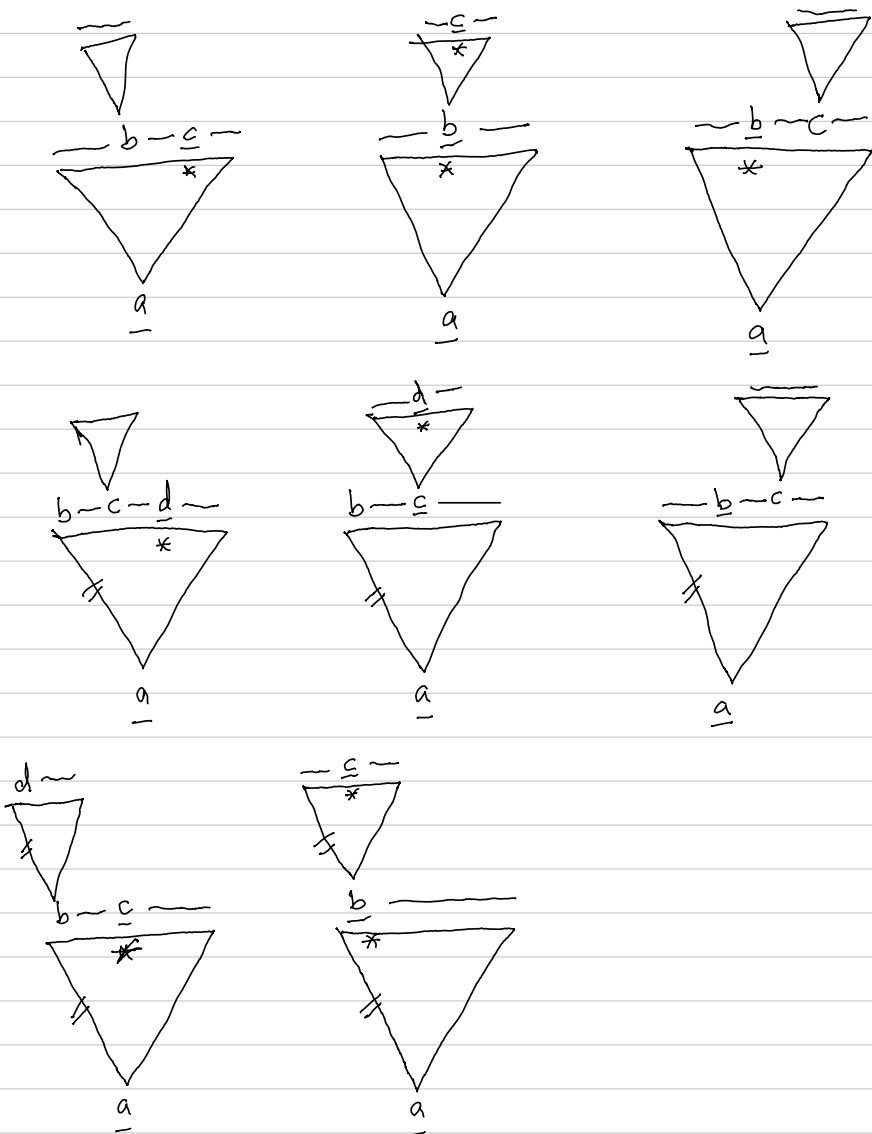
a set of loose morphisms $\vec{a} \xrightarrow{\quad} \underline{c} \xrightarrow{\quad} \vec{b}$ written $\mathcal{E}(\vec{a}, \underline{c}, \vec{b}; \underline{d})$

and a set of tight morphisms $\vec{a} \xrightarrow{\quad} \underline{c} \xrightarrow{\quad} \vec{b}$ written $\mathcal{E}([\vec{a}, \underline{c}, \vec{b}], \underline{d})$

- loosening operations $\mathcal{E}([\vec{a}, \underline{c}, \vec{b}], \underline{d}) \rightarrow \mathcal{E}(\vec{a}, \underline{c}, \vec{b}; \underline{d})$ written $f \mapsto \hat{f}$

- for each \mathcal{E} -object \underline{a} , a tight identity morphism $\underline{a} \xrightarrow{\quad} \underline{a}$

- eight kinds of composition



satisfying the following equations

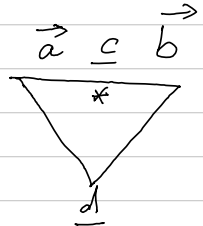
- composition commutes with loosening
- left-unitality for each kind of composition
- right-unitality for each kind of morphism
- associativity
- interchange

Let \mathcal{C} be a left-skew multicategory
and \mathcal{D} a right-skew multicategory.

A biskew $(\mathcal{C}, \mathcal{D})$ -multicategory \mathcal{E} ;

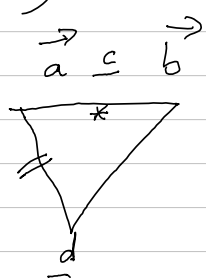
- a set of objects
- for each \mathcal{C} -object list \vec{a} and \mathcal{D} -objects \vec{b}
and \mathcal{E} -objects $\underline{c}, \underline{d}$

a set of loose-loose morphisms



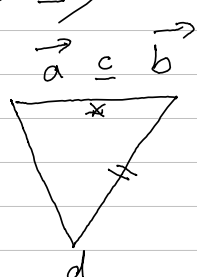
written $\mathcal{E}(\vec{a}, \underline{c}, \vec{b}; \underline{d})$

a set of tight-loose morphisms



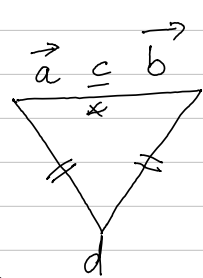
written $\mathcal{E}([\vec{a}, \underline{c}, \vec{b}; \underline{d})$

a set of loose-tight morphisms



written $\mathcal{E}(\vec{a}, \underline{c}, \vec{b}]; \underline{d})$

a set of tight-tight morphisms



written $\mathcal{E}([\vec{a}, \underline{c}, \vec{b}]; \underline{d})$

- left-loosening operations

$$\mathcal{E}([\vec{a}, \underline{c}, \vec{b}; \underline{d}) \rightarrow \mathcal{E}(\vec{a}, \underline{c}, \vec{b}; \underline{d})$$

$$\mathcal{E}([\vec{a}, \underline{c}, \vec{b}]; \underline{d}) \rightarrow \mathcal{E}(\vec{a}, \underline{c}, \vec{b}; \underline{d})$$

- right-loosening operations

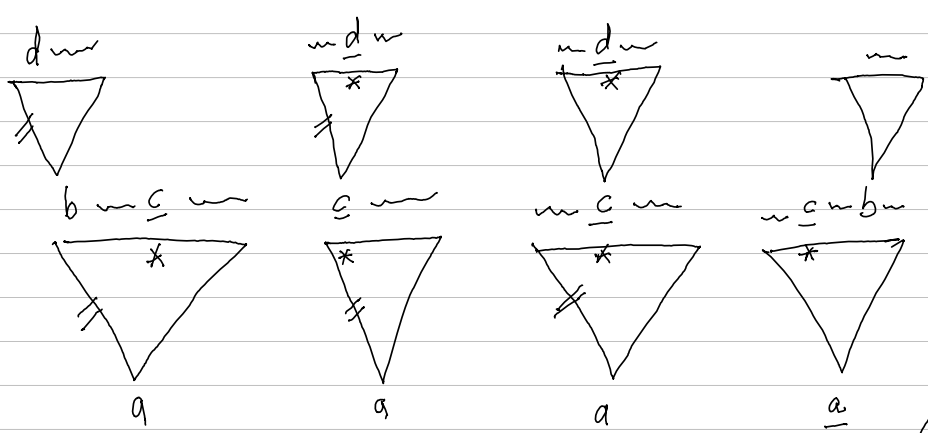
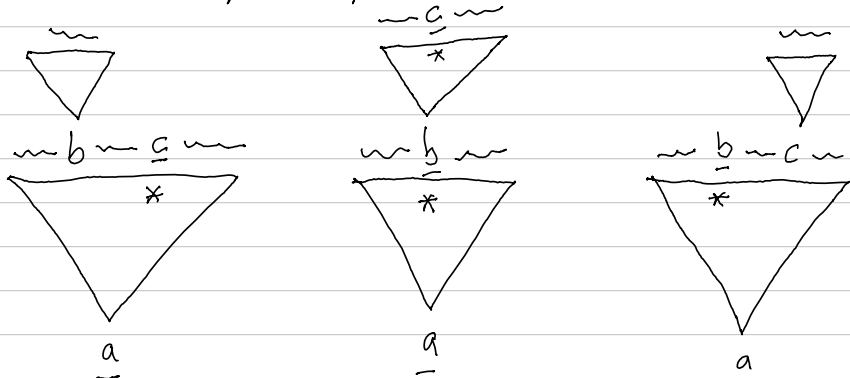
$$\mathcal{E}(\vec{a}, \underline{c}, \vec{b}]; \underline{d}) \rightarrow \mathcal{E}(\vec{a}, \underline{c}, \vec{b}; \underline{d})$$

$$\mathcal{E}([\vec{a}, \underline{c}, \vec{b}]; \underline{d}) \rightarrow \mathcal{E}([\vec{a}, \underline{c}, \vec{b}; \underline{d})$$

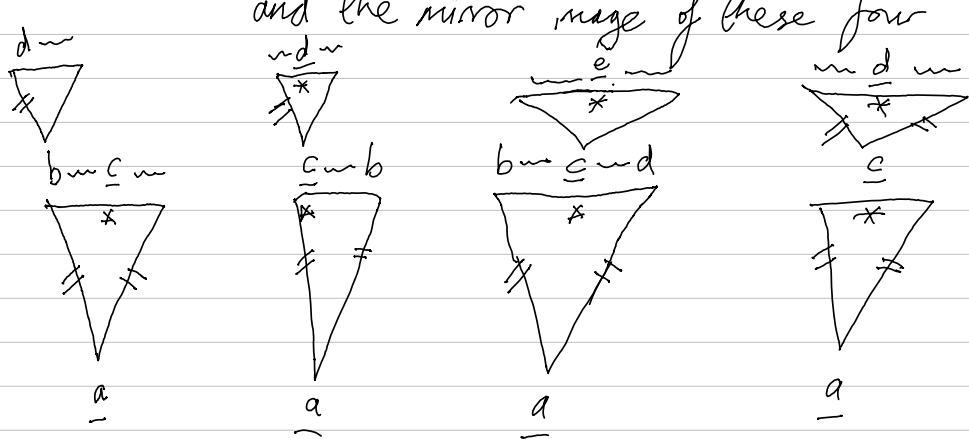
- for each \mathcal{E} -object \underline{a} , a tight-tight identity



- seventeen kinds of composition



and the mirror image of these four



and the mirror
image of these two

satisfying the following equations

- left-loosening & right-loosening commute
- loosening commutes with composition
- left unitality
- right unitality
- associativity
- interchange

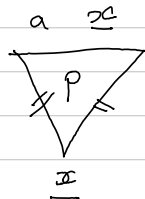
many cases

Let \mathcal{C} be a left-skew multicategory
 and \mathcal{D} a right-skew multicategory
 and \mathcal{E} a biskew $(\mathcal{C}, \mathcal{D})$ -multicategory.

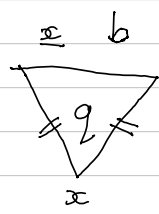
Let $M = (a, u, m)$ be a monoid in \mathcal{C}
 and $N = (b, v, n)$ a monoid in \mathcal{D} .

An (M, N) -set is

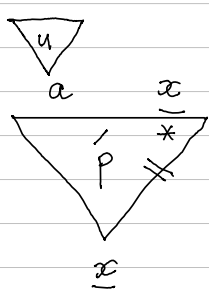
- an \mathcal{E} -object \underline{x}
- a tight-tight morphism



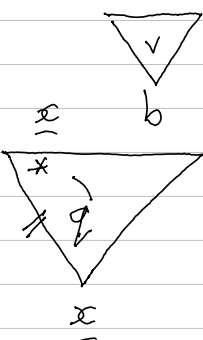
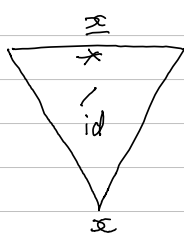
- a tight-tight morphism



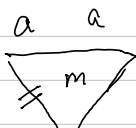
satisfying the following 5 equations.



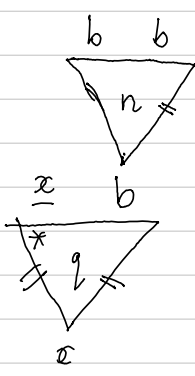
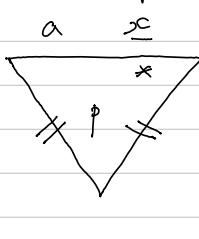
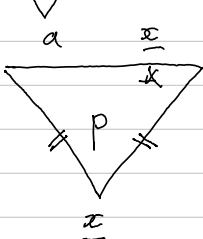
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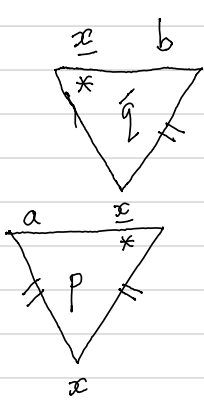
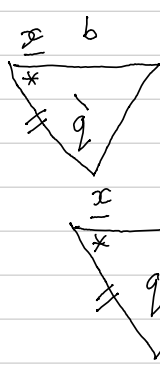
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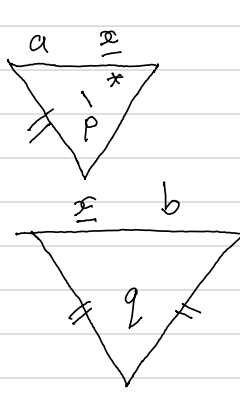
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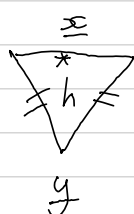


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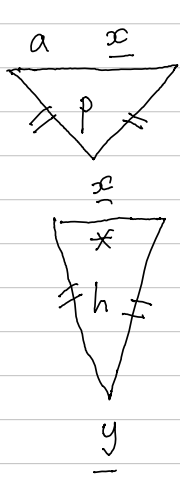


A morphism $(\underline{x}, p, q) \longrightarrow (\underline{y}, r, s)$

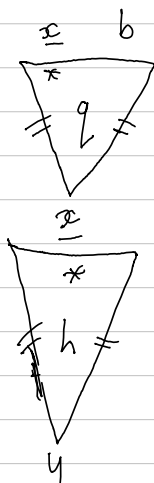
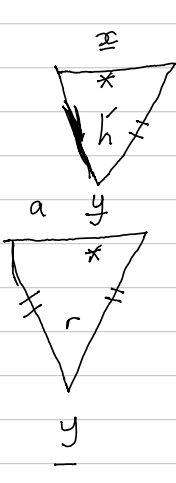
is a tight-tight morphism



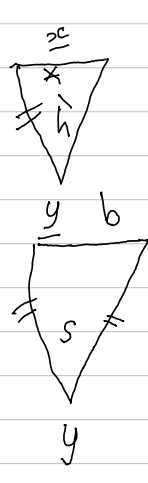
satisfying



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Monoid homomorphisms $f: M \longrightarrow M'$
 $(a, u, m) \quad (a', u', m')$

and

$g: N \longrightarrow N'$
 $(b, v, n) \quad (b', v', n')$

give a faithful functor (M', N') -set \longrightarrow (M, N) -set

It sends (\underline{x}, p, q) to $(\underline{x}, \begin{array}{c} a \\ \text{triangle} \\ a' \end{array}, \underline{x})$, $(\begin{array}{c} \underline{x} \\ \text{triangle} \\ x \end{array}, b, \underline{x})$

and $h: (\underline{x}, p, q) \longrightarrow (\underline{y}, r, s)$ to h .