



Universität  
Zürich<sup>UZH</sup>

Blockchain & Distributed Ledger Technologies

**UZH**  
Blockchain  
Center

# Statistical Properties of Networks

*Network Science '21: Session 2*

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## Lecture Objectives

1. Which kinds of distributions will we see?
2. We will go beyond degree distributions to *higher order network properties*: Assortativity and clustering
3. We will see which networks display different of these properties
4. We will start to understand the necessity for benchmark models to compare our data with



# Outlook



2

### 3 From Local to Global Properties

Shortest path

### 4 Directed Networks

Reciprocity

Network Motifs

### 5 Undirected networks

Clustering coefficient

### 6 Assortativity

### 7 Network randomisation



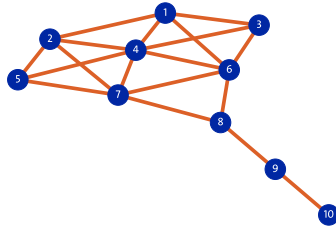
# Recap Lecture 1

## Network representation

Points to...

	1	2	3	4	5	6	7	8	9	10
1		1	1	1		1				
2	1				1	1		1		
3	1			1		1				
4	1	1	1			1	1	1		
5		1		1				1		
6	1		1	1				1	1	
7		1		1	1	1		1		
8						1	1		1	
9								1		1
10								1		

Points from ...



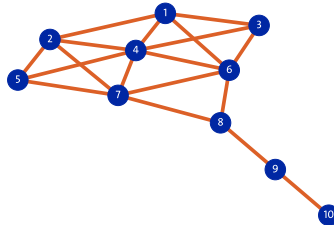
- + The adjacency matrix  $\mathbb{A} = \{a_{ij}\}_{i,j=1}^N$  is a representation of the network
- + Symmetric for undirected networks

## Node degree

Points to...

	1	2	3	4	5	6	7	8	9	10
1		1	1	1		1				
2	1			1	1		1			
3	1			1		1				
4	1	1	1		1	1	1			
5		1		1			1			
6	1		1	1				1		
7		1		1	1	1		1		
8						1	1		1	
9								1		1
10								1		

Points from node 4



The degree of a node can be computed as a sum over the adjacency matrix

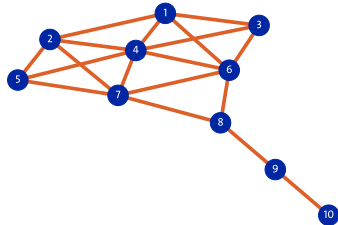
$$k_i = \sum_{j=1}^N a_{ij}$$

## Node degree

Points to...

	1	2	3	4	5	6	7	8	9	10
1		1	1	1		1				
2	1			1	1		1			
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4	1	1	1		1	1	1			
5		1		1			1			
6	1		1	1				1		
7		1		1	1	1		1		
8						1	1		1	
9								1		1
10								1		

Points from node 4



The total number of links  $L = |V|$  can be computed as a sum over the adjacency matrix,

$$L = \sum_{i < j}^N a_{ij}$$





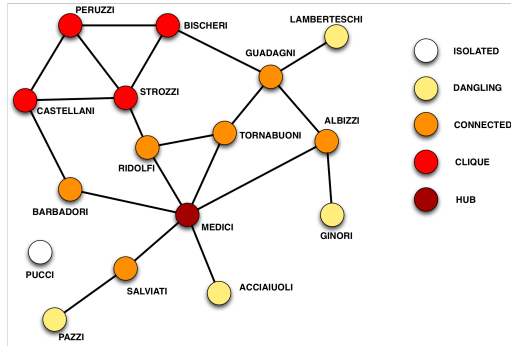
*Degree is a first-order  
property of the nodes*

$$k_i = \sum_{j=1}^N a_{ij}$$



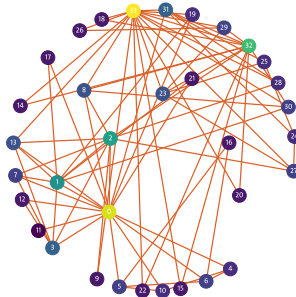
*Node degree may hint at the  
importance of specific vertices*

## Node degree: interpretation



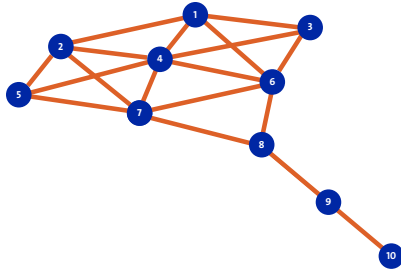
*Graph of the relationships of the Florence families in in the Renaissance*

## Node degree: interpretation



*Zachary Karate Club*

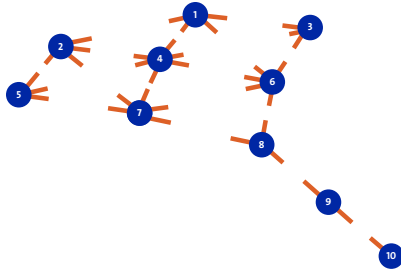
## Degree distribution



### *Definition*

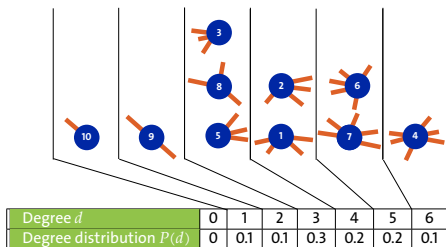
Denoted by  $p(k)$ . It is a function of the degree. Its value is the fraction of nodes with degree  $k$

## Degree distribution



*This is the information we keep from the network*

## Degree distribution



*These are the values of the degree distributions for the Kite network*

## Degree distribution

$$p(k) = \text{Prob}(k_i = k)$$

In many contexts of interest, the degree distribution does not have a characteristic scale

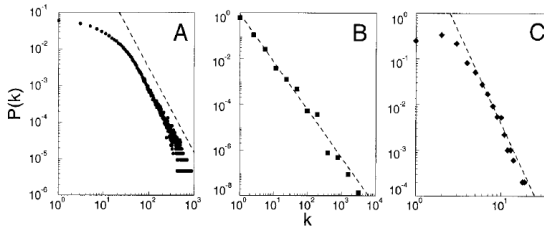


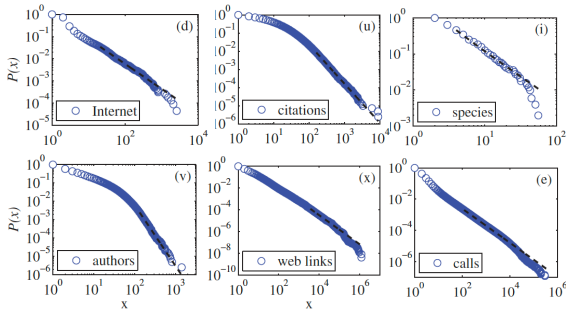
Fig. 1. The distribution function of connectivities for various large networks. (A) Actor collaboration graph with  $N = 212,250$  vertices and average connectivity  $\langle k \rangle = 28.78$ . (B) WWW,  $N = 325,729$ ,  $\langle k \rangle = 5.46$  (6). (C) Power grid data,  $N = 4941$ ,  $\langle k \rangle = 2.67$ . The dashed lines have slopes (A)  $\gamma_{\text{actor}} = 2.3$ , (B)  $\gamma_{\text{www}} = 2.1$  and (C)  $\gamma_{\text{power}} = 4$ .



## Degree distribution

$$p(k) = \text{Prob}(k_i = k)$$

In many contexts of interest, the degree distribution does not have a characteristic scale





*The reason for this  
universality will be explained  
in Session 6*



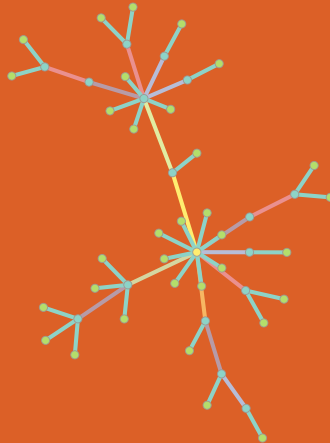
## A note on nomenclature

*The notation  $\langle \cdot \rangle$  represents the average of the argument.*

If the index is obvious, it is omitted



So far we have seen networks ...





... like this





# From Local to Global Properties

## Some examples

### + Average network degree:

$$\langle k \rangle = \frac{1}{N} \sum_{i=1}^N k_i$$

### + Average link strength (for weighted networks):

$$\langle s \rangle = \frac{1}{N^2} \sum_{i=1}^N \sum_{j=1}^N s_{ij}$$

Remember that  $N$  represents always the number of nodes

## Shortest path

- + Shortest path  $\iff$  distance between two nodes;
- + By definition, shortest path *between a node and itself* is zero;
- + If there is *no path* connecting two nodes, shortest path is defined to be *infinite*
- +  $A^r$  gives the number of walks of length  $r$  between each pair of nodes

Shortest path

$$d_{i,j} = \operatorname{argmin}_r ([A^r]_{i,j} \neq 0) \quad (1)$$





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### Shortest path

$$d_{i,j} = \operatorname{argmin}_r ([\mathbb{A}^r]_{i,j} \neq 0) \quad (1)$$

## Related concepts

- + **Average distance** in a network is the average length of the shortest path between two nodes.

$$\langle d_{i,j} \rangle = \frac{1}{(N-1)^2} \sum_{i,j=1}^N d_{i,j}$$

- + **Eccentricity** of a vertex  $e(v)$  is the greatest distance between  $v$  and any other vertex
- + **Diameter** of a graph is maximum eccentricity of any vertex in the graph (or the largest distance between pairs of nodes)

$$D = \max_{i,j} (d_{i,j})$$

- + A **peripheral vertex** is the one whose eccentricity achieves the diameter

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# Directed Networks

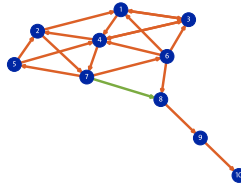


*Are relationships mutual? If  $i$  claims to be friend of  $j$ , does  $j$  think the same about  $i$ ?*

## Reciprocity

For *directed* networks, reciprocity measures the degree of how mutual relationships are

	1	2	3	4	...	to...	5	6	7	8	9	10
1	0	0	1	1	0	0	0	0	0	0	0	0
2	1	0	0	0	0	0	0	0	1	0	0	0
3	1	0	0	1	0	0	0	0	0	0	0	0
4	0	1	1	0	0	0	1	0	0	0	0	0
5	0	1	0	1	0	0	0	0	0	0	0	0
6	1	0	1	1	0	0	0	1	0	0	0	0
7	0	0	0	0	1	1	0	1	0	1	0	0
8	0	0	0	0	0	0	0	0	0	0	1	0
9	0	0	0	0	0	0	0	0	0	0	1	1
10	0	0	0	0	0	0	0	0	0	0	0	0

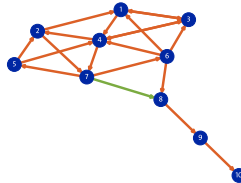


$$R = \frac{\sum_{u,v \in V} a_{u,v} a_{v,u}}{2 \sum_{u,v \in V} a_{u,v}} \quad (2)$$

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1	0	0	1	1	0	0	0	0	0	0
2	1	0	0	0	0	0	0	1	0	0
3	1	0	0	1	0	0	0	0	0	0
4	0	1	1	0	0	0	1	0	0	0
5	0	1	0	1	0	0	0	0	0	0
6	1	0	1	1	0	0	0	1	0	0
7	0	0	0	0	1	1	0	1	0	0
8	0	0	0	0	0	0	0	0	1	0
9	0	0	0	0	0	0	0	0	0	1
10	0	0	0	0	0	0	0	0	0	0



$$R = \frac{\sum_{u,v \in V} a_{u,v} a_{v,u}}{2 \sum_{u,v \in V} a_{u,v}} \quad (2)$$



## Why does it work?

+  $a_{u,v}a_{v,u}$  is only equal to 1 if both edges exist

*Reciprocity is a second-order property because it involves two nodes at the same time*

Conversely, degree is a first order property



*How is friendship data  
collected? Surveys!*

## Empirical results

Reciprocity by Gender and Race (percentages)

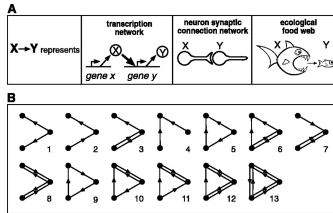
	Males		Females	
	Reciprocal	Nonreciprocal	Reciprocal	Nonreciprocal
Interracial	48.91%	51.09%	65.29%	34.71%
Intraracial	58.87	41.13	72.61	27.61
White	61.00	39.00	75.78	24.22
Black	46.17	53.83	62.39	37.61
Hispanic	48.75	51.25	65.35	34.65
Asian	59.44	40.56	69.09	30.91
Native American	52.22	47.78	59.04	40.96
Overall	56.98	43.02	71.49	28.51
N	10,635	8,030	15,743	6,278

*“Friendship Reciprocity and Its Effects on School Outcomes among Adolescents”* [E.

Vaquera, G. Kao, *Social Science Research* (2008)]

- + Many friendship ties are non-reciprocal
- + Reciprocity decreases across self-perceived groups
- + In organisations, reciprocity is higher within individuals of the same rank, much lower across ranks

## Network motifs: reconstructing the network



*It is possible to decompose the network into smaller sub-graphs representing different smaller-scale processes*





# Undirected networks



*How likely that two of my  
friends are friends among  
themselves?*

## Clustering coefficient

*How likely is it that if two nodes are common neighbours of another, they are connected to each other?*

Each node  $i$  is endowed with a clustering coefficient  $C_i$

$$C_i = \frac{\sum_{u,v \in V} a_{u,i} a_{v,i} a_{u,v}}{\sum_{u,v \in \mathcal{N}_i} a_{u,i} a_{v,i}} = \frac{2 \sum_{u,v \in \mathcal{N}_i} a_{u,v}}{k_i(k_i - 1)} \quad (3)$$

- +  $\mathcal{N}_i$  is the set of neighbours of node  $i$
- + *number of closed triangles divided over number of all possible triangles*

## Clustering coefficient

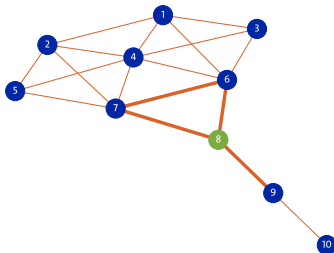
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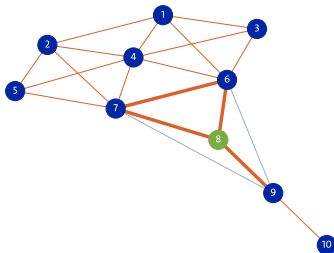
## Clustering coefficient



For node 8, there is 1 triangle realised from 3 triangles possible

$$C_8 = \frac{1}{3}$$

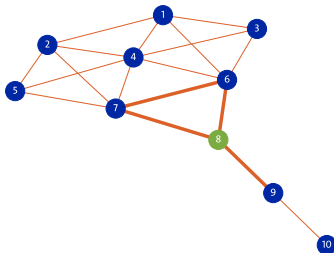
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## Clustering coefficient



The list of all nodes' clustering

Node $i$	1	2	3	4	5	6	7	8	9	10
Clustering $C_i$	2/3	2/3	1	8/15	1/2	1	1/2	1/3	0	0

## Global clustering coefficient

It gives a measure of the average value of the clustering coefficient in all the network

$$\langle C \rangle = \frac{1}{N} \sum_i C_i \quad (4)$$

For the example Kite network,  $\langle C \rangle = 0.52$



# Empirical values of the clustering coefficient

TABLE I. The general characteristics of several real networks. For each network we have indicated the number of nodes, the average degree  $\langle k \rangle$ , the average path length  $\ell$ , and the clustering coefficient  $C$ . For a comparison we have included the average path length  $\ell_{rand}$  and clustering coefficient  $C_{rand}$  of a random graph of the same size and average degree. The numbers in the last column are keyed to the symbols in Figs. 8 and 9.

Network	Size	$\langle k \rangle$	$\ell$	$\ell_{rand}$	$C$	$C_{rand}$	Reference	Nr.
WWW, site level, undir.	153 127	35.21	3.1	3.35	0.1078	0.00023	Adamic, 1999	1
Internet, domain level	3015–6209	3.52–4.11	3.7–3.76	6.36–6.18	0.18–0.3	0.001	Yook <i>et al.</i> , 2001a, Pastor-Satorras <i>et al.</i> , 2001	2
Movie actors	225 226	61	3.65	2.99	0.79	0.00027	Watts and Strogatz, 1998	3
LANL co-authorship	52 909	9.7	5.9	4.79	0.43	$1.8 \times 10^{-4}$	Newman, 2001a, 2001b, 2001c	4
MEDLINE co-authorship	1 520 251	18.1	4.6	4.91	0.066	$1.1 \times 10^{-5}$	Newman, 2001a, 2001b, 2001c	5
SPIRES co-authorship	56 627	173	4.0	2.12	0.726	0.003	Newman, 2001a, 2001b, 2001c	6
NCSTRL co-authorship	11 994	3.59	9.7	7.34	0.496	$3 \times 10^{-4}$	Newman, 2001a, 2001b, 2001c	7
Math. co-authorship	70 975	3.9	9.5	8.2	0.59	$5.4 \times 10^{-5}$	Barabási <i>et al.</i> , 2001	8
Neurosci. co-authorship	209 293	11.5	6	5.01	0.76	$5.5 \times 10^{-5}$	Barabási <i>et al.</i> , 2001	9
<i>E. coli</i> , substrate graph	282	7.35	2.9	3.04	0.32	0.026	Wagner and Fell, 2000	10
<i>E. coli</i> , reaction graph	315	28.3	2.62	1.98	0.59	0.09	Wagner and Fell, 2000	11
Ythan estuary food web	134	8.7	2.43	2.26	0.22	0.06	Montoya and Solé, 2000	12
Silwood Park food web	154	4.75	3.40	3.23	0.15	0.03	Montoya and Solé, 2000	13
Words, co-occurrence	460 902	70.13	2.67	3.03	0.437	0.0001	Ferrer i Cancho and Solé, 2001	14
Words, synonyms	22 311	13.48	4.5	3.84	0.7	0.0006	Yook <i>et al.</i> , 2001b	15
Power grid	4941	2.67	18.7	12.4	0.08	0.005	Watts and Strogatz, 1998	16
<i>C. Elegans</i>	282	14	2.65	2.25	0.28	0.05	Watts and Strogatz, 1998	17



## Empirical findings

- + Real-world networks tend to have a large values of clustering



# Assortativity



*Homophily in social sciences  
refers to the propensity of an  
individual to link to others  
that are alike*



## Assortativity

- + *Is is there homophily in the network? (nodes connect to others similar to them?)*
- + Are nodes of low-degree connected with other low-degree nodes or are they mostly connected with high-degree nodes?
- + Are nodes of high-degree connected with other low-degree nodes or are they mostly connected with high-degree nodes?

## Assortativity

Assortativity is usually computed as the Pearson correlation coefficient of degree between pairs of linked nodes

$$r = \text{Corr}(k_i, k_j) = \frac{\text{Cov}(k_i, k_j)}{\text{Var}(k_i)^2} = \frac{\langle k_i k_j \rangle - \langle k_i \rangle \langle k_j \rangle}{\sigma^2} \quad (5)$$

where

$$\langle k_i k_j \rangle - \langle k_i \rangle \langle k_j \rangle = \sum_{(i,j) \in V} k_i k_j (P(k_i, k_j) - P(k_i)P(k_j)) \quad (6)$$

and

$$\sigma^2 = \sum_k k^2 P(k) - \left( \sum_k k P(k) \right)^2 \quad (7)$$

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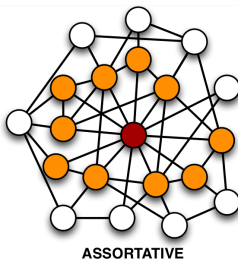
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and

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## Interpretation: Assortativity



*if  $0 < r < 1$ , the network is said to be assortative: nodes with large (resp. low) degree have neighbours with large (resp. low) degree*



## Interpretation: Dissortativity



*if  $-1 < r < 0$ , the network is said to be dissortative:  
nodes with large (resp. low) degree have neighbours with  
low (resp. high) degree*



## Empirical values

Some real-world examples

Network	$n$	$r$
Physics coauthorship (a)	52 909	0.363
Biology coauthorship (a)	1 520 251	0.127
Mathematics coauthorship (b)	253 339	0.120
Film actor collaborations (c)	449 913	0.208
Company directors (d)	7 673	0.276
Internet (e)	10 697	-0.189
World-Wide Web (f)	269 504	-0.065
Protein interactions (g)	2 115	-0.156
Neural network (h)	307	-0.163
Marine food web (i)	134	-0.247
Freshwater food web (j)	92	-0.276



## Empirical findings

- + Technological networks → *dissortative*
- + Social networks → *assortative*

## Nearest neighbours' degree

*Another means of capturing the degree correlation is by examining the properties of neighbour connectivity*

$$k_{nn}(k_i) = \frac{1}{k_i} \sum_j a_{i,j} k_j \quad (8)$$

## Nearest neighbours' degree

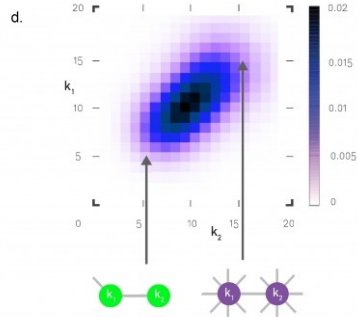
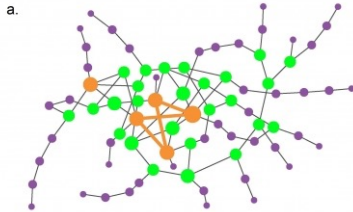
*Another means of capturing the degree correlation is by examining the properties of neighbour connectivity*

It is then averaged over all nodes with degree  $k$

$$\langle k_{nn}(k) \rangle = \sum_q q P(q|k) \quad (8)$$

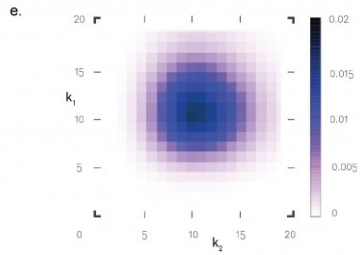
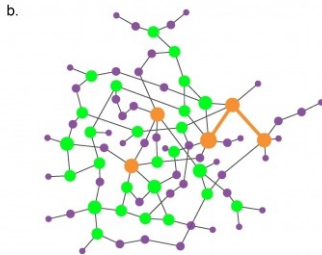
where  $P(q|k)$  is the conditional probability that, given that the node degree is  $k$ , one of its edges points to a node whose degree is  $q$

# Assortativity



If  $\langle k_{nn}(k) \rangle$  is increasing on  $k$ , the network is assortative

# Assortativity



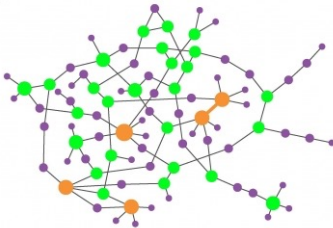
c.

f.

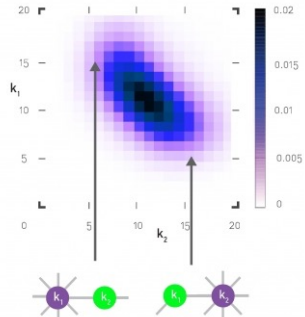
If  $\langle k_{nn}(k) \rangle$  does not have a trend, the network is not assortative nor dissortative

## Assortativity

c.



f.



If  $\langle k_{nn}(k) \rangle$  is decreasing on  $k$ , the network is dissortative





# Network randomisation



## The typical research cycle

- + Data collection
- + Data analysis (e.g. measure clustering or assortativity)
- + What is the meaning of the results? Is it relevant (statistically significant)?
- + What is the origin of the result?



*The approach taken is:  
compare the results with those  
coming from a benchmark  
model*



## The typical research cycle

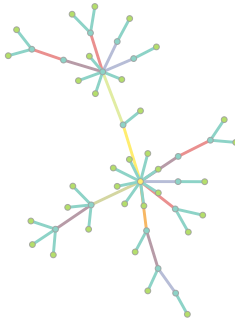
- + Are the results of the benchmark model the same as in the data? **Results are not significant**
- + Are the results of the benchmark model different from those in the data? **Results are significant**

## Degree-preserving network randomisation

*This approach asks the question if the results are purely the effect of the degree distribution*

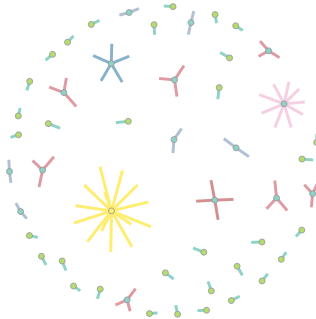
- + If we change the network completely, but leave the nodes' degrees unaltered... do we obtain the same results for the property we are measuring?
- + If the result of our observation *does not change*... then, *the result stems from the degree distribution*
- + If the result of our observation *does change*... then, *the result comes from higher order network properties*

## Turning a network into pieces



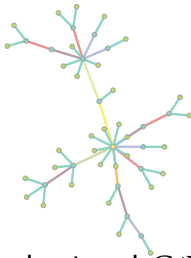
- + Given a network  $G(V, E)$  we cut the edges
- + Each isolated node  $i$  has  $k_i$  stubs
- + Randomly connecting pairs of stubs, we can build other networks with the same degree distribution...

## Turning a network into pieces



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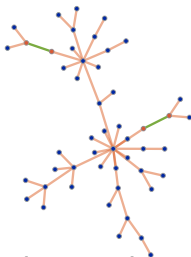
## A non-destructive approach



1. Start from the original network  $G(V, E)$
2. Randomly select two edges
3. Exchange ends
4. Repeat steps 2-3 many times (at least, until each edge was selected - on average - once). We obtain  $G_{rnd}$

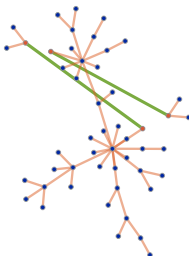


## A non-destructive approach



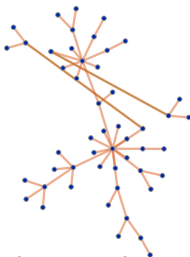
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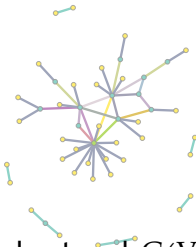
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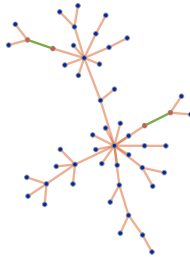
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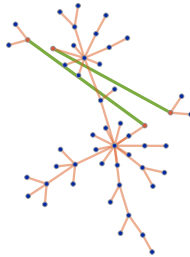
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## No change in the degree distribution!



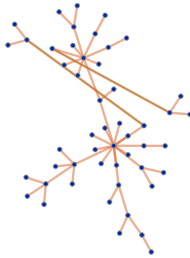
*During an edge swap, the nodes affected do not change their degree!*

## No change in the degree distribution!



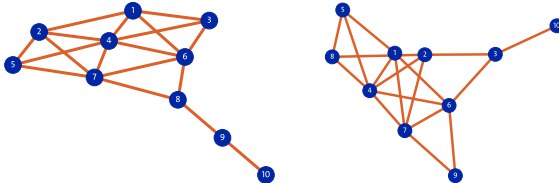
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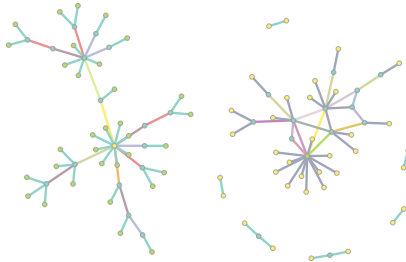
## Example in the kite network



*Kite network before (left) and after (right) randomisation*



## Comparison of the results



Measure a property in  $G$ . Randomise. Measure the property in  $G_{rnd}$  and check if they are different

*To attain statistical significance, multiple realisations of  $G_{rnd}$  are necessary*



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