

Phillip booth
Question 31

$(\text{length } (\text{reverse } xs)) = (\text{length } xs)$

Base Case: xs is nil

= {substitute simple-reverse}

$(\text{length } (\text{if null? '() '() } (\text{append } (\text{simple-reverse } (\text{cdr '()})) (\text{list1 } (\text{car '()}))))))$

= {null-empty law}

$(\text{length } (\text{if null? \#t '() } (\text{append } (\text{simple-reverse } (\text{cdr '()})) (\text{list1 } (\text{car '()}))))))$

= {if-#t law}

(length '())

$(\text{length } xs)$ when xs is null

Inductive Step: xs is guaranteed not to be null

$(\text{length } (\text{simple-reverse } xs))$ where $xs = (\text{cons } z \text{ } zs)$

= {using inductive step to show that xs does not equal nil}

$(\text{length } (\text{simple-reverse } (\text{cons } z \text{ } zs)))$

= {substitute in simple reverse}

$(\text{length } (\text{if null? } (\text{cons } z \text{ } zs) \text{ '() } (\text{append } (\text{simple-reverse } zs) (\text{list1 } z))))))$

= {if-#f law}

$(\text{length } (\text{append } (\text{simple-reverse } zs) (\text{list1 } z))))))$

= {length-append law}

$(+ 1 (\text{length } (\text{simple-reverse } zs)) (\text{length } (\text{list1 } z)))$

= {length-list1}

$(+ 1 (\text{length } (\text{simple-reverse } zs)))$

= {induction hypothesis}

$(+ 1 (\text{length } zs))$

$=\{\text{cons-length law}\}$

$(\text{length } (\text{cons } z \text{ } zs))$

$(\text{length } xs)$