

# **EFFICIENT ANSWERING OF SET CONTAINMENT QUERIES FOR SKEWED ITEM DISTRIBUTIONS**

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# THE PROBLEM

- We have huge collections of transactional data
  - Retail store transaction logs
  - Web logs
  - Biomedical databases etc.
- We address the efficient evaluation of boolean containment queries
  - In which transactions were products 'a' and 'b' sold together?
  - Which users visited only the main page or the download page of our site?
- The problem of indexing set values in secondary storage remains a challenge still
  - Inverted files are the state-of-the-art but long lists can slow them significantly
  - Skewed distributions cause long lists
  - Most of the recent efforts are concerned with more complicated queries or main memory solutions
  - The problem still exists
- *We propose the Ordered-Inverted file (OIF) index*

# QUERIES - SUBSET

<i>tid</i>	products	<i>tid</i>	products
1	{f,a}	9	{a,e}
2	{a,d,c}	10	{g,c,a}
3	{c,b,a}	11	{b,a,e}
4	{f,a,c}	12	{b,d,c}
5	{c,g}	13	{c,f,a,d,b}
6	{a,b,g,c,d,e}	14	{b,d}
7	{a,d,b}	15	{e}
8	{a,e,b}	16	{b,f,a}

■ Find all transactions that contain 'a', 'b' and 'd' (subset)

# QUERIES - EQUALITY

<i>tid</i>	products	<i>tid</i>	products
1	{f,a}	9	{a,e}
2	{a,d,c}	10	{g,c,a}
3	{c,b,a}	11	{b,a,e}
4	{f,a,c}	12	{b,d,c}
5	{c,g}	13	{c,f,a,d,b}
6	{a,b,g,c,d,e}	14	{b,d}
7	{a,d,b}	15	{e}
8	{a,e,b}	16	{b,f,a}

■ Find all transactions that contain 'a', 'b' and 'd' (subset)

■ Find all transactions that contain exactly 'a', 'b' and 'd' (equality)

# QUERIES - SUPERSET

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1	{f,a}	9	{a,e}
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■ Find all transactions that contain 'a', 'b' and 'd' (subset)

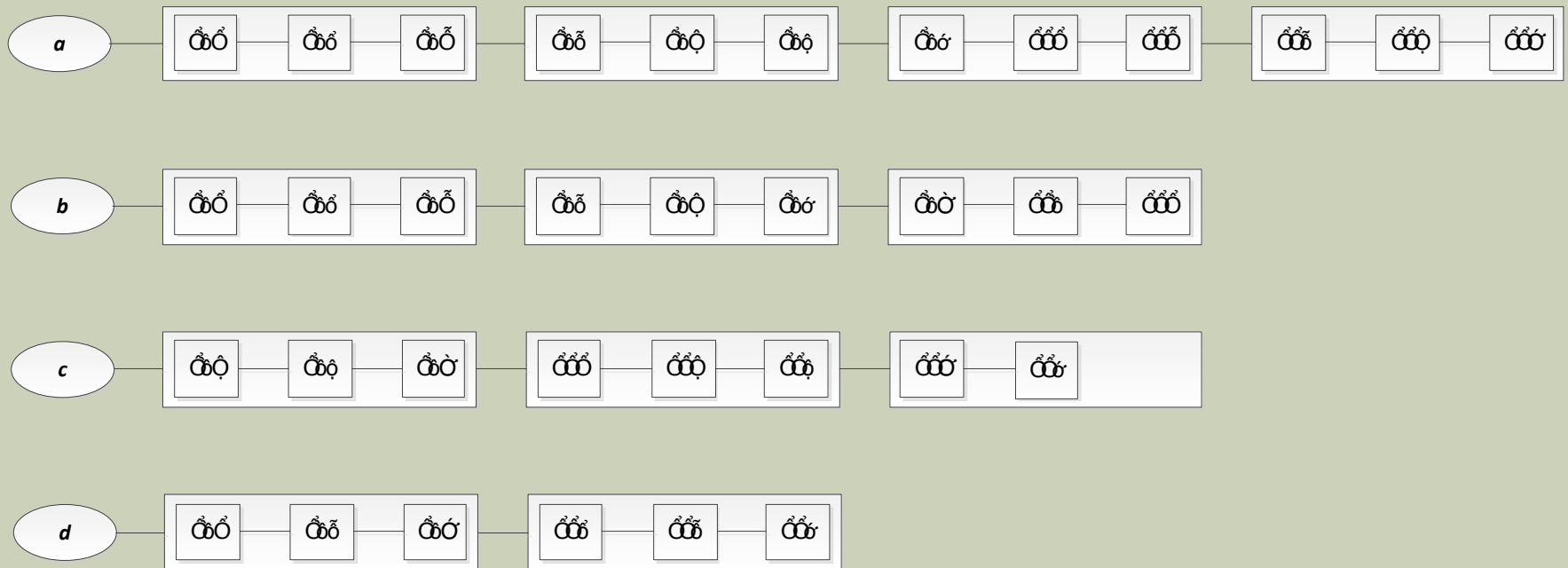
■ Find all transactions that contain exactly 'a', 'b' and 'd' (equality)

■ Find all transactions that contain only items from 'a', 'b' and 'd' (superset)

# ORIGINAL DATA

id	Items	id	Items	id	Items
101	{g, b, a, d}	107	{d, h}	113	{a}
102	{a, e, b}	108	{b, a, f}	114	{a, d}
103	{f, e, a, b}	109	{b, c}	115	{j, c, a}
104	{d, b, a}	110	{j, b, g}	116	{i, c}
105	{a, b, f, c}	111	{a, c, b }	117	{a, c, h}
106	{c, a}	112	{i, d}	118	{d, c}

# SIMPLE INVERTED FILE



# MOTIVATION

- Lists grow large
  - When the (database size) / (domain size) ratio grows big
  - When the item distribution is skewed
- Using indexing on lists (e.g., skip lists, b-tree) does not offer by itself a big advantage
  - Only if the selectivities are extremely small
- Most frequent items are the key
  - Most frequent items cause most problems
  - Most frequent items are most frequently involved in queries
- **Basic Idea**
  - **Introduce a global order to records**
  - **Records that contain common frequent items are placed close to each other**
  - **Postings in lists follow the same order**

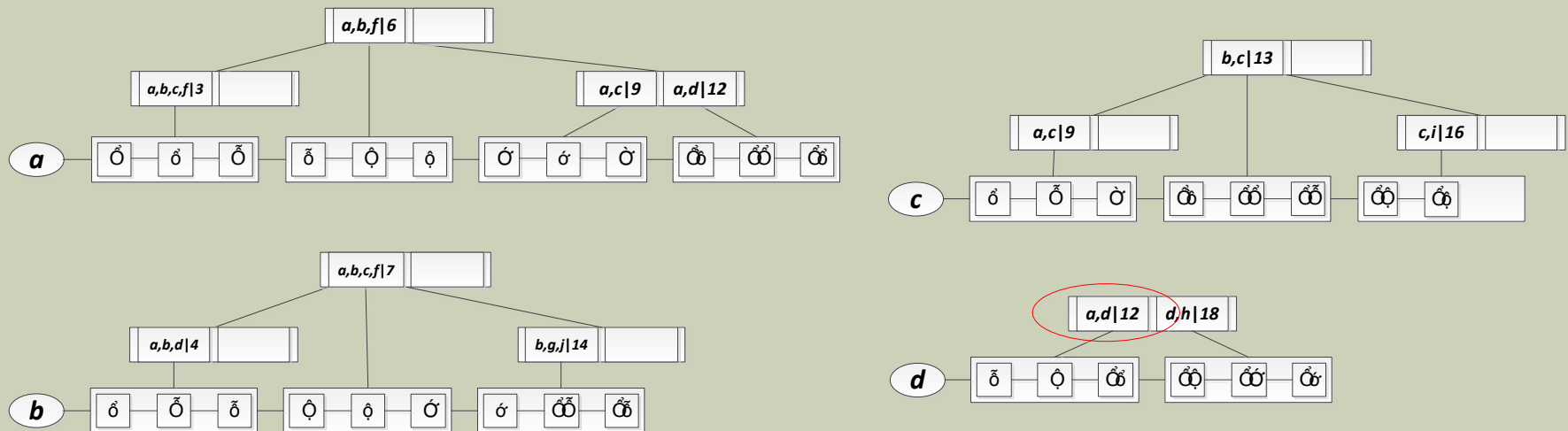


# ORDERED DATA

id	Items	Id	Items	id	Items
1	{a}	7	{a, b, f, e}	13	{b, c}
2	{a, b, c }	8	{a, b, e}	14	{b, g, j}
3	{a, b, c, f}	9	{a, c}	15	{c, d}
4	{a, b, d}	10	{a, c, h}	16	{c, i}
5	{a, b, d, g}	11	{a, c, j}	17	{d, i}
6	{a, b, f}	12	{a, d}	18	{d, h}

- A global order for records
- Records sorted internally by frequency (most frequent first)
- Dataset sorted lexicographically (items are compared based on their support)
- New ids reflecting the new order are assigned (storage might not reflect the new order)

# ORDERED INVERTED FILE

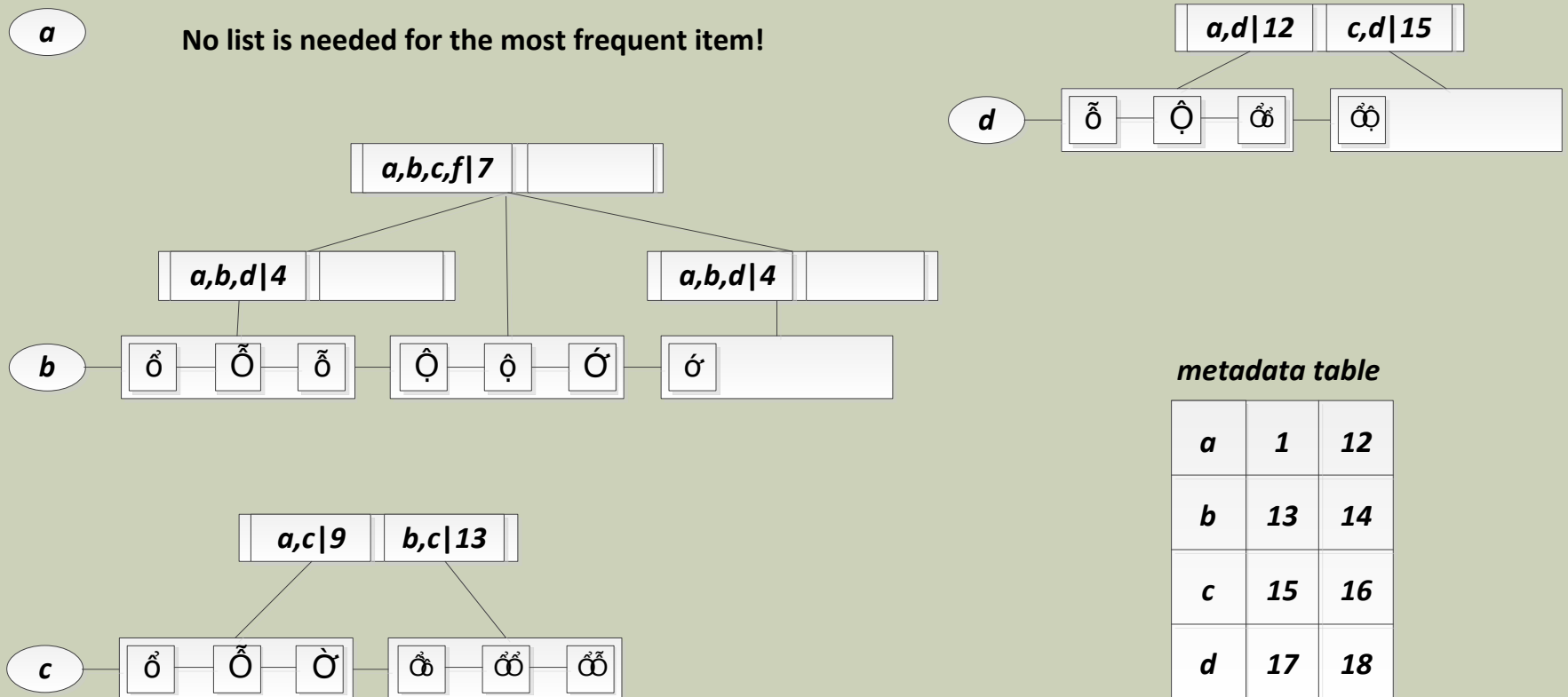


- Lists are created over the new ids
- Indexed by a sparse B-tree
- One key per page

key example

**a,d | 12**  
record content | record id

# ORDERED INVERTED FILE METADATA TABLE



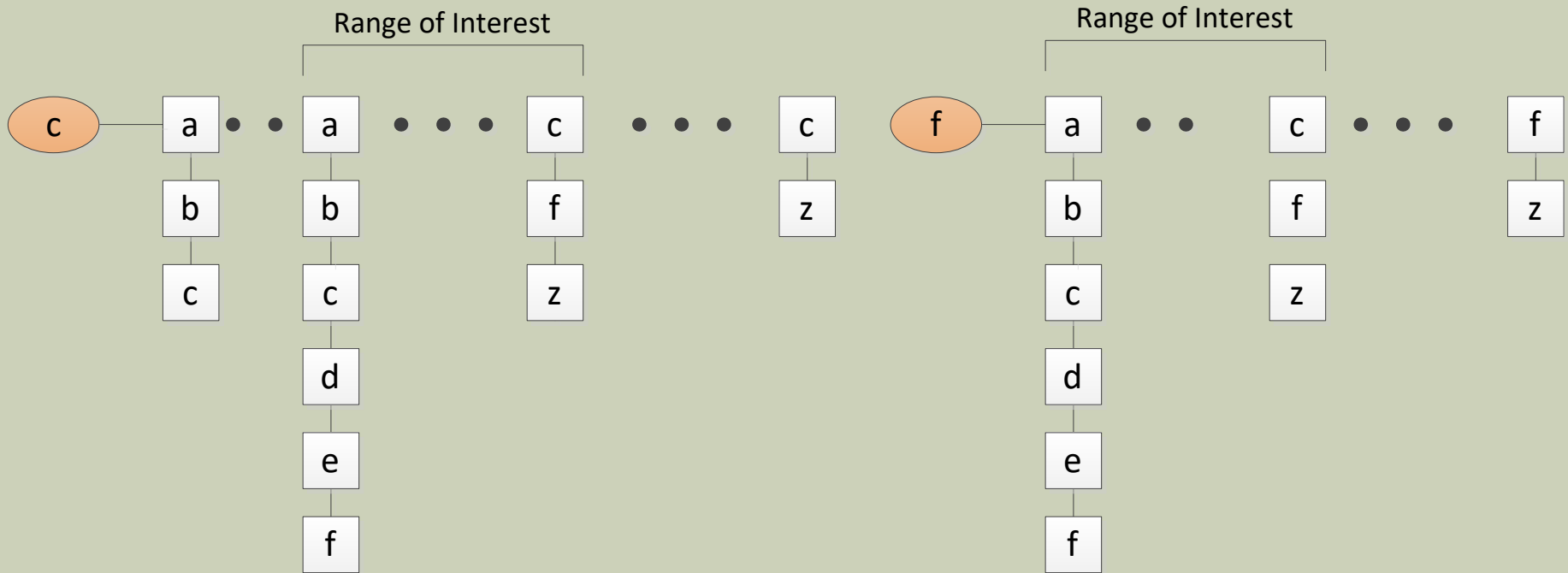
- In the list of item  $x$  we do not need to keep explicit postings for records whose most frequent item is  $x$
- We save  $1/(\text{average record size})$  of the total postings

# QUERY EVALUATION

- The evaluation of containment on OIF is similar to the evaluation on IF
- The difference is that instead of intersecting the whole lists, we intersect only a part of them
- 1. Find the bounds that contain possible answers in each list
  - Range of Interest (RoI) of each query
  - The Range of Interest are calculated using only the query set
- 2. Use the B-tree to access the block that contains the lower bound
- 3. Intersect only the part of the lists that is between the bounds
- Observations:
  - We use the metadata for verifying some results
  - Equality query becomes a value retrieval query on the B-tree

# SUBSET QUERIES

$QS = \{c, f\}$



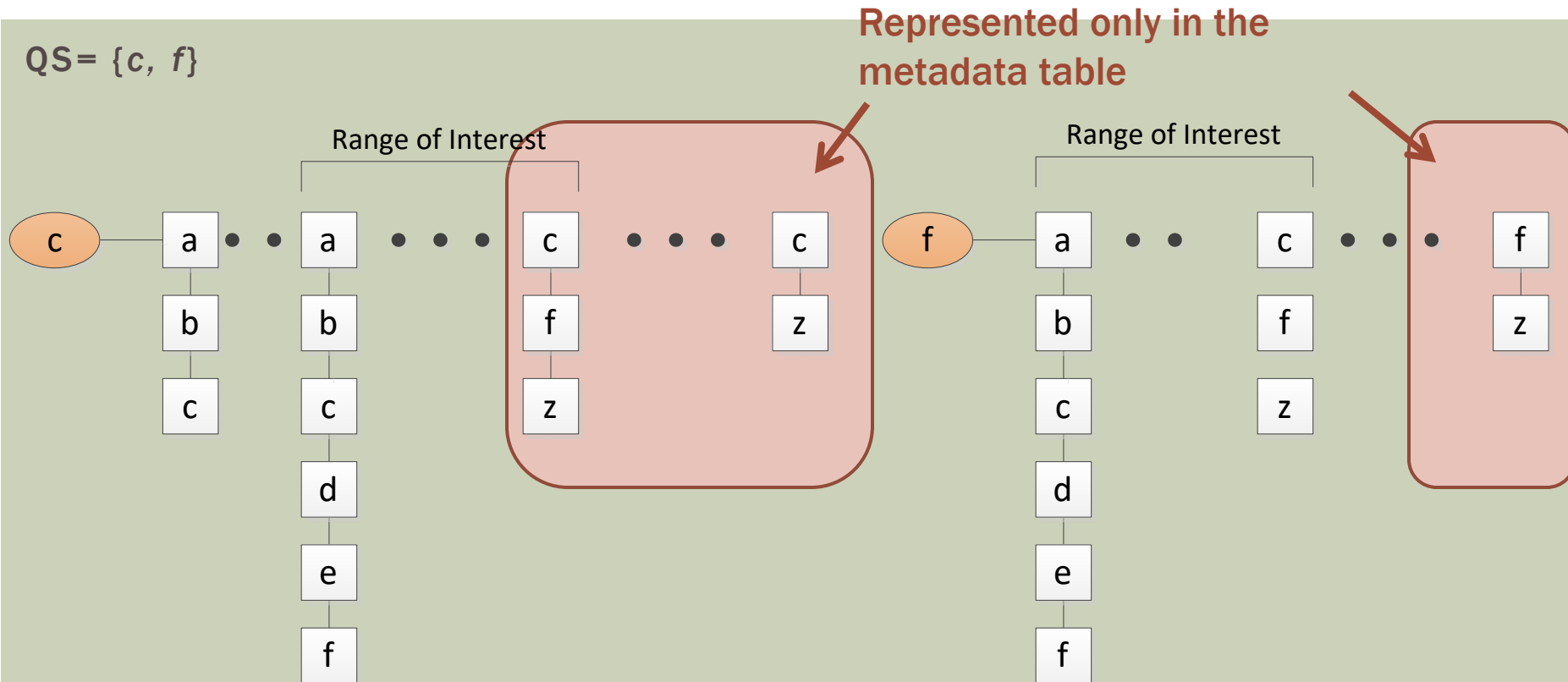
In  $QS = o_{q1}, \dots, o_{qn}$

Lower bound =  $a, b, \dots, o_{qn}$

upper bound =  $o_{q1}, \dots, o_{qn}, z$

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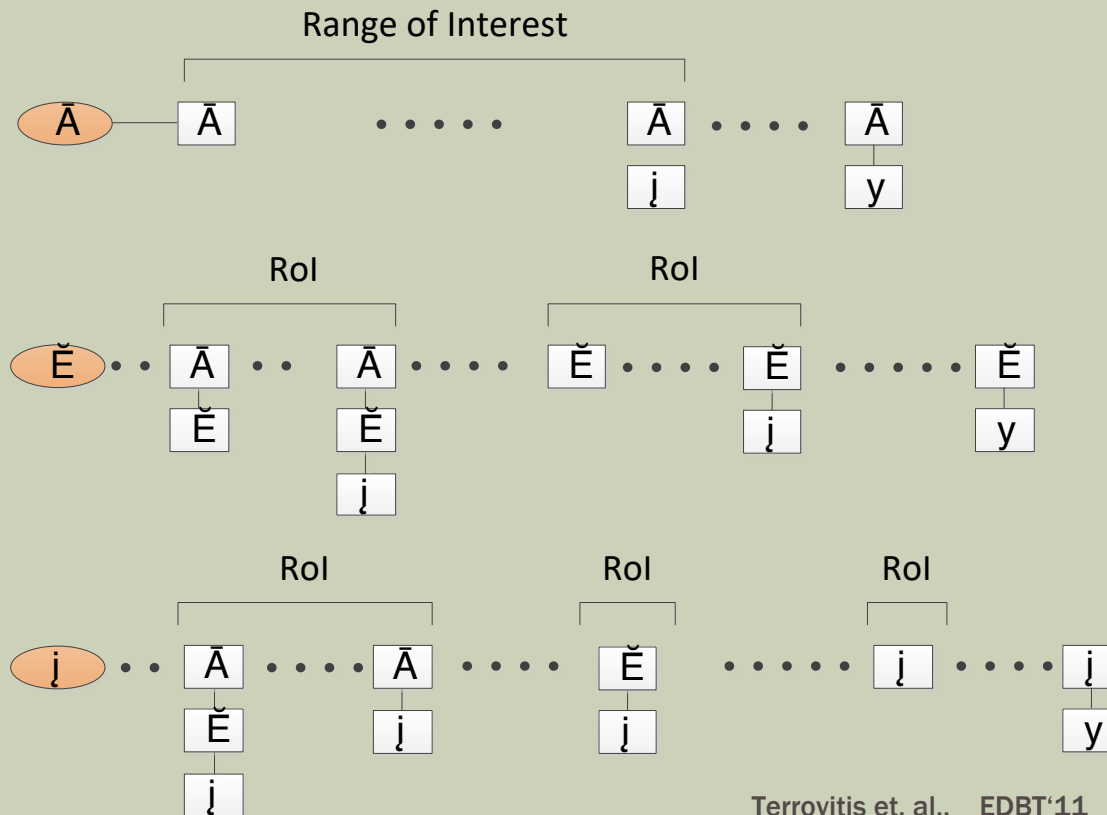
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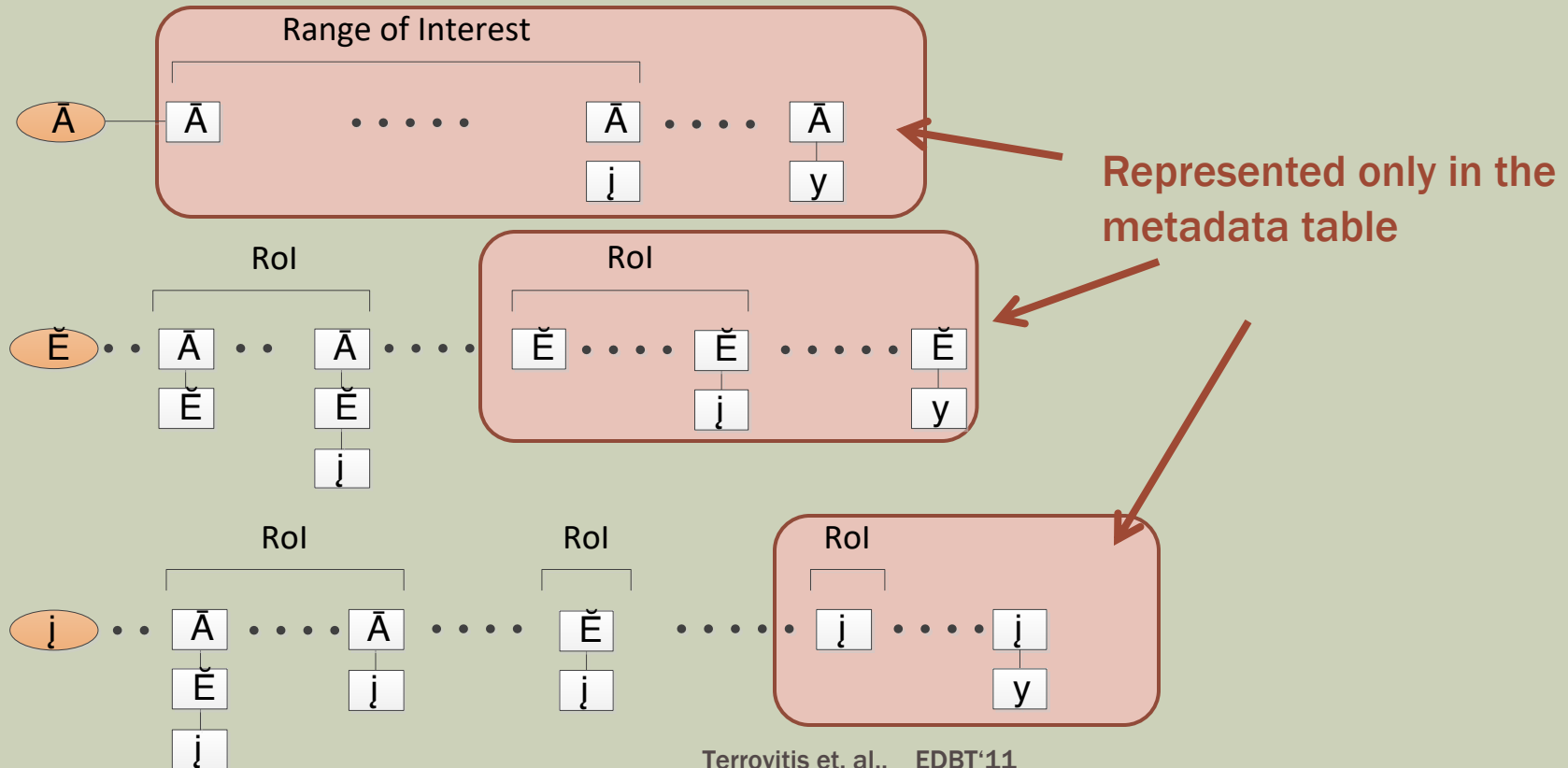
# SUPERSET QUERIES

- $QS = \{a, c, f\}$
- Its  $2^{|QS|}-1$  equality queries
- We group them by the most frequent item



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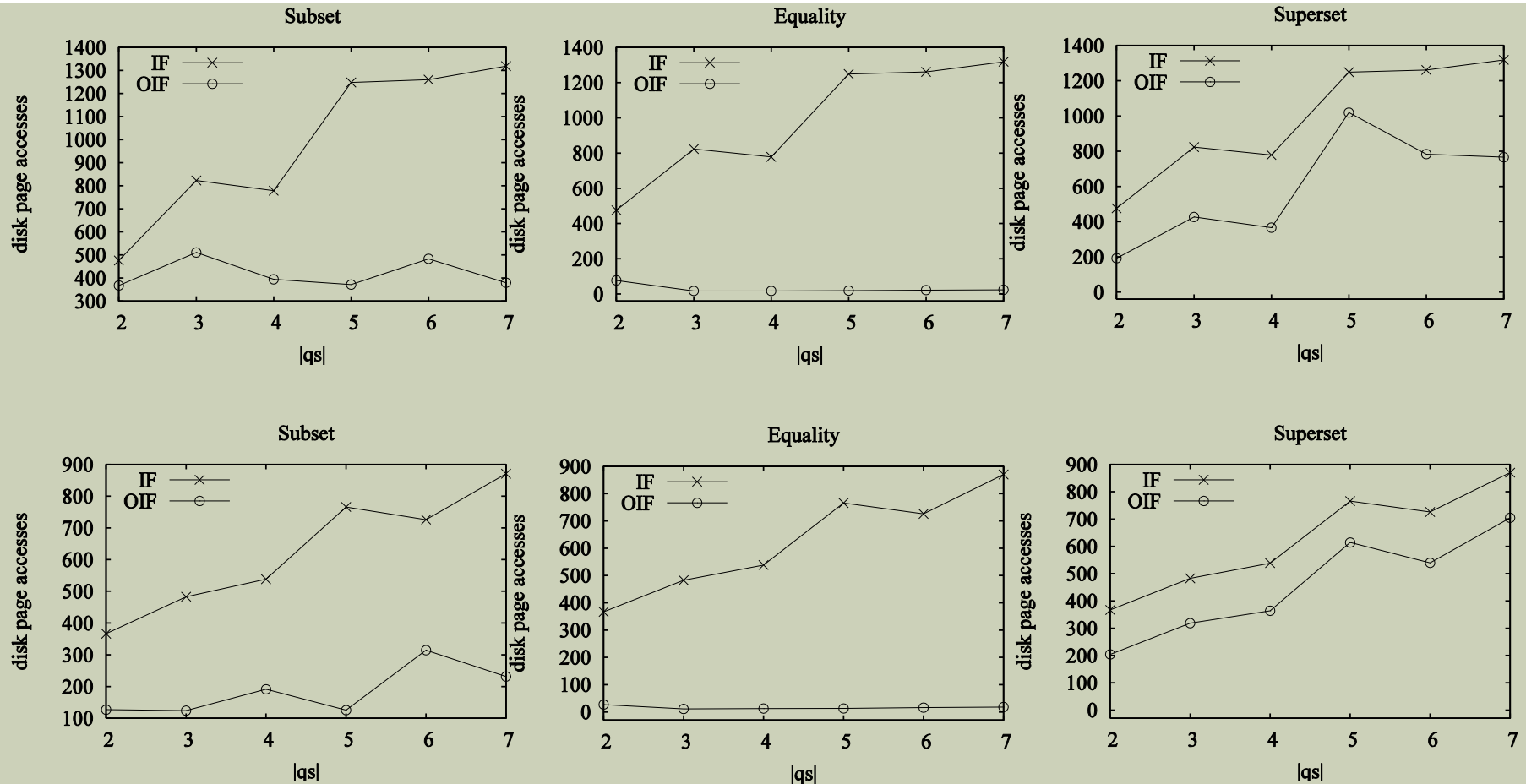
# UPDATES

- Updates are usually insertions
- Strategies for updating IF are basically bulk update strategies
  - Incoming postings are kept in a in-memory smaller index
  - Queries have to be evaluated using both indices
  - When buffer becomes full indices are merged
  - ..or merge follows the retrievals of lists in queries
- OIF can be updated in the same way
  - Incoming postings are kept in a in-memory smaller index
  - Before merging the new indices the records have to be sorted
  - Changes in frequencies of items can be ignored if they are small or estimated by sampling

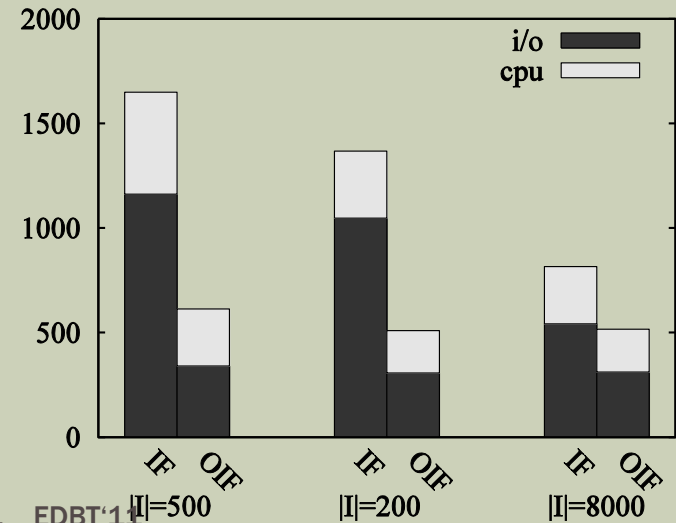
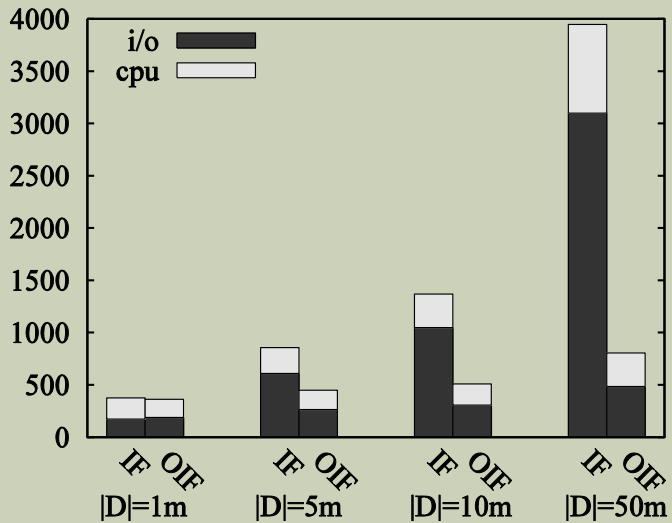
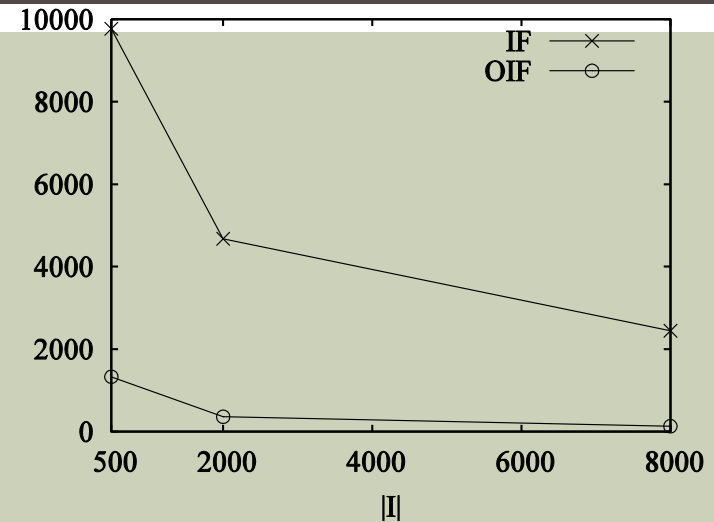
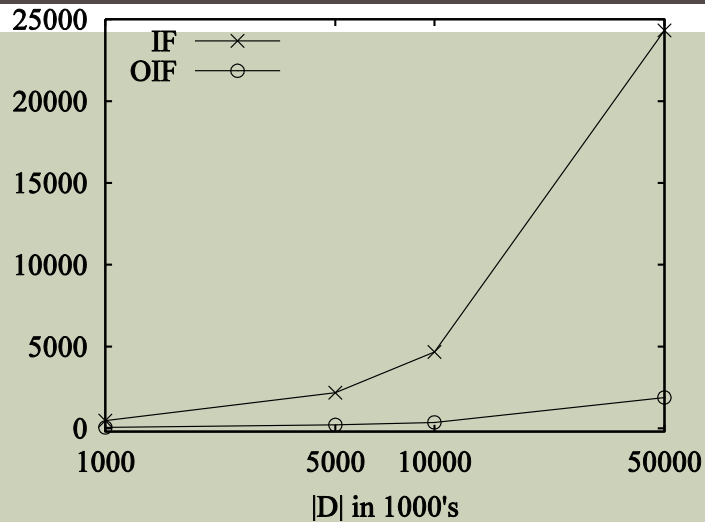
# EXPERIMENTS

- BerkeleyDB used as storage engine both for OIF and IF
- IO was traced in terms of cache misses
- We traced real execution time and CPU time
- All queries that had at least one answer by using existing records
- Datasets:
  - ms-web: 320k records, 297 items (web log)
  - Ms-nbc: 990k records, 17 items
  - Synthetic datasets default values: 10M records, 5000 items, zipf order = 0.8

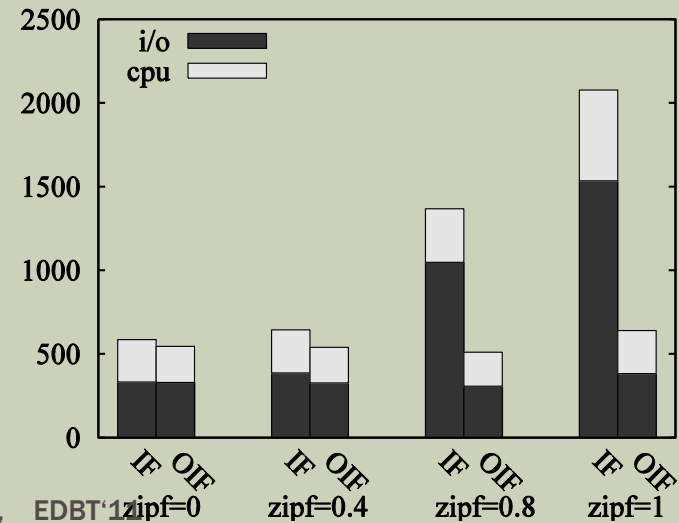
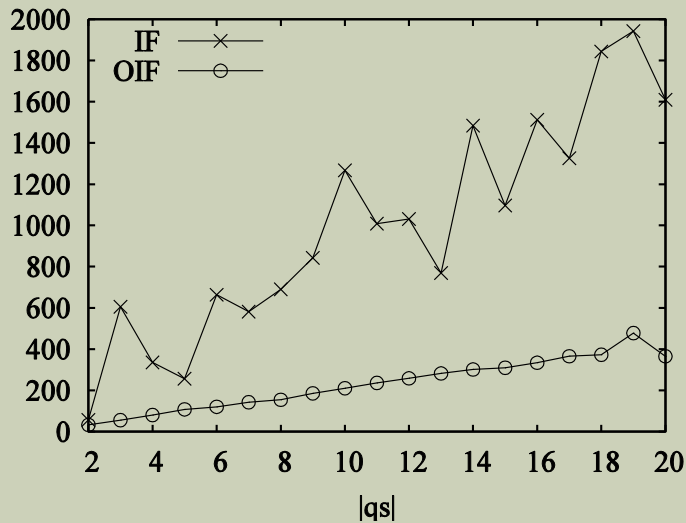
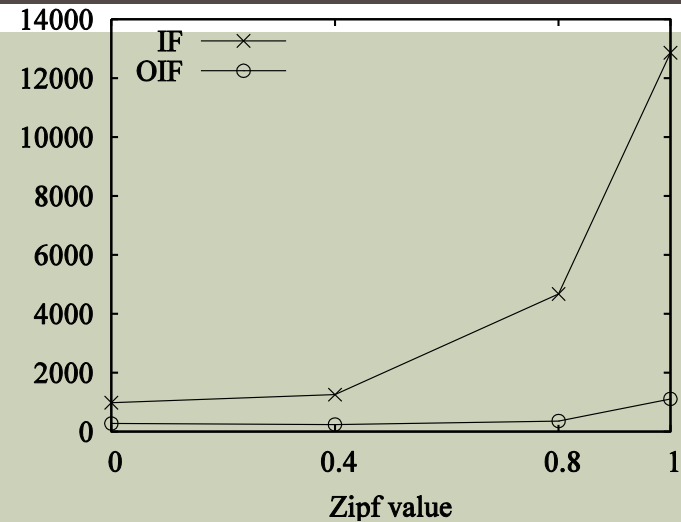
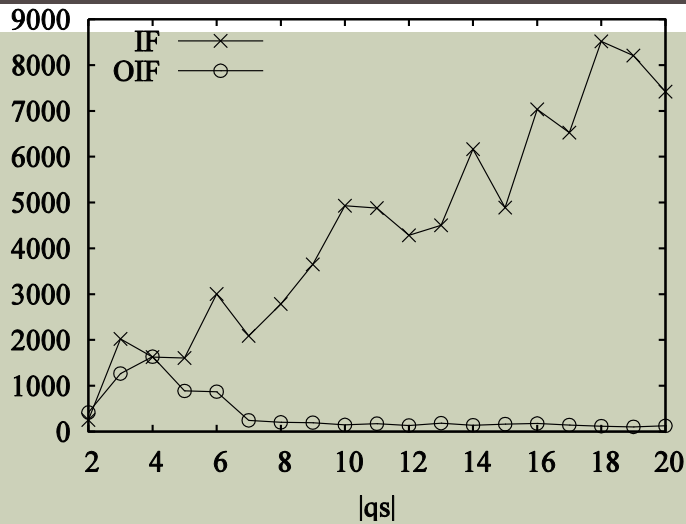
# EXPERIMENTS ON REAL WEB LOGS: MS-WEB, MS-NBC



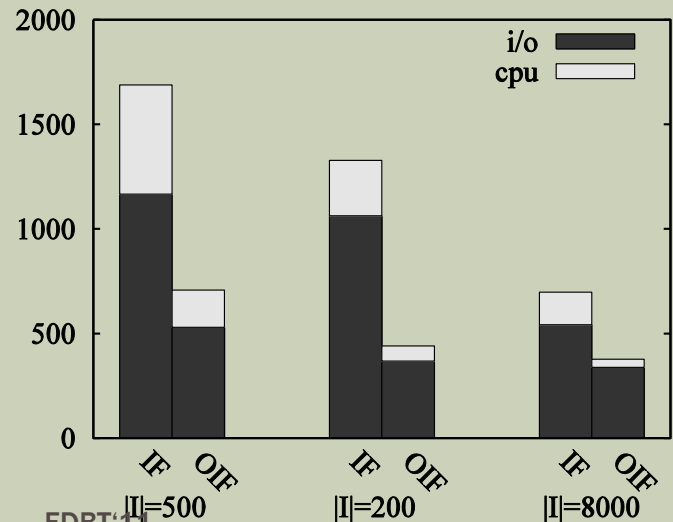
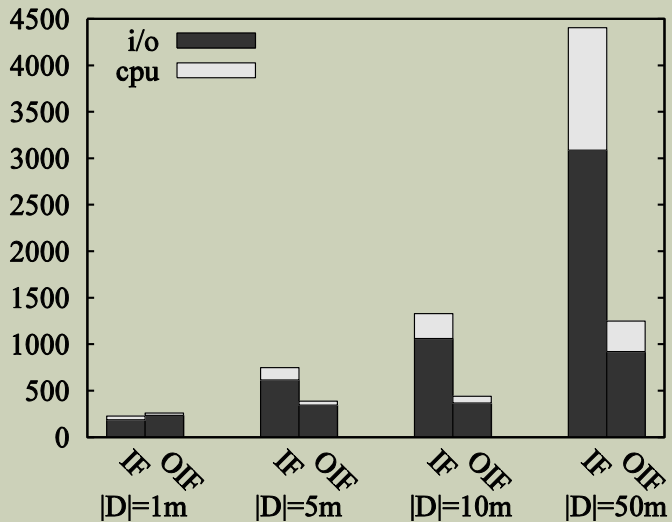
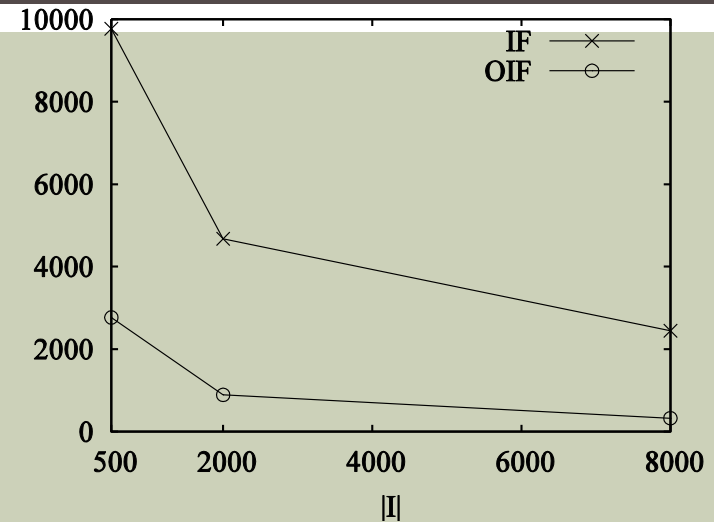
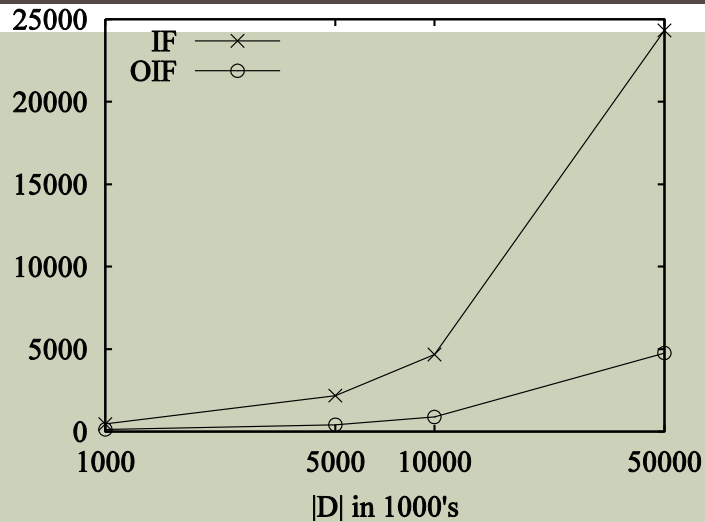
# EXPERIMENTS ON SYNTHETIC DATASETS (SUBSET)



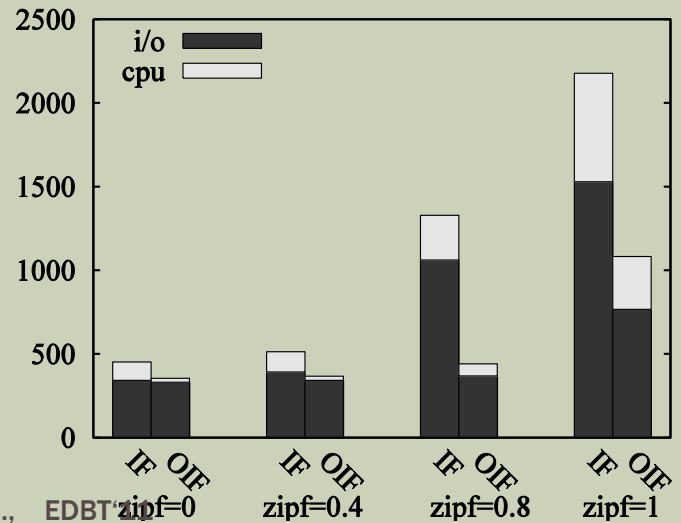
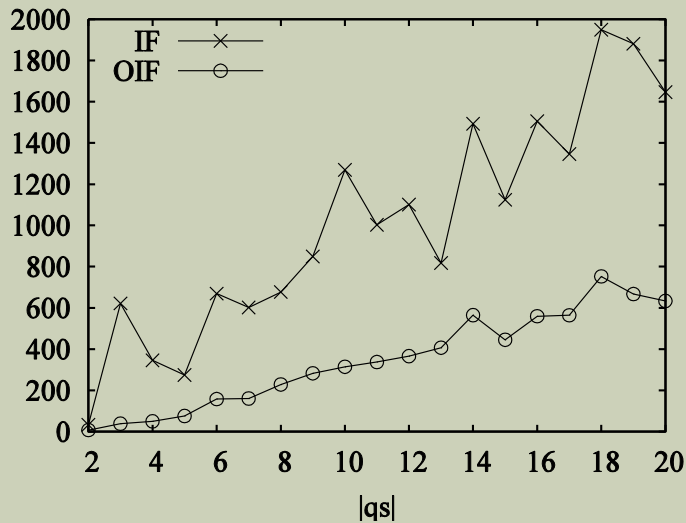
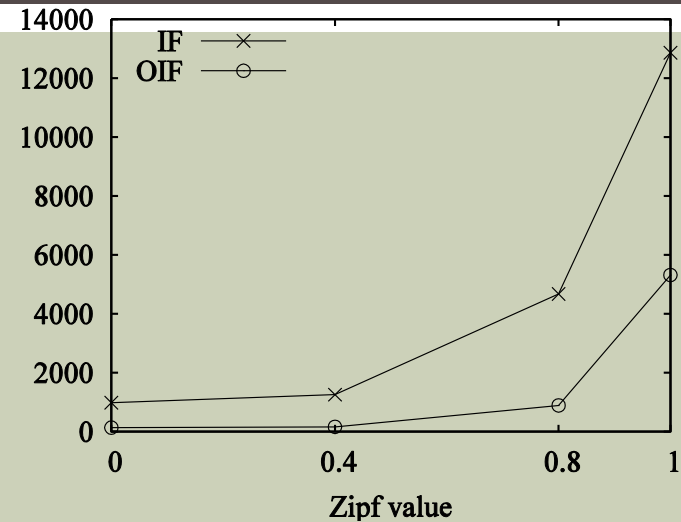
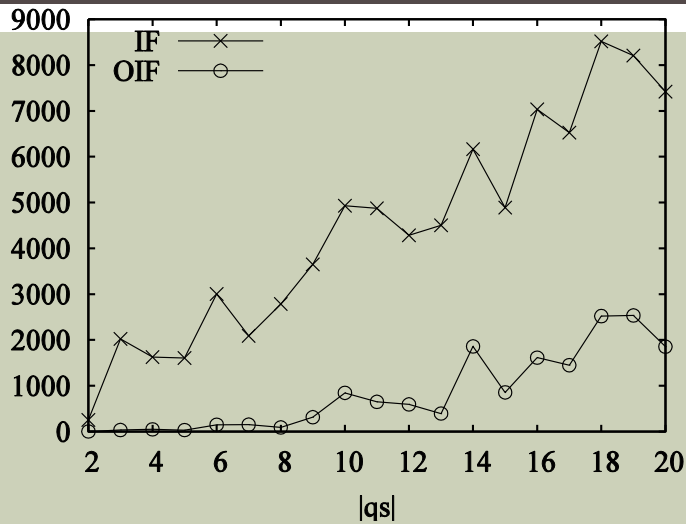
# EXPERIMENTS ON SYNTHETIC DATASETS (SUBSET)



# EXPERIMENTS ON SYNTHETIC DATASETS (SUPERSET)



# EXPERIMENTS ON SYNTHETIC DATASETS (SUPERSET)



# PERFORMANCE SUMMARY

- When the ratio of database size/domain grows or when the data are skewed OIF has substantially superior performance w.r.t IF
- OIF has increased space requirements w.r.t IF
- Updates are more expensive
  - IF has an advantage when updates dominate the workload



# QUESTIONS?

- Thank you!