## EFFICIENT NEAREST NEIGHBOR QUERIES ON NON-POINT DATA

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## Motivation **Two-Layer Partitioning** ☐ Nearest Neighbor search -> fundamental operation ☐ An SOP first-layer, e.g., a grid ☐ Decompose every first-layer partition into classes [1] ☐ GIS, data analytics, scientific applications etc. $\Box$ Find the 3 nearest cities to a query point q☐ Use the object *projections* to the axes Inside Starts In Ends In Covers 4 cases per axis ☐ Overall, *16 classes* represented by a *4-bit* encoding Tile T-1 Tile T Tile T+1 Starts In Inside: The line starts and ends inside the same tile T. Starts In: The line starts inside a tile Ends In T and ends after this. 1001 1010 Ends In: The line ends inside a tile T and starts before this. Covers Covers: The line covers the whole 1100 1101 1110 tile T. **Query Processing** Tile examination [2] Tile categorization *Incremental* search [4] *k-NN* search ☐ *Splitting* the grid into conceptual *levels* 9 types representing a set of classes each ■ Min-heap Q $\square$ k is given ☐ Process types based on their *direction* to q Level 0: the tile containing query q ☐ Storing tile-groups, ☐ *Min-heap* Q (tile- $\square \forall Level n \in [1, N]: 8n \text{ tiles } (4 \text{ tile-groups})$ ☐ *Duplicate* avoidance, instead of elimination [3] tiles and objects groups and tiles) $\square$ Examining groups by distance to q☐ *Max-heap H* (objects) ☐ Reducing comparisons Type 6 Type 7 Type 8 Sets of Classes Tile Types Level 2 Type 0 {0000,0010,1000,1010} Level 3 Type 1 (1) [{0000,0010,0100,0110,1000,1010,1100,1110}] Level 4 Type 2 {0000,0010,0100,0110} **■** 0\* Type 3 Type 4 Type 5 . . . Type 3 $\Diamond$ [{0000,0001,0010,0011,1000,1001,1010,1011}] Type 4 O ALL CLASSES {0000,0001,0010,0011,0100,0101,0110,0111} Type 0 Type 1 Type 2 Type 6 $\diamondsuit$ {0000,0001,1000,1001} Type 7 [ [ {0000,0001,0100,0101,1000,1001,1100,1101 Type 8 {0000,0001,0100,0101} Tile-group ← **Experiments ROADS EDGES ROADS EDGES** Setup **Incremental search Tuning grid granularity** ☐ All data in main memory ☐ 2-layer16 using a regular grid as SOP Throughput [quer. 10<sup>2</sup> ☐ Focus on filtering phase ☐ 2 Tiger 2015 datasets 2-layer16 —— 2-layer16 —— R-tree **ROADS** 2 3 4 5 6 7 8 9 10 2 3 4 5 6 7 8 9 10 number of browsed neighbors number of browsed neighbors partitions per dimension [x1000] partitions per dimension [x1000] ☐ Containing 19 millions *linestrings* **Indexing cost** K-NN search ☐ Occupying *538* MB 2-layer16 —— **EDGES** 2-layer16 —— 2-layer4 [1] 2-layer4 [1] ——— ☐ Containing 69 millions polygons 2-layer16 — 2-layer16 —— R-tree — R-tree ☐ Occupying 1.6 GB 4 5 6 7 8 9 10 50 100 500 1000 50 100 500 1000 partitions per dimension [x1000] partitions per dimension [x1000] number of desired neighbors k number of desired neighbors k

## References

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