Replay Attack in TLS 1.3 0-RTT Handshake: Countermeasure Techniques

Network Security (933II)

M.Sc. Cybersecurity

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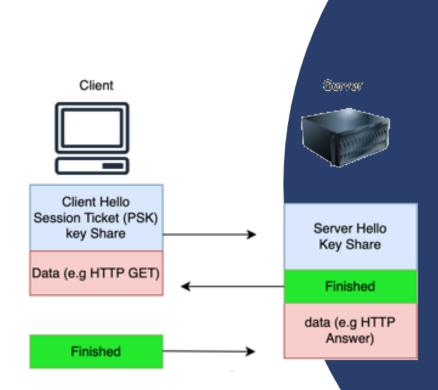
The Paper

- Conference: 2023 IEEE 6th International Conference on Electrical, Electronics and System Engineering (ICEESE)
- Authors: M.E Abdelhafez (Malaysia), Sureswaran Ramadass (Malaysia), Mohammed S. M. Gismallab (Saudi Arabia)
- Goal: review anti-replay protection techniques
- Keywords: TLS 1.3, replay attack, 0-RTT, handshake



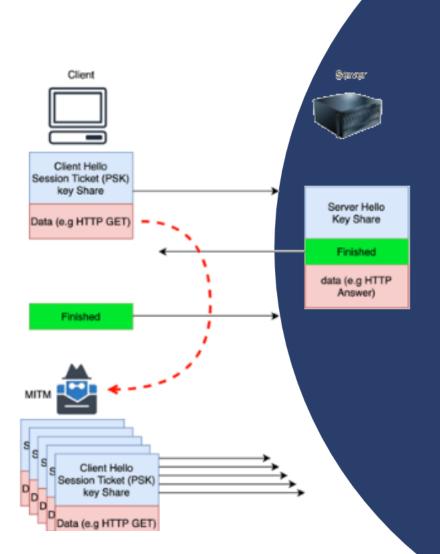
Context

- TLS resumable connections
- TLS 1.3 introduced 0-RTT resume mode, based on a Session-Ticket key created during the initial full handshake
- 0-RTT obtained by sending a single message that contains both the ClientHello (with a known Session-Ticket key) and Application
 Data (also known as Early Data)



Attack Scenarios

- Replay attack
- Attacker intercepts and replays ClientHello messages with Early Data
- The replayed message is valid because the ClientHello contains a Session-Ticket key recognized by the server
- ALTERNATIVE SCENARIO: the attacker
 performs a MITM and makes the client to
 believe that the 0-RTT message wasn't
 received, triggering a resending



Freshness check

Reject ClientHello messages whose gmt_unix_time too much in the past

- PROS: simple implementation
- CONS: can be inconvenient and there is an exploitable time window for attackers

ClientHello Recording

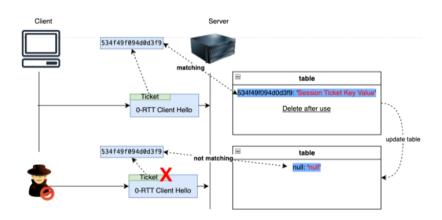
The server keeps a list of received **ClientHello** messages and uses it to detect and discard replays

- PROS: can block all replay attacks
- CONS: complex setup in distributed environments

Single-Use Tickets

The server **deletes** the "session ticket key" used to decrypt the early data after the first 0-RTT resume, making it impossible to decrypt replayed messages.

- PROS: prevents 100% of replay attacks
- CONS: requires synchronization between nodes in distributed environments



Application Profile

Each application should implement a specific **profile** that specifies under which conditions it will use 0-RTT (e.g. HTTP GET).

- PROS: flexibility
- CONS: not 100% safe, requires intervention at application level

Separate API

Both client and servers use libraries that make 0-RTT usage **explicit**, rather than implicit and automatic.

- PROS: explicit behaviour
- CONS: requires TLS libs restructuring and programmers attention

Puncture Pseudorandom Function (PPRF)

By using PPRF the server can decrypt 0-RTT early data only once.

Example approach: a server maintains a session ticket encryption key (STEK) k that can decrypt any session ticket. Then it uses it to decrypt a ticket t and it generates a STEK k' that can decrypt all session tickets but t and so on...

- PROS: forward secrecy
- CONS: long processing time, cannot be used in distributed environments

Universal SSL

Introduced by **Cloudflare** in 2015 (doesn't support TLS 1.3), Universal SSL stores negotiated sessions into multiple **Memcached** instances. Each session is indexed and encrypted by **Session ID**.

- PROS: great performance
- CONS: Memcached servers are synchronized only within each Cloudflare PoP



Just-in-Time Shared Keys (JIT-SK)

Based on a **synchronized PRNG**, dynamically changes keys for each session to secure 0-RTT messages (the same key cannot be reused multiple times, so "blind replaying" is impossible)

- PROS: prevents replay attacks and provides forward secrecy
- CONS: doesn't support distributed environments

Conclusions

- **0-RTT is here to stay**: the performance improvements are real (the paper stats that 0-RTT resume is 44.7% than 1-RTT) and the percentage of resumed TLS connections is also quite high (40% ins some applications)
- **0-RTT anti-reply protection requires trade offs:** the evaluated protections introduce overheads and/or inconveniences, especially in distributed environments (e.g. CDNs), therefore 0-RTT replay protection is still an open research topic

