Math XXXX – Independent Study: Manifolds, Category Theory– Summer 2025

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Chapter 1

Categories, Functors, and Natural Transformations

1.1 Categories

Thinking of a category as atomic objects with some kind of internal structure, A, B. Two such structures are 'morphed' into another object with a separate structure, ϕ. A special operation called 'composition' (defined in special cases) such that all morphisms can be composed to make more morphisms of the same object type. All internal structures to objects A, B are invariant under this morphism.

$$\begin{tikzcd} A \arrow[dr] & \\ & \phi \\ B \arrow[ur] & \end{tikzcd}$$

There exists an identity morphism, as well, for each object $Id_A : A \to A$, thus $Id_A \circ \phi = \phi$ and $\phi \circ Id_B = \phi$. Note: it is wrong to write $\phi : A \to B$. It is more accurate to write $\phi : A \times B \to C$ where C is an object that can be composed. That is to say ϕ, θ composed together to make a third object of type C, $\psi = \phi \circ \theta$ with some other initiating objects E, F thus we similarly define $\psi : E \times F \to \mathbb{C}$ where E, F are objects in the same sense as A, B.

Example 1.1.1 (Many Object-One Morphism: Matrix Category over field k). Given a field k. The objects are natural numbers, n, m. The morphism is $\mathbf{Mat} : \mathbb{N} \times \mathbb{N} \to \mathcal{M}_{n\times m}(k)$. Composition is defined as matrix multiplication. Thus, $A \in \mathbf{Mat}(l, m)$ and $B \in \mathbf{Mat}(m, n)$ can be composed to make $AB \in \mathbf{Mat}(l, n)$. Keeping in mind that l, m, n are the objects and A, B, AB are the morphisms. The left identity for B is the identity matrix $I \in \mathbf{M}(m, m)$.

Things to notice:

- Multiple objects that act more as defining characteristics (in this case, the dimensions of the matrices).
- a single morphism which generates a different object (in thise case, not a number but a set of matrices).
- Composition is defined by the generated object and there is more than one identity (In this case identity matrices of specific degree).

Example 1.1.2 (One object-many morpisms: Group Category). The object is characterized by $*$. Each group element is characterized by a morphism. That is, $g : * \to *$ and we define the composition between group elements to represent the group. For example, $\mathbb{Z}/\mathbb{Z}4$ has elements $\{0, 1, 2, 3\}$. Each of these elements is a morphism, that is $1 : * \to *$ but the composition is defined explicitly in a table

$\circ$	0	1	2	3
0	0	1	2	3
1	1	2	3	0
2	2	3	0	1
3	3	0	1	2

or intuitively as addition mod 4.

Things to notice:

- A single object that acts more like a place holder.
- Morphisms that are identical except for how they interact in the composition function.
- A composition function which identifies the behavior. In this case, there is only one identity, namely, the 0 morphism.

Example 1.1.3 (Group elements as objects, $\mathbb{Z}/5\mathbb{Z}$).

- Objects: Let the objects be the natural numbers $\{0, 1, 2, 3, 4\}$.
- Morphisms: The set $\text{Hom}(a, b) = b - a \mod 5$.
- Composition: $f, g \in \text{Hom}(a, b)$ then $f : a \to b$ and $g : b \to c$ then $f \circ g = f + g = (b-a)+(b-c) = a - c$, hence $f \circ g : a \to c$.
- Identity: 0, of course.
- Associaitivty: follows.

Note: to me this is a more intuitive understanding of Group as a Category. The clear point here is that this example and the one show that the group can be represented as a EITHER a single-element element category (the element being a placeholder) with the structure defined through the morphisms and composition OR a multi-object category with the structure defined by the underlying structure of the object (i.e., in this case the objects can add together to make another object).

1.1.1 Exercises

1.1.12 Find three examples of categories not mentioned above.

Groups, Rings, Fields, Vector Spaces, Modules, Representations, Manifolds, Topological Spaces

1.1.13 Show that a map in a category can have at most one inverse. That is, given a map $f : A \to B$ there is at most one map $g : B \to A$ such that $gf = \mathbb{I}_A$ and $fg = \mathbb{I}_B$.

1.1.14 Let $\mathscr{A}$ and $\mathscr{B}$ be categories. Construct 1.1.11 defined by the product category $\mathscr{A} \times \mathscr{B}$, except that the definitions of composition and identities in $\mathscr{A} \times \mathscr{B}$ are not given. There is only one sensible way to define the: write it down.

1.1.15 There is a category call **Toph** whose objects are topological spaces and whose maps $X \to Y$ are homotopy classes of continuous maps X to Y. What do we need to know about homotopy in order to prove that **Toph** is a category? What does it mean in pure topological terms for two objects of **Toph** to be isomorphic?

1.2 Functors

1.2.1 Exercises

1.2.20 Find three examples of functors not mentioned above.

1.2.21 Show that functors preserve isomorphism. That is, prove that if $F : \mathscr{A} \to \mathscr{B}$ is a functor and $A, A' \in \mathscr{A}$ with $A \cong A'$, then $F(A) \cong F(A')$

Let $f : A \to A'$ be an isomorphism of the category $\mathscr{A}$.

$$\begin{aligned} F(\mathbb{I}) &= F(f \circ f^{-1}) = F(f) \circ F(f^{-1}) \\ F(f) &: F(A) \to F(A') \\ F(f(ka+b)) &= F(f(ka) + f(b)) = F(kf(a)) + F(f(b)) = k(F(f(a)) + F(f(b)) \\ &\text{Let } a \in \ker(f) \\ F(f(a)) &= F(0) \implies F(\ker(f)) = \ker F(f) \end{aligned}$$

Thus $F(f)$ is an isomorphism.

1.2.22 Prove the assertion made in Example 1.2.9. In other words, give ordered sets A and B, and denoted by $\mathscr{A}$ and $\mathscr{B}$ the corresponding categories, show that a functor $F : \mathscr{A} \to \mathscr{B}$ amounts to an order-preserving map $A \to B$.

That is, given $a \geq a'$ prove that $F(a) \geq F(a')$ for any functor $F : A \to B$.

$$\begin{aligned} g(a) &= \{b \in A : b \geq a\} \\ a \geq b &\implies g(a) \supseteq g(b) \\ \text{NTS } F(g(a)) &\supseteq F(g(b)) \\ x \in g(b) &\implies F(x) \in F(g(b)) \text{ and} \\ x \in g(a) &\implies F(x) \in F(g(a)) \\ \therefore F(g(b)) &\subseteq F(g(a)) \end{aligned}$$

1.2.23 Two categories $\mathscr{A}$ and $\mathscr{B}$ are **isomorphic**, wrtitten as $\mathscr{A} \cong \mathscr{B}$, if they are isomorphic as objects in **CAT**.

(a) Let G be a group, regarded as a one-object category all of whose maps are isomorphisms. then its opposite G^{op} is also a one-obejct category all of whose maps are isomorphisms, and can therefore be regarded as a group too. What is G^{op}, in purely group-theoretic terms? Prove that G is isomorphic to G^{op}.

The ONLY difference betwee G and G^{op} is that all of the morphisms (isomorphisms) originate from the left instead of the right. But, such homomorphisms are commutative, thus G and G^{op} are essentially the same thing.

(b) Find a monoid not isomorphic to its opposite.

Find a non-commutative group. Let $C \subset S_3$, $C = \left\{ \begin{pmatrix} 1 & 2 & 3 \end{pmatrix}, \begin{pmatrix} 2 & 3 & 1 \end{pmatrix}, \begin{pmatrix} 3 & 1 & 2 \end{pmatrix} \right\}$

1.2.24 Is there a functor Z : **Grp**$\to$ **Grp** with property that $Z(G)$ is the center of G for all groups G?

Consider $\mathbb{Z}/\mathbb{Z}12, \mathbb{Z}/\mathbb{Z}7 \in$ **Grp**. Then $Z : \{0, 2, 3, 4, 6\} \to \{0\}$. We can see that $Z(2+6) = Z(2) + Z(6) = 0 \neq Z(8)$ which does not exist.

1.2.25 Sometimes we meet functors whose domain is a product $\mathscr{A} \times \mathscr{B}$ of categories. Here you will show that such a functor can be regarded as an inter-locking pair of families of functors, one defined on $\mathscr{A}$ and the other defined on $\mathscr{B}$. (This is very like the situation for bilinear and linear maps.)

(a) Let $F : \mathscr{A} \times \mathscr{B} \to \mathscr{C}$ be a functor. Prove that for each $A \in \mathscr{A}$, there is a functor $F^A : \mathscr{B} \to \mathscr{C}$ defined on objects $B \in \mathscr{B}$ by $F^A(B) = F(A, B)$ and on maps g in $\mathscr{B}$ by $F^A(g) = F(1_A, g)$. Prove that for each $B \in \mathscr{B}$, there is a functor $F_B : \mathscr{A} \to \mathscr{C}$ defined similarly.

(b) Let $F : \mathscr{A} \times \mathscr{B} \to \mathscr{C}$ be a functor. With notation in (a), show that the families of functors $(F^A)_{A\in\mathscr{A}}$ and $(F^B)_{B\in\mathscr{B}}$ statisfy the following two conditions:

- if $A \in \mathscr{A}$ and $B \in \mathscr{B}$ then $F_A(B) - F_B(A)$;
- if $f : A \to A'$ in $\mathscr{A}$ and $g : B \to B'$ in $\mathscr{B}$ then $F^{A'}(g) \circ F_B(f) = F_{B'}(f) \circ F^A(g)$.

(c) Now take categories $\mathscr{A}, \mathscr{B}$ and $\mathscr{C}$, and take familes of functors $(F^A)_{A\in\mathscr{A}}$ and $(F^B)_{B\in\mathscr{B}}$ satisfying the two conditions in (b). Prove that there is a unique functor $F : \mathscr{A} \times \mathscr{B} \to \mathscr{C}$ satisfying the equation in (a). ('There is a unique functor' means in particualr that there *is* a functor, so you have to prove existence as well as uniqueness.)

1.2.26 Fill in the details of Eample 1.2.11, thus constructing a functor $C : \mathbf{Top}^{\mathrm{op}} \to \mathbf{Ring}$.

1.2.27 Find an example of a functor $F : \mathscr{A} \to \mathscr{B}$ such that F is faithful but there exists distinct maps f_1 and f_2 with $F(f_1) = F(f_2)$.

1.2.28 (a) Of the examples of functors appearing in this section, which are faithful and which are full?

(b) Write down one example of a functor that is both full and faithful, one that is full but not faithful, one that is faithful but not full, and one that is neither.

1.2.29 (a) What are the subcategories of an ordered set? which are full?

(b) What are the subcategories of a group? (careful!) Which are full?

1.3 Natural Transformations

Natural Transformations are morphisms between functors – just as functors are morphisms between categories, and morphisms exist between objects in a categories ... the next level

	Element	**Connection**	**Operation**
Category	Object	Morphism	Composition
Functor-Level	Category	Functor	Composition
Natural Transformation	Functor-Family	Natural Transformation	Composition

Transformation Diagrams

$$\alpha : F \to G \equiv \Big(F(A) \xrightarrow{\alpha_A} G(A) \Big)_{A\in\mathscr{A}} \in (B)$$

$$\begin{tikzcd}
F(A) \arrow[r, "F(f)"] \arrow[d, "\alpha_A"'] & F(A') \arrow[d, "\alpha'_A"] \\
G(A) \arrow[r, "G(f)"'] & G(A')
\end{tikzcd}$$

Defining α

$$\begin{tikzcd}
\mathscr{A} \arrow[r, bend left, "F"] \arrow[r, bend right, "G"'] \arrow[r, phantom, "\updownarrow\alpha"] & \mathscr{B}
\end{tikzcd}$$

Defining $\beta \circ \alpha$

$$\begin{tikzcd} \mathscr{A} \arrow[r, bend left=60, "F"] \arrow[r, "G" description] \arrow[r, bend right=60, "H"'] \arrow[r, phantom, shift left=3, "\Downarrow\alpha"] \arrow[r, phantom, shift right=3, "\Downarrow\beta"] & \mathscr{B} \end{tikzcd} = \begin{tikzcd} \mathscr{A} \arrow[r, bend left=50, "F"] \arrow[r, bend right=50, "H"'] \arrow[r, phantom, "\Downarrow\beta\circ\alpha"] & \mathscr{B} \end{tikzcd}$$

1.3.1 Exercises

1.3.25 Find three examples of Natural Transformations not mentioned above.

(a) A group G and its opposite G^{op}. Both span the same set, but the left operations of G are the right operations of G^{op}. The function $F : G \to Gop$ is the identity for the sets $F(A) = A$ along with the reverse homomorphism $F \circ f : G^{\text{op}} \to G$. That is, where $f(gh) = f(g)f(h)$, $F(f)(gh) = F(f(h))F(f(g))$.

1.3.26 Prove Lemma 1.3.11.

1.3.27 Let $\mathscr{A}$ and $\mathscr{B}$ be categories. Prove that $[\mathscr{A}^{\text{op}}, \mathscr{B}^{\text{op}}] \cong [\mathscr{A}, \mathscr{B}]^{\text{op}}$

1.3.28 Let A and B be sets and denote B^A the set of all functions from $A \to B$. Write down:

(a) A canonical function $A \times B^A \to B$.
Let $\phi : A \times B^A \to B$ then $\phi(a, f(a)) = f(a)$.

(b) A canonical function $A \to B^{(B^A)}$.
(Although in principle there could be many such canonical function , in both cases there is only one).
Let $\theta \in B^{(B^A)}$ which means $\theta : B^A \to B$. Thus, $\theta(f(a)) \in B$ for some $f : A \to B$. Let $\phi : A \to B^{(B^A)}$ then $\phi(a) = \theta(f(a))$ for some θ and some f, which must be in B.

1.3.29 Here we consider natural transformations between functors whoe domain is a product of categories $\mathscr{A} \times \mathscr{B}$. Your task is to show that naturality in two variables simultaneously is equivalent to naturality in each variable separately.

Take functors $F, G : \mathscr{A} \times \mathscr{B} \to \mathscr{C}$. For each $A \in \mathscr{A}$, there are functors $F^A, G^A\ \mathscr{B} \to \mathscr{C}$, as in Exercise 1.2.25. Similarly, for each $B \in \mathscr{B}$, there are functors $F_B, G_B : \mathscr{A} \to \mathscr{C}$.

Let $(\alpha_{A,B} : F(A,B) \to G(A,B))_{A \in \mathscr{A}, B \in \mathscr{B}}$ be a family of maps. Show that this family is a natural transformation $F \to G$ if and only if it satisfies the following two conditions

- for each $A \in \mathscr{A}$, the family $(\alpha_{A,B} : F^A(B) \to G^A(B))_{B \in \mathscr{B}}$ is a natural transformation $F^A \to G^A$;
- for each $B \in \mathscr{B}$, the family $(\alpha_{A,B} : F_B(A) \to G_B(A))_{A \in \mathscr{A}}$ is a natural transformation $F_B \to G_B$.

1.3.30 Let G be a group. For each $g \in G$, there is a unique homomorphism $\phi : \mathbb{Z} \to G$ satisfying $\phi(1) = g$. Thus, elements of G are essentially the same thing as homomorphisms $Z \to G$. When groups are regarded as one-object categories, homomorphisms $Z \to G$ are in turn the same a functors $\mathbb{Z} \to G$. natural isomorphism defines an equivalence relation on the set of functors $\mathbb{Z} \to G$, and, therefore, an equivalence relation on G itself. What is this equivalence relation, in purely group-theoretic terms?

(First have a guess. For a general group G, what equivalence relations on G can you think of?)

1.3.31 A **permutation** of a set X is a bijection $X \to X$. Write $\mathbf{Sym}(X)$ for the set of permutations of X. A **total order** on a set X is an odrer $\leq$ such that for all $x, y \in X$, either $x \leq y$ or $y \leq x$; so a total order on a finite set amounts to a way of placing its elements in sequence. Write $\mathbf{Ord}(X)$ for the set of total orders on X.

Let $\mathscr{B}$ deonte the category of finite sets and bijections.

(a) Give a definition of $\mathbf{Sym}$ on maps in $\mathscr{B}$ in such a way that $\mathbf{Sym}$ becomes a functor $\mathscr{B} \to \mathbf{Set}$. Do the same for $\mathbf{Ord}$. Both your definitions should be canonical (no arbitrary choices).

(b) Show that there is no natural transformation $\mathbf{Sym} \to \mathbf{Ord}$. (Hint: consider identity permutations.)

(c) For n-element set X, how many elements of the sets $\mathbf{Sym}(X)$ and $\mathbf{Ord}$ have?

Conclude that $\mathbf{Sym}(X) \cong \mathbf{Ord}(X)$ for all $X \in \mathscr{B}$, but not *naturally* in $X \in \mathscr{B}$. (The moral is that each finite set X, there are exactly as many permutations of X as there are total orders on X, but there is no natural way of matching them up.)

1.3.32 In this exercise, you will prove Proposition 1.3.18. Let $F : \mathscr{A} \to \mathscr{B}$ be a functor.

(a) Suppose that F is an equivalence. Prove that F is full, faithful and essentially surjective on objects. (Hint: prove faithfulness before fullness.)

(b) Now suppose instead that F is full, faithful and essentially surjective on objects. For each $B \in \mathscr{B}$, choose an object $G(B)$ of $\mathscr{A}$ and an isomorphism $\epsilon_B : F(G(B)) \to B$. Prove that G extends to a functor in such a way that $(\epsilon_B)_{B \in \mathscr{B}}$ is a natural isomorphism $FG \to 1_{\mathscr{B}}$. Then construct a natural isomorphism $1_{\mathscr{A}} \to GF$, thus proving that F is an equivalence.

1.3.33 This exercise makes precise the idea that linear algebra can equivalently be done with matrices or with linear maps.

Fix a field k. Let $\mathbf{Mat}$ be the category whose objects are the natural numbers and with

$$\mathbf{Mat}(m, n) = \{ n \times m : \text{ matrices over } k \} .$$

Prove that $\mathbf{Mat}$ is equivalent to $\mathbf{FDVect}$, the category of finite-dimensional vector spaces over k. does you equivalence involve *a canonical* functor from $\mathbf{MAT}$ to $\mathbf{FDVect}$, or from $\mathbf{FDVect}$ to $\mathbf{Mat}$?

(Part of the exercise is to work out what composition in the cagtegory $\mathbf{Mat}$ is supposed to be; there is only on sensible possibility. Propostition 1.3.18 makes the exercise easier.)

1.3.34 Show that equivalence of categories is an equivalence of relation. (Not as obvious as it looks).

Chapter 2

Adjoints

2.1 Definitions and Examples

Definition 2.1.1 (Bar Notation)**.** Given an ***adjunction*** between F and G (i.e., a natural isomorphism) we define a ***transpose*** of morphism, that is, we call "$\bar{f}$ the transpose of f, and simlilarly for g, as in

$$\left(F(A) \xrightarrow{g} B\right) \mapsto \left(A \xrightarrow{\bar{g}} G(B)\right),$$
$$\left(F(A) \xrightarrow{\bar{f}} B\right) \leftarrow \left(A \xrightarrow{f} G(B)\right)$$

i.e., $g : B \to B \implies \bar{g} : A \to A$

Definition 2.1.2 (The Naturality Axioms)**.** The ***naturality axioms*** have two parts. Given an adjuction between F and G, that is

$$\mathscr{B}(F(A), B) \cong \mathscr{A}(A, G(B)) \tag{2.1}$$

that is being "naturally" isomorphic. That is given $A \in \mathscr{A}$ and $B \in \mathscr{B}$ between maps $F(A) \to B$ and $G(B) \to A$ denoted by a horizontal bar in both directions:

$$\overline{\left(F(A) \xrightarrow{g} B \xrightarrow{q} B'\right)} = \left(A \xrightarrow{\bar{g}} G(B) \xrightarrow{G(q)} G(B')\right) \tag{2.2}$$

(that is, $\overline{q \circ g} = G(q) \circ \bar{g}$) for all g and q, and

$$\overline{\left(A' \xrightarrow{p} A \xrightarrow{f} G(B)\right)} = \left(F(A') \xrightarrow{F(p)} F(A) \xrightarrow{\bar{f}} B\right) \tag{2.3}$$

(that is, $\overline{f \circ p} = \bar{f} \circ F(p)$) for all p and f. It makes no difference whether we put the long bar over the left or the right of these equations, since bar is self-inverse.

Remark 2.1.3. Even though there is an adjuction between F and G (i.e., a natural isomorphism) the morphisms $f \in \mathscr{A}$ implies $\bar{f} \in \mathscr{B}$ and distinctly separate from $g \in \mathscr{B}$ and $\bar{g} \in \mathscr{A}$ and yet $\bar{\bar{f}} = f$. That is, the bar operation relates f through the isomorphism (2.1) (being 1-to-1 and onto) to its counterpart.

Remark 2.1.4 (From ChatGPT). There exists an adjunction between $F : \mathscr{C} \to \mathscr{D}$ and $G : \mathscr{D} \to \mathscr{C}$ if

$$\mathrm{Hom}_{\mathscr{D}}(F(X), Y) \cong \mathrm{Hom}_{\mathscr{C}}(X, G(Y))$$

for all $X \in \mathscr{C}$ and $Y \in \mathscr{D}$. That is, the set of all homomorphisms is isometric between the categories.

Intuitively speaking:

1. **Left Adjoint F:** "Frees" up structure.
2. **Right Adjoint G:** "Forgets" or "Extracts" structure.

Remark 2.1.5 (ChatGPT). **Whats the difference between an "adjuction between F and G" and "G is the inverse of F"?**

- **Inverse:** $F \circ G = \mathrm{Id}_F$ and $G \circ F = \mathrm{Id}_G$.

  Meaning: F and G define an **isomorphism of categories**. They are structure-preserving bijections at both the object and morphism levels.

- **Adjunction:** The *natural isomorphism of Hom-sets*

$$\mathrm{Hom}_{\mathscr{D}}(F(X), Y) \cong \mathrm{Hom}_{\mathscr{C}}(X, G(Y))$$

  This expresses a **universal property**.

  That is: F is a left adjoint of G if each $X \in \mathscr{C}$ AND $Y \in \mathscr{D}$ maps out of $F(X)$ correspond naturally to maps out of X to $G(Y)$.

2.1.1 Exercises

2.1.12 Find three examples of adjoint functors not mentioned above. Do the same for intial and terminal objects.

2.1.13 What can be said about adjunctions between discrete categories?

2.1.14 Show that the naturality equation (2.2) and (2.3) can equivalently be replaced by the single equation

$$\overline{\left(A' \xrightarrow{p} A \xrightarrow{f} G(B) \xrightarrow{G(q)} G(B')\right)} = \left(F(A') \xrightarrow{F(p)} F(A) \xrightarrow{\bar{f}} B \xrightarrow{q} B'\right)$$

for all p, f, and q.

2.1.15 Show that the left adjoints preserve intitial objects: that is, $\mathscr{A} \underset{\xleftarrow[G]{}}{\overset{\xrightarrow{F}}{\perp}} \mathscr{B}$ and I is the initial object of $\mathscr{A}$, then $F(I)$ is the initial object of $\mathscr{B}$. Dually show that right adjoints preserve terminal objects.

2.1.16 Let G be a group

(a) What interesting functors are there (in either direction) between **Set** and the category $[G, \mathbf{G}]$ for left G-sets? Which of those functors are adjoint to which?

(b) Similarly, what interesting functors are there between $\mathbf{Vect}_k$ and category $[G, \mathbf{Vect}_k]$ of k-linear representations of G, and what adjunction are there between those functors?

2.1.17 Fix a topological space X, and write $\mathscr{O}(S)$ for the poset of open subsets of X, ordered by inclusion. Let

$$\Delta : \mathbf{Set} \to [\,\mathscr{O}(X)^{\mathrm{op}}, \mathbf{Set}\,]$$

be the functor assigned to a set A the presheaf ΔA with constant value A. Exhibit a chain of adjoint functors

$$\Lambda \dashv \Pi \dashv \Delta \dashv \Gamma \dashv \nabla.$$

2.2 adjunction via units and counits

2.2.10 Let $A \underset{T}{\overset{B}{\rightleftarrows}} B$ be ordered-preserving maps between ordered sets. Prove *directly* that the folloiwng conditions are equivalent:

(a) for all $a \in A$ and $b \in B$.

$$f(a) \leq b \iff a \leq g(b);$$

(b) $a \leq g(f(a))$ for all $a \in A$ and $f(g(b)) \leq b$ for all $b \in B$.
(Both conditions state that $f \dashv g$; see Example 2.2.7.)

2.2.11 (a) Let $\mathscr{A} \underset{G}{\overset{F}{\perp}} \mathscr{B}$ be an adjunction with unit η and counit ε. Write $\mathbf{Fix}(GF)$ for the full subcaegory of $\mathscr{A}$ whose objects are those $A \in \mathscr{A}$ such that η_A is an isomorphsim, and duality $\mathbf{Fix}(FG) \subseteq \mathscr{B}$. Prove that the adjunction $(F, G, \eta, \varepsilon)$ resticts to an equivalence $(F', G', \eta', \varepsilon')$ between $\mathbf{Fix}(GF)$ and $\mathbf{(}FG)$.

(b) Part (a) shows taht every adjuction restricts to an equivalence between full subcategories in a canonical way. Take some examples fo adjunctions and work out what this equivalence is.

2.2.12 (a) Show that for any adjuction, the right adjoint is full and faithful if and only if the counit is an isomorphism

(b) An adjuction satisfying teh quivalent condtiions of part (a) is called **reflection**. (Campare Example 2.3(d).) Of teh examples of adjunctions given in this chapter, which are reflections?

2.2.13 (a) Let $f : K \to L$ be a map of sets, and enote by $f^* : \mathscr{P}(L) \to \mathscr{P}(K)$ the map sending a subseet S to L to its inverse image $f^{-1}S \subseteq K$. Then f^* is order-preserving with respect to the inclusion orderings on $\mathscr{P}(K)$ and $\mathscr{P}(L)$, and so can be seen as a functor. Find left adn right adjoints to f^*.

(b) Now let X and Y be sets, and write $p : X \times Y \to X$ for first projection. REgard a subset S of X as predicate $S(x)$ in one variable $x \in X$, and similarly a subset R of $X \times Y$ as predicate $R(x, y)$ in two variables. What, in terms of predicats, are the left and right adjoints to p^*? For each of the adjuctions, interpret the unit and counit as logical implications. (Hint: The left adjoint to p^* is often writen $\exists_Y$, adn the right adjoint as $\forall_Y$.)

2.2.14 Given a functor $F : \mathscr{A} \to \mathscr{B}$ and a category $\mathscr{S}$, there is a functor $F^* : [\mathscr{B} < \mathscr{S}] \to [\mathscr{A}, \mathscr{S}]$ defined on objects $Y \in [\mathscr{B}, \mathscr{S}]$ by $F^((Y) = Y \circ F$ adn on maps α by $F^*(\alpha) = \alpha F$. Show that any adjuction $\mathscr{A} \underset{G}{\overset{F}{\perp}} \mathscr{B}$ and category $\mathscr{S}$ give rise to an adjunction

$$[\mathscr{A}, \mathscr{S}] \underset{F'}{\overset{G'}{\perp}} [\mathscr{B}, \mathscr{S}].$$

2.3 Adjuctions via intial objects

2.3.1 Exercises

2.3.8 What can be said about adjunctions between groups (regarded as one-object categories)?

2.3.9 State the dual Corollay 2.3.7. How would you prove your dual statement?

2.3.10 let $(F, G, \eta, \varepsilon)$ be an equivalence of categoris, as in Definition1.3.15. Prove taht F is left adjoint to G (heeding thewaring in Rempark 2.2.8).

2.3.11 Let $\mathscr{A} \underset{F}{\overset{U}{\underset{\longleftarrow}{\overset{\longrightarrow}{\perp}}}} \mathbf{Set}$ be an adjunction. Suppose that for at least one $A \in \mathscr{A}$, the set $U(A)$ has at least two elements. Prove that for each set S, the unit map $\eta_S : S \to UF(S)$ is injective. What does this mean in the case the usual adjunction between **Grp** and **Set**?

2.3.12 Given sets A and B in **partial function** from A to B is a pair (S, f) consisint of a subset $S \subseteq A$ adn a function $S \to B$. (Think of it as like a function from A to B, but undefined at certain elements of A.) Let **Par** be teh category of sets adn partial functions.

Show that **Par** is equivalent to Set_*, the category of sets equipped with distinguished element adn functions preserving distinguished elements. Show also that $\mathbf{Set}_*$ can be described as a coslice category in a simple way.