

Problem 1. *A simple observation.* A nitwit argues that by increasing the edge-weights by any value Δ does not change the minimum spanning tree. Prove or give a counterexample.

Problem 2. Give an efficient algorithm $O((|V| + |E|) \log |V|)$ to compute a maximum spanning tree of an undirected graph. Prove the correctness and time complexity of the algorithm. (*Hint:* Prove an analog of the “Cut-Property”.)

Problem 3. Consider an undirected connected graph, all of whose edge weights are distinct. Prove that it has a unique minimum spanning tree.

Problem 4. State with proof or counterexample as applicable, whether each of the following statements is true or false. Assume that $G = (V, E)$ is a weighted, undirected, connected graph.

1. If G has some cycle with a unique heaviest edge e , then, e cannot be part of any MST.
2. Let e be any edge of minimum weight in G . Then e must be part of some MST.
3. If the lightest edge of the graph is unique, then it must be part of every MST.
4. Suppose e is an edge of some MST. Then it must be the minimum cut edge of some cut of G .
5. Prim’s algorithm works correctly when there are negative weight edges.
6. The shortest path tree computed by Dijkstra’s algorithm is always an MST.
7. If G contains a path from s to t consisting of edges all of whose weights $< r$, then every MST of G must contain a path from s to t consisting of edges, all of whose weights $< r$.

Problem 5. Consider the disjoint sets data structure. Give a sequence of (in total) m *union* and *find* operations on n elements that take $\Omega(m \log n)$ time.

Problem 6. Give a linear-time algorithm that takes as input a tree and determines whether it has a *perfect matching*, namely a set of edges that touches each node exactly once.

Problem 7. A *feedback edge set* of an undirected graph $G = (V, E)$ is a subset of edges $E' \subset E$ that intersects every cycle of the graph. Thus removing the edges in E' from the graph will render the graph acyclic. Given an undirected weighted graph $G = (V, E)$ with positive edge weights w_e , output a feedback edge set $E' \subset E$ of minimum total weight $\sum_{e \in E'} w_e$.