

First Draft

Topics : Dynamic Programming

Algorithm 1 Find the maximum profit that you can earn by doing at most one transaction

Ensure: One Based Indexing

```
1: function BUY_SELL(Arr, len)
  ▷ maxRight[i] denotes the max element to the right of the i – th element, including it
2:   maxRight[len] ← a[len]
3:   for i = len – 1 : 1 do
4:     maxRight[i] ← max(maxRight[i + 1], a[i])
  ▷ profit[i] denotes the max profit that you can earn if you can only buy at the i – th day
5:   for i = len : 1 do
6:     profit[i] ← (maxRight[i] – a[i])
7:   return Maximum element of Profit array
```

Algorithm 2 Find the length of the subarray with the maximum sum

Ensure: One Based Indexing

```
1: function KADANE(Arr, len)
  ▷ start[i] denotes the length of the maximum sum subarray which starts at the i – th index
2:   start[len] ← a[len]
3:   for i = len – 1 : 1 do
4:     start[i] ← max(a[i] + start[i + 1], a[i])
5:   return Maximum element of Start array
```

Algorithm 3 Find the length of the longest increasing subsequence

Ensure: One Based Indexing

```

1: function LIS(Arr, len)
  ▷ lis[i] denotes the length of the LIS which starts at the i – th index
2:   for i = len : 1 do
3:     lis[i] ← 1
4:   for i = len : 1 do
5:     for j = i + 1 : len do
6:       if a[j] > a[i] then
7:         lis[i] ← max(lis[i], 1 + lis[j])
8:   return Maximum element of LIS array

```

Algorithm 4 Find the minimum edit distance to convert *str*₁ to *str*₂

Ensure: Zero Based Indexing

```

1: function EDIT_DISTANCE(str1, str2)
  ▷ edit[i][j] denotes the edit distance of the first i characters of str1 and the first j characters
  of str2
2:   edit[i][0] ← i                                     ∀i                                     ▷ Delete i characters
3:   edit[0][j] ← j                                     ∀j                                     ▷ Insert j characters
4:   for i = 1 : str1.len do
5:     for j = 1 : str2.len do
6:       insert ← 1 + edit[i][j – 1]
7:       delete ← 1 + edit[i – 1][j]
8:       replace ← 1 + edit[i – 1][j – 1]
9:       match ← edit[i – 1][j – 1]
10:      if str1[i] == str2[j] then
11:        edit[i][j] ← match
12:      else
13:        edit[i][j] ← min(insert, delete, replace)
14:   return edit[str1.len][str2.len]

```

Algorithm 5 Find the length of the longest common subsequence of 2 strings

Ensure: Zero Based Indexing

```

1: function LIS( $str_1, str_2$ )
   $\triangleright lcs[i][j]$  denotes the longest common subsequence of the first  $i$  characters of  $str_1$  and the first
   $j$  characters of  $str_2$ 
2:    $lcs[i][0] \leftarrow 0$   $\forall i$   $\triangleright$  The first string is empty
3:    $lcs[0][j] \leftarrow 0$   $\forall j$   $\triangleright$  The second string is empty
4:   for  $i = 1 : str_1.len$  do
5:     for  $j = 1 : str_2.len$  do
6:        $exclude\_top \leftarrow lcs[i-1][j]$ 
7:        $exclude\_bot \leftarrow lcs[i][j-1]$ 
8:        $exclude\_both \leftarrow lcs[i-1][j-1]$ 
9:        $match \leftarrow 1 + lcs[i-1][j-1]$ 
10:      if  $str_1[i] == str_2[j]$  then
11:         $lcs[i][j] \leftarrow match$ 
12:      else
13:         $lcs[i][j] \leftarrow \max(exclude\_top, exclude\_bot, exclude\_both)$ 
14:  return  $lcs[str_1.len][str_2.len]$ 

```
