

## 21 – ARIES Database Recovery

### 1. ARIES

A. Algorithms for Recovery and Isolation Exploiting Semantics

B. Main Ideas

- i. Write-Ahead Logging
- ii. Repeating History During redo
- iii. Logging Changes during undo

### 2. Log Sequence Numbers

A. Log identifier and other information

B. Type

- i. FlushedLSN : in Memory, last LSN in log on disk
- ii. pageLSN : in each page, newest update to the page
- iii. recLSN : in each page, oldest update to the page
- iv. lastLSN : in each txn, latest log record of the txn
- v. MasterRecord : in disk, LSN of last checkpoint
- vi. prevLSN : in every log, previous LSN in txn

C. If page is flushed to disk, log should be flushed to disk first  
So flushedLSN  $\geq$  pagLSN

### 3. Normal Commit & Abort Operations

#### A. Commit

- i. Write COMMIT log
- ii. Flush log (up to COMMIT log)
- iii. Write TXN-END log

#### B. Abort

- i. Compensation Log Records :  
log records for reverse operations of some previous operations initially, same with update log, but has "nextUndo" field.
- ii. Process
  1. Write ABORT log
  2. Rollback txn in reverse order, write CLR also
  3. Write TXN-END log

### 4. Fuzzy Checkpointing

#### A. Non-Fuzzy Checkpointing

- i. Wait all Txns finishes, Halt all new-beginning txns
- ii. Flush all buffer

#### B. Pause Checkpointing

- i. Pause modifying txns  
Need Active Transaction Table and Dirty Page Table
- ii. Flush all buffer

#### C. Fuzzy Checkpointing

- i. Allow all txns running until checkpointing
- ii. Write ATT and DPT when writing CHECKPOINT-END log  
new beginning txn during checkpointing is not included in ATT at this time

## 5. Recovery Algorithm

### A. Analysis

Find txns committed or failed after last checkpoint

### B. Redo

repeat all operation

### C. Undo

remove changes of failed txn