

18 – Timestamp Ordering Concurrency Control

1. Basic Timestamp Ordering Protocol

A. Read situation

If record's write-timestamp is bigger than txn's timestamp, txn should be abort, because of reading future value.

B. Write situation

if record's read-timestamp is bigger than txn's time stamp, txn should be abort, because of modifying future value

C. Pros and cons

- i. Pros : no wait, so no deadlock
- ii. Cons : cause starvation for long txn, if short txn make conflict a lot.
recoverable schedule issue : after commit T2, if T1 aborts?

2. Optimistic Concurrency Control

A. Txn's timestamp is assigned at its validation phase, not beginning of txn.

B. Process

- i. Each txn has private workspace,
all operation of txn is executed at this workspace
- ii. if txn commits, check if workspace and DB has some conflicts.
- iii. If there are no conflicts, txn's workspace committed into DB

C. Validation

- i. Backwards : check txn's read-write sets intersects any txn that already committed
- ii. Forwards : check txn's read-write sets intersects any txn that have not committed
- iii. Ti's write set should not intersect with Tj's read set
(if Ti's timestamp is older than Tj)
- iv. Ti's read section should be completes before Tj's read section's end
and Ti's write section should not intersect with Tj's read-write set

3. Partition-based Timestamp Ordering

A. Process

- i. Partition DB into disjoint subset
- ii. If Txn tries to access and read a partition that it have a lock, allow it if not having lock, abort and restart txn.
- iii. If Txn tries to access and write a partition that it have a lock, allow it and maintain buffer to rollback if txn abort if not having lock, abort and restart txn

B. Issue

- i. Multi-partition txn can cause idle partition when execute

4. What is it?

A. Thomas Write Rule

B. Multi-partition Timestamp Ordering (interesting)

5. Introduced Papers

A. On Optimistic Methods for Concurrency Control