

06 – OLTP Database 1

1. In-Memory T-Tree

A. Based on AVL Tree

Storing pointer instead of key

B. Has many pointers, 2 value in node

i. Pointer : left/right child, parent pointer, data pointers

ii. Value : minimum range, maximum range

C. Advantage / Disadvantage

i. Advantages

1. Less memory : no keys are stored in the index
2. Can evaluate predicate and access tuple at same time

ii. Disadvantages

1. Rebalancing is difficult
2. Thread-safe implementation is difficult
3. Pointer chasing hurts cache locality

2. Latch-Free Bw-Tree

A. Cannot build Latch-Free B+Tree because split/merge operations update multiple pointers

if there are indirection layer, it can be possible (Bw-Tree)

B. Delta Chains

- i. Each update make new delta record
delta record points to base page
- ii. Replace mapped pointer in mapping table
with new delta records pointer
- iii. If contention occurs,
mapped delta wins, losed updates should retry or abort
- iv. Consolidate delta chain to copy of base page,
and replace mapped pointer with this

C. Mapping Table

- i. Map page ID and its physical pointer
- ii. All page can replace pointer to other pages with its page ID

D. Garbage collection

- i. Reference counting
 1. Maintain counter to indicate # of threads that accessing it
 2. If counter become zero, it can be GCed.
- ii. Epoch based reclamation
 1. Maintain global epoch
 2. Maintain object-local epoch that indicates when object deleted
 3. If all threads left that epoch, the object can be reclaimed.

E. Modifications

- i. Split delta record – delta for B Tree split operation
 - 1. Delay to update mapping table with split operation
- ii. Separator delta record – delta for search shortcut