

COMP6016 Assembly Language Programming III

```
call
jmp
15CA
call
mov
xor
mov
mov
bx
selection
call
totall
total
```

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Arrays

- •Arrays are implemented as fixed size memory buffers
 •It is up to the programmer to manage the size of the
 array elements and to deal with iteration through the
 array
- •Use [] to specify that something is an address, not a value. Some (limited) arithmetic can be done on this.
- •Care *MUST* be taken over what is a value and what is an address (think pointer arithmetic in C)



Arrays - addressing

```
mov eax, [ebx] ; Move the 4 bytes in memory at the
address contained in EBX into EAX
mov [var], ebx; Move the contents of EBX into the 4
bytes at memory address var.
                   ; (Note, var is a 32-bit constant).
mov eax, [esi-4] ; Move 4 bytes at memory address ESI + (-
4) into EAX
mov [esi+eax], cl ; Move the contents of CL into the byte
at address ESI+EAX
mov edx, [esi+4*ebx] ; Move the 4 bytes of data at
address EST+4*EBX into EDX
; The following address calculations are FORBIDDEN!
mov eax, [ebx-ecx]
                          ; Can only add register
values
mov [eax+esi+edi], ebx ; At most 2 registers in address
computation
```



Arrays – addressing - sizes

•Most of the time, the assembler can work out the memory to use when we move to a memory address but this may not always be possible, especially with constants -

```
mov [rax], 6; how long is the 6? 1,2,4, or 8 bytes?
If there is any ambiguity, we need to specify the size of the destination -
mov BYTE [rax], 2; Move 2 into the single byte at the
address stored in RAX.
mov WORD [rax], 2; Move the 16-bit integer representation
of 2 into the 2
                                  ; bytes starting at the
address in RAX.
mov DWORD [rax], 2 ; Move the 32-bit integer
representation of 2 into the
                                         ; 4 bytes starting
at the address in RAX.
mov QWORD [rax], 2; Move the 64-bit integer
representation of 2 into the
                                         ; 8 bytes starting
at the address in RAX.
```



Arrays – times and lea

•Two useful instructions for use with arrays are times and lea •times – repeats data or instructions the specified number of times (NB – will not work with macros)

```
zerobuf: times 64 db 0 buffer: db 'hello, world' times 64-$+buffer db ''
```

•lea – load effective address – put the address specified by the second operand into the register specified by the first operand lea parameters are in the form base register, index register, scalar multiplier (must be 1,2,4, or 8), and offset.

```
lea rdx, [rax+ rbx*4+67]
```

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```
; Test programme for assembler arrays
; nasm -g -I libasm io-master/include -f
elf64 array2.asm
; gcc -g array2.o -lasm io -o array2
%include
"/home/malware/asm/joey lib io v6 releas
e.asm"
global main
```



section .data

```
echo_welcome: db "Hello, this is a character storage programme",0
echo_bye: db "Goodbye!",0
characters: times 26 db 0 ; an array big enough for the letters
```



section .text

```
; This is our procedure for putting the
characters into the array
     populate:
          lea rax, [characters]
          loop1:
               mov [rax], rbx
                inc rbx
                inc rax
                cmp rbx, 91
               jne loop1
          ret
```



```
; This prints them out
     display:
          loop2:
                movzx rdi, BYTE
[characters+rbx]
                call print char new
                           ; PRINT CHAR rdi
                           ; NEWLINE
                call print nl new
                inc rbx
                cmp rbx, 26
                jne loop2
          ret
```



```
main:
    mov rbp, rsp; for correct debugging
             push rbp
             mov rbp, rsp
              sub rsp, 32
             mov rdi, QWORD echo welcome
                 call print string new
                 ; PRINT STRING [rdi]
                 call print nl new
                 ; NEWLINE
             mov rbx, 65; ASCII character for A
              call populate
              xor rbx, rbx
              call display
```



```
mov rdi, QWORD echo_bye
call print_string_new ;PRINT_STRING [rdi]
call print_nl_new ;NEWLINE

add rsp, 32

pop rbp
ret
```



```
; Test programme for assembler arrays
; nasm -g -I libasm_io-master/include -f
elf64 array3.asm
; gcc -g array3.o -lasm_io -o array3
%include
"/home/malware/asm/joey_lib_io_v6_release.asm
"
global main
```



section .data

```
echo welcome: db "Hello, this is a student ID storage
programme", 0
echo instruction: db "Please enter a student ID: ",0
echo message: db "A valid ID is :",0
echo bye: db "Goodbye!", 0
       dq 0 ; a loop counter
counter:
characters: times 90 db 0 ; an array big enough for
10 student IDs (8 chars + null)
```



```
section .text
; This is our procedure for reading the IDs
populate:
     lea rbx, [characters]
     mov QWORD [counter], 0
     loop1:
          mov rdi, QWORD echo instruction
               call print string new
               call print nl new
               call read string new
```



; We can't move a memory address to another memory address using mov so we use rcx to store the address

```
mov rcx,[rax]
mov [rbx], rcx
add rbx,9
add QWORD [counter],9
cmp QWORD [counter], 27
jne loop1
ret
```



```
; This prints them out
     display:
         loop2:
           mov rdi, QWORD echo message
            call print string new; set rdi to be
the start of the buffer
           mov rdi, characters ; and add the
number of bytes to get to the current record
            add rdi, rbx
            call print string new
            call print nl new
            add rbx, 9
            cmp rbx, 27
            jne loop2
```

ret



```
main:
    mov rbp, rsp; for correct debugging
          push rbp
          mov rbp, rsp
          sub rsp,32
          mov rdi, QWORD echo welcome
          call print string new
          call print nl new
          call populate
          sub rbx, rbx
          call display
```



```
mov rdi, QWORD echo_bye
call print_string_new
call print_nl_new

add rsp, 32

pop rbp

ret
```



Strings

- Strings are arrays of bytes, however there are instructions to help us more efficiently
- Need to consider if we have sentinel style strings (like
 C) where there is a known character at the end (normally NULL)
- •Or fixed length strings (like Pascal) where the length is encoded as part of the string
- •For assembler, most of the time we prefer sentinel strings, but fixed length have their uses
- •Also need to be away that many of the strings that we use are essentially byte blocks and the string optimisations are applicable elsewhere



```
; Test programme for assembler arrays
; nasm -g -I libasm io-master/include -f elf64 strings.asm
; gcc -g strings1.o -lasm io -o strings
%include "/home/malware/asm/joey lib io v6 release.asm"
qlobal main
section .data
      echo welcome: db "Hello, this is a long message", 0
      len equ $ - echo welcome
      echo message: db "It is this many characters :",0
      echo bye: db "Goodbye!",0
```



```
section .text
        main:
                 push rbp
                 mov rbp, rsp
                 sub rsp, 32
                 mov rdi, QWORD echo welcome
                 call print string new
                 call print nl new
                 mov rdi, QWORD echo message
                 call print string new
                 mov rdi, len
                 dec rdi
                 call print uint new
                 call print nl new
                 mov rdi, QWORD echo bye
                 call print string new
                 call print nl new
                 add rsp, 32
                 pop rbp
```

ret



Strings – stosb & movsb

- movs(b/w/d/q) copy the (byte/word/double word/quad word) from memory address stored in rsi to memory address stored in rdi
- stos(b/w/d/q) copies the data from
- (AL/AX/EAX/RAX) to memory address stored in rdi
- •We use movs? to copy from memory location to memory location. This is used to copy blocks of memory
- •We use stos? to copy from a register to a memory location this is used to initialise an area of memory



Strings - rep

- •The instruction rep is used to allow us to repeat stos? and movs? multiple times, enabling fast memory copies and initialisation.
- •To use rep, we need to specify -
- •The direction of the transfer (high to low (std) or low to high (cld))
- •For stos?, the value to copy in which is stored in AL, AX, EAX, or RAX
- •For movs?, the location to copy from which is stored in ESI or RDI
- •The location to copy to which is stored in EDI or RDI
- •The number of bytes to copy which is stored in CL, CX, ECX, or RCX



```
; Test programme for assembler arrays
; nasm -g -I libasm io-master/include -f elf64 strings2.asm
; gcc -g strings1.o -lasm io -o strings2
%include "/home/malware/asm/joey lib io v6 release.asm"
qlobal main
section .data
      echo welcome: db "Hello, this is a long message", 0
      len equ $ - echo welcome
      echo message: db "This is it changed to be all As",0
      echo bye: db "Goodbye!",0
```



```
section .text
main:
        push rbp
        mov rbp, rsp
sub rsp,32
        mov rdi, QWORD echo welcome
        call print string new
        call print nl new
        mov rdi, QWORD echo message
        call print_string_new
        call print nl new
       Cld ; set the direction mov al, 'A' ; we want to overwrite with 'A' mov edi, echo_welcome; we want to overwrite the
                                        echo welcome string
        mov ecx, len+1; we need t\bar{o} add 1 here as the
                              string starts at 0
        rep stosb
                               ; repeat
        mov byte [echo welcome+len+1], 0 ; add the null
                                                terminator
```



```
mov rdi, QWORD echo welcome call print string new call print_nl_new

mov rdi, QWORD echo bye call print string new call print_nl_new

add rsp, 32

pop rbp

ret
```



Strings – cmps

•We can compare strings using cmps(b/w/d/q) – this will compare the values at the addresses pointed to by (DI, EDI, or RDI) and (SI, ESI, or RSI) respectively cmps can be combined with rep, repe (repeat while a condition is zero and CX!=0), or repne (repeat while a condition is not zero or CX!=0) to check blocks of memory



Strings – scas

- •We can search for a character, or set of characters using the instruction scas(b/w/d/q)
- •The register (AL, AX, EAX, or RAX) contains the item we want to search for, and the string to search is at the location stored in (DI, EDI, or RDI)
- •The register (CX, ECX, or RCX) is used to limit how far we will search.
- •When used in conjunction with repne we can scan through a string and when we have found the substring we are looking for, the register (DI, EDI, or RDI) will contain the position in the string of the substring in memory (take off the start of the string to find out how far into the string something is)



Strings – length of a variable string

```
; Test programme for assembler arrays
 nasm -g -I libasm io-master/include -f elf64 strings3.asm
; gcc -g strings3.o -lasm io -o strings3
%include "/usr/local/include/libasm io.inc"
global main
section .data
                           "Please enter a string", 0
      echo welcome: db
                           "The string is this long: ",0
      echo message: db
      echo bye:
                           "Goodbye!", 0
                    db
```



Strings – length of a variable string

```
section .text
          main:
                    push rbp
                    mov rbp, rsp
                    sub rsp,32
                    mov rdi, QWORD echo welcome
                    call print string
                    call print nl
                    call read string
                    mov rbx, rax
                                                                        ; move the address of
the string
                                                                                  ; somewhere
safe
                    mov rdi, QWORD echo message
                    call print string
                                                              ; this sets the max size to
                    sub rcx, rcx
look for to be
                                                   ; the maximum memory size
                    not rcx
```



Strings – length of a variable string

```
sub al, al
                                                   ; We want to look for the null terminator
                    mov rdi, rbx
                                                   ; set the start of the string
                    cld
                    repne scasb
                                                             ; search
                                                   ; we need to take off the start of the
                    sub rdi, rbx
string
                                                   ; and allow for the null terminator
                    dec rdi
                    call print uint
                    call print nl
                    mov rdi, QWORD echo bye
                    call print string
                    call print nl
                    add rsp, 32
                    pop rbp
                    ret
```

Understanding C Code construct in Assembly



- •Malware analyst need to be able to obtain a high-level picture of the code functionality.
- .Skill that need time to develop.
- •Typically malware is developed using high-level language, commonly C.
- •A code construct is a code abstraction level that defines a functional property (not detail implementation).
- •Eg: loops, conditional statement and so on.
- Discuss on popular C code construct.
- Malware analyst need to be able to go from disassembly to high level constructs,
- For help with C language, have a look look at the classic *The C Programming Language* by Brian Kernighan and Dennis Ritchie (Prentice-Hall, 1988)
- •Remember that the goal is to understand the overall functionality of a program, not to analyze every single instruction.



Global vs Local Variables

•Global variables can be accessed and used by any function in a program.

•Local variables can be accessed only by the function in which they are defined. Both global and local variables are declared similarly in C, but they look completely different in assembly.

```
int x = 1;
int y = 2;

void main()
{
    x = x+y;
    printf("Total = %d\n", x);
}
```

Figure: A simple program with two global variables



Global vs Local Variables cont

•Global variables can be accessed and used by any function in a program.

•Local variables can be accessed only by the function in which they are defined. Both global and local variables are declared similarly in C, but they look completely different in assembly.

```
int x = 1;
int y = 2;

void main()
{
    x = x+y;
    printf("Total = %d\n", x);
}
```

Figure 1: A simple program with two global variables

```
void main()
{
   int x = 1;
   int y = 2;

   x = x+y;
   printf("Total = %d\n", x);
}
```

Figure 2: A simple program with two global variables



Global vs Local Variables cont

00401003	mov	eax, dword_40CF60
00401008	add	eax, dword_40C000
0040100E	mov	dword_40CF60, eax ❶
00401013	mov	ecx, dword_40CF60
00401019	push	ecx
0040101A	push	offset aTotalD ;"total = %d\n"
0040101F	call	printf
		•

00401006	mov	dword ptr [ebp-4], 0
0040100D	mov	dword ptr [ebp-8], 1
00401014	mov	eax, [ebp-4]
00401017	add	eax, [ebp-8]
0040101A	mov	[ebp-4], eax
0040101D	mov	ecx, [ebp-4]
00401020	push	ecx
00401021	push	offset aTotalD ; "total = %d\n"
00401026	call	printf

Fig 3: Assembly code for the global variable example in Fig 1

Fig 4: Assembly code for the local variable example in Fig 2

- The global variables are referenced by memory addresses, and the local variables are referenced by the stack addresses.
- In Listing Fig 3, the global variable x is signified by dword_40CF60, a memory location at 0x40CF60. Notice that x is changed in memory when eax is moved into dword_40CF60 at (1). All subsequent functions that utilize this variable will be impacted.
- In Fig 4, memory location [ebp-4] is used consistently throughout this function to reference the local variable x.



Disassembling Arithmetic Operations

```
int a = 0;
int b = 1;
a = a + 11;
a = a - b;
a--;
b++;
b = a % 3;
```

Fig 5: C code with two variables and a variety of arithmetic

Figure 5 shows the C code for two variables and a variety of arithmetic operations. Two of these are the -- and ++ operations, which are used to dec- rement by 1 and increment by 1, respectively. The % operation performs the *modulo* between the two variables, which is the remainder after performing a division operation.



Disassembling Arithmetic Operations

00401006	mov	[ebp+var_4], 0
0040100D	mov	[ebp+var_8], 1
00401014	mov	eax, [ebp+var_4] 🛛
00401017	add	eax, OBh
0040101A	mov	<pre>[ebp+var_4], eax</pre>
0040101D	mov	ecx, [ebp+var_4]
00401020	sub	ecx, [ebp+var_8] ②
00401023	mov	<pre>[ebp+var_4], ecx</pre>
00401026	mov	edx, [ebp+var_4]
00401029	sub	edx, 1 3
0040102C	mov	<pre>[ebp+var_4], edx</pre>
0040102F	mov	<pre>eax, [ebp+var_8]</pre>
00401032	add	eax, 1 4
00401035	mov	<pre>[ebp+var_8], eax</pre>
00401038	mov	<pre>eax, [ebp+var_4]</pre>
0040103B	cdq	
0040103C	mov	ecx, 3
00404044		
00401041	idiv	ecx
00401043	mov	[ebp+var_8], edx ❺

Fig 6: Assembly code for the arithmetic example in Fig 5

- In this example, a and b are local variables because they are referenced by the stack. IDA Pro has labeled a as var_4 and b as var 8.
- First, var_4 and var_8 are initialized to 0 and 1, respectively. a is moved into eax (1), and then 0x0b is added to eax, thereby incrementing a by 11. b is then subtracted from a (2). (The compiler decided to use the sub and add instructions (3) and (4), instead of the inc and dec functions.)
- The final five assembly instructions implement the modulo. When performing the div or idiv instruction (5), we are dividing edx:eax by the operand and storing the result in eax and the remainder in edx. That is why edx is moved into var 8 (5).



Recognizing if Statements

- If statement use by programmer to alter program execution based on certain conditions. if statements are common in C code and disassembly.
- Figure 7 displays a simple if statement in C with the assembly for this code shown in Figure 8. Notice the conditional jump jnz at (2). There must be a conditional jump for an if statement, but not all conditional jumps correspond to if statements.

```
int x = 1;
int y = 2;

if(x == y){
    printf("x equals y.\n");
}else{
    printf("x is not equal to y.\n");
}
```

```
00401006
                        [ebp+var 8], 1
                mov
                        [ebp+var 4], 2
0040100D
                mov
                        eax, [ebp+var 8]
00401014
                mov
                        eax, [ebp+var 4] 0
00401017
                cmp
                        short loc 40102B 🕹
0040101A
                jnz
                        offset aXEqualsY ; "x equals y.\n"
0040101C
                push
                call
                        printf
00401021
00401026
                add
                        esp, 4
                        short loc 401038 6
00401029
                jmp
0040102B loc 40102B:
0040102B
                push
                        offset aXIsNotEqualToY; "x is not equal to y.\n"
                call
00401030
                        printf
```

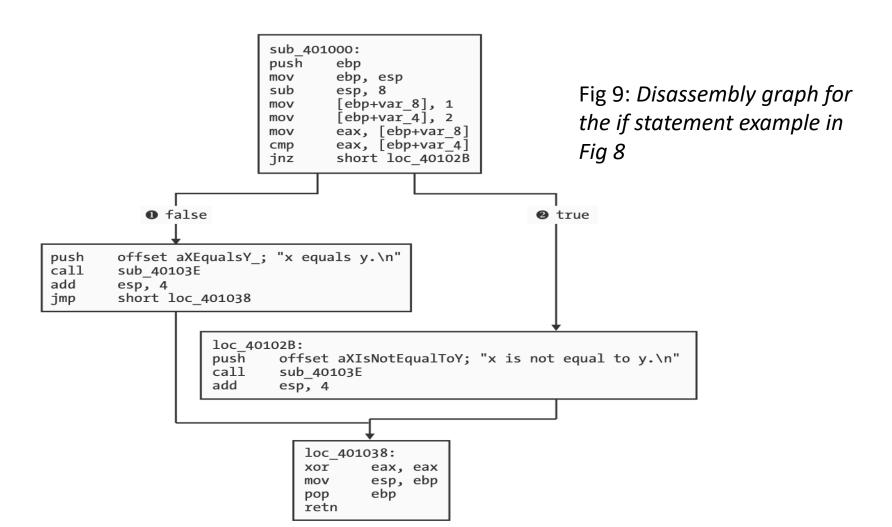
Fig 7: C code if statement example

Fig 8: Assembly code for the if statement example in Fig 7

Graphic function with IDA PRO



• IDA Pro has a graphing tool that is useful in recognizing constructs, as shown in Figure 9. This feature is the default view for analyzing functions.



Recognizing for Loops



The for loop is a basic looping mechanism used in C programming. for loops always have four components: initialization, comparison, execution instructions, and the increment or decrement.

```
int i;
for(i=0; i<100; i++)
{
    printf("i equals %d\n", i);
}</pre>
```

Fig 10: *C code for a for loop*

- The initialization sets i to 0 (zero), and the comparison checks to see if i is less than 100.
- If i is less than 100, the printf instruction will execute, the increment will add 1 to i, and the process will check to see if i is less than 100.
- These steps will repeat until i is greater than or equal to 100



```
[ebp+var 4], 0 ①
00401004
                mov
                        short loc 401016 2
0040100B
                jmp
0040100D loc 40100D:
0040100D
                        eax, [ebp+var 4] 3
                mov
                        eax, 1
00401010
                add
                        [ebp+var 4], eax 4
00401013
                mov
00401016 loc 401016:
00401016
                        [ebp+var 4], 64h ⑤
                cmp
                        short loc 40102F 6
0040101A
                jge
                        ecx, [ebp+var 4]
0040101C
                mov
0040101F
                push
                        ecx
                        offset aID ; "i equals %d\n"
                push
00401020
                call
                        printf
00401025
0040102A
                add
                        esp, 8
                        short loc 40100D €
0040102D
                jmp
```

Fig 11: Assembly code for the for loop example in Fig 10

- In assembly, the for loop can be recognized by locating the four components—initialization, comparison, execution instructions, and increment/decrement.
- For example, in Figure above, (1) corresponds to the initialization step. The code between
 (3) and (4) corresponds to the increment that is initially jumped over at (2) with a jump
 instruction. The comparison occurs at (5), and at (6), the decision is made by the
 conditional jump.
- If the jump is not taken, the printf instruction will execute, and an unconditional jump occurs at (7), which causes the increment to occur.

Loops Graph Mode



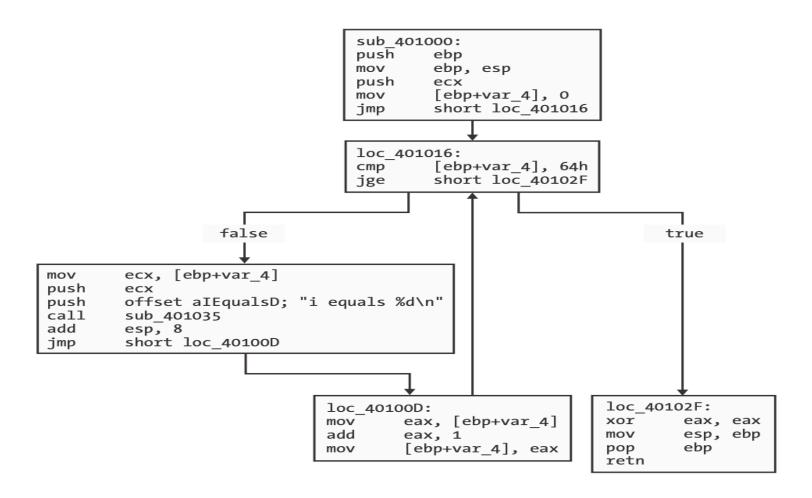


Figure 12: Disassembly graph for the for loop example in Fig 11

Recognizing While Loops



The while loop is frequently used by malware authors to loop until a condition is met, such as receiving a packet or command. while loops look similar to for loops in assembly,

but they are easier to understand.

```
int status=0;
int result = 0;

while(status == 0){
    result = performAction();
    status = checkResult(result);
}
```

Fig 13: C code for a while loop

```
[ebp+var 4], 0
00401036
                mov
                         [ebp+var 8], 0
0040103D
                mov
00401044 loc 401044:
00401044
                         [ebp+var 4], 0
                cmp
                         short loc 401063 1
00401048
                jnz
                         performAction
                call
0040104A
                         [ebp+var 8], eax
0040104F
                mov
                         eax, [ebp+var 8]
00401052
                mov
00401055
                push
                         eax
                call
                         checkResult
00401056
0040105B
                add
                         esp, 4
                         [ebp+var 4], eax
0040105E
                mov
                         short loc 401044 @
00401061
                jmp
```

Fig 14: Assembly code for the while loop example in Fig 13

- The while loop in Fig 13 will continue to loop until the status returned from checkResult is
 0.
- A conditional jump occurs at (1) and an unconditional jump at (2), but the only way for this
 code to stop executing repeatedly is for that conditional jump to occur.
- Revision: Read PMA page 119-132