« Presentation of the Abstraction project proposal »

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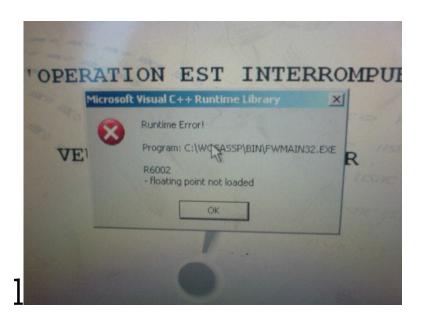
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2. The Problem: The Design of Safe and Secure Computer-Based Systems

Software is Everywhere

- exponential growth of hardware since 1975
- ⇒ exponential growth of software (favored by software engineering methods)
- mainly manual activity \Rightarrow bugs are everywhere





Guaranteeing the Reliability and Security of Software-Intensive Systems

- an objective of the INRIA strategic plan
- an industrial categorical imperative, in particular for safety and security critical software (validation can account for up to 60% of software development costs)









Validation/Formal Methods

- bug-finding methods: unit, integration, and system testing, dynamic verification, bounded model-checking,
 ...
- absence of bug proving methods: formally prove that the semantics of a program satisfies a specification
 - theorem-proving & proof checking
 - model-checking
 - abstract interpretation
- in practice : complementary methods are used,
 very difficult to scale up

3. Abstract Interpretation

The Theory of Abstract Interpretation

- a theory of sound approximation of mathematical structures, in particular those involved in the behavior of computer systems
- systematic derivation of sound methods and algorithms for approximating undecidable or highly complex problems in various areas of computer science
- main current application is on the safety and security of complex hardware and software computer systems

Applications of Abstract Interpretation

- Static Program Analysis [119], [124], [120] including Dataflow Analysis; [120], [123], Set-based Analysis [122], Predicate Abstraction [7], ...
- Grammar Analysis and Parsing [14];
- Hierarchies of Semantics and Proof Methods [121], [10];
- Typing & Type Inference [118];
- (Abstract) Model Checking [123];

Applications of Abstract Interpretation (Cont'd)

- Program Transformation [33];
- Software Watermarking [44];
- Bisimulations [129];
- Language-based security [125];
- Semantic-based obfuscated malware detection [128].

All these techniques involve sound approximations that can be formalized by abstract interpretation

4. An Example of Theoretical Application : Semantics of the Eager λ -calculus

[1] P. Cousot & R. Cousot. Bi-inductive structural semantics. Februray 15th, 2007. Submitted.



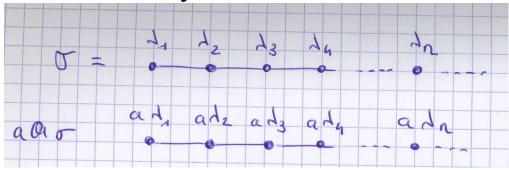
Syntax of the Eager λ -calculus

Traces

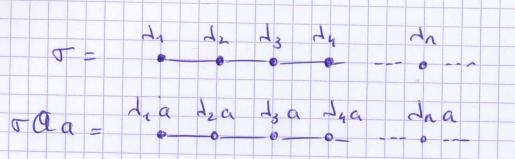
- $-\mathbb{T}^*$ (resp. \mathbb{T}^+ , \mathbb{T}^ω , \mathbb{T}^∞ and \mathbb{T}^∞) be the set of finite (resp. nonempty finite, infinite, finite or infinite, and nonempty finite or infinite) sequences of terms
- $\text{ If } \sigma \in \mathbb{T}^+ \text{ then } |\sigma| > 0 \text{ and } \sigma = \sigma_0 \bullet \sigma_1 \bullet \ldots \bullet \sigma_{|\sigma|-1}.$
- $-\text{ If }\sigma\in\mathbb{T}^\omega\text{ then }|\sigma|=\omega\text{ and }\sigma=\sigma_0\bullet\ldots\bullet\sigma_n\bullet\ldots$
- Given $S, T \in \wp(\mathbb{T}^{\infty})$, we define $S^{+} \triangleq S \cap \mathbb{T}^{+}$, $S^{\omega} \triangleq S \cap \mathbb{T}^{\omega}$ and $S \sqsubseteq T \triangleq S^{+} \subseteq T^{+} \wedge S^{\omega} \supseteq T^{\omega}$, so that $\langle \wp(\mathbb{T}^{\infty}), \sqsubseteq, \mathbb{T}^{\omega}, \mathbb{T}^{+}, \sqcup, \sqcap \rangle$ is a complete lattice.

Operations on traces

- For $a \in \mathbb{T}$ and $\sigma \in \mathbb{T}^{\infty}$, we define $a@\sigma$ to be $\sigma' \in \mathbb{T}^{\infty}$ such that $\forall i < |\sigma| : \sigma'_i = a \ \sigma_i$ and,



- similarly $\sigma @ a$ is σ' such that $\forall i < |\sigma| : \sigma'_i = \sigma_i \ a$.



Bifinitary Trace Semantics \vec{S} of the Eager λ -calculus [121]

Note: $a[x \leftarrow b]$ is the capture-avoiding substitution of b for all free occurences of x within a. We let FV(a) be the free variables of a. We define the call-by-value semantics of closed terms (without free variables) $\overline{\mathbb{T}} \triangleq \{a \in \mathbb{T} \mid FV(a) = \varnothing\}.$

Abstraction to the Bifinitary Relational Semantics of the Eager λ -calculus

remember the input/output behaviors, forget about the intermediate computation steps

$$lpha(T) \stackrel{ ext{def}}{=} \{lpha(\sigma) \mid \sigma \in T\}$$
 $lpha(\sigma_0 ullet \sigma_1 ullet \dots ullet \sigma_n) \stackrel{ ext{def}}{=} \langle \sigma_0, \ \sigma_n
angle$ $lpha(\sigma_0 ullet \dots ullet \sigma_n ullet \dots) \stackrel{ ext{def}}{=} \langle \sigma_0, \ oldsymbol{\perp}
angle$

Bifinitary Relational Semantics of the Eager λ -calculus

$$\begin{array}{l} \mathsf{v} \Rightarrow \mathsf{v}, \quad \mathsf{v} \in \mathbb{V} \\ \hline \mathsf{a} \Rightarrow \bot \\ \hline \mathsf{a} \mathsf{b} \Rightarrow \bot \\ \hline \\ \mathsf{c}, \quad \mathsf{v} \in \mathbb{V}, \quad r \in \mathbb{V} \cup \{\bot\} \\ \hline \\ \mathsf{a} \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{v} \mathsf{b} \Rightarrow r \\ \hline \\ \mathsf{a} \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{v} \mathsf{b} \Rightarrow r \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{a} \mathsf{v} \Rightarrow r \\ \hline \\ \mathsf{a} \mathsf{b} \Rightarrow \mathsf{r} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{a} \mathsf{v} \Rightarrow r \\ \hline \\ \mathsf{a} \mathsf{b} \Rightarrow \mathsf{r} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{a} \mathsf{v} \Rightarrow r \\ \hline \\ \mathsf{a} \mathsf{b} \Rightarrow \mathsf{r} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{a} \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v}, \quad \mathsf{b} \Rightarrow \mathsf{v} \\ \hline \\ \mathsf{b} \Rightarrow \mathsf{v} \\$$

Abstraction to the Natural Big-Step Semantics of the Eager λ -calculus

remember the finite input/output behaviors, forget about non-termination

Natural Big-Step Semantics of the Eager λ -calculus [126]

$$egin{aligned} \mathbf{v} &\Longrightarrow \mathbf{v}, \quad \mathbf{v} \in \mathbb{V} \ & \dfrac{\mathbf{a}[\mathbf{x} \leftarrow \mathbf{v}] \Longrightarrow r}{(oldsymbol{\lambda} \, \mathbf{x} \cdot \mathbf{a}) \quad \mathbf{v} \Longrightarrow r} \subseteq, \quad \mathbf{v} \in \mathbb{V}, \ r \in \mathbb{V} \ & \dfrac{\mathbf{a} \Longrightarrow \mathbf{v}, \quad \mathbf{v} \ \mathbf{b} \Longrightarrow r}{\mathbf{a} \ \mathbf{b} \Longrightarrow r} \subseteq, \quad \mathbf{v} \in \mathbb{V}, \ r \in \mathbb{V} \ & \dfrac{\mathbf{b} \Longrightarrow \mathbf{v}, \quad \mathbf{a} \ \mathbf{v} \Longrightarrow r}{\mathbf{a} \ \mathbf{b} \Longrightarrow r} \subseteq, \quad \mathbf{a} \in \mathbb{V}, \ \mathbf{v} \in \mathbb{V}, \ r \in \mathbb{V} \ . \end{aligned}$$

Abstraction to the Small-Step Operational Semantics of the Eager λ -calculus

remember execution steps, forget about their sequencing

$$egin{aligned} lpha(T) \stackrel{ ext{def}}{=} igcup_{\{lpha(\sigma) \mid \sigma \in T\}} \ & \ lpha(\sigma_0 ullet \sigma_1 ullet \ldots ullet \sigma_n) \stackrel{ ext{def}}{=} \{\langle \sigma_i, \ \sigma_{i+1}
angle \mid 0 \leqslant i \land i < n\} \ & \ lpha(\sigma_0 ullet \ldots ullet \sigma_n ullet \ldots) \stackrel{ ext{def}}{=} \{\langle \sigma_i, \ \sigma_{i+1}
angle \mid i \geqslant 0\} \end{aligned}$$

Small-Step Operational Semantics of the Eager λ -calculus [127]

$$((\lambda \times a) \vee) \longrightarrow a[x \leftarrow v]$$

$$a_0 \longrightarrow a_1$$

$$a_0 \longrightarrow a_1 \longrightarrow a_1$$

$$b_0 \longrightarrow b_1$$

$$v \mapsto b_0 \longrightarrow v \mapsto b_1$$

The Abstract Semantics are Correct by Calculational Design

$v \mapsto v, v \in V$ $a[x \leftarrow v] \mapsto \sigma$ $v \in V$	
$(\lambda \times \cdot a) \vee i \Rightarrow (\lambda \times \cdot a) \vee \cdot \sigma$ $a \mapsto \sigma \cdot v, v \mapsto \sigma'$	
$ab \mapsto \sigma \otimes b$ $ab \mapsto (\sigma \otimes b) \cdot \sigma'$ $c, \forall \in V, \sigma \in A$	
$\frac{b \bowtie \sigma}{a, b \bowtie \sigma \otimes \sigma} = c, a \in V, \sigma \in T^{\omega} \frac{b \bowtie \sigma \cdot v, a \vee \bowtie \sigma'}{a, b \bowtie \sigma \otimes \sigma} = c, a, v \in V, \sigma \in T^{\omega}$	
$a \mapsto a \boxtimes \sigma$ $a \in V, \sigma \in T^{\omega}$ $b \mapsto a \boxtimes \sigma$ $a \in V, \sigma \in T^{\omega}$ $a \mapsto a \boxtimes \sigma$ $a \mapsto a \boxtimes \sigma$	1
5.4 Abstraction into the big-step relational semantics of the call-by-value calculus	: λ-
5.4.1 Relational abstraction of traces	
The relational abstraction of sets of traces is	
$\alpha \in \wp(\mathbb{T}^{\infty}) \mapsto \wp(\mathbb{T} \times (\mathbb{T} \cup \{\bot\}))$ $\alpha(S) \triangleq \{(\sigma_0, \sigma_{a-1}) \mid \sigma \in S \land \sigma = n\} \cup \{(\sigma_0, \bot) \mid \sigma \in S \land \sigma = \omega\}$ $\gamma \in \wp(\mathbb{T} \times (\mathbb{T} \cup \{\bot\})) \mapsto \wp(\mathbb{T}^{\infty})$ $\gamma(T) \triangleq \{\sigma \in \mathbb{T}^{\infty} (\sigma = n \land (\sigma_0, \sigma_{a-1}) \in T) \lor (\sigma = \omega \land (\sigma_0, \bot) \in T)$	(4)
so that	•
$\langle \wp(\mathbb{T}^\infty),\subseteq\rangle \xrightarrow[\alpha \]{} \langle \wp(\mathbb{T}\times(\mathbb{T}\cup\{\bot\})),\subseteq\rangle\ .$	(5)
Proof	
$\alpha(S) \subseteq T$	
\iff $\{\langle \sigma_0, \sigma_{n-1} \rangle \mid \sigma \in S \land \sigma = n\} \cup \{\langle \sigma_0, \bot \rangle \mid \sigma \in S \land \sigma = \omega\} \subseteq T$ $\{\text{def.}$	20
\iff $\forall \sigma \in S^+ : (\sigma_0, \ \sigma_{ \sigma -1}) \in T^+ \land \forall \sigma \in S^\omega : (\sigma_0, \ \bot) \in T^\omega$	
$\{\text{def.} \subseteq, S^+ \triangleq S \cap \mathbb{T}^+, \text{ and } S^\omega \triangleq S \cap \mathbb{T}^+ \}$ $\iff S^+ \subseteq \{\sigma \mid \sigma = n \land \langle \sigma_0, \sigma_{n-1} \rangle \in T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq T\} \land S^\omega \subseteq \{\sigma \mid \sigma = \omega \land \langle \sigma_0, \bot \rangle \in T\} \land S^\omega \subseteq T$ \S^\omega \subseteq T\} \land	
S ⊆ {a a = n ∧ (a ₀ , a _{n-1}) ∈ 1 } ∧ S ⊆ {a a = 2 ∧ (a ₀ , ±) ∈	21
18	
$\{\text{def.}\subseteq,T^+\triangleq T\cap(\mathbb{T}\times\mathbb{T}),\text{and}T^\omega\triangleq T\cap(\sigma\mathbb{T}\times\{\bot\})$	2({
$\iff S \subseteq \gamma(T)$ $\{S = S^+ \cup S^\omega \text{ and def. } \gamma(T)\}$	
5.4.2 Bifinitary relational semantics	
The bifinitary relational semantics $\widetilde{\mathbb{S}} \triangleq \alpha(\widetilde{\mathbb{S}}) \in \wp(\mathbb{T} \times (\mathbb{T} \cup \{\bot\}))$ is relational abstraction of the trace semantics mapping an expression to final value or \bot in case of divergence.	the its
5.4.3 Fixpoint big-step bifinitary relational semantics	
The bifinitary relational semantics $\vec{S} \triangleq \alpha(\vec{S}) = \alpha(\mathbf{up}^{\vec{v}} \vec{F})$ can be defined fixpoint form as $\mathbf{up}^{\vec{v}} \vec{F}$ where the big-step transformer $\vec{F} \in \rho(\mathbb{T} \times (\mathbb{T} \cup \{\bot\}))$ is	
$\vec{F}(T) \triangleq \{\langle v, v \rangle \mid v \in V\} \cup$	(6)
$\{ \langle (\boldsymbol{\lambda} \mathbf{x} \boldsymbol{\cdot} \mathbf{a}) \mathbf{v}, r \rangle \mid \mathbf{v} \in \mathbb{V} \wedge \langle \mathbf{a} [\mathbf{x} \leftarrow \mathbf{v}], r \rangle \in T \} \cup $	
$\{(\langle \mathbf{a} \ \mathbf{b}), \perp \rangle \mid \langle \mathbf{a}, \perp \rangle \in T\} \cup$ $\{(\langle \mathbf{a} \ \mathbf{b}), r \rangle \mid \langle \mathbf{a}, \mathbf{v} \rangle \in T^+ \land \mathbf{v} \in V \land \langle (\mathbf{v} \ \mathbf{b}), r \rangle \in T\} \cup$	
$\{((a b), r) \mid (a, v) \in I \land v \in V \land ((v b), r) \in I \} \cup \{((a b), \bot) \mid a \in V \land (b, \bot) \in T \} \cup \{((a b), \bot) \mid a \in V \land (b, \bot) \in T \} \cup \{((a b), \bot) \mid a \in V \land ((b b), \bot) \in T \} \cup \{((a b), \bot) \mid a \in V \land ((a b), \bot) \in T \} \cup \{((a b), \bot) \mid a \in V \land ((a b), \bot) \in T \} \cup \{((a b), \bot) \mid a \in V \land ((a b), \bot) \in T \} \cup \{((a b), \bot) \mid a \in V \land ((a b), \bot) \in T \} \cup \{(a b), \bot ((a b), \bot) \mid a \in V \land ((a b), \bot) \in T \} \cup \{(a b), \bot ((a b), \bot) \mid a \in V \land ((a b), \bot) \in T \} \cup \{(a b), \bot ((a b), \bot) \mid a \in V \land ((a b), \bot) \in T \} \cup \{(a b), \bot ((a b), \bot) \mid a \in V \land ((a b), \bot) ((a b), \bot) \in T \} \cup \{(a b), \bot ((a b), \bot) \mid a \in V \land ((a b), \bot) ((a b), \bot$	
$\{\langle (\mathbf{a}\ \mathbf{b}),\ r\rangle\ \ \mathbf{a},\mathbf{v}\in \mathbb{V}\wedge \langle \mathbf{b},\ \mathbf{v}\rangle \in T^+\wedge \langle (\mathbf{a}\ \mathbf{v}),\ r\rangle \in T\}\ .$	
Lemma 10 $\alpha(\vec{F}(S)) = \vec{F}(\alpha(S))$ PROOF α is a complete \cup -morphism, so we calculate $\alpha(\vec{F}(S))$ by cases.	0
$-\alpha(\{v \in \mathbb{T}^{\infty} \mid v \in \mathbb{V}\})$	
$= \ \{ \langle v, v \rangle \mid v \in \mathbb{V} \} \qquad \qquad \text{$\langle \det \alpha$ and $ v = $}$	- 15
$= \alpha(\{(\lambda \times a) \lor \cdot a[x \leftarrow v] \cdot \sigma \mid v \in \forall \land a[x \leftarrow v] \cdot \sigma \in S\})$ $= \alpha(\{(\lambda \times a) \lor \cdot a[x \leftarrow v] \cdot \sigma \mid v \in \forall \land a[x \leftarrow v] \cdot \sigma \in S^+\}) \cup \alpha(\{(\lambda \times a) \lor \cdot a[x \leftarrow v] \cdot \sigma \in S^+\}) \cup \alpha(\{(\lambda \times a) \lor \cdot a[x \leftarrow v] \cdot \sigma \in S^+\}) \cup \alpha(\{(\lambda \times a) \lor \cdot a[x \leftarrow v] \cdot \sigma \in S^+\})$	·
$v \cdot \sigma v \in V \land a[x \leftarrow v] \cdot \sigma \in S^{\omega} \}$) $\{S = S^+ \cup S^{\omega} \text{ and } \alpha \text{ preserves lu} \}$	ds Ç
= $\{\langle (\mathbf{A} \times \mathbf{a}) \mathbf{v}, r \rangle \mid \mathbf{v} \in V \land \langle \mathbf{a} \mathbf{x} \leftarrow \mathbf{v} , r \rangle \in \alpha(S)^+ \} \cup \{\langle (\mathbf{A} \times \mathbf{a}) \mathbf{v}, \perp \rangle \mid V \land \langle \mathbf{a} \mathbf{x} \leftarrow \mathbf{v} , \perp \rangle \in \alpha(S)^+ \}$	v ∈
$ (\operatorname{def.} = \{ \langle (\lambda \mathbf{x} \cdot \mathbf{a}) \mathbf{v}, r \rangle \mid \mathbf{v} \in V \land \langle \mathbf{a} [\mathbf{x} \leftarrow \mathbf{v}], r \rangle \in \alpha(S) \} $	αζ
$\{\text{def. } T^+ \triangleq T \cap (\mathbb{T} \times \mathbb{T}) \text{ and } T^\omega \triangleq T \cap (\mathbb{T} \times \{\bot\}) \}$	2({
19	

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= \{\langle (\sigma_0 | \mathbf{b}), \perp \rangle \mid \sigma \in S^{\omega} \}
                                                                                                                                                                                                                                                           7def. α and @ (
                                                                                                                                                                                                                              7S \subseteq \mathbb{T}^{\infty} \text{ so } \sigma_0 \in \mathbb{T}^{n}
      = \{\langle (a \ b), \ \bot \rangle \mid \langle a, \ \bot \rangle \in \alpha(S) \}
      = \alpha(\{(\sigma \otimes \mathbf{b}) \cdot (\mathbf{v} \ \mathbf{b}) \cdot \sigma' \mid \sigma \cdot \mathbf{v} \in S^+ \wedge \mathbf{v} \in V \wedge (\mathbf{v} \ \mathbf{b}) \cdot \sigma' \in S\})
    = \begin{array}{l} \alpha(\{(\sigma @ \mathbf{b}).(\mathbf{v}.\mathbf{b}).\sigma' \mid \sigma.\mathbf{v} \in S^+ \land \mathbf{v} \in \mathbb{V} \land (\mathbf{v}.\mathbf{b}).\sigma' \in S^+\}) \cup \alpha(\{(\sigma @ \mathbf{b}).(\mathbf{v}.\mathbf{b}).\sigma' \mid \sigma.\mathbf{v} \in S^+ \land \mathbf{v} \in \mathbb{V} \land (\mathbf{v}.\mathbf{b}).\sigma' \in S^-\}) \\ (S = S^+ \cup S^- \text{ and } \alpha \text{ preserves lubs}) \end{array}
      = \{ ((\sigma_0 \mathsf{\,b}), r) \mid \sigma \cdot \mathsf{v} \in S^+ \land \mathsf{v} \in \forall \land (\langle \mathsf{v} \mathsf{\,b}), r \rangle \in \alpha(S)^+ \} \cup \{ (\langle \sigma \mathsf{\,b}), \bot \rangle \mid \sigma \cdot \mathsf{v} \in S^+ \land \mathsf{v} \in \forall \land (\langle \mathsf{v} \mathsf{\,b}), \bot \rangle \mid \sigma \cdot \mathsf{v} \in S^+ \land \mathsf{v} \in \forall \land (\langle \mathsf{v} \mathsf{\,b}), \bot \rangle \in \alpha(S)^+ \} 
      = \{ \langle (\sigma_0 \mathsf{b}), r \rangle \mid \langle \sigma_0, \mathsf{v} \rangle \in \alpha(S)^+ \land \mathsf{v} \in V \land \langle (\mathsf{v} \mathsf{b}), r \rangle \in \alpha(S) \}
                                                                                              \{\text{def. } T^+ \triangleq T \cap (\mathbb{T} \times \mathbb{T}), T^{\omega} \triangleq T \cap (\mathbb{T} \times \{\bot\}), \text{ and } \alpha\}
    = \quad \{ \langle (\mathsf{a} \; \mathsf{b}), \; r \rangle \; | \; \langle \mathsf{a}, \; \mathsf{v} \rangle \in \alpha(S)^+ \wedge \mathsf{v} \in \mathbb{V} \wedge \langle (\mathsf{v} \; \mathsf{b}), \; r \rangle \in \alpha(S) \}
    — α({a@σ | a ∈ V ∧ σ ∈ S<sup>ω</sup>})
    = \{\langle (a \sigma_0), \perp \rangle \mid a \in V \land \sigma \in S^{\omega} \}
        = \{\langle (\mathbf{a} \ \sigma_0), \ \bot \rangle \mid \mathbf{a} \in \mathbb{V} \land \langle \sigma_0, \ \bot \rangle \in \alpha(S) \} (def. \alpha and T^{\omega} \triangleq T \cap (\mathbb{T} \cup \{\bot\}))
    = \{((a b), \bot) \mid a \in V \land (b, \bot) \in \alpha(S)\}  ?S \subseteq \mathbb{T}^{\infty} \text{ so } \sigma_0 \in \mathbb{T}^{\vee}\}
      \alpha(\{(a@\sigma), (a v), \sigma' | a, v \in V \land \sigma, v \in S^+ \land (a v), \sigma' \in S\})
      = \alpha(\{(\mathbf{a} @ \sigma) \cdot (\mathbf{a} \ \mathbf{v}) \cdot \sigma' \mid \mathbf{a}, \mathbf{v} \in \mathbb{V} \land \sigma \cdot \mathbf{v} \in S^+ \land (\mathbf{a} \ \mathbf{v}) \cdot \sigma' \in S^+\}) \cup \alpha(\{(\mathbf{a} @ \sigma) \cdot (\mathbf{a} \ \mathbf{v}) \cdot \sigma' \mid \mathbf{a}, \mathbf{v} \in \mathbb{V} \land \sigma \cdot \mathbf{v} \in S^+ \land (\mathbf{a} \ \mathbf{v}) \cdot \sigma' \in S^+\}) \\ (S = S^+ \cup S^- \text{ and } \alpha \text{ preserves lubs})
      = \{ ((\mathbf{a} \ \sigma_0), \ r) \mid \mathbf{a}, \mathbf{v} \in \forall \land (\sigma_0, \mathbf{v}) \in \alpha(S)^+ \land ((\mathbf{a} \ \mathbf{v}), \ r) \in \alpha(S)^+ \} \cup \{ ((\mathbf{a} \ \sigma_0), \ \mathbf{v}) \in \alpha(S)^+ \land ((\mathbf{a} \ \mathbf{v}), \ \bot) \in \alpha(S)^+ \} \cup \{ (\mathbf{a} \ \sigma_0), \ \mathbf{v}) \in \alpha(S)^+ \land ((\mathbf{a} \ \mathbf{v}), \ \bot) \in \alpha(S)^+ \}
      = \{\langle (a b), r \rangle \mid a, v \in V \land \langle b, v \rangle \in \alpha(S) \land \langle (a v), r \rangle \in \alpha(S) \}
                                                                                                                T^{\omega} \triangleq T \cap (\mathbb{T} \cup \{\bot\}) \text{ and } S \subseteq \mathbb{T}^{\infty} \text{ so } \sigma_0 \in \mathbb{T}
      Hence, we have the commutation property \alpha(\vec{F}(S)) = \vec{F}(\alpha(S)) when defining
      Theorem 11 \widetilde{S} \triangleq \alpha(\widetilde{S}) = \alpha(y_p^{\text{\tiny C}} \widetilde{F}) = y_p^{\text{\tiny C}} \widetilde{F}.

PROOF By the fixpoint fusion theorem [7, Th. 9] and the asynchronous fix-
        PROOF By the Exponen use on incorem [7, 11. 3] and the asymmetric point iteration theorem [8, Th. 3.3.10] for \overline{S}^c, the fixpoint definition of \overline{S} can be written in the form (S^+\triangleq S\cap (\mathbb{T}\times\mathbb{T}), S^o\triangleq S\cap (\mathbb{T}\times\{\bot\}) so S^+\cap S^o=\varnothing)
                          \begin{cases} \widetilde{\mathbb{S}} &= \widetilde{\mathbb{S}}^+ \cup \widetilde{\mathbb{S}}^{\omega} \\ \widetilde{\mathbb{S}}^+ &= \widetilde{F} \left( \widetilde{\mathbb{S}}^+ \right) = \operatorname{lfp}^{\mathbb{S}} \widetilde{F}^+ & \text{where} \quad \widetilde{F}^+(S) \triangleq \widetilde{F} \left( S^+ \right) \end{cases}
                            \mathbb{S}^{\omega} = F'(\mathbb{S}^+ \cup \mathbb{S}^{\omega}) = \mathbf{gfp}^{\mathbb{C}} \, F^{\omega} \quad \text{where} \quad F^{\omega}(S) \triangleq F'(\mathbb{S}^+ \cup S^{\omega}) \; .
      We have \alpha(\vec{S}) = \alpha(\vec{S}^+ \cup \vec{S}^\omega) = \alpha(\vec{S}^+) \cup \alpha(\vec{S}^\omega) and prove that \alpha(\vec{S}^+) = \vec{S}^+ and \alpha(\vec{S}^-) = \vec{S}^- so \alpha(\vec{S}) = \vec{S}^+ \cup \vec{S}^\omega = \vec{S}.
  To prove that \alpha(\tilde{\mathbb{S}}^+) = \alpha(\mathbf{H}_0^+, \tilde{F}^+) is equal to \mathbf{H}_0^+, \tilde{F}^+ = \tilde{\mathbb{S}}^+, we observe that a preserves \cup and \alpha \circ \tilde{F}^+ = \tilde{F}^+ \circ \alpha by Lem. 10 so \alpha(\mathbf{H}_0^+, \tilde{F}^+) = \mathbf{H}_0^+, \tilde{F}^+ by [7, \text{ Th. } 3].
        We must prove that \alpha(\vec{\mathbb{S}}^\omega) = \alpha(\mathbf{gfp}^{\subseteq} \vec{F}^\omega) is equal to \mathbf{gfp}^{\subseteq} \vec{F}^\omega = \vec{\mathbb{S}}^\omega.
we must prove that \alpha(S)^{-1} = \alpha(g_0^{-1} F)^{-1} is equal to g_0^{-1} F = 2G_0^{-1} F. To prove that \alpha(g_0^{-1} F)^{-1} = g_0^{-1} F^{-1} when k^{-1} + G = 0 and k^{-1} + G = 0 be the respective transfinite interates of \tilde{F}^{-1} and \tilde{F}^{-1} from X^0 = T^{-1} and \tilde{X}^1 = T^{-1} (4) by that A(G_0^{-1} G)^{-1} = 0 when (G_0^{-1} G)^{-1} = 0 and (G_0^{-1} G)^{-1} = 0
           To prove that gfp<sup>□</sup> F̄<sup>□</sup> ⊆ α(gfp<sup>□</sup> F̄<sup>□</sup>), we show that ∀(a, ⊥) ∈ gfp<sup>□</sup> F̄<sup>□</sup>
        \exists \sigma \in gfp^{\mathbb{S}} \vec{F}^{\omega} : \sigma_0 = a. To do so for any (a, \perp) \in gfp^{\mathbb{S}} \vec{F}^{\omega}, we prove by transfinite induction on \delta that
                          \forall \delta \in \mathcal{O} > 0: \forall \langle \mathsf{a}, \ \bot \rangle \in \mathsf{gfp}^{^{\mathbb{G}}} \, \overrightarrow{F}^{^{\omega}}: \exists \sigma \in \mathbb{T}^{\omega}: \sigma_0 = \mathsf{a} \wedge \sigma \in \bigcap X^{\beta} \, .
      For \delta=1,\,\bigcap_{\beta<\delta}X^\beta=X^0=\mathbb{T}^\omega and \mathbf{a}\in\mathbb{T}.
      Assume by induction hypothesis, that \exists \sigma \in \Gamma^\omega : \sigma_0 = a \wedge \forall \eta \in \Omega : 0 < \eta < \delta : \sigma \in \bigcap_{1 \leq c_0} X^\beta. We have \sigma \in \bigcap_{1 \leq c_0} \bigcap_{1 \leq c_0} X^\beta = \bigcap_{1 \leq c_0} X^\beta : dv we must show that \exists \sigma \in \Gamma^\omega : \sigma_0 = a \wedge \sigma \in X^\delta = \bar{F}^\omega \cap \bigcap_{1 \leq c_0} X^\beta. Because the iterates X^\delta, \delta \in O are decreasing, this implies \exists \sigma \in \Gamma^\omega : \sigma_0 = a \wedge \sigma \in \bigcap_{1 \leq c_0} X^\beta.
    It remains to show, by structural case analysis on a, that if \sigma \in S: \sigma_0 = a.
```

```
then \exists x' \in \tilde{F}(S): x'_0 = a where S = \bigcap_{l \in J} X^l.

If a \in V then (a, 1) \notin \mathfrak{gh}^{-1} \tilde{F}^{-1}.

If a \in V then (a, 1) \notin \mathfrak{gh}^{-1} \tilde{F}^{-1}, where a \in V is the a \in V then a \in V is a \in V.

If a = (A \times x^2) \vee_v \in V then (a, 1) \in \mathfrak{gh}^{-1} \tilde{F}^{-1} = \tilde{F}^{-1} (\mathfrak{gh}^{-1} \tilde{F}^{-1}) so by (6), (x'_1|x_1 - y_1) = y_1^2 + y_2^2 + y_1^2 + y_2^2 + y_2
```

The big-step bifinitary relational semantics \Rightarrow is defined as $\mathbf{a}\Rightarrow r\triangleq \langle \mathbf{a},r\rangle\in\alpha(\mathbb{S}[\mathbf{a}])$ where $\mathbf{a}\in\mathbb{T}$ and $r\in\mathbb{T}\cup\{\bot\}$. It is

5.4.4 Rule-based biq-step bifinitary relational semantics

$$\begin{array}{ll} v \Longrightarrow V, \quad v \in \mathbb{V} & & \dfrac{a[x \leftarrow V] \Longrightarrow r}{(\lambda x \cdot a)} :_{V} \quad v \in \mathbb{V}, \ r \in \mathbb{V} \cup \{\bot\} \\ \\ \dfrac{a \Longrightarrow \bot}{a \ b \Longrightarrow \bot} c & \dfrac{a \Longrightarrow V, \quad v \ b \Longrightarrow r}{a \ b \Longrightarrow r} :_{V} \quad v \in \mathbb{V}, \ r \in \mathbb{V} \cup \{\bot\} \end{array}$$

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$$\cfrac{b \Rightarrow \bot}{a \ b \Rightarrow \bot} c, \quad a \in V \quad \cfrac{b \Rightarrow v, \quad a \ v \Rightarrow r}{a \ b \Rightarrow r} c, \quad a \in V, \ v \in V, \ r \in V \cup \{\bot\} \ .$$

Again this should neither be understood as a structural induction (since $2|x-v|\neq (\lambda x\cdot z)\cdot v$) nor as action induction (because of infinite behavios). The abstraction of 10^{-2} Tr (1 Tr) yields the classical antana issmatics |1| (where all rules with \bot are eliminated and \Box becomes \Box in the remaining ones). The abstraction of $|1\rangle^2$ Tr ((± 1)) yields the divergence semantics (keeping only the rules with \bot , \Box is \bigcirc , and $z\to \bot$ is written a $\overset{\infty}{\Longrightarrow}$ in [18].

Observe that both the maximal trace semantics of Sec. 5.3.1 and the above bifinitary relational semantics of Sec. 5.4 define the semantics of a term that "goes wrong" as empty.

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5. An Example of Practical Application: A Demo of Astrée

6. Long-Term Research Program

Objectives²

- a list of problems on which progress is necessary
- provides a flavor of our general research directions
- hard problems, difficult to predict if and when solutions will be found
- ambitious objectives are necessary for stimulation and progress
- long term / short term objectives will be considered in parallel

² Project membership dependant!

Abstract Formalization of Computations

- semantics: for real-life languages
- abstract properties and specifications: safety, liveness,
 security, probabilistic behaviors . . . and beyond
- time abstraction: continuous to discrete, scheduling, performance properties

Abstraction of Computational Paradigms

- abstraction of data structures
- abstraction of control structures: imperative, functional, procedural, logical, synchronous, parallel, distributed, and mobile control paradigms
- abstraction of program structures: procedures, modules, objects, classes, ...
- abstraction of communication and cooperation structures: synchronous/asynchronous lossy/lossless channels, events, semaphores, mobile communications, . . .

- abstraction of hardware structures: memory caches,
 pipelines, branch prediction . . . at the assembler level,
 hardware description languages
- abstraction of biological systems: abstraction of agentbased descriptions of biological systems

Abstraction Validation

- abstraction translation: translation of abstractions while translating models
- verified abstractions: beyond toy examples

Abstraction Automatization

- imprecision localization: origin of false alarms
- automatic refinement: automatic design of abstract domains to eliminate false alarms
- automatic abstraction: too precise abstractions are costly

7. The Research Program for the 4 Next Years

Objectives

- a list of problems on which, thanks to our past experience, progress is expected in the short/mid term
- hard problems, difficult to predict if the proposed solutions will scale-up
- ambitious program, should find end-users
- strongly project membership dependant!



Software Verification with no False Alarm³

- industrialization of ASTRÉE for synchronous programs
 (2/3 years)
- extension of the scope of sequential analyzes (data structures, separate analyzes?), including translation validation (2/4 years)
- universal libraries for numerical/symbolic abstract domains (2/4 years)

³ Strongly project membership dependant!

Analysis of Parallel Applications 4

- foundations (and prototype?) for analyzing quasi-synchronous programs (3/4 years)
- foundations and prototype for analyzing asynchronous
 programs (4/6 years)⁵

⁴ Strongly project membership dependant!

⁵ ASTRÉE started Nov. 2001!

Verification of Security Protocols⁶

- development of an effective cryptographic protocol certifier in the computational model (3/4 years)

⁶ Strongly project membership dependant!

8. Current Projects of the Team

International & European Projects

- JST France-Japan
- ESA ITI "Space software validation using abstract interpretation" (2007–2008)⁷
- ITEA 2 ES_PASS "Embedded Software Product-based ASSurance" (2007–2009)⁸

Academic partners: École Normale Supérieure, CNRS FéRIA federation (INPT-IRIT and ONERA-DTIM), Saarland University, Technical University of Munich, Tel-Aviv University, Universidad Politécnica de Madrid; Industrial partners: AbsInt GmbH, Airbus France, CS Systèmes d'Information, Esterel Technologies, PolySpace Technologies, Thales Avionics, . . .



⁷ Astrium Space Transportation (David Lesens), the CEA LIST (Éric Goubault, Coordinator), the École Normale Supérieure (Patrick Cousot), and the École Polytechnique (Radhia Cousot), in order to verify safety properties of a C version of the Monitoring and Safing Unit (MSU) criticality level A software of the Automated Transfer Vehicle (ATV)

Institutional Projects

- APRON⁹ (2005–2008): numerical public-domain abstract domain library
- ANR/ARA SSIA ¹⁰/CONTROVERT ¹¹: analyze a full control-command system from its mathematical design to its computer-based implementation

¹¹ École normale supérieure (Patrick Cousot, coordinator), the CNRS (Radhia Cousot), the ONERA (Pierre Apkarian), and the University Paul Sabatier of Toulouse (Dominikus Noll).



⁹ CRI/ENSMP (François IRIGOIN, coordinator), the École Normale Supérieure (Patrick Cousot), the École Polytechnique (Radhia Cousot), Vérimag (Nicolas Halbwachs), and INRIA Alpes (Bertrand Jeannet)

¹⁰ Sécurité des Systèmes embarqués & Intelligence Ambiante

Institutional Projects (Cont'd)

- ANR/ARA SSIA/FORMACRYPT¹² (2005–2008): convergence of the computational formal and Dolev-Yao models
- RNTL/THESÉE ¹³ (2005–2008): analysis of asynchronous (control/command and communication) software

¹² École Normale Supérieure (Bruno BLANCHET, coordinator), the École Normale Supérieure de Cachan (Jean Goubault-Larrecq), and the LORIA (Véronique Cortier

¹³ École Normale Supérieure (Patrick Cousot, coordinator), the CNRS (Radhia Cousot), EDF (Alain Ourghanlian) and Airbus France (Jean Souyris

Industrial Projects

- ASBAPROD (Assurance Basée Produit, 2005–2008) 14
- aims at "introducing abstract-interpretation-based verification methods, technologies and tools to master the development of avionic embedded synchronous and asynchronous software".

¹⁴ École Normale Supérieure (Patrick Cousot) and Airbus France EYY (Jean Souyris)



9. Technology Transfer

Research/Development/Transfer Cycles

Three overlapping activities on each software development:

- 1. Fundamental research and experiments (2 years)
- 2. Prototypes development and validation (2 years)
- 3. Industrial transfer (2 years).
- \Rightarrow simultaneously pursue three activities, each one in a different phase.

Current situation

- Fundamental research: analysis of quasi-synchronous systems
- Prototype development: analysis of asynchronous programs
- Industrial transfer: ASTRÉE

10. Necessary Means

Necessary Means

- Stabilizing brilliant young researchers:
 - 3 "chargé de recherche 2" or "maître de conférences"
 - 1 "directeur de recherche" or "professor"
- Software development and technological transfer support:
 - 1 "Research Engineer" ¹⁵ (software industrialisation, contribution to new software developments)
- Administrative support:
 - 1 "Project assistant".

¹⁵ in the context of an ODL Opérations de développement logiciel (Software Development Support)?



11. Conclusion

Objectives of the creation of Abstraction

- an internationally recognized research team;
- ensure the durability of the investment on the static analysis of synchronous programs for control/command (ASTRÉE);
- support the technological transfer of ASTRÉE to the industry;
- support the development of new analysis and verification techniques for asynchronous applications;
- support the development of abstract interpretation theory and practice in the long-term.



THE END, THANK YOU



12. Publications by the Project Members

Publications of the project members between 2002 and 2006 ¹⁶.

Theses

- [2] L. Mauborgne. Analyse statique et domaines abstraits symboliques. ThËse, Mémoire d'habilitation à diriger les recherches en informatique, Université de Paris Dauphine, 12 February 2007.
- [3] A. Miné. Domaines numériques abstraits faiblement relationnels. Thèse de doctorat en informatique, École polytechnique, Palaiseau, France, 6 December 2004.
- [4] J. Feret. Analyse de systèmes mobiles par interprétation abstraite. Thèse de doctorat en informatique, École polytechnique, Palaiseau, France, 25 February 2005.
- [5] X. Rival. Analyse statique et transformations de programmes dans le cadre de l'interprétation abstraite. Thèse de doctorat en informatique, École polytechnique, Palaiseau, France, 21 October 2005.

The titles of the publications are clickable references to their web location, whenever available.



Invited Book Chapters

- [6] B. Blanchet, P. Cousot, R. Cousot, J. Feret, L. Mauborgne, A. Miné, D. Monniaux and X. Rival. Design and Implementation of a Special-Purpose Static Program Analyzer for Safety-Critical Real-Time Embedded Software, invited chapter. In: The Essence of Computation: Complexity, Analysis, Transformation. Essays Dedicated to Neil D. Jones, edited by T. Mogensen, D. Schmidt and I. Sudborough, pp. 85–108. Springer, Berlin, Germany, 2002, Lecture Notes in Computer Science 2566.
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- [85] P. Cousot. Automated Verification of Infinite-State Systems by Abstract Interpretation, invited talk. In: Third International Workshop on Automated Verification of Infinite-State Systems (AVIS'04), Barcelona, Spain, 3–4 April 2004.



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- [87] P. Cousot. A Lagrangian relaxation and mathematical programming framework for static analysis and verification, invited talk. In: International Symposium on Static Analysis, SAS '04 & on Logic Program Synthesis and Transformation, LOPSTR '04, Verona, Italy, 28 August 2004.
- [88] P. Cousot. Software Verification by Abstract Interpretation: Current Trends and Perspectives, invited talk. In: IV Jornadas de Programación y Lenguajes, Málaga, Spain, 11–12 November 2004.
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- [103] B. Blanchet. Automated security proofs with sequences of games. In: Seminar, Université de Caen, October 2006. (joint work with D. Pointcheval).
- [104] B. Blanchet. Automated verification of selected equivalences for security protocols. In: Seminar, PPS, Université de Paris VII, January 2006. – (joint work with M. Abadi and C. Fournet).
- [105] B. Blanchet. An automatic security protocol verifier based on resolution theorem proving. In: Seminar, IRMAR, Rennes, January 2006.

- [106] B. Blanchet. A computationally sound mechanized prover for cryptographic protocols. In: Cryptography Seminar, École Normale Supérieure, January 2006.
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- [108] P. Cousot. Abstract Interpretation & Applications. In: AA & EECS Seminar, MIT, Cambridge, Massachusetts, United States, 3 April 2006.
- [109] P. Cousot. Application of Abstract Interpretation to the Static Verification of Safety Critical Code. In: Seminar, IBM Thomas J. Watson Research Center, Hawthorne, New York, United States, 20 January 2006.
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- [111] P. Cousot. Program Termination Proofs by Parametric Abstraction, Lagrangian Relaxation and Semi-Definite Programming. In: Specialised Talk, Seminar Series, Department of Computing and Information Sciences, Kansas State University, Manhattan, Kansas, United States, 6 september 2006.
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