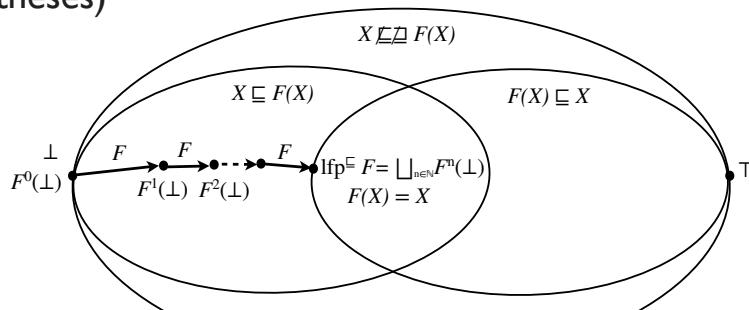


Abstracting Induction by Extrapolation and Interpolation

Mumbai, India
January 12th, 2015

Patrick Cousot

pcousot@cs.nyu.edu cs.nyu.edu/~pcousot



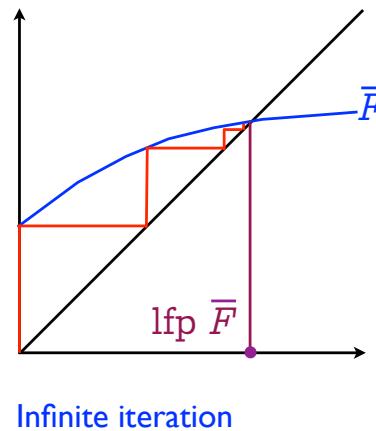
Abstract Interpreters

- **Transitional abstract interpreters:** proceed by induction on program steps
- **Structural abstract interpreters:** proceed by induction on the program syntax
- **Common main problem:** over/under-approximate fixpoints in non-Noetherian^(*) abstract domains^(**)

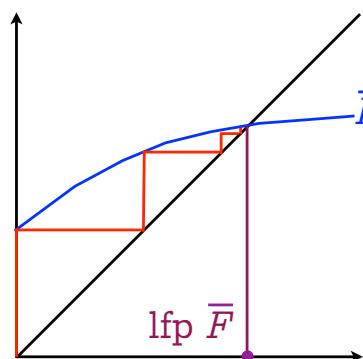
(*) Iterative fixpoint computations may not converge in finitely many steps

(**) Or convergence may be guaranteed but slow.

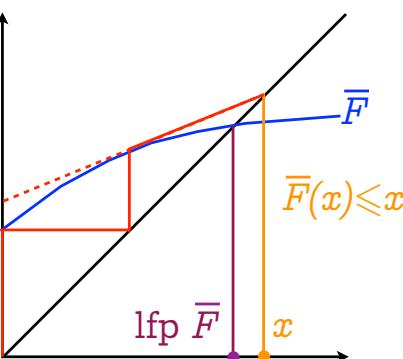
Convergence acceleration with widening



Convergence acceleration with widening



Infinite iteration



Accelerated iteration with widening
(e.g. with a widening based on the derivative
as in Newton-Raphson method^(*))

^(*) Javier Esparza, Stefan Kiefer, Michael Luttenberger: Newtonian program analysis. J. ACM 57(6): 33 (2010)

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The oldest widenings

- Primitive widening [1,2]

$(x \bar{v} y) = \text{cas } x \in V_a, y \in V_b \text{ dans}$
 $\begin{cases} \square, ? \Rightarrow y \downarrow \\ ?, \square \Rightarrow x \downarrow \\ [a_1, a_2], [b_1, b_2] \Rightarrow \\ \quad \text{[if } a_2 < a_1 \text{ then } -\infty \text{ else } a_1 \text{ fi}, \\ \quad \text{[if } b_2 > b_1 \text{ then } +\infty \text{ else } b_1 \text{ fi]} \\ \text{fincas;} \end{cases}$

$[a_1, b_1] \bar{v} [a_2, b_2] =$
 $\begin{array}{c} \text{[if } a_2 < a_1 \text{ then } -\infty \text{ else } a_1 \text{ fi}, \\ \text{[if } b_2 > b_1 \text{ then } +\infty \text{ else } b_1 \text{ fi]} \end{array}$

- Widening with thresholds [3]

$\forall x \in \bar{L}_2, \perp \nabla_2(j) x = x \nabla_2(j) \perp = x$
 $[l_1, u_1] \nabla_2(j) [l_2, u_2]$
 $= [\text{if } 0 \leq l_2 < l_1 \text{ then } 0 \text{ elseif } l_2 < l_1 \text{ then } -b - 1 \text{ else } l_1 \text{ fi},$
 $\quad \text{if } u_1 < u_2 \leq 0 \text{ then } 0 \text{ elseif } u_1 < u_2 \text{ then } b \text{ else } u_1 \text{ fi}]$

[1] Patrick Cousot, Radhia Cousot: Vérification statique de la cohérence dynamique des programmes. Rapport du contrat IRIA-SESORI No. 75-032, 23 septembre 1975.

[2] Patrick Cousot, Radhia Cousot: Abstract Interpretation: A Unified Lattice Model for Static Analysis of Programs by Construction or Approximation of Fixpoints. POPL 1977: 238-252

[3] Patrick Cousot, Semantic foundations of program analysis, Ch. 10 of Program flow analysis: theory and practice, N. Jones & S. Muchnick (eds), Prentice Hall, 1981.

Extrapolation by Widening

- $X^0 = \perp$ (increasing iterates with widening)

$$X^{n+1} = X^n \nabla F(X^n) \quad \text{when } F(X^n) \not\subseteq X^n$$

$$X^{n+1} = X^n \quad \text{when } F(X^n) \subseteq X^n$$

- Widening ∇ :

$$Y \sqsubseteq X \nabla Y$$

(extrapolation)

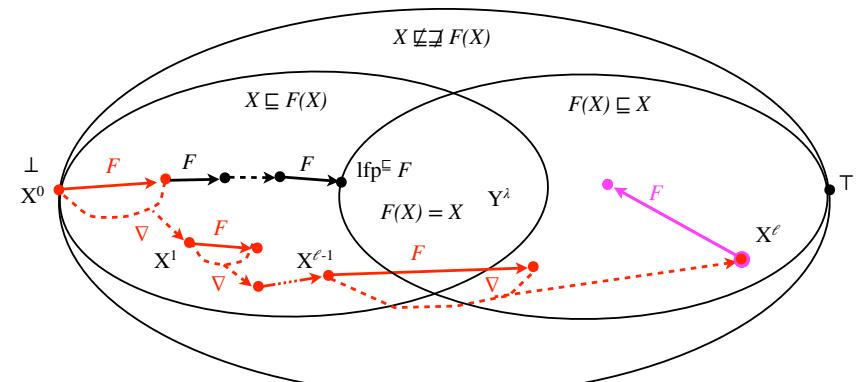
- Enforces convergence of increasing iterates with widening (to a limit X^ℓ)

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Extrapolation with widening



Interpolation with narrowing

- $Y^0 = X^\ell$ (decreasing iterates with narrowing)

$$Y^{n+1} = Y^n \Delta F(Y^n) \quad \text{when } F(Y^n) \sqsubset Y^n$$

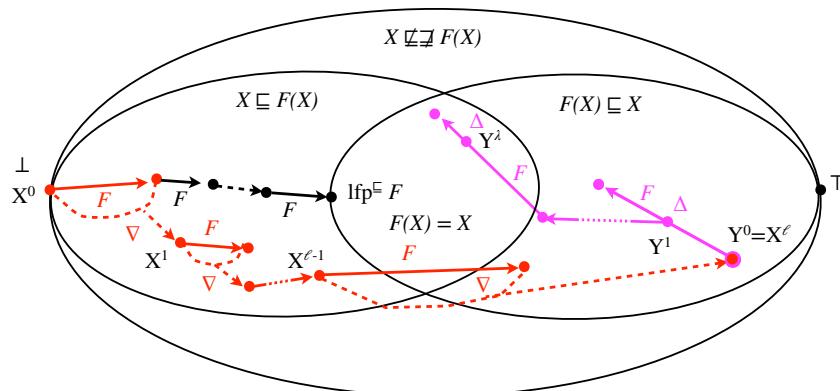
$$Y^{n+1} = Y^n \quad \text{when } F(Y^n) = Y^n$$

- Narrowing Δ :

- $Y \sqsubseteq X \implies Y \sqsubseteq X \Delta Y \sqsubseteq X$ (interpolation)

- Enforces convergence of decreasing iterates with narrowing (to a limit Y^λ)

Interpolation with narrowing



Could stop when $F(X) \not\subseteq X \wedge F(F(X)) \subseteq F(X)$ but not the current practice.

The oldest narrowing

- [2]

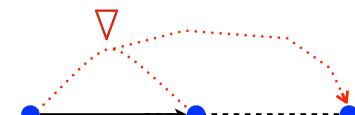
```
[a1,b1] Δ [a2,b2] =
[if a1 = -∞ then a2 else MIN(a1,a2),
 if b1 = +∞ then b2 else MAX(b1,b2)]
```

Duality

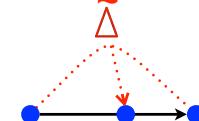
	Convergence above the limit	Convergence below the limit
Increasing iteration	Widening ∇	Dual-narrowing $\tilde{\Delta}$
Decreasing iteration	Narrowing Δ	Dual widening $\tilde{\nabla}$

Extrapolators ($\nabla, \tilde{\nabla}$) and interpolators ($\Delta, \tilde{\Delta}$)

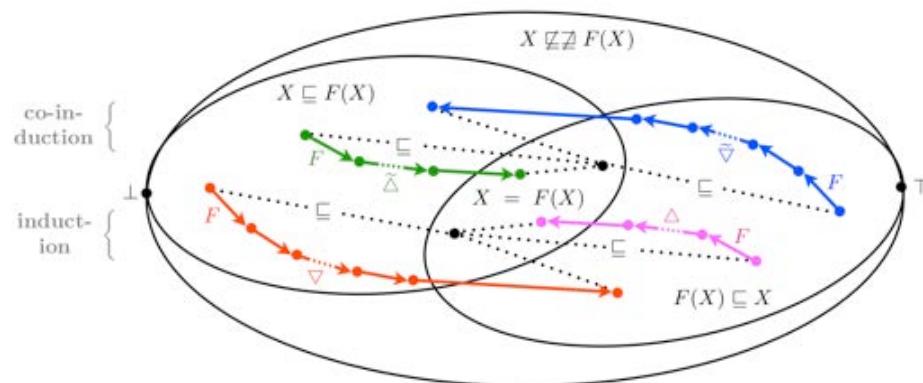
- Extrapolators:



- Interpolators:



Extrapolators, Interpolators, and Duals



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Interpolation with dual narrowing

- $Z^0 = \perp$ (increasing iterates with dual-narrowing)
- $Z^{n+1} = F(Z^n) \tilde{\Delta} Y^\lambda$ when $F(Z^n) \not\subseteq Z^n$
- $Z^{n+1} = Z^n$ when $F(Z^n) \subseteq Z^n$
- Dual-narrowing $\tilde{\Delta}$:
 - $X \sqsubseteq Y \implies X \sqsubseteq X \tilde{\Delta} Y \sqsubseteq Y$ (interpolation)
 - Enforces convergence of increasing iterates with dual-narrowing

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Example of dual-narrowing

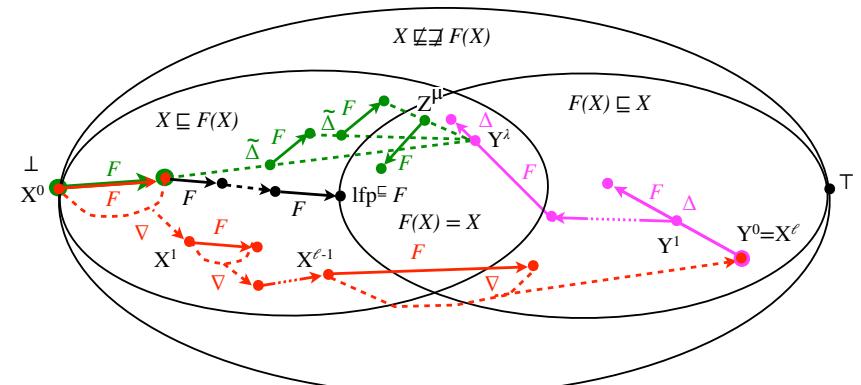
- $[a, b] \tilde{\Delta} [c, d] \triangleq [[c = -\infty \wedge a \leq (a+c)/2], [d = \infty \wedge b \geq (b+d)/2]]$
- The first method we tried in the late 70's with Radhia
 - Slow
 - Does not easily generalize (e.g. to polyhedra)

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Interpolation with dual-narrowing



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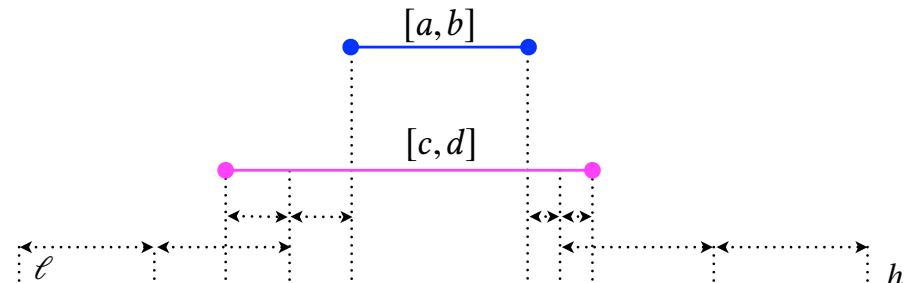
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Relationship between narrowing and dual-narrowing

- $\tilde{\Delta} = \Delta^{-1}$
- $Y \sqsubseteq X \implies Y \sqsubseteq X \Delta Y \sqsubseteq X$ (narrowing)
- $Y \sqsubseteq X \implies Y \sqsubseteq Y \tilde{\Delta} X \sqsubseteq X$ (dual-narrowing)
- Example: Craig interpolation
- Why not use a bounded widening (bounded by B)?
 - $F(X) \sqsubseteq B \implies F(X) \sqsubseteq F(X) \tilde{\Delta} B \sqsubseteq B$ (dual-narrowing)
 - $X \sqsubseteq F(X) \sqsubseteq B \implies F(X) \sqsubseteq X \nabla_B F(X) \sqsubseteq B$ (bounded widening)

Example of widenings (cont'd)

- Bounded widening (in $[\ell, h]$):



$$[a,b] \nabla_{[\ell,h]} [c,d] \triangleq \left[\frac{c+a-2\ell}{2}, \frac{b+d+2h}{2} \right]$$

More in the paper...

Widenings

Widenings are not increasing

- A well-known fact

$$[1,1] \subseteq [1,2] \text{ but } [1,1] \nabla [1,2] = [1,\infty] \subseteq [1,2] \nabla [1,2] = [1,2]$$

- A widening cannot both:

- Be increasing in its first parameter
- Enforce termination of the iterates
- Avoid useless over-approximations as soon as a solution is found^(*)

(*) A counter-example is $x \nabla y = \top$

Soundness

Soundness

- In the paper, the fixpoint approximation soundness theorems are expressed with minimalist hypotheses:
- No need for complete lattices, complete partial orders (CPO's):
 - The concrete domain is a poset
 - The abstract domain is a pre-order
 - The concretization is defined for the abstract iterates only.

Soundness (cont'd)

- No need for increasingness/monotony hypotheses for fixpoint theorems (Tarski, Kleene, etc)
 - The concrete transformer is increasing and the limit of the iterations does exist in the concrete domain
 - No hypotheses on the abstract transformer (no need for fixpoints in the abstract)
 - Soundness hypotheses on the extrapolators/interpolators with respect to the concrete
- In addition, termination hypotheses on the extrapolators/interpolators ensure convergence in finitely many steps

Soundness (cont'd)

- No need for increasingness/monotony hypotheses for fixpoint theorems (Tarski, Kleene, etc)
 - The concrete transformer is increasing and the limit of the iterations does exist in the concrete domain
 - No hypotheses on the abstract transformer (no need for fixpoints in the abstract)
 - Soundness hypotheses on the extrapolators/interpolators with respect to the concrete

Examples of interpolators

Craig interpolation

- Craig interpolation:

Given $P \Rightarrow Q$ find I such that $P \Rightarrow I \Rightarrow Q$ with
 $\text{var}(I) \subseteq \text{var}(P) \cap \text{var}(Q)$

is a **dual narrowing** (already observed by Vijay D'Silva and Leopold Haller as an inversed narrowing)

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is a **dual narrowing** (already observed by Vijay D'Silva and Leopold Haller as an inverted narrowing)

- This evidence looked very controversial to some reviewers
- The generalization of an idea does not diminish in any way the merits and originality of this idea

Conclusion

Conclusion

- Abstract interpretation in infinite domains is traditionally by **iteration with widening/narrowing**.
- We have shown how to use **iteration with dual-narrowing** (alone or after widening/narrowing).
- These ideas of the 70's **generalize Craig interpolation from logic to arbitrary abstract domains**.

The End, Thank You