Reinhard Wilhelm

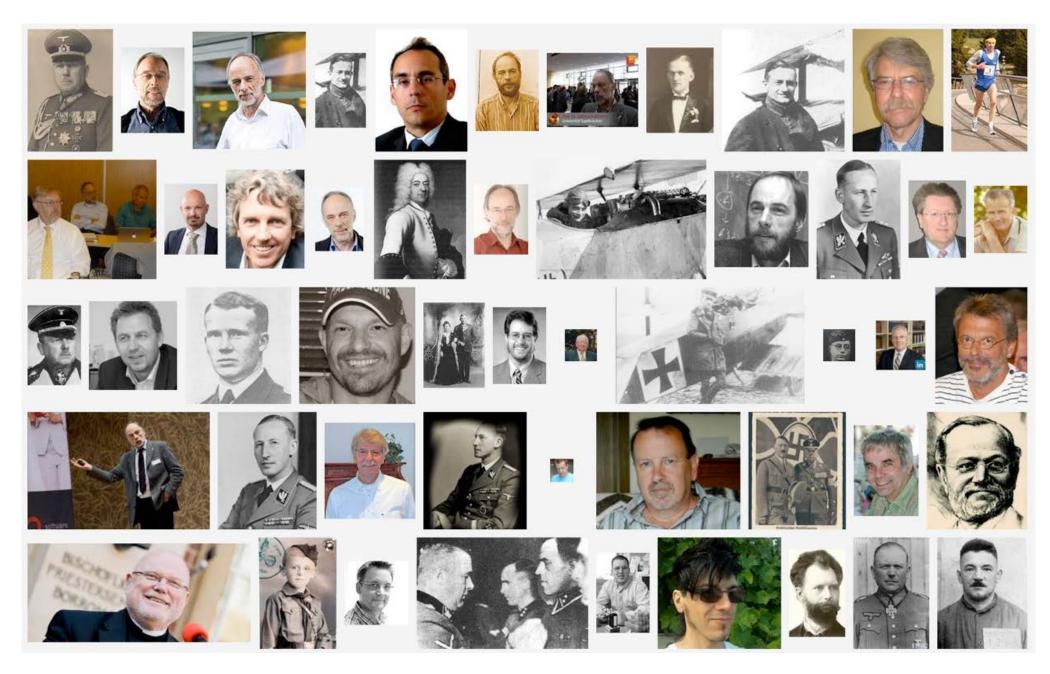
Patrick Cousot

cims.nyu.edu/~pcousot pcousot@cims.nyu.edu

Farewell Colloquium on the Occasion of Reinhard Wilhelm's 68th birthday

Saarbrüchen, November, 28th 2014

You said Reinhard Wilhelm?





You have undestood the limitations of "Big data" and "Advanced machine learning"

You have undestood the limitations of "Big data" and "Advanced machine learning"

2. This is THE Reinhard Wilhelm:

Reinhard Wilhelm

Reinhard Wilhelm is a German computer scientist. Wikipedia

Born: June 5, 1946 (age 68), Finnentrop, Germany

Education: University of Münster



- You have undestood the limitations of "Big data" and "Advanced machine learning"
- 2. This is THE Reinhard Wilhelm:

Reinhard Wilhelm

Reinhard Wilhelm is a German computer scientist. Wikipedia

Born: June 5, 1946 (age 68), Finnentrop, Germany

Education: University of Münster



sorry, this was 2 months ago on Wikipedia, thanks to the true Reinhard Wilhelm for updating his picture last month!

- You have undestood the limitations of "Big data" and "Advanced machine learning"
- 2. This is THE Prof. em. Dr. Dr. h.c. Reinhard Wilhelm:



5 June 1946 (age 68) Born Finnentrop, Germany Fields Computer Scientist Institutions Saarland University Alma mater University of Münster, Stanford University, Technical University Munich Known for compiler technology Konrad Zuse Medal (2009) Notable awards Merit Cross on Ribbon (2010) **ACM Distinguished Service** Award (2011)

There is only one, the proof is by Google

There is only one, the proof is by Google

Images for Prof. em. Dr. Dr. h.c. Reinhard Wilhelm Report images



More images for Prof. em. Dr. Dr. h.c. Reinhard Wilhelm

And more ...



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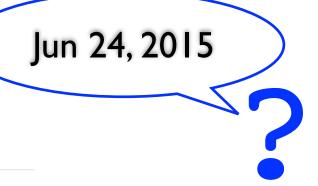
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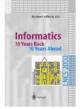
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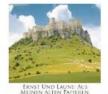
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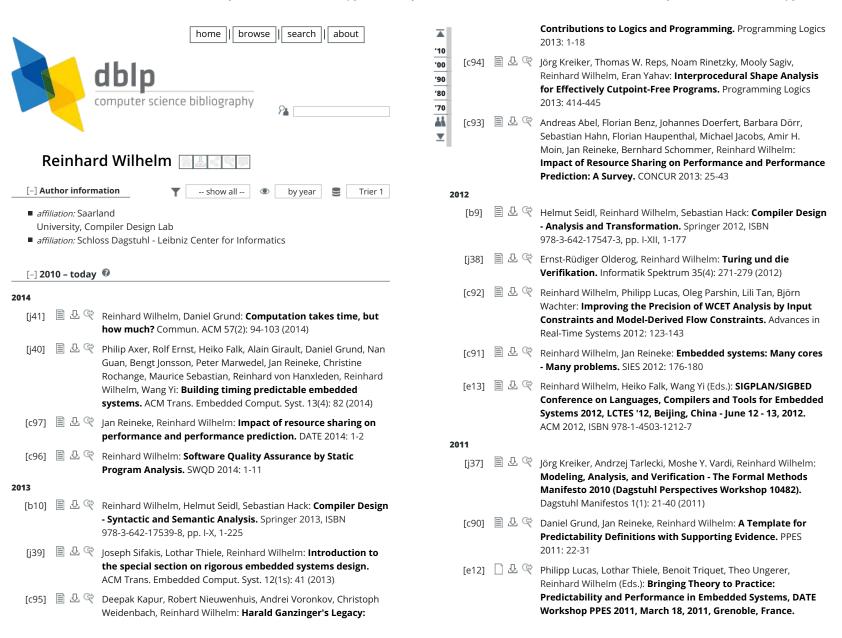


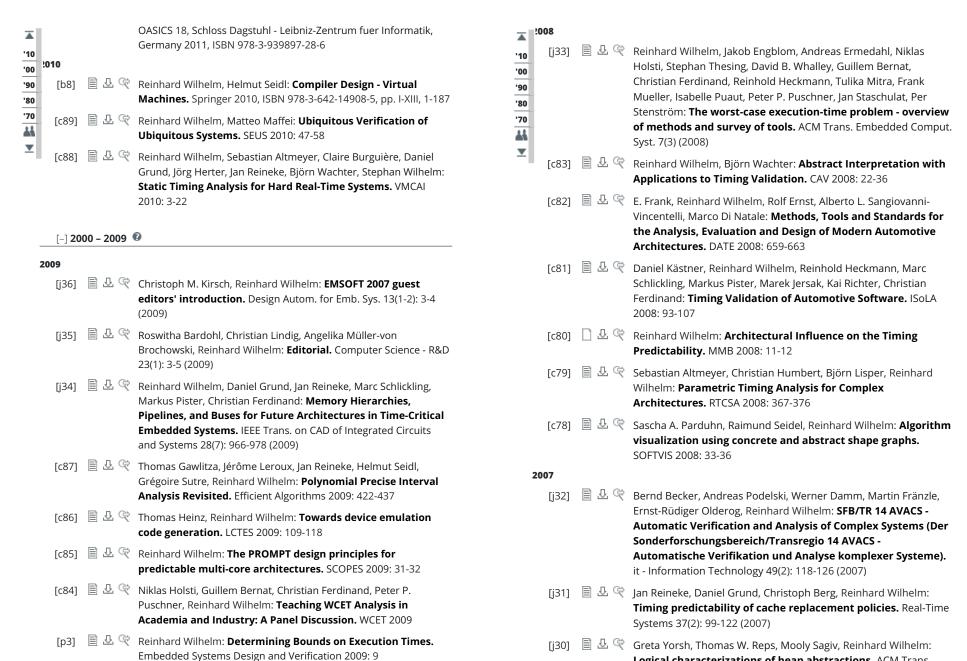
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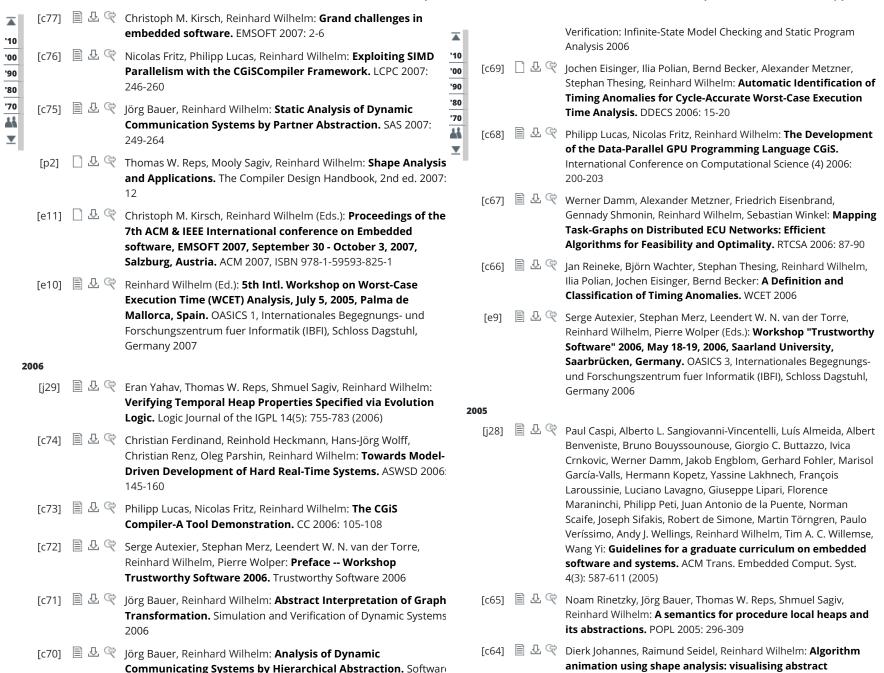
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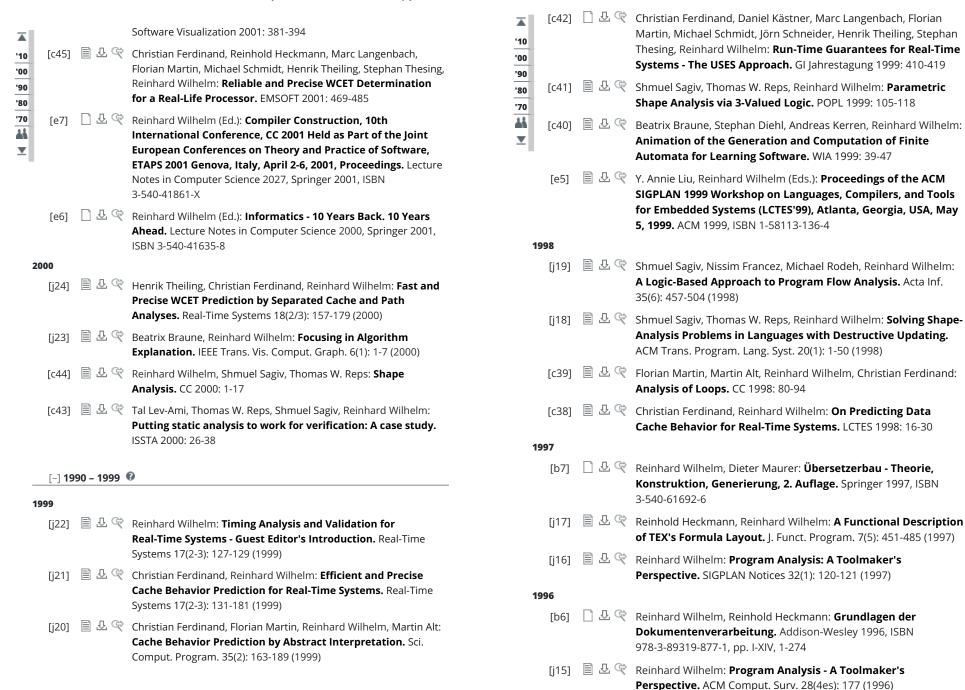
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7 of 16 11/20/14 10:17 AM

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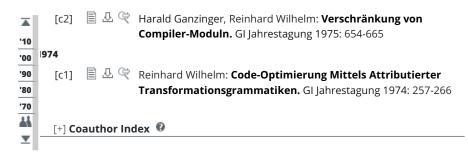
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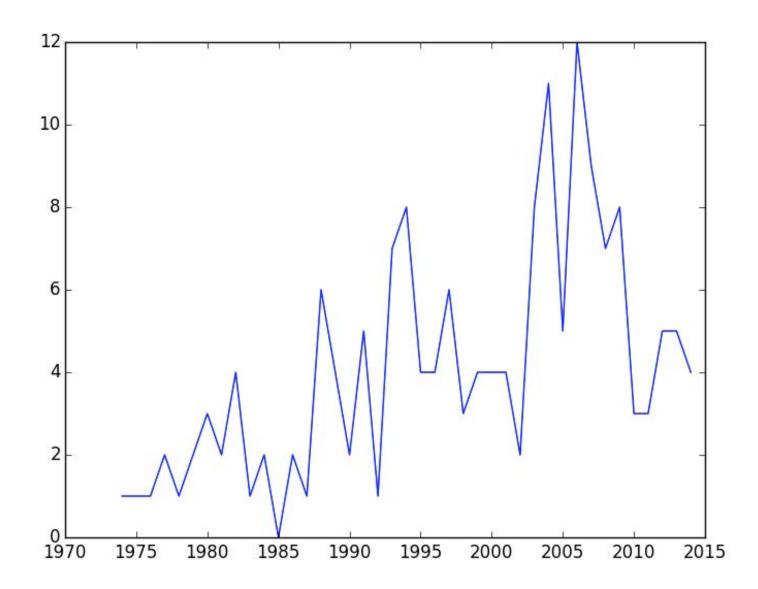


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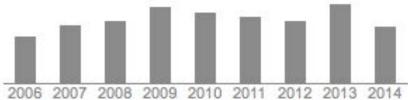
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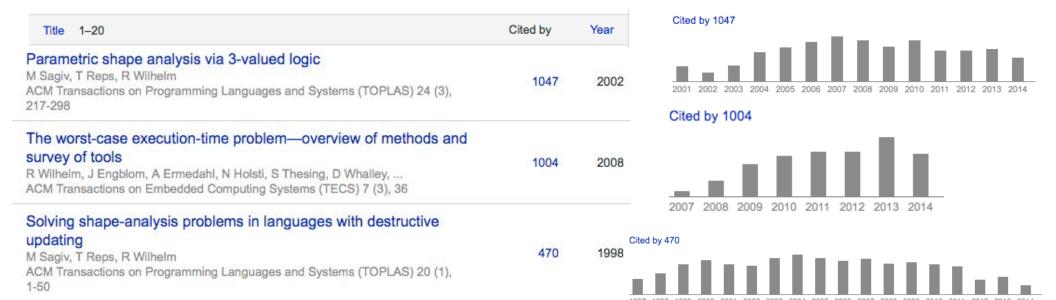


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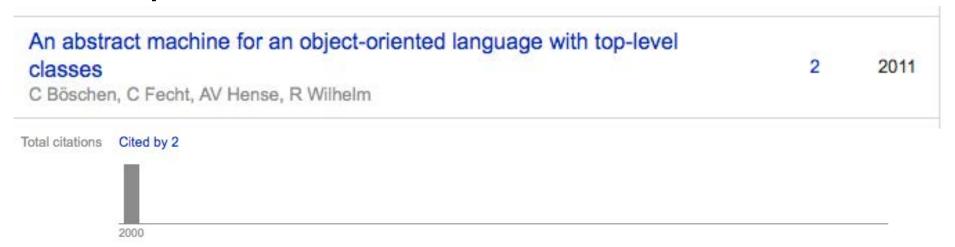


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Science

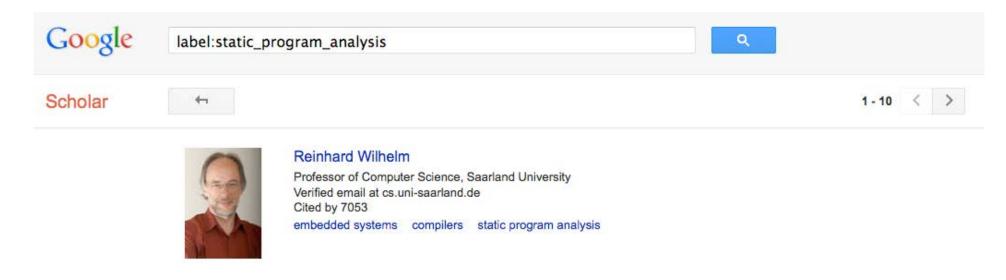
Main contributions



Reinhard Wilhelm

Professor of Computer Science, Saarland University embedded systems, compilers, static program analysis Verified email at cs.uni-saarland.de

Coming number one in static analysis, world-wide:



What is static analysis?

Static program analysis

From Wikipedia, the free encyclopedia

Static program analysis is the analysis of computer software that is performed without actually executing programs (analysis performed on executing programs is known as dynamic analysis).

at least the static analyzer must execute!

by a computer

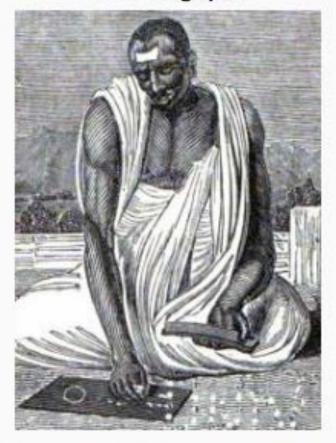
A short introduction to static analysis

The very first static analysis

Brahmagupta (Sanskrit: ब्रह्मगुप्त;

listen (help-info)) (598–c.670 CE) was an Indian mathematician and astronomer who wrote two important works on Mathematics and Astronomy: the *Brāhmasphuṭasiddhānta* (Extensive Treatise of Brahma) (628), a theoretical treatise, and the *Khaṇḍakhādyaka*, a more practical text.

Brahmagupta



Born 598 CE

Died c.670 CE

Fields Mathematics, Astronomy

Known for Zero, modern Number system

18.30. [The sum] of two positives is positives, of two negatives negative;

The abstraction is that you do not (always)
need to known the absolute value of the
arguments to know the sign of the result;

- The abstraction is that you do not (always)
 need to known the absolute value of the
 arguments to know the sign of the result;
- Sometimes imprecise (don't know the sign of the sum of a positive and a negative)

- The abstraction is that you do not (always)
 need to known the absolute value of the
 arguments to know the sign of the result;
- Sometimes imprecise (don't know the sign of the sum of a positive and a negative)
- Useful in practice (if you know what to do when you don't know the sign)

- The abstraction is that you do not (always)
 need to known the absolute value of the
 arguments to know the sign of the result;
- Sometimes imprecise (don't know the sign of the sum of a positive and a negative)
- Useful in practice (if you know what to do when you don't know the sign)
- e.g. in compilation: do not optimize (a division by 2 into a shift when positive*)

^(*) Unless processor uses 2's complement and can shift the sign.

18.30. [The sum] of two positives is positives, of two negatives negative; [...]

18.32. A negative minus zero is negative, a positive [minus zero] positive; zero [minus zero] is zero. When a positive is to be subtracted from a negative or a negative from a positive, then it is to be added.

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18.34. A positive divided by a positive or a negative divided by a negative is positive; a zero divided by a zero is zero; a positive divided by a negative is negative; a negative divided by a positive is [also] negative.

wrong

The rule of signs by Michel Sintzoff (1972)

```
For example, a×a+b×b yields the value 25
when a is 3 and b is -4, and when + and \times are
the arithmetic multiplication and addition.
But axa+bxb yields always the object "pos" when
a and b are the objects "pos" or "neg", and when
the valuation is defined as follows:
                                                                                                                            pos×pos=pos
pos+pos=pos
pos+neg=pos,neg
                                                                                                                         pos×neg=neg
neg+pos=pos,neg
                                                                                                                          neg×pos=neg
                                                                                                                         neg×neg=pos
neg+neg=neg
                                                                                                                            V(p \times q) = V(p) \times V(q)
V(p+q)=V(p)+V(q)
V(0)=V(1)=...=pos
V(-1)=V(-2)=...=neg
The valuation of axa+bxb yields "pos" by the
 following computations:
                                                                                                            V(b)=pos.neg
V(a) = pos, neg
V(a \times a) = pos \times pos, neg \times neg = V(b \times b) = pos \times pos, neg \times neg = v(b \times b) = pos \times pos, neg \times neg = v(b \times b) = pos \times pos, neg \times neg = v(b \times b) = pos \times pos, neg \times neg = v(b \times b) = pos \times pos, neg \times neg = v(b \times b) = v(b
                                                                                                                                                 =pos,pos=pos
                            zpos,poszpos
 V(a\times a+b\times b)=V(a\times a)+V(b\times b)=pos+pos=pos
                        This valuation proves that the result of
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 compute its square root without any preliminary
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The rule of signs by Reinhard Wilhelm (2012/13)

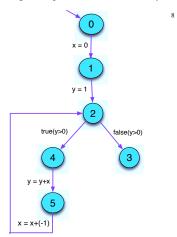
http://www.rw.cdl.uni-saarland.de/teaching/dses I 2/slides/lecture6 static analysis.pdf

2 Example — Rules-of-Sign Analysis

Problem: Determine at each program point the sign of the values of all variables of numeric type.

Example program:

Program representation as control-flow graphs



The analysis should "bind" program variables to elements in Signs.

 $\perp \in \mathbb{D}$ is the function mapping all arguments to $\{\}$.

The partial order on \mathbb{D} is $D_1 \sqsubseteq D_2$ iff

$$D_1 \sqsubseteq D_2$$
 iff

$$o_1 = \bot$$
 o_2

$$D_1 x \subseteq D_2 x \quad (x \in Vars)$$

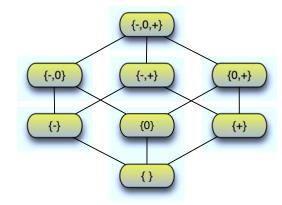
Intuition?

How did we analyze the program?

All the ingredients:

- a set of information elements, each a set of possible signs,
- a partial order, "\subsection", on these elements, specifying the "relative strength" of two information elements,
- these together form the abstract domain, a lattice,
- functions describing how signs of variables change by the execution of a statement, abstract edge effects,
- these need an abstract arithmetic, an arithmetic on signs.

We construct the abstract domain for single variables starting with the lattice $Signs = 2^{\{-,0,+\}}$ with the relation " \square " =" \square ".



So, the abstract domain is $\mathbb{D} = (Vars \rightarrow Signs)_{\perp}$, a Sign-environment.

all order on
$$\mathbb{D}$$
 is $D_1 \subseteq D_2$

$$D_1 = \bot$$

$$D_1 x \subseteq D_2 x \quad (x \in Vars)$$

In particular, how did we walk the lattice for y at program point 5? {-,0,+}

The analysis should "bind" program variables to elements in Signs.

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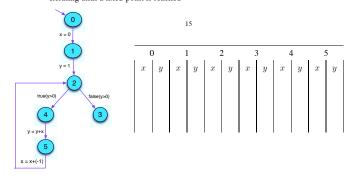
The partial order on \mathbb{D} is $D_1 \sqsubseteq D_2$ iff

$$D_1 = \bot$$
 or

$$D_1 x \subseteq D_2 x \quad (x \in Vars)$$

 D_1 is at least as precise as D_2 since D_2 admits at least as many signs as

How is a solution found? Iterating until a fixed-point is reached



Idea:

- We want to determine the sign of the values of expressions.
- For some sub-expressions, the analysis may yield $\{+, -, 0\}$ which means, it couldn't find out.

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The rule of signs by Reinhard Wilhelm (2012/13)

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- We replace the concrete operators □ working on values by abstract operators □[‡] working on signs:
- The abstract operators allow to define an abstract evaluation of expressions:

$$\llbracket e \rrbracket^{\sharp} : (Vars \to Signs) \to Signs$$

Determining the sign of expressions in a Sign-environment works as follows:

$$\begin{bmatrix} c \end{bmatrix}^{\sharp} D \qquad \stackrel{\mathcal{I}}{=} \qquad \begin{cases} \{+\} & \text{if } c > 0 \\ \{-\} & \text{if } c < 0 \\ \{0\} & \text{if } c = 0 \end{cases} \\
 \begin{bmatrix} v \end{bmatrix}^{\sharp} \qquad = \qquad D(v) \\
 \begin{bmatrix} e_1 \square e_2 \end{bmatrix}^{\sharp} D \qquad = \qquad \begin{bmatrix} e_1 \end{bmatrix}^{\sharp} D \square^{\sharp} \begin{bmatrix} e_2 \end{bmatrix}^{\sharp} D \\
 \begin{bmatrix} \square e \end{bmatrix}^{\sharp} D \qquad = \qquad \square^{\sharp} [e]^{\sharp} D$$

Abstract operators working on signs (Addition)

+#	{0}	{+}	{-}	{-,0}	{-,+}	$\{0, +\}$	{-,0,+}
{0}	{0}	{+}		22			
{+}							
{-}							
$\{-,0\}$							
{-,+}							
$\{0, +\}$							
{-,0,+}	{-,0,+}						

Abstract operators working on signs (Multiplication)

×#	{0}	{+}	{-}	{-,0}	{-,+}	$\{0, +\}$	{-,0,+}
{0}	{0}	{0}					
{+}							
{-}							
$\{-,0\}$							
{-,+}							
$\{0, +\}$							
{-,0,+}	{0}						

Abstract operators working on signs (unary minus)

						{-,0,+}
 {0}	{-}	{+}	{+,0}	{-,+}	{0, -}	{-,0,+}

Working an example: ${}^{24}D = \{x \mapsto \{+\}, y \mapsto \{+\}\}$

$$[[x+7]]^{\sharp} D = [[x]]^{\sharp} D + [[7]]^{\sharp} D$$

$$= \{+\} + [+]^{\sharp} \{+\}$$

$$= \{+\}$$

$$[[x+(-y)]]^{\sharp} D = \{+\} + [(-\sharp[y]]^{\sharp} D)$$

$$= \{+\} + [(-\sharp\{+\})]$$

$$= \{+\} + [-]$$

$$= \{+, -, 0\}$$

Thus, we obtain the following effects of edges $\lceil lab \rceil^{\sharp}$:

... whenever

Attention to details

 $D \neq \bot$

49

The rule of signs by Reinhard Wilhelm (2012/13)

Idea:

- We want to determine the signs of the values of expressions.
- For some sub-expressions, the analysis may yield $\{+,-,0\}$, which means, it couldn't find out.
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Abstract operators working on signs (Multiplication)

_	$\times^{\#}$	{0}	{+}	{-}	{-,0}	{-,+}	$\{0, +\}$	{-,0,+}
_	{0}	{0}	{0}					
	{+}							
	{-}							
	{-,0} {-,+}							
	{-,+}							
	$\{0, +\}$ $\{-, 0, +\}$							
_	{-,0,+}	{0}						

Abstract operators working on signs (unary minus)

_#	{0}	{+}	{-}	{-,0}	{-,+}	$\{0, +\}$	{-,0,+}
	{0}	{-}	{+}	$\{+,0\}$	{-,+}	{0, -}	{-,0,+}

Working an example: ${}^{2}D = \{x \mapsto \{+\}, y \mapsto \{+\}\}$

$$[x+7]^{\sharp} D = [x]^{\sharp} D + {\sharp} [7]^{\sharp} D$$

$$= \{+\} + {\sharp} \{+\}$$

$$= \{+\}$$

$$[x+(-y)]^{\sharp} D = \{+\} + {\sharp} (-{\sharp} [y]^{\sharp} D)$$

$$= \{+\} + {\sharp} (-{\sharp} \{+\})$$

$$= \{+\} + {\sharp} \{-\}$$

$$= \{+-0\}$$

Thus, we obtain the following effects of edges $\lceil lab \rceil^{\sharp}$:

... whenever

 $D \neq \bot$

if the program
does not
terminate isn't it
correct to say that
x is 0 upon its
termination?

Attention to details

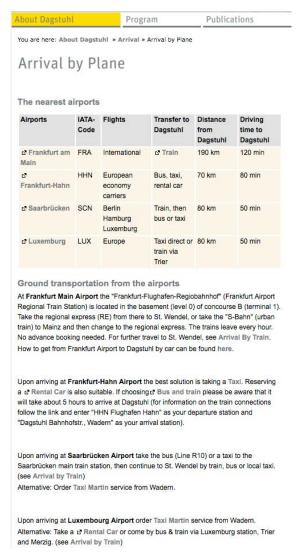
That's where you recognize a great scientist: make simple what is complicated!

Suggestions for an happy retirement

Have ambitious objectives!

Have ambitious objectives!

 Move Dagstuhl close to an airport (or an airport close to Dagstuhl)

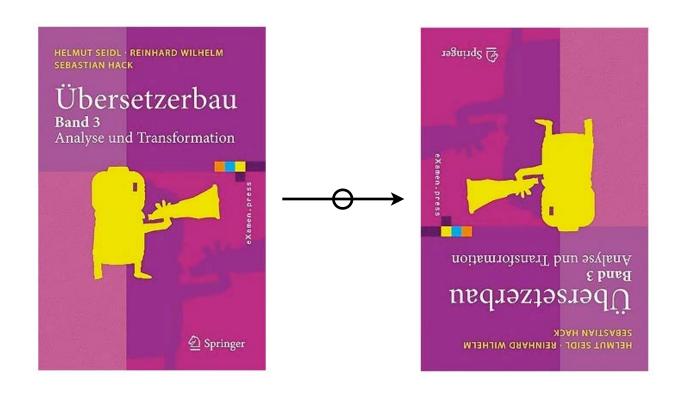


Start working on cyberimbedded systems

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- Start working on cyberimbedded systems
- Consider dynamic methods for static analysis
- Write a book on decompilation, by duality



Time for a serious conclusion

Thanks a lot for 30 years of friendship

Thanks a lot for 30 years of friendship, with lots of problems!

Thanks a lot for 30 years of friendship, with lots of problems!

shared scientific

The End,

thank you



