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Unit 3

Interface Between Assembly and C Languages

MICROPROCESSOR-BASED SYSTEMS

**Degree in Computer Science Engineering
Double Degree in Computer Engineering and Mathematics**

EPS - UAM

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3. Interface between assembly and C languages.

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3.1. General features

- Many applications written in high-level languages require critical parts written in assembly language (real time execution, use of multimedia instructions such as MMX, etc.)
- Need to be able to call assembly routines from programs written in compatible high-level languages.
- Also possible to call routines written in high-level compilable languages from assembly routines.
- Feasible if assembly programs follow the conventions (nomenclature, passing of parameters and results, ...) of the high-level languages.

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3.2. The example of C language (I)

- Most applications requiring interaction with assembly language are written in C (and C++).
- C language has typical high-level constructs (loops, structured types, recursion, ...), but also allows control at a very low level (access to I/O ports, bit manipulation, ...).
- C compilers allow linkage with assembly programs only if they follow the **same conventions** applied by the compiler.
- C program compiled into object file, assembly program assembled into object file and linker of the C compiler generates executable file by linking the object files.

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3.2. The example of C language (II)

- C language conventions related to:
 - Use of near or far addresses for accessing data (variables) and/or procedures: memory model.
 - Nomenclature of segments, variables and procedures.
 - Parameter passing to procedures and return of results.

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3.3. The different models of C language (I)

- A memory model must be chosen when a C program is compiled.
- Every model determines the memory location of the logic segments (code, data and stack) and if near or far addresses are used for accessing them.
- Six memory models in Turbo C:
 - TINY
 - SMALL
 - MEDIUM
 - COMPACT
 - LARGE
 - HUGE

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3.3. The different models of C language (II)

• TINY

- Minimum memory occupation.
- The four segment registers (**CS**, **SS**, **DS** y **ES**) are identical.
- The program occupies up to 64 KB.
- Code, data and stack on the same physical segment.
- Programs compiled in this model can be converted to .COM (*drivers*) through the EXE2BIN utility of DOS or using the linker option /t.
- Near pointers for code and data.

• SMALL

- Small programs in which minimum memory occupation is not necessary.
- A physical segment for code (up to 64 KB) and another segment for data and stack (up to 64 KB).
- Near pointers for code and data.

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3.3. The different models of C language (III)

• MEDIUM

- Big programs that use few data.
- Several physical segments for code (up to 1 MB) and a segment for data and stack (up to 64 KB).
- Far pointers for code and near pointers for data.

• COMPACT

- Small programs that use many data.
- A physical segment for code (up to 64 KB) and several segments for data and stack (up to 1 MB).
- Near pointers for code and far pointers for data.

• LARGE

- Big programs that use many data.
- Several physical segments for code (up to 1 MB) and for data and stack (up to 1 MB). Not possible to exceed 1 MB in total.
- Far pointers for code and data.

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3.3. The different models of C language (IV)

• HUGE

- Similar to LARGE with some advantages and disadvantages.
- Normalized pointers (offset < 16).
- Static global variables can exceed 64 KB (possible to manipulate data blocks larger than 64 KB).
- Compiler inserts code that automatically updates data segment registers (data pointers always normalized).
- Most costly model in terms of execution time.

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3.3. The different models of C language (V)

Model	Segments			Pointers	
	Code	Data	Stack	Code	Data
Tiny	64 KB			NEAR	NEAR
Small	64 KB	64 KB		NEAR	NEAR
Medium	1 MB	64 KB		FAR	NEAR
Compact	64 KB	1 MB		NEAR	FAR
Large	1 MB	1 MB		FAR	FAR
Huge	1 MB	1 MB (blocks larger than 64 KB)		FAR	FAR

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3.4. Conventions on nomenclature, parameter passing and return of results (I)

Nomenclature

- The C compiler always gives the same name to the logical segments it uses:
 - The code segment is called **_TEXT**.
 - Segment **_DATA** contains the initialized global variables.
 - Segment **_BSS** contains the uninitialized global variables.
 - The stack segment is defined and initialized by the C compiler in the *main* function.
 - In the small data models (*tiny*, *small* and *medium*), all data segments are grouped under the name **DGROUP**:

DGROUP GROUP _DATA, _BSS

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3.4. Conventions on nomenclature, parameter passing and return of results (II)

Nomenclature

- The C compiler adds **_** right before all names of variables and procedures:

- Example:

```
int a = 12345;  
char b = 'A';  
char c[] = 'Hello world';  
int d = 12;
```

- It is compiled as:

```
_DATA SEGMENT WORD PUBLIC 'DATA'  
PUBLIC _a, _b, _c, _d  
_a DW 12345  
_b DB 'A'  
_c DB 'Hello world', 0  
_d DW 12  
_DATA ENDS
```

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3.4. Conventions on nomenclature, parameter passing and return of results (III)

Nomenclature

- The C compiler adds **_** right before all names of variables and procedures:

- Example:

```
main()
{
    function();
}
```

- It is compiled as:

```
_TEXT SEGMENT BYTE PUBLIC 'CODE'

_main PROC FAR
    CALL _function
    RET
_main ENDP

_TEXT ENDS
```

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3.4. Conventions on nomenclature, parameter passing and return of results (IV)

Nomenclature

- The assembly variables and procedures accessed from C programs must be preceded by `_`, which does not appear in C.
- The C language distinguishes between lowercase and uppercase: `function()` and `FUNCTION()` are different procedures.
- It is necessary that the assembler also distinguishes between lowercase and uppercase.
- In TASM, this is done by assembling with options `/mx` (it forces distinction for public symbols) or `/ml` (it forces distinction for all symbols).

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3.4. Conventions on nomenclature, parameter passing and return of results (V)

Parameter passing

- In C language, a procedure that calls another stacks its parameters before executing **CALL**.
- The assembly procedures that call C functions must also stack their parameters.
- Parameters are stacked in reversed order with respect to their position in the C call: starting with the last and ending with the first.
- After returning from the subroutine, parameters are extracted from the stack by adding their size in bytes to register **SP**.
- Single-byte parameters (char) are stacked using two bytes (the most significant byte is 0).

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3.4. Conventions on nomenclature, parameter passing and return of results (VI)

Parameter passing

- Parameters are stacked in *little endian* format: least significant word in lowest address and least significant byte in lowest address.
- In order to pass pointers to functions or data, it is necessary to know the memory model in which the C program is being compiled, so as to stack the segment register (far model) or not (near model).

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3.4. Conventions on nomenclature, parameter passing and return of results (VII)

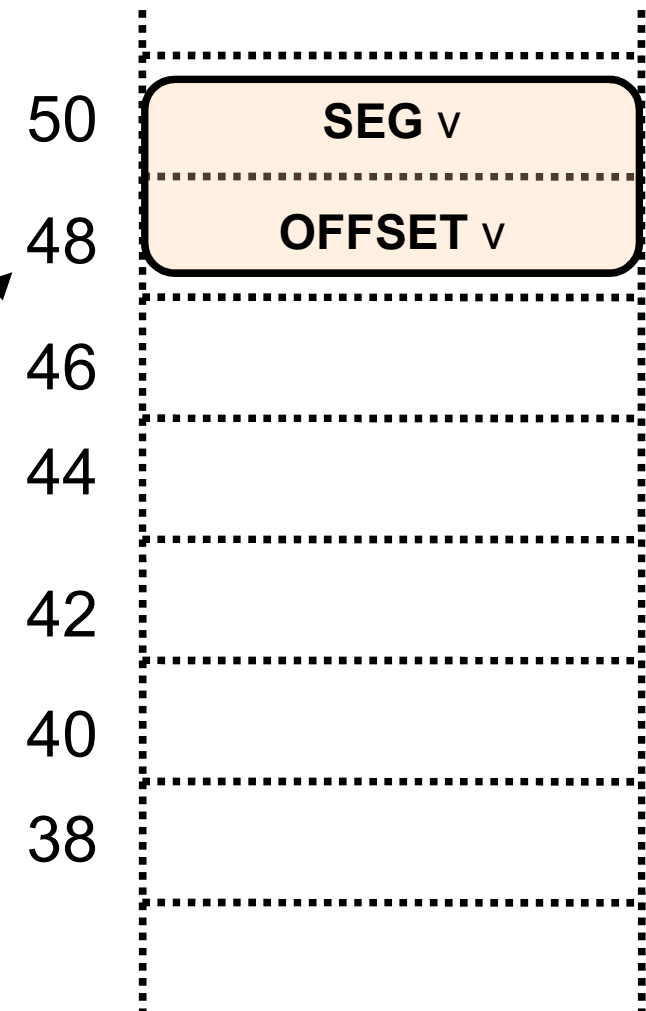
Parameter passing (example)

- Function prototype:

```
void function ( char, long int, void * );
```

- Call in far model (**FAR** pointers for data and code)

```
function ( 'P', 0x23, &v );
```



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3.4. Conventions on nomenclature, parameter passing and return of results (VIII)

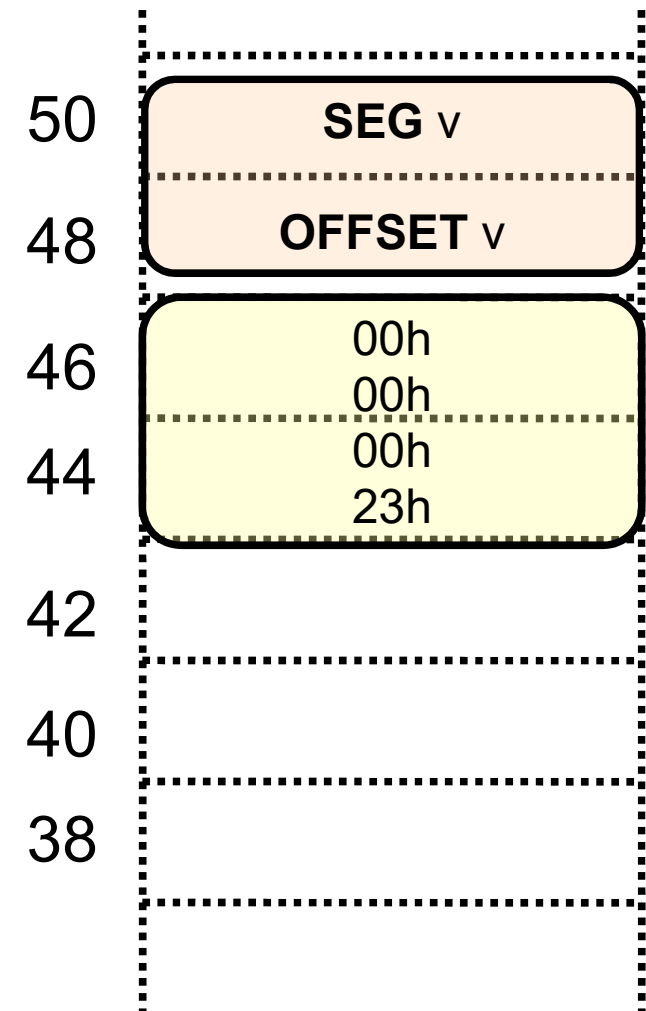
Parameter passing (example)

- Function prototype:

```
void function ( char, long int, void * );
```

- Call in far model (**FAR** pointers for data and code)

```
function ( 'P', 0x23, &v );
```



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3.4. Conventions on nomenclature, parameter passing and return of results (IX)

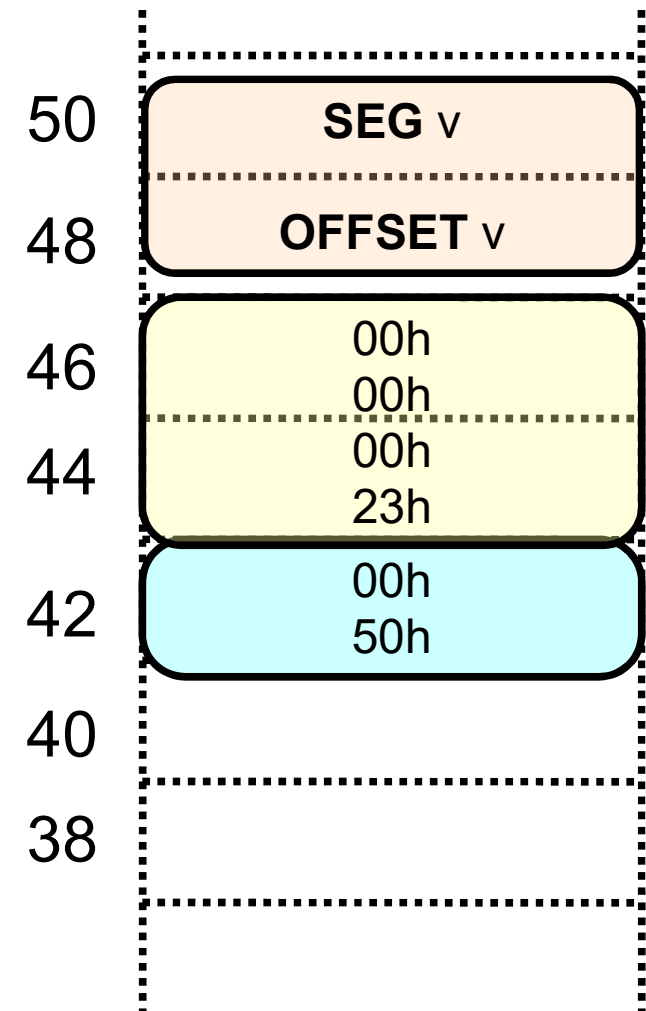
Parameter passing (example)

- Function prototype:

```
void function ( char, long int, void * );
```

- Call in far model (**FAR** pointers for data and code)

```
function ( 'P', 0x23, &v );
```



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3.4. Conventions on nomenclature, parameter passing and return of results (X)

Parameter passing (example)

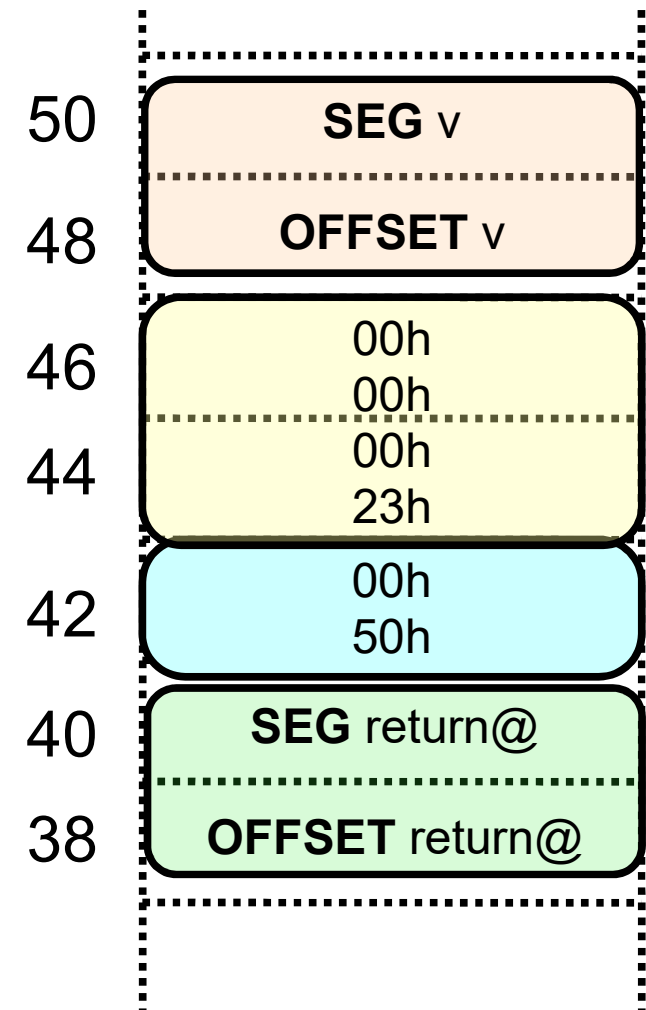
- Function prototype:

`void function (char, long int, void *);`

- Call in far model (**FAR** pointers for data and code)

`function ('P', 0x23, &v);`

`call _function`

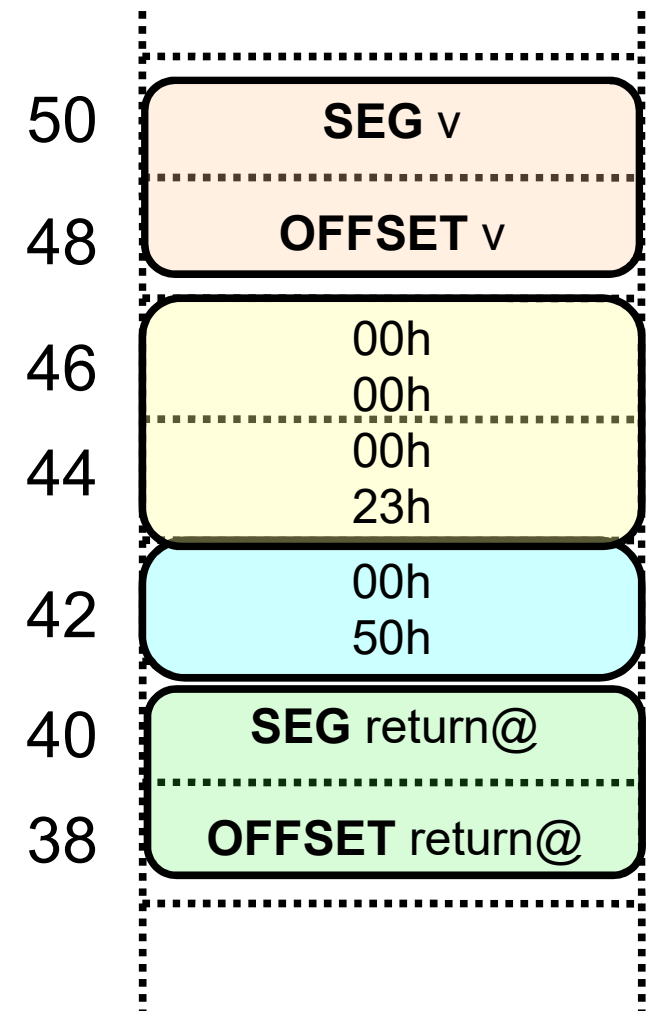
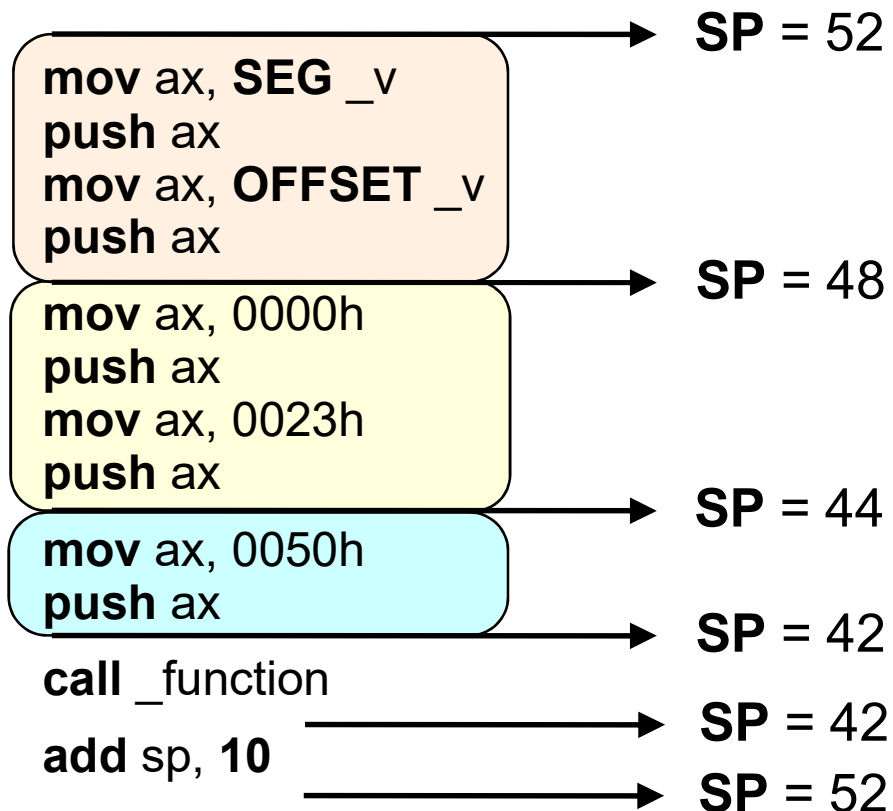


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3.4. Conventions on nomenclature, parameter passing and return of results (XI)

Parameter passing (example)

```
void function ( char, long int, void * );  
function ( 'P', 0x23, &v );
```



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3.4. Conventions on nomenclature, parameter passing and return of results (XII)

Parameter passing (example)

```
void function ( char, long int, void * );
```

```
_function PROC FAR
```

```
    push bp
```

```
    mov bp, sp
```

bx := OFFSET v

es := SEG v

```
    les bx, [ bp + 12 ]
```

```
    mov dx, [ bp + 10 ]
```

```
    mov ax, [ bp + 8 ]
```

```
    mov cx, [ bp + 6 ]
```

```
    ...
```

```
    ret
```

```
_function ENDP
```

SP = BP

50

SEG v

48

OFFSET v

46

00h

00h

44

00h

23h

42

00h

50h

40

SEG return@

38

OFFSET return@

36

initial BP

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3.4. Conventions on nomenclature, parameter passing and return of results (XIII)

Return of results

- The return variables of the function with length of 16 bits are returned to the calling procedure through **AX** and the ones of 32 bits through **DX:AX**.

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3.4. Conventions on nomenclature, parameter passing and return of results (XIV)

Example 1

/ Visible variable (external) from the assembly routine */*

int variable_c;

/ Variable defined as public in the assembly program */*

extern int data_as;

/ Declaration of the assembly function (it could be within an **include**) */*

int function (**int** a, **char far** *p, **char** b);

main()

{

int a = 123;

/ Declaration of C local variables */*

char b = 'F';

char far *p;

/ Call the function and store in variable_c the value returned through AX */*

variable_c = function (a, p, b);

}

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3.4. Conventions on nomenclature, parameter passing and return of results (XV)

```
DGROUP GROUP _DATA, _BSS
_DATA SEGMENT WORD PUBLIC 'DATA'
    PUBLIC _data_as
    _data_as DW 0
_DATA ENDS
```

; Data segments are grouped
; Public data segment
; Declaration of _data_as as public
; Allocation of _data_as and initialization

```
_BSS SEGMENT WORD PUBLIC 'BSS'
    EXTRN _variable_c : WORD

    _hidden_as DW ?
_BSS ENDS
```

; Public data segment
; Declaration of _variable_c as external and of
; type WORD (defined in the C module)
; Variable not accessible from the C module

```
_TEXT SEGMENT BYTE PUBLIC 'CODE'
    ASSUME CS:_TEXT, DS:DGROUP
PUBLIC _function
_function PROC NEAR
    PUSH BP
    MOV BP, SP
    MOV BX, [ BP+4 ]
    LDS SI, [ BP+6 ]
    MOV CX, [ BP+10 ]
    ...

    MOV AX, CX
    POP BP
    RET
_function ENDP
_TEXT ENDS
END
```

; Definition of the code segment

; Make function accessible from C
; It is function() in C
; For using BP in order to address the stack,
; it is loaded with pointer to the top
; Store a in BX , BX=123
; Store p in DS:SI
; Store b in CX, CL='F', CH =0

; Value returned through AX since function is int
; Restore BP
; Return to calling procedure

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3.4. Conventions on nomenclature, parameter passing and return of results (XVI)

Example 2

`char * strchr (char *string, int character);`

data **SEGMENT**

 mystring **DB** "This is an ASCIIZ string", 0

data **ENDS**

code **SEGMENT**

mov ax,'a'

/* Sought character is stacked */

push ax

mov ax, **SEG** mystring

/* String is stacked */

push ax

mov ax, **OFFSET** mystring

push ax

call **FAR** _strchr

/* Procedure call */

add sp, 6

/* Stack balancing */

mov ds, dx

/* Pointer returned through **DX:AX** */

mov si, ax

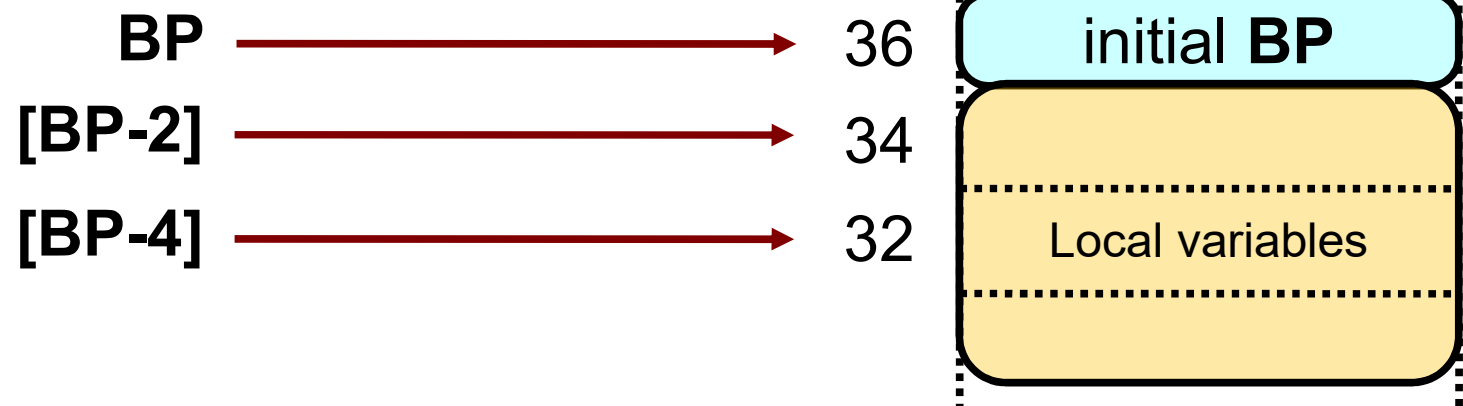
code **ENDS**

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3.4. Conventions on nomenclature, parameter passing and return of results (XVII)

Definition of local variables

- C routines define the local variables they require on the stack and above BP.
- Local variables are introduced in the same order as they are declared.
- They are accessed through [BP-2], [BP-4],



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3.4. Conventions on nomenclature, parameter passing and return of results (XVIII)

- Instruction `asm` allows inserting assembly code within a C program (*inline assembly*).

```
main
{
    int d1 = 5, d2 = 4, result;
    asm {
        push cx
        push ax
        mov ax, 0
        mov cx, d2
        cmp cx, 0
        jz final
    }
    mult:
    asm {
        add ax, d1
        dec cx
        jnz mult
    }
    final:
    mov result, ax
    pop ax
    pop cx
}
printf("result %d\n", result); }
```

BP must not be used, since it is used by the compiler to access local variables.