

lema: sea $l^{(k)} = (0 \dots 0 l_{k+1}^{(k)} \dots l_n^{(k)})^t$, por $k=1 \dots n-1$

definimos $L_k = I + l^{(k)} \otimes e_k$ $[(e_k)_j = \delta_{k,j}]$

\Rightarrow i) $L_k^{-1} = I - l^{(k)} \otimes e_k$

ii) $L_k L_{k+1} = I + l^{(k)} \otimes e_k + l^{(k+1)} \otimes e_k$

demo:

i) $(I + l^{(k)} \otimes e_k)(I - l^{(k)} \otimes e_k)$

$= I + \cancel{l^{(k)} \otimes e_k} - \cancel{l^{(k)} \otimes e_k} - \underbrace{(l^{(k)} \otimes e_k)^2}_0$

ejercicio: demostrar que este término es 0

ii) $(I + l^{(k)} \otimes e_k)(I + l^{(k+1)} \otimes e_{k+1})$

$= I + l^{(k)} \otimes e_k + l^{(k+1)} \otimes e_{k+1} + T$

$T = \underbrace{(l^{(k)} \otimes e_k)}_M \underbrace{(l^{(k+1)} \otimes e_{k+1})}_N$, $t_{ij} = \sum_f m_{if} n_{fj}$

$t_{ij} = \sum_{\substack{f \\ f \neq k}} l_i^{(k)} \underbrace{(e_k)_f}_{\delta_{k,f}} l_f^{(k+1)} (e_{k+1})_j$

$= l_i^{(k)} \underbrace{l_k^{(k+1)}}_0 \delta_{j, k+1} = 0 \quad \forall i, j = 1 \dots n$

alternativa: $T = \underbrace{\begin{pmatrix} l^{(k)} \\ 1 \end{pmatrix} \begin{pmatrix} -e_k - \end{pmatrix} \begin{pmatrix} l^{(k+1)} \\ 1 \end{pmatrix} \begin{pmatrix} -e_{k+1} - \end{pmatrix}}_0$

$\langle e_k, l^{(k+1)} \rangle = l_n^{(k+1)} = 0$

demostración (teorema LU)

$$\bullet \quad U = U^{(m-1)} = L_{m-1}^{-1} U^{(m-2)} = \dots = L_{m-1}^{-1} L_{m-2}^{-1} \dots L_1^{-1} \underbrace{U^{(0)}}_A$$

$$\Rightarrow A = \underbrace{L_1 L_2 \dots L_{m-1}} U$$

$$\text{ii) lema : } = I + \sum_{k=1}^{m-1} l^{(k)} \otimes e_k$$

tenemos que demostrar que U triangular alta

$$U^{(k)} = L_k^{-1} U^{(k-1)}$$

$$u_{ij}^{(k)} \stackrel{\text{i) lema}}{=} \sum_j (I - l^{(k)} \otimes e_k)_{ij} u_{jj}^{(k-1)}$$

$$= u_{ij}^{(k-1)} - \sum_j l_i^{(k)} \delta_{jk} u_{jj}^{(k-1)}$$

$$= u_{ij}^{(k-1)} - l_i^{(k)} u_{k,j}^{(k-1)}$$

$$\Rightarrow u_{ij}^{(k)} = \begin{cases} u_{ij}^{(k-1)}, & i = 1 \dots k \\ u_{ij}^{(k-1)} - \frac{u_{ik}^{(k-1)}}{u_{kk}^{(k-1)}} u_{kj}^{(k-1)}, & i = k+1 \dots m \end{cases}$$

las primeras k filas no cambian

ver luego lo que se hace en el algoritmo

$$\boxed{S_i \quad j = k} \quad u_{ik}^{(k)} = \begin{cases} u_{ik}^{(k-1)}, & i = 1 \dots k \\ 0, & i = k+1 \dots m \end{cases}$$

en la columna k de la k -ésima iteración, los elementos de las filas $k+1 \dots m$ son 0

$$\left(\begin{array}{ccc|ccc} \circ & \circ & \circ & \circ & \circ & \circ \\ \circ & \circ & \circ & \circ & \circ & \circ \\ \circ & \circ & \circ & \circ & \circ & \circ \\ \circ & \circ & \circ & \circ & \circ & \circ \\ \circ & \circ & \circ & \circ & \circ & \circ \end{array} \right) \Rightarrow \text{la matriz } U^{(m-1)} \text{ es triangular alta}$$

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Algoritmo: input A matriz $n \times n$
outputs L, U "

$U = A; L = I;$

for $k = 1 : n-1$

for $i = k+1 : n$

$$L(i, k) = \frac{U(i, k)}{U(k, k)};$$

⚠ no hay control sobre dividir por 0

for $j = k : n$

pregunta: qué pasa si ponemos $j = k+1 : n$?

$$U(i, j) = U(i, j) - L(i, k) U(k, j);$$

end

end

end