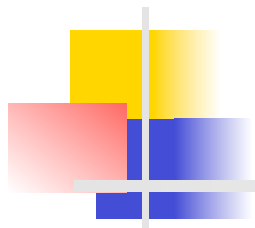


# COMP3411-9814- Artificial Intelligence



## Prolog Terminology & Syntax

2019 – Summer Term

Tatjana Zrimec



# Outline

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- ◆ Terminology
- ◆ Syntax and semantics
- ◆ Data objects
- ◆ Structures
- ◆ Matching
- ◆ Declarative Meaning



# Terminology

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## Terminology

- Prolog program - a set of *clauses*
- *Clauses* - facts, rules, questions
- *Fact* – things that are always, unconditionally true.
- Rules declare things that are true given condition

```
rule(X,Y) :- part1(X), part2(X,Y).
```

Diagram illustrating the structure of a Prolog rule:

- The **head** of the rule is `rule(X,Y)`.
- The **body** of the rule is `part1(X), part2(X,Y).`

- Variables - X, Y, B1, X12...
- Constants – numbers or atoms (a1, tom ....)



# Terminology - examples

## Constants

### Numbers:

1    -2    3.14

### Atoms:

tigger

'100 Acre Wood'

## Variables

X   A\_variable   \_

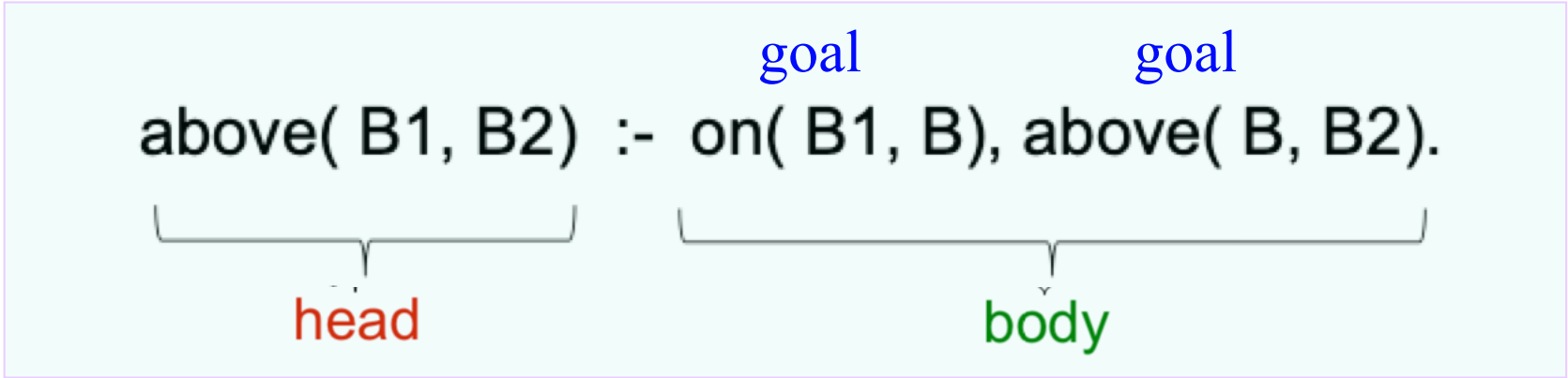
## Compound terms

likes(pooh\_bear,honey)  
plus(4,mult(3,plus(1,9)))



## Terminology – examples - Rules

- ◆ Example of a rule:



above( B1, B2) :- on( B1, B ), above( B, B2 ).

head

body

goal

goal

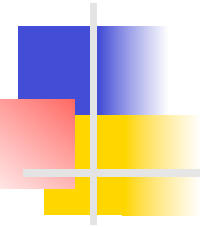
- ◆ The *head* is true if the *first goal* and the *second goal* are true.



## The term *atom*

---

- ◆ The term **atom** is used to denote a fundamental data type that cannot be made up from other data types.
- ◆ For example:
  - numbers and words are atoms,
  - lists are not atoms.



# Prolog – syntax and semantics

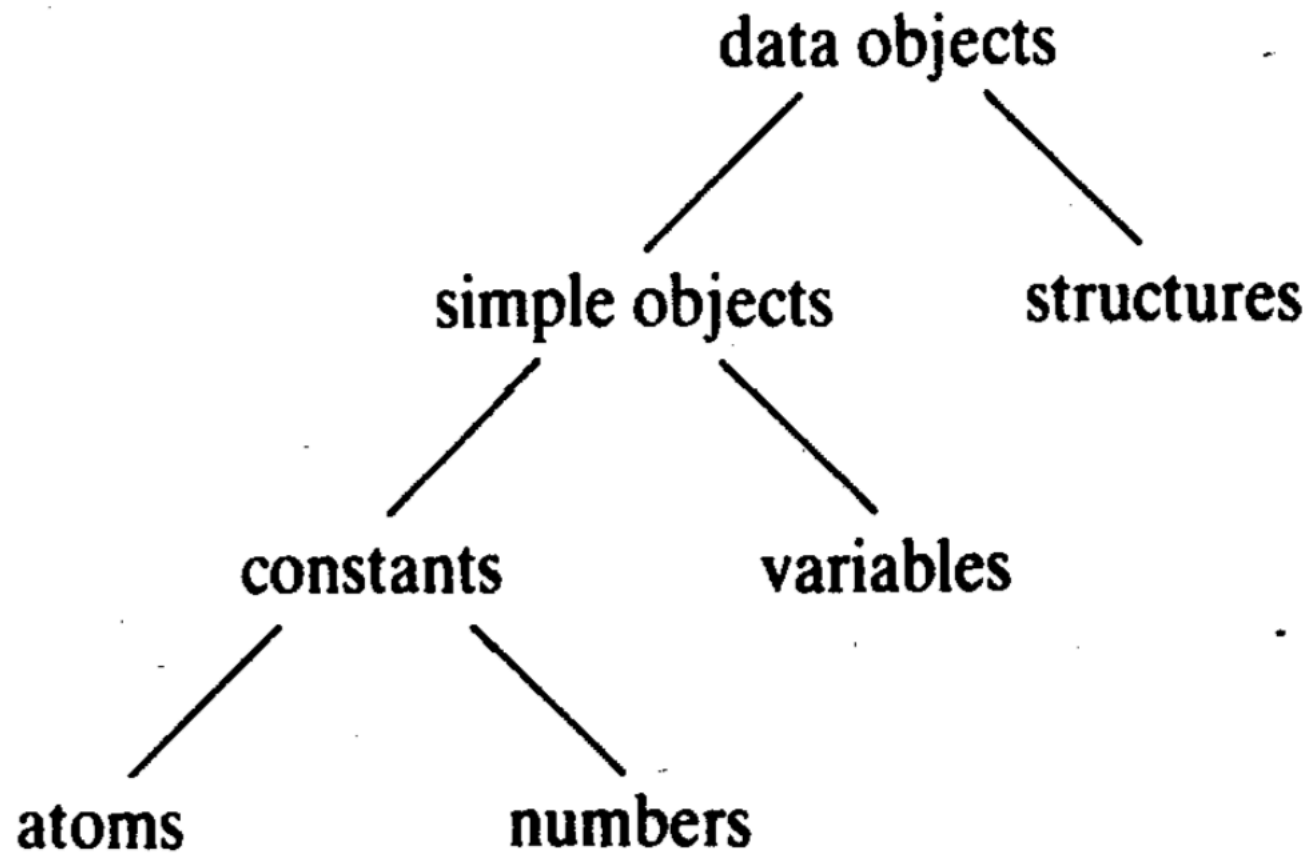
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## Data objects in Prolog

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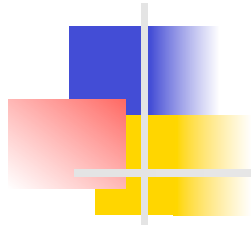




# Object Syntax

---

- ◆ The type of object is always recognizable from a syntactic form



## Three Syntactic Forms for Atoms

---

- (1) Strings of letters, digits and the underscore character “-”, starting with lower case letter

x      x15      x\_15      aBC\_CBa7

alpha\_beta\_algorithm      taxi\_35

peter      missJones      miss\_Jones2



## Three Syntactic Forms for Atoms

---

### (2) Strings of special characters

--->

<==>

<<

.

<

>

+

++

!

..

...

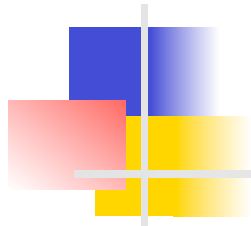
::=

[]

### (3) Strings of characters enclosed in single quotes

'X\_35'   'Peter'   'Britney Spears'

This is useful if we want an atom to start with a capital letter



# Numbers

---

- Strings of special characters

1      1313      0      -55

- Real numbers

3.14      -0.0045      1.34E-21      1.34e-21

Real numbers not much used in Prolog



# Variables

---

- ◆ Variable are strings of letters, digits and underscore character :

X    Results    Object2B    Participant\_list

\_x35    \_335

- ◆ The lexical range of variable names is one clause.



# Anonymous Variables

---

```
visible_block( B) :-  
    see( B, _, _).
```

It is equivalent to :

```
visible_block( B) :-  
    see( B, X, Y).
```



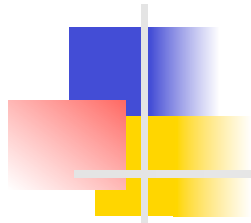
## Anonymous Variables

---

```
visible_block( B) :-  
    see( B, _, _).
```

- ◆ Each occurrence of the underscore character's appearing alone means: *I don't care what '\_' matches so long as it matches something.*
- ◆ Multiple occurrences of the character can be matched to different values.
- ◆ The '\_' character is used when the value of a variable is not needed in the evaluation of a clause.





# Structures

- ◆ Struktury so objekti z več komponentami
  - Npr.: datum je struktura s tremi komponentami
  - Datum 5. marec 2017:

`date( 5, march, 2017)`

*functor*      *arguments*

- ◆ The argument can be any object, including the structure

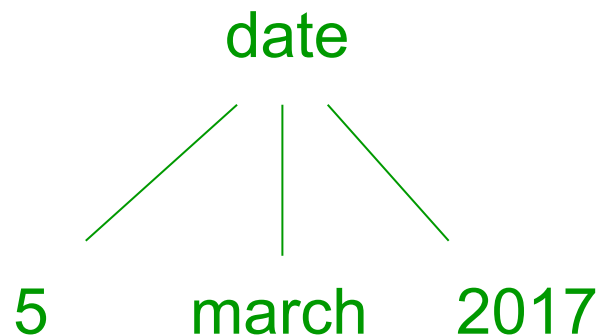


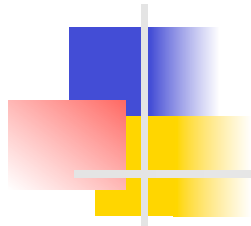
## Tree representation of structures

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- ◆ Structures are sometimes illustrated as trees:

date( 5, march, 2017)

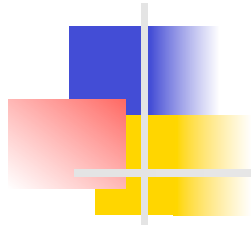




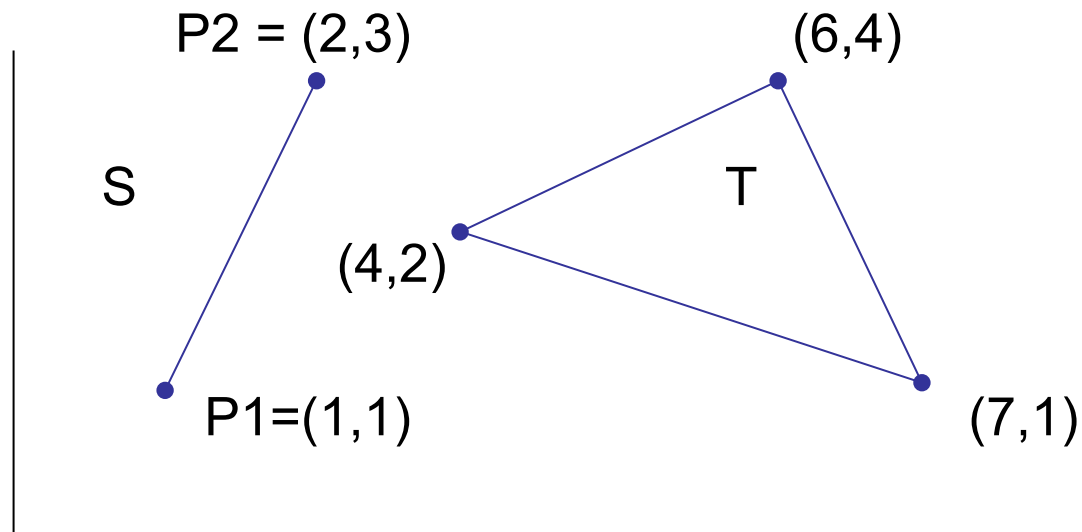
## Structure

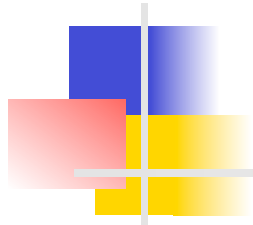
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- ◆ All structured objects in the prolog can be illustrated by trees
  - This is the only way of constructing structures in a Prolog
- ◆ Syntactically all abject in Prolog are “*terms*”

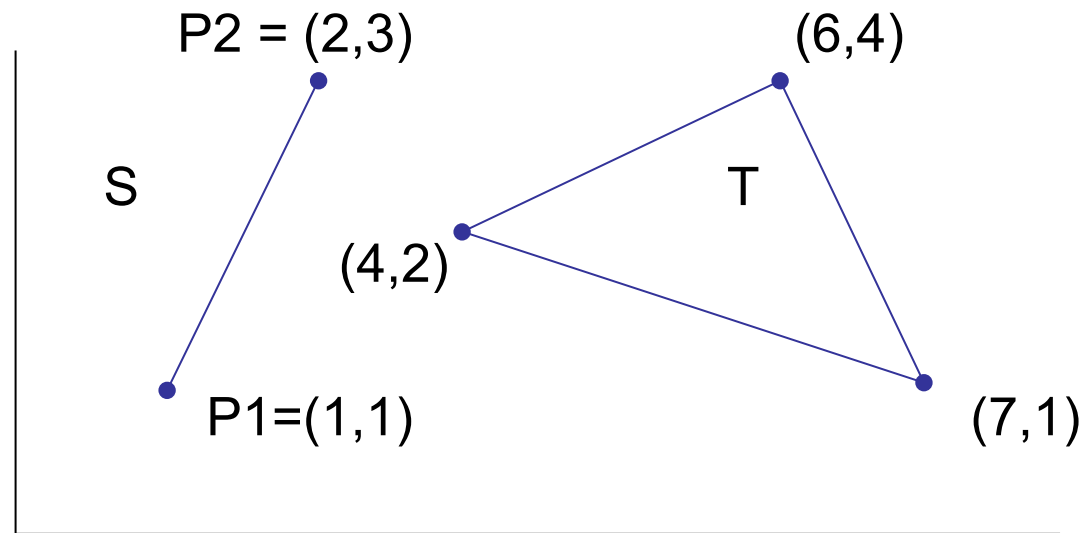


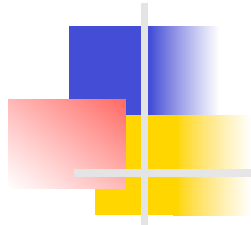
## Some simple geometric objects



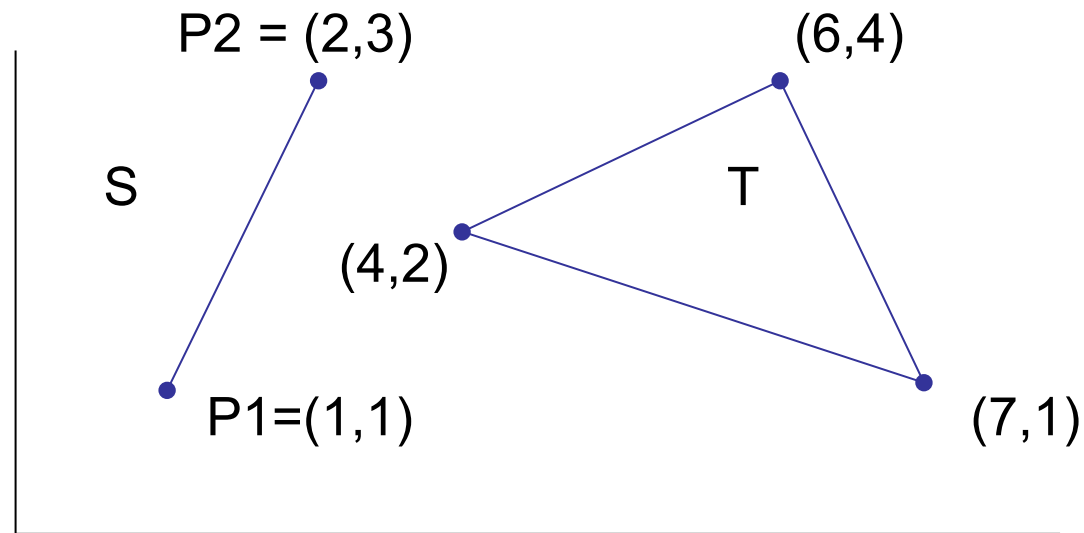


## Some simple geometric objects



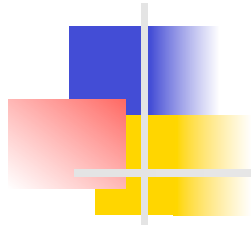


## Some simple geometric objects

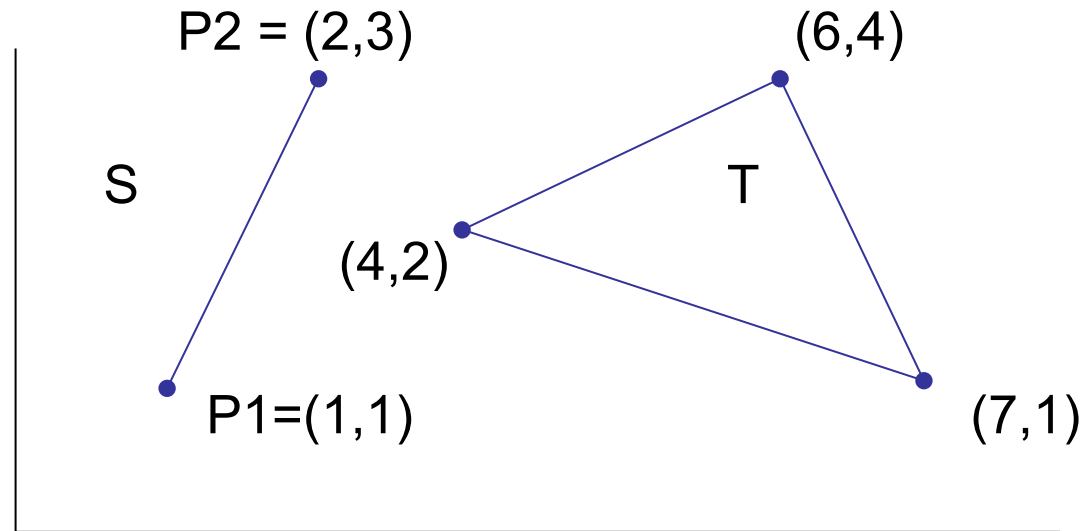


$P1 = \text{point}(1, 1)$

$P2 = \text{point}(2, 3)$



## Some simple geometric objects



$P1 = \text{point}(1, 1)$

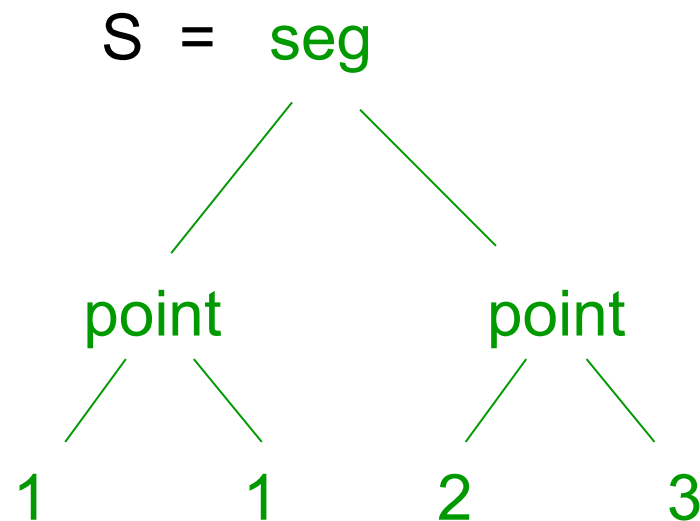
$P2 = \text{point}(2, 3)$

$S = \text{seg}(P1, P2) = \text{seg}(\text{point}(1,1), \text{point}(2,3))$

$T = \text{triangle}(\text{point}(4,2), \text{point}(6,4), \text{point}(7,1))$



$S = \text{seg}(\text{point}(1,1), \text{point}(2,3))$

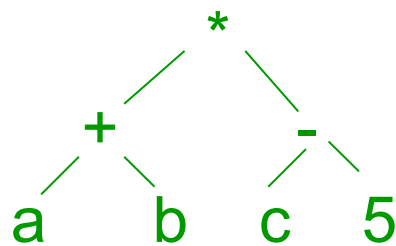


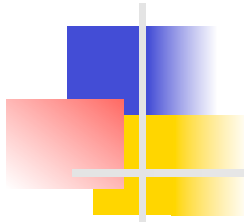


# The Arithmetic Expressions are also Trees

- ◆ For example:  $(a + b) * (c - 5)$
- ◆ Written as an expression with the functors:

$*(+(a, b), -(c, 5))$





# Matching

---

- ◆ Matching is an operation on *terms*.
- ◆ Two structures can be customized if:
  - (1) They are identical, or
  - (2) We can make them identical with appropriate definition of variables

## definition of variable

- the variable gets the value =
- instantiation of variable



# Matching

---

- ◆ Matching is an operation on *terms*.

Given two *terms*, they match if:

- (1) They are identical, or
- (2) The variable in both terms can be instantiated to objects in such a way that after the substitution of variables by these objects in the terms become identical

- substitution of variable
  - the variable gets a value = instantiation of variable



## Examples of *Matching*

---

- ◆ *Matching of dates:*

$\text{date}(D1, M1, 2006) = \text{date}(D2, \text{june}, Y2)$

- ◆ One instantiation that make both terms identical

$D1 = D2$

$M1 = \text{june}$

$Y2 = 2006$

- ◆ This is the most general instantiation, there are others that are less general...

- ◆ For matching using the operator “=”



## *Matching*- most general instantiation

---

- ◆ Prolog always returns the most general instantiation.
- ◆ With this instantiation leaves grater freedom for further instantiation if further *Matching* is required

?- date( D1, M1, 2006) = date( D2, june, Y2),  
date( D1, M1, 2006) = date( 17, M3, Y3).

D1 = 17, D2 = 17, M1 = june, M3 = june,  
Y2 = 2006, Y3 = 2006



## Matching

- ◆ *Matching* succeeds or fails.

Two terms  $S$  and  $T$  match:

1. If  $S$  and  $T$  are constants, then they match only if they are identical
2. If  $S$  is a variable, the *Matching* succeeds,  $S$  becomes equal to  $T$ .
3. If  $S$  and  $T$  are structure, then they are adjusted only if:
  - a) they have the same principal functor and
  - b) all their corresponding components match.

The resulting instantiation is determined by the matching of the components



## Example

---

- ◆ Prolog – finding answers



## Prolog – finding answers

---

- ◆ Prolog uses depth first search to find answers !

```
a(1) .  
a(2) .  
a(3) .  
b(1) .  
b(2) .  
b(3) .  
c(A,B) :- a(A), b(B) .
```



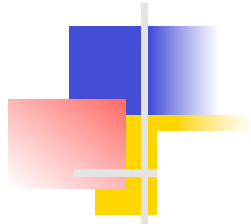


## Prolog – finding answers

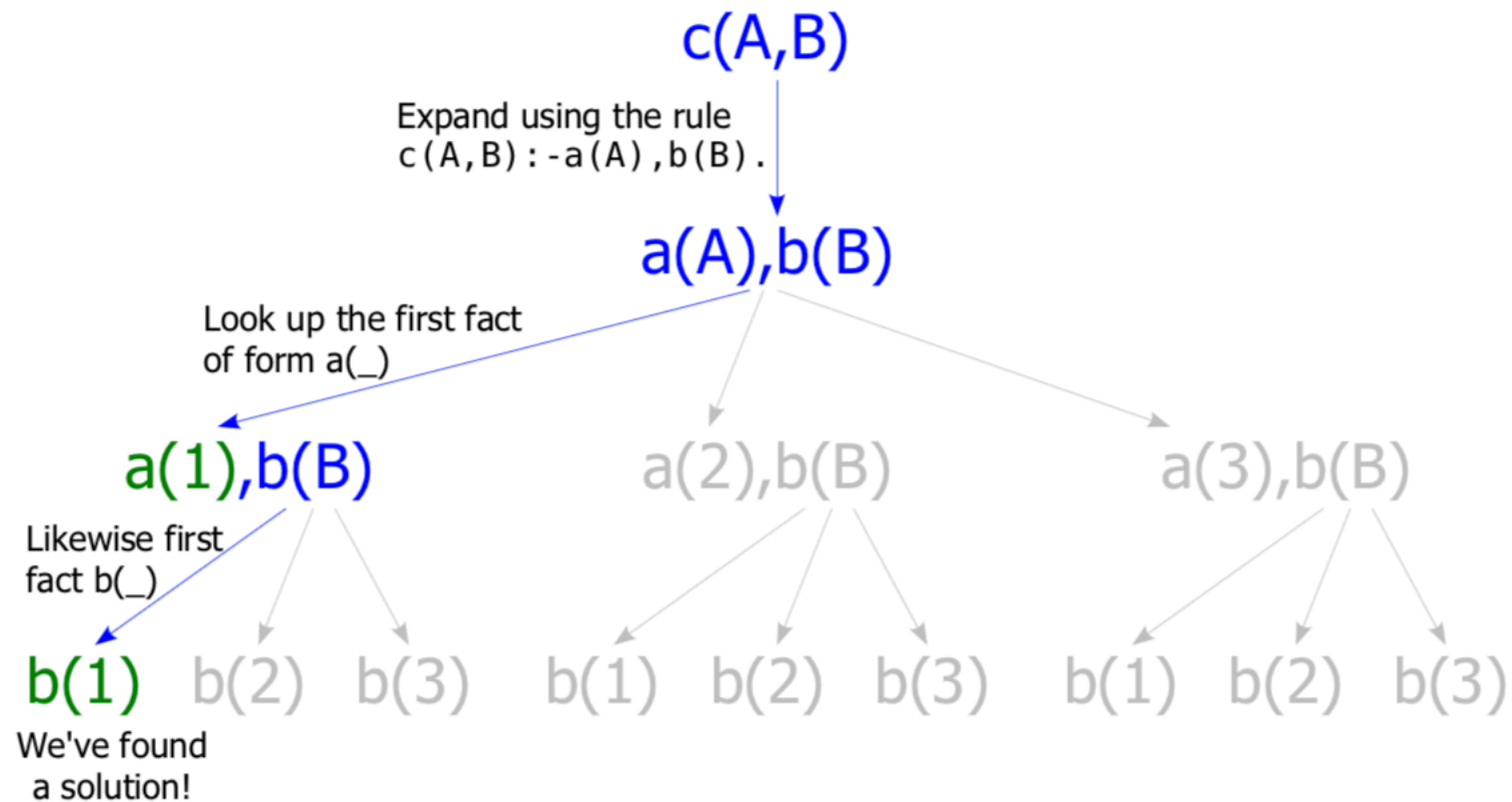
- ◆ Prolog uses depth first search to find answers !

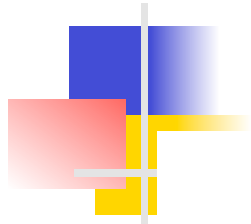
```
a(1) .  
a(2) .  
a(3) .  
b(1) .  
b(2) .  
b(3) .  
c(A,B) :- a(A), b(B) .
```

What does Prolog do when given this query ? `c(A,B)`

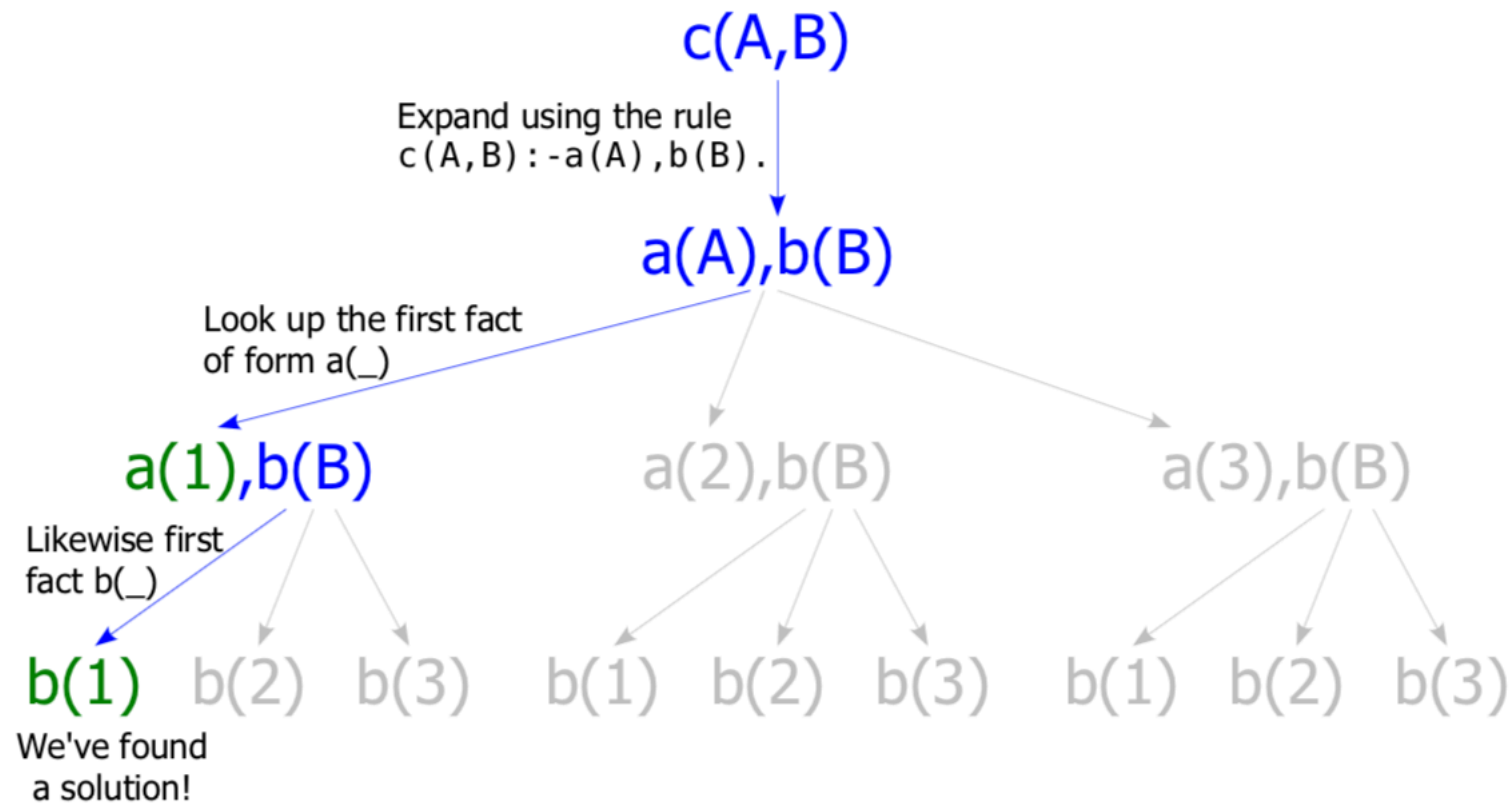


## Depth-first solution of query $c(A,B)$

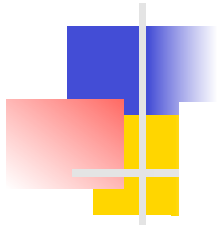




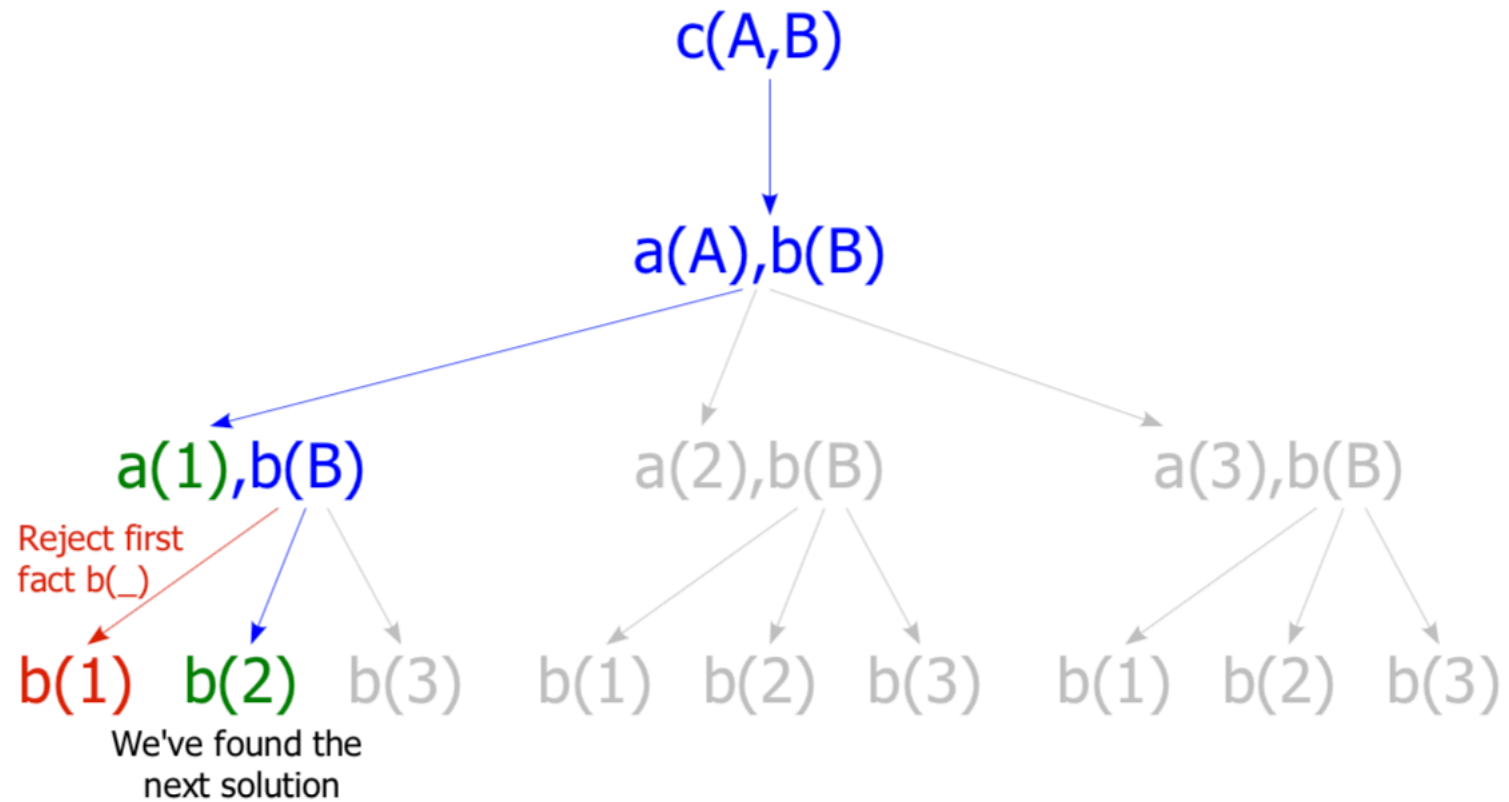
## Depth-first solution of query $c(A,B)$

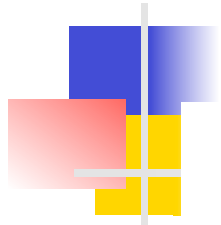


Variable bindings :  $A=1, B=1$

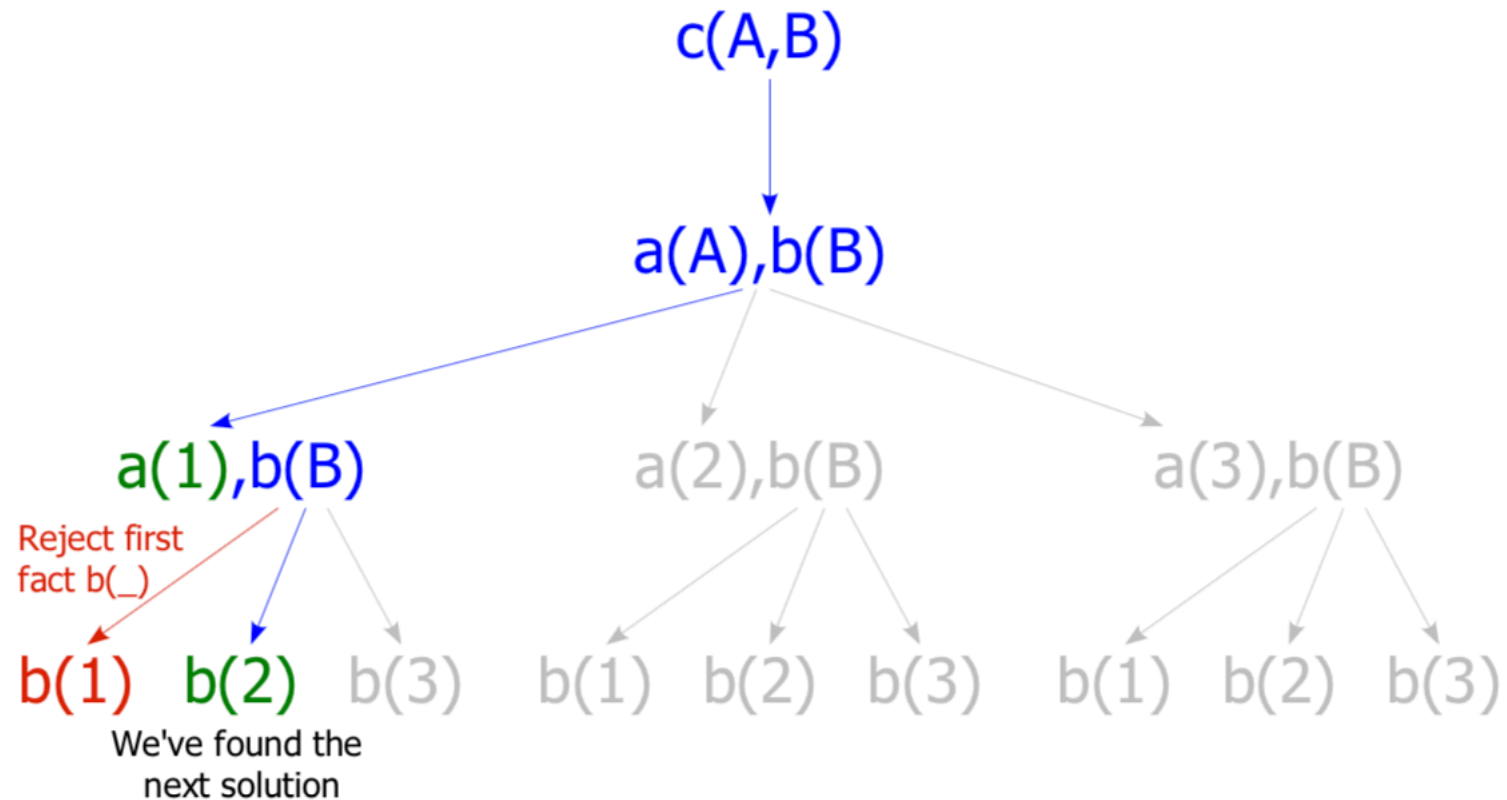


## Backtrack to find another solution



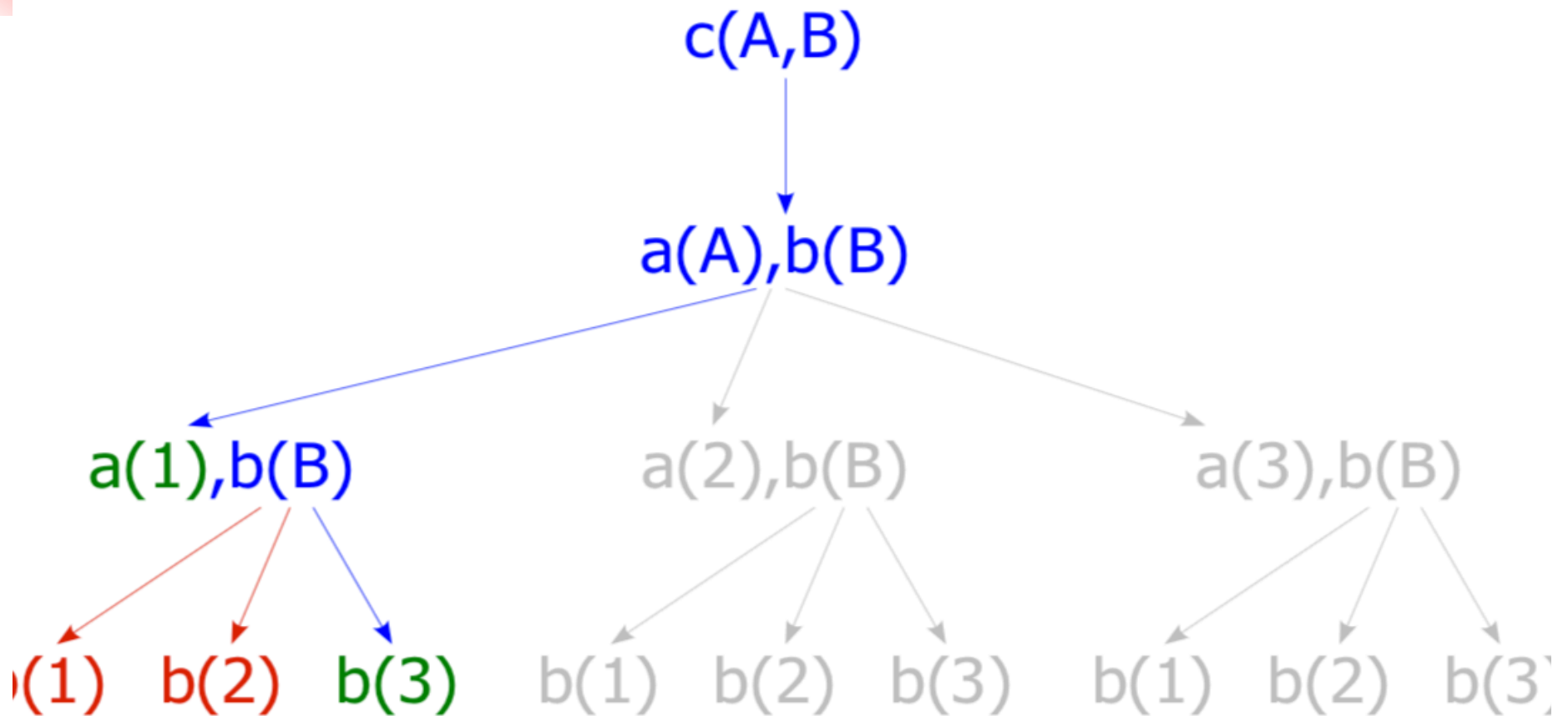


## Backtrack to find another solution

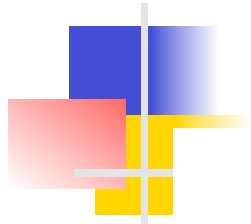


Variable bindings :  $A = 1, B = 2$

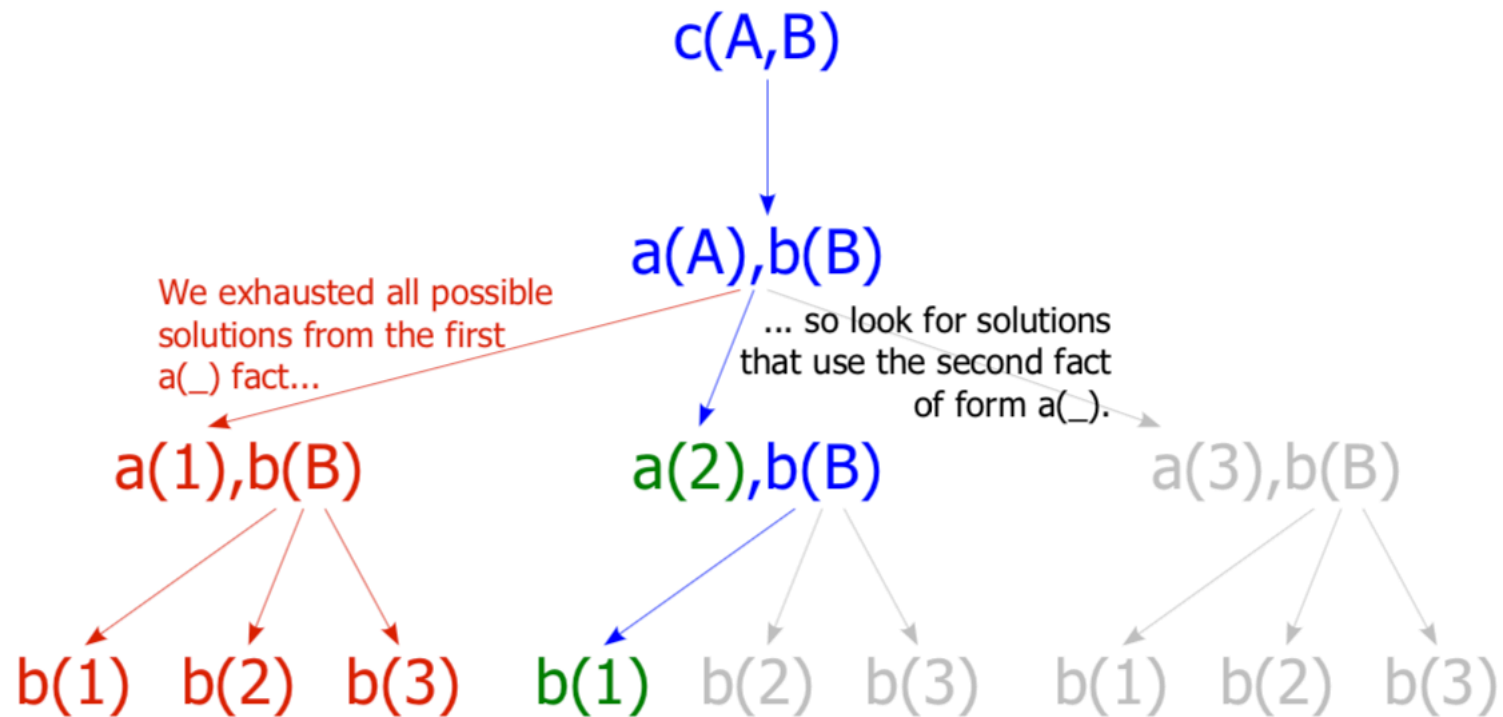
## Backtrack to find another solution



Variable bindings :  $A= 1, B=3$



## Backtrack to find another solution



Variable bindings :  $A=2, B=1$



## Declarative Meaning

- ◆ Let  $P$  be program and target  $G$
- ◆ A goal  $G$  is true (this is a logical follow from  $P$ ), if and only if:
  - (1) There is a clause  $C$  in  $P$  that is valid
  - (2) There is clause instance  $I$  of  $C$  such that
    - ✓ (a) the head of  $I$  is identical to  $G$ , and
    - ✓ (b) all the goals in the body of  $I$  are true
- ◆ In general, a question to Prolog is a *list* of goals separated by comas. A list of goals is true if all the goals in the list are true for some instantiation of variables.
- ◆ The values of the variables result from the most general instantiation.





# Declarative and Procedural Meaning of programs

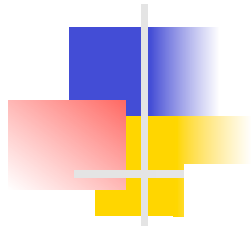
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- ◆ Let look at the clause:  
 $P :- Q, R.$
- ◆ Declarative reading of the clause:
  - P is true if Q and R are true.
  - From Q and R follows P.
- ◆ Procedural reading:
  - To solve the problem P, solve Q and then R.
  - To prove P, first prove Q and then R.

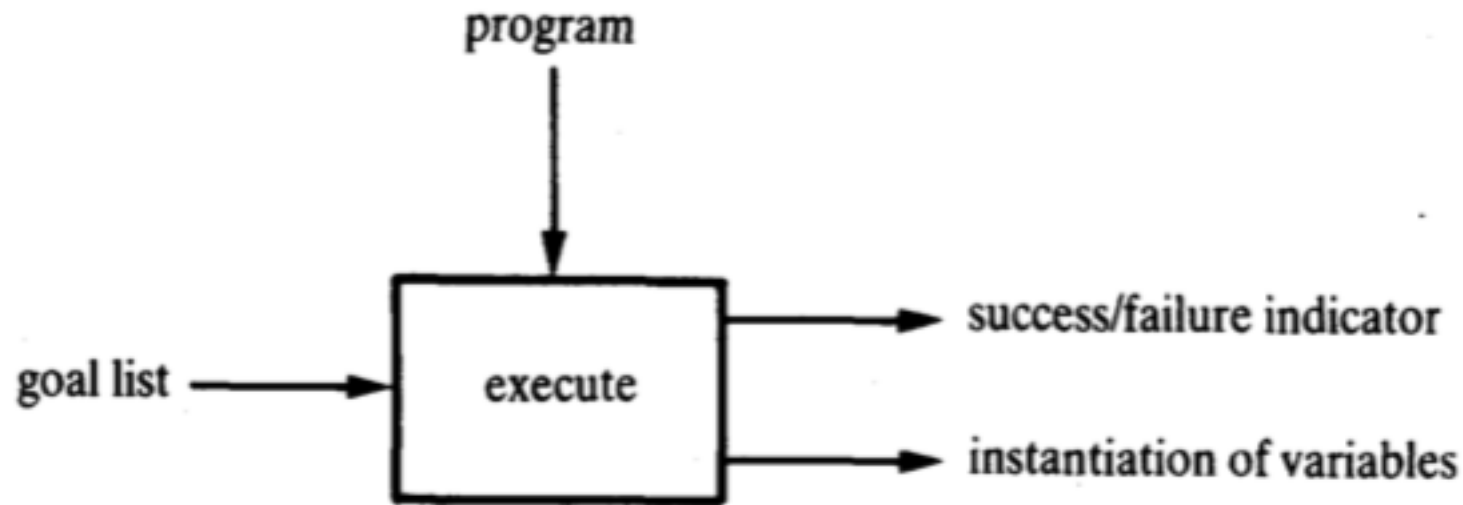


# Declarative and Procedural Meaning of programs

- ◆ A & B is logically equivalent to B & A
- ◆ Declarative meaning - only the relations defined by the program - *What* will be the output of a program
- ◆ The order of the goals in the clauses ***does not influence*** the declarative meaning
- ◆ The procedural meaning – how the relations are actually derived by the Prolog system
  - The algorithm
- ◆ The order of the goals in the clauses **influence** the procedural meaning



## Procedural meaning



Input/output view of the procedure that executes a list of goals.

It shows how Prolog prove goals

A procedural meaning is an algorithm for executing a list of goals with respect to a given program