PROJECT-2 NETWORK TOPOLOGY DESIGN

ALGORITHMIC ASPECTS OF TELECOMMUNICATION NETWORKS

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OBJECTIVE:

Given n (>=15) nodes on a plane and their coordinates, the key objective of this project is to create and implement two heuristic algorithms (different) of a network topology design, and experiment with it.

The network topology generated has the following properties:

- 1. The graph is connected, i.e. it contains all the given nodes.
- 2. The degree of every vertex in the graph is at least 3, i.e. each node is connected to 3 other nodes
- 3. The diameter of the graph is at most 4, i.e. any node can be reached from any other node in at most 4 hops. It is to be noted that this does not depend on the geometric distance.
- 4. The network topology has the lowest potential overall cost. The entire geometric length of all the links is used to compute this overall cost.

Only the fourth property is achieved by the heuristic algorithm. Every other characteristic is a constraint on the graph, requiring that it be present in the resultant network topology.

"Heuristic algorithms which prioritize speed over optimality and precision to achieve a solution that is quicker and efficient than the trivial techniques, are commonly used to tackle the decision issue class known as NP-complete problems. Heuristic approaches are typically employed when approximate answers are adequate but exact solutions are computationally expensive."[2] Numerous general purpose heuristic optimization algorithms are available, including Tabu search, Simulated Annealing, Greedy local search, and others.

HEURISTIC ALGORITHM & PSEUDOCODE – I

Greedy Local Search Algorithm:

Iteratively eliminating the edges with the highest cost while periodically determining if the resulting graph still meets the three (aforementioned) constraints is the basic notion behind this method. Calculate the whole cost once again if it does. This will continue until the conditions are no longer met by the graph.[1]

The program keeps eliminating the heavy weighted edges and ultimately gives the final optimum (least) total cost as the output.

- a. The original graph will serve as the basis for the current solution, S, and the total cost of the solution, C(S).
- b. Pick an edge that has the most weight and attempt to remove it.
- c. Verify that the new solution S' continues to fulfill all the constraints.
- d. If "yes," change C(S) to C(S') as the current total cost.
- e. Loop until the conditions are no longer met by the answer.

Pseudocode:

```
Cal_cost = sum(costs)
Sort(costs)
pointr = 0
while(len(costs)){
    Eliminate costs [pointr]
    if constraints satisfy
    Calculate Opt_cost
    if opt_cost<cal_cost
        then cal_cost = opt_cost
        pointr++
}
Else
    then add back removed edges
    Iterate from the next max weighed edge</pre>
```

HEURISTIC ALGORITHM & PSEUDOCODE – II

Heuristic Search Algorithm:

A sub-graph with a minimal cost path to reach a node from another node by traveling via every node in the inputted graph will be built in this case. If this sub-graph satisfies the constraints, the network's lowest overall cost will be determined based on its total cost. Otherwise, the topology will be updated with the subsequent lowest cost edge, and the conditions will be checked once more, and so forth until a sub-graph is created that meets the requirements.[2]

- a. Creating a sub-graph from the main graph by finding the minimum path between a source and destination that passes through every node in the network.
- b. If all the 3 conditions against the sub-graph are fulfilled, then the overall cost of the sub-graph will be determined.
- c. Otherwise, add the subsequent minimum-weighted edges to the sub-graph repeatedly until all the requirements are met.
- d. Substitute the newly calculated cost for the optimized cost now.

Pseudocode:

```
Visited = [0]
Unvisited = [1, 2, ..., 18];
Sub = [];
while ( Unvisited != [] ){
    Find edge e = (x, y) such that:
        1. x in Visited
        2. y in Unvisited
        3. ec edge with min cost
    Sub = Sub.append(ec)
    Visited = Visited.append(y)
    Unvisited = Unvisited.remove(y)
}
Iterate{
    if (all constraints):
        calculate total cost
```

```
Else:
    sub_graph.append(next min. wgt)
}
```

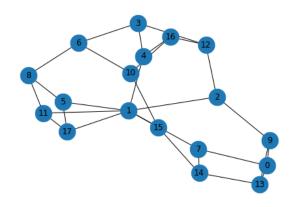
APPROACH:

- 1. To achieve, we construct a network of 18 nodes, and two nodes are chosen at random to create the edges connecting them.
- 2. We then determine if the resulting graph complies with the constraints and regenerate the graph otherwise.
- 3. Verify the nodes' connectivity first using the BFS algorithm.
- 4. Then verify if each node has a degree greater than 3 by counting the # of ones in each row of the adjacency matrix.
- 5. Then measure the diameter of each node, which should not exceed 4 using the built-in Python "dijkstra" function.
- 6. For randomly chosen coordinates of each node, the weights of each edge will be determined using the built-in Python function for calculating Euclidean distance. The weights of all the edges are added to determine the network's total cost.
- 7. We then implement Greedy Search and Heuristic Search algorithms to find the optimum cost for the randomly generated coordinates for five iterations.
- 8. Coordinates are generated randomly in a range of 80 and all these values are displayed.
- 9. For every set of 18 nodes, the optimum costs and runtimes for both the heuristic algorithms are calculated and displayed along with the respective network topologies.

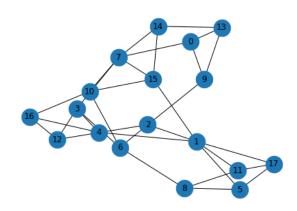
OUTPUTS/ GRAPHS:

Iteration-1:

Coordinates are [[2, 56], [32, 9], [78, 10], [37, 38], [46, 50], [55, 33], [55, 0], [4, 51], [40, 51], [73, 40], [38, 32], [33, 0], [32, 47], [7, 9], [49, 63], [37, 35], [5, 52], [18, 9]]



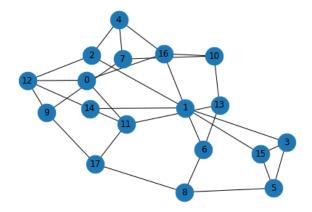
Greedy Search Optimum-cost 1053.017 Greedy Search Runtime 3.297ms



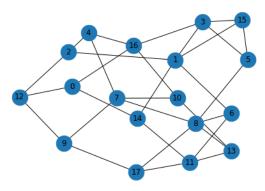
Heuristic Search Optimum-cost 1208.358 Heuristic Search Runtime 0.075ms

Iteration -2:

Coordinates are [[75, 42], [19, 25], [54, 27], [31, 48], [79, 36], [51, 9], [23, 39], [49, 23], [63, 9], [49, 6], [4, 72], [2, 8], [32, 33], [60, 38], [52, 8], [18, 12], [34, 32], [34, 15]]



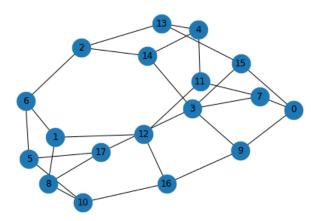
Greedy Search Optimum-cost 1022.027 Greedy Search Runtime 4.905ms



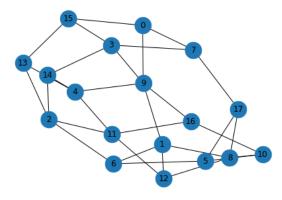
Heuristic Search Optimum-cost 1341.484 Heuristic Search Runtime 0.099ms

Iteration – 3:

Coordinates are [[21, 52], [62, 58], [27, 54], [6, 0], [26, 54], [39, 12], [78, 66], [12, 51], [32, 31], [51, 24], [72, 53], [31, 0], [49, 67], [73, 12], [26, 76], [39, 50], [6, 44], [36, 40]]



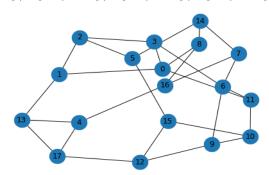
Greedy Search Optimum-cost 1120.041 Greedy Search Runtime 5.358ms



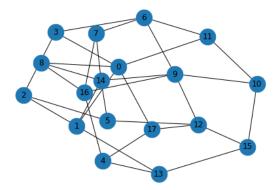
Heuristic Search Optimum-cost 1148.809 Heuristic Search Runtime 0.093ms

Iteration – 4:

Coordinates are [[17, 40], [8, 47], [41, 71], [65, 66], [44, 1], [78, 72], [29, 10], [51, 59], [79, 45], [53, 42], [62, 17], [48, 65], [24, 27], [77, 45], [17, 31], [55, 33], [69, 11], [3, 29]]



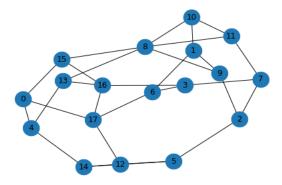
Greedy Search Optimum-cost 1243.921 Greedy Search Runtime 9.81ms



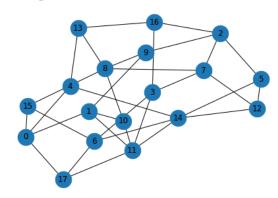
Heuristic Search Optimum-cost 1399.655 Heuristic Search Runtime 0.099ms

Iteration -5:

Coordinates are [[20, 29], [26, 34], [62, 25], [53, 16], [16, 40], [72, 31], [36, 56], [45, 78], [27, 6], [21, 74], [2, 10], [55, 63], [50, 36], [15, 10], [36, 43], [1, 55], [4, 70], [26, 20]]



Greedy Search Optimum-cost 880.148 Greedy Search Runtime 7.84ms



Heuristic Search Optimum-cost 1413.136 Heuristic Search Runtime 0.099ms

CONCLUSION:

In Greedy Search algorithm, as the network iterates, its total cost reduces. Based on the graph, the number of iterations required to attain the lowest cost differs.[1] Some graphs may require more iterations than others to get the lowest total cost.

In the heuristic search algorithm, on the other hand, the graph is at its lowest total cost as soon as all three requirements are met.[2] Thus, only the lowest output is displayed.

As can be witnessed from the examples above, almost always, the greedy search algorithm yields a more optimized graph than the heuristic algorithm. Since the heuristic technique is not iterative, it has a substantially shorter run time than the greedy search algorithm. In other words, the number of iterations affects run time. One can see that the runtime required by the algorithm for a graph with a higher starting total cost is somewhat longer than for networks that have a slight bit lower initial cost.

REFERENCES

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 5.pdf
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- 3. Display Graphs from adjacency matrix: https://stackoverflow.com/questions/29572623/plot-networkx-graph-from-adjacency-matrix-in-csv-file
- 4. BFSearch: https://www.geeksforgeeks.org/breadth-first-search-or-bfs-for-a-graph/
- 5. Graph Tools: https://graph-tool.skewed.de/static/doc/quickstart.html
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- 10. Professor's learning modules

APPENDIX:

BFSearch.py

```
import numpy as np
import networkx as ntx
import matplotlib.pyplot as mplt
class BFSearch:
    def BFSearch(adj mat):
        check = False
        visit list = [False for i in range(18)]
        nodes_visited = []
        visited count = 0
        for i in range(len(adj mat)):
            for j in range(18):
                # to verify edge between any 2 vertices
                if (not(visit list[j]) and adj mat[i][j] == 1
                        and (j not in nodes visited)):
                        nodes visited.append(j)
                        visited count += 1
                        visit list[j] = True
        # if all nodes are visited
        if (visited count == 18):
            check = True
        return check
    def display graph(adj mat):
        rows, cols = np.where(adj mat == 1)
        edges = zip(rows.tolist(), cols.tolist())
        graph = ntx.Graph()
        all rows = range(0, adj mat.shape[0])
        for n in all rows:
            graph.add node(n)
        graph.add edges from(edges)
        ntx.draw(graph, node size=600, with labels=True)
        mplt.show()
main.py
```

```
import random
import numpy as np
import timeit
from operator import itemgetter
from BFSearch import BFSearch
```

```
from graph tools import *
from enum import Flag
import math
import sys
sys.setrecursionlimit(1 000 000)
nodes list = list(range(0, 18))
adj_matrx = [[0] * 18 for _ in range(18)]
# to create an adjacent matrix 18x18 and a graph of 18 nodes
global graph
graph = Graph(directed=True)
for i in nodes list:
    graph.add vertex(i)
def set adj matrix():
    nodes list = list(range(18))
    adj_mat = [[0] * 18 for _ in range(18)]
    cand list = []
    for curr cand in nodes list:
        while True:
            random cand = random.choice(nodes list)
            if (random cand != curr cand):
                if (random cand not in cand list):
                    cand list.append(curr cand)
                    # updating the adjacency matrix
                    adj mat[curr cand][random cand] = 1
                    adj_mat[random_cand][curr_cand] = 1
                    # adding the edges generated randomely
                    graph.add edge(curr cand, random cand)
                    graph.add edge(random cand, curr cand)
            if (len(cand list) == 3):
                cand list = []
                break
    first constraint check(adj mat)
def first constraint check(adj mat, is optimal=False):
    # constraint to check if graph is a complete digraph
    if not BFSearch.BFSearch(adj mat):
        if not is optimal:
            set_adj_matrix()
        else:
            return False
    # constraint to check if the diameter is atleast 4
    for i in nodes list:
        dis, pre = graph.dijkstra(i)
        if any (4 < value for value in dis.values()):
            if not is optimal:
```

```
set adj matrix()
            else:
                return False
    # constraint to check if degree of graph is at least 3
    for i in adj mat:
        if i.count(1) < 3:
            if not is optimal:
                set adj matrix()
            else:
                return False
    if not is optimal:
        gen node coordinates (adj mat)
    else:
        return True
def second constraint check(adj mat):
    # constraint to check if graph is a complete digraph
    if not BFSearch.BFSearch(adj mat):
        return False
    # constraint to check if the diameter is atleast 4
    for i in nodes list:
        dis, pre = graph.dijkstra(i)
        if any (4 < value for value in dis.values()):
            return False
    # constraint to check if degree of graph is at least 3
    for i in adj mat:
        if i.count(1) < 3:
            return False
    return True
# this generates the (x,y) coordinates
# of every node on the graph
def gen node coordinates(adj mat):
    # selects coordinates randomly in range of 0-80
    coord range = list(range(0, 80))
    coord list = []
    while (len(coord list) < 18):
        x coord = random.choice(coord range)
        y coord = random.choice(coord range)
        x y = [x coord, y coord]
        if x y not in coord list:
            coord list.append(x_y)
    # prints all coordinates list
    print('Coordinates are {}\n'.format(coord list))
    total cost = total costs(adj mat, coord list)
# greedy search
```

```
def greedy search (edges cost, opt cost,
                  adj mat, coord list):
    # the edges with max weight are removed from the graph
    edg costs sorted = sorted(edges cost,
                              key=itemgetter(2), reverse=True)
    for i in range(len(edg costs sorted)):
        coord 1 = edg costs sorted[i][0]
        coord 2 = edg costs sorted[i][1]
        node 1 = coord list.index(coord 1)
        node 2 = coord list.index(coord 2)
        adj mat[node 1][node 2] = 0
        adj mat[node 2][node 1] = 0
        # the resultant graph is checked for constraints
        if first constraint check(adj mat, True):
            # if yes, calculate the total cost
            o cost = total costs(adj mat, coord list, True)
            # check if this new cost is less than the earlier cost
            if o cost < opt cost:</pre>
                opt cost = o cost
                # Loop through all edges to remove any edges
            for j in range(len(edg costs sorted)):
                coord 11 = edg costs sorted[j][0]
                coord 22 = edg_costs_sorted[j][1]
                node 11 = coord list.index(coord 11)
                node 22 = coord list.index(coord 22)
                adj_mat[node_11][node_22] = 0
                adj_mat[node 22][node 11] = 0
                # the resultant graph is checked for constraints
                if first constraint check(adj mat, True):
                    # if yes, calculate the total cost
                    o cost = total costs(adj mat, coord list, True)
                    if o cost < opt_cost:</pre>
                        opt cost = o cost
                    else:
                # if the total cost of resultant is not optimum
                # revert the removed edge
                    adj mat[node 11][node 22] = 1
                    adj_mat[node_22][node 11] = 1
        else:
            adj_mat[node_1][node_2] = 1
            adj mat[node 2][node 1] = 1
        adj mat[node 1][node 2] = 1
        adj mat[node 2][node 1] = 1
    # display the graph of the adjacent matrix
    mat array = []
    mat array = np.array(adj mat)
    BFSearch.display graph (mat array)
    print('Greedy Search Optimum-cost {}'.format(opt cost))
```

```
# heuristic
def heuristic search (edges cost, final total cost,
                     adj mat, cost mat, coord list):
    # Let the max weight
    max wgt = float('inf')
    # list the selected nodes to avoid redundant selections
    node 1 = [False for i in range(18)]
    res mat = [[0] * 18 \text{ for } in \text{ range}(18)]
    count = 0
    while (False in node 1):
        min wgt = max wgt
        begin = 0
        last = 0
        for x in range (18):
            if node 1[x]:
                for y in range(18):
                     # avoid cycles in the graph
                    if (not node 1[y] and cost mat[x][y] > 0):
                         if (cost mat[x][y] < min wgt):</pre>
                             min wgt = cost mat[x][y]
                             begin, last = x, y
        node 1[last] = True
        res mat[begin][last] = min wgt
        if min wgt == max wgt:
            res mat[begin][last] = 0
        count += 1
        # resultant matrix should have path with min cost possible
        res mat[last][begin] = res mat[begin][last]
    adj mat_2 = [[0] * 18 for _ in range(18)]
    for i in range(18):
        for j in range(18):
            if res mat[i][j] != 0:
                adj mat 2[i][j] = 1
    flg = True
    # check for constraints on the resultant graph
    while not flg == second constraint check(adj mat 2):
        thisflg = list(map(sum, adj mat 2))
        for i in range(18):
            for j in range(18):
                if adj mat[i][j] != adj mat 2[i][j] and thisflg [i] < 3:
                     adj mat 2[i][j] = 1
                    adj mat 2[j][i] = 1
                     thisflg = list(map(sum, adj mat 2))
```

```
# now calculate the total cost of resultant graph
    o cost = total costs(adj mat 2, coord list, True)
    # display graph of resultant adjacent matrix
    mat array = []
    mat array = np.array(adj mat 2)
    BFSearch.display graph(mat array)
    print('Heuristic Search Optimum-cost {}\n'.format(o cost))
def total costs(adj mat, coord list, is optimal=False):
    cost mat = [[0] * 18 \text{ for } in \text{ range}(18)]
    # coordinate list
    temp_coords = [[x_i, y_i] for x_i, row in enumerate(adj mat)
                       for y i, i in enumerate(row) if i == 1]
    temps = [tuple(sorted(x)) for x in temp coords]
    distinct coords = [list(x) for x in {*temps}]
    # to calculate Euclidean distance between any two nodes
    edges = []
    for item in distinct coords:
        coord = []
        for j in item:
            coord.append(coord list[j])
        edges.append(coord)
    # to get edge weight
    edges_cost = []
    for i in edges:
        a = np.array(i[0])
        b = np.array(item[1])
        dis = round(np.linalg.norm(a - b), 3)
        edges cost.append([i[0], i[1], dis])
    # sum up the edge cost
    tot cost = 0
    for i in edges cost:
        tot cost += i[-1]
    final totalcost = round(tot cost, 3)
    # matrix with edge costs
    x = 0
    for i in range(18):
        for j in range(18):
            if i < j and adj mat[i][j] == 1:</pre>
                cost mat[i][j] = edges cost[x][2]
                cost mat[j][i] = edges cost[x][2]
                x += 1
    if not is optimal:
        tot cost1 = timeit.timeit(lambda: greedy search(
            edges cost, final totalcost, adj mat, coord list),
```

READ ME:

install python 3.9.12

- 1. pip install graphtools
- 2. pip install networkx = = 2.7.1

Steps to run the program:

- Copy the code in appendix and open files in Terminal/IDE to paste and save them.
 Maintain the same mentioned names for the files and both the files are supposed to be in same folder.
- Run the "main.py" file and the results will be displayed.
- It takes a little while for all the iterations to complete. The results will be displayed sequentially.